5 Attribute declarations and specifications

2 5.1 General

Every data object has a type and rank and may have type parameters and other attributes that determine
the uses of the object. Collectively, these properties are the **attributes** of the object. The type of a
named data object is either specified explicitly in a type declaration statement or determined implicitly
by the first letter of its name (5.5). All of its attributes may be specified in a type declaration statement
or individually in separate specification statements.

8 A function has a type and rank and may have type parameters and other attributes that determine the 9 uses of the function. The type, rank, and type parameters are the same as those of its result variable.

10 A subroutine does not have a type, rank, or type parameters, but may have other attributes that 11 determine the uses of the subroutine.

12 5.2 Type declaration statements

13 5.2.1 Syntax

14 R501 type-declaration-stmt

is declaration-type-spec [[, attr-spec]...:] entity-decl-list

15 The type declaration statement specifies the type of the entities in the entity declaration list. The type 16 and type parameters are those specified by *declaration-type-spec*, except that the character length type 17 parameter may be overridden for an entity by the appearance of * *char-length* in its *entity-decl*.

18	R502	attr-spec	\mathbf{is}	access-spec
19			\mathbf{or}	ALLOCATABLE
20			or	ASYNCHRONOUS
21			\mathbf{or}	DIMENSION (array-spec)
22			or	EXTERNAL
23			\mathbf{or}	INTENT (intent-spec)
24			or	INTRINSIC
25			\mathbf{or}	language-binding-spec
26			\mathbf{or}	OPTIONAL
27			\mathbf{or}	PARAMETER
28			\mathbf{or}	POINTER
29			or	PROTECTED
30			or	SAVE
31			or	TARGET
32			or	VALUE
33			or	VOLATILE
34				
35	C501	(R501) The same $attr-spec$ sh	nall	not appear more than once in a given <i>type-declaration-stmt</i> .
36 37	C502	(R501) If a <i>language-binding</i> consist of a single <i>entitu-decl</i> .	-	ec with a NAME= specifier appears, the $entity$ - $decl$ - $list$ shall

C503 (R501) If a *language-binding-spec* is specified, the *entity-decl-list* shall not contain any procedure
 names.

The type declaration statement also specifies the attributes whose keywords appear in the *attr-spec*,
 except that the DIMENSION attribute may be specified or overridden for an entity by the appearance
 of *array-spec* in its *entity-decl*.

4 5	R503	entity-decl	<pre>is object-name [(array-spec)] [* char-length] [initialization] or function-name [* char-length]</pre>		
6	C504	(R503) If the entity is not of	of type character, * <i>char-length</i> shall not appear.		
7	C505	(R501) If initialization appo	ears, a double-colon separator shall appear before the <i>entity-decl-list</i> .		
8 9 10 11	an object in a named common block unless the type declaration is in a block data program unit, an object in blank common, an allocatable variable, an external function, an intrinsic function,				
12	C507	(R503) An <i>initialization</i> sh	all appear if the entity is a named constant $(5.3.11)$.		
13 14	C508	(R503) The <i>function-name</i> shall be the name of an external function, an intrinsic function, a function dummy procedure, or a statement function.			
15	R504	object-name	is name		
16	C509	(R504) The <i>object-name</i> sh	all be the name of a data object.		
17	R505	initialization	is = initialization-expr		
18 19	R506	null-init	or => null-init is function-reference		
20 21	C510	. ,	<i>nitialization</i> , the entity shall have the POINTER attribute. If $=$ e entity shall not have the POINTER attribute.		
22 23					
24 25 26	25 explicit type declaration statement is not required; however, it is permitted. Specifying a type for a				

27 generic properties from that function.

An automatic data object is a nondummy data object with a type parameter or array bound that depends on the value of a *specification-expr* that is not an initialization expression.

NOTE 5.1 An automatic object shall not have the SAVE attribute and shall not appear in a common block.

30 If a type parameter in a *declaration-type-spec* or in a *char-length* in an *entity-decl* is defined by an 31 expression that is not an initialization expression, the type parameter value is established on entry to 32 the procedure and is not affected by any redefinition or undefinition of the variables in the expression 33 during execution of the procedure.

34 5.2.2 Initialization

35 The appearance of *initialization* in an *entity-decl* for an entity without the PARAMETER attribute

36 specifies that the entity is a variable with explicit initialization. Explicit initialization alternatively 37 may be specified in a DATA statement unless the variable is of a derived type for which default initial-

ization is specified. If *initialization* is =*initialization-expr*, the variable is initially defined with the value

specified by the *initialization-expr*; if necessary, the value is converted according to the rules of intrinsic
 assignment (7.4.1.3) to a value that agrees in type, type parameters, and shape with the variable. A
 variable, or part of a variable, shall not be explicitly initialized more than once in a program. If the
 variable is an array, it shall have its shape specified in either the type declaration statement or a previous
 attribute specification statement in the same scoping unit.

6 If *initialization* is =>null-init, the variable shall be a pointer, and its initial association status is disas-7 sociated.

8 Explicit initialization of a variable that is not in a common block implies the SAVE attribute, which9 may be confirmed by explicit specification.

10 5.2.3 Examples of type declaration statements

NOTE 5.2

```
REAL A (10)

LOGICAL, DIMENSION (5, 5) :: MASK1, MASK2

COMPLEX :: CUBE_ROOT = (-0.5, 0.866)

INTEGER, PARAMETER :: SHORT = SELECTED_INT_KIND (4)

INTEGER (SHORT) K ! Range at least -9999 to 9999.

REAL (KIND (0.0D0)) A

REAL (KIND = 2) B

COMPLEX (KIND = KIND (0.0D0)) :: C

CHARACTER (LEN = 10, KIND = 2) A

CHARACTER B, C *20

TYPE (PERSON) :: CHAIRMAN

TYPE (NODE), POINTER :: HEAD => NULL ()

TYPE (humongous_matrix (k=8, d=1000)) :: mat
```

(The last line above uses a type definition from Note 4.25.)

11 5.3 Attributes

12 5.3.1 Constraints

An attribute may be explicitly specified by an *attr-spec* in a type declaration statement or by an attribute
specification statement (5.4). The following constraints apply to attributes.

15 C512 An entity shall not be explicitly given any attribute more than once in a scoping unit.

- 16 C513 An array-spec for a function result that does not have the ALLOCATABLE or POINTER
 17 attribute shall be an explicit-shape-spec-list.
- 18 C514 The ALLOCATABLE, POINTER, or OPTIONAL attribute shall not be specified for a dummy argument of a procedure that has a *proc-language-binding-spec*.

20 5.3.2 Accessibility attribute

21 The **accessibility attribute** specifies the accessibility of an entity via a particular identifier.

1	R507	access-spec	is	PUBLIC
2			or	PRIVATE

3 C515 (R507) An access-spec shall appear only in the specification-part of a module.

4 Identifiers that are specified in a module or accessible in that module by use association have either
5 the PUBLIC or PRIVATE attribute. Identifiers for which an *access-spec* is not explicitly specified in
6 that module have the default accessibility attribute for that module. The default accessibility attribute
7 for a module is PUBLIC unless it has been changed by a PRIVATE statement (5.4.1). Only identifiers
8 that have the PUBLIC attribute in that module are available to be accessed from that module by use

9 association.

NOTE 5.3

In order for an identifier to be accessed by use association, it must have the PUBLIC attribute in the module from which it is accessed. It can nonetheless have the PRIVATE attribute in a module in which it is accessed by use association, and therefore not be available for use association from a module where it is PRIVATE.

NOTE 5.4

An example of an accessibility specification is:

REAL, PRIVATE :: X, Y, Z

10 5.3.3 ALLOCATABLE attribute

11 An entity with the **ALLOCATABLE attribute** is a variable for which space is allocated by an AL-12 LOCATE statement (6.3.1) or by an intrinsic assignment statement (7.4.1.3).

13 5.3.4 ASYNCHRONOUS attribute

An entity with the ASYNCHRONOUS attribute is a variable that may be subject to asynchronous
 input/output.

16 The base object of a variable shall have the ASYNCHRONOUS attribute in a scoping unit if

- (1) the variable appears in an executable statement or specification expression in that scoping
 unit and
- any statement of the scoping unit is executed while the variable is a pending I/O storage
 sequence affector (9.5.1.4).

Use of a variable in an asynchronous input/output statement can imply the ASYNCHRONOUS attribute;
 see subclause (9.5.1.4).

An object may have the ASYNCHRONOUS attribute in a particular scoping unit without necessarily
having it in other scoping units (11.2.1, 16.4.1.3). If an object has the ASYNCHRONOUS attribute,
then all of its subobjects also have the ASYNCHRONOUS attribute.

NOTE 5.5

The ASYNCHRONOUS attribute specifies the variables that might be associated with a pending input/output storage sequence (the actual memory locations on which asynchronous input/output is being performed) while the scoping unit is in execution. This information could be used by the compiler to disable certain code motion optimizations.

The ASYNCHRONOUS attribute is similar to the VOLATILE attribute. It is intended to facilitate

NOTE 5.5 (cont.)

traditional code motion optimizations in the presence of asynchronous input/output.

1 5.3.5 BIND attribute for data entities

2 The BIND attribute for a variable or common block specifies that it is capable of interoperating with a

- 3 C variable that has external linkage (15.3).
- 4 R508 language-binding-spec is BIND (C [, NAME = scalar-char-initialization-expr])
- 5 C516 An entity with the BIND attribute shall be a common block, variable, or procedure.
- 6 C517 A variable with the BIND attribute shall be declared in the specification part of a module.
- 7 C518 A variable with the BIND attribute shall be interoperable (15.2).
- 8 C519 Each variable of a common block with the BIND attribute shall be interoperable.
- 9 C520 (R508) The scalar-char-initialization-expr shall be of default character kind. If the value of the scalar-char-initialization-expr after discarding leading and trailing blanks has nonzero length,
 11 it shall be valid as an identifier on the companion processor.

NOTE 5.6

The C International Standard provides a facility for creating C identifiers whose characters are not restricted to the C basic character set. Such a C identifier is referred to as a universal character name (6.4.3 of the C International Standard). The name of such a C identifier might include characters that are not part of the representation method used by the processor for type default character. If so, the C entity cannot be referenced from Fortran.

12 The BIND attribute for a variable or common block implies the SAVE attribute, which may be confirmed13 by explicit specification.

14 5.3.6 DIMENSION attribute

15 5.3.6.1 General

16 The DIMENSION attribute specifies that an entity is an array. The rank or rank and shape is17 specified by its *array-spec*.

18	R509	array- $spec$	is	explicit-shape-spec-list
19			or	$assumed\-shape\-spec\-list$
20			or	deferred-shape-spec-list
21			or	assumed-size-spec

22 C521 (R509) The maximum rank is seven.

NOTE 5.7

```
Examples of DIMENSION attribute specifications are:

SUBROUTINE EX (N, A, B)

REAL, DIMENSION (N, 10) :: W ! Automatic explicit-shape array

REAL A (:), B (0:) ! Assumed-shape arrays

REAL, POINTER :: D (:, :) ! Array pointer

REAL, DIMENSION (:), POINTER :: P ! Array pointer
```

NOTE 5.7 (cont.)

REAL, ALLOCATABLE, DIMENSION (:) :: E ! Allocatable array

1 5.3.6.2 Explicit-shape array

2 An explicit-shape array is a named array that is declared with an *explicit-shape-spec-list*. This specifies
3 explicit values for the bounds in each dimension of the array.

4	R510	explicit-shape-spec	is	[lower-bound:]	upper- $bound$
---	------	---------------------	----	----------------	----------------

5 R511 lower-bound is specification-expr

6 R512 upper-bound is specification-expr

7 C522 (R510) An *explicit-shape-spec* whose bounds are not initialization expressions shall appear only
 8 in a subprogram or interface body.

9 If an explicit-shape array has bounds that are not initialization expressions, the bounds, and hence
10 shape, are determined at entry to the procedure by evaluating the bounds expressions. The bounds of
11 such an array are unaffected by the redefinition or undefinition of any variable during execution of the
12 procedure.

13 The values of each *lower-bound* and *upper-bound* determine the bounds of the array along a particular 14 dimension and hence the extent of the array in that dimension. The value of a lower bound or an upper 15 bound may be positive, negative, or zero. The subscript range of the array in that dimension is the set 16 of integer values between and including the lower and upper bounds, provided the upper bound is not 17 less than the lower bound. If the upper bound is less than the lower bound, the range is empty, the 18 extent in that dimension is zero, and the array is of zero size. If the *lower-bound* is omitted, the default 19 value is 1. The rank is equal to the number of *explicit-shape-specs*.

20 **5.3.6.3** Assumed-shape array

- An assumed-shape array is a nonpointer dummy argument array that takes its shape from the associated actual argument array.
- 23 R513 assumed-shape-spec is [lower-bound]:
- 24 The rank is equal to the number of colons in the *assumed-shape-spec-list*.

The extent of a dimension of an assumed-shape array dummy argument is the extent of the corresponding dimension of the associated actual argument array. If the lower bound value is d and the extent of the

corresponding dimension of the associated actual argument array is s, then the value of the upper bound is s + d - 1. If *lower-bound* appears it specifies the lower bound; otherwise the lower bound is 1.

28 is s + a - 1. If *tower-bound* appears it specifies the lower bound, otherwise the lower bound is 1.

29 5.3.6.4 Deferred-shape array

30 A deferred-shape array is an allocatable array or an array pointer.

An allocatable array is an array that has the ALLOCATABLE attribute and a specified rank, but its
 bounds, and hence shape, are determined by allocation or argument association.

An array pointer is an array with the POINTER attribute and a specified rank. Its bounds, and hence
shape, are determined when it is associated with a target.

- 35 R514 deferred-shape-spec is :
- 36 C523 An array that has the POINTER or ALLOCATABLE attribute shall have an array-spec that is
 37 a deferred-shape-spec-list.

 $1 \quad {\rm The \ rank \ is \ equal \ to \ the \ number \ of \ colons \ in \ the \ deferred-shape-spec-list.}$

2 The size, bounds, and shape of an unallocated allocatable array or a disassociated array pointer are

3 undefined. No part of such an array shall be referenced or defined; however, the array may appear as an

- 4 argument to an intrinsic inquiry function as specified in 13.1.
- 5 The bounds of each dimension of an allocatable array are those specified when the array is allocated.
- 6 The bounds of each dimension of an array pointer may be specified in two ways:
- 7 (1) in an ALLOCATE statement (6.3.1) when the target is allocated;
- 8 (2) by pointer assignment (7.4.2).

9 The bounds of the array pointer or allocatable array are unaffected by any subsequent redefinition or10 undefinition of variables on which the bounds' expressions depend.

11 5.3.6.5 Assumed-size array

12 An assumed-size array is a dummy argument array whose size is assumed from that of an associated 13 actual argument. The rank and extents may differ for the actual and dummy arrays; only the size of the 14 actual array is assumed by the dummy array. An assumed-size array is declared with an *assumed-size-*15 *spec*.

- 16 R515 assumed-size-spec is [explicit-shape-spec-list,][lower-bound:]*
- 17 C524 An assumed-size-spec shall not appear except as the declaration of the array bounds of a dummy data argument.
- 19C525An assumed-size array with INTENT (OUT) shall not be of a type for which default initialization20is specified.
- 21 The size of an assumed-size array is determined as follows.
- (1) If the actual argument associated with the assumed-size dummy array is an array of any
 type other than default character, the size is that of the actual array.
- 24 (2) If the actual argument associated with the assumed-size dummy array is an array element 25 of any type other than default character with a subscript order value of r (6.2.2.2) in an 26 array of size x, the size of the dummy array is x - r + 1.
- 27 (3) If the actual argument is a default character array, default character array element, or a 28 default character array element substring (6.1.1), and if it begins at character storage unit t29 of an array with c character storage units, the size of the dummy array is MAX (INT ((c - t + 1)/e), 0), where e is the length of an element in the dummy character array.
- 31 (4) If the actual argument is of type default character and is a scalar that is not an array element 32 or array element substring designator, the size of the dummy array is MAX (INT (l/e), 0), 33 where *e* is the length of an element in the dummy character array and *l* is the length of the 34 actual argument.
- 35 The rank is equal to one plus the number of *explicit-shape-specs*.

An assumed-size array has no upper bound in its last dimension and therefore has no extent in its last
dimension and no shape. An assumed-size array name shall not be written as a whole array reference
except as an actual argument in a procedure reference for which the shape is not required.

39 If an *explicit-shape-spec-list* appears, it specifies the bounds of the first rank -1 dimensions. If *lower-*40 *bound* appears it specifies the lower bound of the last dimension; otherwise that lower bound is 1. An 41 assumed-size array may be subscripted or sectioned (6.2.2.3). The upper bound shall not be omitted 42 from a subscript triplet in the last dimension

 $42 \quad {\rm from \ a \ subscript \ triplet \ in \ the \ last \ dimension.}$

1 If an assumed-size array has bounds that are not initialization expressions, the bounds are determined

at entry to the procedure. The bounds of such an array are unaffected by the redefinition or undefinitionof any variable during execution of the procedure.

4 5.3.7 EXTERNAL attribute

5 The **EXTERNAL attribute** specifies that an entity is an external procedure, dummy procedure, 6 procedure pointer, or block data subprogram.

7 C526 An entity shall not have both the EXTERNAL attribute and the INTRINSIC attribute.

8 In an external subprogram, the EXTERNAL attribute shall not be specified for a procedure defined by9 the subprogram.

10 If an external procedure or dummy procedure is used as an actual argument or is the target of a procedure11 pointer assignment, it shall be declared to have the EXTERNAL attribute.

12 A procedure that has both the EXTERNAL and POINTER attributes is a procedure pointer.

13 **5.3.8 INTENT attribute**

14 The INTENT attribute specifies the intended use of a dummy argument. An INTENT (IN) dummy 15 argument is suitable for receiving data from the invoking scoping unit, an INTENT (OUT) dummy 16 argument is suitable for returning data to the invoking scoping unit, and an INTENT (INOUT) dummy 17 argument is suitable for use both to receive data from and to return data to the invoking scoping unit.

18	R516	intent-spec	is	IN
19			or	OUT
20			or	INOUT

21C527An entity with the INTENT attribute shall be a dummy data object or a dummy procedure22pointer.

C528 (R516) A nonpointer object with the INTENT (IN) attribute shall not appear in a variable
definition context (16.5.7).

C529 A pointer with the INTENT (IN) attribute shall not appear in a pointer association context
 (16.5.8).

The INTENT (IN) attribute for a nonpointer dummy argument specifies that it shall neither be defined nor become undefined during the execution of the procedure. The INTENT (IN) attribute for a pointer dummy argument specifies that during the execution of the procedure its association shall not be changed except that it may become undefined if the target is deallocated other than through the pointer (16.4.2.1.3).

The INTENT (OUT) attribute for a nonpointer dummy argument specifies that the dummy argument becomes undefined on invocation of the procedure, except for any subcomponents that are defaultinitialized (4.5.4.5). Any actual argument that becomes associated with such a dummy argument shall be definable. The INTENT (OUT) attribute for a pointer dummy argument specifies that on invocation of the procedure the pointer association status of the dummy argument becomes undefined. Any actual argument associated with such a pointer dummy shall be a pointer variable.

NOTE 5.8

If the actual argument is finalizable it will be finalized before undefinition and default initialization of the dummy argument (4.5.6).

- $1 \quad {\rm The \ INTENT \ (INOUT) \ attribute \ for \ a \ nonpointer \ dummy \ argument \ specifies \ that \ any \ actual \ argument$
- 2 $\,$ associated with the dummy argument shall be definable. The INTENT (INOUT) attribute for a pointer $\,$
- 3 dummy argument specifies that any actual argument associated with the dummy argument shall be a
- 4 pointer variable.

The INTENT attribute for an allocatable dummy argument applies to both the allocation status and the definition status. An actual argument associated with an INTENT(OUT) allocatable dummy argument is deallocated on procedure invocation (6.3.3.1).

5 If no INTENT attribute is specified for a dummy argument, its use is subject to the limitations of the 6 associated actual argument (12.4.1.2, 12.4.1.3, 12.4.1.4).

NOTE 5.10

An example of INTENT specification is:

SUBROUTINE MOVE (FROM, TO) USE PERSON_MODULE TYPE (PERSON), INTENT (IN) :: FROM TYPE (PERSON), INTENT (OUT) :: TO

7 If an object has an INTENT attribute, then all of its subobjects have the same INTENT attribute.

NOTE 5.11

If a dummy argument is a derived-type object with a pointer component, then the pointer as a pointer is a subobject of the dummy argument, but the target of the pointer is not. Therefore, the restrictions on subobjects of the dummy object apply to the pointer in contexts where it is used as a pointer, but not in contexts where it is dereferenced to indicate its target. For example, if X is a dummy argument of derived type with an integer pointer component P, and X has INTENT(IN), then the statement

X%P => NEW_TARGET

is prohibited, but

X%P = 0

is allowed (provided that X%P is associated with a definable target).

Similarly, the INTENT restrictions on pointer dummy arguments apply only to the association of the dummy argument; they do not restrict the operations allowed on its target.

NOTE 5.12

Argument intent specifications serve several purposes in addition to documenting the intended use of dummy arguments. A processor can check whether an INTENT (IN) dummy argument is used in a way that could redefine it. A slightly more sophisticated processor could check to see whether an INTENT (OUT) dummy argument could possibly be referenced before it is defined. If the procedure's interface is explicit, the processor can also verify that actual arguments corresponding to INTENT (OUT) or INTENT (INOUT) dummy arguments are definable. A more sophisticated processor could use this information to optimize the translation of the referencing scoping unit by taking advantage of the fact that actual arguments corresponding to INTENT (IN) dummy arguments will not be changed and that any prior value of an actual argument corresponding to

NOTE 5.12 (cont.)

an INTENT (OUT) dummy argument will not be referenced and could thus be discarded.

INTENT (OUT) means that the value of the argument after invoking the procedure is entirely the result of executing that procedure. If there is any possibility that an argument should retain its current value rather than being redefined, INTENT (INOUT) should be used rather than INTENT (OUT), even if there is no explicit reference to the value of the dummy argument. Because an INTENT(OUT) variable is considered undefined on entry to the procedure, any default initialization specified for its type will be applied.

INTENT (INOUT) is not equivalent to omitting the INTENT attribute. The actual argument corresponding to an INTENT (INOUT) dummy argument is always required to be definable, while an argument corresponding to a dummy argument without an INTENT attribute need be definable only if the dummy argument is actually redefined.

1 5.3.9 INTRINSIC attribute

2 The INTRINSIC attribute specifies that the entity is an intrinsic procedure. It may be a generic
3 procedure (13.5), a specific procedure (13.6), or both.

4 If the specific name of an intrinsic procedure (13.6) is used as an actual argument, the name shall be
5 explicitly specified to have the INTRINSIC attribute. An intrinsic procedure whose specific name is
6 marked with a bullet (•) in 13.6 shall not be used as an actual argument.

7 C530 If the name of a generic intrinsic procedure is explicitly declared to have the INTRINSIC attribute, and it is also the generic name of one or more generic interfaces (12.3.2.1) accessible in the same scoping unit, the procedures in the interfaces and the specific intrinsic procedures shall
all be functions or all be subroutines, and the characteristics of the specific intrinsic procedures and the procedures in the interfaces shall differ as specified in 16.2.3.

12 5.3.10 OPTIONAL attribute

13 The **OPTIONAL attribute** specifies that the dummy argument need not be associated with an actual 14 argument in a reference to the procedure (12.4.1.6). The PRESENT intrinsic function can be used to 15 determine whether an optional dummy argument is associated with an actual argument.

16 C531 An entity with the OPTIONAL attribute shall be a dummy argument.

17 5.3.11 PARAMETER attribute

18 The PARAMETER attribute specifies that an entity is a named constant. The entity has the value
19 specified by its *initialization-expr*, converted, if necessary, to the type, type parameters and shape of the
20 entity.

21 C532 An entity with the PARAMETER attribute shall not be a variable or a procedure.

A named constant shall not be referenced unless it has been defined previously in the same statement,
defined in a prior statement, or made accessible by use or host association.

NOTE 5.13

Examples of declarations with a PARAMETER attribute are: REAL, PARAMETER :: ONE = 1.0, Y = 4.1 / 3.0 INTEGER, DIMENSION (3), PARAMETER :: ORDER = (/ 1, 2, 3 /) NOTE 5.13 (cont.)

TYPE(NODE), PARAMETER :: DEFAULT = NODE(0, NULL ())

1 5.3.12 POINTER attribute

2 Entities with the **POINTER attribute** can be associated with different data objects or procedures
3 during execution of a program. A pointer is either a data pointer or a procedure pointer. Procedure
4 pointers are described in 12.3.2.3.

- 5 C533 An entity with the POINTER attribute shall not have the ALLOCATABLE, INTRINSIC, or
 6 TARGET attribute.
- 7 C534 A procedure with the POINTER attribute shall have the EXTERNAL attribute.

8 A data pointer shall not be referenced unless it is pointer associated with a target object that is defined.

9 A data pointer shall not be defined unless it is pointer associated with a target object that is definable.

10 If a data pointer is associated, the values of its deferred type parameters are the same as the values of11 the corresponding type parameters of its target.

12 A procedure pointer shall not be referenced unless it is pointer associated with a target procedure.

NOTE 5.14

Examples of POINTER attribute specifications are:

TYPE (NODE), POINTER :: CURRENT, TAIL REAL, DIMENSION (:, :), POINTER :: IN, OUT, SWAP

For a more elaborate example see C.2.1.

13 5.3.13 PROTECTED attribute

- 14 The **PROTECTED attribute** imposes limitations on the usage of module entities.
- 15 C535 The PROTECTED attribute shall be specified only in the specification part of a module.
- 16 C536 An entity with the PROTECTED attribute shall be a procedure pointer or variable.
- 17 C537 An entity with the PROTECTED attribute shall not be in a common block.

C538 A nonpointer object that has the PROTECTED attribute and is accessed by use association
 shall not appear in a variable definition context (16.5.7) or as the *data-target* or *proc-target* in
 a *pointer-assignment-stmt*.

C539 A pointer that has the PROTECTED attribute and is accessed by use association shall not
 appear in a pointer association context (16.5.8).

Other than within the module in which an entity is given the PROTECTED attribute, or within any ofits descendant submodules,

- 25 (1) if it is a nonpointer object, it is not definable, and
- (2) if it is a pointer, its association status shall not be changed except that it may become undefined if its target is deallocated other than through the pointer (16.4.2.1.3) or if its target becomes undefined by execution of a RETURN or END statement.

1 If an object has the PROTECTED attribute, all of its subobjects have the PROTECTED attribute.

NOTE 5.15

```
An example of the PROTECTED attribute:
MODULE temperature
REAL, PROTECTED :: temp_c, temp_f
CONTAINS
SUBROUTINE set_temperature_c(c)
REAL, INTENT(IN) :: c
temp_c = c
temp_f = temp_c*(9.0/5.0) + 32
END SUBROUTINE
END MODULE
```

The PROTECTED attribute ensures that the variables temp_c and temp_f cannot be modified other than via the set_temperature_c procedure, thus keeping them consistent with each other.

2 5.3.14 SAVE attribute

3 The SAVE attribute specifies that a variable retains its association status, allocation status, definition
4 status, and value after execution of a RETURN or END statement unless it is a pointer and its target
5 becomes undefined (16.4.2.1.3(4)). If it is a local variable of a subprogram it is shared by all instances
6 (12.5.2.3) of the subprogram.

(12.3.2.3) of the subprogram.

7 $\,$ Giving a common block the SAVE attribute confers the attribute on all variables in the common block.

8 C540 An entity with the SAVE attribute shall be a common block, variable, or procedure pointer.

9 C541 The SAVE attribute shall not be specified for a dummy argument, a function result, an automatic
 10 data object, or an object that is in a common block.

11 A saved entity is an entity that has the SAVE attribute. An unsaved entity is an entity that does not 12 have the SAVE attribute.

13 The SAVE attribute has no effect on entities declared in a main program. If a common block has the 14 SAVE attribute in any scoping unit that is not a main program, it shall have the SAVE attribute in 15 every scoping unit that is not a main program.

16 5.3.15 TARGET attribute

17 The TARGET attribute specifies that a data object may have a pointer associated with it (7.4.2).18 An object without the TARGET attribute shall not have an accessible pointer associated with it.

19 C542 An entity with the TARGET attribute shall be a variable.

20 C543 An entity with the TARGET attribute shall not have the POINTER attribute.

NOTE 5.16

In addition to variables explicitly declared to have the TARGET attribute, the objects created by allocation of pointers (6.3.1.2) have the TARGET attribute.

21 If an object has the TARGET attribute, then all of its nonpointer subobjects also have the TARGET

1 attribute.

NOTE 5.17

Examples of TARGET attribute specifications are:

TYPE (NODE), TARGET :: HEAD REAL, DIMENSION (1000, 1000), TARGET :: A, B

For a more elaborate example see C.2.2.

NOTE 5.18

Every object designator that starts from a target object will have either the TARGET or POINTER attribute. If pointers are involved, the designator might not necessarily be a subobject of the original target object, but because pointers may point only to targets, there is no way to end up at a nonpointer that is not a target.

2 5.3.16 VALUE attribute

- 3 The VALUE attribute specifies a type of argument association (12.4.1.2) for a dummy argument.
- 4 C544 An entity with the VALUE attribute shall be a scalar dummy data object.
- 5 C545 An entity with the VALUE attribute shall not have the ALLOCATABLE, INTENT(INOUT),
 6 INTENT(OUT), POINTER, or VOLATILE attributes.
- 7 C546 If an entity has the VALUE attribute, any length type parameter value in its declaration shall
 8 be omitted or specified by an initialization expression.

9 5.3.17 VOLATILE attribute

10 The VOLATILE attribute specifies that an object may be referenced, defined, or become undefined,
11 by means not specified by the program.

12 C547 An entity with the VOLATILE attribute shall be a variable that is not an INTENT(IN) dummy argument.

14 An object may have the VOLATILE attribute in a particular scoping unit without necessarily having 15 it in other scoping units (11.2.1, 16.4.1.3). If an object has the VOLATILE attribute, then all of its 16 subobjects also have the VOLATILE attribute.

NOTE 5.19

The Fortran processor should use the most recent definition of a volatile object when a value is required. Likewise, it should make the most recent Fortran definition available. It is the programmer's responsibility to manage any interaction with non-Fortran processes.

- 17 A pointer with the VOLATILE attribute may additionally have its association status, dynamic type and
- 18 type parameters, and array bounds changed by means not specified by the program.

NOTE 5.20

If the target of a pointer is referenced, defined, or becomes undefined, by means not specified by the program, while the pointer is associated with the target, then the pointer shall have the VOLATILE attribute. Usually a pointer should have the VOLATILE attribute if its target has the VOLATILE attribute. Similarly, all members of an EQUIVALENCE group should have the NOTE 5.20 (cont.)

VOLATILE attribute if one member has the VOLATILE attribute.

An allocatable object with the VOLATILE attribute may additionally have its allocation status, dynamic
 type and type parameters, and array bounds changed by means not specified by the program.

3 5.4 Attribute specification statements

4 5.4.1 Accessibility statements

5	R517	access-stmt	is	access-spec [[:::] access-id-list]
6	R518	access-id	\mathbf{is}	use-name
7			or	generic-spec

8 C548 (R517) An access-stmt shall appear only in the specification-part of a module. Only one accessibility statement with an omitted access-id-list is permitted in the specification-part of a module.
10 module.

C549 (R518) Each use-name shall be the name of a named variable, procedure, derived type, named
 constant, or namelist group.

13 An access-stmt with an access-id-list specifies the accessibility attribute (5.3.2), PUBLIC or PRIVATE,

14 of each access-id in the list. An access-stmt without an access-id list specifies the default accessibility

15 that applies to all potentially accessible identifiers in the *specification-part* of the module. The statement

16 PUBLIC

17 specifies a default of public accessibility. The statement

18 PRIVATE

specifies a default of private accessibility. If no such statement appears in a module, the default is publicaccessibility.

NOTE 5.21

Examples of accessibility statements are: MODULE EX PRIVATE PUBLIC :: A, B, C, ASSIGNMENT (=), OPERATOR (+)

21 5.4.2 ALLOCATABLE statement

22	R519	$allocatable\hbox{-}stmt$	is	ALLOCATABLE $[::] \blacksquare$
23				\blacksquare object-name [(array-spec)]
24				$\blacksquare [, object\text{-}name [(array\text{-}spec)]] \dots$

25 This statement specifies the ALLOCATABLE attribute (5.3.3) for a list of objects.

NOTE 5.22

An example of an ALLOCATABLE statement is:

REAL A, B (:), SCALAR

NOTE 5.22 (cont.)

ALLOCATABLE :: A (:, :), B, SCALAR

1 5.4.3 ASYNCHRONOUS statement

- 2 R520 asynchronous-stmt is ASYNCHRONOUS [::] object-name-list
- 3 The ASYNCHRONOUS statement specifies the ASYNCHRONOUS attribute (5.3.4) for a list of objects.

4 5.4.4 BIND statement

5	R521	bind- $stmt$	is	$language-binding-spec \ [\ :: \] \ bind-entity-list$
6	R522	bind- $entity$	\mathbf{is}	entity-name
7			\mathbf{or}	/ common-block-name /

8 C550 (R521) If the language-binding-spec has a NAME= specifier, the bind-entity-list shall consist of
9 a single bind-entity.

10 The BIND statement specifies the BIND attribute (5.3.5) for a list of variables and common blocks.

11 5.4.5 DATA statement

12 R523 data-stmt is DATA data-stmt-set [[,] data-stmt-set]...

13 This statement specifies explicit initialization (5.2.2).

A variable, or part of a variable, shall not be explicitly initialized more than once in a program. If a
nonpointer object has been specified with default initialization in a type definition, it shall not appear
in a data-stmt-object-list.

A variable that appears in a DATA statement and has not been typed previously may appear in a
subsequent type declaration only if that declaration confirms the implicit typing. An array name,
array section, or array element that appears in a DATA statement shall have had its array properties
established by a previous specification statement.

Except for variables in named common blocks, a named variable has the SAVE attribute if any part ofit is initialized in a DATA statement, and this may be confirmed by explicit specification.

23	R524	data- $stmt$ - set	\mathbf{is}	data-stmt-object-list / data-stmt-value-list /
24	R525	$data\mathchar`stmt\mathchar`object$	is	variable
25			or	$data\-implied\-do$
26	R526	$data\-implied\-do$	\mathbf{is}	$(data-i-do-object-list , data-i-do-variable = \blacksquare$
27				■ scalar-int-expr , scalar-int-expr [, scalar-int-expr])
28	R527	data-i-do-object	\mathbf{is}	array-element
29			\mathbf{or}	scalar- $structure$ - $component$
30			\mathbf{or}	data-implied-do
31	R528	$data\-i\-do\-variable$	is	scalar-int-variable
32	C551	(R525) In a <i>variable</i> that is	a da	ta-stmt-object, any subscript, section subscript, substring start-
33				point shall be an initialization expression.

C552 (R525) A variable whose designator appears as a *data-stmt-object* or a *data-i-do-object* shall not be a dummy argument, made accessible by use association or host association, in a named common block unless the DATA statement is in a block data program unit, in a blank common

1		block, a function name, a function result name, an automatic object, or an allocatable variable.				
2 3	C553	(R525) A <i>data-i-do-object</i> or a <i>variable</i> that appears as a <i>data-stmt-object</i> shall not be an object designator in which a pointer appears other than as the entire rightmost <i>part-ref</i> .				
4	C554	(R528) The <i>data-i-do-variable</i> shall be a named variable.				
5 6 7	C555	(R526) A <i>scalar-int-expr</i> of a <i>data-implied-do</i> shall involve as primaries only constants, subobjects of constants, or DO variables of the containing <i>data-implied-dos</i> , and each operation shall be intrinsic.				
8	C556	(R527) The <i>array-element</i> shall be a variable.				
9	C557	(R527) The <i>scalar-structure-component</i> shall be a variable.				
10 11	C558	(R527) The scalar-structure-component shall contain at least one part-ref that contains a subscript-list.				
12 13 14 15	C559	59 (R527) In an <i>array-element</i> or a <i>scalar-structure-component</i> that is a <i>data-i-do-object</i> , any sub- script shall be an expression whose primaries are either constants, subobjects of constants, or DO variables of this <i>data-implied-do</i> or the containing <i>data-implied-dos</i> , and each operation shall be intrinsic.				
16 17 18	R529 R530	data-stmt-valueis[data-stmt-repeat *] data-stmt-constantdata-stmt-repeatisscalar-int-constantorscalar-int-constant-subobject				
19 20 21	C560	(R530) The <i>data-stmt-repeat</i> shall be positive or zero. If the <i>data-stmt-repeat</i> is a named constant, it shall have been declared previously in the scoping unit or made accessible by use association or host association.				
22 23 24 25 26 27	R531	data-stmt-constant is scalar-constant or scalar-constant-subobject or signed-int-literal-constant or signed-real-literal-constant or null-init or structure-constructor				
28 29 30	C561	(R531) If a DATA statement constant value is a named constant or a structure constructor, the named constant or derived type shall have been declared previously in the scoping unit or made accessible by use or host association.				
31	C562	(R531) If a <i>data-stmt-constant</i> is a <i>structure-constructor</i> , it shall be an initialization expression.				
32	R532	int-constant-subobject is $constant-subobject$				
33	C563	(R532) <i>int-constant-subobject</i> shall be of type integer.				
34	R533	constant-subobject is designator				
35	C564	(R533) constant-subobject shall be a subobject of a constant.				
36 37	C565	$(\mathbf{R533})$ Any subscript, substring starting point, or substring ending point shall be an initialization expression.				
38 39 40	The <i>data-stmt-object-list</i> is expanded to form a sequence of pointers and scalar variables, referred to as "sequence of variables" in subsequent text. A nonpointer array whose unqualified name appears					

40 as a *data-stmt-object* or *data-i-do-object* is equivalent to a complete sequence of its array elements in

array element order (6.2.2.2). An array section is equivalent to the sequence of its array elements in
 array element order. A *data-implied-do* is expanded to form a sequence of array elements and structure

3 components, under the control of the *data-i-do-variable*, as in the DO construct (8.1.6.4).

4 The data-stmt-value-list is expanded to form a sequence of data-stmt-constants. A data-stmt-repeat

5 indicates the number of times the following *data-stmt-constant* is to be included in the sequence; omission
6 of a *data-stmt-repeat* has the effect of a repeat factor of 1.

7 A zero-sized array or a *data-implied-do* with an iteration count of zero contributes no variables to the 8 expanded sequence of variables, but a zero-length scalar character variable does contribute a variable 9 to the expanded sequence. A *data-stmt-constant* with a repeat factor of zero contributes no *data-stmt-*

10 *constants* to the expanded sequence of scalar *data-stmt-constants*.

 $\label{eq:constants} \mbox{ 11 } \mbox{ The expanded sequences of variables and } data-stmt-constants \mbox{ are in one-to-one correspondence. Each } \mbox{ Each } \mbox{ 11 } \mbox{ Constants } \mbox{ are in one-to-one correspondence. Each } \mbox{ 11 } \mbox{ Constants } \mbox{ are in one-to-one correspondence. Each } \mbox{ Constants } \mbox{ Const$

12 *data-stmt-constant* specifies the initial value or *null-init* for the corresponding variable. The lengths of 13 the two expanded sequences shall be the same.

A data-stmt-constant shall be null-init if and only if the corresponding data-stmt-object has the POINTER attribute. The initial association status of a pointer data-stmt-object is disassociated.

A data-stmt-constant other than null-init shall be compatible with its corresponding variable according
to the rules of intrinsic assignment (7.4.1.2). The variable is initially defined with the value specified by
the data-stmt-constant; if necessary, the value is converted according to the rules of intrinsic assignment
(7.4.1.3) to a value that agrees in type, type parameters, and shape with the variable.

20 If a *data-stmt-constant* is a *boz-literal-constant*, the corresponding variable shall be of type integer. The

21 boz-literal-constant is treated as if it were an *int-literal-constant* with a kind-param that specifies the

22 representation method with the largest decimal exponent range supported by the processor.

NOTE 5.23

Examples of DATA statements are:

CHARACTER (LEN = 10) NAME INTEGER, DIMENSION (0:9) :: MILES REAL, DIMENSION (100, 100) :: SKEW TYPE (NODE), POINTER :: HEAD_OF_LIST TYPE (PERSON) MYNAME, YOURNAME DATA NAME / 'JOHN DOE' /, MILES / 10 * 0 / DATA ((SKEW (K, J), J = 1, K), K = 1, 100) / 5050 * 0.0 / DATA ((SKEW (K, J), J = K + 1, 100), K = 1, 99) / 4950 * 1.0 / DATA HEAD_OF_LIST / NULL() / DATA MYNAME / PERSON (21, 'JOHN SMITH') / DATA YOURNAME % AGE, YOURNAME % NAME / 35, 'FRED BROWN' /

The character variable NAME is initialized with the value JOHN DOE with padding on the right because the length of the constant is less than the length of the variable. All ten elements of the integer array MILES are initialized to zero. The two-dimensional array SKEW is initialized so that the lower triangle of SKEW is zero and the strict upper triangle is one. The structures MYNAME and YOURNAME are declared using the derived type PERSON from Note 4.18. The pointer HEAD_OF_LIST is declared using the derived type NODE from Note 4.37; it is initially disassociated. MYNAME is initialized by a structure constructor. YOURNAME is initialized by supplying a separate value for each component.

1 5.4.6 DIMENSION statement

dimension-stmt

2 R534

is DIMENSION [:::] array-name (array-spec) ■ ■ [, array-name (array-spec)]...

4 This statement specifies the DIMENSION attribute (5.3.6) for a list of objects.

NOTE 5.24

An example of a DIMENSION statement is:

DIMENSION A (10), B (10, 70), C (:)

5 5.4.7 INTENT statement

6 R535 intent-stmt is INTENT (intent-spec) [::] dummy-arg-name-list

7 This statement specifies the INTENT attribute (5.3.8) for the dummy arguments in the list.

NOTE 5.25

An example of an INTENT statement is: SUBROUTINE EX (A, B) INTENT (INOUT) :: A, B

8 5.4.8 OPTIONAL statement

9 R536 optional-stmt is OPTIONAL [::] dummy-arg-name-list

10 This statement specifies the OPTIONAL attribute (5.3.10) for the dummy arguments in the list.

NOTE 5.26

An example of an OPTIONAL statement is: SUBROUTINE EX (A, B) OPTIONAL :: B

11 5.4.9 PARAMETER statement

12 The **PARAMETER statement** specifies the PARAMETER attribute (5.3.11) and the values for the 13 named constants in the list.

14	R537	parameter-stmt	\mathbf{is}	PARAMETER (named-constant-def-list)
15	R538	$named\mathchar`-def$	\mathbf{is}	named-constant = initialization-expr

16 If a named constant is defined by a PARAMETER statement, it shall not be subsequently declared to
17 have a type or type parameter value that differs from the type and type parameters it would have if
18 declared implicitly (5.5). A named array constant defined by a PARAMETER statement shall have its
19 shape specified in a prior specification statement.

The value of each named constant is that specified by the corresponding initialization expression; if necessary, the value is converted according to the rules of intrinsic assignment (7.4.1.3) to a value that agrees in type, type parameters, and shape with the named constant.

An example of a PARAMETER statement is:

PARAMETER (MODULUS = MOD (28, 3), NUMBER_OF_SENATORS = 100)

1 5.4.10 POINTER statement

2	R539	pointer- $stmt$	is	POINTER [::] pointer-decl-list
3	R540	pointer-decl	is	object-name [(deferred-shape-spec-list)]
4			or	proc-entity-name

5 This statement specifies the POINTER attribute (5.3.12) for a list of entities.

NOTE 5.28

An example of a POINTER statement is:

TYPE (NODE) :: CURRENT POINTER :: CURRENT, A (:, :)

6 5.4.11 PROTECTED statement

- 7 R541 protected-stmt is PROTECTED [::] entity-name-list
- 8 The **PROTECTED statement** specifies the PROTECTED attribute (5.3.13) for a list of entities.

9 5.4.12 SAVE statement

10	R542	save-stmt	\mathbf{is}	SAVE [[:::] saved-entity-list]
11	R543	saved- $entity$	\mathbf{is}	object-name
12			\mathbf{or}	proc- $pointer$ - $name$
13			\mathbf{or}	$/ \ common-block-name \ /$
14	R544	proc- $pointer$ - $name$	\mathbf{is}	name

15 C566 (R542) If a SAVE statement with an omitted saved entity list appears in a scoping unit, no
 other appearance of the SAVE *attr-spec* or SAVE statement is permitted in that scoping unit.

A SAVE statement with a saved entity list specifies the SAVE attribute (5.3.14) for a list of entities. A
SAVE statement without a saved entity list is treated as though it contained the names of all allowed
items in the same scoping unit.

NOTE 5.29

An example of a SAVE statement is:

SAVE A, B, C, / BLOCKA /, D

20 5.4.13 TARGET statement

- 21 R545 target-stmt
- 21 F

- TARGET [::] object-name [(array-spec)] \blacksquare \blacksquare [, object-name [(array-spec)]] ...
- 23 This statement specifies the TARGET attribute (5.3.15) for a list of objects.

is

An example of a TARGET statement is:

TARGET :: A (1000, 1000), B

1 5.4.14 VALUE statement

- 2 R546 value-stmt is VALUE [::] dummy-arg-name-list
- 3 The VALUE statement specifies the VALUE attribute (5.3.16) for a list of dummy arguments.

4 5.4.15 VOLATILE statement

- 5 R547 volatile-stmt is VOLATILE [::] object-name-list
- 6 The VOLATILE statement specifies the VOLATILE attribute (5.3.17) for a list of objects.

7 5.5 IMPLICIT statement

8 In a scoping unit, an IMPLICIT statement specifies a type, and possibly type parameters, for all
9 implicitly typed data entities whose names begin with one of the letters specified in the statement.
10 Alternatively, it may indicate that no implicit typing rules are to apply in a particular scoping unit.

11	R548	implicit- $stmt$	\mathbf{is}	IMPLICIT <i>implicit-spec-list</i>
12			\mathbf{or}	IMPLICIT NONE
13	R549	implicit-spec	\mathbf{is}	declaration-type-spec (letter-spec-list)
14	R550	letter-spec	\mathbf{is}	letter [-letter]

- C567 (R548) If IMPLICIT NONE is specified in a scoping unit, it shall precede any PARAMETER
 statements that appear in the scoping unit and there shall be no other IMPLICIT statements
 in the scoping unit.
- 18 C568 (R550) If the minus and second *letter* appear, the second letter shall follow the first letter
 alphabetically.

A *letter-spec* consisting of two *letters* separated by a minus is equivalent to writing a list containing all of the letters in alphabetical order in the alphabetic sequence from the first letter through the second letter. For example, A–C is equivalent to A, B, C. The same letter shall not appear as a single letter, or be included in a range of letters, more than once in all of the IMPLICIT statements in a scoping unit.

In each scoping unit, there is a mapping, which may be null, between each of the letters A, B, ..., Z and a type (and type parameters). An IMPLICIT statement specifies the mapping for the letters in its *letter-spec-list*. IMPLICIT NONE specifies the null mapping for all the letters. If a mapping is not specified for a letter, the default for a program unit or an interface body is default integer if the letter is I, J, ..., or N and default real otherwise, and the default for an internal or module procedure is the mapping in the host scoping unit.

Any data entity that is not explicitly declared by a type declaration statement, is not an intrinsic function, and is not made accessible by use association or host association is declared implicitly to be of the type (and type parameters) mapped from the first letter of its name, provided the mapping is not null. The mapping for the first letter of the data entity shall either have been established by a prior HIMPLICIT statement or be the default mapping for the letter. The mapping may be to a derived type that is inaccessible in the local scope if the derived type is accessible in the host scope. The data entity is treated as if it were declared in an explicit type declaration in the outermost scoping unit in which it

- 1 appears. An explicit type specification in a FUNCTION statement overrides an IMPLICIT statement
- $2 \ \ \, {\rm for \ the \ name \ of \ the \ result \ variable \ \, of \ that \ \, {\rm function \ \, subprogram.}}$

```
The following are examples of the use of IMPLICIT statements:
MODULE EXAMPLE_MODULE
   IMPLICIT NONE
   . . .
  INTERFACE
      FUNCTION FUN (I) ! Not all data entities need to
        INTEGER FUN ! be declared explicitly
     END FUNCTION FUN
  END INTERFACE
CONTAINS
  FUNCTION JFUN (J)
                       ! All data entities need to
     INTEGER JFUN, J ! be declared explicitly.
      . . .
  END FUNCTION JFUN
END MODULE EXAMPLE_MODULE
SUBROUTINE SUB
   IMPLICIT COMPLEX (C)
  C = (3.0, 2.0) ! C is implicitly declared COMPLEX
   . . .
CONTAINS
  SUBROUTINE SUB1
     IMPLICIT INTEGER (A, C)
     C = (0.0, 0.0) ! C is host associated and of
                      ! type complex
     Z = 1.0
                     ! Z is implicitly declared REAL
      A = 2
                     ! A is implicitly declared INTEGER
     CC = 1
                     ! CC is implicitly declared INTEGER
      . . .
  END SUBROUTINE SUB1
  SUBROUTINE SUB2
     Z = 2.0
                      ! Z is implicitly declared REAL and
                      ! is different from the variable of
                      ! the same name in SUB1
      . . .
  END SUBROUTINE SUB2
  SUBROUTINE SUB3
      USE EXAMPLE_MODULE ! Accesses integer function FUN
                          ! by use association
      Q = FUN (K)
                          ! Q is implicitly declared REAL and
                          ! K is implicitly declared INTEGER
      . . .
  END SUBROUTINE SUB3
END SUBROUTINE SUB
```

NOTE 5.32

The following is an example of a mapping to a derived type that is inaccessible in the local scope:

PROGRAM MAIN IMPLICIT TYPE(BLOB) (A) NOTE 5.32 (cont.)

```
TYPE BLOB
INTEGER :: I
END TYPE BLOB
TYPE(BLOB) :: B
CALL STEVE
CONTAINS
SUBROUTINE STEVE
INTEGER :: BLOB
..
AA = B
..
END SUBROUTINE STEVE
END PROGRAM MAIN
```

In the subroutine STEVE, it is not possible to explicitly declare a variable to be of type BLOB because BLOB has been given a different meaning, but implicit mapping for the letter A still maps to type BLOB, so AA is of type BLOB.

1 5.6 NAMELIST statement

2 A **NAMELIST statement** specifies a group of named data objects, which may be referred to by a 3 single name for the purpose of data transfer (9.5, 10.10).

 4
 R551
 namelist-stmt
 is
 NAMELIST ■

 5
 □
 / namelist-group-name / namelist-group-object-list ■

 6
 □
 [[,]] / namelist-group-name / ■

 7
 □
 namelist-group-object-list] ...

- 8 C569 (R551) The *namelist-group-name* shall not be a name accessed by use association.
- 9 R552 namelist-group-object is variable-name
- 10 C570 (R552) A namelist-group-object shall not be an assumed-size array.

C571 (R551) A namelist-group-object shall not have the PRIVATE attribute if the namelist-group name has the PUBLIC attribute.

13 The order in which the variables are specified in the NAMELIST statement determines the order in14 which the values appear on output.

15 Any namelist-group-name may occur more than once in the NAMELIST statements in a scoping unit.

16 The namelist-group-object-list following each successive appearance of the same namelist-group-name in

17 a scoping unit is treated as a continuation of the list for that *namelist-group-name*.

18 $\,$ A namelist group object may be a member of more than one namelist group.

A namelist group object shall either be accessed by use or host association or shall have its type, type parameters, and shape specified by previous specification statements or the procedure heading in the same scoping unit or by the implicit typing rules in effect for the scoping unit. If a namelist group object is typed by the implicit typing rules, its appearance in any subsequent type declaration statement shall

23 confirm the implied type and type parameters.

An example of a NAMELIST statement is:

NAMELIST /NLIST/ A, B, C

1 5.7 Storage association of data objects

2 5.7.1 EQUIVALENCE statement

3 5.7.1.1 General

4 An **EQUIVALENCE statement** is used to specify the sharing of storage units by two or more objects 5 in a scoping unit. This causes storage association (16.4.3) of the objects that share the storage units.

6 If the equivalenced objects have differing type or type parameters, the EQUIVALENCE statement does
7 not cause type conversion or imply mathematical equivalence. If a scalar and an array are equivalenced,
8 the scalar does not have array properties and the array does not have the properties of a scalar.

9 10 11 12 13	R553 R554 R555	equivalence-stmt equivalence-set equivalence-object	 is EQUIVALENCE equivalence-set-list is (equivalence-object, equivalence-object-list) is variable-name or array-element or substring 				
14 15 16 17 18 19	C572	(R555) An <i>equivalence-object</i> shall not be a designator with a base object that is a dummy argument, a pointer, an allocatable variable, a derived-type object that has an allocatable ultimate component, an object of a nonsequence derived type, an object of a derived type that has a pointer at any level of component selection, an automatic object, a function name, an entry name, a result name, a variable with the BIND attribute, a variable in a common block that has the BIND attribute, or a named constant.					
20	C573	(R555) An <i>equivalence-object</i> shall not be a designator that has more than one <i>part-ref</i> .					
21	C574	(R555) An <i>equivalence-object</i> shall not have the TARGET attribute.					
22 23	C575	(R555) Each subscript or substring range expression in an <i>equivalence-object</i> shall be an integer initialization expression (7.1.7).					
24 25 26	C576	(R554) If an <i>equivalence-object</i> is of type default integer, default real, double precision real, default complex, default logical, or numeric sequence type, all of the objects in the equivalence set shall be of these types.					
27 28	C577	(R554) If an <i>equivalence-object</i> is of type default character or character sequence type, all of the objects in the equivalence set shall be of these types.					
29 30 31	C578	(R554) If an <i>equivalence-object</i> is of a sequence derived type that is not a numeric sequence or character sequence type, all of the objects in the equivalence set shall be of the same type with the same type parameter values.					
32 33 34	C579	double precision real, default	ect is of an intrinsic type other than default integer, default real complex, default logical, or default character, all of the objects in of the same type with the same kind type parameter value.				
35	C580	(R555) If an equivalence-obje	ct has the PROTECTED attribute, all of the objects in the equiv-				

- 1 alence set shall have the PROTECTED attribute.
- 2 C581 (R555) The name of an *equivalence-object* shall not be a name made accessible by use association.
- 3 C582 (R555) A substring shall not have length zero.

The EQUIVALENCE statement allows the equivalencing of sequence structures and the equivalencing of objects of intrinsic type with nondefault type parameters, but there are strict rules regarding the appearance of these objects in an EQUIVALENCE statement.

A structure that appears in an EQUIVALENCE statement shall be a sequence structure. If a sequence structure is not of numeric sequence type or of character sequence type, it shall be equivalenced only to objects of the same type with the same type parameter values.

A structure of a numeric sequence type shall be equivalenced only to another structure of a numeric sequence type, an object of default integer type, default real type, double precision real type, default complex type, or default logical type such that components of the structure ultimately become associated only with objects of these types.

A structure of a character sequence type shall be equivalenced only to an object of default character type or another structure of a character sequence type.

An object of intrinsic type with nondefault kind type parameters shall not be equivalenced to objects of different type or kind type parameters.

Further rules on the interaction of EQUIVALENCE statements and default initialization are given in 16.4.3.3.

4 5.7.1.2 Equivalence association

5 An EQUIVALENCE statement specifies that the storage sequences (16.4.3.1) of the data objects specified
6 in an *equivalence-set* are storage associated. All of the nonzero-sized sequences in the *equivalence-set*, if

7 any, have the same first storage unit, and all of the zero-sized sequences in the *equivalence-set*, if any,

8 are storage associated with one another and with the first storage unit of any nonzero-sized sequences.

9 This causes the storage association of the data objects in the *equivalence-set* and may cause storage10 association of other data objects.

10 association of other data objects.

11 5.7.1.3 Equivalence of default character objects

A data object of type default character shall not be equivalenced to an object that is not of type default
character and not of a character sequence type. The lengths of the equivalenced character objects need
not be the same.

An EQUIVALENCE statement specifies that the storage sequences of all the default character data objects specified in an *equivalence-set* are storage associated. All of the nonzero-sized sequences in the *equivalence-set*, if any, have the same first character storage unit, and all of the zero-sized sequences in the *equivalence-set*, if any, are storage associated with one another and with the first character storage unit of any nonzero-sized sequences. This causes the storage association of the data objects in the *equivalence-set* and may cause storage association of other data objects.

NOTE 5.35

For example, using the declarations: CHARACTER (LEN = 4) :: A, B CHARACTER (LEN = 3) :: C (2)

NOTE 5.35 (cont.)

EQUIVAL	ENCE (A,	C (1)), (B, (C (2))			
the assoc	iation of	А, В, а	and C car	n be illust	rated g	raphically as:	
					-		
1	2	3	4	5	6	7	
	A		4 				
				B			
	C(1)			C(2)			

1 5.7.1.4 Array names and array element designators

2 For a nonzero-sized array, the use of the array name unqualified by a subscript list as an equivalence-

3 *object* has the same effect as using an array element designator that identifies the first element of the 4 array.

5 5.7.1.5 Restrictions on EQUIVALENCE statements

6 An EQUIVALENCE statement shall not specify that the same storage unit is to occur more than once7 in a storage sequence.

NOTE 5.36

For example: REAL, DIMENSION (2) :: A REAL :: B EQUIVALENCE (A (1), B), (A (2), B) ! Not standard conforming is prohibited, because it would specify the same storage unit for A (1) and A (2).

8 An EQUIVALENCE statement shall not specify that consecutive storage units are to be nonconsecutive.

NOTE 5.37

```
For example, the following is prohibited:

REAL A (2)

DOUBLE PRECISION D (2)

EQUIVALENCE (A (1), D (1)), (A (2), D (2)) ! Not standard conforming
```

9 5.7.2 COMMON statement

10 5.7.2.1 General

11 The **COMMON statement** specifies blocks of physical storage, called **common blocks**, that can be 12 accessed by any of the scoping units in a program. Thus, the COMMON statement provides a global 13 data facility based on storage association (16.4.3).

The common blocks specified by the COMMON statement may be named and are called named com mon blocks, or may be unnamed and are called blank common.

16	R556	common-stmt	is	COMMON ■
17				$\blacksquare \ [\ / \ [\ common-block-name \] \ / \] \ common-block-object-list \blacksquare$
18				$\blacksquare [[,]/[common-block-name]/\blacksquare$

1 2 3	R557	common-block-objectiscommon-block-object-list]isvariable-name [(array-spec)]orproc-pointer-name			
4	C583	(R557) An array-spec in a common-block-object shall be an explicit-shape-spec-list.			
5 6	C584	(R557) Only one appearance of a given <i>variable-name</i> or <i>proc-pointer-name</i> is permitted in all <i>common-block-object-lists</i> within a scoping unit.			
7 8 9	C585	(R557) A <i>common-block-object</i> shall not be a dummy argument, an allocatable variable, a derived-type object with an ultimate component that is allocatable, an automatic object, a function name, an entry name, a variable with the BIND attribute, or a result name.			
10 11	C586	(R557) If a <i>common-block-object</i> is of a derived type, it shall be a sequence type $(4.5.2)$ or a type with the BIND attribute and it shall have no default initialization.			
12 13	C587	(R557) A variable-name or proc-pointer-name shall not be a name made accessible by use association.			
14 15		COMMON statement, the data objects whose names appear in a common block object list ag a common block name are declared to be in that common block. If the first common block			

15 following a common block name are declared to be in that common block. If the first common block 16 name is omitted, all data objects whose names appear in the first common block object list are specified to 17 be in blank common. Alternatively, the appearance of two slashes with no common block name between 18 them declares the data objects whose names appear in the common block object list that follows to be 19 in blank common.

Any common block name or an omitted common block name for blank common may occur more than once in one or more COMMON statements in a scoping unit. The common block list following each successive appearance of the same common block name in a scoping unit is treated as a continuation of the list for that common block name. Similarly, each blank common block object list in a scoping unit is treated as a continuation of blank common.

25 The form *variable-name* (*array-spec*) specifies the DIMENSION attribute for that variable.

26 If derived-type objects of numeric sequence type (4.5.2) or character sequence type (4.5.2) appear in
27 common, it is as if the individual components were enumerated directly in the common list.

NOTE 5.38

Examples of COMMON statements are: COMMON /BLOCKA/ A, B, D (10, 30) COMMON I, J, K

28 5.7.2.2 Common block storage sequence

29 For each common block in a scoping unit, a **common block storage sequence** is formed as follows:

- A storage sequence is formed consisting of the sequence of storage units in the storage sequences (16.4.3.1) of all data objects in the common block object lists for the common block. The order of the storage sequences is the same as the order of the appearance of the common block object lists in the scoping unit.
- 34 (2) The storage sequence formed in (1) is extended to include all storage units of any storage
 35 sequence associated with it by equivalence association. The sequence shall be extended only
 36 by adding storage units beyond the last storage unit. Data objects associated with an entity
 37 in a common block are considered to be in that common block.

Only COMMON statements and EQUIVALENCE statements appearing in the scoping unit contribute
 to common block storage sequences formed in that scoping unit.

3 5.7.2.3 Size of a common block

4 The size of a common block is the size of its common block storage sequence, including any extensions
5 of the sequence resulting from equivalence association.

6 5.7.2.4 Common association

Within a program, the common block storage sequences of all nonzero-sized common blocks with the 7 same name have the same first storage unit, and the common block storage sequences of all zero-sized 8 common blocks with the same name are storage associated with one another. Within a program, the 9 10 common block storage sequences of all nonzero-sized blank common blocks have the same first storage unit and the storage sequences of all zero-sized blank common blocks are associated with one another and 11 with the first storage unit of any nonzero-sized blank common blocks. This results in the association of 12 objects in different scoping units. Use association or host association may cause these associated objects 13 14 to be accessible in the same scoping unit.

A nonpointer object of default integer type, default real type, double precision real type, default complex
type, default logical type, or numeric sequence type shall be associated only with nonpointer objects of
these types.

18 A nonpointer object of type default character or character sequence type shall be associated only with19 nonpointer objects of these types.

A nonpointer object of a derived type that is not a numeric sequence or character sequence type shall
be associated only with nonpointer objects of the same type with the same type parameter values.

A nonpointer object of intrinsic type other than default integer, default real, double precision real,
 default complex, default logical, or default character shall be associated only with nonpointer objects of

24 $\,$ the same type and type parameters.

A data pointer shall be storage associated only with data pointers of the same type and rank. Data pointers that are storage associated shall have deferred the same type parameters; corresponding nondeferred type parameters shall have the same value. A procedure pointer shall be storage associated only with another procedure pointer; either both interfaces shall be explicit or both interfaces shall be implicit. If the interfaces are explicit, the characteristics shall be the same. If the interfaces are implicit, either both shall be subroutines or both shall be functions with the same type and type parameters.

An object with the TARGET attribute shall be storage associated only with another object that has the TARGET attribute and the same type and type parameters.

NOTE 5.39

A common block is permitted to contain sequences of different storage units, provided each scoping unit that accesses the common block specifies an identical sequence of storage units for the common block. For example, this allows a single common block to contain both numeric and character storage units.

Association in different scoping units between objects of default type, objects of double precision real type, and sequence structures is permitted according to the rules for equivalence objects (5.7.1).

1 5.7.2.5 Differences between named common and blank common

2 A blank common block has the same properties as a named common block, except for the following.

- 3 (1) Execution of a RETURN or END statement may cause data objects in a named common 4 block to become undefined unless the common block has the SAVE attribute, but never 5 causes data objects in blank common to become undefined (16.5.6).
- 6 (2) Named common blocks of the same name shall be of the same size in all scoping units of a 7 program in which they appear, but blank common blocks may be of different sizes.
- 8 (3) A data object in a named common block may be initially defined by means of a DATA
 9 statement or type declaration statement in a block data program unit (11.3), but objects in
 10 blank common shall not be initially defined.

11 5.7.3 Restrictions on common and equivalence

12 An EQUIVALENCE statement shall not cause the storage sequences of two different common blocks to13 be associated.

14 Equivalence association shall not cause a derived-type object with default initialization to be associated

15 $\,$ with an object in a common block.

16 Equivalence association shall not cause a common block storage sequence to be extended by adding 17 storage units preceding the first storage unit of the first object specified in a COMMON statement for 19 the common block

18 the common block.

NOTE 5.40

For example, the following is not permitted: COMMON /X/ A REAL B (2) EQUIVALENCE (A, B (2)) ! Not standard conforming