

WD 1539-1

J3/16-007r1  
(F2015 Working Document)

1st May 2016 16:07

This is an internal working document of J3 and WG5.

**NOTE:** This working document is only available as a PDF file.

This page intentionally left nonblank.

# Contents

Foreword . . . . .	xvii
Introduction . . . . .	xviii
1 Overview . . . . .	1
1.1 Scope . . . . .	1
1.2 Normative references . . . . .	1
1.3 Terms and definitions . . . . .	2
1.4 Notation, symbols and abbreviated terms . . . . .	22
1.4.1 Syntax rules . . . . .	22
1.4.2 Constraints . . . . .	23
1.4.3 Assumed syntax rules . . . . .	24
1.4.4 Syntax conventions and characteristics . . . . .	24
1.4.5 Text conventions . . . . .	24
1.5 Conformance . . . . .	24
1.6 Compatibility . . . . .	25
1.6.1 Previous Fortran standards . . . . .	25
1.6.2 New intrinsic procedures . . . . .	26
1.6.3 Fortran 2008 compatibility . . . . .	26
1.6.4 Fortran 2003 compatibility . . . . .	26
1.6.5 Fortran 95 compatibility . . . . .	27
1.6.6 Fortran 90 compatibility . . . . .	27
1.6.7 FORTRAN 77 compatibility . . . . .	27
1.7 Deleted and obsolescent features . . . . .	28
1.7.1 General . . . . .	28
1.7.2 Nature of deleted features . . . . .	28
1.7.3 Nature of obsolescent features . . . . .	28
2 Fortran concepts . . . . .	29
2.1 High level syntax . . . . .	29
2.2 Program unit concepts . . . . .	32
2.2.1 Program units and scoping units . . . . .	32
2.2.2 Program . . . . .	32
2.2.3 Procedure . . . . .	32
2.2.4 Module . . . . .	33
2.2.5 Submodule . . . . .	33
2.3 Execution concepts . . . . .	33
2.3.1 Statement classification . . . . .	33
2.3.2 Statement order . . . . .	33
2.3.3 The END statement . . . . .	34
2.3.4 Program execution . . . . .	34
2.3.5 Execution sequence . . . . .	35
2.3.6 Image execution states . . . . .	35
2.3.7 Termination of execution . . . . .	36
2.4 Data concepts . . . . .	36
2.4.1 Type . . . . .	36
2.4.2 Data value . . . . .	37
2.4.3 Data entity . . . . .	37

2.4.4	Definition of objects and pointers . . . . .	38
2.4.5	Reference . . . . .	39
2.4.6	Array . . . . .	39
2.4.7	Coarray . . . . .	39
2.4.8	Established coarrays . . . . .	40
2.4.9	Pointer . . . . .	40
2.4.10	Allocatable variables . . . . .	40
2.4.11	Storage . . . . .	40
2.5	Fundamental concepts . . . . .	41
2.5.1	Names and designators . . . . .	41
2.5.2	Statement keyword . . . . .	41
2.5.3	Other keywords . . . . .	41
2.5.4	Association . . . . .	41
2.5.5	Intrinsic . . . . .	41
2.5.6	Operator . . . . .	41
2.5.7	Companion processors . . . . .	41
3	Lexical tokens and source form . . . . .	43
3.1	Processor character set . . . . .	43
3.1.1	Characters . . . . .	43
3.1.2	Letters . . . . .	43
3.1.3	Digits . . . . .	43
3.1.4	Underscore . . . . .	43
3.1.5	Special characters . . . . .	43
3.1.6	Other characters . . . . .	44
3.2	Low-level syntax . . . . .	44
3.2.1	Tokens . . . . .	44
3.2.2	Names . . . . .	44
3.2.3	Constants . . . . .	45
3.2.4	Operators . . . . .	45
3.2.5	Statement labels . . . . .	46
3.2.6	Delimiters . . . . .	46
3.3	Source form . . . . .	47
3.3.1	Program units, statements, and lines . . . . .	47
3.3.2	Free source form . . . . .	47
3.3.3	Fixed source form . . . . .	48
3.4	Including source text . . . . .	49
4	Types . . . . .	51
4.1	Characteristics of types . . . . .	51
4.1.1	The concept of type . . . . .	51
4.1.2	Type classification . . . . .	51
4.1.3	Set of values . . . . .	51
4.1.4	Constants . . . . .	51
4.1.5	Operations . . . . .	51
4.2	Type parameters . . . . .	52
4.3	Types, type specifiers, and values . . . . .	52
4.3.1	Relationship of types and values to objects . . . . .	52
4.3.2	Type specifiers and type compatibility . . . . .	53
4.4	Intrinsic types . . . . .	55
4.4.1	Classification and specification . . . . .	55
4.4.2	Intrinsic operations on intrinsic types . . . . .	55
4.4.3	Numeric intrinsic types . . . . .	55
4.4.4	Character type . . . . .	59
4.4.5	Logical type . . . . .	62
4.5	Derived types . . . . .	62

4.5.1	Derived type concepts . . . . .	62
4.5.2	Derived-type definition . . . . .	63
4.5.3	Derived-type parameters . . . . .	66
4.5.4	Components . . . . .	68
4.5.5	Type-bound procedures . . . . .	75
4.5.6	Final subroutines . . . . .	77
4.5.7	Type extension . . . . .	79
4.5.8	Derived-type values . . . . .	81
4.5.9	Derived-type specifier . . . . .	81
4.5.10	Construction of derived-type values . . . . .	82
4.5.11	Derived-type operations and assignment . . . . .	84
4.6	Enumerations and enumerators . . . . .	84
4.7	Binary, octal, and hexadecimal literal constants . . . . .	86
4.8	Construction of array values . . . . .	86
5	Attribute declarations and specifications . . . . .	89
5.1	Attributes of procedures and data objects . . . . .	89
5.2	Type declaration statement . . . . .	89
5.3	Automatic data objects . . . . .	91
5.4	Initialization . . . . .	91
5.5	Attributes . . . . .	91
5.5.1	Attribute specification . . . . .	91
5.5.2	Accessibility attribute . . . . .	92
5.5.3	ALLOCATABLE attribute . . . . .	92
5.5.4	ASYNCHRONOUS attribute . . . . .	92
5.5.5	BIND attribute for data entities . . . . .	93
5.5.6	CODIMENSION attribute . . . . .	93
5.5.7	CONTIGUOUS attribute . . . . .	95
5.5.8	DIMENSION attribute . . . . .	96
5.5.9	EXTERNAL attribute . . . . .	99
5.5.10	INTENT attribute . . . . .	99
5.5.11	INTRINSIC attribute . . . . .	101
5.5.12	OPTIONAL attribute . . . . .	102
5.5.13	PARAMETER attribute . . . . .	102
5.5.14	POINTER attribute . . . . .	102
5.5.15	PROTECTED attribute . . . . .	103
5.5.16	SAVE attribute . . . . .	103
5.5.17	TARGET attribute . . . . .	104
5.5.18	VALUE attribute . . . . .	104
5.5.19	VOLATILE attribute . . . . .	104
5.6	Attribute specification statements . . . . .	105
5.6.1	Accessibility statement . . . . .	105
5.6.2	ALLOCATABLE statement . . . . .	106
5.6.3	ASYNCHRONOUS statement . . . . .	106
5.6.4	BIND statement . . . . .	106
5.6.5	CODIMENSION statement . . . . .	107
5.6.6	CONTIGUOUS statement . . . . .	107
5.6.7	DATA statement . . . . .	107
5.6.8	DIMENSION statement . . . . .	109
5.6.9	INTENT statement . . . . .	110
5.6.10	OPTIONAL statement . . . . .	110
5.6.11	PARAMETER statement . . . . .	110
5.6.12	POINTER statement . . . . .	110
5.6.13	PROTECTED statement . . . . .	111
5.6.14	SAVE statement . . . . .	111
5.6.15	TARGET statement . . . . .	111

5.6.16	VALUE statement . . . . .	112
5.6.17	VOLATILE statement . . . . .	112
5.7	IMPLICIT statement . . . . .	112
5.8	NAMelist statement . . . . .	115
5.9	Storage association of data objects . . . . .	115
5.9.1	EQUIVALENCE statement . . . . .	115
5.9.2	COMMON statement . . . . .	117
5.9.3	Restrictions on common and equivalence . . . . .	119
6	Use of data objects . . . . .	121
6.1	Designator . . . . .	121
6.2	Variable . . . . .	121
6.3	Constants . . . . .	122
6.4	Scalars . . . . .	122
6.4.1	Substrings . . . . .	122
6.4.2	Structure components . . . . .	122
6.4.3	Coindexed named objects . . . . .	124
6.4.4	Complex parts . . . . .	124
6.4.5	Type parameter inquiry . . . . .	124
6.5	Arrays . . . . .	125
6.5.1	Order of reference . . . . .	125
6.5.2	Whole arrays . . . . .	125
6.5.3	Array elements and array sections . . . . .	125
6.5.4	Simply contiguous array designators . . . . .	128
6.6	Image selectors . . . . .	129
6.7	Dynamic association . . . . .	130
6.7.1	ALLOCATE statement . . . . .	130
6.7.2	NULLIFY statement . . . . .	134
6.7.3	DEALLOCATE statement . . . . .	134
6.7.4	STAT= specifier . . . . .	136
6.7.5	ERRMSG= specifier . . . . .	137
7	Expressions and assignment . . . . .	139
7.1	Expressions . . . . .	139
7.1.1	Expression semantics . . . . .	139
7.1.2	Form of an expression . . . . .	139
7.1.3	Precedence of operators . . . . .	143
7.1.4	Evaluation of operations . . . . .	145
7.1.5	Intrinsic operations . . . . .	145
7.1.6	Defined operations . . . . .	152
7.1.7	Evaluation of operands . . . . .	153
7.1.8	Integrity of parentheses . . . . .	154
7.1.9	Type, type parameters, and shape of an expression . . . . .	154
7.1.10	Conformability rules for elemental operations . . . . .	155
7.1.11	Specification expression . . . . .	156
7.1.12	Constant expression . . . . .	157
7.2	Assignment . . . . .	159
7.2.1	Assignment statement . . . . .	159
7.2.2	Pointer assignment . . . . .	163
7.2.3	Masked array assignment – WHERE . . . . .	167
7.2.4	FORALL . . . . .	169
8	Execution control . . . . .	173
8.1	Executable constructs containing blocks . . . . .	173
8.1.1	Blocks . . . . .	173
8.1.2	Rules governing blocks . . . . .	173

8.1.3	ASSOCIATE construct . . . . .	174
8.1.4	BLOCK construct . . . . .	175
8.1.5	CHANGE TEAM construct . . . . .	177
8.1.6	CRITICAL construct . . . . .	179
8.1.7	DO construct . . . . .	180
8.1.8	IF construct and statement . . . . .	187
8.1.9	SELECT CASE construct . . . . .	188
8.1.10	SELECT RANK construct . . . . .	191
8.1.11	SELECT TYPE construct . . . . .	193
8.1.12	EXIT statement . . . . .	196
8.2	Branching . . . . .	196
8.2.1	Branch concepts . . . . .	196
8.2.2	GO TO statement . . . . .	196
8.2.3	Computed GO TO statement . . . . .	196
8.3	CONTINUE statement . . . . .	197
8.4	STOP and ERROR STOP statements . . . . .	197
8.5	FAIL IMAGE statement . . . . .	197
8.6	Image execution control . . . . .	198
8.6.1	Image control statements . . . . .	198
8.6.2	Segments . . . . .	199
8.6.3	SYNC ALL statement . . . . .	200
8.6.4	SYNC IMAGES statement . . . . .	201
8.6.5	SYNC MEMORY statement . . . . .	203
8.6.6	SYNC TEAM statement . . . . .	204
8.6.7	EVENT POST statement . . . . .	205
8.6.8	EVENT WAIT statement . . . . .	205
8.6.9	FORM TEAM statement . . . . .	206
8.6.10	LOCK and UNLOCK statements . . . . .	207
8.6.11	STAT= and ERRMSG= specifiers in image control statements . . . . .	208
9	Input/output statements . . . . .	211
9.1	Input/output concepts . . . . .	211
9.2	Records . . . . .	211
9.2.1	Definition of a record . . . . .	211
9.2.2	Formatted record . . . . .	211
9.2.3	Unformatted record . . . . .	211
9.2.4	Endfile record . . . . .	212
9.3	External files . . . . .	212
9.3.1	External file concepts . . . . .	212
9.3.2	File existence . . . . .	212
9.3.3	File access . . . . .	213
9.3.4	File position . . . . .	215
9.3.5	File storage units . . . . .	216
9.4	Internal files . . . . .	217
9.5	File connection . . . . .	217
9.5.1	Referring to a file . . . . .	217
9.5.2	Connection modes . . . . .	218
9.5.3	Unit existence . . . . .	219
9.5.4	Connection of a file to a unit . . . . .	219
9.5.5	Preconnection . . . . .	220
9.5.6	OPEN statement . . . . .	220
9.5.7	CLOSE statement . . . . .	224
9.6	Data transfer statements . . . . .	225
9.6.1	Form of input and output statements . . . . .	225
9.6.2	Control information list . . . . .	226
9.6.3	Data transfer input/output list . . . . .	230

9.6.4	Execution of a data transfer input/output statement . . . . .	233
9.6.5	Termination of data transfer statements . . . . .	243
9.7	Waiting on pending data transfer . . . . .	244
9.7.1	Wait operation . . . . .	244
9.7.2	WAIT statement . . . . .	244
9.8	File positioning statements . . . . .	245
9.8.1	Syntax . . . . .	245
9.8.2	BACKSPACE statement . . . . .	245
9.8.3	ENDFILE statement . . . . .	246
9.8.4	REWIND statement . . . . .	246
9.9	FLUSH statement . . . . .	247
9.10	File inquiry statement . . . . .	247
9.10.1	Forms of the INQUIRE statement . . . . .	247
9.10.2	Inquiry specifiers . . . . .	248
9.10.3	Inquire by output list . . . . .	254
9.11	Error, end-of-record, and end-of-file conditions . . . . .	254
9.11.1	Occurrence of input/output conditions . . . . .	254
9.11.2	Error conditions and the ERR= specifier . . . . .	254
9.11.3	End-of-file condition and the END= specifier . . . . .	255
9.11.4	End-of-record condition and the EOR= specifier . . . . .	255
9.11.5	IOSTAT= specifier . . . . .	256
9.11.6	IOMSG= specifier . . . . .	256
9.12	Restrictions on input/output statements . . . . .	256
10	Input/output editing . . . . .	259
10.1	Format specifications . . . . .	259
10.2	Explicit format specification methods . . . . .	259
10.2.1	FORMAT statement . . . . .	259
10.2.2	Character format specification . . . . .	259
10.3	Form of a format item list . . . . .	260
10.3.1	Syntax . . . . .	260
10.3.2	Edit descriptors . . . . .	260
10.3.3	Fields . . . . .	262
10.4	Interaction between input/output list and format . . . . .	262
10.5	Positioning by format control . . . . .	264
10.6	Decimal symbol . . . . .	264
10.7	Data edit descriptors . . . . .	264
10.7.1	Purpose of data edit descriptors . . . . .	264
10.7.2	Numeric editing . . . . .	265
10.7.3	Logical editing . . . . .	272
10.7.4	Character editing . . . . .	272
10.7.5	Generalized editing . . . . .	273
10.7.6	User-defined derived-type editing . . . . .	274
10.8	Control edit descriptors . . . . .	274
10.8.1	Position edit descriptors . . . . .	274
10.8.2	Slash editing . . . . .	275
10.8.3	Colon editing . . . . .	275
10.8.4	SS, SP, and S editing . . . . .	276
10.8.5	P editing . . . . .	276
10.8.6	BN and BZ editing . . . . .	276
10.8.7	RU, RD, RZ, RN, RC, and RP editing . . . . .	276
10.8.8	DC and DP editing . . . . .	277
10.9	Character string edit descriptors . . . . .	277
10.10	List-directed formatting . . . . .	277
10.10.1	Purpose of list-directed formatting . . . . .	277
10.10.2	Values and value separators . . . . .	277

10.10.3	List-directed input . . . . .	278
10.10.4	List-directed output . . . . .	280
10.11	Namelist formatting . . . . .	281
10.11.1	Purpose of namelist formatting . . . . .	281
10.11.2	Name-value subsequences . . . . .	281
10.11.3	Namelist input . . . . .	282
10.11.4	Namelist output . . . . .	285
11	Program units . . . . .	287
11.1	Main program . . . . .	287
11.2	Modules . . . . .	287
11.2.1	Module syntax and semantics . . . . .	287
11.2.2	The USE statement and use association . . . . .	288
11.2.3	Submodules . . . . .	291
11.3	Block data program units . . . . .	291
12	Procedures . . . . .	293
12.1	Concepts . . . . .	293
12.2	Procedure classifications . . . . .	293
12.2.1	Procedure classification by reference . . . . .	293
12.2.2	Procedure classification by means of definition . . . . .	293
12.3	Characteristics . . . . .	294
12.3.1	Characteristics of procedures . . . . .	294
12.3.2	Characteristics of dummy arguments . . . . .	294
12.3.3	Characteristics of function results . . . . .	294
12.4	Procedure interface . . . . .	295
12.4.1	Interface and abstract interface . . . . .	295
12.4.2	Implicit and explicit interfaces . . . . .	295
12.4.3	Specification of the procedure interface . . . . .	296
12.5	Procedure reference . . . . .	307
12.5.1	Syntax of a procedure reference . . . . .	307
12.5.2	Actual arguments, dummy arguments, and argument association . . . . .	309
12.5.3	Function reference . . . . .	320
12.5.4	Subroutine reference . . . . .	320
12.5.5	Resolving named procedure references . . . . .	320
12.5.6	Resolving type-bound procedure references . . . . .	323
12.6	Procedure definition . . . . .	323
12.6.1	Intrinsic procedure definition . . . . .	323
12.6.2	Procedures defined by subprograms . . . . .	323
12.6.3	Definition and invocation of procedures by means other than Fortran . . . . .	329
12.6.4	Statement function . . . . .	329
12.7	Pure procedures . . . . .	330
12.8	Elemental procedures . . . . .	332
12.8.1	Elemental procedure declaration and interface . . . . .	332
12.8.2	Elemental function actual arguments and results . . . . .	332
12.8.3	Elemental subroutine actual arguments . . . . .	332
13	Intrinsic procedures and modules . . . . .	333
13.1	Classes of intrinsic procedures . . . . .	333
13.2	Arguments to intrinsic procedures . . . . .	333
13.2.1	General rules . . . . .	333
13.2.2	The shape of array arguments . . . . .	334
13.2.3	Mask arguments . . . . .	334
13.2.4	DIM arguments and reduction functions . . . . .	334
13.3	Bit model . . . . .	334
13.3.1	General . . . . .	334

13.3.2	Bit sequence comparisons . . . . .	335
13.3.3	Bit sequences as arguments to INT and REAL . . . . .	335
13.4	Numeric models . . . . .	335
13.5	Atomic subroutines . . . . .	336
13.6	Collective subroutines . . . . .	337
13.7	Standard generic intrinsic procedures . . . . .	338
13.8	Specific names for standard intrinsic functions . . . . .	344
13.9	Specifications of the standard intrinsic procedures . . . . .	345
13.9.1	General . . . . .	345
13.10	Standard intrinsic modules . . . . .	432
13.10.1	General . . . . .	432
13.10.2	The ISO_FORTRAN_ENV intrinsic module . . . . .	433
14	Exceptions and IEEE arithmetic . . . . .	439
14.1	Overview of IEEE arithmetic support . . . . .	439
14.2	Derived types and constants defined in the modules . . . . .	440
14.3	The exceptions . . . . .	440
14.4	The rounding modes . . . . .	442
14.5	Underflow mode . . . . .	443
14.6	Halting . . . . .	443
14.7	The floating-point modes and status . . . . .	443
14.8	Exceptional values . . . . .	443
14.9	IEEE arithmetic . . . . .	444
14.10	Summary of the procedures . . . . .	445
14.11	Specifications of the procedures . . . . .	447
14.11.1	General . . . . .	447
14.12	Examples . . . . .	469
15	Interoperability with C . . . . .	473
15.1	General . . . . .	473
15.2	The ISO_C_BINDING intrinsic module . . . . .	473
15.2.1	Summary of contents . . . . .	473
15.2.2	Named constants and derived types in the module . . . . .	473
15.2.3	Procedures in the module . . . . .	474
15.3	Interoperability between Fortran and C entities . . . . .	479
15.3.1	General . . . . .	479
15.3.2	Interoperability of intrinsic types . . . . .	479
15.3.3	Interoperability with C pointer types . . . . .	480
15.3.4	Interoperability of derived types and C struct types . . . . .	480
15.3.5	Interoperability of scalar variables . . . . .	482
15.3.6	Interoperability of array variables . . . . .	482
15.3.7	Interoperability of procedures and procedure interfaces . . . . .	483
15.4	C descriptors . . . . .	485
15.5	The source file ISO_Fortran_binding.h . . . . .	486
15.5.1	Summary of contents . . . . .	486
15.5.2	The CFI_dim_t structure type . . . . .	486
15.5.3	The CFI_cdesc_t structure type . . . . .	486
15.5.4	Macros and typedefs in ISO_Fortran_binding.h . . . . .	487
15.5.5	Functions declared in ISO_Fortran_binding.h . . . . .	490
15.6	Restrictions on C descriptors . . . . .	497
15.7	Restrictions on formal parameters . . . . .	497
15.8	Restrictions on lifetimes . . . . .	498
15.9	Interoperation with C global variables . . . . .	498
15.9.1	General . . . . .	498
15.9.2	Binding labels for common blocks and variables . . . . .	499
15.10	Interoperation with C functions . . . . .	500

15.10.1	Definition and reference of interoperable procedures . . . . .	500
15.10.2	Binding labels for procedures . . . . .	500
15.10.3	Exceptions and IEEE arithmetic procedures . . . . .	501
15.10.4	Asynchronous communication . . . . .	501
16	Scope, association, and definition . . . . .	503
16.1	Scopes, identifiers, and entities . . . . .	503
16.2	Global identifiers . . . . .	503
16.3	Local identifiers . . . . .	504
16.3.1	Classes of local identifiers . . . . .	504
16.3.2	Local identifiers that are the same as common block names . . . . .	505
16.3.3	Function results . . . . .	505
16.3.4	Components, type parameters, and bindings . . . . .	505
16.3.5	Argument keywords . . . . .	505
16.4	Statement and construct entities . . . . .	506
16.5	Association . . . . .	507
16.5.1	Name association . . . . .	507
16.5.2	Pointer association . . . . .	511
16.5.3	Storage association . . . . .	514
16.5.4	Inheritance association . . . . .	516
16.5.5	Establishing associations . . . . .	516
16.6	Definition and undefinition of variables . . . . .	517
16.6.1	Definition of objects and subobjects . . . . .	517
16.6.2	Variables that are always defined . . . . .	517
16.6.3	Variables that are initially defined . . . . .	517
16.6.4	Variables that are initially undefined . . . . .	518
16.6.5	Events that cause variables to become defined . . . . .	518
16.6.6	Events that cause variables to become undefined . . . . .	520
16.6.7	Variable definition context . . . . .	522
16.6.8	Pointer association context . . . . .	523
Annex A	(informative) Processor Dependencies . . . . .	525
A.1	Unspecified Items . . . . .	525
A.2	Processor Dependencies . . . . .	525
Annex B	(informative) Deleted and obsolescent features . . . . .	531
B.1	Deleted features from Fortran 90 . . . . .	531
B.2	Deleted features from Fortran 2008 . . . . .	532
B.3	Obsolescent features . . . . .	532
B.3.1	General . . . . .	532
B.3.2	Alternate return . . . . .	532
B.3.3	Computed GO TO statement . . . . .	533
B.3.4	Statement functions . . . . .	533
B.3.5	DATA statements among executables . . . . .	533
B.3.6	Assumed character length functions . . . . .	533
B.3.7	Fixed form source . . . . .	533
B.3.8	CHARACTER* form of CHARACTER declaration . . . . .	534
B.3.9	ENTRY statements . . . . .	534
B.3.10	Label DO statement . . . . .	534
B.3.11	COMMON and EQUIVALENCE statements and the block data program unit . . . . .	534
B.3.12	Specific names for intrinsic functions . . . . .	534
B.3.13	FORALL construct and statement . . . . .	534
Annex C	(informative) Extended notes . . . . .	535
C.1	Clause 4 notes . . . . .	535
C.1.1	Selection of the approximation methods (4.4.3.2) . . . . .	535

C.1.2	Type extension and component accessibility (4.5.2.2, 4.5.4)	535
C.1.3	Generic type-bound procedures (4.5.5)	536
C.1.4	Abstract types (4.5.7.1)	537
C.1.5	Structure constructors and generic names (4.5.10)	538
C.1.6	Final subroutines (4.5.6, 4.5.6.2, 4.5.6.3, 4.5.6.4)	540
C.2	Clause 5 notes: The VOLATILE attribute (5.5.19)	542
C.3	Clause 6 notes	542
C.3.1	Structure components (6.4.2)	542
C.3.2	Allocation with dynamic type (6.7.1)	544
C.4	Clause 7 notes	544
C.4.1	Evaluation of function references (7.1.7)	544
C.4.2	Pointers in expressions (7.1.9.2)	544
C.4.3	Pointers in variable definition contexts (7.2.1.3, 16.6.7)	545
C.5	Clause 8 notes	545
C.5.1	The SELECT CASE construct (8.1.9)	545
C.5.2	Loop control (8.1.7)	545
C.5.3	Examples of DO constructs (8.1.7)	545
C.5.4	Examples of invalid DO constructs (8.1.7)	546
C.5.5	Simple example using events	547
C.5.6	Example using three teams	548
C.5.7	Accessing coarrays in sibling teams	549
C.5.8	Example involving failed images	550
C.5.9	EVENT_QUERY example that tolerates image failure	552
C.6	Clause 9 notes	554
C.6.1	External files (9.3)	554
C.6.2	Nonadvancing input/output (9.3.4.2)	555
C.6.3	OPEN statement (9.5.6)	556
C.6.4	Connection properties (9.5.4)	558
C.6.5	Asynchronous input/output (9.6.2.5)	558
C.7	Clause 10 notes	559
C.7.1	Number of records (10.4, 10.5, 10.8.2)	559
C.7.2	List-directed input (10.10.3)	560
C.8	Clause 11 notes	561
C.8.1	Main program and block data program unit (11.1, 11.3)	561
C.8.2	Dependent compilation (11.2)	561
C.8.3	Examples of the use of modules (11.2.1)	563
C.8.4	Modules with submodules (11.2.3)	569
C.9	Clause 12 notes	573
C.9.1	Portability problems with external procedures (12.4.3.6)	573
C.9.2	Procedures defined by means other than Fortran (12.6.3)	574
C.9.3	Abstract interfaces and procedure pointer components (12.4, 4.5)	574
C.9.4	Pointers and targets as arguments (12.5.2.4, 12.5.2.6, 12.5.2.7)	576
C.9.5	Polymorphic Argument Association (12.5.2.9)	577
C.9.6	Rules ensuring unambiguous generics (12.4.3.5.5)	578
C.10	Clause 13 notes	583
C.10.1	Atomic memory consistency	583
C.10.2	EVENT_QUERY example	584
C.11	Clause 13 notes	586
C.11.1	Collective subroutine examples	586
C.11.2	Atomic memory consistency	587
C.12	Clause 15 notes	589
C.12.1	Runtime environments (15.1)	589
C.12.2	Example of Fortran calling C (15.3)	589
C.12.3	Example of C calling Fortran (15.3)	591
C.12.4	Example of calling C functions with noninteroperable data (15.10)	592

C.12.5	Example of opaque communication between C and Fortran (15.3)	593
C.12.6	Using assumed type to interoperate with C	594
C.12.7	Using assumed-type variables in Fortran	596
C.12.8	Simplifying interfaces for arbitrary rank procedures	597
C.12.9	Processing assumed-shape arrays in C	598
C.12.10	Creating a contiguous copy of an array	599
C.12.11	Changing the attributes of an array	601
C.12.12	Creating an array section in C using CFI_section	602
C.12.13	Use of CFI_setpointer	604
C.12.14	Mapping of MPI interfaces to Fortran	605
C.13	Clause 16 notes : Examples of host association (16.5.1.4)	607
Index		609

(Blank page)

## List of Tables

1.3	Previous editions of the Fortran International Standard . . . . .	26
2.1	Requirements on statement ordering . . . . .	33
2.2	Statements allowed in scoping units . . . . .	34
3.1	Special characters . . . . .	44
3.2	Adjacent keywords where separating blanks are optional . . . . .	47
6.1	Subscript order value . . . . .	126
7.1	Categories of operations and relative precedence . . . . .	143
7.2	Type of operands and results for intrinsic operators . . . . .	146
7.3	Interpretation of the numeric intrinsic operators . . . . .	147
7.4	Interpretation of the character intrinsic operator <code>//</code> . . . . .	149
7.5	Interpretation of the logical intrinsic operators . . . . .	150
7.6	The values of operations involving logical intrinsic operators . . . . .	150
7.7	Interpretation of the relational intrinsic operators . . . . .	151
7.8	Type conformance for the intrinsic assignment statement . . . . .	159
7.9	Numeric conversion and the assignment statement . . . . .	161
10.1	E and D exponent forms . . . . .	268
10.2	EN exponent forms . . . . .	269
10.3	ES exponent forms . . . . .	269
13.1	Standard generic intrinsic procedure summary . . . . .	339
13.2	Unrestricted specific intrinsic functions . . . . .	344
13.3	Restricted specific intrinsic functions . . . . .	345
13.4	Default BOUNDARY values for EOSHIFT . . . . .	371
13.5	Characteristics of the result of NULL ( ) . . . . .	407
14.1	IEEE_ARITHMETIC module procedure summary . . . . .	445
14.2	IEEE_EXCEPTIONS module procedure summary . . . . .	446

15.1 Names of C characters with special semantics . . . . .	474
15.2 Interoperability between Fortran and C types . . . . .	479
15.3 ISO_Fortran_binding.h macros for attribute codes . . . . .	488
15.4 ISO_Fortran_binding.h macros for type codes . . . . .	488
15.5 ISO_Fortran_binding.h macros for error codes . . . . .	489

## Foreword

- 1 ISO (the International Organization for Standardization) and IEC (the International Electrotechnical Commission) form the specialized system for worldwide standardization. National bodies that are members of ISO or IEC participate in the development of International Standards through technical committees established by the respective organization to deal with particular fields of technical activity. ISO and IEC technical committees collaborate in fields of mutual interest. Other international organizations, governmental and nongovernmental, in liaison with ISO and IEC, also take part in the work. In the field of information technology, ISO and IEC have established a joint technical committee, ISO/IEC JTC 1.
- 2 International Standards are drafted in accordance with the rules given in the ISO/IEC Directives, Part 2.
- 3 The main task of the joint technical committee is to prepare International Standards. Draft International Standards adopted by the joint technical committee are circulated to national bodies for voting. Publication as an International Standard requires approval by at least 75 % of the national bodies casting a vote.
- 4 Attention is drawn to the possibility that some of the elements of this document may be the subject of patent rights. ISO and IEC shall not be held responsible for identifying any or all such patent rights.
- 5 ISO/IEC 1539-1 was prepared by Joint Technical Committee ISO/IEC JTC 1, *Information technology*, Subcommittee SC 22, *Programming languages, their environments and system software interfaces*.
- 6 This fourth edition cancels and replaces the third edition (ISO/IEC 1539-1:2010), which has been technically revised. It also incorporates the Technical Corrigenda ISO/IEC 1539-1:2010/Cor. 1:2012, ISO/IEC 1539-1:2010/Cor. 2:2013, and ISO/IEC 1539-1:2010/Cor. 3:2014, and the Technical Specifications ISO/IEC TS 29113:2012 and ISO/IEC TS 18508:2015.
- 7 ISO/IEC 1539 consists of the following parts, under the general title *Information technology — Programming languages — Fortran*:
  - 8 — *Part 1: Base language*
  - 9 — *Part 2: Varying length character strings*

## Introduction

- 1 This part of ISO/IEC 1539 comprises the specification of the base Fortran language, informally known as Fortran 2015. With the limitations noted in 1.6.4, the syntax and semantics of Fortran 2008 are contained entirely within Fortran 2015. Therefore, any standard-conforming Fortran 2008 program not affected by such limitations is a standard-conforming Fortran 2015 program. New features of Fortran 2015 can be compatibly incorporated into such Fortran 2008 programs, with any exceptions indicated in the text of this part of ISO/IEC 1539.
- 2 Fortran 2015 contains several extensions to Fortran 2008; these are listed below.
  - Data declaration:  
Constant properties of an object declared in its *entity-decl* can be used in its *initialization*.
  - Data usage and computation:  
Labeled **DO** loops have been redundant since Fortran 90 and are now specified to be *obsolescent*. The *arithmetic IF statement* has been deleted. The **EQUIVALENCE** and **COMMON** statements and the *block data program unit* have been redundant since Fortran 90 and are now specified to be *obsolescent*. The *nonblock DO construct* has been deleted. **FORALL** is now specified to be *obsolescent*. The type and kind of an implied DO variable in an *array constructor* or *DATA statement* can be specified within the constructor or statement. The locality of a variable used in a *DO CONCURRENT construct* can be explicitly specified. The **SELECT RANK construct** provides structured access to the elements of an *assumed-rank array*.
  - Input/output:  
The *SIZE= specifier* can be used with advancing input. It is no longer prohibited to open a file on more than one unit. The value assigned by the *RECL= specifier* in an *INQUIRE statement* has been standardized. The values assigned by the *POS=* and *SIZE=* specifiers in an *INQUIRE statement* for a unit that has pending asynchronous operations have been standardized. The *G0.d edit descriptor* can be used for list items of type *Integer*, *Logical*, and *Character*. The *D*, *E*, *EN*, and *ES* edit descriptors can have a field width of zero, analogous to the *F* edit descriptor. The exponent width *e* in a *data edit descriptor* can be zero, analogous to a field width of zero. Floating-point formatted input accepts hexadecimal-significand numbers that conform to ISO/IEC/IEEE 60559:2011. The *EX edit descriptor* provides hexadecimal-significand formatted output conforming to ISO/IEC/IEEE 60559:2011. An error condition occurs if unacceptable characters are presented for logical or numeric editing during execution of a formatted input statement.
  - Execution control:  
The stop code in a **STOP** or **ERROR STOP** statement can be nonconstant. Output of the stop code and exception summary from the **STOP** and **ERROR STOP** statements can be controlled.
  - Intrinsic procedures and modules:  
In references to the intrinsic functions **ALL**, **ANY**, **FINDLOC**, **IALL**, **IANY**, **IPARITY**, **MAXLOC**, **MAXVAL**, **MINLOC**, **MINVAL**, **NORM2**, **PARITY**, **PRODUCT**, **SUM**, and **THIS\_IMAGE**, the *actual argument* for **DIM** can be a present optional *dummy argument*. In a reference of the intrinsic function **CMPLX** with an actual argument of type complex, no keyword is needed for a **KIND** argument. The new intrinsic function **COSHAPE** returns the coshape of a *coarray*. The new intrinsic function **OUT\_OF\_RANGE** tests whether a numeric value can be safely converted to a different type or kind. The new intrinsic subroutine **RANDOM\_INIT** establishes the initial state of the pseudorandom number generator used by **RANDOM\_NUMBER**. The new intrinsic function **REDUCE** performs user-specified array reductions. A processor is required to report use of a *nonstandard intrinsic* procedure, use of a *nonstandard intrinsic* module, and use of a non-standard procedure from a *standard intrinsic* module. Integer and logical arguments to intrinsic procedures and intrinsic module procedures that were previously required to be of default kind no longer have that requirement, except for **RANDOM\_SEED**. *Specific names for intrinsic functions* are now deemed *obsolescent*. All standard procedures in the intrinsic module **ISO\_C\_BINDING**, other than **C\_F\_POINTER**, are now *pure*. The arguments to the intrinsic function **SIGN** can be of different kind. Nonpolymorphic pointer arguments to the intrinsic functions **EXTENDS\_TYPE\_OF** and **SAME\_TYPE\_AS** need not have defined pointer association status. The effects of invoking the intrinsic procedures **COMMAND\_ARGUMENT\_COUNT**, **GET\_COMMAND**, and **GET\_COMMAND\_ARGUMENT**, on *images* other than *image* one, are no longer processor dependent. Access to error messages from the intrinsic subroutines **GET\_COMMAND**, **GET\_COMMAND\_ARGUMENT**, and **GET\_ENVIRONMENT\_VARIABLE** is provided by an optional **ERRMSG** argument.

- Program units and procedures:

The **IMPORT statement** can appear in a contained subprogram or **BLOCK construct**, and can restrict access via host association; diagnosis of violation of the **IMPORT** restrictions is required. The **GENERIC statement** can be used to declare generic interfaces. The **ERROR STOP statement** can appear in a **pure subprogram**. The number of procedure arguments is used in **generic resolution**. In a module, the **default accessibility** of entities accessed from another module can be controlled separately from the default accessibility of entities declared in the using module. An **IMPLICIT NONE statement** can require explicit declaration of the **EXTERNAL attribute** throughout a **scoping unit** and its contained **scoping units**. A defined operation need not specify **INTENT (IN)** for a **dummy argument** with the **VALUE attribute**. A defined assignment need not specify **INTENT (IN)** for the second **dummy argument** if it has the **VALUE attribute**. Procedures, including elemental procedures, can be invoked recursively by default; the **RECURSIVE** keyword is advisory only. The **NON\_RECURSIVE** keyword specifies that a procedure is not recursive. A **dummy argument** of a **pure function** is permitted in a variable definition context, if it has the **VALUE attribute**.

- Features previously described by ISO/IEC TS 29113:2012:

A **dummy data object** can **assume its rank** from its **effective argument**. A **dummy data object** can **assume the type** from its **effective argument**, without having the ability to perform type selection. An **interoperable** procedure can have **dummy arguments** that are **assumed-type** and/or **assumed-rank**. An **interoperable** procedure can have **dummy data objects** that are **allocatable**, **assumed-shape**, **optional**, or **pointers**. The **character length** of a **dummy data object** of an **interoperable** procedure can be assumed. The argument to **C\_LOC** can be a noninteroperable array. The FPTR argument to **C\_F\_POINTER** can be a noninteroperable array pointer. The argument to **C\_FUNLOC** can be a noninteroperable procedure. The FPTR argument to **C\_F\_PROCPONTER** can be a noninteroperable procedure pointer.

- Changes to the intrinsic modules IEEE\_ARITHMETIC, IEEE\_EXCEPTIONS, and IEEE\_FEATURES for conformance with ISO/IEC/IEEE 60559:2011:

There is a new, optional, rounding mode **IEEE\_AWAY**. The new type **IEEE\_MODES\_TYPE** encapsulates all floating-point modes. Features associated with subnormal numbers can be accessed with functions and types named ...SUBNORMAL... (the old ...DENORMAL... names remain). The standard intrinsic relational operations on IEEE numbers provide the compareSignalingrelation operations. The new function **IEEE\_FMA** performs fused multiply-add operations. The function **IEEE\_INT** performs rounded conversions to integer type. The new functions **IEEE\_MAX\_NUM**, **IEEE\_MAX\_NUM\_MAG**, **IEEE\_MIN\_NUM**, and **IEEE\_MIN\_NUM\_MAG** calculate maximum and minimum numeric values. The new functions **IEEE\_NEXT\_DOWN** and **IEEE\_NEXT\_UP** return the adjacent machine numbers. The new functions **IEEE\_QUIET\_EQ**, **IEEE\_QUIET\_GE**, **IEEE\_QUIET\_GT**, **IEEE\_QUIET\_LE**, **IEEE\_QUIET\_LT**, and **IEEE\_QUIET\_NE** perform quiet comparisons. The decimal rounding mode can be inquired and set independently of the binary rounding mode, using the RADIX argument to **IEEE\_GET\_ROUNDING\_MODE** and **IEEE\_SET\_ROUNDING\_MODE**. The new function **IEEE\_REAL** performs rounded conversions to real type. The function **IEEE\_REM** now requires its arguments to have the same radix. The function **IEEE\_RINT** now has a **ROUND** argument to perform specific rounding. The new function **IEEE\_SIGNBIT** tests the sign bit of an IEEE number.

- Features previously described by ISO/IEC TS 18508:2015:

The **CRITICAL statement** has optional **ERRMSG=** and **STAT=** specifiers. The intrinsic subroutines **ATOMIC\_DEFINE** and **ATOMIC\_REF** have an optional **STAT** argument. The new intrinsic subroutines **ATOMIC\_ADD**, **ATOMIC\_AND**, **ATOMIC\_CAS**, **ATOMIC\_FETCH\_ADD**, **ATOMIC\_FETCH\_AND**, **ATOMIC\_FETCH\_OR**, **ATOMIC\_FETCH\_XOR**, **ATOMIC\_OR**, and **ATOMIC\_XOR** perform atomic operations. The new intrinsic functions **FAILED\_IMAGES** and **STOPPED\_IMAGES** return images known to have **failed** or **stopped** respectively. The new intrinsic function **IMAGE\_STATUS** returns the image execution status of an image. The intrinsic subroutine **MOVE\_ALLOC** has optional **ERRMSG** and **STAT** arguments. The intrinsic functions **IMAGE\_INDEX** and **NUM\_IMAGES** have additional forms with a **TEAM** or **TEAM\_NUMBER** argument. The intrinsic function **THIS\_IMAGE** has an optional **TEAM** argument. The **EVENT POST** and **EVENT WAIT** statements, the intrinsic subroutine **EVENT\_QUERY**, and the type **EVENT\_TYPE** provide an event facility for one-sided segment ordering. The **CHANGE TEAM construct**, derived type **TEAM\_TYPE**, **FORM TEAM** and **SYNC TEAM** statements, and the **STAT=** and **TEAM=** specifiers on image selectors, provide a team facility for a subset of the program's images to act in concert as if it were the set of all images. The new intrinsic subroutines **CO\_BROADCAST**, **CO\_MAX**, **CO\_MIN**, **CO\_REDUCE**, and **CO\_SUM** perform collective reduction operations on the images.

of the [current team](#). The concept of failed images, the [FAIL IMAGE statement](#), and the named constant [STAT\\_FAILED\\_IMAGE](#) provide support for fault-tolerant parallel execution.

- 3 This part of ISO/IEC 1539 is organized in 16 clauses, dealing with 8 conceptual areas. These 8 areas, and the clauses in which they are treated, are:

High/low level concepts	Clauses <a href="#">1</a> , <a href="#">2</a> , <a href="#">3</a>
Data concepts	Clauses <a href="#">4</a> , <a href="#">5</a> , <a href="#">6</a>
Computations	Clauses <a href="#">7</a> , <a href="#">13</a> , <a href="#">14</a>
Execution control	Clause <a href="#">8</a>
Input/output	Clauses <a href="#">9</a> , <a href="#">10</a>
Program units	Clauses <a href="#">11</a> , <a href="#">12</a>
Interoperability with C	Clause <a href="#">15</a>
Scoping and association rules	Clause <a href="#">16</a>

- 4 It also contains the following nonnormative material:

Processor dependencies	Annex <a href="#">A</a>
Deleted and obsolescent features	Annex <a href="#">B</a>
Extended notes	Annex <a href="#">C</a>
Index	Index

# Information technology — Programming languages — Fortran —

## Part 1: Base language

### 1 Overview

#### 1.1 Scope

- 1 This part of ISO/IEC 1539 specifies the form and establishes the interpretation of programs expressed in the base Fortran language. The purpose of this part of ISO/IEC 1539 is to promote portability, reliability, maintainability, and efficient execution of Fortran programs for use on a variety of computing systems.
- 2 This part of ISO/IEC 1539 specifies
  - the forms that a program written in the Fortran language may take,
  - the rules for interpreting the meaning of a program and its data,
  - the form of the input data to be processed by such a program, and
  - the form of the output data resulting from the use of such a program.
- 3 Except where stated otherwise, requirements and prohibitions specified by this part of ISO/IEC 1539 apply to programs rather than processors.
- 4 This part of ISO/IEC 1539 does not specify
  - the mechanism by which programs are transformed for use on computing systems,
  - the operations required for setup and control of the use of programs on computing systems,
  - the method of transcription of programs or their input or output data to or from a storage medium,
  - the program and processor behavior when this part of ISO/IEC 1539 fails to establish an interpretation except for the processor detection and reporting requirements in items (2) to (10) of 1.5,
  - the maximum number of [images](#), or the size or complexity of a program and its data that will exceed the capacity of any particular computing system or the capability of a particular processor,
  - the mechanism for determining the number of [images](#) of a program,
  - the physical properties of an [image](#) or the relationship between [images](#) and the computational elements of a computing system,
  - the physical properties of the representation of quantities and the method of rounding, approximating, or computing numeric values on a particular processor, except by reference to ISO/IEC/IEEE 60559:2011 under conditions specified in Clause 14,
  - the physical properties of input/output records, files, and units, or
  - the physical properties and implementation of storage.

#### 1.2 Normative references

- 1 The following referenced standards are indispensable for the application of this document. For dated references, only the edition cited applies. For undated references, the latest edition of the referenced document (including any amendments) applies.

ISO/IEC 646:1991 (International Reference Version), *Information technology—ISO 7-bit coded character set for information interchange*

ISO/IEC 9899:2011, *Programming languages—C*

ISO/IEC 10646, *Information technology—Universal Multiple-Octet Coded Character Set (UCS)*

ISO/IEC/IEEE 60559:2011, *Binary floating-point arithmetic for microprocessor systems*

## 1.3 Terms and definitions

1 For the purposes of this document, the following terms and definitions apply.

### 1.3.1

#### **abstract interface**

set of procedure characteristics with dummy argument names ([12.4.1](#))

### 1.3.2

#### **actual argument**

entity ([R1225](#)) that appears in a [procedure reference](#)

### 1.3.3

#### **allocatable**

having the [ALLOCATABLE attribute](#) ([5.5.3](#))

### 1.3.4

#### **array**

set of scalar data, all of the same type and [type parameters](#), whose individual elements are arranged in a rectangular pattern

#### **1.3.4.1**

##### **array element**

scalar individual element of an array

#### **1.3.4.2**

##### **array pointer**

array with the [POINTER attribute](#) ([5.5.14](#))

#### **1.3.4.3**

##### **array section**

array [subobject](#) designated by [array-section](#), and which is itself an array ([6.5.3.3](#))

#### **1.3.4.4**

##### **assumed-shape array**

nonallocatable nonpointer [dummy argument](#) array that takes its shape from its [effective argument](#) ([5.5.8.3](#))

#### **1.3.4.5**

##### **assumed-size array**

[dummy argument](#) array whose size is assumed from that of its [effective argument](#) ([5.5.8.5](#))

#### **1.3.4.6**

##### **deferred-shape array**

[allocatable](#) array or array pointer, declared with a [deferred-shape-spec-list](#) ([5.5.8.4](#))

#### **1.3.4.7**

##### **explicit-shape array**

array declared with an [explicit-shape-spec-list](#), which specifies explicit values for the [bounds](#) in each dimension of the array ([5.5.8.2](#))

**1.3.5****ASCII character**

character whose representation method corresponds to ISO/IEC 646:1991 (International Reference Version)

**1.3.6****associate name**

name of [construct entity](#) associated with a selector of an [ASSOCIATE](#), [SELECT RANK](#), or [SELECT TYPE](#) construct, or associated with a *coselector-name* of a [CHANGE TEAM](#) construct ([8.1.3](#), [8.1.5](#), [8.1.10](#), [8.1.11](#))

**1.3.7****associating entity**

(in a dynamically-established association) the entity that did not exist prior to the establishment of the association ([16.5.5](#))

**1.3.8****association**

[inheritance association](#) ([16.5.4](#)), [name association](#) ([16.5.1](#)), [pointer association](#) ([16.5.2](#)), or [storage association](#) ([16.5.3](#)).

**1.3.8.1****argument association**

association between an [effective argument](#) and a [dummy argument](#) ([12.5.2](#))

**1.3.8.2****construct association**

association between a selector and an [associate name](#) in an [ASSOCIATE](#), [SELECT RANK](#), or [SELECT TYPE](#) construct, or between a coselector and an [associate name](#) in a [CHANGE TEAM](#) construct ([8.1.3](#), [8.1.5](#), [8.1.10](#), [8.1.11](#), [16.5.1.6](#))

**1.3.8.3****host association**

name association, other than argument association, between entities in a submodule or contained [scoping unit](#) and entities in its host ([16.5.1.4](#))

**1.3.8.4****inheritance association**

association between the [inherited](#) components of an [extended type](#) and the components of its [parent component](#)

**1.3.8.5****linkage association**

association between a variable or common block with the [BIND attribute](#) and a C global variable ([15.9](#), [16.5.1.5](#))

**1.3.8.6****name association**

[argument association](#), [construct association](#), [host association](#), [linkage association](#), or [use association](#) ([16.5.1](#))

**1.3.8.7****pointer association**

association between a [pointer](#) and an entity with the [TARGET attribute](#) ([16.5.2](#))

**1.3.8.8****storage association**

association between storage sequences ([16.5.3](#))

**1.3.8.9****use association**

association between entities in a module and entities in a [scoping unit](#) or construct that references that module, as specified by a [USE statement](#) ([11.2.2](#))

**1.3.9****assumed-rank dummy data object**

[dummy data object](#) that assumes the rank, shape, and size of its [effective argument](#)

**1.3.10****assumed-type**

declared with a `TYPE(*)` type specifier ([4.3.2](#))

**1.3.11****attribute**

property of an entity that determines its uses ([5.1](#))

**1.3.12****automatic data object**

nondummy [data object](#) with a [type parameter](#) or array [bound](#) that depends on the value of a [specification-expr](#) that is not a [constant expression](#)

**1.3.13****base object**

[⟨data-ref⟩](#) object designated by the leftmost *part-name* ([6.4.2](#))

**1.3.14****binding**

[type-bound procedure](#) or [final subroutine](#) ([4.5.5](#))

**1.3.15****binding name**

name given to a specific or generic [type-bound procedure](#) in the type definition ([4.5.5](#))

**1.3.16****binding label**

default character value specifying the name by which a global entity with the [BIND attribute](#) is known to the [companion processor](#) ([15.10.2](#), [15.9.2](#))

**1.3.17****block**

sequence of executable constructs formed by the syntactic class [block](#) and which is treated as a unit by the executable constructs described in [8.1](#)

**1.3.18****bound****array bound**

limit of a dimension of an [array](#)

**1.3.19****branch target statement**

[action-stmt](#), [associate-stmt](#), [end-associate-stmt](#), [if-then-stmt](#), [end-if-stmt](#), [select-case-stmt](#), [end-select-stmt](#), [select-rank-stmt](#), [end-select-rank-stmt](#), [select-type-stmt](#), [end-select-type-stmt](#), [do-stmt](#), [end-do-stmt](#), [block-stmt](#), [end-block-stmt](#), [critical-stmt](#), [end-critical-stmt](#), [forall-construct-stmt](#), or [where-construct-stmt](#) whose statement label appears as a *label* in a [GO TO](#) statement, computed [GO TO](#) statement, [alt-return-spec](#), [END=](#) specifier, [EOR=](#) specifier, or [ERR=](#) specifier ([8.2.1](#))

**1.3.20****C address**

value identifying the location of a [data object](#) or procedure either defined by the [companion processor](#) or which might be accessible to the [companion processor](#)

**NOTE 1.1**

This is the concept that ISO/IEC 9899:2011 calls the address.

**1.3.21****C descriptor**

C structure of type [CFI\\_cdesc\\_t](#) defined in the source file [ISO\\_Fortran\\_binding.h](#) ([15.4](#), [15.5](#))

**1.3.22****character context**

within a character literal constant ([4.4.4](#)) or within a character string edit descriptor ([10.3.2](#))

**1.3.23****characteristics**

⟨[dummy argument](#)⟩ being a [dummy data object](#), [dummy procedure](#), or an asterisk (alternate return indicator)

**1.3.24****characteristics**

⟨[dummy data object](#)⟩ properties listed in [12.3.2.2](#)

**1.3.25****characteristics**

⟨[dummy procedure](#) or [dummy procedure pointer](#)⟩ properties listed in [12.3.2.3](#)

**1.3.26****characteristics**

⟨function result⟩ properties listed in [12.3.3](#)

**1.3.27****characteristics**

⟨procedure⟩ properties listed in [12.3.1](#)

**1.3.28****coarray**

[data entity](#) that has nonzero corank ([2.4.7](#))

**1.3.28.1****established coarray**

[coarray](#) that is accessible using an *[image-selector](#)* ([2.4.7](#))

**1.3.29****cobound**

bound (limit) of a [codimension](#)

**1.3.30****codimension**

dimension of the pattern formed by a set of corresponding [coarrays](#)

**1.3.31****coindexed object**

[data object](#) whose [designator](#) includes an *[image-selector](#)* ([R624](#), [6.6](#))

**1.3.32****collating sequence**

one-to-one mapping from a character set into the nonnegative integers ([4.4.4.4](#))

**1.3.33****common block**

block of physical storage specified by a [COMMON statement](#) ([5.9.2](#))

**1.3.33.1****blank common**

unnamed common block

**1.3.34****companion processor**

processor-dependent mechanism by which global data and procedures may be referenced or defined ([2.5.7](#))

**1.3.35****component**

part of a derived type, or of an object of derived type, defined by a [component-def-stmt](#) ([4.5.4](#))

**1.3.35.1****direct component**

one of the components, or one of the direct components of a nonpointer nonallocatable component ([4.5.1](#))

**1.3.35.2****parent component**

component of an [extended type](#) whose type is that of the [parent type](#) and whose components are [inheritance associated](#) with the [inherited](#) components of the [parent type](#) ([4.5.7.2](#))

**1.3.35.3****potential subobject component**

nonpointer component, or potential subobject component of a nonpointer component ([4.5.1](#))

**1.3.35.4****subcomponent**

[structure](#) [direct component](#) that is a [subobject](#) of the [structure](#) ([6.4.2](#))

**1.3.35.5****ultimate component**

component that is of [intrinsic type](#), a pointer, or allocatable; or an ultimate component of a nonpointer nonallocatable component of derived type

**1.3.36****component order**

ordering of the nonparent components of a derived type that is used for [intrinsic](#) formatted input/output and [structure constructors](#) (where component keywords are not used) ([4.5.4.7](#))

**1.3.37****conformable**

[of two data entities](#) having the same shape, or one being an array and the other being scalar

**1.3.38****connected**

relationship between a [unit](#) and a file: each is connected if and only if the [unit](#) refers to the file ([9.5.4](#))

**1.3.39****constant**

[data object](#) that has a value and which cannot be defined, redefined, or become undefined during execution of a program ([3.2.3](#), [6.3](#))

**1.3.39.1****literal constant**

constant that does not have a name ([R305](#), [4.4](#))

**1.3.39.2****named constant**

named [data object](#) with the [PARAMETER attribute](#) ([5.5.13](#))

**1.3.40****construct entity**

entity whose identifier has the scope of a construct ([16.1](#), [16.4](#))

**1.3.41****constant expression**

expression satisfying the requirements specified in [7.1.12](#), thus ensuring that its value is constant

**1.3.42****contiguous**

⟨array⟩ having array elements in order that are not separated by other data objects, as specified in [5.5.7](#)

**1.3.43****contiguous**

⟨multi-part data object⟩ that the parts in order are not separated by other data objects

**1.3.44****corank**

number of [codimensions](#) of a [coarray](#) (zero for objects that are not coarrays)

**1.3.45****cosubscript**

([R625](#)) scalar integer expression in an *image-selector* ([R624](#))

**1.3.46****data entity**

[data object](#), result of the evaluation of an expression, or the result of the execution of a function reference

**1.3.47****data object****object**

constant ([4.1.4](#)), [variable](#) ([6](#)), or [subobject](#) of a constant ([2.4.3.2.3](#))

**1.3.48****decimal symbol**

character that separates the whole and fractional parts in the decimal representation of a real number in a file ([10.6](#))

**1.3.49****declaration**

specification of attributes for various program entities

**NOTE 1.2**

Often this involves specifying the type of a named [data object](#) or specifying the shape of a named array [object](#).

**1.3.50****default initialization**

mechanism for automatically initializing pointer components to have a defined pointer association status, and nonpointer components to have a particular value (4.5.4.6)

**1.3.51****default-initialized**

⟨subcomponent⟩ subject to a [default initialization](#) specified in the type definition for that [component](#) (4.5.4.6)

**1.3.52****definable**

capable of [definition](#) and permitted to become [defined](#)

**1.3.53****defined**

⟨data object⟩ has a valid value

**1.3.54****defined**

⟨pointer⟩ has a pointer association status of associated or [disassociated](#) (16.5.2.2)

**1.3.55****defined assignment**

assignment defined by a procedure (7.2.1.4, 12.4.3.5.3)

**1.3.56****defined input/output**

input/output defined by a procedure and accessed via a [defined-io-generic-spec](#) (R1209, 9.6.4.8)

**1.3.57****defined operation**

operation defined by a procedure (7.1.6.1, 12.4.3.5.2)

**1.3.58****definition**

⟨data object⟩ process by which the data object becomes defined (16.6.5)

**1.3.59****definition**

⟨derived type (4.5.2), enumeration (4.6), or procedure (12.6)⟩ specification of the type, enumeration, or procedure

**1.3.60****descendant**

⟨module or submodule⟩ submodule that extends that module or submodule or that extends another descendant thereof

**1.3.61****designator**

name followed by zero or more component selectors, complex part selectors, array section selectors, array element selectors, image selectors, and substring selectors (6.1)

**1.3.61.1****complex part designator**

designator that designates the real or imaginary part of a complex [data object](#), independently of the other part (6.4.4)

**1.3.61.2****object designator****data object designator**

designator for a data object

**NOTE 1.3**

An object name is a special case of an object designator.

**1.3.61.3****procedure designator**

designator for a procedure

**1.3.62****disassociated**

⟨pointer association⟩ pointer association status of not being associated with any target and not being undefined (16.5.2.2)

**1.3.63****disassociated**

⟨pointer⟩ has a pointer association status of disassociated

**1.3.64****dummy argument**

entity whose identifier appears in a dummy argument list (R1237) in a FUNCTION, SUBROUTINE, ENTRY, or statement function statement, or whose name can be used as an argument keyword in a reference to an intrinsic procedure or a procedure in an intrinsic module

**1.3.64.1****dummy data object**

dummy argument that is a data object

**1.3.64.2****dummy function**

dummy procedure that is a function

**1.3.65****effective argument**

entity that is argument-associated with a dummy argument (12.5.2.3)

**1.3.66****effective item**

scalar object resulting from the application of the rules in 9.6.3 to an input/output list

**1.3.67****elemental**

independent scalar application of an action or operation to elements of an array or corresponding elements of a set of conformable arrays and scalars, or possessing the capability of elemental operation

**NOTE 1.4**

Combination of scalar and array operands or arguments combine the scalar operand(s) with each element of the array operand(s).

**1.3.67.1****elemental assignment**

assignment that operates elementally

**1.3.67.2****elemental operation**

operation that operates elementally

**1.3.67.3****elemental operator**

operator in an elemental operation

**1.3.67.4****elemental procedure**

elemental [intrinsic](#) procedure or procedure defined by an elemental subprogram

**1.3.67.5****elemental reference**

reference to an elemental procedure with at least one array actual argument

**1.3.67.6****elemental subprogram**

subprogram with the [ELEMENTAL](#) prefix

**1.3.68****END statement**

*end-block-data-stmt*, *end-function-stmt*, *end-module-stmt*, *end-mp-subprogram-stmt*, *end-program-stmt*, *end-submodule-stmt*, or *end-subroutine-stmt*

**1.3.69****explicit initialization**

initialization of a data object by a specification statement ([5.4](#), [5.6.7](#))

**1.3.70****explicit interface**

interface of a procedure that includes all the characteristics of the procedure and names for its dummy arguments except for asterisk dummy arguments ([12.4.2](#))

**1.3.71****extent**

number of elements in a single dimension of an [array](#)

**1.3.72****external file**

file that exists in a medium external to the program ([9.3](#))

**1.3.73****external unit****external input/output unit**

entity that can be [connected](#) to an [external file](#)

**1.3.74****file storage unit**

unit of storage in a [stream file](#) or an unformatted [record file](#) ([9.3.5](#))

**1.3.75****final subroutine**

subroutine whose name appears in a [FINAL statement](#) ([4.5.6](#)) in a type definition, and which can be automatically invoked by the processor when an object of that type is finalized ([4.5.6.2](#))

**1.3.76****finalizable**

⟨type⟩ has a final subroutine or a nonpointer nonallocatable component of finalizable type

**1.3.77****finalizable**

⟨nonpointer data entity⟩ of finalizable type

**1.3.78****finalization**

process of calling [final subroutines](#) when one of the events listed in [4.5.6.3](#) occurs

**1.3.79****function**

procedure that is invoked by an expression

**1.3.80****function result**

entity that returns the value of a function

**1.3.81****generic identifier**

[lexical token](#) that identifies a generic set of procedures, [intrinsic](#) operations, and/or [intrinsic](#) assignments

**1.3.82****host instance**

⟨[internal procedure](#), or [dummy procedure](#) or [procedure pointer](#) associated with an [internal procedure](#)⟩ instance of the host procedure that supplies the host environment of the [internal procedure](#) ([12.6.2.4](#))

**1.3.83****host scoping unit****host**

[scoping unit](#) immediately surrounding another [scoping unit](#), or the [scoping unit](#) extended by a submodule

**1.3.84****IEEE infinity**

ISO/IEC/IEEE 60559:2011 conformant infinite floating-point value

**1.3.85****IEEE NaN**

ISO/IEC/IEEE 60559:2011 conformant floating-point datum that does not represent a number

**1.3.86****image**

instance of a Fortran program ([2.3.4](#))

**1.3.86.1****active image**

[image](#) that has not failed or stopped ([2.3.6](#))

**1.3.86.2****failed image**

[image](#) that has not initiated termination but which has ceased to participate in program execution ([2.3.6](#))

**1.3.86.3****stopped image**

[image](#) that has initiated [normal termination](#) ([2.3.6](#))

**1.3.87****image index**

integer value identifying an [image](#) within a [team](#)

**1.3.88****image control statement**

statement that affects the execution ordering between [images](#) (8.6)

**1.3.89****implicit interface**

interface of a procedure that includes only the type and type parameters of a function result (12.4.2, 12.4.3.9)

**1.3.90****inclusive scope**

nonblock [scoping unit](#) plus every [block scoping unit](#) whose [host](#) is that [scoping unit](#) or that is nested within such a [block scoping unit](#)

**NOTE 1.5**

That is, inclusive scope is the scope as if [BLOCK constructs](#) were not [scoping units](#).

**1.3.91****inherit**

⟨[extended type](#)⟩ acquire entities ([components](#), [type-bound procedures](#), and [type parameters](#)) through type extension from the parent type

**1.3.92****inquiry function**

[intrinsic](#) function, or function in an [intrinsic](#) module, whose result depends on the properties of one or more of its arguments instead of their values

**1.3.93****interface**

⟨procedure⟩ name, procedure characteristics, dummy argument names, binding label, and generic identifiers (12.4.1)

**1.3.93.1****generic interface**

set of procedure interfaces identified by a [generic identifier](#)

**1.3.93.2****specific interface**

[interface](#) identified by a nongeneric name

**1.3.94****interface block**

[abstract interface block](#), [generic interface block](#), or [specific interface block](#) (12.4.3.2)

**1.3.94.1****abstract interface block**

interface block with the [ABSTRACT](#) keyword; collection of [interface bodies](#) that specify named [abstract interfaces](#)

**1.3.94.2****generic interface block**

interface block with a [generic-spec](#); collection of [interface bodies](#) and procedure statements that are to be given that generic identifier

**1.3.94.3****specific interface block**

interface block with no *generic-spec* or **ABSTRACT** keyword; collection of *interface bodies* that specify the interfaces of procedures

**1.3.95****interoperable**

⟨Fortran entity⟩ equivalent to an entity defined by or definable by the *companion processor* (15.3)

**1.3.96****intrinsic**

type, procedure, module, assignment, operator, or input/output operation defined in this part of ISO/IEC 1539 and accessible without further definition or specification, or a procedure or module provided by a processor but not defined in this part of ISO/IEC 1539

**1.3.96.1****standard intrinsic**

⟨procedure or module⟩ defined in this part of ISO/IEC 1539 (13)

**1.3.96.2****nonstandard intrinsic**

⟨procedure or module⟩ provided by a processor but not defined in this part of ISO/IEC 1539

**1.3.97****internal file**

character variable that is *connected* to an *internal unit* (9.4)

**1.3.98****internal unit**

*input/output unit* that is *connected* to an *internal file* (9.5.4)

**1.3.99****ISO 10646 character**

character whose representation method corresponds to UCS-4 in ISO/IEC 10646

**1.3.100****keyword**

statement keyword, argument keyword, type parameter keyword, or component keyword

**1.3.100.1****argument keyword**

word that identifies the corresponding *dummy argument* in an *actual argument* list

**1.3.100.2****component keyword**

word that identifies a *component* in a *structure constructor*

**1.3.100.3****statement keyword**

word that is part of the syntax of a statement (2.5.2)

**1.3.100.4****type parameter keyword**

word that identifies a *type parameter* in a type parameter list

**1.3.101****lexical token**

keyword, name, literal constant other than a complex literal constant, operator, label, delimiter, comma, =, =>, :, ::, :::, or % (3.2)

**1.3.102****line**

sequence of zero or more characters

**1.3.103****main program**

program unit that is not a subprogram, module, submodule, or block data program unit (11.1)

**1.3.104****masked array assignment**

assignment statement in a WHERE statement or WHERE construct (7.2.3)

**1.3.105****module**

program unit containing (or accessing from other modules) definitions that are to be made accessible to other program units (11.2)

**1.3.106****name**

identifier of a program constituent, formed according to the rules given in 3.2.2

**1.3.107****NaN**

Not a Number, a symbolic floating-point datum (ISO/IEC/IEEE 60559:2011)

**1.3.108****operand**

data value that is the subject of an operator

**1.3.109****operator**

*intrinsic-operator*, *defined-unary-op*, or *defined-binary-op* (R308, R703, R723)

**1.3.110****passed-object dummy argument**

dummy argument of a type-bound procedure or procedure pointer component that becomes associated with the object through which the procedure is invoked (4.5.4.5)

**1.3.111****pointer**

data pointer (1.3.111.1) or procedure pointer (1.3.111.2)

**1.3.111.1****data pointer**

data entity with the POINTER attribute (5.5.14)

**1.3.111.2****procedure pointer**

procedure with the EXTERNAL and POINTER attributes (5.5.9, 5.5.14)

**1.3.111.3****local procedure pointer**

[procedure pointer](#) that is part of a [local variable](#), or a named [procedure pointer](#) that is not a [dummy argument](#) or accessed by use or host association

**1.3.112****pointer assignment**

association of a pointer with a target, by execution of a [pointer assignment statement](#) (7.2.2) or an [intrinsic assignment statement](#) (7.2.1.2) for a derived-type object that has the pointer as a subobject

**1.3.113****polymorphic**

⟨data entity⟩ able to be of differing [dynamic types](#) during program execution (4.3.2.3)

**1.3.114****preconnected**

⟨file or [unit](#)⟩ [connected](#) at the beginning of execution of the program (9.5.5)

**1.3.115****procedure**

entity encapsulating an arbitrary sequence of actions that can be invoked directly during program execution

**1.3.115.1****dummy procedure**

procedure that is a [dummy argument](#) (12.2.2.3)

**1.3.115.2****external procedure**

procedure defined by an external subprogram (R203) or by means other than Fortran (12.6.3)

**1.3.115.3****internal procedure**

procedure defined by an internal subprogram (R212)

**1.3.115.4****module procedure**

procedure that is defined by a module subprogram (R1108)

**1.3.115.5****pure procedure**

procedure declared or defined to be pure according to the rules in 12.7

**1.3.115.6****type-bound procedure**

procedure that is bound to a derived type and referenced via an object of that type (4.5.5)

**1.3.116****processor**

combination of a computing system and mechanism by which programs are transformed for use on that computing system

**1.3.117****processor dependent**

not completely specified in this part of ISO/IEC 1539, having methods and semantics determined by the processor

**1.3.118****program**

set of Fortran [program units](#) and entities defined by means other than Fortran that includes exactly one [main program](#)

**1.3.119****program unit**

[main program](#), [external subprogram](#), [module](#), [submodule](#), or [block data program unit](#) (2.2.1)

**1.3.120****rank**

number of array dimensions of a [data entity](#) (zero for a scalar entity)

**1.3.121****record**

sequence of values or characters in a file (9.2)

**1.3.122****record file**

file composed of a sequence of records (9.1)

**1.3.123****reference**

[data object](#) reference, [procedure](#) reference, or [module](#) reference

**1.3.123.1****data object reference**

appearance of a [data object designator](#) (6.1) in a context requiring its value at that point during execution

**1.3.123.2****function reference**

appearance of the [procedure designator](#) for a function, or operator symbol in a context requiring execution of the function during expression evaluation (12.5.3)

**1.3.123.3****module reference**

appearance of a module name in a [USE statement](#) (11.2.2)

**1.3.123.4****procedure reference**

appearance of a [procedure designator](#), operator symbol, or assignment symbol in a context requiring execution of the procedure at that point during execution; or occurrence of defined input/output (10.7.6) or derived-type [finalization](#) (4.5.6.2)

**1.3.124****saved**

having the [SAVE attribute](#) (5.5.16)

**1.3.125****scalar**

[data entity](#) that can be represented by a single value of the type and that is not an array (6.5)

**1.3.126****scoping unit**

[BLOCK construct](#), derived-type definition, [interface body](#), [program unit](#), or subprogram, excluding all nested scoping units in it

**1.3.126.1****block scoping unit**

scoping unit of a [BLOCK construct](#)

**1.3.127****sequence**

set of elements ordered by a one-to-one correspondence with the numbers 1, 2, to  $n$

**1.3.128****sequence structure**

[scalar data object](#) of a [sequence type](#) (4.5.2.3)

**1.3.129****sequence type**

derived type with the [SEQUENCE attribute](#) (4.5.2.3)

**1.3.129.1****character sequence type**

sequence type with no [type parameters](#), no [allocatable](#) or [pointer components](#), and whose [components](#) are all default character or of another character sequence type

**1.3.129.2****numeric sequence type**

sequence type with no [type parameters](#), no [allocatable](#) or [pointer components](#), and whose [components](#) are all default complex, default integer, default logical, default real, double precision real, or of another numeric sequence type

**1.3.130****shape**

array dimensionality of a data entity, represented as a rank-one array whose size is the [rank](#) of the data entity and whose elements are the extents of the data entity

**NOTE 1.6**

Thus the shape of a scalar data entity is an array with rank one and size zero.

**1.3.131****simply contiguous**

[array designator](#) or [variable](#) satisfying the conditions specified in 6.5.4

**NOTE 1.7**

These conditions are simple ones which make it clear that the designator or variable designates a [contiguous](#) array.

**1.3.132****size**

[array](#) total number of elements in the array

**1.3.133****specification expression**

expression satisfying the requirements specified in 7.1.11, thus being suitable for use in specifications

**1.3.134****specific name**

name that is not a generic name

**1.3.135****standard-conforming program**

program that uses only those forms and relationships described in, and has an interpretation according to, this part of ISO/IEC 1539

**1.3.136****statement**

sequence of one or more complete or partial lines satisfying a syntax rule that ends in *-stmt* (3.3)

**1.3.136.1****executable statement**

statement that is a member of the syntactic class *executable-construct*, excluding those in the *block-specification-part* of a **BLOCK construct**

**1.3.136.2****nonexecutable statement**

statement that is not an *executable statement*

**1.3.137****statement entity**

entity whose identifier has the scope of a statement or part of a statement (16.1, 16.4)

**1.3.138****statement label****label**

unsigned positive number of up to five digits that refers to an individual statement (3.2.5)

**1.3.139****storage sequence**

contiguous sequence of *storage units* (16.5.3.2)

**1.3.140****storage unit**

*character storage unit*, *numeric storage unit*, *file storage unit*, or *unspecified storage unit* (16.5.3.2)

**1.3.140.1****character storage unit**

unit of storage that holds a default character value (16.5.3.2)

**1.3.140.2****numeric storage unit**

unit of storage that holds a default real, default integer, or default logical value (16.5.3.2)

**1.3.140.3****unspecified storage unit**

unit of storage that holds a value that is not default character, default real, double precision real, default logical, or default complex (16.5.3.2)

**1.3.141****stream file**

file composed of a sequence of file storage units (9.1)

**1.3.142****structure**

*scalar data object* of *derived type* (4.5)

**1.3.142.1****structure component**

[component](#) of a structure

**1.3.142.2****structure constructor**

syntax (*structure-constructor*, [4.5.10](#)) that specifies a structure value or creates such a value

**1.3.143****submodule**

[program unit](#) that extends a [module](#) or another [submodule](#) ([11.2.3](#))

**1.3.144****subobject**

portion of [data object](#) that can be referenced, and if it is a [variable](#) defined, independently of any other portion

**1.3.145****subprogram**

*function-subprogram* ([R1229](#)) or *subroutine-subprogram* ([R1235](#))

**1.3.145.1****external subprogram**

subprogram that is not contained in a [main program](#), [module](#), [submodule](#), or another subprogram

**1.3.145.2****internal subprogram**

subprogram that is contained in a [main program](#) or another subprogram

**1.3.145.3****module subprogram**

subprogram that is contained in a [module](#) or [submodule](#) but is not an internal subprogram

**1.3.146****subroutine**

procedure invoked by a [CALL statement](#), by [defined assignment](#), or by some operations on derived-type entities

**1.3.146.1****atomic subroutine**

[intrinsic](#) subroutine that performs an action on its ATOM argument atomically

**1.3.146.2****collective subroutine**

[intrinsic](#) subroutine that performs a calculation on a [team](#) of images without requiring synchronization

**1.3.147****target**

entity that is [pointer associated](#) with a [pointer](#) ([16.5.2.2](#)), entity on the right-hand-side of a [pointer assignment statement](#) ([R733](#)), or entity with the [TARGET attribute](#) ([5.5.17](#))

**1.3.148****team**

ordered set of images created by execution of [FORM TEAM statements](#), or the initial ordered set of all images

**1.3.148.1****current team**

team specified by the most recently executed [CHANGE TEAM statement](#) of a [CHANGE TEAM construct](#) that has not completed execution ([8.1.5](#)), or initial team if no [CHANGE TEAM construct](#) is being executed

**1.3.148.2****initial team**

team existing at the beginning of program execution, consisting of all images

**1.3.148.3****parent team**

⟨team except for initial team⟩ current team at time of execution of the [FORM TEAM statement](#) that created the team ([8.6.9](#))

**1.3.148.4****team number**

⟨team except for initial team⟩ integer value that identifies a team within its parent team

**1.3.149****transformational function**

[intrinsic](#) function, or function in an [intrinsic](#) module, that is neither [elemental](#) nor an [inquiry function](#)

**1.3.150****type****data type**

named category of data characterized by a set of values, a syntax for denoting these values, and a set of operations that interpret and manipulate the values ([4.1](#))

**1.3.150.1****abstract type**

type with the [ABSTRACT attribute](#) ([4.5.7.1](#))

**1.3.150.2****declared type**

type that a data entity is declared to have, either explicitly or implicitly ([4.3.2](#), [7.1.9](#))

**1.3.150.3****derived type**

type defined by a type definition ([4.5](#)) or by an [intrinsic](#) module

**1.3.150.4****dynamic type**

type of a data entity at a particular point during execution of a program ([4.3.2.3](#), [7.1.9](#))

**1.3.150.5****extended type**

type with the [EXTENDS attribute](#) ([4.5.7.1](#))

**1.3.150.6****extensible type**

type that may be extended using the EXTENDS clause ([4.5.7.1](#))

**1.3.150.7****extension type**

⟨of one type with respect to another⟩ is the same type or is an extended type whose parent type is an extension type of the other type

**1.3.150.8****intrinsic type**

type defined by this part of ISO/IEC 1539 that is always accessible ([4.4](#))

**1.3.150.9****numeric type**

one of the types integer, real, and complex

**1.3.150.10****parent type**

⟨[extended type](#)⟩ type named in the EXTENDS clause

**1.3.150.11****type compatible**

compatibility of the type of one entity with respect to another for purposes such as [argument association](#), [pointer association](#), and allocation ([4.3.2](#))

**1.3.150.12****type parameter**

value used to parameterize a type ([4.2](#))

**1.3.150.12.1****assumed type parameter**

[length type parameter](#) that assumes the type parameter value from another entity

**NOTE 1.8**

The other entity is

- the selector for an [associate name](#),
- the [constant-expr](#) for a [named constant](#) of type character, or
- the [effective argument](#) for a [dummy argument](#).

**1.3.150.12.2****deferred type parameter**

[length type parameter](#) whose value can change during execution of a program and whose [type-param-value](#) is a colon

**1.3.150.12.3****kind type parameter**

type parameter whose value is required to be defaulted or given by a [constant expression](#)

**1.3.150.12.4****length type parameter**

type parameter whose value is permitted to be [assumed](#), [deferred](#), or given by a [specification expression](#)

**1.3.150.12.5****type parameter inquiry**

syntax ([type-param-inquiry](#)) that is used to inquire the value of a type parameter of a data object ([6.4.5](#))

**1.3.150.12.6****type parameter order**

ordering of the type parameters of a type ([4.5.3.2](#)) used for derived-type specifiers ([derived-type-spec](#), [4.5.9](#))

**1.3.151****ultimate argument**

nondummy entity with which a [dummy argument](#) is associated via a chain of argument associations ([12.5.2.3](#))

**1.3.152****undefined**

⟨[data object](#)⟩ does not have a valid value

**1.3.153****undefined**

⟨*pointer*⟩ does not have a pointer association status of associated or *disassociated* (16.5.2.2)

**1.3.154****unit****input/output unit**

means, specified by an *io-unit*, for referring to a file (9.5.1)

**1.3.155****unlimited polymorphic**

able to have any *dynamic type* during program execution (4.3.2.3)

**1.3.156****unsaved**

not having the *SAVE attribute* (5.5.16)

**1.3.157****variable**

*data entity* that can be *defined* and redefined during execution of a program

**1.3.157.1****event variable**

scalar variable of type *EVENT\_TYPE* (13.10.2.10) from the intrinsic module *ISO\_FORTRAN\_ENV*

**1.3.157.2****local variable**

variable in a *scoping unit* that is not a *dummy argument* or part thereof, is not a global entity or part thereof, and is not an entity or part of an entity that is accessible outside that *scoping unit*

**1.3.157.3****lock variable**

scalar variable of type *LOCK\_TYPE* (13.10.2.19) from the intrinsic module *ISO\_FORTRAN\_ENV*

**1.3.157.4****team variable**

scalar variable of type *TEAM\_TYPE* (13.10.2.31) from the intrinsic module *ISO\_FORTRAN\_ENV*

**1.3.158****vector subscript**

*section-subscript* that is an array (6.5.3.3.2)

**1.3.159****whole array**

array component or array name without further qualification (6.5.2)

**1.4 Notation, symbols and abbreviated terms****1.4.1 Syntax rules**

- 1 Syntax rules describe the forms that Fortran *lexical tokens*, statements, and constructs may take. These syntax rules are expressed in a variation of Backus-Naur form (BNF) with the following conventions.

- Characters from the Fortran character set (3.1) are interpreted literally as shown, except where otherwise noted.
  - Lower-case italicized letters and words (often hyphenated and abbreviated) represent general syntactic classes for which particular syntactic entities shall be substituted in actual statements.
- Common abbreviations used in syntactic terms are:

<i>arg</i>	for	argument	<i>attr</i>	for	attribute
<i>decl</i>	for	declaration	<i>def</i>	for	definition
<i>desc</i>	for	descriptor	<i>expr</i>	for	expression
<i>int</i>	for	integer	<i>op</i>	for	operator
<i>spec</i>	for	specifier	<i>stmt</i>	for	statement

- The syntactic metasympols used are:

<b>is</b>	introduces a syntactic class definition
<b>or</b>	introduces a syntactic class alternative
[ ]	encloses an optional item
[ ] ...	encloses an optionally repeated item that may occur zero or more times
■	continues a syntax rule

- Each syntax rule is given a unique identifying number of the form R<sub>snn</sub>, where s is a one- or two-digit clause number and nn is a two-digit sequence number within that clause. The syntax rules are distributed as appropriate throughout the text, and are referenced by number as needed. Some rules in Clauses 2 and 3 are more fully described in later clauses; in such cases, the clause number s is the number of the later clause where the rule is repeated.
- The syntax rules are not a complete and accurate syntax description of Fortran, and cannot be used to generate a Fortran parser automatically; where a syntax rule is incomplete, it is restricted by corresponding constraints and text.

#### NOTE 1.9

An example of the use of the syntax rules is:

*digit-string*                      **is** *digit* [ *digit* ] ...

The following are examples of forms for a digit string allowed by the above rule:

*digit*  
*digit digit*  
*digit digit digit digit*  
*digit digit digit digit digit digit digit digit*

If particular entities are substituted for *digit*, actual digit strings might be:

4  
67  
1999  
10243852

## 1.4.2 Constraints

- Each constraint is given a unique identifying number of the form C<sub>snn</sub>, where s is a one or two digit clause number and nn is a two or three digit sequence number within that clause.
- Often a constraint is associated with a particular syntax rule. Where that is the case, the constraint is annotated with the syntax rule number in parentheses. A constraint that is associated with a syntax rule constitutes part of the definition of the syntax term defined by the rule. It thus applies in all places where the syntax term appears.
- Some constraints are not associated with particular syntax rules. The effect of such a constraint is similar to that of a restriction stated in the text, except that a processor is required to have the capability to detect and report violations of constraints (1.5). In some cases, a broad requirement is stated in text and a subset of the same requirement is also stated as a constraint. This indicates that a standard-conforming program is required to

adhere to the broad requirement, but that a standard-conforming processor is required only to have the capability of diagnosing violations of the constraint.

### 1.4.3 Assumed syntax rules

- 1 In order to minimize the number of additional syntax rules and convey appropriate constraint information, the following rules are assumed.

R101 *xyz-list* is *xyz* [ , *xyz* ] ...

R102 *xyz-name* is *name*

R103 *scalar-xyz* is *xyz*

C101 (R103) *scalar-xyz* shall be scalar.

- 2 The letters “*xyz*” stand for any syntactic class phrase. An explicit syntax rule for a term overrides an assumed rule.

### 1.4.4 Syntax conventions and characteristics

- 1 Any syntactic class name ending in “-*stmt*” follows the source form statement rules: it shall be delimited by end-of-line or semicolon, and may be labeled unless it forms part of another statement (such as an **IF** or **WHERE** statement). Conversely, everything considered to be a source form statement is given a “-*stmt*” ending in the syntax rules.
- 2 The rules on statement ordering are described rigorously in the definition of *program-unit* (R202). Expression hierarchy is described rigorously in the definition of *expr* (R722).
- 3 The suffix “-*spec*” is used consistently for specifiers, such as input/output statement specifiers. It also is used for type declaration attribute specifications (for example, “*array-spec*” in R515), and in a few other cases.
- 4 Where reference is made to a type parameter, including the surrounding parentheses, the suffix “-*selector*” is used. See, for example, “*kind-selector*” (R406) and “*length-selector*” (R422).

### 1.4.5 Text conventions

- 1 In descriptive text, an equivalent English word is frequently used in place of a syntactic term. Particular statements and attributes are identified in the text by an upper-case keyword, e.g., “END statement”. The descriptions of obsolescent features appear in a smaller type size.

#### NOTE 1.10

This sentence is an example of the type size used for obsolescent features.

## 1.5 Conformance

- 1 A **program** (2.2.2) is a **standard-conforming program** if it uses only those forms and relationships described herein and if the program has an interpretation according to this part of ISO/IEC 1539. A **program unit** (2.2.1) conforms to this part of ISO/IEC 1539 if it can be included in a **program** in a manner that allows the **program** to be standard conforming.
- 2 A **processor** conforms to this part of ISO/IEC 1539 if:
  - (1) it executes any **standard-conforming program** in a manner that fulfills the interpretations herein, subject to any limits that the **processor** may impose on the size and complexity of the **program**;
  - (2) it contains the capability to detect and report the use within a submitted **program unit** of a form designated herein as obsolescent, insofar as such use can be detected by reference to the numbered syntax rules and constraints;

- (3) it contains the capability to detect and report the use within a submitted **program unit** of an additional form or relationship that is not permitted by the numbered syntax rules or constraints, including the deleted features described in Annex B
- (4) it contains the capability to detect and report the use within a submitted **program unit** of an intrinsic type with a **kind type parameter** value not supported by the processor (4.4);
- (5) it contains the capability to detect and report the use within a submitted **program unit** of source form or characters not permitted by Clause 3;
- (6) it contains the capability to detect and report the use within a submitted **program** of name usage not consistent with the scope rules for names, labels, operators, and assignment symbols in Clause 16;
- (7) it contains the capability to detect and report the use within a submitted **program unit** of a **non-standard intrinsic** procedure (including one with the same name as a **standard intrinsic** procedure but with different requirements);
- (8) it contains the capability to detect and report the use within a submitted **program unit** of a **non-standard intrinsic** module;
- (9) it contains the capability to detect and report the use within a submitted **program unit** of a procedure from a **standard intrinsic** module, if the procedure is not defined by this part of ISO/IEC 1539 or the procedure has different requirements from those specified by this part of ISO/IEC 1539; and
- (10) it contains the capability to detect and report the reason for rejecting a submitted **program**.

3 However, in a format specification that is not part of a **FORMAT statement** (10.2.1), a processor need not detect or report the use of deleted or obsolescent features, or the use of additional forms or relationships.

4 A standard-conforming processor may allow additional forms and relationships provided that such additions do not conflict with the standard forms and relationships. However, a standard-conforming processor may allow additional intrinsic procedures even though this could cause a conflict with the name of a procedure in a standard-conforming program. If such a conflict occurs and involves the name of an **external procedure**, the processor is permitted to use the intrinsic procedure unless the name has the **EXTERNAL attribute** (5.5.9) where it is used. A standard-conforming program shall not use **nonstandard intrinsic** procedures or modules that have been added by the processor.

5 Because a standard-conforming program may place demands on a processor that are not within the scope of this part of ISO/IEC 1539 or may include standard items that are not portable, such as **external procedures** defined by means other than Fortran, conformance to this part of ISO/IEC 1539 does not ensure that a program will execute consistently on all or any standard-conforming processors.

6 The semantics of facilities that are identified as **processor dependent** are not completely specified in this part of ISO/IEC 1539. They shall be provided, with methods or semantics determined by the processor.

7 The **processor** should be accompanied by documentation that specifies the limits it imposes on the size and complexity of a **program** and the means of reporting when these limits are exceeded, that defines the additional forms and relationships it allows, and that defines the means of reporting the use of additional forms and relationships and the use of deleted or obsolescent forms. In this context, the use of a deleted form is the use of an additional form.

8 The **processor** should be accompanied by documentation that specifies the methods or semantics of processor-dependent facilities.

## 1.6 Compatibility

### 1.6.1 Previous Fortran standards

1 Table 1.3 lists the previous editions of the Fortran International Standard, along with their informal names.

Table 1.3: Previous editions of the Fortran International Standard

Official designation	Informal name
ISO R 1539-1972	FORTTRAN 66
ISO 1539-1980	FORTTRAN 77
ISO/IEC 1539:1991	Fortran 90
ISO/IEC 1539-1:1997	Fortran 95
ISO/IEC 1539-1:2004	Fortran 2003
ISO/IEC 1539-1:2010	Fortran 2008

## 1.6.2 New intrinsic procedures

- Each Fortran International Standard since ISO 1539:1980 (FORTTRAN 77), defines more intrinsic procedures than the previous one. Therefore, a Fortran program conforming to an older standard might have a different interpretation under a newer standard if it invokes an external procedure having the same name as one of the new standard intrinsic procedures, unless that procedure is specified to have the [EXTERNAL attribute](#).

## 1.6.3 Fortran 2008 compatibility

- Except as identified in this subclause, this part of ISO/IEC 1539 is an upward compatible to the preceding Fortran International Standard, ISO/IEC 1539-1:2010 (Fortran 2008). Except for the deleted features noted in Annex [B.2](#), any standard-conforming Fortran 2008 program remains standard-conforming under this part of ISO/IEC 1539.
- Fortran 2008 specifies that the [IOSTAT=](#) variable shall be set to a processor-dependent negative value if the flush operation is not supported for the unit specified. This part of ISO/IEC 1539 specifies that the processor-dependent negative integer value must be different from the [named constants](#) [IOSTAT\\_EOR](#) or [IOSTAT\\_END](#) from the intrinsic module [ISO\\_FORTRAN\\_ENV](#).

## 1.6.4 Fortran 2003 compatibility

- Except as identified in this subclause, this part of ISO/IEC 1539 is an upward compatible extension to ISO/IEC 1539-1:2004 (Fortran 2003). Except as identified in this subclause, any standard-conforming Fortran 2003 program remains standard-conforming under this part of ISO/IEC 1539.
- Fortran 2003 permitted a [sequence type](#) to have [type parameters](#); that is not permitted by this part of ISO/IEC 1539.
- Fortran 2003 specified that array constructors and structure constructors of [finalizable](#) type are finalized. This part of ISO/IEC 1539 specifies that these constructors are not finalized.
- The form produced by the G edit descriptor for some values and some input/output rounding modes differs from that specified by Fortran 2003.
- Fortran 2003 required an [explicit interface](#) only for a procedure that was actually referenced in the scope, not merely passed as an [actual argument](#). This part of ISO/IEC 1539 requires an [explicit interface](#) for a procedure under the conditions listed in [12.4.2.2](#), regardless of whether the procedure is [referenced](#) in the scope.
- Fortran 2003 permitted the [function result](#) of a [pure](#) function to be a [polymorphic](#) allocatable variable, and to be [finalizable](#) by an impure [final subroutine](#). These are not permitted by this part of ISO/IEC 1539.
- Fortran 2003 permitted an [INTENT \(OUT\)](#) argument of a [pure](#) subroutine to be [polymorphic](#); that is not permitted by this part of ISO/IEC 1539.

### 1.6.5 Fortran 95 compatibility

- 1 Except as identified in this subclause, this part of ISO/IEC 1539 is an upward compatible extension to ISO/IEC 1539-1:1997 (Fortran 95). Except as identified in this subclause, any standard-conforming Fortran 95 program remains standard-conforming under this part of ISO/IEC 1539.
- 2 Fortran 95 permitted defined assignment between character strings of the same rank and different kinds. This part of ISO/IEC 1539 does not permit that if both of the different kinds are ASCII, ISO 10646, or default kind.
- 3 The following Fortran 95 features might have different interpretations in this part of ISO/IEC 1539.
  - Earlier Fortran standards had the concept of printing, meaning that column one of formatted output had special meaning for a processor-dependent (possibly empty) set of [external files](#). This could be neither detected nor specified by a standard-specified means. The interpretation of the first column is not specified by this part of ISO/IEC 1539.
  - This part of ISO/IEC 1539 specifies a different output format for real zero values in list-directed and namelist output.
  - If the processor distinguishes between positive and negative real zero, this part of ISO/IEC 1539 requires different returned values for [ATAN2\(Y,X\)](#) when  $X < 0$  and  $Y$  is negative real zero and for [LOG\(X\)](#) and [SQRT\(X\)](#) when  $X$  is complex with [REAL\(X\) < 0](#) and negative zero imaginary part.
  - This part of ISO/IEC 1539 has fewer restrictions on [constant expressions](#) than Fortran 95; this might affect whether a variable is considered to be an [automatic data object](#).
  - The form produced by the G edit descriptor with  $d$  equal to zero differs from that specified by Fortran 95 for some values.

### 1.6.6 Fortran 90 compatibility

- 1 Except for the deleted features noted in Annex [B.1](#), and except as identified in this subclause, this part of ISO/IEC 1539 is an upward compatible extension to ISO/IEC 1539:1991 (Fortran 90). Any standard-conforming Fortran 90 program that does not use one of the deleted features remains standard-conforming under this part of ISO/IEC 1539.
- 2 The [PAD= specifier](#) in the [INQUIRE statement](#) in this part of ISO/IEC 1539 returns the value UNDEFINED if there is no connection or the connection is for unformatted input/output. Fortran 90 specified YES.
- 3 Fortran 90 specified that if the second argument to [MOD](#) or [MODULO](#) was zero, the result was processor dependent. This part of ISO/IEC 1539 specifies that the second argument shall not be zero.
- 4 Fortran 90 permitted defined assignment between character strings of the same rank and different kinds. This part of ISO/IEC 1539 does not permit that if both of the different kinds are ASCII, ISO 10646, or default kind.
- 5 The following Fortran 90 features have different interpretations in this part of ISO/IEC 1539:
  - if the processor distinguishes between positive and negative real zero, the result value of the intrinsic function [SIGN](#) when the second argument is a negative real zero;
  - formatted output of negative real values (when the output value is zero);
  - whether an expression is a [constant expression](#) (thus whether a variable is considered to be an [automatic data object](#));
  - the G edit descriptor with  $d$  equal to zero for some values.

### 1.6.7 FORTRAN 77 compatibility

- 1 Except for the deleted features noted in Annex [B.1](#), and except as identified in this subclause, this part of ISO/IEC 1539 is an upward compatible extension to ISO 1539:1980 (FORTRAN 77). Any standard-conforming FORTRAN 77 program that does not use one of the deleted features noted in Annex [B.1](#) and that does not depend on the differences specified here remains standard-conforming under this part of ISO/IEC 1539. This part of ISO/IEC 1539 restricts the behavior for some features that were processor dependent in FORTRAN 77. Therefore,

a standard-conforming FORTRAN 77 program that uses one of these processor-dependent features might have a different interpretation under this part of ISO/IEC 1539, yet remain a standard-conforming program. The following FORTRAN 77 features might have different interpretations in this part of ISO/IEC 1539.

- FORTRAN 77 permitted a processor to supply more precision derived from a default real constant than can be represented in a default real datum when the constant is used to initialize a double precision real data object in a [DATA statement](#). This part of ISO/IEC 1539 does not permit a processor this option.
- If a named variable that was not in a [common block](#) was initialized in a [DATA statement](#) and did not have the [SAVE attribute](#) specified, FORTRAN 77 left its [SAVE attribute](#) processor dependent. This part of ISO/IEC 1539 specifies (5.6.7) that this named variable has the [SAVE attribute](#).
- FORTRAN 77 specified that the number of characters required by the input list was to be less than or equal to the number of characters in the record during formatted input. This part of ISO/IEC 1539 specifies (9.6.4.5.3) that the input record is logically padded with blanks if there are not enough characters in the record, unless the [PAD= specifier](#) with the value 'NO' is specified in an appropriate [OPEN](#) or [READ](#) statement.
- A value of 0 for a list item in a formatted [output statement](#) will be formatted in a different form for some G edit descriptors. In addition, this part of ISO/IEC 1539 specifies how rounding of values will affect the output field form, but FORTRAN 77 did not address this issue. Therefore, some FORTRAN 77 processors might produce an output form different from the output form produced by Fortran 2003 processors for certain combinations of values and G edit descriptors.
- FORTRAN 77 did not permit a processor to distinguish between positive and negative real zero; if the processor does so distinguish, the result will differ for the intrinsic function [SIGN](#) when the second argument is negative real zero, and formatted output of negative real zero will be different.

## 1.7 Deleted and obsolescent features

### 1.7.1 General

- 1 This part of ISO/IEC 1539 protects the users' investment in existing software by including all but six of the language elements of Fortran 90 that are not processor dependent. This part of ISO/IEC 1539 identifies two categories of outmoded features. The first category, deleted features, consists of features considered to have been redundant in FORTRAN 77 and largely unused in Fortran 90. Those in the second category, obsolescent features, are considered to have been redundant in Fortran 90 and Fortran 95, but are still frequently used.

### 1.7.2 Nature of deleted features

- 1 Better methods existed in FORTRAN 77 for each deleted feature. These features were not included in Fortran 95 or Fortran 2003, and are not included in this revision of Fortran.

### 1.7.3 Nature of obsolescent features

- 1 Better methods existed in Fortran 90 and Fortran 95 for each obsolescent feature. It is recommended that programmers use these better methods in new programs and convert existing code to these methods.
- 2 The obsolescent features are identified in the text of this part of ISO/IEC 1539 by a distinguishing type font (1.4.5).
- 3 A future revision of this part of ISO/IEC 1539 might delete an obsolescent feature if its use has become insignificant.

## 2 Fortran concepts

### 2.1 High level syntax

- 1 This subclause introduces the syntax associated with [program units](#) and other Fortran concepts above the construct, statement, and expression levels and illustrates their relationships.

#### NOTE 2.1

Constraints and other information related to the rules that do not begin with R2 appear in the appropriate clause.

R201	<i>program</i>	is	<i>program-unit</i> [ <i>program-unit</i> ] ...
R202	<i>program-unit</i>	is	<i>main-program</i> or <i>external-subprogram</i> or <i>module</i> or <i>submodule</i> or <i>block-data</i>
R1101	<i>main-program</i>	is	[ <i>program-stmt</i> ] [ <i>specification-part</i> ] [ <i>execution-part</i> ] [ <i>internal-subprogram-part</i> ] <i>end-program-stmt</i>
R203	<i>external-subprogram</i>	is	<i>function-subprogram</i> or <i>subroutine-subprogram</i>
R1229	<i>function-subprogram</i>	is	<i>function-stmt</i> [ <i>specification-part</i> ] [ <i>execution-part</i> ] [ <i>internal-subprogram-part</i> ] <i>end-function-stmt</i>
R1235	<i>subroutine-subprogram</i>	is	<i>subroutine-stmt</i> [ <i>specification-part</i> ] [ <i>execution-part</i> ] [ <i>internal-subprogram-part</i> ] <i>end-subroutine-stmt</i>
R1104	<i>module</i>	is	<i>module-stmt</i> [ <i>specification-part</i> ] [ <i>module-subprogram-part</i> ] <i>end-module-stmt</i>
R1116	<i>submodule</i>	is	<i>submodule-stmt</i> [ <i>specification-part</i> ] [ <i>module-subprogram-part</i> ] <i>end-submodule-stmt</i>
R1120	<i>block-data</i>	is	<i>block-data-stmt</i> [ <i>specification-part</i> ]

1			<i>end-block-data-stmt</i>
2	R204	<i>specification-part</i>	is [ <i>use-stmt</i> ] ...
3			[ <i>import-stmt</i> ] ...
4			[ <i>implicit-part</i> ]
5			[ <i>declaration-construct</i> ] ...
6	R205	<i>implicit-part</i>	is [ <i>implicit-part-stmt</i> ] ...
7			<i>implicit-stmt</i>
8	R206	<i>implicit-part-stmt</i>	is <i>implicit-stmt</i>
9			or <i>parameter-stmt</i>
10			or <i>format-stmt</i>
11			or <i>entry-stmt</i>
12	R207	<i>declaration-construct</i>	is <i>specification-construct</i>
13			or <i>data-stmt</i>
14			or <i>format-stmt</i>
15			or <i>entry-stmt</i>
16			or <i>stmt-function-stmt</i>
17	R208	<i>specification-construct</i>	is <i>derived-type-def</i>
18			or <i>enum-def</i>
19			or <i>generic-stmt</i>
20			or <i>interface-block</i>
21			or <i>parameter-stmt</i>
22			or <i>procedure-declaration-stmt</i>
23			or <i>other-specification-stmt</i>
24			or <i>type-declaration-stmt</i>
25	R209	<i>execution-part</i>	is <i>executable-construct</i>
26			[ <i>execution-part-construct</i> ] ...
27	R210	<i>execution-part-construct</i>	is <i>executable-construct</i>
28			or <i>format-stmt</i>
29			or <i>entry-stmt</i>
30			or <i>data-stmt</i>
31	R211	<i>internal-subprogram-part</i>	is <i>contains-stmt</i>
32			[ <i>internal-subprogram</i> ] ...
33	R212	<i>internal-subprogram</i>	is <i>function-subprogram</i>
34			or <i>subroutine-subprogram</i>
35	R1107	<i>module-subprogram-part</i>	is <i>contains-stmt</i>
36			[ <i>module-subprogram</i> ] ...
37	R1108	<i>module-subprogram</i>	is <i>function-subprogram</i>
38			or <i>subroutine-subprogram</i>
39			or <i>separate-module-subprogram</i>
40	R1239	<i>separate-module-subprogram</i>	is <i>mp-subprogram-stmt</i>
41			[ <i>specification-part</i> ]
42			[ <i>execution-part</i> ]
43			[ <i>internal-subprogram-part</i> ]
44			<i>end-mp-subprogram-stmt</i>
45	R213	<i>other-specification-stmt</i>	is <i>access-stmt</i>
46			or <i>allocatable-stmt</i>

1		or	<i>asynchronous-stmt</i>
2		or	<i>bind-stmt</i>
3		or	<i>codimension-stmt</i>
4		or	<i>contiguous-stmt</i>
5		or	<i>dimension-stmt</i>
6		or	<i>external-stmt</i>
7		or	<i>intent-stmt</i>
8		or	<i>intrinsic-stmt</i>
9		or	<i>namelist-stmt</i>
10		or	<i>optional-stmt</i>
11		or	<i>pointer-stmt</i>
12		or	<i>protected-stmt</i>
13		or	<i>save-stmt</i>
14		or	<i>target-stmt</i>
15		or	<i>volatile-stmt</i>
16		or	<i>value-stmt</i>
17		or	<i>common-stmt</i>
18		or	<i>equivalence-stmt</i>
19	R214	<i>executable-construct</i>	is <i>action-stmt</i>
20			or <i>associate-construct</i>
21			or <i>block-construct</i>
22			or <i>case-construct</i>
23			or <i>change-team-construct</i>
24			or <i>critical-construct</i>
25			or <i>do-construct</i>
26			or <i>if-construct</i>
27			or <i>select-rank-construct</i>
28			or <i>select-type-construct</i>
29			or <i>where-construct</i>
30			or <i>forall-construct</i>
31	R215	<i>action-stmt</i>	is <i>allocate-stmt</i>
32			or <i>assignment-stmt</i>
33			or <i>backspace-stmt</i>
34			or <i>call-stmt</i>
35			or <i>close-stmt</i>
36			or <i>continue-stmt</i>
37			or <i>cycle-stmt</i>
38			or <i>deallocate-stmt</i>
39			or <i>end-function-stmt</i>
40			or <i>end-mp-subprogram-stmt</i>
41			or <i>end-program-stmt</i>
42			or <i>end-subroutine-stmt</i>
43			or <i>endfile-stmt</i>
44			or <i>error-stop-stmt</i>
45			or <i>event-post-stmt</i>
46			or <i>event-wait-stmt</i>
47			or <i>exit-stmt</i>
48			or <i>fail-image-stmt</i>
49			or <i>flush-stmt</i>
50			or <i>form-team-stmt</i>
51			or <i>goto-stmt</i>
52			or <i>if-stmt</i>
53			or <i>inquire-stmt</i>
54			or <i>lock-stmt</i>

or *nullify-stmt*  
 or *open-stmt*  
 or *pointer-assignment-stmt*  
 or *print-stmt*  
 or *read-stmt*  
 or *return-stmt*  
 or *rewind-stmt*  
 or *stop-stmt*  
 or *sync-all-stmt*  
 or *sync-images-stmt*  
 or *sync-memory-stmt*  
 or *sync-team-stmt*  
 or *unlock-stmt*  
 or *wait-stmt*  
 or *where-stmt*  
 or *write-stmt*  
 or *computed-goto-stmt*  
 or *forall-stmt*

C201 (R209) An *execution-part* shall not contain an *end-function-stmt*, *end-mp-subprogram-stmt*, *end-program-stmt*, or *end-subroutine-stmt*.

## 2.2 Program unit concepts

### 2.2.1 Program units and scoping units

- 1 **Program units** are the fundamental components of a Fortran program. A **program unit** is a **main program**, an **external subprogram**, a **module**, a **submodule**, or a block data program unit.
- 2 A **subprogram** is a function subprogram or a subroutine subprogram. A **module** contains definitions that are to be made accessible to other **program units**. A **submodule** is an extension of a **module**; it may contain the definitions of procedures declared in a **module** or another **submodule**. A block data program unit is used to specify initial values for **data objects** in named **common blocks**.
- 3 Each type of **program unit** is described in Clause 11 or 12.
- 4 A **program unit** consists of a set of nonoverlapping **scoping units**.

#### NOTE 2.2

The module or submodule containing a **module subprogram** is the **host scoping unit** of the **module subprogram**. The containing **main program** or **subprogram** is the **host scoping unit** of an **internal subprogram**.

An **internal procedure** is local to its **host** in the sense that its name is accessible within the **host scoping unit** and all its other **internal procedures** but is not accessible elsewhere.

### 2.2.2 Program

- 1 A **program** shall consist of exactly one **main program**, any number (including zero) of other kinds of **program units**, any number (including zero) of **external procedures**, and any number (including zero) of other entities defined by means other than Fortran. The **main program** shall be defined by a Fortran *main-program program-unit* or by means other than Fortran, but not both.

### 2.2.3 Procedure

- 1 A procedure is either a function or a subroutine. Invocation of a function in an expression causes a value to be computed which is then used in evaluating the expression.

1 2 A procedure that is not pure might change the program state by changing the value of accessible [data objects](#) or  
2 [procedure pointers](#).

3 3 Procedures are described further in Clause [12](#).

## 4 2.2.4 Module

5 1 A [module](#) contains (or accesses from other modules) definitions that are to be made accessible to other [program](#)  
6 [units](#). These definitions include [data object declarations](#), type definitions, procedure definitions, and [interface](#)  
7 [blocks](#). Modules are further described in Clause [11](#).

## 8 2.2.5 Submodule

9 1 A [submodule](#) extends a [module](#) or another [submodule](#).

10 2 It may provide definitions ([12.6](#)) for procedures whose [interfaces](#) are declared ([12.4.3.2](#)) in an ancestor module  
11 or submodule. It may also contain declarations and definitions of other entities, which are accessible in its  
12 [descendants](#). An entity declared in a submodule is not accessible by [use association](#) unless it is a module procedure  
13 whose [interface](#) is declared in the ancestor module. Submodules are further described in Clause [11](#).

### NOTE 2.3

A submodule has access to entities in its parent module or submodule by [host association](#).

## 14 2.3 Execution concepts

### 15 2.3.1 Statement classification

16 1 Each Fortran statement is classified as either an [executable statement](#) or a [nonexecutable statement](#).

17 2 An [executable statement](#) is an instruction to perform or control an action. Thus, the executable statements of a  
18 [program unit](#) determine the behavior of the [program unit](#).

19 3 [Nonexecutable statements](#) are used to configure the program environment in which actions take place.

### 20 2.3.2 Statement order

Table 2.1: **Requirements on statement ordering**

PROGRAM, FUNCTION, SUBROUTINE, MODULE, SUBMODULE, or BLOCK DATA statement		
USE statements		
IMPORT statements		
FORMAT and ENTRY statements	IMPLICIT NONE	
	PARAMETER statements	IMPLICIT statements
	PARAMETER and DATA statements	Derived-type definitions, interface blocks, type declaration statements, enumeration definitions, procedure declarations, specification statements, and statement function statements
	DATA statements	Executable constructs

## Requirements on statement ordering (cont.)

CONTAINS statement
Internal subprograms or module subprograms
END statement

1 The syntax rules of subclause 2.1 specify the statement order within **program units** and **subprograms**. These  
 2 rules are illustrated in Table 2.1 and Table 2.2. Table 2.1 shows the ordering rules for statements and applies  
 3 to all **program units**, **subprograms**, and **interface bodies**. Vertical lines delineate varieties of statements that  
 4 may be interspersed and horizontal lines delineate varieties of statements that shall not be interspersed. **Internal**  
 5 or **module subprograms** shall follow a **CONTAINS statement**. Between **USE** and **CONTAINS** statements in a  
 6 **subprogram**, nonexecutable statements generally precede executable statements, although the **ENTRY statement**,  
 7 **FORMAT statement**, and **DATA statement** may appear among the executable statements. Table 2.2 shows which  
 8 statements are allowed in some kinds of **scoping units**.

Table 2.2: Statements allowed in scoping units

Statement type	Kind of scoping unit						
	Main program	Module or submodule	Block data	External subprogram	Module subprogram	Internal subprogram	Interface body
<b>USE</b>	Yes	Yes	Yes	Yes	Yes	Yes	Yes
<b>IMPORT</b>	No	No	No	No	No	No	Yes
<b>ENTRY</b>	No	No	No	Yes	Yes	No	No
<b>FORMAT</b>	Yes	No	No	Yes	Yes	Yes	No
Misc. decl.s <sup>1</sup>	Yes	Yes	Yes	Yes	Yes	Yes	Yes
<b>DATA</b>	Yes	Yes	Yes	Yes	Yes	Yes	No
Derived-type	Yes	Yes	Yes	Yes	Yes	Yes	Yes
Interface	Yes	Yes	No	Yes	Yes	Yes	Yes
Executable	Yes	No	No	Yes	Yes	Yes	No
<b>CONTAINS</b>	Yes	Yes	No	Yes	Yes	No	No
<b>Statement function</b>	Yes	No	No	Yes	Yes	Yes	No

(1) Miscellaneous declarations are **PARAMETER statements**, **IMPLICIT statements**, **type declaration statements**, enumeration definitions, **procedure declaration statements**, and specification statements.

### 2.3.3 The END statement

1 Each **program unit**, **module subprogram**, and **internal subprogram** shall have exactly one **END statement**. The  
 2 *end-program-stmt*, *end-function-stmt*, *end-subroutine-stmt*, and *end-mp-subprogram-stmt* statements are execut-  
 3 able, and may be **branch target statements** (8.2). Executing an *end-program-stmt* initiates **normal termination**.  
 4 Executing an *end-function-stmt*, *end-subroutine-stmt*, or *end-mp-subprogram-stmt* is equivalent to executing a  
 5 *return-stmt* with no *scalar-int-expr*.

6 The *end-module-stmt*, *end-submodule-stmt*, and *end-block-data-stmt* statements are nonexecutable.

### 2.3.4 Program execution

1 Execution of a program consists of the asynchronous execution of a fixed number (which may be one) of its **images**.  
 2 Each **image** has its own execution state, floating-point status (14.7), and set of **data objects**, **input/output units**,  
 3 and procedure pointers. The **image index** that identifies an **image** is an integer value in the range one to the  
 4 number of **images**.

- 1 2 A team is an ordered set of [images](#) that is either the initial team, consisting of all images, or a subset of a parent  
2 team formed by execution of a [FORM TEAM statement](#). The initial team has no parent; every other team has a  
3 unique parent team. Within the parent team, each new team is identified by its team number; this is the integer  
4 value that was specified in the [FORM TEAM statement](#).
- 5 3 During execution, each image has a current team, which is only changed by execution of [CHANGE TEAM](#)  
6 and [END TEAM](#) statements. Image indices, and thus coindexing of variable names with an *image-selector*, are  
7 relative to the current team unless a different team is specified. Initially, the current team is the initial team.

**NOTE 2.4**

Fortran control constructs ([8.1](#), [8.2](#)) control the progress of execution in each [image](#). [Image control statements](#) ([8.6.1](#)) affect the relative progress of execution between [images](#). [Coarrays](#) ([2.4.7](#)) provide a mechanism for accessing data on one [image](#) from another [image](#).

**NOTE 2.5**

A processor might allow the number of [images](#) to be chosen at compile time, link time, or run time. It might be the same as the number of CPUs but this is not required. Compiling for a single [image](#) might permit the optimizer to eliminate overhead associated with parallel execution. A program that makes assumptions about the exact number of [images](#) is unlikely to be portable.

## 8 2.3.5 Execution sequence

- 9 1 Following the creation of a fixed number of instances of the program, execution begins on each [image](#). [Image](#)  
10 execution is a sequence, in time, of actions. Actions take place during execution of the statement that performs  
11 them (except when explicitly stated otherwise). Segments ([8.6.2](#)) executed by a single [image](#) are totally ordered,  
12 and segments executed by separate [images](#) are partially ordered by [image control statements](#) ([8.6.1](#)).
- 13 2 If the program contains a Fortran [main program](#), each [image](#) begins execution with the first executable construct  
14 of the [main program](#). The execution of a [main program](#) or [subprogram](#) involves execution of the executable  
15 constructs within its [scoping unit](#). When a Fortran procedure is invoked, the [specification expressions](#) within  
16 the *specification-part* of the invoked procedure, if any, are evaluated in a processor dependent order. Thereafter,  
17 execution proceeds to the first executable construct appearing within the [scoping unit](#) of the procedure after the  
18 invoked entry point. With the following exceptions, the effect of execution is as if the executable constructs are  
19 executed in the order in which they appear in the [main program](#) or [subprogram](#) until a [STOP](#), [ERROR STOP](#),  
20 [RETURN](#), or [END](#) statement is executed.
  - 21 • Execution of a branching statement ([8.2](#)) changes the execution sequence. These statements explicitly  
22 specify a new starting place for the execution sequence.
  - 23 • [DO constructs](#), [IF constructs](#), [SELECT CASE constructs](#), [SELECT RANK constructs](#), and [SELECT TYPE](#)  
24 [constructs](#) contain an internal statement structure and execution of these constructs involves implicit in-  
25 ternal transfer of control. See Clause 8 for the detailed semantics of each of these constructs.
  - 26 • A [BLOCK construct](#) may contain [specification expressions](#); see [8.1.4](#) for detailed semantics of this construct.
  - 27 • An [END=](#), [ERR=](#), or [EOR=](#) specifier ([9.11](#)) might result in a branch.
  - 28 • An alternate return might result in a branch.

## 29 2.3.6 Image execution states

- 30 1 There are three image execution states: active, stopped, and failed. An image that has initiated [normal termin-](#)  
31 [ation](#) of execution is a [stopped image](#). An image that has ceased participating in program execution but has not  
32 initiated termination is a [failed image](#). All other images are active.
- 33 2 A [failed image](#) remains failed for the remainder of the execution of the program. The conditions that cause an  
34 image to fail are processor dependent.
- 35 3 Defining a coindexed object on a [failed image](#) has no effect other than defining the *stat-variable*, if one appears,

with the value `STAT_FAILED_IMAGE` (13.10.2.26). The value of a reference to a coindexed object on a `failed image` is processor dependent. Execution continues after such a reference.

When an image fails during the execution of a segment, a data object on a nonfailed image becomes undefined if it is not a `lock variable` and it might be defined or become undefined by execution of a statement of the segment other than an invocation of an `atomic subroutine` with the object as an `actual argument` corresponding to the ATOM `dummy argument`.

## 2.3.7 Termination of execution

Termination of execution of a program is either normal termination or error termination. Normal termination occurs only when all `images` initiate normal termination and occurs in three steps: initiation, synchronization, and completion. In this case, all `images` synchronize execution at the second step so that no image starts the completion step until all `images` have finished the initiation step. Error termination occurs when any `image` initiates error termination. Once error termination has been initiated on an `image`, error termination is initiated on all `images` that have not already initiated error termination. Termination of execution of the program occurs when all `images` have terminated execution.

Normal termination of execution of an `image` is initiated when a `STOP statement` or `end-program-stmt` is executed. Normal termination of execution of an `image` also may be initiated during execution of a procedure defined by a `companion processor` (ISO/IEC 9899:2011 5.1.2.2.3 and 7.22.4.4). If normal termination of execution is initiated within a Fortran `program unit` and the program incorporates procedures defined by a `companion processor`, the process of execution termination shall include the effect of executing the C `exit()` function (ISO/IEC 9899:2011 7.22.4.4) during the completion step.

Error termination of execution of an `image` is initiated if an `ERROR STOP statement` is executed or as specified elsewhere in this part of ISO/IEC 1539. When error termination on an `image` has been initiated, the processor should initiate error termination on other `images` as quickly as possible.

If the processor supports the concept of a process exit status, it is recommended that error termination initiated other than by an `ERROR STOP statement` supplies a processor-dependent nonzero value as the process exit status.

### NOTE 2.6

As well as in the circumstances specified in this part of ISO/IEC 1539, error termination might be initiated by means other than Fortran.

### NOTE 2.7

If an `image` has initiated normal termination, its data remain available for possible reference or definition by other `images` that are still executing.

## 2.4 Data concepts

### 2.4.1 Type

A `type` is a named categorization of data that, together with its `type parameters`, determines the set of values, syntax for denoting these values, and the set of operations that interpret and manipulate the values. This central concept is described in 4.1.

A `type` is either an `intrinsic type` or a `derived type`.

#### 2.4.1.1 Intrinsic type

The `intrinsic types` are integer, real, complex, character, and logical. The properties of `intrinsic types` are described in 4.4.

1 2 All **intrinsic types** have a **kind type parameter** called KIND, which determines the representation method for the  
 2 specified type. The **intrinsic type** character also has a **length type parameter** called LEN, which determines the  
 3 length of the character string.

#### 4 2.4.1.2 Derived type

5 1 **Derived types** may be parameterized. A scalar **object** of **derived type** is a **structure**; assignment of structures  
 6 is defined intrinsically (7.2.1.3), but there are no **intrinsic** operations for structures. For each **derived type**, a  
 7 **structure constructor** is available to create values (4.5.10). In addition, **objects** of **derived type** may be used as  
 8 procedure arguments and function results, and may appear in input/output lists. If additional operations are  
 9 needed for a **derived type**, they shall be defined by procedures (7.1.6).

10 2 **Derived types** are described further in 4.5.

### 11 2.4.2 Data value

12 1 Each **intrinsic type** has associated with it a set of values that a datum of that type may take, depending on the  
 13 values of the type parameters. The values for each **intrinsic type** are described in 4.4. The values that **objects** of  
 14 a **derived type** may assume are determined by the type definition, type parameter values, and the sets of values  
 15 of its components.

### 16 2.4.3 Data entity

#### 17 2.4.3.1 General

18 1 A **data entity** has a type and type parameters; it might have a data value (an exception is an undefined variable).  
 19 Every **data entity** has a **rank** and is thus either a scalar or an array.

20 2 A **data entity** that is the result of the execution of a **function reference** is called the function result.

#### 21 2.4.3.2 Data object

22 1 A **data object** is either a constant, variable, or a **subobject** of a constant. The type and type parameters of a  
 23 named **data object** may be specified explicitly (5.2) or implicitly (5.7).

24 2 **Subobjects** are portions of **data objects** that may be referenced and defined (variables only) independently of the  
 25 other portions.

26 3 These include portions of arrays (array elements and array sections), portions of character strings (substrings),  
 27 portions of complex **objects** (real and imaginary parts), and portions of **structures** (components). **Subobjects**  
 28 are themselves **data objects**, but **subobjects** are referenced only by **object designators** or **intrinsic** functions. A  
 29 **subobject** of a variable is a variable. **Subobjects** are described in Clause 6.

30 4 The following **objects** are referenced by a name:

- 31 • a named scalar (a scalar object);
- a named array (an array object).

32 5 The following **subobjects** are referenced by an **object designator**:

- an array element (a scalar subobject);
- an array section (an array subobject);
- 33 • a **complex part designator** (the real or imaginary part of a complex object);
- a structure component (a scalar or an array subobject);
- a substring (a scalar subobject).

#### 34 2.4.3.2.1 Variable

35 1 A **variable** can have a value or be undefined; during execution of a program it can be defined and redefined.

- 1 2 A [local variable](#) of a [module](#), [submodule](#), [main program](#), [subprogram](#), or [BLOCK construct](#) is accessible only in  
2 that [scoping unit](#) or construct and in any contained [scoping units](#) and constructs.

**NOTE 2.8**

A [subobject](#) of a [local variable](#) is also a [local variable](#).

A [local variable](#) cannot be in COMMON or have the [BIND attribute](#), because [common blocks](#) and [variables](#) with the [BIND attribute](#) are global entities.

3 **2.4.3.2.2 Constant**

- 4 1 A constant is either a [named constant](#) or a [literal constant](#).  
5 2 Named constants are defined using the [PARAMETER attribute](#) (5.5.13, 5.6.11). The syntax of literal constants  
6 is described in 4.4.

7 **2.4.3.2.3 Subobject of a constant**

- 8 1 A [subobject](#) of a [constant](#) is a portion of a [constant](#).  
9 2 In an [object designator](#) for a [subobject](#) of a [constant](#), the portion referenced may depend on the value of a [variable](#).

**NOTE 2.9**

For example, given:

```
CHARACTER (LEN = 10), PARAMETER :: DIGITS = '0123456789'
CHARACTER (LEN = 1)           :: DIGIT
INTEGER :: I
...
DIGIT = DIGITS (I:I)
```

DIGITS is a [named constant](#) and DIGITS (I:I) designates a [subobject](#) of the [constant](#) DIGITS.

10 **2.4.3.3 Expression**

- 11 1 An expression (7.1) produces a [data entity](#) when evaluated. An expression represents either a [data object reference](#)  
12 or a computation; it is formed from operands, operators, and parentheses. The type, type parameters, value, and  
13 [rank](#) of an expression result are determined by the rules in Clause 7.

14 **2.4.3.4 Function reference**

- 15 1 A [function reference](#) produces a [data entity](#) when the function is executed during expression evaluation. The  
16 type, type parameters, and [rank](#) of a function result are determined by the [interface](#) of the function (12.3.3). The  
17 value of a function result is determined by execution of the function.

18 **2.4.4 Definition of objects and pointers**

- 19 1 When an [object](#) is given a valid value during program execution, it becomes [defined](#). This is often accomplished  
20 by execution of an [assignment](#) or [input](#) statement. When a variable does not have a predictable value, it is  
21 [undefined](#).  
22 2 Similarly, when a pointer is associated with a [target](#) or nullified, its [pointer association](#) status becomes [defined](#).  
23 When the association status of a pointer is not predictable, its [pointer association](#) status is [undefined](#).  
24 3 Clause 16 describes the ways in which variables become [defined](#) and [undefined](#) and the association status of  
25 pointers becomes [defined](#) and [undefined](#).

## 2.4.5 Reference

- 1 A **data object** is **referenced** when its value is required during execution. A procedure is **referenced** when it is executed.
- 2 The appearance of a **data object designator** or **procedure designator** as an **actual argument** does not constitute a **reference** to that **data object** or procedure unless such a **reference** is necessary to complete the specification of the **actual argument**.

## 2.4.6 Array

- 1 An array may have up to fifteen dimensions, and any **extent** in any dimension. The **size** of an array is the total number of elements, which is equal to the product of the **extents**. An array may have zero size. The **shape** of an array is determined by its **rank** and its **extent** in each dimension, and is represented as a rank-one array whose elements are the extents. All named arrays shall be declared, and the **rank** of a named array is specified in its declaration. The **rank** of a named array, once declared, is constant; the extents may be constant or may vary during execution.
- 2 Any **intrinsic** operation defined for scalar **objects** may be applied to **conformable objects**. Such operations are performed **elementally** to produce a resultant array **conformable** with the array operands. If an **elemental** operation is intrinsically pure or is implemented by a pure **elemental function** (12.8), the element operations may be performed simultaneously or in any order.
- 3 A rank-one array may be constructed from scalars and other arrays and may be reshaped into any allowable array shape (4.8).
- 4 Arrays may be of any type and are described further in 6.5.

## 2.4.7 Coarray

- 1 A **coarray** is a **data entity** that has nonzero **corank**; it can be directly referenced or defined by other **images**. It may be a scalar or an array.
- 2 For each **coarray** on an **image**, there is a corresponding **coarray** with the same type, type parameters, and **bounds** on every other **image** that is in the same **team**.
- 3 The set of corresponding **coarrays** on all **images** in a **team** is arranged in a rectangular pattern. The dimensions of this pattern are the **codimensions**; the number of **codimensions** is the **corank**. The bounds for each **codimension** are the **cobounds**.

### NOTE 2.10

If the total number of **images** is not a multiple of the product of the sizes of each but the rightmost of the **codimensions**, the rectangular pattern will be incomplete.

- 4 A **coarray** on any **image** can be accessed directly by using **cosubscripts**. On its own **image**, a **coarray** can also be accessed without use of **cosubscripts**.
- 5 A **subobject** of a **coarray** is a **coarray** if it does not have any **cosubscripts**, **vector subscripts**, **allocatable** component selection, or pointer component selection.
- 6 For a **coindexed object**, its **cosubscript** list determines the **image index** in the same way that a subscript list determines the subscript order value for an **array element** (6.5.3.2).
- 7 **Intrinsic** procedures are provided for mapping between an **image index** and a list of **cosubscripts**.

**NOTE 2.11**

The mechanism for an **image** to reference and define a **coarray** on another **image** might vary according to the hardware. On a shared-memory machine, a **coarray** on an **image** and the corresponding **coarrays** on other **images** could be implemented as a sequence of arrays with evenly spaced starting addresses. On a distributed-memory machine with separate physical memory for each **image**, a processor might store a **coarray** at the same virtual address in each physical memory.

**NOTE 2.12**

Except in contexts where **coindexed objects** are disallowed, accessing a **coarray** on its own **image** by using a set of **cosubscripts** that specify that **image** has the same effect as accessing it without **cosubscripts**. In particular, the segment ordering rules (8.6.2) apply whether or not **cosubscripts** are used to access the **coarray**.

**2.4.8 Established coarrays**

- 1 A nonallocatable **coarray** that is not a **dummy argument** or **host associated** with a **dummy argument**, and not a **local variable** of a **subprogram** or **BLOCK construct**, is established in the **initial team**.
- 2 An allocated allocatable **coarray** is established in the **team** in which it was allocated. An unallocated allocatable **coarray** is not established.
- 3 A **coarray** that is an associating entity in a **coarray-association** of a **CHANGE TEAM statement** is established in the **team** of its **CHANGE TEAM construct**.
- 4 A nonallocatable **coarray** that is a **dummy argument** or **host associated** with a **dummy argument** is established in the **team** in which the procedure was invoked. A **coarray dummy argument** is not established in any **ancestor team** even if the corresponding **actual argument** is established in one or more of them.
- 5 A nonallocatable **coarray** that is a **local variable** of a **subprogram** or **host associated** with a **local variable** of a **subprogram** is established in the **team** in which the procedure was invoked. A nonallocatable **coarray** declared in a **BLOCK construct** is established in the **team** in which the **BLOCK statement** was executed.

**2.4.9 Pointer**

- 1 A **pointer** has an association status which is either associated, **disassociated**, or undefined (16.5.2.2).
- 2 A **pointer** that is not associated shall not be referenced or defined.
- 3 If a **data pointer** is an **array**, the **rank** is declared, but the **bounds** are determined when it is associated with a **target**.

**2.4.10 Allocatable variables**

- 1 The allocation status of an **allocatable** variable is either allocated or unallocated. An **allocatable** variable becomes allocated as described in 6.7.1.3. It becomes unallocated as described in 6.7.3.2.
- 2 An unallocated **allocatable** variable shall not be **referenced** or **defined**.
- 3 If an **allocatable** variable is an **array**, the **rank** is declared, but the **bounds** are determined when it is allocated. If an **allocatable** variable is a **coarray**, the **corank** is declared, but the **cobounds** are determined when it is allocated.

**2.4.11 Storage**

- 1 Many of the facilities of this part of ISO/IEC 1539 make no assumptions about the physical storage characteristics of **data objects**. However, **program units** that include **storage association** dependent features shall observe the storage restrictions described in 16.5.3.



- 1 2 If a [procedure](#) is defined by means of a [companion processor](#) that is not the [Fortran processor](#) itself, this part of  
2 ISO/IEC 1539 refers to the C function that defines the [procedure](#), although the [procedure](#) need not be defined  
3 by means of the C programming language.

**NOTE 2.14**

A [companion processor](#) might or might not be a mechanism that conforms to the requirements of ISO/IEC 9899:2011.

For example, a [processor](#) might allow a [procedure](#) defined by some language other than Fortran or C to be invoked if it can be described by a C prototype as defined in 6.7.6.3 of ISO/IEC 9899:2011.

## 3 Lexical tokens and source form

### 3.1 Processor character set

#### 3.1.1 Characters

1 The processor character set is processor dependent. Each character in a processor character set is either a control character or a graphic character. The set of graphic characters is further divided into letters (3.1.2), digits (3.1.3), underscore (3.1.4), special characters (3.1.5), and other characters (3.1.6).

2 The letters, digits, underscore, and special characters make up the Fortran character set. Together, the set of letters, digits, and underscore define the syntax class *alphanumeric-character*.

R301 *alphanumeric-character* is *letter*  
 or *digit*  
 or *underscore*

3 Except for the currency symbol, the graphics used for the characters shall be as given in 3.1.2, 3.1.3, 3.1.4, and 3.1.5. However, the style of any graphic is not specified.

#### 3.1.2 Letters

1 The twenty-six letters are:

2 A B C D E F G H I J K L M N O P Q R S T U V W X Y Z

3 The set of letters defines the syntactic class *letter*. The processor character set shall include lower-case and upper-case letters. A lower-case letter is equivalent to the corresponding upper-case letter in *program units* except in a *character context* (1.3.22).

##### NOTE 3.1

The following statements are equivalent:

```
CALL BIG_COMPLEX_OPERATION (NDATE)
call big_complex_operation (ndate)
Call Big_Complex_Operation (NDate)
```

#### 3.1.3 Digits

1 The ten digits are:

2 0 1 2 3 4 5 6 7 8 9

3 The ten digits define the syntactic class *digit*.

#### 3.1.4 Underscore

R302 *underscore* is `_`

#### 3.1.5 Special characters

1 The special characters are shown in Table 3.1.

Table 3.1: Special characters

Character	Name of character	Character	Name of character
	Blank	;	Semicolon
=	Equals	!	Exclamation point
+	Plus	"	Quotation mark or quote
-	Minus	%	Percent
*	Asterisk	&	Ampersand
/	Slash	~	Tilde
\	Backslash	<	Less than
(	Left parenthesis	>	Greater than
)	Right parenthesis	?	Question mark
[	Left square bracket	'	Apostrophe
]	Right square bracket	`	Grave accent
{	Left curly bracket	^	Circumflex accent
}	Right curly bracket		Vertical line
,	Comma	\$	Currency symbol
.	Decimal point or period	#	Number sign
:	Colon	@	Commercial at

- 1 2 Some of the special characters are used for operator symbols, bracketing, and various forms of separating and  
2 delimiting other lexical tokens.

### 3 3.1.6 Other characters

- 4 1 Additional characters may be representable in the processor, but may appear only in comments (3.3.2.3, 3.3.3.2),  
5 character constants (4.4.4), input/output records (9.2.2), and character string edit descriptors (10.3.2).

## 6 3.2 Low-level syntax

### 7 3.2.1 Tokens

- 8 1 The low-level syntax describes the fundamental lexical tokens of a [program unit](#). A [lexical token](#) is a keyword,  
9 name, literal constant other than a complex literal constant, operator, statement label, delimiter, comma, =, =>  
10 :, ::, ; or %.

### 11 3.2.2 Names

- 12 1 [Names](#) are used for various entities such as [variables](#), [program units](#), [dummy arguments](#), [named constants](#), and  
13 derived types.

14 R303 *name* is letter [ [alphanumeric-character](#) ] ...

15 C301 (R303) The maximum length of a [name](#) is 63 characters.

#### NOTE 3.2

Examples of names:

```
A1
NAME_LENGTH      (single underscore)
S_P_R_E_A_D__O_U_T (two consecutive underscores)
TRAILER_         (trailing underscore)
```

**NOTE 3.3**

The word “name” always denotes this particular syntactic form. The word “identifier” is used where entities can be identified by other syntactic forms or by values; its particular meaning depends on the context in which it is used.

**3.2.3 Constants**

R304	<i>constant</i>	is	<i>literal-constant</i>
		or	<i>named-constant</i>
R305	<i>literal-constant</i>	is	<i>int-literal-constant</i>
		or	<i>real-literal-constant</i>
		or	<i>complex-literal-constant</i>
		or	<i>logical-literal-constant</i>
		or	<i>char-literal-constant</i>
		or	<i>boz-literal-constant</i>
R306	<i>named-constant</i>	is	<i>name</i>
R307	<i>int-constant</i>	is	<i>constant</i>
C302	(R307) <i>int-constant</i> shall be of type integer.		

**3.2.4 Operators**

R308	<i>intrinsic-operator</i>	is	<i>power-op</i>
		or	<i>mult-op</i>
		or	<i>add-op</i>
		or	<i>concat-op</i>
		or	<i>rel-op</i>
		or	<i>not-op</i>
		or	<i>and-op</i>
		or	<i>or-op</i>
		or	<i>equiv-op</i>
R707	<i>power-op</i>	is	<b>**</b>
R708	<i>mult-op</i>	is	<b>*</b>
		or	<b>/</b>
R709	<i>add-op</i>	is	<b>+</b>
		or	<b>−</b>
R711	<i>concat-op</i>	is	<b>//</b>
R713	<i>rel-op</i>	is	<b>.EQ.</b>
		or	<b>.NE.</b>
		or	<b>.LT.</b>
		or	<b>.LE.</b>
		or	<b>.GT.</b>
		or	<b>.GE.</b>
		or	<b>==</b>
		or	<b>/=</b>
		or	<b>&lt;</b>
		or	<b>&lt;=</b>
		or	<b>&gt;</b>
		or	<b>&gt;=</b>

1	R718	<i>not-op</i>	is	<i>.NOT.</i>
2	R719	<i>and-op</i>	is	<i>.AND.</i>
3	R720	<i>or-op</i>	is	<i>.OR.</i>
4	R721	<i>equiv-op</i>	is	<i>.EQV.</i>
5			or	<i>.NEQV.</i>
6	R309	<i>defined-operator</i>	is	<i>defined-unary-op</i>
7			or	<i>defined-binary-op</i>
8			or	<i>extended-intrinsic-op</i>
9	R703	<i>defined-unary-op</i>	is	<i>. letter [ letter ] ... .</i>
10	R723	<i>defined-binary-op</i>	is	<i>. letter [ letter ] ... .</i>
11	R310	<i>extended-intrinsic-op</i>	is	<i>intrinsic-operator</i>

### 3.2.5 Statement labels

- 1 A *statement label* provides a means of referring to an individual statement.
- R311 *label* is *digit [ digit [ digit [ digit [ digit ] ] ] ]*
- C303 (R311) At least one digit in a *label* shall be nonzero.
- 2 If a statement is labeled, the statement shall contain a nonblank character. The same statement label shall not be given to more than one statement in its scope. Leading zeros are not significant in distinguishing between statement labels. There are 99999 possible unique statement labels and a processor shall accept any of them as a statement label. However, a processor may have a limit on the total number of unique statement labels in one *program unit*.

#### NOTE 3.4

For example:

99999  
10  
010

are all statement labels. The last two are equivalent.

- 3 Any statement that is not part of another statement, and that is not preceded by a semicolon in fixed form, may begin with a statement label, but the labels are used only in the following ways.
- The label on a *branch target statement* (8.2) is used to identify that statement as the possible destination of a branch.
  - The label on a *FORMAT statement* (10.2.1) is used to identify that statement as the format specification for a *data transfer statement* (9.6).
  - In some forms of the *DO construct* (8.1.7), the terminal statement of the construct is identified by a label.

### 3.2.6 Delimiters

- 1 A *lexical token* that is a delimiter is a (, ), /, [, ], (/, or /).

## 3.3 Source form

### 3.3.1 Program units, statements, and lines

- 1 A Fortran **program unit** is a sequence of one or more **lines**, organized as Fortran statements, comments, and **INCLUDE lines**. A **line** is a sequence of zero or more characters. **Lines** following a **program unit END statement** are not part of that **program unit**. A Fortran **statement** is a sequence of one or more complete or partial **lines**.
- 2 A comment may contain any character that may occur in any **character context**.
- 3 There are two source forms. Subclause 3.3.2 applies only to free form source. Subclause 3.3.3 applies only to fixed source form. Free form and fixed form shall not be mixed in the same program unit. The means for specifying the source form of a **program unit** are processor dependent.

### 3.3.2 Free source form

#### 3.3.2.1 Free form line length

- 1 In free source form there are no restrictions on where a statement (or portion of a statement) may appear within a **line**. A **line** may contain zero characters. If a **line** consists entirely of characters of default kind (4.4.4), it may contain at most 132 characters. If a **line** contains any character that is not of default kind, the maximum number of characters allowed on the **line** is processor dependent.

#### 3.3.2.2 Blank characters in free form

- 1 In free source form blank characters shall not appear within lexical tokens other than in a **character context** or in a format specification. Blanks may be inserted freely between tokens to improve readability; for example, blanks may occur between the tokens that form a complex literal constant. A sequence of blank characters outside of a **character context** is equivalent to a single blank character.
- 2 A blank shall be used to separate names, constants, or labels from adjacent keywords, names, constants, or labels.

#### NOTE 3.5

For example, the blanks after REAL, READ, 30, and DO are required in the following:

```
REAL X
READ 10
30 DO K=1,3
```

- 3 One or more blanks shall be used to separate adjacent keywords except in the following cases, where blanks are optional:

Table 3.2: Adjacent keywords where separating blanks are optional

BLOCK DATA	END ENUM	END SELECT
DOUBLE PRECISION	END FILE	END SUBMODULE
ELSE IF	END FORALL	END SUBROUTINE
ELSE WHERE	END FUNCTION	END TYPE
END ASSOCIATE	END IF	END WHERE
END BLOCK	END INTERFACE	GO TO
END BLOCK DATA	END MODULE	IN OUT
END CRITICAL	END PROCEDURE	SELECT CASE
END DO	END PROGRAM	SELECT TYPE

### 3.3.2.3 Free form commentary

- 1 The character “!” initiates a comment except where it appears within a [character context](#). The comment extends to the end of the [line](#). If the first nonblank character on a [line](#) is an “!”, the [line](#) is a comment line. [Lines](#) containing only blanks or containing no characters are also comment lines. Comments may appear anywhere in a [program unit](#) and may precede the first statement of a [program unit](#) or may follow the last statement of a [program unit](#). Comments have no effect on the interpretation of the [program unit](#).

#### NOTE 3.6

This part of ISO/IEC 1539 does not restrict the number of consecutive comment lines.

### 3.3.2.4 Free form statement continuation

- 1 The character “&” is used to indicate that the statement is continued on the next [line](#) that is not a comment line. Comment lines cannot be continued; an “&” in a comment has no effect. Comments may occur within a continued statement. When used for continuation, the “&” is not part of the statement. No [line](#) shall contain a single “&” as the only nonblank character or as the only nonblank character before an “!” that initiates a comment.
- 2 If a non[character context](#) is to be continued, an “&” shall be the last nonblank character on the [line](#), or the last nonblank character before an “!”. There shall be a later [line](#) that is not a comment; the statement is continued on the next such [line](#). If the first nonblank character on that [line](#) is an “&”, the statement continues at the next character position following that “&”; otherwise, it continues with the first character position of that [line](#).
- 3 If a lexical token is split across the end of a [line](#), the first nonblank character on the first following noncomment [line](#) shall be an “&” immediately followed by the successive characters of the split token.
- 4 If a [character context](#) is to be continued, an “&” shall be the last nonblank character on the [line](#) and shall not be followed by commentary. There shall be a later [line](#) that is not a comment; an “&” shall be the first nonblank character on the next such [line](#) and the statement continues with the next character following that “&”.

### 3.3.2.5 Free form statement termination

- 1 If a statement is not continued, a comment or the end of the [line](#) terminates the statement.
- 2 A statement may alternatively be terminated by a “;” character that appears other than in a [character context](#) or in a comment. The “;” is not part of the statement. After a “;” terminator, another statement may appear on the same [line](#), or begin on that [line](#) and be continued. A sequence consisting only of zero or more blanks and one or more “;” terminators, in any order, is equivalent to a single “;” terminator.

### 3.3.2.6 Free form statements

- 1 A label may precede any statement not forming part of another statement.

#### NOTE 3.7

No Fortran statement begins with a digit.

- 2 A statement shall not have more than 255 continuation [lines](#).

## 3.3.3 Fixed source form

### 3.3.3.1 General

- 1 In fixed source form, there are restrictions on where a statement may appear within a [line](#). If a source [line](#) contains only characters of default kind, it shall contain exactly 72 characters; otherwise, its maximum number of characters is processor dependent.
- 2 Except in a [character context](#), blanks are insignificant and may be used freely throughout the program.

### 3.3.3.2 Fixed form commentary

- 1 The character “!” initiates a comment except where it appears within a [character context](#) or in character position 6. The comment extends to the end of the [line](#). If the first nonblank character on a [line](#) is an “!” in any character position other than character position 6, the [line](#) is a comment line. [Lines](#) beginning with a “C” or “\*” in character position 1 and lines containing only blanks are also comment lines. Comments may appear anywhere in a [program unit](#) and may precede the first statement of the [program unit](#) or may follow the last statement of a [program unit](#). Comments have no effect on the interpretation of the [program unit](#).

#### NOTE 3.8

This part of ISO/IEC 1539 does not restrict the number of consecutive comment lines.

### 3.3.3.3 Fixed form statement continuation

- 1 Except within commentary, character position 6 is used to indicate continuation. If character position 6 contains a blank or zero, the [line](#) is the initial [line](#) of a new statement, which begins in character position 7. If character position 6 contains any character other than blank or zero, character positions 7–72 of the [line](#) constitute a continuation of the preceding noncomment [line](#).

#### NOTE 3.9

An “!” or “;” in character position 6 is interpreted as a continuation indicator unless it appears within commentary indicated by a “C” or “\*” in character position 1 or by an “!” in character positions 1–5.

- 2 Comment lines cannot be continued. Comment lines may occur within a continued statement.

### 3.3.3.4 Fixed form statement termination

- 1 If a statement is not continued, a comment or the end of the [line](#) terminates the statement.
- 2 A statement may alternatively be terminated by a “;” character that appears other than in a [character context](#), in a comment, or in character position 6. The “;” is not part of the statement. After a “;” terminator, another statement may begin on the same [line](#), or begin on that [line](#) and be continued. A “;” shall not appear as the first nonblank character on an initial [line](#). A sequence consisting only of zero or more blanks and one or more “;” terminators, in any order, is equivalent to a single “;” terminator.

### 3.3.3.5 Fixed form statements

- 1 A label, if it appears, shall occur in character positions 1 through 5 of the first [line](#) of a statement; otherwise, positions 1 through 5 shall be blank. Blanks may appear anywhere within a label. A statement following a “;” on the same [line](#) shall not be labeled. Character positions 1 through 5 of any continuation [lines](#) shall be blank. A statement shall not have more than 255 continuation [lines](#). The [program unit END statement](#) shall not be continued. A statement whose initial [line](#) appears to be a [program unit END statement](#) shall not be continued.

## 3.4 Including source text

- 1 Additional text may be incorporated into the source text of a [program unit](#) during processing. This is accomplished with the `INCLUDE`, which has the form

```
INCLUDE char-literal-constant
```

- 3 The *char-literal-constant* shall not have a kind type parameter value that is a *named-constant*.

- 4 An `INCLUDE` line is not a Fortran statement.

- 5 An `INCLUDE` line shall appear on a single source [line](#) where a statement may appear; it shall be the only nonblank text on this [line](#) other than an optional trailing comment. Thus, a statement label is not allowed.

- 6 The effect of the `INCLUDE` line is as if the referenced source text physically replaced the `INCLUDE` line prior to program processing. Included text may contain any source text, including additional `INCLUDE` lines; such nested `INCLUDE` lines are similarly replaced with the specified source text. The maximum depth of nesting of any nested `INCLUDE` lines is processor dependent. Inclusion of the source text referenced by an `INCLUDE` line shall not, at any level of nesting, result in inclusion of the same source text.

- 1 7 When an INCLUDE line is resolved, the first included statement *line* shall not be a continuation *line* and the last  
2 included statement *line* shall not be continued.
- 3 8 The interpretation of *char-literal-constant* is processor dependent. An example of a possible valid interpretation  
4 is that *char-literal-constant* is the name of a file that contains the source text to be included.

**NOTE 3.10**

In some circumstances, for example where source code is maintained in an INCLUDE file for use in programs whose source form might be either fixed or free, observing the following rules allows the code to be used with either source form.

- Confine statement labels to character positions 1 to 5 and statements to character positions 7 to 72.
- Treat blanks as being significant.
- Use only the exclamation mark (!) to indicate a comment, but do not start the comment in character position 6.
- For continued statements, place an ampersand (&) in both character position 73 of a continued *line* and character position 6 of a continuation *line*.

## 4 Types

### 4.1 Characteristics of types

#### 4.1.1 The concept of type

1 Fortran provides an abstract means whereby data can be categorized without relying on a particular physical representation. This abstract means is the concept of type.

2 A type has a name, a set of valid values, a means to denote such values (constants), and a set of operations to manipulate the values.

#### 4.1.2 Type classification

1 A type is either an [intrinsic type](#) or a [derived type](#).

2 This part of ISO/IEC 1539 defines five intrinsic types: integer, real, complex, character, and logical.

3 A derived type is one that is defined by a derived-type definition ([4.5.2](#)) or by an intrinsic module. It shall be used only where it is accessible ([4.5.2.2](#)). An intrinsic type is always accessible.

#### 4.1.3 Set of values

1 For each type, there is a set of valid values. The set of valid values for logical is completely determined by this part of ISO/IEC 1539. The sets of valid values for integer, character, and real are processor dependent. The set of valid values for complex consists of the set of all the combinations of the values of the individual components. The set of valid values for a derived type is as defined in [4.5.8](#).

#### 4.1.4 Constants

1 The syntax for denoting a value indicates the type, type parameters, and the particular value.

2 The syntax for literal constants of each intrinsic type is specified in [4.4](#).

3 A [structure constructor](#) ([4.5.10](#)) that is a [constant expression](#) ([7.1.12](#)) denotes a scalar constant value of derived type. An array constructor ([4.8](#)) that is a [constant expression](#) denotes a constant array value of intrinsic or derived type.

4 A constant value can be named ([5.5.13](#), [5.6.11](#)).

#### 4.1.5 Operations

1 For each of the intrinsic types, a set of operations and corresponding operators is defined intrinsically. These are described in Clause [7](#). The intrinsic set can be augmented with operations and operators defined by functions with the [OPERATOR](#) interface ([12.4.3.2](#)). Operator definitions are described in Clauses [7](#) and [12](#).

2 For derived types, there are no intrinsic operations. Operations on derived types can be defined by the program ([4.5.11](#)).

## 4.2 Type parameters

1 A type might be parameterized. In this case, the set of values, the syntax for denoting the values, and the set of operations on the values of the type depend on the values of the parameters.

2 A [type parameter](#) is either a [kind type parameter](#) or a [length type parameter](#). All type parameters are of type integer. A [kind type parameter](#) participates in generic resolution ([12.5.5.2](#)), but a [length type parameter](#) does not.

3 Each intrinsic type has a [kind type parameter](#) named KIND. The intrinsic character type has a [length type parameter](#) named LEN. A derived type can have type parameters.

4 A type parameter value may be specified by a type specification ([4.4](#), [4.5.9](#)).

```
R401  type-param-value      is  scalar-int-expr
                                or  *
                                or  :
```

(R401) The [type-param-value](#) for a kind type parameter shall be a [constant expression](#).

(R401) A colon shall not be used as a [type-param-value](#) except in the declaration of an entity or component that has the [POINTER](#) or [ALLOCATABLE](#) attribute.

5 A colon as a [type-param-value](#) specifies a [deferred type parameter](#).

6 The values of the [deferred type parameters](#) of an object are determined by successful execution of an [ALLOCATE statement](#) ([6.7.1](#)), execution of an [intrinsic assignment statement](#) ([7.2.1.3](#)), execution of a [pointer assignment statement](#) ([7.2.2](#)), or by [argument association](#) ([12.5.2](#)).

### NOTE 4.1

[Deferred type parameters](#) of functions, including function procedure pointers, have no values. Instead, they indicate that those [type parameters](#) of the function result will be determined by execution of the function, if it returns an allocated [allocatable](#) result or an associated pointer result.

7 An asterisk as a [type-param-value](#) specifies that a [length type parameter](#) is an [assumed type parameter](#). It is used for a [dummy argument](#) to assume the type parameter value from the [effective argument](#), for an [associate name](#) in a [SELECT TYPE construct](#) to assume the type parameter value from the corresponding selector, and for a [named constant](#) of type character to assume the character length from the [constant-expr](#).

### NOTE 4.2

The value of a [kind type parameter](#) is always known at compile time. Some parameterizations that involve multiple representation forms need to be distinguished at compile time for practical implementation and performance. Examples include the multiple precisions of the intrinsic real type and the possible multiple character sets of the intrinsic character type.

The adjective “length” is used for type parameters other than kind type parameters because they often specify a length, as for intrinsic character type. However, they can be used for other purposes. The important difference from kind type parameters is that their values need not be known at compile time and might change during execution.

## 4.3 Types, type specifiers, and values

### 4.3.1 Relationship of types and values to objects

1 The name of a type serves as a type specifier and may be used to declare objects of that type. A declaration specifies the type of a named object. A data object may be declared explicitly or implicitly. A data object has

1 *attributes* in addition to its type. Clause 5 describes the way in which a data object is declared and how its type  
2 and other attributes are specified.

3 2 Scalar data of any intrinsic or derived type may be shaped in a rectangular pattern to compose an array of the  
4 same type and type parameters. An array object has a type and type parameters just as a scalar object does.

5 3 A variable is a data object. The type and type parameters of a variable determine which values that variable  
6 may take. Assignment (7.2) provides one means of changing the value of a variable of any type.

7 4 The type of a variable determines the operations that may be used to manipulate the variable.

## 8 4.3.2 Type specifiers and type compatibility

### 9 4.3.2.1 Type specifier syntax

10 1 A type specifier specifies a type and type parameter values. It is either a *type-spec* or a *declaration-type-spec*.

11 R402 *type-spec* is *intrinsic-type-spec*  
12 or *derived-type-spec*

13 C403 (R402) The *derived-type-spec* shall not specify an *abstract type* (4.5.7).

14 R403 *declaration-type-spec* is *intrinsic-type-spec*  
15 or TYPE ( *intrinsic-type-spec* )  
16 or TYPE ( *derived-type-spec* )  
17 or CLASS ( *derived-type-spec* )  
18 or CLASS ( \* )  
19 or TYPE ( \* )

20 C404 (R403) In a *declaration-type-spec*, every *type-param-value* that is not a colon or an asterisk shall be a  
21 *specification-expr*.

22 C405 (R403) In a *declaration-type-spec* that uses the CLASS keyword, *derived-type-spec* shall specify an *ex-*  
23 *tensible type* (4.5.7).

24 C406 (R403) TYPE(*derived-type-spec*) shall not specify an *abstract type* (4.5.7).

25 C407 (R402) In TYPE(*intrinsic-type-spec*) the *intrinsic-type-spec* shall not end with a comma.

26 C408 An entity declared with the CLASS keyword shall be a *dummy argument* or have the *ALLOCATABLE*  
27 or *POINTER attribute*.

28 2 An *intrinsic-type-spec* specifies the named intrinsic type and its type parameter values. A *derived-type-spec*  
29 specifies the named derived type and its type parameter values.

#### NOTE 4.3

A *type-spec* is used in an array constructor, a SELECT TYPE construct, or an ALLOCATE statement. An *integer-type-spec* is used in a DO CONCURRENT or FORALL statement. Elsewhere, a *declaration-type-spec* is used.

### 30 4.3.2.2 TYPE type specifier

31 1 A TYPE type specifier is used to declare entities that are *assumed-type*, or of an intrinsic or derived type.

32 2 A *derived-type-spec* in a TYPE type specifier in a *type declaration statement* shall specify a previously defined  
33 derived type. If the data entity is a function result, the derived type may be specified in the *FUNCTION*  
34 *statement* provided the derived type is defined within the body of the function or is accessible there by use or  
35 *host* association. If the derived type is specified in the *FUNCTION statement* and is defined within the body  
36 of the function, it is as if the *function result* were declared with that derived type immediately following the

1 *derived-type-def* of the specified derived type.

2 3 An entity that is declared using the TYPE(\*) type specifier is **assumed-type** and is an **unlimited polymorphic**  
3 entity. Its **dynamic type** and **type parameters** are assumed from its **effective argument**.

4 C409 An **assumed-type** entity shall be a **dummy data object** that does not have the **ALLOCATABLE**, **CODI-**  
5 **MENTION**, **INTENT (OUT)**, **POINTER**, or **VALUE** attribute and is not an **explicit-shape array**.

6 C410 An **assumed-type** variable name shall not appear in a designator or expression except as an **actual**  
7 **argument** corresponding to a **dummy argument** that is **assumed-type**, or as the first argument to the  
8 **intrinsic** function **IS\_CONTIGUOUS**, **LBOUND**, **PRESENT**, **RANK**, **SHAPE**, **SIZE**, or **UBOUND**, or  
9 the function **C\_LOC** from the intrinsic module **ISO\_C\_BINDING**.

10 C411 An **assumed-type actual argument** that corresponds to an **assumed-rank dummy argument** shall be  
11 **assumed-shape** or **assumed-rank**.

#### 12 4.3.2.3 CLASS type specifier

13 1 The **CLASS** type specifier is used to declare **polymorphic** entities. A **polymorphic** entity is a data entity that is  
14 able to be of differing **dynamic types** during program execution.

15 2 The **declared type** of a **polymorphic** entity is the specified type if the **CLASS** type specifier contains a type name.

16 3 An entity declared with the **CLASS(\*)** specifier is an **unlimited polymorphic** entity. An **unlimited polymorphic**  
17 entity is not declared to have a type. It is not considered to have the same **declared type** as any other entity,  
18 including another **unlimited polymorphic** entity.

19 4 A nonpolymorphic entity is **type compatible** only with entities of the same **declared type**. A **polymorphic** entity  
20 that is not an **unlimited polymorphic** entity is **type compatible** with entities of the same **declared type** or any of  
21 its **extensions**. Even though an **unlimited polymorphic** entity is not considered to have a **declared type**, it is **type**  
22 **compatible** with all entities. An entity is **type compatible** with a type if it is **type compatible** with entities of that  
23 type.

#### NOTE 4.4

Given

```
TYPE TROOT
...
TYPE,EXTENDS(TROOT) :: TEXTENDED
...
CLASS(TROOT) A
CLASS(TEXTENDED) B
...
```

A is **type compatible** with B but B is not **type compatible** with A.

24 5 A **polymorphic allocatable** object may be allocated to be of any type with which it is **type compatible**. A  
25 **polymorphic** pointer or **dummy argument** may, during program execution, be associated with objects with which  
26 it is **type compatible**.

27 6 The **dynamic type** of an allocated **allocatable polymorphic** object is the type with which it was allocated. The  
28 **dynamic type** of an associated **polymorphic** pointer is the **dynamic type** of its **target**. The **dynamic type** of a  
29 nonallocatable nonpointer **polymorphic dummy argument** is the **dynamic type** of its **effective argument**. The  
30 **dynamic type** of an unallocated **allocatable** object or a **disassociated** pointer is the same as its **declared type**. The  
31 **dynamic type** of an entity identified by an **associate name** (8.1.3) is the **dynamic type** of the selector with which  
32 it is associated. The **dynamic type** of an object that is not **polymorphic** is its **declared type**.

## 4.4 Intrinsic types

### 4.4.1 Classification and specification

1 Each intrinsic type is classified as a [numeric type](#) or a nonnumeric type. The [numeric types](#) are integer, real, and complex. The nonnumeric intrinsic types are character and logical.

2 Each intrinsic type has a [kind type parameter](#) named KIND; this [type parameter](#) is of type integer with default kind.

```
R404  intrinsic-type-spec      is  integer-type-spec
                                     or  REAL [ kind-selector ]
                                     or  DOUBLE PRECISION
                                     or  COMPLEX [ kind-selector ]
                                     or  CHARACTER [ char-selector ]
                                     or  LOGICAL [ kind-selector ]
```

```
R405  integer-type-spec      is  INTEGER [ kind-selector ]
```

```
R406  kind-selector          is  ( [ KIND = ] scalar-int-constant-expr )
```

C412 (R406) The value of *scalar-int-constant-expr* shall be nonnegative and shall specify a representation method that exists on the processor.

### 4.4.2 Intrinsic operations on intrinsic types

1 Intrinsic numeric operations are defined as specified in [7.1.5.2.1](#) for the numeric intrinsic types. Relational intrinsic operations are defined as specified in [7.1.5.5](#) for numeric and character intrinsic types. The intrinsic concatenation operation is defined as specified in [7.1.5.3](#) for the character type. Logical intrinsic operations are defined as specified in [7.1.5.4](#) for the logical type.

### 4.4.3 Numeric intrinsic types

#### 4.4.3.1 Integer type

1 The set of values for the integer type is a subset of the mathematical integers. The processor shall provide one or more representation methods that define sets of values for data of type integer. Each such method is characterized by a value for the [kind type parameter](#) KIND. The [kind type parameter](#) of a representation method is returned by the intrinsic function [KIND](#) ([13.9.108](#)). The decimal exponent range of a representation method is returned by the intrinsic function [RANGE](#) ([13.9.158](#)). The intrinsic function [SELECTED\\_INT\\_KIND](#) ([13.9.169](#)) returns a kind value based on a specified decimal exponent range requirement. The integer type includes a zero value, which is considered to be neither negative nor positive. The value of a signed integer zero is the same as the value of an unsigned integer zero.

2 The processor shall provide at least one representation method with a decimal exponent range greater than or equal to 18.

3 The type specifier for the integer type uses the keyword INTEGER.

4 The keyword INTEGER with no [kind-selector](#) specifies type integer with default kind; the [kind type parameter](#) value is equal to [KIND](#) (0). The decimal exponent range of default integer shall be at least 5.

5 Any integer value may be represented as a [signed-int-literal-constant](#).

```
R407  signed-int-literal-constant is  [ sign ] int-literal-constant
```

```
R408  int-literal-constant       is  digit-string [ _ kind-param ]
```

```
R409  kind-param                 is  digit-string
```

or *scalar-int-constant-name*

R410 *signed-digit-string* is [ *sign* ] *digit-string*

R411 *digit-string* is *digit* [ *digit* ] ...

R412 *sign* is +

or -

C413 (R409) A *scalar-int-constant-name* shall be a **named constant** of type integer.

C414 (R409) The value of *kind-param* shall be nonnegative.

C415 (R408) The value of *kind-param* shall specify a representation method that exists on the processor.

6 The optional **kind type parameter** following *digit-string* specifies the **kind type parameter** of the integer constant; if it is does not appear, the constant is default integer.

7 An integer constant is interpreted as a decimal value.

#### NOTE 4.5

Examples of signed integer literal constants are:

```
473
+56
-101
21_2
21_SHORT
1976354279568241_8
```

where SHORT is a scalar integer **named constant**.

#### 4.4.3.2 Real type

1 The real type has values that approximate the mathematical real numbers. The processor shall provide two or more approximation methods that define sets of values for data of type real. Each such method has a representation method and is characterized by a value for the **kind type parameter** KIND. The **kind type parameter** of an approximation method is returned by the intrinsic function **KIND** (13.9.108).

2 The decimal precision, decimal exponent range, and radix of an approximation method are returned by the intrinsic functions **PRECISION** (13.9.151), **RADIX** (13.9.154) and **RANGE** (13.9.158). The intrinsic function **SELECTED\_REAL\_KIND** (13.9.170) returns a kind value based on specified precision, range, and radix requirements.

#### NOTE 4.6

See C.1.1 for remarks concerning selection of approximation methods.

3 The real type includes a zero value. Processors that distinguish between positive and negative zeros shall treat them as mathematically equivalent

- in all intrinsic relational operations, and
- as **actual arguments** to intrinsic procedures other than those for which it is explicitly specified that negative zero is distinguished.

#### NOTE 4.7

On a processor that distinguishes between 0.0 and -0.0,

**NOTE 4.7 (cont.)**

( X >= 0.0 )

evaluates to true if X = 0.0 or if X = −0.0,

( X < 0.0 )

evaluates to false for X = −0.0.

In order to distinguish between 0.0 and −0.0, a program can use the intrinsic function **SIGN**. SIGN(1.0,X) will return −1.0 if X < 0.0 or if the processor distinguishes between 0.0 and −0.0 and X has the value −0.0.

- 1 4 The type specifier for the real type uses the keyword REAL. The keyword DOUBLE PRECISION is an alternative  
2 specifier for one kind of real type.
- 3 5 If the type keyword REAL is used without a **kind type parameter**, the real type with default real kind is specified  
4 and the kind value is **KIND** (0.0). The type specifier DOUBLE PRECISION specifies type real with double  
5 precision kind; the kind value is **KIND** (0.0D0). The decimal precision of the double precision real approximation  
6 method shall be greater than that of the default real method.
- 7 6 The decimal precision of double precision real shall be at least 10, and its decimal exponent range shall be at  
8 least 37. It is recommended that the decimal precision of default real be at least 6, and that its decimal exponent  
9 range be at least 37.
- 10 R413 *signed-real-literal-constant* is [ *sign* ] *real-literal-constant*
- 11 R414 *real-literal-constant* is *significand* [ *exponent-letter exponent* ] [ *\_\_ kind-param* ]  
12 or *digit-string exponent-letter exponent* [ *\_\_ kind-param* ]
- 13 R415 *significand* is *digit-string* . [ *digit-string* ]  
14 or . *digit-string*
- 15 R416 *exponent-letter* is E  
16 or D
- 17 R417 *exponent* is *signed-digit-string*
- 18 C416 (R414) If both *kind-param* and *exponent-letter* appear, *exponent-letter* shall be E.
- 19 C417 (R414) The value of *kind-param* shall specify an approximation method that exists on the processor.
- 20 7 A real literal constant without a **kind type parameter** is a default real constant if it is without an exponent part  
21 or has exponent letter E, and is a double precision real constant if it has exponent letter D. A real literal constant  
22 written with a **kind type parameter** is a real constant with the specified **kind type parameter**.
- 23 8 The exponent represents the power of ten scaling to be applied to the significand or digit string. The meaning of  
24 these constants is as in decimal scientific notation.
- 25 9 The significand may be written with more digits than a processor will use to approximate the value of the constant.  
26

**NOTE 4.8**

Examples of signed real literal constants are:

−12.78

+1.6E3

2.1

−16.E4\_8

**NOTE 4.8 (cont.)**

```

0.45D-4
10.93E7_QUAD
.123
3E4

```

where QUAD is a scalar integer [named constant](#).

**4.4.3.3 Complex type**

1 The complex type has values that approximate the mathematical complex numbers. The values of a complex type are ordered pairs of real values. The first real value is called the real part, and the second real value is called the imaginary part.

2 Each approximation method used to represent data entities of type real shall be available for both the real and imaginary parts of a data entity of type complex. The (default integer) [kind type parameter](#) KIND for a complex entity specifies for both parts the real approximation method characterized by this kind type parameter value. The [kind type parameter](#) of an approximation method is returned by the intrinsic function [KIND](#) (13.9.108).

3 The type specifier for the complex type uses the keyword COMPLEX. There is no keyword for double precision complex. If the type keyword COMPLEX is used without a [kind type parameter](#), the complex type with default complex kind is specified, the kind value is [KIND](#) (0.0), and both parts are default real.

R418 *complex-literal-constant* is ( *real-part* , *imag-part* )

R419 *real-part* is [signed-int-literal-constant](#)  
or [signed-real-literal-constant](#)  
or [named-constant](#)

R420 *imag-part* is [signed-int-literal-constant](#)  
or [signed-real-literal-constant](#)  
or [named-constant](#)

C418 (R418) Each [named constant](#) in a complex literal constant shall be of type integer or real.

4 If the real part and the imaginary part of a complex literal constant are both real, the [kind type parameter](#) value of the complex literal constant is the [kind type parameter](#) value of the part with the greater decimal precision; if the precisions are the same, it is the [kind type parameter](#) value of one of the parts as determined by the processor. If a part has a [kind type parameter](#) value different from that of the complex literal constant, the part is converted to the approximation method of the complex literal constant.

5 If both the real and imaginary parts are integer, they are converted to the default real approximation method and the constant is default complex. If only one of the parts is an integer, it is converted to the approximation method selected for the part that is real and the [kind type parameter](#) value of the complex literal constant is that of the part that is real.

**NOTE 4.9**

Examples of complex literal constants are:

```

(1.0, -1.0)
(3, 3.1E6)
(4.0_4, 3.6E7_8)
( 0., PI)      ! where PI is a previously declared named real constant.

```

## 4.4.4 Character type

### 4.4.4.1 Character sets

- 1 The character type has a set of values composed of character strings. A character string is a sequence of characters, numbered from left to right 1, 2, 3, ... up to the number of characters in the string. The number of characters in the string is called the length of the string. The length is a type parameter; its kind is processor dependent and its value is greater than or equal to zero.
- 2 The processor shall provide one or more representation methods that define sets of values for data of type character. Each such method is characterized by a value for the (default integer) *kind type parameter* *KIND*. The *kind type parameter* of a representation method is returned by the intrinsic function *KIND* (13.9.108). The intrinsic function *SELECTED\_CHAR\_KIND* (13.9.168) returns a kind value based on the name of a character type. Any character of a particular representation method representable in the processor may occur in a character string of that representation method.
- 3 The character set specified in ISO/IEC 646:1991 (International Reference Version) is referred to as the *ASCII character* set and its corresponding representation method is *ASCII character* kind. The character set UCS-4 as specified in ISO/IEC 10646 is referred to as the *ISO 10646 character* set and its corresponding representation method is the *ISO 10646 character* kind.

### 4.4.4.2 Character type specifier

- 1 The type specifier for the character type uses the keyword *CHARACTER*.
- 2 If the type keyword *CHARACTER* is used without a *kind type parameter*, the character type with default character kind is specified and the kind value is *KIND* ('A').
- 3 The default character kind shall support a character set that includes the characters in the Fortran character set (3.1). By supplying nondefault character kinds, the processor may support additional character sets. The characters available in nondefault character kinds are not specified by this part of ISO/IEC 1539, except that one character in each nondefault character set shall be designated as a blank character to be used as a padding character.

R421 *char-selector*                      is *length-selector*  
    or ( *LEN* = *type-param-value* , ■  
    ■ *KIND* = *scalar-int-constant-expr* )  
    or ( *type-param-value* , ■  
    ■ [ *KIND* = ] *scalar-int-constant-expr* )  
    or ( *KIND* = *scalar-int-constant-expr* ■  
    ■ [ , *LEN* = *type-param-value* ] )

R422 *length-selector*                      is ( [ *LEN* = ] *type-param-value* )  
    or \* *char-length* [ , ]

R423 *char-length*                          is ( *type-param-value* )  
    or *int-literal-constant*

C419 (R421) The value of *scalar-int-constant-expr* shall be nonnegative and shall specify a representation method that exists on the processor.

C420 (R423) The *int-literal-constant* shall not include a *kind-param*.

C421 (R423) A *type-param-value* in a *char-length* shall be a colon, asterisk, or *specification-expr*.

C422 (R421 R422 R423) A *type-param-value* of \* shall be used only

- to declare a *dummy argument*,
- to declare a *named constant*,

- in the *type-spec* of an **ALLOCATE** statement wherein each *allocate-object* is a dummy argument of type CHARACTER with an assumed character length,
- in the *type-spec* or *derived-type-spec* of a type guard statement (8.1.11), or
- in an external function, to declare the character length parameter of the function result.

C423 A function name shall not be declared with an asterisk *type-param-value* unless it is of type CHARACTER and is the name of a dummy function or the name of the result of an external function.

C424 A function name declared with an asterisk *type-param-value* shall not be an array, a pointer, elemental, recursive, or pure.

(R422) C425 The optional comma in a *length-selector* is permitted only in a *declaration-type-spec* in a *type-declaration-stmt*.

(R422) C426 The optional comma in a *length-selector* is permitted only if no double-colon separator appears in the *type-declaration-stmt*.

(R421) C427 The length specified for a character statement function or for a statement function dummy argument of type character shall be a constant expression.

4 The *char-selector* in a CHARACTER *intrinsic-type-spec* and the \* *char-length* in an *entity-decl* or in a *component-decl* of a type definition specify character length. The \* *char-length* in an *entity-decl* or a *component-decl* specifies an individual length and overrides the length specified in the *char-selector*, if any. If a \* *char-length* is not specified in an *entity-decl* or a *component-decl*, the *length-selector* or *type-param-value* specified in the *char-selector* is the character length. If the length is not specified in a *char-selector* or a \* *char-length*, the length is 1.

5 If the character length parameter value evaluates to a negative value, the length of character entities declared is zero. A character length parameter value of : indicates a deferred type parameter (4.2). A *char-length* type parameter value of \* has the following meanings.

- If used to declare a dummy argument of a procedure, the dummy argument assumes its length from its effective argument.
- If used to declare a named constant, the length is that of the constant value.
- If used in the *type-spec* of an **ALLOCATE** statement, each *allocate-object* assumes its length from its effective argument.
- If used in the *type-spec* of a type guard statement, the associating entity assumes its length from the selector.
- If used to specify the character length parameter of a function result, any scoping unit invoking the function or passing it as an actual argument shall declare the function name with a character length parameter value other than \* or access such a definition by argument, host, or use association. When the function is invoked, the length of the function result is assumed from the value of this type parameter.

#### 4.4.4.3 Character literal constant

1 The syntax of a character literal constant is given by R424.

R424 *char-literal-constant* is [ *kind-param* \_ ] ' [ *rep-char* ] ... '  
or [ *kind-param* \_ ] " [ *rep-char* ] ... "

(R424) C428 The value of *kind-param* shall specify a representation method that exists on the processor.

2 The optional kind type parameter preceding the leading delimiter specifies the kind type parameter of the character constant; if it does not appear, the constant is default character.

3 For the type character with kind *kind-param*, if it appears, and for default character otherwise, a representable character, *rep-char*, is defined as follows.

- In free source form, it is any graphic character in the processor-dependent character set.
- In fixed source form, it is any character in the processor-dependent character set. A processor may restrict the occurrence of some or all of the control characters.

- 1 4 The delimiting apostrophes or quotation marks are not part of the value of the character literal constant.
- 2 5 An apostrophe character within a character constant delimited by apostrophes is represented by two consecutive  
3 apostrophes (without intervening blanks); in this case, the two apostrophes are counted as one character. Sim-  
4 ilarly, a quotation mark character within a character constant delimited by quotation marks is represented by  
5 two consecutive quotation marks (without intervening blanks) and the two quotation marks are counted as one  
6 character.
- 7 6 A zero-length character literal constant is represented by two consecutive apostrophes (without intervening blanks)  
8 or two consecutive quotation marks (without intervening blanks) outside of a [character context](#).

**NOTE 4.10**

Examples of character literal constants are:

```
"DON'T"
'DON' 'T'
```

both of which have the value DON'T and

```
''
```

which has the zero-length character string as its value.

**NOTE 4.11**

An example of a nondefault character literal constant, where the processor supports the corresponding character set, is:

```
NIHONGO_ '彼女なしでは何もできない。'
```

where NIHONGO is a [named constant](#) whose value is the kind type parameter for Nihongo (Japanese) characters. This means “Without her, nothing is possible”.

9 **4.4.4.4 Collating sequence**

- 10 1 The processor defines a [collating sequence](#) for the character set of each kind of character. The [collating sequence](#)  
11 is an isomorphism between the character set and the set of integers  $\{I : 0 \leq I < N\}$ , where  $N$  is the number of  
12 characters in the set. The intrinsic functions [CHAR](#) (13.9.44) and [ICHAR](#) (13.9.95) provide conversions between  
13 the characters and the integers according to this mapping.

**NOTE 4.12**

For example:

```
ICHAR ( 'X' )
```

returns the integer value of the character 'X' according to the [collating sequence](#) of the processor.

- 14 2 The [collating sequence](#) of the default character kind shall satisfy the following constraints.
- 15 • ICHAR ('A') < ICHAR ('B') < ... < ICHAR ('Z') for the twenty-six upper-case letters.
  - 16 • ICHAR ('0') < ICHAR ('1') < ... < ICHAR ('9') for the ten digits.
  - 17 • ICHAR (' ') < ICHAR ('0') < ICHAR ('9') < ICHAR ('A') or  
18 ICHAR (' ') < ICHAR ('A') < ICHAR ('Z') < ICHAR ('0').
  - 19 • ICHAR ('a') < ICHAR ('b') < ... < ICHAR ('z') for the twenty-six lower-case letters.
  - 20 • ICHAR (' ') < ICHAR ('0') < ICHAR ('9') < ICHAR ('a') or  
21 ICHAR (' ') < ICHAR ('a') < ICHAR ('z') < ICHAR ('0').



- 1 6 The [components](#), [direct components](#), [potential subobject components](#), and [ultimate components](#) of an object of  
 2 derived type are the [components](#), [direct components](#), [potential subobject components](#), and [ultimate components](#)  
 3 of its type, respectively.
- 4 7 By default, no [storage sequence](#) is implied by the order of the component definitions. However, a storage order  
 5 is implied for a [sequence type](#) (4.5.2.3). If the derived type has the [BIND attribute](#), the [storage sequence](#) is that  
 6 required by the [companion processor](#) (2.5.7, 15.3.4).
- 7 8 A scalar entity of derived type is a [structure](#). If a derived type has the [SEQUENCE attribute](#), a scalar entity of  
 8 the type is a [sequence structure](#).

**NOTE 4.14**

The [ultimate components](#) of an object of the derived type `kids` defined below are `name`, `age`, and `other_kids`. The [direct components](#) of such an object are `name`, `age`, `other_kids`, and `oldest_child`.

```

type :: person
  character(len=20) :: name
  integer :: age
end type person

type :: kids
  type(person) :: oldest_child
  type(person), allocatable, dimension(:) :: other_kids
end type kids

```

## 9 4.5.2 Derived-type definition

### 10 4.5.2.1 Syntax

- 11 R426 *derived-type-def* is *derived-type-stmt*  
 12 [ *type-param-def-stmt* ] ...  
 13 [ *private-or-sequence* ] ...  
 14 [ *component-part* ]  
 15 [ *type-bound-procedure-part* ]  
 16 *end-type-stmt*
- 17 R427 *derived-type-stmt* is TYPE [ [ , *type-attr-spec-list* ] :: ] *type-name* ■  
 18 ■ [ ( *type-param-name-list* ) ]
- 19 R428 *type-attr-spec* is ABSTRACT  
 20 or *access-spec*  
 21 or BIND (C)  
 22 or EXTENDS ( *parent-type-name* )
- 23 C430 (R427) A derived type *type-name* shall not be [DOUBLEPRECISION](#) or the same as the name of any  
 24 intrinsic type defined in this part of ISO/IEC 1539.
- 25 C431 (R427) The same *type-attr-spec* shall not appear more than once in a given *derived-type-stmt*.
- 26 C432 (R427) The same *type-param-name* shall not appear more than once in a *derived-type-stmt*.
- 27 C433 (R428) A *parent-type-name* shall be the name of a previously defined [extensible type](#) (4.5.7).
- 28 C434 (R426) If the type definition contains or [inherits](#) (4.5.7.2) a deferred [type-bound procedure](#) (4.5.5), [AB-](#)  
 29 [STRACT](#) shall appear.
- 30 C435 (R426) If [ABSTRACT](#) appears, the type shall be [extensible](#).

- 1 C436 (R426) If EXTENDS appears, SEQUENCE shall not appear.
- 2 C437 (R426) If EXTENDS appears and the type being defined has a [coarray ultimate component](#), its [parent](#)  
3 [type](#) shall have a [coarray ultimate component](#).
- 4 C438 (R426) If EXTENDS appears and the type being defined has a [potential subobject component](#) of type  
5 [EVENT\\_TYPE](#) or [LOCK\\_TYPE](#) from the intrinsic module ISO\_FORTRAN\_ENV, its [parent type](#)  
6 shall be [EVENT\\_TYPE](#) or [LOCK\\_TYPE](#) or have a [potential subobject component](#) of type [EVENT\\_](#)  
7 [TYPE](#) or [LOCK\\_TYPE](#).
- 8 R429 *private-or-sequence* is [private-components-stmt](#)  
9 or [sequence-stmt](#)
- 10 C439 (R426) The same [private-or-sequence](#) shall not appear more than once in a given [derived-type-def](#).
- 11 R430 *end-type-stmt* is END TYPE [ *type-name* ]
- 12 C440 (R430) If END TYPE is followed by a *type-name*, the *type-name* shall be the same as that in the  
13 corresponding [derived-type-stmt](#).
- 14 1 Derived types with the [BIND attribute](#) are subject to additional constraints as specified in [15.3.4](#).

**NOTE 4.15**

An example of a derived-type definition is:

```
TYPE PERSON
  INTEGER AGE
  CHARACTER (LEN = 50) NAME
END TYPE PERSON
```

An example of declaring a variable CHAIRMAN of type PERSON is:

```
TYPE (PERSON) :: CHAIRMAN
```

15 **4.5.2.2 Accessibility**

- 16 1 The accessibility of a type name is determined as specified in [5.5.2](#). The accessibility of a type name does not  
17 affect, and is not affected by, the accessibility of its [components](#) and [type-bound procedures](#).
- 18 2 If a derived type is defined in the scoping unit of a module, and its name is private in that module, then the type  
19 name, and thus the structure constructor ([4.5.10](#)) for the type, are accessible only within that module and its  
20 [descendants](#).

**NOTE 4.16**

An example of a type with a private name is:

```
TYPE, PRIVATE :: AUXILIARY
  LOGICAL :: DIAGNOSTIC
  CHARACTER (LEN = 20) :: MESSAGE
END TYPE AUXILIARY
```

Such a type would be accessible only within the module in which it is defined, and within its [descendants](#).

21 **4.5.2.3 Sequence type**

- 22 R431 *sequence-stmt* is SEQUENCE
- 23 C441 (R426) If SEQUENCE appears, the type shall have at least one component, each data component shall

be declared to be of an intrinsic type or of a **sequence type**, the derived type shall not have any **type parameter**, and a **type-bound-procedure-part** shall not appear.

- 1 If the SEQUENCE statement appears, the type has the **SEQUENCE attribute** and is a **sequence type**. The order of the component definitions in a **sequence type** specifies a **storage sequence** for objects of that type. The type is a **numeric sequence type** if there are no **pointer** or **allocatable components**, and each **component** is default integer, default real, double precision real, default complex, default logical, or of **numeric sequence type**. The type is a **character sequence type** if there are no **pointer** or **allocatable** components, and each component is default character or of **character sequence type**.

#### NOTE 4.17

An example of a **numeric sequence type** is:

```
TYPE NUMERIC_SEQ
  SEQUENCE
  INTEGER :: INT_VAL
  REAL    :: REAL_VAL
  LOGICAL :: LOG_VAL
END TYPE NUMERIC_SEQ
```

#### NOTE 4.18

A structure resolves into a sequence of components. Unless the structure includes a SEQUENCE statement, the use of this terminology in no way implies that these components are stored in this, or any other, order. Nor is there any requirement that **contiguous** storage be used. The sequence merely refers to the fact that in writing the definitions there will necessarily be an order in which the components appear, and this will define a sequence of components. This order is of limited significance because a component of an object of derived type will always be accessed by a component name except in the following contexts: the sequence of expressions in a derived-type value constructor, intrinsic assignment, the data values in namelist input data, and the inclusion of the structure in an input/output list of a formatted data transfer, where it is expanded to this sequence of components. Provided the processor adheres to the defined order in these cases, it is otherwise free to organize the storage of the components for any nonsequence structure in memory as best suited to the particular architecture.

#### 4.5.2.4 Determination of derived types

- 1 Derived-type definitions with the same type name may appear in different **scoping units**, in which case they might be independent and describe different derived types or they might describe the same type.
- 2 Two data entities have the same type if they are declared with reference to the same derived-type definition. Data entities also have the same type if they are declared with reference to different derived-type definitions that specify the same type name, all have the **SEQUENCE attribute** or all have the **BIND attribute**, have no components with **PRIVATE** accessibility, and have components that agree in order, name, and attributes. Otherwise, they are of different derived types. A data entity declared using a type with the **SEQUENCE attribute** or with the **BIND attribute** is not of the same type as an entity of a type that has any components that are **PRIVATE**.

#### NOTE 4.19

An example of declaring two entities with reference to the same derived-type definition is:

```
TYPE POINT
  REAL X, Y
END TYPE POINT
TYPE (POINT) :: X1
CALL SUB (X1)
...
CONTAINS
```

**NOTE 4.19 (cont.)**

```

SUBROUTINE SUB (A)
  TYPE (POINT) :: A
  ...
END SUBROUTINE SUB

```

The definition of derived type POINT is known in subroutine SUB by host association. Because the declarations of X1 and A both reference the same derived-type definition, X1 and A have the same type. X1 and A also would have the same type if the derived-type definition were in a module and both SUB and its containing [program unit](#) referenced the module.

**NOTE 4.20**

An example of data entities in different [scoping units](#) having the same type is:

```

PROGRAM PGM
  TYPE EMPLOYEE
    SEQUENCE
    INTEGER          ID_NUMBER
    CHARACTER (50) NAME
  END TYPE EMPLOYEE
  TYPE (EMPLOYEE) PROGRAMMER
  CALL SUB (PROGRAMMER)
  ...
END PROGRAM PGM
SUBROUTINE SUB (POSITION)
  TYPE EMPLOYEE
    SEQUENCE
    INTEGER          ID_NUMBER
    CHARACTER (50) NAME
  END TYPE EMPLOYEE
  TYPE (EMPLOYEE) POSITION
  ...
END SUBROUTINE SUB

```

The [actual argument](#) PROGRAMMER and the [dummy argument](#) POSITION have the same type because they are declared with reference to a derived-type definition with the same name, the [SEQUENCE attribute](#), and components that agree in order, name, and [attributes](#).

Suppose the component name ID\_NUMBER was ID\_NUM in the subroutine. Because all the component names are not identical to the component names in derived type EMPLOYEE in the main program, the [actual argument](#) PROGRAMMER would not be of the same type as the [dummy argument](#) POSITION. Thus, the program would not be standard-conforming.

**NOTE 4.21**

The requirement that the two types have the same name applies to the *type-names* in the respective derived type definitions, not to local names introduced via renaming in [USE statements](#).

1 **4.5.3 Derived-type parameters**2 **4.5.3.1 Type parameter definition statement**

3 R432 *type-param-def-stmt* is *integer-type-spec, type-param-attr-spec* :: ■  
 4 ■ *type-param-decl-list*

5 R433 *type-param-decl* is *type-param-name* [ = *scalar-int-constant-expr* ]

- C442     A *type-param-name* in a *derived-type-def-stmt* in a *derived-type-def* shall be one of the *type-param-names* in the *derived-type-stmt* of that *derived-type-def*.

C443     (R432) Each *type-param-name* in the *derived-type-stmt* in a *derived-type-def* shall appear exactly once as a *type-param-name* in a *type-param-def-stmt* in that *derived-type-def*.

R434     *type-param-attr-spec*                **is**   KIND  
   **or**   LEN

  - 1 The derived type is parameterized if the *derived-type-stmt* has any *type-param-names*.
  - 2 Each type parameter is itself of type integer. If its kind selector is omitted, the **kind type parameter** is default integer.
  - 3 The *type-param-attr-spec* explicitly specifies whether a type parameter is a kind parameter or a length parameter.
  - 4 If a *type-param-decl* has a *scalar-int-constant-expr*, the type parameter has a default value which is specified by the expression. If necessary, the value is converted according to the rules of intrinsic assignment (7.2.1.3) to a value of the same kind as the type parameter.
  - 5 A **type parameter** may be used as a primary in a **specification expression** (7.1.11) in the *derived-type-def*. A **kind type parameter** may also be used as a primary in a **constant expression** (7.1.12) in the *derived-type-def*.

### NOTE 4.22

The following example uses derived-type parameters.

```

TYPE humongous_matrix(k, d)
  INTEGER, KIND :: k = kind(0.0)
  INTEGER(selected_int_kind(12)), LEN :: d
  !-- Specify a nondefault kind for d.
  REAL(k) :: element(d,d)
END TYPE

```

In the following example, `dim` is declared to be a kind parameter, allowing generic overloading of procedures distinguished only by `dim`.

```
TYPE general_point(dim)
  INTEGER, KIND :: dim
  REAL :: coordinates(dim)
END TYPE
```

#### 4.5.3.2 Type parameter order

- 17    1    **Type parameter order** is an ordering of the type parameters of a derived type; it is used for derived-type specifiers.
- 18    2    The **type parameter order** of a nonextended type is the order of the type parameter list in the derived-type
- 19    definition. The **type parameter order** of an **extended type** (4.5.7) consists of the **type parameter order** of its
- 20    **parent type** followed by any additional type parameters in the order of the type parameter list in the derived-type
- 21    definition.

### NOTE 4.23

Given

```
TYPE :: t1(k1,k2)
    INTEGER,KIND :: k1,k2
    REAL(k1) a(k2)
END TYPE
TYPE,EXTENDS(t1) :: t2(k3)
```

## NOTE 4.23 (cont.)

```

    INTEGER,KIND :: k3
    LOGICAL(k3) flag
END TYPE

```

the type parameter order for type `t1` is `k1` then `k2`, and the type parameter order for type `t2` is `k1` then `k2` then `k3`.

## 4.5.4 Components

## 4.5.4.1 Component definition statement

R435 *component-part* is [ *component-def-stmt* ] ...

R436 *component-def-stmt* is *data-component-def-stmt*  
or *proc-component-def-stmt*

R437 *data-component-def-stmt* is *declaration-type-spec* [ [ , *component-attr-spec-list* ] :: ] ■  
■ *component-decl-list*

R438 *component-attr-spec* is *access-spec*  
or ALLOCATABLE  
or CODIMENSION *lbracket coarray-spec rbracket*  
or CONTIGUOUS  
or DIMENSION ( *component-array-spec* )  
or POINTER

R439 *component-decl* is *component-name* [ ( *component-array-spec* ) ] ■  
■ [ *lbracket coarray-spec rbracket* ] ■  
■ [ \* *char-length* ] [ *component-initialization* ]

R440 *component-array-spec* is *explicit-shape-spec-list*  
or *deferred-shape-spec-list*

C444 (R437) No *component-attr-spec* shall appear more than once in a given *component-def-stmt*.

C445 (R437) If neither the **POINTER** nor the **ALLOCATABLE** attribute is specified, the *declaration-type-spec* in the *component-def-stmt* shall specify an intrinsic type or a previously defined derived type.

C446 (R437) If the **POINTER** or **ALLOCATABLE** attribute is specified, each *component-array-spec* shall be a *deferred-shape-spec-list*.

C447 (R437) If a *coarray-spec* appears, it shall be a *deferred-coshape-spec-list* and the component shall have the **ALLOCATABLE** attribute.

C448 (R437) If a *coarray-spec* appears, the component shall not be of type **C\_PTR** or **C\_FUNPTR** from the intrinsic module **ISO\_C\_BINDING** (15.3.3).

C449 A data component whose type has a *coarray ultimate component* shall be a nonpointer nonallocatable scalar and shall not be a *coarray*.

C450 (R437) If neither the **POINTER** nor the **ALLOCATABLE** attribute is specified, each *component-array-spec* shall be an *explicit-shape-spec-list*.

C451 (R440) Each *bound* in the *explicit-shape-spec* shall be a *specification expression* in which there are no references to specification functions or the intrinsic functions **ALLOCATED**, **ASSOCIATED**, **EXTENDS\_TYPE\_OF**, **PRESENT**, or **SAME\_TYPE\_AS**, every specification inquiry reference is a *constant ex-*

1 **pression**, and the value does not depend on the value of a variable.

2 C452 (R437) A component shall not have both the **ALLOCATABLE** and **POINTER** attributes.

3 C453 (R437) If the **CONTIGUOUS attribute** is specified, the component shall be an array with the **POINTER**  
4 **attribute**.

5 C454 (R439) The \* *char-length* option is permitted only if the component is of type character.

C455 (R436) Each *type-param-value* within a *component-def-stmt* shall be a colon or a **specification expression** in which there are no references to specification functions or the intrinsic functions **ALLOCATED**, **ASSOCIATED**, **EXTENDS\_TYPE\_OF**, **PRESENT**, or **SAME\_TYPE\_AS**, every specification inquiry reference is a **constant expression**, and the value does not depend on the value of a variable.

**NOTE 4.24**

Because a type parameter is not an object, a *type-param-value* or a *bound* in an *explicit-shape-spec* can contain a *type-param-name*.

10        R441    *proc-component-def-stmt*    is   PROCEDURE ( [ *proc-interface* ] ) , ■  
11        ■ *proc-component-attr-spec-list* :: *proc-decl-list*

### NOTE 4.25

See 12.4.3.7 for definitions of *proc-interface* and *proc-decl*.

```

12      R442  proc-component-attr-spec  is  POINTER
13                                     or  PASS [ (arg-name) ]
14                                     or  NOPASS
15                                     or  access-spec

```

16 C456 (R441) The same *proc-component-attr-spec* shall not appear more than once in a given *proc-component-*  
17 *def-stmt*.

18 C457 (R441) POINTER shall appear in each *proc-component-attr-spec-list*.

19 C458 (R441) If the procedure pointer component has an **implicit interface** or has no arguments, **NOPASS** shall  
20 be specified.

21 C459 (R441) If **PASS** (*arg-name*) appears, the **interface** of the procedure pointer component shall have a **dummy**  
22 **argument** named *arg-name*.

23 C460 (R441) **PASS** and **NOPASS** shall not both appear in the same *proc-component-attr-spec-list*.

1 The *declaration-type-spec* in the *data-component-def-stmt* specifies the type and type parameters of the components in the *component-decl-list*, except that the character length parameter may be specified or overridden for a component by the appearance of \* *char-length* in its *entity-decl*. The *component-attr-spec-list* in the *data-component-def-stmt* specifies the attributes whose keywords appear for the components in the *component-decl-list*, except that the **DIMENSION** attribute may be specified or overridden for a component by the appearance of a *component-array-spec* in its *component-decl*, and the **CODIMENSION** attribute may be specified or overridden for a component by the appearance of a *coarray-spec* in its *component-decl*.

#### 31 4.5.4.2 Array components

1 A data component is an array if its *component-decl* contains a *component-array-spec* or its *data-component-def-stmt* contains a DIMENSION clause. If the *component-decl* contains a *component-array-spec*, it specifies the array *rank*, and if the array is explicit shape (5.5.8.2), the *array bounds*; otherwise, the *component-array-spec* in the DIMENSION clause specifies the array *rank*, and if the array is explicit shape, the *array bounds*.

**NOTE 4.26**

An example of a derived type definition with an array component is:

```

TYPE LINE
  REAL, DIMENSION (2, 2) :: COORD      !
                                     ! COORD(:,1) has the value of [X1, Y1]
                                     ! COORD(:,2) has the value of [X2, Y2]
  REAL                      :: WIDTH    ! Line width in centimeters
  INTEGER                   :: PATTERN  ! 1 for solid, 2 for dash, 3 for dot
END TYPE LINE

```

An example of declaring a variable LINE\_SEGMENT to be of the type LINE is:

```

TYPE (LINE)          :: LINE_SEGMENT

```

The scalar variable LINE\_SEGMENT has a component that is an array. In this case, the array is a subobject of a scalar. The double colon in the definition for COORD is required; the double colon in the definition for WIDTH and PATTERN is optional.

**NOTE 4.27**

An example of a derived type definition with an [allocatable](#) component is:

```

TYPE STACK
  INTEGER          :: INDEX
  INTEGER, ALLOCATABLE :: CONTENTS (:)
END TYPE STACK

```

For each scalar variable of type STACK, the shape of the component CONTENTS is determined by execution of an [ALLOCATE statement](#) or [assignment statement](#), or by [argument association](#).

**NOTE 4.28**

[Default initialization](#) of an [explicit-shape array](#) component can be specified by a [constant expression](#) consisting of an array constructor (4.8), or of a single scalar that becomes the value of each array element.

#### 1 4.5.4.3 Coarray components

- 2 1 A data component is a [coarray](#) if its [component-decl](#) contains a [coarray-spec](#) or its [data-component-def-stmt](#)  
 3 contains a CODIMENSION clause. If the [component-decl](#) contains a [coarray-spec](#) it specifies the [corank](#); otherwise,  
 4 the [coarray-spec](#) in the CODIMENSION clause specifies the [corank](#).

**NOTE 4.29**

An example of a derived type definition with a coarray component is:

```

TYPE GRID_TYPE
  REAL, ALLOCATABLE, CODIMENSION[:, :, :] :: GRID(:, :, :)
END TYPE GRID_TYPE

```

An object of type grid\_type cannot be an array, an [allocatable](#) object, a [coarray](#), or a [pointer](#).

#### 5 4.5.4.4 Pointer components

- 6 1 A component is a pointer (2.4.9) if its [component-attr-spec-list](#) contains the [POINTER attribute](#). A pointer  
 7 component may be a data pointer or a procedure pointer.

**NOTE 4.30**

An example of a derived type definition with a pointer component is:

**TYPE REFERENCE**

```

    INTEGER                                :: VOLUME, YEAR, PAGE
    CHARACTER (LEN = 50)                   :: TITLE
    PROCEDURE (printer_interface), POINTER :: PRINT => NULL()
    CHARACTER, DIMENSION (:), POINTER     :: SYNOPSIS
END TYPE REFERENCE

```

Any object of type REFERENCE will have the four nonpointer components VOLUME, YEAR, PAGE, and TITLE, the procedure pointer PRINT, which has an [explicit interface](#) the same as printer\_interface, plus a pointer to an array of characters holding SYNOPSIS. The size of this [target](#) array will be determined by the length of the synopsis. The space for the [target](#) could be allocated ([6.7.1](#)) or the pointer component could be associated with a target by a [pointer assignment statement](#) ([7.2.2](#)).

**4.5.4.5 The passed-object dummy argument**

1 A [passed-object dummy argument](#) is a distinguished [dummy argument](#) of a procedure pointer component or [type-bound procedure](#). It affects procedure overriding ([4.5.7.3](#)) and [argument association](#) ([12.5.2.2](#)).

2 If [NOPASS](#) is specified, the procedure pointer component or [type-bound procedure](#) has no [passed-object dummy argument](#).

3 If neither [PASS](#) nor [NOPASS](#) is specified or [PASS](#) is specified without *arg-name*, the first [dummy argument](#) of a procedure pointer component or [type-bound procedure](#) is its [passed-object dummy argument](#).

4 If [PASS](#) (*arg-name*) is specified, the [dummy argument](#) named *arg-name* is the [passed-object dummy argument](#) of the procedure pointer component or named [type-bound procedure](#).

C461 The [passed-object dummy argument](#) shall be a scalar, nonpointer, nonallocatable [dummy data object](#) with the same [declared type](#) as the type being defined; all of its [length type parameters](#) shall be assumed; it shall be [polymorphic](#) ([4.3.2.3](#)) if and only if the type being defined is [extensible](#) ([4.5.7](#)). It shall not have the [VALUE](#) attribute.

**NOTE 4.31**

If a procedure is bound to several types as a [type-bound procedure](#), different [dummy arguments](#) might be the [passed-object dummy argument](#) in different contexts.

**4.5.4.6 Default initialization for components**

1 [Default initialization](#) provides a means of automatically initializing pointer components to be [disassociated](#) or associated with specific [targets](#), and nonpointer nonallocatable components to have a particular value. [Allocatable](#) components are always initialized to unallocated.

2 A pointer variable or component is data-pointer-initialization compatible with a [target](#) if the pointer is [type compatible](#) with the [target](#), they have the same [rank](#), all nondeferred [type parameters](#) of the pointer have the same values as the corresponding [type parameters](#) of the [target](#), and the [target](#) is [contiguous](#) if the pointer has the [CONTIGUOUS](#) attribute.

```

R443  component-initialization    is = constant-expr
                                     or => null-init
                                     or => initial-data-target

```

```

R444  initial-data-target        is designator

```

C462 (R437) If [component-initialization](#) appears, a double-colon separator shall appear before the [component-](#)

- 1 *decl-list*.
- 2 C463 (R437) If *component-initialization* appears, every type parameter and *array bound* of the component  
3 shall be a colon or *constant expression*.
- 4 C464 (R437) If => appears in *component-initialization*, POINTER shall appear in the *component-attr-spec-*  
5 *list*. If = appears in *component-initialization*, neither POINTER nor ALLOCATABLE shall appear in  
6 the *component-attr-spec-list*.
- 7 C465 (R443) If *initial-data-target* appears, *component-name* shall be data-pointer-initialization compatible with  
8 it.
- 9 C466 (R444) The *designator* shall designate a nonallocatable variable that has the TARGET and SAVE attrib-  
10 utes and does not have a *vector subscript*. Every subscript, section subscript, substring starting point,  
11 and substring ending point in *designator* shall be a *constant expression*.
- 12 3 If *null-init* appears for a pointer component, that component in any object of the type has an initial association  
13 status of *disassociated* (1.3.62) or becomes disassociated as specified in 16.5.2.4.
- 14 4 If *initial-data-target* appears for a data pointer component, that component in any object of the type is initially  
15 associated with the *target* or becomes associated with the *target* as specified in 16.5.2.3.
- 16 5 If *initial-proc-target* (12.4.3.7) appears in *proc-decl* for a procedure pointer component, that component in any  
17 object of the type is initially associated with the *target* or becomes associated with the *target* as specified in  
18 16.5.2.3.
- 19 6 If *constant-expr* appears for a nonpointer component, that component in any object of the type is initially defined  
20 (16.6.3) or becomes defined as specified in 16.6.5 with the value determined from *constant-expr*. If necessary,  
21 the value is converted according to the rules of intrinsic assignment (7.2.1.3) to a value that agrees in type, type  
22 parameters, and shape with the component. If the component is of a type for which *default initialization* is  
23 specified for a component, the *default initialization* specified by *constant-expr* overrides the *default initialization*  
24 specified for that component. When one *initialization* overrides another it is as if only the overriding *initialization*  
25 were specified (see Note 4.33). *Explicit initialization* in a *type declaration statement* (5.2) overrides *default*  
26 *initialization* (see Note 4.32). Unlike *explicit initialization*, *default initialization* does not imply that the object  
27 has the *SAVE attribute*.
- 28 7 A *subcomponent* (6.4.2) is *default-initialized* if the type of the object of which it is a component specifies *default*  
29 *initialization* for that component, and the *subcomponent* is not a subobject of an object that is *default-initialized*  
30 or *explicitly initialized*.
- 31 8 A type has *default initialization* if *component-initialization* is specified for any *direct component* of the type. An  
32 object has *default initialization* if it is of a type that has *default initialization*.

**NOTE 4.32**

It is not required that *initialization* be specified for each component of a derived type. For example:

```

TYPE DATE
  INTEGER DAY
  CHARACTER (LEN = 5) MONTH
  INTEGER :: YEAR = 2008          ! Partial default initialization
END TYPE DATE

```

In the following example, the default initial value for the YEAR component of TODAY is overridden by *explicit initialization* in the *type declaration statement*:

```

TYPE (DATE), PARAMETER :: TODAY = DATE (21, "Feb.", 2009)

```

**NOTE 4.33**

The default initial value of a component of derived type can be overridden by [default initialization](#) specified in the definition of the type. Continuing the example of Note 4.32:

```
TYPE SINGLE_SCORE
  TYPE (DATE) :: PLAY_DAY = TODAY
  INTEGER SCORE
  TYPE (SINGLE_SCORE), POINTER :: NEXT => NULL ( )
END TYPE SINGLE_SCORE
TYPE (SINGLE_SCORE) SETUP
```

The PLAY\_DAY component of SETUP receives its initial value from TODAY, overriding the [initialization](#) for the YEAR component.

**NOTE 4.34**

Arrays of structures can be declared with elements that are partially or totally initialized by default. Continuing the example of Note 4.33 :

```
TYPE MEMBER (NAME_LEN)
  INTEGER, LEN :: NAME_LEN
  CHARACTER (LEN = NAME_LEN) :: NAME = ''
  INTEGER :: TEAM_NO, HANDICAP = 0
  TYPE (SINGLE_SCORE), POINTER :: HISTORY => NULL ( )
END TYPE MEMBER
TYPE (MEMBER(9)) LEAGUE (36)      ! Array of partially initialized elements
TYPE (MEMBER(9)) :: ORGANIZER = MEMBER (9) ("I. Manage",1,5,NULL ( ))
```

ORGANIZER is [explicitly initialized](#), overriding the [default initialization](#) for an object of type MEMBER.

Allocated objects can also be initialized partially or totally. For example:

```
ALLOCATE (ORGANIZER % HISTORY)    ! A partially initialized object of type
                                ! SINGLE_SCORE is created.
```

**NOTE 4.35**

A pointer component of a derived type can have as its [target](#) an object of that derived type. The type definition can specify that in objects declared to be of this type, such a pointer is default initialized to [disassociated](#). For example:

```
TYPE NODE
  INTEGER :: VALUE = 0
  TYPE (NODE), POINTER :: NEXT_NODE => NULL ( )
END TYPE
```

A type such as this can be used to construct linked lists of objects of type NODE. Linked lists can also be constructed using [allocatable](#) components.

**NOTE 4.36**

A pointer component of a derived type can be default initialized to have an initial [target](#).

```
TYPE NODE
  INTEGER :: VALUE = 0
  TYPE (NODE), POINTER :: NEXT_NODE => SENTINEL
END TYPE
```

**NOTE 4.36 (cont.)**

```
TYPE(NODE), SAVE, TARGET :: SENTINEL
```

**4.5.4.7 Component order**

- 1 **Component order** is an ordering of the nonparent components of a derived type; it is used for intrinsic formatted input/output and **structure constructors** where **component keywords** are not used. **Parent components** are excluded from the **component order** of an **extended type** (4.5.7).
- 2 The **component order** of a nonextended type is the order of the declarations of the components in the derived-type definition. The component order of an **extended type** consists of the **component order** of its **parent type** followed by any additional components in the order of their declarations in the extended derived-type definition.

**NOTE 4.37**

Given the same type definitions as in Note 4.23, the **component order** of type T1 is just A (there is only one component), and the **component order** of type T2 is A then FLAG. The **parent component** (T1) does not participate in the **component order**.

**4.5.4.8 Component accessibility**

R445 *private-components-stmt* is PRIVATE

C467 (R445) A *private-components-stmt* is permitted only if the type definition is within the specification part of a module.

- 1 The default accessibility for the components that are declared in a type's *component-part* is private if the type definition contains a *private-components-stmt*, and public otherwise. The accessibility of a component may be explicitly declared by an *access-spec*; otherwise its accessibility is the default for the type definition in which it is declared.
- 2 If a component is private, that component name is accessible only within the module containing the definition, and within its **descendants**.

**NOTE 4.38**

Type parameters are not components. They are effectively always public.

**NOTE 4.39**

The accessibility of the components of a type is independent of the accessibility of the type name. It is possible to have all four combinations: a public type name with a public component, a private type name with a private component, a public type name with a private component, and a private type name with a public component.

**NOTE 4.40**

An example of a public type with private components is:

```
TYPE, PUBLIC :: POINT
  PRIVATE
  REAL :: X, Y
END TYPE POINT
```

Such a type definition can be accessed by use association; however, the components X and Y are accessible only within the module and its **descendants**.

**NOTE 4.41**

The following example illustrates the use of an individual component *access-spec* to override the default accessibility:

```

TYPE MIXED
  PRIVATE
    INTEGER :: I
    INTEGER, PUBLIC :: J
END TYPE MIXED

TYPE (MIXED) :: M

```

The component M%J is accessible in any *scoping unit* where M is accessible; M%I is accessible only within the module containing the TYPE MIXED definition, and within its *descendants*.

**4.5.5 Type-bound procedures**

- R446 *type-bound-procedure-part* is *contains-stmt*  
 [ *binding-private-stmt* ]  
 [ *type-bound-proc-binding* ] ...
- R447 *binding-private-stmt* is PRIVATE
- C468 (R446) A *binding-private-stmt* is permitted only if the type definition is within the specification part of a module.
- R448 *type-bound-proc-binding* is *type-bound-procedure-stmt*  
 or *type-bound-generic-stmt*  
 or *final-procedure-stmt*
- R449 *type-bound-procedure-stmt* is PROCEDURE [ [ , *binding-attr-list* ] :: ] *type-bound-proc-decl-list*  
 or PROCEDURE ( *interface-name* ), *binding-attr-list* :: *binding-name-list*
- R450 *type-bound-proc-decl* is *binding-name* [ => *procedure-name* ]
- C469 (R449) If => *procedure-name* appears in a *type-bound-proc-decl*, the double-colon separator shall appear.
- C470 (R449) The *procedure-name* shall be the name of an accessible module procedure or an *external procedure* that has an *explicit interface*.
- 1 If neither => *procedure-name* nor *interface-name* appears in a *type-bound-proc-decl*, it is as though => *procedure-name* had appeared with a procedure name the same as the *binding name*.
- R451 *type-bound-generic-stmt* is GENERIC [ , *access-spec* ] :: *generic-spec* => *binding-name-list*
- C471 (R451) Within the *specification-part* of a module, each *type-bound-generic-stmt* shall specify, either implicitly or explicitly, the same accessibility as every other *type-bound-generic-stmt* with that *generic-spec* in the same derived type.
- C472 (R451) Each *binding-name* in *binding-name-list* shall be the name of a specific *binding* of the type.
- C473 (R451) If *generic-spec* is not *generic-name*, each of its specific *bindings* shall have a *passed-object dummy argument* (4.5.4.5).
- C474 (R451) If *generic-spec* is OPERATOR ( *defined-operator* ), the *interface* of each *binding* shall be as specified in 12.4.3.5.2.
- C475 (R451) If *generic-spec* is ASSIGNMENT ( = ), the *interface* of each *binding* shall be as specified in

## 12.4.3.5.3.

(R451) If *generic-spec* is *defined-io-generic-spec*, the *interface* of each *binding* shall be as specified in 9.6.4.8. The type of the *dtv* argument shall be *type-name*.

*binding-attr*                      **is** PASS [ (*arg-name*) ]  
    **or** NOPASS  
    **or** NON\_OVERRIDABLE  
    **or** DEFERRED  
    **or** *access-spec*

(R452) The same *binding-attr* shall not appear more than once in a given *binding-attr-list*.

(R449) If the *interface* of the *binding* has no *dummy argument* of the type being defined, **NOPASS** shall appear.

(R449) If **PASS** (*arg-name*) appears, the *interface* of the *binding* shall have a *dummy argument* named *arg-name*.

(R452) **PASS** and **NOPASS** shall not both appear in the same *binding-attr-list*.

(R452) **NON\_OVERRIDABLE** and **DEFERRED** shall not both appear in the same *binding-attr-list*.

(R452) **DEFERRED** shall appear if and only if *interface-name* appears.

(R449) An overriding *binding* (4.5.7.3) shall have the **DEFERRED** attribute only if the *binding* it overrides is deferred.

(R449) A *binding* shall not override an *inherited binding* (4.5.7.2) that has the **NON\_OVERRIDABLE** attribute.

2 A type-bound procedure statement declares one or more specific *type-bound procedures*. A specific *type-bound procedure* can have a *passed-object dummy argument* (4.5.4.5). A *type-bound procedure* with the **DEFERRED** attribute is a deferred type-bound procedure. The **DEFERRED** keyword shall appear only in the definition of an *abstract type*.

3 A **GENERIC** statement declares a generic *type-bound procedure*, which is a type-bound *generic interface* for its specific *type-bound procedures*.

4 A *binding* of a type is a *type-bound procedure* (specific or generic), a *generic type-bound interface*, or a *final subroutine*. These are referred to as specific bindings, generic bindings, and final bindings respectively.

5 A *type-bound procedure* may be identified by a *binding name* in the scope of the type definition. This name is the *binding-name* for a specific *type-bound procedure*, and the *generic-name* for a generic binding whose *generic-spec* is *generic-name*. A final binding, or a generic binding whose *generic-spec* is not *generic-name*, has no *binding name*.

6 The *interface* of a specific *type-bound procedure* is that of the procedure specified by *procedure-name* or the *interface* specified by *interface-name*.

**NOTE 4.42**

An example of a type and a *type-bound procedure* is:

```
TYPE POINT
  REAL :: X, Y
CONTAINS
  PROCEDURE, PASS :: LENGTH => POINT_LENGTH
END TYPE POINT
...
```

**NOTE 4.42 (cont.)**

and in the *module-subprogram-part* of the same module:

```
REAL FUNCTION POINT_LENGTH (A, B)
  CLASS (POINT), INTENT (IN) :: A, B
  POINT_LENGTH = SQRT ( (A%X - B%X)**2 + (A%Y - B%Y)**2 )
END FUNCTION POINT_LENGTH
```

- 1 7 The same *generic-spec* may be used in several GENERIC statements within a single derived-type definition. Each  
2 additional GENERIC statement with the same *generic-spec* extends the *generic interface*.

**NOTE 4.43**

Unlike the situation with generic procedure names, a generic *type-bound procedure* name is not permitted to be the same as a specific *type-bound procedure* name in the same type (16.3).

- 3 8 The default accessibility for the *type-bound procedures* of a type is private if the type definition contains a *binding-private-stmt*, and public otherwise. The accessibility of a *type-bound procedure* may be explicitly declared by an  
4 *access-spec*; otherwise its accessibility is the default for the type definition in which it is declared.  
5  
6 9 A public *type-bound procedure* is accessible via any accessible object of the type. A private *type-bound procedure*  
7 is accessible only within the module containing the type definition, and within its *descendants*.

**NOTE 4.44**

The accessibility of a *type-bound procedure* is not affected by a *PRIVATE* statement in the *component-part*; the accessibility of a data component is not affected by a *PRIVATE* statement in the *type-bound-procedure-part*.

## 8 4.5.6 Final subroutines

### 9 4.5.6.1 FINAL statement

10 R453 *final-procedure-stmt* is FINAL [ :: ] *final-subroutine-name-list*

11 C485 (R453) A *final-subroutine-name* shall be the name of a module procedure with exactly one *dummy*  
12 *argument*. That argument shall be nonoptional and shall be a noncoarray, nonpointer, nonallocatable,  
13 nonpolymorphic variable of the derived type being defined. All length type parameters of the *dummy*  
14 *argument* shall be assumed. The *dummy argument* shall not have the *INTENT (OUT)* or *VALUE*  
15 attribute.

16 C486 (R453) A *final-subroutine-name* shall not be one previously specified as a *final subroutine* for that type.

17 C487 (R453) A *final subroutine* shall not have a *dummy argument* with the same *kind type parameters* and  
18 *rank* as the *dummy argument* of another *final subroutine* of that type.

19 1 The FINAL statement specifies that each procedure it names is a *final subroutine*. A *final subroutine* might be  
20 executed when a data entity of that type is finalized (4.5.6.2).

21 2 A derived type is *finalizable* if and only if it has a *final subroutine* or a nonpointer, nonallocatable component of  
22 *finalizable* type. A nonpointer data entity is *finalizable* if and only if it is of *finalizable* type. No other entity is  
23 *finalizable*.

**NOTE 4.45**

*Final subroutines* are effectively always “accessible”. They are called for entity *finalization* regardless of the accessibility of the type, its other *type-bound procedures*, or the subroutine name itself.

**NOTE 4.46**

Final subroutines are not inherited through type extension and cannot be overridden. The final subroutines of the parent type are called after any additional final subroutines of an extended type are called.

**4.5.6.2 The finalization process**

- 1 Only finalizable entities are finalized. When an entity is finalized, the following steps are carried out in sequence.
  - (1) If the dynamic type of the entity has a final subroutine whose dummy argument has the same kind type parameters and rank as the entity being finalized, it is called with the entity as an actual argument. Otherwise, if there is an elemental final subroutine whose dummy argument has the same kind type parameters as the entity being finalized, it is called with the entity as an actual argument. Otherwise, no subroutine is called at this point.
  - (2) All finalizable components that appear in the type definition are finalized in a processor-dependent order. If the entity being finalized is an array, each finalizable component of each element of that entity is finalized separately.
  - (3) If the entity is of extended type and the parent type is finalizable, the parent component is finalized.
- 2 If several entities are to be finalized as a consequence of an event specified in 4.5.6.3, the order in which they are finalized is processor dependent. During this process, execution of a final subroutine for one of these entities shall not reference or define any of the other entities that have already been finalized.
- 3 If an object is not finalized, it retains its definition status and does not become undefined.

**NOTE 4.47**

An implementation might need to ensure that when an event causes more than one coarray to be deallocated, they are deallocated in the same order on all images in the current team.

**4.5.6.3 When finalization occurs**

- 1 When an intrinsic assignment statement is executed (7.2.1.3), if the variable is not an unallocated allocatable variable, it is finalized after evaluation of *expr* and before the definition of the variable. If the variable is an allocated allocatable variable that would be deallocated by intrinsic assignment, the finalization occurs before the deallocation.
- 2 When a pointer is deallocated its target is finalized. When an allocatable entity is deallocated, it is finalized unless it is the variable in an intrinsic assignment statement or a component thereof. If an error condition occurs during deallocation, it is processor dependent whether finalization occurs.
- 3 A nonpointer, nonallocatable object that is not a dummy argument or function result is finalized immediately before it would become undefined due to execution of a RETURN or END statement (16.6.6, item (3)).
- 4 A nonpointer nonallocatable local variable of a BLOCK construct is finalized immediately before it would become undefined due to termination of the BLOCK construct (16.6.6, item (23)).
- 5 If an executable construct references a nonpointer function, the result is finalized after execution of the innermost executable construct containing the reference.
- 6 If a specification expression in a scoping unit references a function, the result is finalized before execution of the executable constructs in the scoping unit.
- 7 When a procedure is invoked, a nonpointer, nonallocatable, INTENT (OUT) dummy argument of that procedure is finalized before it becomes undefined. The finalization caused by INTENT (OUT) is considered to occur within the invoked procedure; so for elemental procedures, an INTENT (OUT) argument will be finalized only if a scalar or elemental final subroutine is available, regardless of the rank of the actual argument.
- 8 If an object is allocated via pointer allocation and later becomes unreachable due to all pointers associated with

1 that object having their [pointer association](#) status changed, it is processor dependent whether it is finalized. If it  
 2 is finalized, it is processor dependent as to when the [final subroutines](#) are called.

#### NOTE 4.48

If [finalization](#) is used for storage management, it often needs to be combined with defined assignment.

### 3 4.5.6.4 Entities that are not finalized

4 1 If image execution is [terminated](#), either by an error (e.g. an allocation failure) or by execution of a [stop-stmt](#),  
 5 [error-stop-stmt](#), or [end-program-stmt](#), entities existing immediately prior to [termination](#) are not finalized.

#### NOTE 4.49

A nonpointer, nonallocatable object that has the [SAVE attribute](#) is never finalized as a direct consequence of the execution of a [RETURN](#) or [END statement](#).

## 6 4.5.7 Type extension

### 7 4.5.7.1 Extensible, extended, and abstract types

8 1 A derived type, other than the type [C\\_PTR](#) or [C\\_FUNPTR](#) from the intrinsic module [ISO\\_C\\_BINDING](#), that  
 9 does not have the [BIND attribute](#) or the [SEQUENCE attribute](#) is an [extensible type](#).

10 2 A type with the [EXTENDS attribute](#) is an [extended type](#); its [parent type](#) is the type named in the [EXTENDS](#)  
 11 [type-attr-spec](#).

#### NOTE 4.50

The name of the [parent type](#) might be a local name introduced via renaming in a [USE statement](#).

12 3 An [extensible type](#) that does not have the [EXTENDS attribute](#) is an [extension type](#) of itself only. An [extended](#)  
 13 [type](#) is an [extension](#) of itself and of all types for which its [parent type](#) is an [extension](#).

14 4 An [abstract type](#) is a type that has the [ABSTRACT attribute](#).

#### NOTE 4.51

The [DEFERRED attribute](#) (4.5.5) defers the implementation of a [type-bound procedure](#) to [extensions](#) of the type; it can appear only in an [abstract type](#). The [dynamic type](#) of an object cannot be [abstract](#); therefore, a deferred [type-bound procedure](#) cannot be invoked. An [extension](#) of an [abstract type](#) need not be [abstract](#) if it has no deferred [type-bound procedures](#). A short example of an [abstract type](#) is:

```
TYPE, ABSTRACT :: FILE_HANDLE
CONTAINS
  PROCEDURE(OPEN\_FILE), DEFERRED, PASS(HANDLE) :: OPEN
  ...
END TYPE
```

For a more elaborate example see [C.1.4](#).

### 15 4.5.7.2 Inheritance

16 1 An [extended type](#) includes all of the type parameters, all of the components, and the nonoverridden (4.5.7.3)  
 17 [type-bound procedures](#) of its [parent type](#). These are [inherited](#) by the [extended type](#) from the [parent type](#). They  
 18 retain all of the attributes that they had in the [parent type](#). Additional type parameters, components, and  
 19 procedure [bindings](#) may be declared in the derived-type definition of the [extended type](#).

**NOTE 4.52**

Inaccessible components and [bindings](#) of the [parent type](#) are also [inherited](#), but they remain inaccessible in the [extended type](#). Inaccessible entities occur if the type being extended is accessed via [use association](#) and has a private entity.

**NOTE 4.53**

A derived type is not required to have any components, [bindings](#), or parameters; an [extended type](#) is not required to have more components, [bindings](#), or parameters than its [parent type](#).

- 1 2 An [extended type](#) has a scalar, nonpointer, nonallocatable, [parent component](#) with the type and type parameters of the [parent type](#). The name of this component is the [parent type](#) name. If the [extended type](#) is defined in a module, the [parent component](#) has the accessibility of the [parent type](#) in the module in which the [parent type](#) was defined. Components of the [parent component](#) are [inheritance associated](#) (16.5.4) with the corresponding components [inherited](#) from the [parent type](#). An ancestor component of a type is the [parent component](#) of the type or an ancestor component of the [parent component](#).
- 7 3 If a generic [binding](#) specified in a type definition has the same [generic-spec](#) as an [inherited binding](#), it extends the [generic interface](#) and shall satisfy the requirements specified in 12.4.3.5.5.

**NOTE 4.54**

A component or type parameter declared in an [extended type](#) shall not have the same name as any accessible component or type parameter of its [parent type](#).

**NOTE 4.55**

Examples:

```

TYPE POINT                                ! A base type
  REAL :: X, Y
END TYPE POINT

TYPE, EXTENDS(POINT) :: COLOR_POINT      ! An extension of TYPE(POINT)
  ! Components X and Y, and component name POINT, inherited from parent
  INTEGER :: COLOR
END TYPE COLOR_POINT

```

### 9 4.5.7.3 Type-bound procedure overriding

- 10 1 If a specific [type-bound procedure](#) specified in a type definition has the same [binding name](#) as an accessible [type-bound procedure](#) from the [parent type](#) then the [binding](#) specified in the type definition overrides the one from the [parent type](#).
- 13 2 The overriding and overridden [type-bound procedures](#) shall satisfy the following conditions.
  - 14 • Either both shall have a [passed-object dummy argument](#) or neither shall.
  - 15 • If the overridden [type-bound procedure](#) is pure then the overriding one shall also be pure.
  - 16 • Either both shall be [elemental](#) or neither shall.
  - 17 • They shall have the same number of [dummy arguments](#).
  - 18 • [Passed-object dummy arguments](#), if any, shall correspond by name and position.
  - 19 • [Dummy arguments](#) that correspond by position shall have the same names and [characteristics](#), except for the type of the [passed-object dummy arguments](#).
  - 20 • Either both shall be subroutines or both shall be functions having the same result [characteristics](#) (12.3.3).
  - 21 • If the overridden [type-bound procedure](#) is [PUBLIC](#) then the overriding one shall not be [PRIVATE](#).
  - 22

- 1 3 A **binding** of a type and a **binding** of an **extension** of that type correspond if the latter **binding** is the same **binding**  
 2 as the former, overrides a corresponding **binding**, or is an **inherited** corresponding **binding**.

**NOTE 4.56**

The following is an example of procedure overriding, expanding on the example in Note 4.42.

```
TYPE, EXTENDS (POINT) :: POINT_3D
  REAL :: Z
CONTAINS
  PROCEDURE, PASS :: LENGTH => POINT_3D_LENGTH
END TYPE POINT_3D
...
```

and in the *module-subprogram-part* of the same module:

```
REAL FUNCTION POINT_3D_LENGTH ( A, B )
  CLASS (POINT_3D), INTENT (IN) :: A
  CLASS (POINT), INTENT (IN) :: B
  SELECT TYPE(B)
    CLASS IS(POINT_3D)
      POINT_3D_LENGTH = SQRT( (A%X-B%X)**2 + (A%Y-B%Y)**2 + (A%Z-B%Z)**2 )
      RETURN
    END SELECT
  PRINT *, 'In POINT_3D_LENGTH, dynamic type of argument is incorrect.'
  STOP
END FUNCTION POINT_3D_LENGTH
```

### 3 4.5.8 Derived-type values

- 4 1 The component value of
- 5 • a pointer component is its **pointer association**,
  - 6 • an **allocatable** component is its allocation status and, if it is allocated, its **dynamic type** and type parameters,
  - 7 **bounds** and value, and
  - 8 • a nonpointer nonallocatable component is its value.
- 9 2 The set of values of a particular derived type consists of all possible sequences of the component values of its  
 10 components.

### 11 4.5.9 Derived-type specifier

- 12 1 A derived-type specifier is used in several contexts to specify a particular derived type and type parameters.
- 13 R454 *derived-type-spec* is *type-name* [ ( *type-param-spec-list* ) ]
- 14 R455 *type-param-spec* is [ *keyword* = ] *type-param-value*
- 15 C488 (R454) *type-name* shall be the name of an accessible derived type.
- 16 C489 (R454) *type-param-spec-list* shall appear only if the type is parameterized.
- 17 C490 (R454) There shall be at most one *type-param-spec* corresponding to each parameter of the type. If a  
 18 type parameter does not have a default value, there shall be a *type-param-spec* corresponding to that  
 19 type parameter.
- 20 C491 (R455) The *keyword*= may be omitted from a *type-param-spec* only if the *keyword*= has been omitted  
 21 from each preceding *type-param-spec* in the *type-param-spec-list*.

- 1 C492 (R455) Each *keyword* shall be the name of a parameter of the type.
- 2 C493 (R455) An asterisk may be used as a *type-param-value* in a *type-param-spec* only in the declaration of a  
3 *dummy argument* or *associate name* or in the allocation of a *dummy argument*.
- 4 2 Type parameter values that do not have *type parameter keywords* specified correspond to type parameters in type  
5 parameter order (4.5.3.2). If a *type parameter keyword* appears, the value corresponds to the type parameter  
6 named by the keyword. If necessary, the value is converted according to the rules of intrinsic assignment (7.2.1.3)  
7 to a value of the same kind as the type parameter.
- 8 3 The value of a type parameter for which no *type-param-value* has been specified is its default value.

#### 9 4.5.10 Construction of derived-type values

- 10 1 A derived-type definition implicitly defines a corresponding *structure constructor* that allows construction of  
11 scalar values of that derived type. The type and type parameters of a constructed value are specified by a derived  
12 type specifier.
- 13 R456 *structure-constructor* is *derived-type-spec* ( [ *component-spec-list* ] )
- 14 R457 *component-spec* is [ *keyword* = ] *component-data-source*
- 15 R458 *component-data-source* is *expr*  
16 or *data-target*  
17 or *proc-target*
- 18 C494 (R456) The *derived-type-spec* shall not specify an *abstract type* (4.5.7).
- 19 C495 (R456) At most one *component-spec* shall be provided for a component.
- 20 C496 (R456) If a *component-spec* is provided for an ancestor component, a *component-spec* shall not be provided  
21 for any component that is *inheritance associated* with a *subcomponent* of that ancestor component.
- 22 C497 (R456) A *component-spec* shall be provided for a nonallocatable component unless it has *default initializ-*  
23 *ation* or is *inheritance associated* with a *subcomponent* of another component for which a *component-spec*  
24 is provided.
- 25 C498 (R457) The *keyword*= may be omitted from a *component-spec* only if the *keyword*= has been omitted  
26 from each preceding *component-spec* in the constructor.
- 27 C499 (R457) Each *keyword* shall be the name of a component of the type.
- 28 C4100 (R456) The type name and all components of the type for which a *component-spec* appears shall be  
29 accessible in the *scoping unit* containing the *structure constructor*.
- 30 C4101 (R456) If *derived-type-spec* is a type name that is the same as a generic name, the *component-spec-list*  
31 shall not be a valid *actual-arg-spec-list* for a function reference that is resolvable as a generic reference to  
32 that name (12.5.5.2).
- 33 C4102 (R458) A *data-target* shall correspond to a data pointer component; a *proc-target* shall correspond to a  
34 procedure pointer component.
- 35 C4103 (R458) A *data-target* shall have the same *rank* as its corresponding component.

#### NOTE 4.57

The form 'name(...)' is interpreted as a generic *function-reference* if possible; it is interpreted as a *structure-constructor* only if it cannot be interpreted as a generic *function-reference*.

- 36 2 In the absence of a *component keyword*, each *component-data-source* is assigned to the corresponding component

1 in [component order](#) (4.5.4.7). If a [component keyword](#) appears, the *expr* is assigned to the component named  
 2 by the keyword. For a nonpointer component, the [declared type](#) and type parameters of the component and  
 3 *expr* shall conform in the same way as for a [variable](#) and *expr* in an [intrinsic assignment statement](#) (7.2.1.2), as  
 4 specified in Table 7.8. If necessary, each value of intrinsic type is converted according to the rules of intrinsic  
 5 assignment (7.2.1.3) to a value that agrees in type and type parameters with the corresponding component of  
 6 the derived type. For a nonpointer nonallocatable component, the shape of the expression shall conform with the  
 7 shape of the component.

8 3 If a component with [default initialization](#) has no corresponding [component-data-source](#), then the [default initial-](#)  
 9 [ization](#) is applied to that component. If an [allocatable](#) component has no corresponding [component-data-source](#),  
 10 then that component has an allocation status of unallocated.

#### NOTE 4.58

Because no [parent components](#) appear in the defined [component ordering](#), a value for a [parent component](#) can be specified only with a [component keyword](#). Examples of equivalent values using types defined in Note 4.55:

```
! Create values with components x = 1.0, y = 2.0, color = 3.
TYPE(POINT) :: PV = POINT(1.0, 2.0)      ! Assume components of TYPE(POINT)
                                           ! are accessible here.
...
COLOR_POINT( point=point(1,2), color=3)  ! Value for parent component
COLOR_POINT( point=PV, color=3)          ! Available even if TYPE(point)
                                           ! has private components
COLOR_POINT( 1, 2, 3)                    ! All components of TYPE(point)
                                           ! need to be accessible.
```

11 4 A [structure constructor](#) shall not appear before the referenced type is defined.

#### NOTE 4.59

This example illustrates a derived-type constant expression using a derived type defined in Note 4.15:

```
PERSON (21, 'JOHN SMITH')
```

This could also be written as

```
PERSON (NAME = 'JOHN SMITH', AGE = 21)
```

#### NOTE 4.60

An example constructor using the derived type GENERAL\_POINT defined in Note 4.22 is

```
general_point(dim=3) ( [ 1., 2., 3. ] )
```

12 5 For a pointer component, the corresponding [component-data-source](#) shall be an allowable [data-target](#) or [proc-](#)  
 13 [target](#) for such a pointer in a [pointer assignment statement](#) (7.2.2). If the component data source is a pointer,  
 14 the association of the component is that of the pointer; otherwise, the component is [pointer associated](#) with the  
 15 component data source.

#### NOTE 4.61

For example, if the variable TEXT were declared (5.2) to be

```
CHARACTER, DIMENSION (1:400), TARGET :: TEXT
```

and BIBLIO were declared using the derived-type definition REFERENCE in Note 4.30

**NOTE 4.61** (cont.)

TYPE (REFERENCE) :: BIBLIO

the statement

```
BIBLIO = REFERENCE (1, 1987, 1, "This is the title of the referenced &
&paper", SYNOPSIS=TEXT)
```

is valid and associates the pointer component SYNOPSIS of the object BIBLIO with the [target](#) object TEXT. The keyword SYNOPSIS is required because the fifth component of the type REFERENCE is a procedure pointer component, not a data pointer component of type character. It is not necessary to specify a [proc-target](#) for the procedure pointer component because it has [default initialization](#).

- 6 If a component of a derived type is **allocatable**, the corresponding constructor expression shall either be a reference  
to the intrinsic function **NULL** with no arguments, an **allocatable** entity of the same **rank**, or shall evaluate to an  
entity of the same **rank**. If the expression is a reference to the intrinsic function **NULL**, the corresponding com-  
ponent of the constructor has a status of unallocated. If the expression is an **allocatable** entity, the corresponding  
component of the constructor has the same allocation status as that **allocatable** entity and, if it is allocated, the  
same **dynamic type**, **bounds**, and value; if a length parameter of the component is **deferred**, its value is the same  
as the corresponding parameter of the expression. Otherwise the corresponding component of the constructor  
has an allocation status of allocated and has the same **bounds** and value as the expression.

**NOTE 4.62**

When the constructor is an **actual argument**, the allocation status of the **allocatable** component is available through the associated **dummy argument**.

### 9 4.5.11 Derived-type operations and assignment

- 1 Intrinsic assignment of derived-type entities is described in 7.2.1. This part of ISO/IEC 1539 does not specify  
 11 any intrinsic operations on derived-type entities. Any operation on derived-type entities or defined assignment  
 12 (7.2.1.4) for derived-type entities shall be defined explicitly by a function or a subroutine, and a [generic interface](#)  
 13 (4.5.5, 12.4.3.2).

## 4.6 Enumerations and enumerators

- 1 An enumeration is a set of enumerators. An enumerator is a named integer constant. An enumeration definition  
16 specifies the enumeration and its set of enumerators of the corresponding integer kind.

```

17      R459      enum-def                is  enum-def-stmt
18                                     enumerator-def-stmt
19                                     [ enumerator-def-stmt ] ...
20                                     end-enum-stmt

```

```

21      R460  enum-def-stmt      is  ENUM, BIND(C)

```

```
22      R461  enumerator-def-stmt      is  ENUMERATOR [ :: ] enumerator-list
```

23      R462    *enumerator*                    is    *named-constant* [ = *scalar-int-constant-expr* ]

```

24      R463  end-enum-stmt          is  END ENUM

```

25 C4104 (R461) If = appears in an *enumerator*, a double-colon separator shall appear before the *enumerator-list*.

- 2 For an enumeration, the kind is selected such that an integer type with that kind is [interoperable \(15.3.2\)](#) with the  
corresponding C enumeration type. The corresponding C enumeration type is the type that would be declared  
by a C enumeration specifier (6.7.2.2 of ISO/IEC 9899:2011) that specified C enumeration constants with the

- 1 same values as those specified by the *enum-def*, in the same order as specified by the *enum-def*.
- 2 3 The *companion processor* (2.5.7) shall be one that uses the same representation for the types declared by all C
- 3 enumeration specifiers that specify the same values in the same order.

**NOTE 4.63**

If a *companion processor* uses an unsigned type to represent a given enumeration type, the Fortran processor will use the signed integer type of the same width for the enumeration, even though some of the values of the enumerators cannot be represented in this signed integer type. The types of any such enumerators will be *interoperable* with the type declared in the C enumeration.

**NOTE 4.64**

ISO/IEC 9899:2011 guarantees the enumeration constants fit in a C int (6.7.2.2 of ISO/IEC 9899:2011). Therefore, the Fortran processor can evaluate all enumerator values using the integer type with kind parameter C\_INT, and then determine the kind parameter of the integer type that is *interoperable* with the corresponding C enumerated type.

**NOTE 4.65**

ISO/IEC 9899:2011 specifies that two enumeration types are compatible only if they specify enumeration constants with the same names and same values in the same order. This part of ISO/IEC 1539 further requires that a C processor that is to be a *companion processor* of a Fortran processor use the same representation for two enumeration types if they both specify enumeration constants with the same values in the same order, even if the names are different.

- 4 4 An enumerator is treated as if it were explicitly declared with the *PARAMETER attribute*. The enumerator is
- 5 defined in accordance with the rules of intrinsic assignment (7.2) with the value determined as follows.
- 6 (1) If *scalar-int-constant-expr* is specified, the value of the enumerator is the result of *scalar-int-constant-*
- 7 *expr*.
- 8 (2) If *scalar-int-constant-expr* is not specified and the enumerator is the first enumerator in *enum-def*,
- 9 the enumerator has the value 0.
- 10 (3) If *scalar-int-constant-expr* is not specified and the enumerator is not the first enumerator in *enum-*
- 11 *def*, its value is the result of adding 1 to the value of the enumerator that immediately precedes it
- 12 in the *enum-def*.

**NOTE 4.66**

Example of an enumeration definition:

```
ENUM, BIND(C)
  ENUMERATOR :: RED = 4, BLUE = 9
  ENUMERATOR YELLOW
END ENUM
```

The *kind type parameter* for this enumeration is processor dependent, but the processor is required to select a kind sufficient to represent the values 4, 9, and 10, which are the values of its enumerators. The following declaration might be equivalent to the above enumeration definition.

```
INTEGER(SELECTED_INT_KIND(2)), PARAMETER :: RED = 4, BLUE = 9, YELLOW = 10
```

An entity of the same *kind type parameter* value can be declared using the intrinsic function *KIND* with one of the enumerators as its argument, for example

```
INTEGER(KIND(RED)) :: X
```

**NOTE 4.67**

There is no difference in the effect of declaring the enumerators in multiple ENUMERATOR statements or in a single ENUMERATOR statement. The order in which the enumerators in an enumeration definition are declared is significant, but the number of ENUMERATOR statements is not.

## 4.7 Binary, octal, and hexadecimal literal constants

- 1 A binary, octal, or hexadecimal constant (*boz-literal-constant*) is a sequence of digits that represents an ordered sequence of bits. Such a constant has no type.

R464 *boz-literal-constant*            **is** *binary-constant*  
    **or** *octal-constant*  
    **or** *hex-constant*

R465 *binary-constant*            **is** B ' *digit* [ *digit* ] ... '  
    **or** B " *digit* [ *digit* ] ... "

C4105 (R465) *digit* shall have one of the values 0 or 1.

R466 *octal-constant*            **is** O ' *digit* [ *digit* ] ... '  
    **or** O " *digit* [ *digit* ] ... "

C4106 (R466) *digit* shall have one of the values 0 through 7.

R467 *hex-constant*            **is** Z ' *hex-digit* [ *hex-digit* ] ... '  
    **or** Z " *hex-digit* [ *hex-digit* ] ... "

R468 *hex-digit*            **is** *digit*  
    **or** A  
    **or** B  
    **or** C  
    **or** D  
    **or** E  
    **or** F

- 2 The *hex-digits* A through F represent the numbers ten through fifteen, respectively; they may be represented by their lower-case equivalents. Each digit of a *boz-literal-constant* represents a sequence of bits, according to its numerical interpretation, using the model of 13.3, with *z* equal to one for binary constants, three for octal constants or four for hexadecimal constants. A *boz-literal-constant* represents a sequence of bits that consists of the concatenation of the sequences of bits represented by its digits, in the order the digits are specified. The positions of bits in the sequence are numbered from right to left, with the position of the rightmost bit being zero. The length of a sequence of bits is the number of bits in the sequence. The processor shall allow the position of the leftmost nonzero bit to be at least  $z - 1$ , where *z* is the maximum value that could result from invoking the intrinsic function `STORAGE_SIZE` (13.9.184) with an argument that is a real or integer scalar of any kind supported by the processor.

C4107 (R464) A *boz-literal-constant* shall appear only as a *data-stmt-constant* in a **DATA statement**, or where explicitly allowed in subclause 13.9 as an actual argument of an **intrinsic** procedure.

## 4.8 Construction of array values

- 1 An array constructor constructs a rank-one array value from a sequence of scalar values, array values, and implied DO loops.

R469 *array-constructor*            **is** (/ *ac-spec* /)  
    **or** *lbracket ac-spec rbracket*

- 1 R470 *ac-spec* is *type-spec* ::  
2 or [*type-spec* ::] *ac-value-list*
- 3 R471 *lbracket* is [  
4 R472 *rbracket* is ]
- 5 R473 *ac-value* is *expr*  
6 or *ac-implied-do*
- 7 R474 *ac-implied-do* is ( *ac-value-list* , *ac-implied-do-control* )
- 8 R475 *ac-implied-do-control* is [ *integer-type-spec* :: ] *ac-do-variable* = *scalar-int-expr* , ■  
9 ■ *scalar-int-expr* [ , *scalar-int-expr* ]
- 10 R476 *ac-do-variable* is *do-variable*
- 11 C4108 (R470) If *type-spec* is omitted, each *ac-value* expression in the *array-constructor* shall have the same  
12 declared type and kind type parameters.
- 13 C4109 (R470) If *type-spec* specifies an intrinsic type, each *ac-value* expression in the *array-constructor* shall be  
14 of an intrinsic type that is in type conformance with a variable of type *type-spec* as specified in Table 7.8.
- 15 C4110 (R470) If *type-spec* specifies a derived type, the declared type of each *ac-value* expression in the *array-*  
16 *constructor* shall be that derived type and shall have the same kind type parameter values as specified  
17 by *type-spec*.
- 18 C4111 (R473) An *ac-value* shall not be unlimited polymorphic.
- 19 C4112 (R473) The declared type of an *ac-value* shall not be abstract.
- 20 C4113 (R474) The *ac-do-variable* of an *ac-implied-do* that is in another *ac-implied-do* shall not appear as the  
21 *ac-do-variable* of the containing *ac-implied-do*.
- 22 2 If *type-spec* is omitted, corresponding length type parameters of the declared type of each *ac-value* expression  
23 shall have the same value; in this case, the declared type and type parameters of the array constructor are those  
24 of the *ac-value* expressions.
- 25 3 If *type-spec* appears, it specifies the declared type and type parameters of the array constructor. Each *ac-value*  
26 expression in the *array-constructor* shall be compatible with intrinsic assignment to a variable of this type and  
27 type parameters. Each value is converted to the type and type parameters of the *array-constructor* in accordance  
28 with the rules of intrinsic assignment (7.2.1.3).
- 29 4 The dynamic type of an array constructor is the same as its declared type.
- 30 5 The character length of an *ac-value* in an *ac-implied-do* whose iteration count is zero shall not depend on the  
31 value of the *ac-do-variable* and shall not depend on the value of an expression that is not a constant expression.
- 32 6 If an *ac-value* is a scalar expression, its value specifies an element of the array constructor. If an *ac-value* is  
33 an array expression, the values of the elements of the expression, in array element order (6.5.3.2), specify the  
34 corresponding sequence of elements of the array constructor. If an *ac-value* is an *ac-implied-do*, it is expanded to  
35 form a sequence of elements under the control of the *ac-do-variable*, as in the DO construct (8.1.7.4).
- 36 7 For an *ac-implied-do*, the loop initialization and execution is the same as for a DO construct.
- 37 8 An empty sequence forms a zero-sized array.

**NOTE 4.68**

A one-dimensional array can be reshaped into any allowable array shape using the intrinsic function **RESHAPE** (13.9.163). An example is:

**NOTE 4.68 (cont.)**

```
X = (/ 3.2, 4.01, 6.5 /)
Y = RESHAPE (SOURCE = [ 2.0, [ 4.5, 4.5 ], X ], SHAPE = [ 3, 2 ])
```

This results in Y having the  $3 \times 2$  array of values:

```
2.0    3.2
4.5    4.01
4.5    6.5
```

**NOTE 4.69**

Examples of array constructors containing an implied DO are:

```
(/ (I, I = 1, 1075) /)
```

and

```
[ 3.6, (3.6 / I, I = 1, N) ]
```

**NOTE 4.70**

Using the type definition for PERSON in Note 4.15, an example of the construction of a derived-type array value is:

```
[ PERSON (40, 'SMITH'), PERSON (20, 'JONES') ]
```

**NOTE 4.71**

Using the type definition for LINE in Note 4.26, an example of the construction of a derived-type scalar value with a rank-2 array component is:

```
LINE (RESHAPE ( [ 0.0, 0.0, 1.0, 2.0 ], [ 2, 2 ] ), 0.1, 1)
```

The intrinsic function **RESHAPE** is used to construct a value that represents a solid line from (0, 0) to (1, 2) of width 0.1 centimeters.

**NOTE 4.72**

Examples of zero-size array constructors are:

```
[ INTEGER :: ]
[ ( I, I = 1, 0) ]
```

**NOTE 4.73**

An example of an array constructor that specifies a length type parameter:

```
[ CHARACTER(LEN=7) :: 'Takata', 'Tanaka', 'Hayashi' ]
```

In this constructor, without the type specification, it would have been necessary to specify all of the constants with the same character length.

## 5 Attribute declarations and specifications

### 5.1 Attributes of procedures and data objects

- 1 Every data object has a type and *rank* and may have *type parameters* and other properties that determine the uses of the object. Collectively, these properties are the *attributes* of the object. The *declared type* of a named data object is either specified explicitly in a *type declaration statement* or determined implicitly by the first letter of its name (5.7). All of its *attributes* may be specified in a *type declaration statement* or individually in separate specification statements.
- 2 A function has a type and *rank* and may have type parameters and other *attributes* that determine the uses of the function. The type, *rank*, and type parameters are the same as those of the *function result*.
- 3 A subroutine does not have a type, *rank*, or type parameters, but may have other *attributes* that determine the uses of the subroutine.

### 5.2 Type declaration statement

R501 *type-declaration-stmt* is *declaration-type-spec* [ [ , *attr-spec* ] ... :: ] *entity-decl-list*

- 1 The type declaration statement specifies the *declared type* of the entities in the entity declaration list. The type and type parameters are those specified by *declaration-type-spec*, except that the character length type parameter may be overridden for an entity by the appearance of \* *char-length* in its *entity-decl*.

R502 *attr-spec* is *access-spec*  
or ALLOCATABLE  
or ASYNCHRONOUS  
or CODIMENSION *lbracket coarray-spec rbracket*  
or CONTIGUOUS  
or DIMENSION ( *array-spec* )  
or EXTERNAL  
or INTENT ( *intent-spec* )  
or INTRINSIC  
or *language-binding-spec*  
or OPTIONAL  
or PARAMETER  
or POINTER  
or PROTECTED  
or SAVE  
or TARGET  
or VALUE  
or VOLATILE

C501 (R501) The same *attr-spec* shall not appear more than once in a given *type-declaration-stmt*.

C502 (R501) If a *language-binding-spec* with a NAME= specifier appears, the *entity-decl-list* shall consist of a single *entity-decl*.

C503 (R501) If a *language-binding-spec* is specified, the *entity-decl-list* shall not contain any procedure names.

- 2 The type declaration statement also specifies the *attributes* whose keywords appear in the *attr-spec*, except that

the **DIMENSION** attribute may be specified or overridden for an entity by the appearance of *array-spec* in its *entity-decl*, and the **CODIMENSION** attribute may be specified or overridden for an entity by the appearance of *coarray-spec* in its *entity-decl*.

R503 *entity-decl* is *object-name* [( *array-spec* )] ■  
 ■ [ *lbracket coarray-spec rbracket* ] ■  
 ■ [ \* *char-length* ] [ *initialization* ]  
 or *function-name* [ \* *char-length* ]

C504 (R503) If the entity is not of type character, \* *char-length* shall not appear.

C505 (R501) If *initialization* appears, a double-colon separator shall appear before the *entity-decl-list*.

C506 (R501) If the PARAMETER keyword appears, *initialization* shall appear in each *entity-decl*.

C507 (R503) An *initialization* shall not appear if *object-name* is a **dummy argument**, a function result, an object in a named **common block** unless the type declaration is in a block data program unit, an object in **blank common**, an **allocatable** variable, or an **automatic data object**.

C508 (R503) The *function-name* shall be the name of an external function, an intrinsic function, a **dummy function**, a **procedure pointer**, or a statement function.

R504 *object-name* is *name*

C509 (R504) The *object-name* shall be the name of a data object.

R505 *initialization* is = *constant-expr*  
 or => *null-init*  
 or => *initial-data-target*

R506 *null-init* is *function-reference*

C510 (R503) If => appears in *initialization*, the entity shall have the **POINTER** attribute. If = appears in *initialization*, the entity shall not have the **POINTER** attribute.

C511 (R503) If *initial-data-target* appears, *object-name* shall be data-pointer-initialization compatible with it (4.5.4.6).

C512 (R506) The *function-reference* shall be a reference to the intrinsic function **NULL** with no arguments.

3 A name that identifies a specific intrinsic function has a type as specified in 13.8. An explicit type declaration statement is not required; however, it is permitted. Specifying a type for a generic intrinsic function name in a type declaration statement has no effect.

- 4 If *initialization* appears for a nonpointer entity,
- its type and **type parameters** shall conform as specified for intrinsic assignment (7.2.1.2);
  - if the entity has implied shape, the rank of *initialization* shall be the same as the rank of the entity;
  - if the entity does not have implied shape, *initialization* shall either be scalar or have the same shape as the entity.

#### NOTE 5.1

Examples of type declaration statements:

```
REAL A (10)
LOGICAL, DIMENSION (5, 5) :: MASK1, MASK2
COMPLEX :: CUBE_ROOT = (-0.5, 0.866)
INTEGER, PARAMETER :: SHORT = SELECTED_INT_KIND (4)
INTEGER (SHORT) K      ! Range at least -9999 to 9999.
```

## NOTE 5.1 (cont.)

```

REAL (KIND (0.0D0)) B1
REAL (KIND = 2) B2
COMPLEX (KIND = KIND (0.0D0)) :: C
CHARACTER (LEN = 10, KIND = 2) TEXT2
CHARACTER CHAR, STRING *20
TYPE (PERSON) :: CHAIRMAN
TYPE(NODE), POINTER :: HEAD => NULL ( )
TYPE (humongous_matrix (k=8, d=1000)) :: MAT

```

(The last line above uses a type definition from Note 4.22.)

## 5.3 Automatic data objects

- 1 An [automatic data object](#) is a [nondummy data object](#) with a [type parameter](#) or [array bound](#) that depends on the value of a [specification-expr](#) that is not a [constant expression](#).

C513 An [automatic data object](#) shall not have the [SAVE attribute](#).

- 2 If a type parameter in a [declaration-type-spec](#) or in a [char-length](#) in an [entity-decl](#) for a [local variable](#) of a subprogram or [BLOCK construct](#) is defined by an expression that is not a [constant expression](#), the type parameter value is established on entry to a procedure defined by the subprogram, or on execution of the [BLOCK statement](#), and is not affected by any redefinition or undefinition of the variables in the expression during execution of the procedure or [BLOCK construct](#).

## 5.4 Initialization

- 1 The appearance of [initialization](#) in an [entity-decl](#) for an entity without the [PARAMETER attribute](#) specifies that the entity is a variable with [explicit initialization](#). [Explicit initialization](#) alternatively may be specified in a [DATA statement](#) unless the variable is of a derived type for which [default initialization](#) is specified. If [initialization](#) is [= constant-expr](#), the variable is initially defined with the value specified by the [constant-expr](#); if necessary, the value is converted according to the rules of intrinsic assignment (7.2.1.3) to a value that agrees in type, type parameters, and shape with the variable. A variable, or part of a variable, shall not be [explicitly initialized](#) more than once in a program. If the variable is an array, it shall have its shape specified in either the type declaration statement or a previous attribute specification statement in the same [scoping unit](#).
- 2 If [null-init](#) appears, the initial association status of the object is [disassociated](#). If [initial-data-target](#) appears, the object is initially associated with the [target](#).
- 3 [Explicit initialization](#) of a variable that is not in a [common block](#) implies the [SAVE attribute](#), which may be confirmed by explicit specification.

## 5.5 Attributes

### 5.5.1 Attribute specification

- 1 An [attribute](#) may be explicitly specified by an [attr-spec](#) in a [type declaration statement](#) or by an attribute specification statement (5.6). The following constraints apply to [attributes](#).

C514 An entity shall not be explicitly given any [attribute](#) more than once in a [scoping unit](#).

C515 An [array-spec](#) for a nonallocatable nonpointer function result shall be an [explicit-shape-spec-list](#).

C516 The **ALLOCATABLE** or **POINTER** attribute shall not be specified for a **default-initialized dummy argument** of a procedure that has a *proc-language-binding-spec*.

### 5.5.2 Accessibility attribute

1 The accessibility attribute specifies the accessibility of an entity via a particular identifier.

R507 *access-spec* is PUBLIC  
or PRIVATE

C517 An *access-spec* shall appear only in the *specification-part* of a module.

2 An *access-spec* in a type declaration statement specifies the accessibility of the names of all the entities declared by that statement. An *access-spec* in a *derived-type-stmt* specifies the accessibility of the derived type name. Accessibility can also be specified by an *access-stmt*.

3 An identifier that is specified in a module or is accessible in a module by use association has either the **PUBLIC attribute** or **PRIVATE attribute**. An identifier whose accessibility is not explicitly specified has default accessibility (5.6.1).

4 The default accessibility attribute for a module is **PUBLIC** unless it has been changed by a **PRIVATE** statement. Only an identifier that has the **PUBLIC** attribute in that module is available to be accessed from that module by use association.

#### NOTE 5.2

In order for an identifier to be accessed by **use association**, it must have the **PUBLIC attribute** in the module from which it is accessed. It can nonetheless have the **PRIVATE attribute** in a module in which it is accessed by **use association**, and therefore not be available for **use association** from that module.

#### NOTE 5.3

An example of an accessibility specification is:

```
REAL, PRIVATE :: X, Y, Z
```

### 5.5.3 ALLOCATABLE attribute

1 An entity with the **ALLOCATABLE attribute** is a variable for which space is allocated by an **ALLOCATE statement** (6.7.1) or by an **intrinsic assignment statement** (7.2.1.2).

### 5.5.4 ASYNCHRONOUS attribute

1 An entity with the **ASYNCHRONOUS attribute** is a variable that may be subject to asynchronous input/output or asynchronous communication.

2 The **base object** of a variable shall have the **ASYNCHRONOUS attribute** in a **scoping unit** if

- the variable appears in an executable statement or **specification expression** in that **scoping unit** and
- any statement of the **scoping unit** is executed while the variable is a pending input/output storage sequence affector (9.6.2.5) or a pending communication affector (15.10.4).

3 Use of a variable in an asynchronous data transfer statement can imply the **ASYNCHRONOUS attribute**; see subclause 9.6.2.5.

4 An object with the **ASYNCHRONOUS attribute** may be associated with an object that does not have the **ASYNCHRONOUS attribute**, including by **use** (11.2.2) or **host** association (16.5.1.4). If an object that is not a **local variable** of a **BLOCK construct** is specified to have the **ASYNCHRONOUS attribute** in the *specification-part* of the construct, the object has the attribute within the construct even if it does not have the attribute

outside the construct. If an object has the **ASYNCHRONOUS attribute**, then all of its subobjects also have the **ASYNCHRONOUS attribute**.

#### NOTE 5.4

The **ASYNCHRONOUS attribute** specifies the variables that might be associated with a pending input/output storage sequence (the actual memory locations on which asynchronous input/output is being performed) while the **scoping unit** is in execution. This information could be used by the compiler to disable certain code motion optimizations.

### 5.5.5 BIND attribute for data entities

1 The **BIND attribute** for a variable or **common block** specifies that it is capable of interoperating with a C variable whose name has external linkage (15.9).

R508 *language-binding-spec* is BIND (C [, NAME = *scalar-default-char-constant-expr* ])

C518 An entity with the **BIND attribute** shall be a **common block**, variable, type, or procedure.

C519 A variable with the **BIND attribute** shall be declared in the specification part of a module.

C520 A variable with the **BIND attribute** shall be **interoperable** (15.3).

C521 Each variable of a **common block** with the **BIND attribute** shall be **interoperable**.

2 If the value of the *scalar-default-char-constant-expr* after discarding leading and trailing blanks has nonzero length, it shall be valid as an identifier on the **companion processor**.

#### NOTE 5.5

ISO/IEC 9899:2011 provides a facility for creating C identifiers whose characters are not restricted to the C basic character set. Such a C identifier is referred to as a universal character name (6.4.3 of ISO/IEC 9899:2011). The name of such a C identifier might include characters that are not part of the representation method used by the processor for default character. If so, the C entity cannot be referenced from Fortran.

3 The **BIND attribute** for a **common block** implies the **SAVE attribute**, which may be confirmed by explicit specification.

### 5.5.6 CODIMENSION attribute

#### 5.5.6.1 General

1 The **CODIMENSION attribute** specifies that an entity is a **coarray**. The *coarray-spec* specifies its **corank** or **corank** and **cobounds**.

R509 *coarray-spec* is *deferred-coshape-spec-list*  
or *explicit-coshape-spec*

C522 The sum of the **rank** and **corank** of an entity shall not exceed fifteen.

C523 A **coarray** shall be a component or a variable that is not a function result.

C524 A **coarray** shall not be of type **C\_PTR** or **C\_FUNPTR** from the intrinsic module **ISO\_C\_BINDING** (15.3.3).

C525 An entity whose type has a **coarray ultimate component** shall be a nonpointer nonallocatable scalar, shall not be a **coarray**, and shall not be a function result.

C526 A **coarray** or an object with a **coarray ultimate component** shall be a **dummy argument** or have the **ALLOCATABLE** or **SAVE** attribute.

C527 A **coarray** shall not be a **dummy argument** of a procedure that has a *proc-language-binding-spec*.

**NOTE 5.6**

A [coarray](#) is permitted to be of a derived type with pointer or [allocatable](#) components. The [target](#) of such a pointer component is always on the same [image](#) as the pointer.

**NOTE 5.7**

This requirement for the [SAVE attribute](#) has the effect that [automatic coarrays](#) are not permitted; for example, the [coarray](#) WORK in the following code fragment is not valid.

```
SUBROUTINE SOLVE3(N,A,B)
  INTEGER :: N
  REAL    :: A(N)[*], B(N)
  REAL    :: WORK(N)[*]    ! Not permitted
```

If this were permitted, it would require an implicit synchronization on entry to the procedure.

[Explicit-shape coarrays](#) that are declared in a subprogram and are not [dummy arguments](#) are required to have the [SAVE attribute](#) because otherwise they might be implemented as if they were [automatic coarrays](#).

**NOTE 5.8**

Examples of CODIMENSION attribute specifications are:

```
REAL W(100,100)[0:2,*]           ! Explicit-shape coarray
REAL, CODIMENSION[*] :: X        ! Scalar coarray
REAL, CODIMENSION[3,*] :: Y(:)   ! Assumed-shape coarray
REAL, CODIMENSION[:,],ALLOCATABLE :: Z(:,,:) ! Allocatable coarray
```

**5.5.6.2 Allocatable coarray**

1 A [coarray](#) with the [ALLOCATABLE attribute](#) has a specified [corank](#), but its [cobounds](#) are determined by  
3 allocation or [argument association](#).

R510 *deferred-coshape-spec* is :

C528 A [coarray](#) with the [ALLOCATABLE attribute](#) shall have a *coarray-spec* that is a *deferred-coshape-spec-list*.

2 The [corank](#) of an [allocatable coarray](#) is equal to the number of colons in its *deferred-coshape-spec-list*.

3 The [cobounds](#) of an unallocated [allocatable coarray](#) are undefined. No part of such a [coarray](#) shall be referenced  
9 or defined; however, the [coarray](#) may appear as an argument to an intrinsic [inquiry function](#) as specified in 13.1.

4 The [cobounds](#) of an allocated [allocatable coarray](#) are those specified when the [coarray](#) is allocated.

5 The [cobounds](#) of an [allocatable coarray](#) are unaffected by any subsequent redefinition or undefinition of the  
12 variables on which the [cobounds](#)' expressions depend.

**5.5.6.3 Explicit-coshape coarray**

1 An explicit-coshape coarray is a named [coarray](#) that has its [corank](#) and [cobounds](#) declared by an *explicit-coshape-spec*.  
15

R511 *explicit-coshape-spec* is [ [ *lower-cobound* : ] *upper-cobound*, ]... ■  
17 ■ [ *lower-cobound* : ] \*

C529 A nonallocatable [coarray](#) shall have a *coarray-spec* that is an *explicit-coshape-spec*.

2 The [corank](#) is equal to one plus the number of *upper-cobounds*.

1 R512 *lower-cobound* is *specification-expr*

2 R513 *upper-cobound* is *specification-expr*

3 C530 (R511) A *lower-cobound* or *upper-cobound* that is not a *constant expression* shall appear only in a sub-  
4 program, *BLOCK construct*, or *interface body*.

5 3 If an explicit-coshape *coarray* is a *local variable* of a subprogram or *BLOCK construct* and has *cobounds* that are  
6 not *constant expressions*, the *cobounds* are determined on entry to a procedure defined by the subprogram, or  
7 on execution of the *BLOCK statement*, by evaluating the *cobounds* expressions. The *cobounds* of such a *coarray*  
8 are unaffected by the redefinition or undefinition of any variable during execution of the procedure or *BLOCK*  
9 *construct*.

10 4 The values of each *lower-cobound* and *upper-cobound* determine the *cobounds* of the *coarray* along a particular  
11 *codimension*. The *cosubscript* range of the *coarray* in that *codimension* is the set of integer values between and  
12 including the lower and upper *cobounds*. If the lower *cobound* is omitted, the default value is 1. The upper  
13 *cobound* shall not be less than the lower *cobound*.

## 14 5.5.7 CONTIGUOUS attribute

15 C531 An entity with the *CONTIGUOUS attribute* shall be an *array pointer*, an *assumed-shape array*, or an  
16 *assumed-rank dummy data object*.

17 1 The *CONTIGUOUS attribute* specifies that an *assumed-shape array* is contiguous, that an *array pointer* can  
18 only be *pointer associated* with a *contiguous target*, or that an *assumed-rank dummy data object* is contiguous.

19 2 An object is *contiguous* if it is

- 20 (1) an object with the *CONTIGUOUS attribute*,
- 21 (2) a nonpointer whole array that is not *assumed-shape*,
- 22 (3) an *assumed-shape array* that is argument associated with an array that is contiguous,
- 23 (4) an *assumed-rank dummy data object* whose *effective argument* is contiguous,
- 24 (5) an array allocated by an *ALLOCATE statement*,
- 25 (6) a *pointer associated* with a contiguous *target*, or
- 26 (7) a nonzero-sized *array section* (6.5.3) provided that
  - 27 (a) its *base object* is *contiguous*,
  - 28 (b) it does not have a *vector subscript*,
  - 29 (c) the array element ordering of the elements of the section is the same as the array element  
30 ordering of those elements of the *base object*,
  - 31 (d) in the array element ordering of the *base object*, every element of the *base object* that is not  
32 an element of the section either precedes every element of the section or follows every element  
33 of the section,
  - 34 (e) if the array is of type character and a *substring-range* appears, the *substring-range* specifies all  
35 of the characters of the *parent-string* (6.4.1),
  - 36 (f) only its final *part-ref* has nonzero *rank*, and
  - 37 (g) it is not the real or imaginary part (6.4.4) of an array of type complex.

38 3 An object is not contiguous if it is an array subobject, and

- 39 • the object has two or more elements,
- 40 • the elements of the object in array element order are not consecutive in the elements of the *base object*,
- 41 • the object is not of type character with length zero, and
- 42 • the object is not of a derived type that has no *ultimate components* other than zero-sized arrays and  
43 characters with length zero.

- 1 4 It is processor dependent whether any other object is contiguous.

#### NOTE 5.9

If a derived type has only one component that is not zero-sized, it is processor dependent whether a structure component of a contiguous array of that type is contiguous. That is, the derived type might contain padding on some processors.

#### NOTE 5.10

The **CONTIGUOUS** attribute makes it easier for a processor to enable optimizations that depend on the memory layout of the object occupying a contiguous block of memory. Examples of **CONTIGUOUS** attribute specifications are:

```
REAL, POINTER, CONTIGUOUS      :: SPTR(:)
REAL, CONTIGUOUS, DIMENSION(:, :) :: D
```

## 2 5.5.8 DIMENSION attribute

### 3 5.5.8.1 General

- 4 1 The **DIMENSION** attribute specifies that an entity is **assumed-rank** or an array. An **assumed-rank dummy data**  
 5 **object** has the rank, shape, and size of its **effective argument**; otherwise, the **rank** or **rank** and shape is specified  
 6 by its **array-spec**.

```
7 R514 dimension-spec      is DIMENSION ( array-spec )
8 R515 array-spec          is explicit-shape-spec-list
9                          or assumed-shape-spec-list
10                         or deferred-shape-spec-list
11                         or assumed-size-spec
12                         or implied-shape-spec
13                         or implied-shape-or-assumed-size-spec
14                         or assumed-rank-spec
```

#### NOTE 5.11

The maximum **rank** of an entity is fifteen minus the **corank**.

#### NOTE 5.12

Examples of **DIMENSION** attribute specifications are:

```
SUBROUTINE EX (N, A, B)
  REAL, DIMENSION (N, 10) :: W      ! Automatic explicit-shape array
  REAL A (:), B (0:)                ! Assumed-shape arrays
  REAL, POINTER :: D (:, :)         ! Array pointer
  REAL, DIMENSION (:), POINTER :: P ! Array pointer
  REAL, ALLOCATABLE, DIMENSION (:) :: E ! Allocatable array
  REAL, PARAMETER :: V(0:*) = [0.1, 1.1] ! Implied-shape array
```

### 15 5.5.8.2 Explicit-shape array

```
16 R516 explicit-shape-spec      is [ lower-bound : ] upper-bound
17 R517 lower-bound              is specification-expr
18 R518 upper-bound              is specification-expr
```

C532 (R516) An *explicit-shape-spec* whose *bounds* are not *constant expressions* shall appear only in a subprogram, derived type definition, *BLOCK construct*, or *interface body*.

1 An *explicit-shape array* is an array whose shape is explicitly declared by an *explicit-shape-spec-list*. The *rank* is equal to the number of *explicit-shape-specs*.

2 An *explicit-shape array* that is a named local variable of a subprogram or *BLOCK construct* may have *bounds* that are not *constant expressions*. The *bounds*, and hence shape, are determined on entry to a procedure defined by the subprogram, or on execution of the *BLOCK statement*, by evaluating the *bounds' expressions*. The *bounds* of such an array are unaffected by the redefinition or undefinition of any variable during execution of the procedure or *BLOCK construct*.

3 The values of each *lower-bound* and *upper-bound* determine the bounds of the array along a particular dimension and hence the extent of the array in that dimension. If *lower-bound* appears it specifies the lower bound; otherwise the lower bound is 1. The value of a lower bound or an upper bound may be positive, negative, or zero. The subscript range of the array in that dimension is the set of integer values between and including the lower and upper bounds, provided the upper bound is not less than the lower bound. If the upper bound is less than the lower bound, the range is empty, the extent in that dimension is zero, and the array is of zero size.

#### 5.5.8.3 Assumed-shape array

1 An *assumed-shape array* is a nonallocatable nonpointer *dummy argument* array that takes its shape from its *effective argument*.

R519 *assumed-shape-spec* is [ *lower-bound* ] :

2 The *rank* is equal to the number of colons in the *assumed-shape-spec-list*.

3 The extent of a dimension of an *assumed-shape array dummy argument* is the extent of the corresponding dimension of its *effective argument*. If the lower bound value is  $d$  and the extent of the corresponding dimension of its *effective argument* is  $s$ , then the value of the upper bound is  $s + d - 1$ . If *lower-bound* appears it specifies the lower bound; otherwise the lower bound is 1.

#### 5.5.8.4 Deferred-shape array

1 A *deferred-shape array* is an *allocatable* array or an *array pointer*. (An *allocatable* array has the *ALLOCATABLE* attribute; an *array pointer* has the *POINTER* attribute.)

R520 *deferred-shape-spec* is :

C533 An array with the *POINTER* or *ALLOCATABLE* attribute shall have an *array-spec* that is a *deferred-shape-spec-list*.

2 The *rank* is equal to the number of colons in the *deferred-shape-spec-list*.

3 The size, bounds, and shape of an unallocated *allocatable* array or a *disassociated array pointer* are undefined. No part of such an array shall be referenced or defined; however, the array may appear as an argument to an intrinsic *inquiry function* as specified in 13.1.

4 The bounds of each dimension of an allocated *allocatable* array are those specified when the array is allocated or, if it is a *dummy argument*, when it is argument associated with an allocated *effective argument*.

5 The bounds of each dimension of an associated *array pointer*, and hence its shape, may be specified

- in an *ALLOCATE statement* (6.7.1) when the *target* is allocated,
- by pointer assignment (7.2.2), or
- if it is a *dummy argument*, by *argument association* with a nonpointer *actual argument* or an associated pointer *effective argument*.

6 The bounds of an [array pointer](#) or [allocatable](#) array are unaffected by any subsequent redefinition or undefinition of variables on which the bounds' expressions depend.

### 5.5.8.5 Assumed-size array

1 An [assumed-size array](#) is a [dummy argument](#) array whose size is assumed from that of its [effective argument](#), or the associate name of a [RANK \( \\* \)](#) block in a [SELECT RANK construct](#). The [rank](#) and extents may differ for the [effective](#) and [dummy](#) arguments; only the size of the [effective argument](#) is assumed by the [dummy argument](#). A [dummy argument](#) is declared to be an [assumed-size array](#) by an [assumed-size-spec](#) or an [implied-shape-or-assumed-size-spec](#).

R521 *assumed-implied-spec*            **is**   [\[ lower-bound : \] \\*](#)

R522 *assumed-size-spec*            **is**   [explicit-shape-spec-list](#), [assumed-implied-spec](#)

C534 An object whose array bounds are specified by an [assumed-size-spec](#) shall be a [dummy data object](#).

C535 An [assumed-size array](#) with the [INTENT \(OUT\)](#) attribute shall not be [polymorphic](#), [finalizable](#), of a type with an [allocatable ultimate component](#), or of a type for which [default initialization](#) is specified.

R523 *implied-shape-or-assumed-size-spec*   **is**   [assumed-implied-spec](#)

C536 An object whose array bounds are specified by an [implied-shape-or-assumed-size-spec](#) shall be a [dummy data object](#) or a [named constant](#).

2 The size of an [assumed-size array](#) is determined as follows.

- If the [effective argument](#) associated with the [assumed-size](#) dummy array is an array of any type other than default character, the size is that of the [effective argument](#).
- If the [actual argument](#) corresponding to the [assumed-size](#) dummy array is an array element of any type other than default character with a subscript order value of  $r$  ([6.5.3.2](#)) in an array of size  $x$ , the size of the dummy array is  $x - r + 1$ .
- If the [actual argument](#) is a default character array, default character array element, or a default character array element substring ([6.4.1](#)), and if it begins at [character storage unit](#)  $t$  of an array with  $c$  [character storage units](#), the size of the dummy array is  $\text{MAX}(\text{INT}((c - t + 1)/e), 0)$ , where  $e$  is the length of an element in the dummy character array.
- If the [actual argument](#) is a default character scalar that is not an array element or array element substring [designator](#), the size of the dummy array is  $\text{MAX}(\text{INT}(l/e), 0)$ , where  $e$  is the length of an element in the dummy character array and  $l$  is the length of the [actual argument](#).

3 The [rank](#) is equal to one plus the number of [explicit-shape-specs](#).

4 An [assumed-size array](#) has no upper bound in its last dimension and therefore has no extent in its last dimension and no shape. An [assumed-size array](#) shall not appear in a context that requires its shape.

5 If a list of [explicit-shape-specs](#) appears, it specifies the bounds of the first [rank](#)–1 dimensions. If [lower-bound](#) appears it specifies the lower bound of the last dimension; otherwise that lower bound is 1. An [assumed-size array](#) may be subscripted or sectioned ([6.5.3.3](#)). The upper bound shall not be omitted from a subscript triplet in the last dimension.

6 If an [assumed-size array](#) has bounds that are not [constant expressions](#), the bounds are determined on entry to the procedure. The bounds of such an array are unaffected by the redefinition or undefinition of any variable during execution of the procedure.

### 5.5.8.6 Implied-shape array

1 An implied-shape array is a [named constant](#) that takes its shape from the [constant-expr](#) in its declaration. A named constant is declared to be an implied-shape array with an [array-spec](#) that is an [implied-shape-or-assumed-](#)

1 *size-spec* or an *implied-shape-spec*.

2 R524 *implied-shape-spec* is *assumed-implied-spec*, *assumed-implied-spec-list*

3 C537 An implied-shape array shall be a **named constant**.

4 2 The **rank** of an implied-shape array is the number of *assumed-implied-specs* in its *array-spec*.

5 3 The extent of each dimension of an implied-shape array is the same as the extent of the corresponding dimension  
6 of the *constant-expr*. The lower bound of each dimension is *lower-bound*, if it appears, and 1 otherwise; the upper  
7 bound is one less than the sum of the lower bound and the extent.

#### 8 5.5.8.7 Assumed-rank entity

9 1 An assumed-rank entity is a **dummy data object** whose **rank** is assumed from its **effective argument**, or the associate  
10 name of a **RANK DEFAULT** block in a **SELECT RANK construct**; this rank can be zero. An assumed-rank  
11 entity is declared with an *array-spec* that is an *assumed-rank-spec*.

12 R525 *assumed-rank-spec* is ..

13 C538 An assumed-rank entity shall be a **dummy data object** that does not have the **CODIMENSION** or **VALUE**  
14 attribute.

15 C539 An assumed-rank variable name shall not appear in a **designator** or expression except as an **actual**  
16 **argument** that corresponds to a **dummy argument** that is assumed-rank, the argument of the function  
17 **C\_LOC** from the intrinsic module **ISO\_C\_BINDING** (15.2.3.6), the first **dummy argument** of an **intrinsic**  
18 **inquiry function**, or the selector of a **SELECT RANK statement**.

19 C540 If an **assumed-size** or nonallocatable nonpointer assumed-rank array is an **actual argument** that corres-  
20 ponds to a **dummy argument** that is an **INTENT (OUT)** assumed-rank array, it shall not be **polymorphic**,  
21 **finalizable**, of a type with an **allocatable ultimate component**, or of a type for which **default initialization**  
22 is specified.

#### 23 5.5.9 EXTERNAL attribute

24 1 The **EXTERNAL attribute** specifies that an entity is an **external procedure**, **dummy procedure**, **procedure pointer**,  
25 or block data subprogram.

26 C541 An entity shall not have both the **EXTERNAL attribute** and the **INTRINSIC attribute**.

27 C542 In an external subprogram, the **EXTERNAL attribute** shall not be specified for a procedure defined by  
28 the subprogram.

29 2 If an **external procedure** or **dummy procedure** is used as an **actual argument** or is the **target** of a procedure **pointer**  
30 **assignment**, it shall be declared to have the **EXTERNAL attribute**.

31 3 A procedure that has both the **EXTERNAL** and **POINTER** attributes is a **procedure pointer**.

#### NOTE 5.13

The **EXTERNAL attribute** can be specified in a **type declaration statement**, by an **interface body** (12.4.3.2),  
by an **EXTERNAL statement** (12.4.3.6), or by a **procedure declaration statement** (12.4.3.7).

#### 32 5.5.10 INTENT attribute

33 1 The **INTENT attribute** specifies the intended use of a **dummy argument**. An **INTENT (IN)** **dummy argument**  
34 is suitable for receiving data from the invoking **scoping unit**, an **INTENT (OUT)** **dummy argument** is suitable  
35 for returning data to the invoking **scoping unit**, and an **INTENT (INOUT)** **dummy argument** is suitable for use  
36 both to receive data from and to return data to the invoking **scoping unit**.

- 1 R526 *intent-spec* is IN  
 2 or OUT  
 3 or INOUT
- 4 C543 An entity with the **INTENT attribute** shall be a **dummy data object** or a dummy procedure pointer.
- 5 C544 (R526) A nonpointer object with the **INTENT (IN) attribute** shall not appear in a variable definition  
 6 context (16.6.7).
- 7 C545 A pointer with the **INTENT (IN) attribute** shall not appear in a pointer association context (16.6.8).
- 8 C546 An **INTENT (OUT) dummy argument** of a nonintrinsic procedure shall not be an **allocatable coarray** or  
 9 have a **subobject** that is an **allocatable coarray**.
- 10 C547 An entity with the **INTENT (OUT) attribute** shall not be of, or have a **subcomponent** of, type **EVENT\_-**  
 11 **TYPE** (13.10.2.10) or **LOCK\_TYPE** (13.10.2.19) from the intrinsic module **ISO\_FORTRAN\_ENV**.
- 12 2 The **INTENT (IN) attribute** for a nonpointer **dummy argument** specifies that it shall neither be defined nor  
 13 become undefined during the invocation and execution of the procedure. The **INTENT (IN) attribute** for a  
 14 pointer **dummy argument** specifies that during the invocation and execution of the procedure its association shall  
 15 not be changed except that it may become undefined if the **target** is deallocated other than through the pointer  
 16 (16.5.2.5).
- 17 3 The **INTENT (OUT) attribute** for a nonpointer **dummy argument** specifies that the **dummy argument** becomes  
 18 undefined on invocation of the procedure, except for any subcomponents that are **default-initialized** (4.5.4.6). Any  
 19 **actual argument** that corresponds to such a **dummy argument** shall be **definable**. The **INTENT (OUT) attribute**  
 20 for a pointer **dummy argument** specifies that on invocation of the procedure the **pointer association** status of the  
 21 **dummy argument** becomes undefined. Any **actual argument** that corresponds to such a pointer dummy shall be  
 22 a pointer variable. Any undefinition or definition implied by association of an **actual argument** with an **INTENT**  
 23 **(OUT) dummy argument** shall not affect any other entity within the statement that invokes the procedure.
- 24 4 The **INTENT (INOUT) attribute** for a nonpointer **dummy argument** specifies that any **actual argument** that  
 25 corresponds to the **dummy argument** shall be **definable**. The **INTENT (INOUT) attribute** for a pointer **dummy**  
 26 **argument** specifies that any **actual argument** that corresponds to the **dummy argument** shall be a pointer variable.

**NOTE 5.14**

The **INTENT attribute** for an **allocatable dummy argument** applies to both the allocation status and the definition status. An **actual argument** that corresponds to an **INTENT (OUT) allocatable dummy argument** is deallocated on procedure invocation (6.7.3.2). To avoid this deallocation for coarrays, **INTENT (OUT)** is not allowed for a **dummy argument** that is an allocatable coarray or has a subobject that is an allocatable coarray.

- 27 5 If no **INTENT attribute** is specified for a **dummy argument**, its use is subject to the limitations of its **effective**  
 28 **argument** (12.5.2).
- 29 6 If an object has an **INTENT attribute**, then all of its subobjects have the same **INTENT attribute**.

**NOTE 5.15**

An example of **INTENT** specification is:

```
SUBROUTINE MOVE (FROM, TO)
  USE PERSON_MODULE
  TYPE (PERSON), INTENT (IN) :: FROM
  TYPE (PERSON), INTENT (OUT) :: TO
```

**NOTE 5.16**

If a **dummy argument** is a derived-type object with a pointer component, then the pointer as a pointer is a subobject of the **dummy argument**, but the **target** of the pointer is not. Therefore, the restrictions on subobjects of the **dummy argument** apply to the pointer in contexts where it is used as a pointer, but not in contexts where it is dereferenced to indicate its **target**. For example, if X is a **dummy argument** of derived type with an integer pointer component P, and X is **INTENT (IN)**, then the statement

```
X%P => NEW_TARGET
```

is prohibited, but

```
X%P = 0
```

is allowed (provided that X%P is associated with a **definable target**).

Similarly, the **INTENT** restrictions on pointer **dummy arguments** apply only to the association of the **dummy argument**; they do not restrict the operations allowed on its **target**.

**NOTE 5.17**

Argument intent specifications serve several purposes in addition to documenting the intended use of **dummy arguments**. A processor can check whether an **INTENT (IN) dummy argument** is used in a way that could redefine it. A slightly more sophisticated processor could check to see whether an **INTENT (OUT) dummy argument** could possibly be referenced before it is defined. If the procedure's **interface** is explicit, the processor can also verify that **actual arguments** corresponding to **INTENT (OUT)** or **INTENT (INOUT) dummy arguments** are **definable**. A more sophisticated processor could use this information to optimize the translation of the referencing **scoping unit** by taking advantage of the fact that **actual arguments** corresponding to **INTENT (IN) dummy arguments** will not be changed and that any prior value of an **actual argument** corresponding to an **INTENT (OUT) dummy argument** will not be referenced and could thus be discarded.

**INTENT (OUT)** means that the value of the argument after invoking the procedure is entirely the result of executing that procedure. If an argument might not be redefined and it is desired to have the argument retain its value in that case, **INTENT (OUT)** cannot be used because it would cause the argument to become undefined; however, **INTENT (INOUT)** can be used, even if there is no explicit reference to the value of the dummy argument.

**INTENT (INOUT)** is not equivalent to omitting the **INTENT attribute**. The **actual argument** corresponding to an **INTENT (INOUT) dummy argument** is always required to be **definable**, while an **actual argument** corresponding to a **dummy argument** without an **INTENT attribute** need be **definable** only if the **dummy argument** is actually redefined.

**5.5.11 INTRINSIC attribute**

- 1 The **INTRINSIC attribute** specifies that the entity is an intrinsic procedure. The procedure name may be a generic name (13.7), a **specific name** (13.8), or both.
- 2 If the **specific name** of an intrinsic procedure (13.8) is used as an **actual argument**, the name shall be explicitly specified to have the **INTRINSIC attribute**. Note that a specific intrinsic procedure listed in Table 13.3 is not permitted to be used as an **actual argument** (C1240).
- C548 If the generic name of an intrinsic procedure is explicitly declared to have the **INTRINSIC attribute**, and it is also the generic name of one or more **generic interfaces** (12.4.3.2) accessible in the same **scoping unit**, the procedures in the interfaces and the generic intrinsic procedure shall all be functions or all be subroutines.

## 5.5.12 OPTIONAL attribute

- 1 The **OPTIONAL attribute** specifies that the **dummy argument** need not have a corresponding **actual argument** in a **reference** to the procedure (12.5.2.12).

C549 An entity with the **OPTIONAL attribute** shall be a **dummy argument**.

### NOTE 5.18

The intrinsic function **PRESENT** (13.9.152) can be used to determine whether an optional **dummy argument** has a corresponding **actual argument**.

## 5.5.13 PARAMETER attribute

- 1 The **PARAMETER attribute** specifies that an entity is a **named constant**. The entity has the value specified by its *constant-expr*, converted, if necessary, to the type, type parameters and shape of the entity.

C550 An entity with the **PARAMETER attribute** shall not be a **variable**, a **coarray**, or a **procedure**.

C551 An expression that specifies a **length type parameter** or array bound of a named constant shall be a constant expression.

- 2 A **named constant** shall not be referenced unless it has been defined previously; it may be defined previously in the same statement.

### NOTE 5.19

Examples of declarations with a **PARAMETER attribute** are:

```
REAL, PARAMETER :: ONE = 1.0, Y = 4.1 / 3.0
INTEGER, DIMENSION (3), PARAMETER :: ORDER = (/ 1, 2, 3 /)
TYPE(NODE), PARAMETER :: DEFAULT = NODE(0, NULL ( ))
```

## 5.5.14 POINTER attribute

- 1 Entities with the **POINTER attribute** can be associated with different data objects or procedures during execution of a program. A pointer is either a data pointer or a procedure pointer.

C552 An entity with the **POINTER attribute** shall not have the **ALLOCATABLE**, **INTRINSIC**, or **TARGET** attribute, and shall not be a **coarray**.

C553 A procedure with the **POINTER attribute** shall have the **EXTERNAL attribute**.

- 2 A data pointer shall not be referenced unless it is **pointer associated** with a **target** object that is defined. A data pointer shall not be defined unless it is **pointer associated** with a **target** object that is **definable**.

- 3 If a data pointer is associated, the values of its **deferred type parameters** are the same as the values of the corresponding type parameters of its **target**.

- 4 A **procedure pointer** shall not be referenced unless it is **pointer associated** with a **target** procedure.

### NOTE 5.20

Examples of **POINTER attribute** specifications are:

```
TYPE (NODE), POINTER :: CURRENT, TAIL
REAL, DIMENSION (:, :), POINTER :: IN, OUT, SWAP
```

### 5.5.15 PROTECTED attribute

1 The **PROTECTED attribute** imposes limitations on the usage of module entities.

C554 The **PROTECTED attribute** shall be specified only in the specification part of a module.

C555 An entity with the **PROTECTED attribute** shall be a procedure pointer or variable.

C556 An entity with the **PROTECTED attribute** shall not be in a **common block**.

C557 A nonpointer object that has the **PROTECTED attribute** and is accessed by **use association** shall not appear in a variable definition context (16.6.7) or as the *data-target* or *proc-target* in a *pointer-assignment-stmt*.

C558 A pointer that has the **PROTECTED attribute** and is accessed by **use association** shall not appear in a pointer association context (16.6.8).

2 Other than within the module in which an entity is given the **PROTECTED attribute**, or within any of its **descendants**,

- if it is a nonpointer object, it is not **definable**, and
- if it is a pointer, its association status shall not be changed except that it may become undefined if its **target** is deallocated other than through the pointer (16.5.2.5) or if its **target** becomes undefined by execution of a **RETURN** or **END** statement.

3 If an object has the **PROTECTED attribute**, all of its subobjects have the **PROTECTED attribute**.

#### NOTE 5.21

An example of the **PROTECTED attribute**:

```
MODULE temperature
  REAL, PROTECTED :: temp_c, temp_f
CONTAINS
  SUBROUTINE set_temperature_c(c)
    REAL, INTENT(IN) :: c
    temp_c = c
    temp_f = temp_c*(9.0/5.0) + 32
  END SUBROUTINE
END MODULE
```

The **PROTECTED attribute** ensures that the variables **temp\_c** and **temp\_f** cannot be modified other than via the **set\_temperature\_c** procedure, thus keeping them consistent with each other.

### 5.5.16 SAVE attribute

1 The **SAVE attribute** specifies that a local variable of a **program unit** or subprogram retains its association status, allocation status, definition status, and value after execution of a **RETURN** or **END** statement unless it is a pointer and its **target** becomes undefined (16.5.2.5(6)). If it is a local variable of a subprogram it is shared by all instances (12.6.2.4) of the subprogram.

2 The **SAVE attribute** specifies that a local variable of a **BLOCK construct** retains its association status, allocation status, definition status, and value after termination of the construct unless it is a pointer and its **target** becomes undefined (16.5.2.5(7)). If the **BLOCK construct** is within a subprogram the variable is shared by all instances (12.6.2.4) of the subprogram.

3 Giving a **common block** the **SAVE attribute** confers the attribute on all entities in the **common block**.

C559 An entity with the **SAVE attribute** shall be a **common block**, variable, or procedure pointer.

1 C560 The **SAVE attribute** shall not be specified for a **dummy argument**, a function result, an **automatic data**  
 2 **object**, or an object that is in a **common block**.

3 4 A variable, **common block**, or procedure pointer declared in the **scoping unit** of a main program, module, or  
 4 submodule implicitly has the **SAVE attribute**, which may be confirmed by explicit specification. If a **common block**  
 5 has the **SAVE attribute** in any other kind of **scoping unit**, it shall have the **SAVE attribute** in every **scoping unit** that is not of a  
 6 main program, module, or submodule.

## 7 5.5.17 TARGET attribute

8 1 The **TARGET attribute** specifies that a data object may have a **pointer associated** with it (7.2.2). An object  
 9 without the **TARGET attribute** shall not have a **pointer associated** with it.

10 C561 An entity with the **TARGET attribute** shall be a variable.

11 C562 An entity with the **TARGET attribute** shall not have the **POINTER attribute**.

12 2 If an object has the **TARGET attribute**, then all of its nonpointer subobjects also have the **TARGET attribute**.

### NOTE 5.22

In addition to variables explicitly declared to have the **TARGET attribute**, the objects created by allocation of pointers (6.7.1.4) have the **TARGET attribute**.

### NOTE 5.23

Examples of **TARGET attribute** specifications are:

```
TYPE (NODE), TARGET :: HEAD
REAL, DIMENSION (1000, 1000), TARGET :: A, B
```

### NOTE 5.24

Every **object designator** that starts from an object with the **TARGET attribute** will have either the **TARGET** or **POINTER** attribute. If pointers are involved, the **designator** might not necessarily be a subobject of the original object, but because a pointer can point only to an entity with the **TARGET attribute**, there is no way to end up at a nonpointer that does not have the **TARGET attribute**.

## 13 5.5.18 VALUE attribute

14 1 The **VALUE attribute** specifies a type of **argument association** (12.5.2.4) for a **dummy argument**.

15 C563 An entity with the **VALUE attribute** shall be a **dummy data object**. It shall not be an **assumed-size**  
 16 **array**, a **coarray**, or a variable with a **coarray ultimate component**.

17 C564 An entity with the **VALUE attribute** shall not have the **ALLOCATABLE**, **INTENT (INOUT)**, **INTENT**  
 18 **(OUT)**, **POINTER**, or **VOLATILE** attributes.

19 C565 A **dummy argument** of a procedure with the **BIND attribute** shall not have both the **OPTIONAL** and  
 20 **VALUE** attributes.

## 21 5.5.19 VOLATILE attribute

22 1 The **VOLATILE attribute** specifies that an object may be referenced, defined, or become undefined, by means  
 23 not specified by the program. A pointer with the **VOLATILE attribute** may additionally have its association  
 24 status, **dynamic type** and type parameters, and array bounds changed by means not specified by the program.  
 25 An **allocatable** object with the **VOLATILE attribute** may additionally have its allocation status, **dynamic type**  
 26 and type parameters, and array bounds changed by means not specified by the program.

- 1 C566 An entity with the **VOLATILE attribute** shall be a variable that is not an **INTENT (IN) dummy argu-**  
 2 **ment**.
- 3 C567 The **VOLATILE attribute** shall not be specified for a **coarray** that is accessed by use (11.2.2) or host  
 4 (16.5.1.4) association.
- 5 C568 Within a **BLOCK construct** (8.1.4), the **VOLATILE attribute** shall not be specified for a **coarray** that is  
 6 not a **construct entity** (16.4) of that construct.
- 7 2 A noncoarray object that has the **VOLATILE attribute** may be associated with an object that does not have  
 8 the **VOLATILE attribute**, including by use (11.2.2) or **host** association (16.5.1.4). If an object that is not a  
 9 **local variable** of a **BLOCK construct** is specified to have the **VOLATILE attribute** in the *specification-part* of  
 10 the construct, the object has the attribute within the construct even if it does not have the attribute outside the  
 11 construct. The relationship between **coarrays**, the **VOLATILE attribute**, and **argument association** is described  
 12 in 12.5.2.8. The relationship between **coarrays**, the **VOLATILE attribute**, and **pointer association** is  
 13 described in 7.2.2.3.
- 14 3 A pointer should have the **VOLATILE attribute** if its target has the **VOLATILE attribute**. If, by means not  
 15 specified by the program, the target is referenced, defined, or becomes undefined, the pointer shall have the  
 16 **VOLATILE attribute**. All members of an EQUIVALENCE group should have the **VOLATILE attribute** if any member has the  
 17 **VOLATILE attribute**.
- 18 4 If an object has the **VOLATILE attribute**, then all of its subobjects also have the **VOLATILE attribute**.
- 19 5 The Fortran processor should use the most recent definition of a volatile object each time its value is required.  
 20 When a volatile object is defined by means of Fortran, it should make that definition available to the non-Fortran  
 21 parts of the program as soon as possible.

## 22 5.6 Attribute specification statements

### 23 5.6.1 Accessibility statement

- 24 R527 *access-stmt* is *access-spec* [ [ :: ] *access-id-list* ]
- 25 R528 *access-id* is *access-name*  
 26 or *generic-spec*
- 27 C569 (R527) An *access-stmt* shall appear only in the *specification-part* of a module. Only one accessibility  
 28 statement with an omitted *access-id-list* is permitted in the *specification-part* of a module.
- 29 C570 (R528) Each *access-name* shall be the name of a module, variable, procedure, derived type, **named**  
 30 **constant**, or namelist group.
- 31 C571 A module whose name appears in an *access-stmt* shall be referenced by a **USE statement** in the **scoping**  
 32 **unit** that contains the *access-stmt*.
- 33 C572 The name of a module shall appear at most once in all of the *access-stmts* in a module.
- 34 1 An *access-stmt* with an *access-id-list* specifies the **accessibility attribute**, **PUBLIC** or **PRIVATE**, of each *access-id*  
 35 in the list that is not a module name. An *access-stmt* without an *access-id* list specifies the default accessibility  
 36 of the identifiers of entities declared in the module, and of entities accessed from a module whose name does  
 37 not appear in any *access-stmt* in the module. If an identifier is accessed from another module and also declared  
 38 locally, it has the default accessibility of a locally declared identifier. The statement  
 39 **PUBLIC**  
 40 specifies a default of public accessibility. The statement  
 41 **PRIVATE**  
 42 specifies a default of private accessibility. If no such statement appears in a module, the default is public  
 43 accessibility.

- 1 2 If an identifier is accessed by [use association](#) and not declared in the module, and the name of every module  
 2 from which it is accessed appears in an [access-stmt](#) in the scoping unit, its default accessibility is PRIVATE if  
 3 the [access-spec](#) in every such [access-stmt](#) is PRIVATE, or PUBLIC if the [access-spec](#) in any such [access-stmt](#) is  
 4 PUBLIC.

**NOTE 5.25**

Examples of accessibility statements are:

```
MODULE EX
  PRIVATE
  PUBLIC :: A, B, C, ASSIGNMENT (=), OPERATOR (+)
```

**NOTE 5.26**

The following is an example of using an accessibility statement on a module name.

```
MODULE m2
  USE m1
  ! We want to use the types and procedures in m1, but we only want to
  ! re-export m_type and our own procedures.
  PRIVATE m1
  PUBLIC m_type
  ... definitions for our own entities and module procedures.
END MODULE
```

## 5 5.6.2 ALLOCATABLE statement

6 R529 *allocatable-stmt* is ALLOCATABLE [ :: ] [allocatable-decl-list](#)

7 R530 *allocatable-decl* is [object-name](#) [ ( [array-spec](#) ) ] ■  
 8 ■ [ [lbracket coarray-spec rbracket](#) ]

- 9 1 The ALLOCATABLE statement specifies the [ALLOCATABLE attribute](#) (5.5.3) for a list of objects.

**NOTE 5.27**

An example of an ALLOCATABLE statement is:

```
REAL A, B (:), SCALAR
ALLOCATABLE :: A (:, :), B, SCALAR
```

## 10 5.6.3 ASYNCHRONOUS statement

11 R531 *asynchronous-stmt* is ASYNCHRONOUS [ :: ] [object-name-list](#)

- 12 1 The ASYNCHRONOUS statement specifies the [ASYNCHRONOUS attribute](#) (5.5.4) for a list of objects.

## 13 5.6.4 BIND statement

14 R532 *bind-stmt* is [language-binding-spec](#) [ :: ] [bind-entity-list](#)

15 R533 *bind-entity* is [entity-name](#)  
 16 or / [common-block-name](#) /

- 17 C573 (R532) If the [language-binding-spec](#) has a NAME= specifier, the [bind-entity-list](#) shall consist of a single  
 18 [bind-entity](#).

- 19 1 The BIND statement specifies the [BIND attribute](#) for a list of variables and [common blocks](#).

## 5.6.5 CODIMENSION statement

R534 *codimension-stmt* is CODIMENSION [ :: ] *codimension-decl-list*

R535 *codimension-decl* is *coarray-name* *lbracket coarray-spec rbracket*

- 1 The CODIMENSION statement specifies the CODIMENSION attribute (5.5.6) for a list of objects.

### NOTE 5.28

An example of a CODIMENSION statement is:

```
CODIMENSION a[*], b[3,*], c[:]
```

## 5.6.6 CONTIGUOUS statement

R536 *contiguous-stmt* is CONTIGUOUS [ :: ] *object-name-list*

- 1 The CONTIGUOUS statement specifies the CONTIGUOUS attribute (5.5.7) for a list of objects.

## 5.6.7 DATA statement

R537 *data-stmt* is DATA *data-stmt-set* [ [ , ] *data-stmt-set* ] ...

- 1 The DATA statement specifies explicit initialization (5.4).

- 2 If a nonpointer variable has default initialization, it shall not appear in a *data-stmt-object-list*.

- 3 A variable that appears in a DATA statement and has not been typed previously may appear in a subsequent type declaration only if that declaration confirms the implicit typing. An array name, array section, or array element that appears in a DATA statement shall have had its array properties established by a previous specification statement.

- 4 Except for variables in named common blocks, a named variable has the SAVE attribute if any part of it is initialized in a DATA statement, and this may be confirmed by explicit specification.

R538 *data-stmt-set* is *data-stmt-object-list* / *data-stmt-value-list* /

R539 *data-stmt-object* is *variable*  
or *data-implied-do*

R540 *data-implied-do* is ( *data-i-do-object-list* , [ *integer-type-spec* :: ] *data-i-do-variable* = ■  
■ *scalar-int-constant-expr* , ■  
■ *scalar-int-constant-expr* ■  
■ [ , *scalar-int-constant-expr* ] )

R541 *data-i-do-object* is *array-element*  
or *scalar-structure-component*  
or *data-implied-do*

R542 *data-i-do-variable* is *do-variable*

- C574 A *data-stmt-object* or *data-i-do-object* shall not be a coindexed variable.

- C575 (R539) A *data-stmt-object* that is a *variable* shall be a *designator*. Each subscript, section subscript, substring starting point, and substring ending point in the variable shall be a constant expression.

- C576 (R539) A variable whose *designator* appears as a *data-stmt-object* or a *data-i-do-object* shall not be a dummy argument, accessed by use or host association, in a named common block unless the DATA statement is in a block data program unit, in blank common, a function name, a function result name, an automatic data

- 1           object, or an *allocatable* variable.
- 2       C577   (R539) A *data-i-do-object* or a *variable* that appears as a *data-stmt-object* shall not be an *object designator*  
3       in which a pointer appears other than as the entire rightmost *part-ref*.
- 4       C578   (R541) The *array-element* shall be a variable.
- 5       C579   (R541) The *scalar-structure-component* shall be a variable.
- 6       C580   (R541) The *scalar-structure-component* shall contain at least one *part-ref* that contains a *subscript-list*.
- 7       C581   (R541) In an *array-element* or *scalar-structure-component* that is a *data-i-do-object*, any subscript shall  
8       be a *constant expression*, and any primary within that subscript that is a *data-i-do-variable* shall be a  
9       DO variable of this *data-implied-do* or of a containing *data-implied-do*.
- 10      R543   *data-stmt-value*                is   [ *data-stmt-repeat* \* ] *data-stmt-constant*
- 11      R544   *data-stmt-repeat*             is   *scalar-int-constant*  
12   or   *scalar-int-constant-subobject*
- 13      C582   (R544) The *data-stmt-repeat* shall be positive or zero. If the *data-stmt-repeat* is a *named constant*, it  
14      shall have been defined previously.
- 15      R545   *data-stmt-constant*           is   *scalar-constant*  
16   or   *scalar-constant-subobject*  
17   or   *signed-int-literal-constant*  
18   or   *signed-real-literal-constant*  
19   or   *null-init*  
20   or   *initial-data-target*  
21   or   *structure-constructor*
- 22      C583   (R545) If a DATA statement constant value is a *named constant* or a *structure constructor*, the *named*  
23      *constant* or derived type shall have been defined previously.
- 24      C584   (R545) If a *data-stmt-constant* is a *structure-constructor*, it shall be a *constant expression*.
- 25      R546   *int-constant-subobject*       is   *constant-subobject*
- 26      C585   (R546) *int-constant-subobject* shall be of type integer.
- 27      R547   *constant-subobject*           is   *designator*
- 28      C586   (R547) *constant-subobject* shall be a subobject of a constant.
- 29      C587   (R547) Any subscript, substring starting point, or substring ending point shall be a *constant expression*.
- 30   5   The *data-stmt-object-list* is expanded to form a sequence of pointers and scalar variables, referred to as “sequence  
31   of variables” in subsequent text. A nonpointer array whose unqualified name appears as a *data-stmt-object* or  
32   *data-i-do-object* is equivalent to a complete sequence of its array elements in array element order (6.5.3.2). An  
33   *array section* is equivalent to the sequence of its array elements in array element order. A *data-implied-do* is  
34   expanded to form a sequence of array elements and *structure components*, under the control of the *data-i-do-*  
35   *variable*, as in the DO construct (8.1.7.4).
- 36   6   The *data-stmt-value-list* is expanded to form a sequence of *data-stmt-constants*. A *data-stmt-repeat* indicates the  
37   number of times the following *data-stmt-constant* is to be included in the sequence; omission of a *data-stmt-repeat*  
38   has the effect of a repeat factor of 1.
- 39   7   A zero-sized array or a *data-implied-do* with an iteration count of zero contributes no variables to the expanded  
40   sequence of variables, but a zero-length scalar character variable does contribute a variable to the expanded  
41   sequence. A *data-stmt-constant* with a repeat factor of zero contributes no *data-stmt-constants* to the expanded

- 1 sequence of scalar *data-stmt-constants*.
- 2 8 The expanded sequences of variables and *data-stmt-constants* are in one-to-one correspondence. Each *data-stmt-constant* specifies the initial value, initial data *target*, or *null-init* for the corresponding variable. The lengths of
- 3 the two expanded sequences shall be the same.
- 4
- 5 9 A *data-stmt-constant* shall be *null-init* or *initial-data-target* if and only if the corresponding *data-stmt-object* has
- 6 the **POINTER** attribute. If *data-stmt-constant* is *null-init*, the initial association status of the corresponding data
- 7 statement object is *disassociated*. If *data-stmt-constant* is *initial-data-target* the corresponding data statement
- 8 object shall be data-pointer-initialization compatible (4.5.4.6) with the initial data *target*; the data statement
- 9 object is initially associated with the *target*.
- 10 10 A *data-stmt-constant* other than *boz-literal-constant*, *null-init*, or *initial-data-target* shall be compatible with its
- 11 corresponding variable according to the rules of intrinsic assignment (7.2.1.2). The variable is initially defined
- 12 with the value specified by the *data-stmt-constant*; if necessary, the value is converted according to the rules of
- 13 intrinsic assignment (7.2.1.3) to a value that agrees in type, type parameters, and shape with the variable.
- 14 11 If a *data-stmt-constant* is a *boz-literal-constant*, the corresponding variable shall be of type integer. The *boz-*
- 15 *literal-constant* is treated as if it were converted by the intrinsic function **INT** (13.9.100) to type integer with the
- 16 kind type parameter of the variable.

**NOTE 5.29**

Examples of DATA statements are:

```
CHARACTER (LEN = 10) NAME
INTEGER, DIMENSION (0:9) :: MILES
REAL, DIMENSION (100, 100) :: SKEW
TYPE (NODE), POINTER :: HEAD_OF_LIST
TYPE (PERSON) MYNAME, YOURNAME
DATA NAME / 'JOHN DOE' /, MILES / 10 * 0 /
DATA ((SKEW (K, J), J = 1, K), K = 1, 100) / 5050 * 0.0 /
DATA ((SKEW (K, J), J = K + 1, 100), K = 1, 99) / 4950 * 1.0 /
DATA HEAD_OF_LIST / NULL() /
DATA MYNAME / PERSON (21, 'JOHN SMITH') /
DATA YOURNAME % AGE, YOURNAME % NAME / 35, 'FRED BROWN' /
```

The character variable NAME is initialized with the value JOHN DOE with padding on the right because the length of the constant is less than the length of the variable. All ten elements of the integer array MILES are initialized to zero. The two-dimensional array SKEW is initialized so that the lower triangle of SKEW is zero and the strict upper triangle is one. The structures MYNAME and YOURNAME are declared using the derived type PERSON from Note 4.15. The pointer HEAD\_OF\_LIST is declared using the derived type NODE from Note 4.35; it is initially *disassociated*. MYNAME is initialized by a *structure constructor*. YOURNAME is initialized by supplying a separate value for each component.

## 17 5.6.8 DIMENSION statement

18 R548 *dimension-stmt* is DIMENSION [ :: ] *array-name* ( *array-spec* ) ■

19 ■ [ , *array-name* ( *array-spec* ) ] ...

- 20 1 The DIMENSION statement specifies the **DIMENSION** attribute (5.5.8) for a list of objects.

**NOTE 5.30**

An example of a DIMENSION statement is:

```
DIMENSION A (10), B (10, 70), C (:)
```

### 5.6.9 INTENT statement

R549 *intent-stmt* is INTENT ( *intent-spec* ) [ :: ] *dummy-arg-name-list*

- 1 The INTENT statement specifies the **INTENT attribute** (5.5.10) for the **dummy arguments** in the list.

#### NOTE 5.31

An example of an INTENT statement is:

```
SUBROUTINE EX (A, B)
  INTENT (INOUT) :: A, B
```

### 5.6.10 OPTIONAL statement

R550 *optional-stmt* is OPTIONAL [ :: ] *dummy-arg-name-list*

- 1 The OPTIONAL statement specifies the **OPTIONAL attribute** (5.5.12) for the **dummy arguments** in the list.

#### NOTE 5.32

An example of an OPTIONAL statement is:

```
SUBROUTINE EX (A, B)
  OPTIONAL :: B
```

### 5.6.11 PARAMETER statement

- 1 The PARAMETER statement specifies the **PARAMETER attribute** (5.5.13) and the values for the **named constants** in the list.

R551 *parameter-stmt* is PARAMETER ( *named-constant-def-list* )

R552 *named-constant-def* is *named-constant* = *constant-expr*

- 2 If a **named constant** is defined by a PARAMETER statement, it shall not be subsequently declared to have a type or type parameter value that differs from the type and type parameters it would have if declared implicitly (5.7). A named array constant defined by a PARAMETER statement shall have its **rank** specified in a prior specification statement.
- 3 The constant expression that corresponds to a named constant shall have type and **type parameters** that conform with the named constant as specified for intrinsic assignment (7.2.1.2). If the named constant has implied shape, the expression shall have the same rank as the named constant; otherwise, the expression shall either be scalar or have the same shape as the named constant.
- 4 The value of each **named constant** is that specified by the corresponding **constant expression**; if necessary, the value is converted according to the rules of intrinsic assignment (7.2.1.3) to a value that agrees in type, type parameters, and shape with the **named constant**.

#### NOTE 5.33

An example of a PARAMETER statement is:

```
PARAMETER (MODULUS = MOD (28, 3), NUMBER_OF_SENATORS = 100)
```

### 5.6.12 POINTER statement

R553 *pointer-stmt* is POINTER [ :: ] *pointer-decl-list*

R554 *pointer-decl* is *object-name* [ ( *deferred-shape-spec-list* ) ]

1 *or proc-entity-name*

2 C588 A *proc-entity-name* shall have the **EXTERNAL** attribute.

- 3 1 The **POINTER** statement specifies the **POINTER** attribute (5.5.14) for a list of entities.

#### NOTE 5.34

An example of a **POINTER** statement is:

```
TYPE (NODE) :: CURRENT
POINTER :: CURRENT, A (:, :)
```

### 5.6.13 PROTECTED statement

R555 *protected-stmt* is **PROTECTED** [ :: ] *entity-name-list*

- 6 1 The **PROTECTED** statement specifies the **PROTECTED** attribute (5.5.15) for a list of entities.

### 5.6.14 SAVE statement

R556 *save-stmt* is **SAVE** [ [ :: ] *saved-entity-list* ]

R557 *saved-entity* is *object-name*  
or *proc-pointer-name*  
or / *common-block-name* /

R558 *proc-pointer-name* is *name*

C589 (R556) If a **SAVE** statement with an omitted saved entity list appears in a **scoping unit**, no other appearance of the **SAVE** *attr-spec* or **SAVE** statement is permitted in that **scoping unit**.

C590 A *proc-pointer-name* shall be the name of a procedure pointer.

- 16 1 A **SAVE** statement with a saved entity list specifies the **SAVE** attribute (5.5.16) for a list of entities. A **SAVE**  
17 statement without a saved entity list is treated as though it contained the names of all allowed items in the same  
18 **scoping unit**.

#### NOTE 5.35

An example of a **SAVE** statement is:

```
SAVE A, B, C, / BLOCKA /, D
```

### 5.6.15 TARGET statement

R559 *target-stmt* is **TARGET** [ :: ] *target-decl-list*

R560 *target-decl* is *object-name* [ ( *array-spec* ) ] ■  
■ [ *lbracket coarray-spec rbracket* ]

- 23 1 The **TARGET** statement specifies the **TARGET** attribute (5.5.17) for a list of objects.

#### NOTE 5.36

An example of a **TARGET** statement is:

```
TARGET :: A (1000, 1000), B
```

## 5.6.16 VALUE statement

R561 *value-stmt* is VALUE [ :: ] *dummy-arg-name-list*

- 1 The VALUE statement specifies the VALUE attribute (5.5.18) for a list of dummy arguments.

## 5.6.17 VOLATILE statement

R562 *volatile-stmt* is VOLATILE [ :: ] *object-name-list*

- 1 The VOLATILE statement specifies the VOLATILE attribute (5.5.19) for a list of objects.

## 5.7 IMPLICIT statement

- 1 In a *scoping unit*, an IMPLICIT statement specifies a type, and possibly type parameters, for all implicitly typed data entities whose names begin with one of the letters specified in the statement. An IMPLICIT NONE statement can indicate that no implicit typing rules are to apply in a particular *scoping unit*, or that *external* and *dummy* procedures need to be explicitly given the EXTERNAL attribute.

R563 *implicit-stmt* is IMPLICIT *implicit-spec-list*  
or IMPLICIT NONE [ ( [ *implicit-none-spec-list* ] ) ]

R564 *implicit-spec* is *declaration-type-spec* ( *letter-spec-list* )

R565 *letter-spec* is *letter* [ – *letter* ]

R566 *implicit-none-spec* is EXTERNAL  
or TYPE

- C591 (R563) If an IMPLICIT NONE statement appears in a *scoping unit*, it shall precede any *PARAMETER statements* that appear in the *scoping unit*. No more than one IMPLICIT NONE statement shall appear in a *scoping unit*.

- C592 If an IMPLICIT NONE statement in a *scoping unit* has an *implicit-none-spec* of TYPE or has no *implicit-none-spec-list*, there shall be no other IMPLICIT statements in the *scoping unit*.

- C593 (R565) If the minus and second *letter* appear, the second letter shall follow the first letter alphabetically.

- C594 If IMPLICIT NONE with an *implicit-none-spec* of EXTERNAL appears within a *scoping unit*, the name of an *external* or *dummy* procedure in that *scoping unit* or in a contained *subprogram* or *BLOCK construct* shall be explicitly declared to have the EXTERNAL attribute.

- 2 A *letter-spec* consisting of two *letters* separated by a minus is equivalent to writing a list containing all of the letters in alphabetical order in the alphabetic sequence from the first letter through the second letter. For example, A–C is equivalent to A, B, C. The same letter shall not appear as a single letter, or be included in a range of letters, more than once in all of the IMPLICIT statements in a *scoping unit*.

- 3 In each *scoping unit*, there is a mapping, which may be null, between each of the letters A, B, ..., Z and a type (and type parameters). An IMPLICIT statement specifies the mapping for the letters in its *letter-spec-list*. IMPLICIT NONE with an *implicit-none-spec* of TYPE or with no *implicit-none-spec-list* specifies the null mapping for all the letters. If a mapping is not specified for a letter, the default for a *program unit* or an *interface body* is default integer if the letter is I, J, ..., or N and default real otherwise, and the default for a *BLOCK construct*, *internal subprogram*, or *module subprogram* is the mapping in the *host scoping unit*.

- 4 Any data entity that is not explicitly declared by a *type declaration statement*, is not an intrinsic function, is not a *component*, and is not accessed by *use* or *host* association is declared implicitly to be of the type (and type parameters) mapped from the first letter of its name, provided the mapping is not null. The mapping for the first letter of the data entity shall either have been established by a prior IMPLICIT statement or be the

1 default mapping for the letter. The data entity is treated as if it were declared in an explicit type declaration; if  
 2 the outermost **inclusive scope** in which it appears is not a type definition, it is declared in that scope, otherwise  
 3 it is declared in the host of that scope. An explicit type specification in a **FUNCTION statement** overrides an  
 4 **IMPLICIT** statement for the **result** of that function.

**NOTE 5.37**

The following are examples of the use of **IMPLICIT** statements:

```

MODULE EXAMPLE_MODULE
  IMPLICIT NONE
  ...
  INTERFACE
    FUNCTION FUN (I)      ! Not all data entities need to
      INTEGER FUN          ! be declared explicitly
    END FUNCTION FUN
  END INTERFACE
CONTAINS
  FUNCTION JFUN (J)       ! All data entities need to
    INTEGER JFUN, J       ! be declared explicitly.
    ...
  END FUNCTION JFUN
END MODULE EXAMPLE_MODULE
SUBROUTINE SUB
  IMPLICIT COMPLEX (C)
  C = (3.0, 2.0)          ! C is implicitly declared COMPLEX
  ...
CONTAINS
  SUBROUTINE SUB1
    IMPLICIT INTEGER (A, C)
    C = (0.0, 0.0)        ! C is host associated and of
                          ! type complex
    Z = 1.0               ! Z is implicitly declared REAL
    A = 2                 ! A is implicitly declared INTEGER
    CC = 1                ! CC is implicitly declared INTEGER
    ...
  END SUBROUTINE SUB1
  SUBROUTINE SUB2
    Z = 2.0               ! Z is implicitly declared REAL and
                          ! is different from the variable of
                          ! the same name in SUB1
    ...
  END SUBROUTINE SUB2
  SUBROUTINE SUB3
    USE EXAMPLE_MODULE    ! Accesses integer function FUN
                          ! by use association
    Q = FUN (K)            ! Q is implicitly declared REAL and
    ...                   ! K is implicitly declared INTEGER
  END SUBROUTINE SUB3
END SUBROUTINE SUB

```

**NOTE 5.38**

The following is an example of a mapping to a derived type that is inaccessible in the local scope:

```

PROGRAM MAIN
  IMPLICIT TYPE(BLOB) (A)

```

**NOTE 5.38 (cont.)**

```

TYPE BLOB
  INTEGER :: I
END TYPE BLOB
TYPE(BLOB) :: B
CALL STEVE
CONTAINS
  SUBROUTINE STEVE
    INTEGER :: BLOB
    ..
    AA = B
    ..
  END SUBROUTINE STEVE
END PROGRAM MAIN

```

In the subroutine STEVE, it is not possible to explicitly declare a variable to be of type BLOB because BLOB has been given a different meaning, but implicit mapping for the letter A still maps to type BLOB, so AA is of type BLOB.

**NOTE 5.39**

Implicit typing is not affected by [BLOCK constructs](#). For example, in

```

SUBROUTINE S(N)
  ...
  IF (N>0) THEN
    BLOCK
      NSQP = CEILING(SQRT(DBLE(N)))
    END BLOCK
  END IF
  ...
  IF (N>0) THEN
    BLOCK
      PRINT *,NSQP
    END BLOCK
  END IF
END SUBROUTINE

```

even if the only two appearances of NSQP are within the [BLOCK constructs](#), the scope of NSQP is the whole subroutine S.

**NOTE 5.40**

In the subprogram

```

SUBROUTINE EXAMPLE (X, Y)
  IMPLICIT NONE (EXTERNAL)
  REAL, EXTERNAL :: G
  REAL :: X, Y
  X = F (Y)           ! Invalid: F lacks the EXTERNAL attribute.
  X = G (Y)           ! Valid: G has the EXTERNAL attribute.
END SUBROUTINE

```

the referenced function F needs to have the [EXTERNAL attribute \(5.5.9\)](#).

## 5.8 NAMELIST statement

- 1 A NAMELIST statement specifies a group of named data objects, which may be referred to by a single name for the purpose of data transfer (9.6, 10.11).

R567 *namelist-stmt* is NAMELIST ■  
 ■ / *namelist-group-name* / *namelist-group-object-list* ■  
 ■ [ [ , ] / *namelist-group-name* / ■  
 ■ *namelist-group-object-list* ] ...

C595 (R567) The *namelist-group-name* shall not be a name accessed by use association.

R568 *namelist-group-object* is *variable-name*

C596 (R568) A *namelist-group-object* shall not be an assumed-size array.

C597 (R567) A *namelist-group-object* shall not have the PRIVATE attribute if the *namelist-group-name* has the PUBLIC attribute.

- 2 The order in which the variables are specified in the NAMELIST statement determines the order in which the values appear on output.

- 3 Any *namelist-group-name* may occur more than once in the NAMELIST statements in a scoping unit. The *namelist-group-object-list* following each successive appearance of the same *namelist-group-name* in a scoping unit is treated as a continuation of the list for that *namelist-group-name*.

- 4 A namelist group object may be a member of more than one namelist group.

- 5 A namelist group object shall either be accessed by use or host association or shall have its declared type, kind type parameters of the declared type, and rank specified by previous specification statements or the procedure heading in the same scoping unit or by the implicit typing rules in effect for the scoping unit. If a namelist group object is typed by the implicit typing rules, its appearance in any subsequent type declaration statement shall confirm the implied type and type parameters.

### NOTE 5.41

An example of a NAMELIST statement is:

```
NAMELIST /NLIST/ A, B, C
```

## 5.9 Storage association of data objects

### 5.9.1 EQUIVALENCE statement

#### 5.9.1.1 General

- 1 An EQUIVALENCE statement is used to specify the sharing of storage units by two or more objects in a scoping unit. This causes storage association (16.5.3) of the objects that share the storage units.

- 2 If the equivalenced objects have differing type or type parameters, the EQUIVALENCE statement does not cause type conversion or imply mathematical equivalence. If a scalar and an array are equivalenced, the scalar does not have array properties and the array does not have the properties of a scalar.

R569 *equivalence-stmt* is EQUIVALENCE *equivalence-set-list*

R570 *equivalence-set* is ( *equivalence-object* , *equivalence-object-list* )

R571 *equivalence-object* is *variable-name*  
 or *array-element*  
 or *substring*

C598 (R571) An *equivalence-object* shall not be a designator with a base object that is a dummy argument, a function result,

a pointer, an [allocatable](#) variable, a derived-type object that has an [allocatable](#) or pointer [ultimate component](#), an object of a nonsequence derived type, an [automatic data object](#), a [coarray](#), a variable with the [BIND attribute](#), a variable in a [common block](#) that has the [BIND attribute](#), or a [named constant](#).

C599 (R571) An [equivalence-object](#) shall not be a [designator](#) that has more than one [part-ref](#).

C5100 (R571) An [equivalence-object](#) shall not have the [TARGET attribute](#).

C5101 (R571) Each subscript or substring range expression in an [equivalence-object](#) shall be an integer [constant expression](#) (7.1.12).

C5102 (R570) If an [equivalence-object](#) is default integer, default real, double precision real, default complex, default logical, or of [numeric sequence type](#), all of the objects in the equivalence set shall be of these types and kinds.

C5103 (R570) If an [equivalence-object](#) is default character or of [character sequence type](#), all of the objects in the equivalence set shall be of these types and kinds.

C5104 (R570) If an [equivalence-object](#) is of a [sequence type](#) that is not a [numeric sequence](#) or [character sequence](#) type, all of the objects in the equivalence set shall be of the same type with the same type parameter values.

C5105 (R570) If an [equivalence-object](#) is of an intrinsic type but is not default integer, default real, double precision real, default complex, default logical, or default character, all of the objects in the equivalence set shall be of the same type with the same kind type parameter value.

C5106 (R571) If an [equivalence-object](#) has the [PROTECTED attribute](#), all of the objects in the equivalence set shall have the [PROTECTED attribute](#).

C5107 (R571) The name of an [equivalence-object](#) shall not be a name made accessible by [use association](#).

C5108 (R571) A [substring](#) shall not have length zero.

#### NOTE 5.42

The EQUIVALENCE statement allows the equivalencing of sequence structures and the equivalencing of objects of intrinsic type with nondefault type parameters, but there are strict rules regarding the appearance of these objects in an EQUIVALENCE statement.

A structure that appears in an EQUIVALENCE statement shall be a sequence structure. If a sequence structure is not of [numeric sequence type](#) or of [character sequence type](#), it shall be equivalenced only to objects of the same type with the same type parameter values.

A structure of a [numeric sequence type](#) shall be equivalenced only to another structure of a [numeric sequence type](#), an object that is default integer, default real, double precision real, default complex, or default logical type such that components of the structure ultimately become associated only with objects of these types and kinds.

A structure of a [character sequence type](#) shall be equivalenced only to a default character object or another structure of a [character sequence type](#).

An object of intrinsic type with nondefault kind type parameters shall not be equivalenced to objects of different type or kind type parameters.

Further rules on the interaction of EQUIVALENCE statements and [default initialization](#) are given in 16.5.3.4.

### 5.9.1.2 Equivalence association

1 An EQUIVALENCE statement specifies that the [storage sequences](#) (16.5.3.2) of the data objects specified in an [equivalence-set](#) are storage associated. All of the nonzero-sized sequences in the [equivalence-set](#), if any, have the same first [storage unit](#), and all of the zero-sized sequences in the [equivalence-set](#), if any, are storage associated with one another and with the first [storage unit](#) of any nonzero-sized sequences. This causes the [storage association](#) of the data objects in the [equivalence-set](#) and may cause [storage association](#) of other data objects.

2 If any data object in an [equivalence-set](#) has the [SAVE attribute](#), all other objects in the [equivalence-set](#) have the [SAVE attribute](#); this may be confirmed by explicit specification.

### 5.9.1.3 Equivalence of default character objects

1 A default character data object shall not be equivalenced to an object that is not default character and not of a [character sequence type](#). The lengths of equivalenced default character objects need not be the same.

2 An EQUIVALENCE statement specifies that the [storage sequences](#) of all the default character data objects specified in an [equivalence-set](#) are storage associated. All of the nonzero-sized sequences in the [equivalence-set](#), if any, have the same first [character storage unit](#),

and all of the zero-sized sequences in the *equivalence-set*, if any, are storage associated with one another and with the first *character storage unit* of any nonzero-sized sequences. This causes the *storage association* of the data objects in the *equivalence-set* and may cause *storage association* of other data objects.

#### NOTE 5.43

For example, using the declarations:

```
CHARACTER (LEN = 4) :: A, B
CHARACTER (LEN = 3) :: C (2)
EQUIVALENCE (A, C (1)), (B, C (2))
```

the association of A, B, and C can be illustrated graphically as:

1	2	3	4	5	6	7
---	--- A ---	---	---			
			---	--- B ---	---	---
---	C(1) ---	---	---	C(2) ---	---	

#### 5.9.1.4 Array names and array element designators

- For a nonzero-sized array, the use of the array name unqualified by a subscript list as an *equivalence-object* has the same effect as using an array element *designator* that identifies the first element of the array.

#### 5.9.1.5 Restrictions on EQUIVALENCE statements

- An EQUIVALENCE statement shall not specify that the same *storage unit* is to occur more than once in a *storage sequence*.
- An EQUIVALENCE statement shall not specify that consecutive *storage units* are to be nonconsecutive.

### 5.9.2 COMMON statement

#### 5.9.2.1 General

- The COMMON statement specifies blocks of physical storage, called *common blocks*, that can be accessed by any of the *scoping units* in a program. Thus, the COMMON statement provides a global data facility based on *storage association* (16.5.3).
- A *common block* that does not have a name is called *blank common*.

R572 *common-stmt* is COMMON ■  
 ■ [ / [ *common-block-name* ] / ] *common-block-object-list* ■  
 ■ [ [ , ] / [ *common-block-name* ] / ■  
 ■ *common-block-object-list* ] ...

R573 *common-block-object* is *variable-name* [ ( *array-spec* ) ]

C5109 (R573) An *array-spec* in a *common-block-object* shall be an *explicit-shape-spec-list*.

C5110 (R573) Only one appearance of a given *variable-name* is permitted in all *common-block-object-lists* within a *scoping unit*.

C5111 (R573) A *common-block-object* shall not be a *dummy argument*, a *function result*, an *allocatable* variable, a derived-type object with an *ultimate component* that is *allocatable*, a *procedure pointer*, an *automatic data object*, a variable with the *BIND* attribute, an *unlimited polymorphic* pointer, or a *coarray*.

C5112 (R573) If a *common-block-object* is of a derived type, the type shall have the *BIND* attribute or the *SEQUENCE* attribute and it shall have no *default initialization*.

C5113 (R573) A *variable-name* shall not be a name made accessible by *use association*.

- In each COMMON statement, the data objects whose names appear in a common block object list following a *common block* name are declared to be in that *common block*. If the first *common block* name is omitted, all data objects whose names appear in the first common block object list are specified to be in *blank common*. Alternatively, the appearance of two slashes with no *common block* name between them declares the data objects whose names appear in the common block object list that follows to be in *blank common*.
- Any *common block* name or an omitted *common block* name for *blank common* may occur more than once in one or more COMMON statements in a *scoping unit*. The common block list following each successive appearance of the same common block name in a *scoping unit* is treated as a continuation of the list for that common block name. Similarly, each blank common block object list in a *scoping unit* is treated as a continuation of *blank common*.

- 1 5 The form *variable-name* (*array-spec*) specifies the **DIMENSION attribute** for that variable.
- 2 6 If derived-type objects of **numeric sequence type** or **character sequence type** (4.5.2.3) appear in **common**, it is as if the individual
- 3 components were enumerated directly in the common list.

#### 4 5.9.2.2 Common block storage sequence

- 5 1 For each **common block** in a **scoping unit**, a common block **storage sequence** is formed as follows:
- 6 (1) A **storage sequence** is formed consisting of the sequence of **storage units** in the **storage sequences** (16.5.3.2) of all data
- 7 objects in the common block object lists for the **common block**. The order of the **storage sequences** is the same as the
- 8 order of the appearance of the common block object lists in the **scoping unit**.
- 9 (2) The **storage sequence** formed in (1) is extended to include all **storage units** of any **storage sequence** associated with it
- 10 by equivalence association. The **sequence** shall be extended only by adding **storage units** beyond the last **storage unit**.
- 11 Data objects associated with an entity in a **common block** are considered to be in that **common block**.
- 12 2 Only COMMON statements and **EQUIVALENCE** statements appearing in the **scoping unit** contribute to common block **storage**
- 13 **sequences** formed in that **scoping unit**.

#### 14 5.9.2.3 Size of a common block

- 15 1 The size of a common block is the size of its common block **storage sequence**, including any extensions of the **sequence** resulting from
- 16 equivalence association.

#### 17 5.9.2.4 Common association

- 18 1 Within a program, the common block **storage sequences** of all nonzero-sized **common blocks** with the same name have the same first
- 19 **storage unit**, and the common block **storage sequences** of all zero-sized **common blocks** with the same name are storage associated
- 20 with one another. Within a program, the common block **storage sequences** of all nonzero-sized **blank common** blocks have the same
- 21 first **storage unit** and the **storage sequences** of all zero-sized **blank common** blocks are associated with one another and with the first
- 22 **storage unit** of any nonzero-sized **blank common** blocks. This results in the association of objects in different **scoping units**. Use or
- 23 **host** association may cause these associated objects to be accessible in the same **scoping unit**.
- 24 2 A nonpointer object that is default integer, default real, double precision real, default complex, default logical, or of **numeric sequence**
- 25 **type** shall be associated only with nonpointer objects of these types and kinds.
- 26 3 A nonpointer object that is default character or of **character sequence type** shall be associated only with nonpointer objects of these
- 27 types and kinds.
- 28 4 A nonpointer object of a derived type that is not a **numeric sequence** or **character sequence type** shall be associated only with
- 29 nonpointer objects of the same type with the same type parameter values.
- 30 5 A nonpointer object of intrinsic type but which is not default integer, default real, double precision real, default complex, default
- 31 logical, or default character shall be associated only with nonpointer objects of the same type and type parameters.
- 32 6 A data pointer shall be storage associated only with data pointers of the same type and **rank**. Data pointers that are storage
- 33 associated shall have **deferred** the same type parameters; corresponding nondeferred **type parameters** shall have the same value.
- 34 7 An object with the **TARGET attribute** shall be storage associated only with another object that has the **TARGET attribute** and the
- 35 same type and type parameters.

##### NOTE 5.44

A **common block** is permitted to contain sequences of different **storage units**, provided each **scoping unit** that accesses the **common block** specifies an identical sequence of **storage units** for the **common block**. For example, this allows a single **common block** to contain both **numeric** and **character storage units**.

Association in different **scoping units** between objects of default type, objects of double precision real type, and sequence structures is permitted according to the rules for equivalence objects (5.9.1).

#### 36 5.9.2.5 Differences between named common and blank common

- 37 1 A **blank common** block has the same properties as a named **common block**, except for the following.
- 38 • Execution of a **RETURN** or **END** statement might cause data objects in a named **common block** to become undefined unless
- 39 the **common block** has the **SAVE attribute**, but never causes data objects in **blank common** to become undefined (16.6.6).
- 40 • Named **common blocks** of the same name shall be of the same size in all **scoping units** of a program in which they appear, but
- 41 **blank common** blocks may be of different sizes.
- 42 • A data object in a named **common block** may be initially defined by means of a **DATA statement** or **type declaration statement**
- 43 in a **block data program unit** (11.3), but objects in **blank common** shall not be initially defined.

### 5.9.3 Restrictions on common and equivalence

- 1
  - 2
  - 3
  - 4
  - 5
  - 6
- 1 An **EQUIVALENCE statement** shall not cause the **storage sequences** of two different **common blocks** to be associated.
  - 2 Equivalence association shall not cause a derived-type object with **default initialization** to be associated with an object in a **common block**.
  - 3 Equivalence association shall not cause a common block **storage sequence** to be extended by adding **storage units** preceding the first **storage unit** of the first object specified in a COMMON statement for the **common block**.

1 (Blank page)

## 6 Use of data objects

### 6.1 Designator

R601 *designator* is *object-name*  
or *array-element*  
or *array-section*  
or *coindexed-named-object*  
or *complex-part-designator*  
or *structure-component*  
or *substring*

- 1 The appearance of a *data object designator* in a context that requires its value is termed a reference.

### 6.2 Variable

R602 *variable* is *designator*  
or *function-reference*

C601 (R602) *designator* shall not be a constant or a subobject of a constant.

C602 (R602) *function-reference* shall have a data pointer result.

- 1 A variable is either the data object denoted by *designator* or the *target* of the pointer resulting from the evaluation of *function-reference*; this pointer shall be associated.
- 2 A reference is permitted only if the variable is defined. A reference to a data pointer is permitted only if the pointer is associated with a *target* object that is defined. A variable becomes defined with a value when events described in 16.6.5 occur.

R603 *variable-name* is *name*

C603 (R603) *variable-name* shall be the name of a variable.

R604 *logical-variable* is *variable*

C604 (R604) *logical-variable* shall be of type logical.

R605 *char-variable* is *variable*

C605 (R605) *char-variable* shall be of type character.

R606 *default-char-variable* is *variable*

C606 (R606) *default-char-variable* shall be default character.

R607 *int-variable* is *variable*

C607 (R607) *int-variable* shall be of type integer.

#### NOTE 6.1

For example, given the declarations:

```
CHARACTER (10)  A, B (10)
```

**NOTE 6.1 (cont.)**

TYPE (PERSON) P ! See Note 4.15

then A, B, B (1), B (1:5), P % AGE, and A (1:1) are all variables.

**6.3 Constants**

- 1 A constant (3.2.3) is a literal constant or a [named constant](#). A literal constant is a scalar denoted by a syntactic form, which indicates its type, type parameters, and value. A [named constant](#) is a constant that has a name; the name has the [PARAMETER attribute](#) (5.5.13, 5.6.11). A reference to a constant is always permitted; redefinition of a constant is never permitted.

**6.4 Scalars****6.4.1 Substrings**

- 1 A substring is a [contiguous](#) portion of a character string (4.4.4).

R608 *substring* is *parent-string* ( *substring-range* )

R609 *parent-string* is *scalar-variable-name*  
or *array-element*  
or *coindexed-named-object*  
or *scalar-structure-component*  
or *scalar-constant*

R610 *substring-range* is [ *scalar-int-expr* ] : [ *scalar-int-expr* ]

C608 (R609) *parent-string* shall be of type character.

- 2 The value of the first *scalar-int-expr* in *substring-range* is the starting point of the substring and the value of the second one is the ending point of the substring. The length of a substring is the number of characters in the substring and is MAX ( $l - f + 1, 0$ ), where  $f$  and  $l$  are the starting and ending points, respectively.
- 3 Let the characters in the parent string be numbered 1, 2, 3, ...,  $n$ , where  $n$  is the length of the parent string. Then the characters in the substring are those from the parent string from the starting point and proceeding in sequence up to and including the ending point. If the starting point is greater than the ending point, the substring has length zero; otherwise, both the starting point and the ending point shall be within the range 1, 2, ...,  $n$ . If the starting point is not specified, the default value is 1. If the ending point is not specified, the default value is  $n$ .

**NOTE 6.2**

Examples of character substrings are:

B(1)(1:5)	array element as parent string
P%NAME(1:1)	structure component as parent string
ID(4:9)	scalar variable name as parent string
'0123456789'(N:N)	character constant as parent string

**6.4.2 Structure components**

- 1 A [structure component](#) is part of an object of derived type; it may be referenced by an [object designator](#). A [structure component](#) may be a scalar or an array.

R611 *data-ref* is *part-ref* [ % *part-ref* ] ...

- 1 R612 *part-ref* is *part-name* [ ( *section-subscript-list* ) ] [ *image-selector* ]
- 2 C609 (R611) Each *part-name* except the rightmost shall be of derived type.
- 3 C610 (R611) Each *part-name* except the leftmost shall be the name of a component of the *declared type* of the  
4 preceding *part-name*.
- 5 C611 (R611) If the rightmost *part-name* is of *abstract type*, *data-ref* shall be *polymorphic*.
- 6 C612 (R611) The leftmost *part-name* shall be the name of a data object.
- 7 C613 (R612) If a *section-subscript-list* appears, the number of *section-subscripts* shall equal the *rank* of *part-*  
8 *name*.
- 9 C614 (R612) If *image-selector* appears, the number of *cosubscripts* shall be equal to the *corank* of *part-name*.
- 10 C615 (R612) If *image-selector* appears and *part-name* is an array, *section-subscript-list* shall appear.
- 11 C616 (R611) If *image-selector* appears, *data-ref* shall not be of type *C\_PTR* or *C\_FUNPTR* from the intrinsic  
12 module *ISO\_C\_BINDING* (15.3.3).
- 13 C617 (R611) Except as an *actual argument* to an intrinsic *inquiry function* or as the *designator* in a type  
14 parameter inquiry, a *data-ref* shall not be a *polymorphic* subobject of a *coindexed object* and shall not  
15 be a *coindexed object* that has a *polymorphic allocatable subcomponent*.
- 16 2 The *rank* of a *part-ref* of the form *part-name* is the *rank* of *part-name*. The *rank* of a *part-ref* that has a section  
17 subscript list is the number of subscript triplets and *vector subscripts* in the list.
- 18 C618 (R611) There shall not be more than one *part-ref* with nonzero *rank*. A *part-name* to the right of a  
19 *part-ref* with nonzero *rank* shall not have the *ALLOCATABLE* or *POINTER* attribute.
- 20 3 The *rank* of a *data-ref* is the *rank* of the *part-ref* with nonzero *rank*, if any; otherwise, the *rank* is zero. The *base*  
21 *object* of a *data-ref* is the data object whose name is the leftmost part name.
- 22 4 The type and type parameters, if any, of a *data-ref* are those of the rightmost part name.
- 23 5 A *data-ref* with more than one *part-ref* is a subobject of its *base object* if none of the *part-names*, except for  
24 possibly the rightmost, is a pointer. If the rightmost *part-name* is the only pointer, then the *data-ref* is a subobject  
25 of its *base object* in contexts that pertain to its *pointer association* status but not in any other contexts.

**NOTE 6.3**

If X is an object of derived type with a pointer component P, then the pointer X%P is a subobject of X when considered as a pointer – that is in contexts where it is not dereferenced.

However the *target* of X%P is not a subobject of X. Thus, in contexts where X%P is dereferenced to refer to the *target*, it is not a subobject of X.

- 26 R613 *structure-component* is *data-ref*
- 27 C619 (R613) There shall be more than one *part-ref* and the rightmost *part-ref* shall not have a *section-subscript-*  
28 *list*.
- 29 6 A *structure component* shall be neither referenced nor defined before the declaration of the *base object*. A  
30 *structure component* is a pointer only if the rightmost part name is defined to have the *POINTER* attribute.

**NOTE 6.4**

Examples of structure components are:

SCALAR\_PARENT%SCALAR\_FIELD                      scalar component of scalar parent

**NOTE 6.4 (cont.)**

ARRAY_PARENT(J)%SCALAR_FIELD	component of array element parent
ARRAY_PARENT(1:N)%SCALAR_FIELD	component of array section parent

For a more elaborate example see C.3.1.

**NOTE 6.5**

The syntax rules are structured such that a *data-ref* that ends in a component name without a following subscript list is a structure component, even when other component names in the *data-ref* are followed by a subscript list. A *data-ref* that ends in a component name with a following subscript list is either an array element or an *array section*. A *data-ref* of nonzero *rank* that ends with a *substring-range* is an *array section*. A *data-ref* of zero *rank* that ends with a *substring-range* is a substring.

**6.4.3 Coindexed named objects**

- 1 A *coindexed-named-object* is a named scalar coarray variable followed by an image selector.

R614 *coindexed-named-object* is *data-ref*

C620 (R614) The *data-ref* shall contain exactly one *part-ref*. The *part-ref* shall contain an *image-selector*. The *part-name* shall be the name of a scalar coarray.

**6.4.4 Complex parts**

R615 *complex-part-designator* is *designator* % RE  
or *designator* % IM

C621 (R615) The *designator* shall be of complex type.

- 1 If *complex-part-designator* is *designator*%RE it designates the real part of *designator*. If it is *designator*%IM it designates the imaginary part of *designator*. The type of a *complex-part-designator* is real, and its kind and shape are those of the *designator*.

**NOTE 6.6**

The following are examples of *complex part designators*:

impedance%re	!-- Same value as REAL(impedance)
fft%im	!-- Same value as AIMAG(fft)
x%im = 0.0	!-- Sets the imaginary part of X to zero

**6.4.5 Type parameter inquiry**

- 1 A *type parameter inquiry* is used to inquire about a type parameter of a data object. It applies to both intrinsic and derived types.

R616 *type-param-inquiry* is *designator* % *type-param-name*

C622 (R616) The *type-param-name* shall be the name of a type parameter of the *declared type* of the object designated by the *designator*.

- 2 A *deferred type parameter* of a pointer that is not associated or of an unallocated *allocatable* variable shall not be inquired about.

**NOTE 6.7**

A *type-param-inquiry* has a syntax like that of a *structure component* reference, but it does not have the same semantics. It is not a variable and thus can never be assigned to. It can be used only as a primary in an expression. It is scalar even if *designator* is an array.

The intrinsic type parameters can also be inquired about by using the intrinsic functions *KIND* and *LEN*.

**NOTE 6.8**

The following are examples of type parameter inquiries:

```

a%kind      !-- A is real.  Same value as KIND(a).
s%len       !-- S is character.  Same value as LEN(s).
b(10)%kind  !-- Inquiry about an array element.
p%dim       !-- P is of the derived type general_point.

```

See Note 4.22 for the definition of the *general\_point* type used in the last example above.

**6.5 Arrays****6.5.1 Order of reference**

- 1 No order of reference to the elements of an array is indicated by the appearance of the array *designator*, except where array element ordering (6.5.3.2) is specified.

**6.5.2 Whole arrays**

- 1 A *whole array* is a named array or a structure component whose final *part-ref* is an array component name; no subscript list is appended.
- 2 The appearance of a *whole array* variable in an executable construct specifies all the elements of the array (2.4.6). The appearance of a whole array *designator* in a nonexecutable statement specifies the entire array except for the appearance of a *whole array* designator in an equivalence set (5.9.1.4). An *assumed-size array* (5.5.8.5) is permitted to appear as a *whole array* in an executable construct or *specification expression* only as an *actual argument* in a *procedure reference* that does not require the shape.

**6.5.3 Array elements and array sections****6.5.3.1 Syntax**

R617 *array-element* is *data-ref*

C623 (R617) Every *part-ref* shall have *rank* zero and the last *part-ref* shall contain a *subscript-list*.

R618 *array-section* is *data-ref* [ ( *substring-range* ) ]  
or *complex-part-designator*

C624 (R618) Exactly one *part-ref* shall have nonzero *rank*, and either the final *part-ref* shall have a *section-subscript-list* with nonzero *rank*, another *part-ref* shall have nonzero *rank*, or the *complex-part-designator* shall be an array.

C625 (R618) If a *substring-range* appears, the rightmost *part-name* shall be of type character.

R619 *subscript* is *scalar-int-expr*

R620 *section-subscript* is *subscript*  
or *subscript-triplet*

or *vector-subscript*

R621 *subscript-triplet* is [ *subscript* ] : [ *subscript* ] [ : *stride* ]

R622 *stride* is *scalar-int-expr*

R623 *vector-subscript* is *int-expr*

C626 (R623) A *vector-subscript* shall be an integer array expression of **rank** one.

C627 (R621) The second subscript shall not be omitted from a *subscript-triplet* in the last dimension of an **assumed-size** array.

1 An array element is a scalar. An **array section** is an array. If a *substring-range* appears in an *array-section*, each element is the designated substring of the corresponding element of the **array section**.

2 The value of a subscript in an array element shall be within the bounds for its dimension.

#### NOTE 6.9

For example, with the declarations:

```
REAL A (10, 10)
```

```
CHARACTER (LEN = 10) B (5, 5, 5)
```

A (1, 2) is an array element, A (1:N:2, M) is a rank-one **array section**, and B (:, :, :) (2:3) is an array of shape (5, 5, 5) whose elements are substrings of length 2 of the corresponding elements of B.

#### NOTE 6.10

Unless otherwise specified, an array element or **array section** does not have an attribute of the **whole array**. In particular, an array element or an **array section** does not have the **POINTER** or **ALLOCATABLE** attribute.

#### NOTE 6.11

Examples of array elements and array sections are:

ARRAY_A(1:N:2)%ARRAY_B(I, J)%STRING(K)(:)	array section
SCALAR_PARENT%ARRAY_FIELD(J)	array element
SCALAR_PARENT%ARRAY_FIELD(1:N)	array section
SCALAR_PARENT%ARRAY_FIELD(1:N)%SCALAR_FIELD	array section

### 6.5.3.2 Array element order

1 The elements of an array form a sequence known as the array element order. The position of an array element in this sequence is determined by the subscript order value of the subscript list designating the element. The subscript order value is computed from the formulas in Table 6.1.

Table 6.1: **Subscript order value**

Rank	Subscript bounds	Subscript list	Subscript order value
1	$j_1:k_1$	$s_1$	$1 + (s_1 - j_1)$
2	$j_1:k_1, j_2:k_2$	$s_1, s_2$	$1 + (s_1 - j_1) + (s_2 - j_2) \times d_1$
3	$j_1:k_1, j_2:k_2, j_3:k_3$	$s_1, s_2, s_3$	$1 + (s_1 - j_1) + (s_2 - j_2) \times d_1 + (s_3 - j_3) \times d_2 \times d_1$
.	.	.	.
.	.	.	.

## Subscript order value

(cont.)

Rank	Subscript bounds	Subscript list	Subscript order value
.	.	.	.
15	$j_1:k_1, \dots, j_{15}:k_{15}$	$s_1, \dots, s_{15}$	$1 + (s_1 - j_1)$ $+ (s_2 - j_2) \times d_1$ $+ (s_3 - j_3) \times d_2 \times d_1$ $+ \dots$ $+ (s_{15} - j_{15}) \times d_{14}$ $\times d_{13} \times \dots \times d_1$
Notes for Table 6.1: 1) $d_i = \max(k_i - j_i + 1, 0)$ is the size of the $i$ th dimension. 2) If the size of the array is nonzero, $j_i \leq s_i \leq k_i$ for all $i = 1, 2, \dots, 15$ .			

## 6.5.3.3 Array sections

- 1 In an *array-section* having a *section-subscript-list*, each *subscript-triplet* and *vector-subscript* in the section subscript list indicates a sequence of subscripts, which may be empty. Each subscript in such a sequence shall be within the bounds for its dimension unless the sequence is empty. The *array section* is the set of elements from the array determined by all possible subscript lists obtainable from the single subscripts or sequences of subscripts specified by each section subscript.
- 2 In an *array-section* with no *section-subscript-list*, the *rank* and shape of the array is the *rank* and shape of the *part-ref* with nonzero *rank*; otherwise, the *rank* of the *array section* is the number of subscript triplets and *vector subscripts* in the section subscript list. The shape is the rank-one array whose  $i$ th element is the number of integer values in the sequence indicated by the  $i$ th subscript triplet or *vector subscript*. If any of these sequences is empty, the *array section* has size zero. The subscript order of the elements of an *array section* is that of the array data object that the *array section* represents.

## 6.5.3.3.1 Subscript triplet

- 1 A subscript triplet designates a regular sequence of subscripts consisting of zero or more subscript values. The *stride* in the subscript triplet specifies the increment between the subscript values. The subscripts and stride of a subscript triplet are optional. An omitted first subscript in a subscript triplet is equivalent to a subscript whose value is the lower bound for the array and an omitted second subscript is equivalent to the upper bound. An omitted stride is equivalent to a stride of 1.
- 2 The stride shall not be zero.
- 3 When the stride is positive, the subscripts specified by a triplet form a regularly spaced sequence of integers beginning with the first subscript and proceeding in increments of the stride to the largest such integer not greater than the second subscript; the sequence is empty if the first subscript is greater than the second.

## NOTE 6.12

For example, suppose an array is declared as A (5, 4, 3). The section A (3 : 5, 2, 1 : 2) is the array of shape (3, 2):

A (3, 2, 1)	A (3, 2, 2)
A (4, 2, 1)	A (4, 2, 2)
A (5, 2, 1)	A (5, 2, 2)

- 4 When the stride is negative, the sequence begins with the first subscript and proceeds in increments of the stride down to the smallest such integer equal to or greater than the second subscript; the sequence is empty if the second subscript is greater than the first.

**NOTE 6.13**

For example, if an array is declared B (10), the section B (9 : 1 : -2) is the array of shape (5) whose elements are B (9), B (7), B (5), B (3), and B (1), in that order.

**NOTE 6.14**

A subscript in a subscript triplet need not be within the declared bounds for that dimension if all values used in selecting the array elements are within the declared bounds.

For example, if an array is declared as B (10), the [array section](#) B (3 : 11 : 7) is the array of shape (2) consisting of the elements B (3) and B (10), in that order.

**6.5.3.3.2 Vector subscript**

- 1 A [vector subscript](#) designates a sequence of subscripts corresponding to the values of the elements of the expression. Each element of the expression shall be defined.
- 2 An [array section](#) with a [vector subscript](#) shall not be [finalized](#) by a nonelemental [final subroutine](#).
- 3 If a [vector subscript](#) has two or more elements with the same value, an [array section](#) with that [vector subscript](#) is not [definable](#) and shall not be [defined](#) or become [undefined](#).

**NOTE 6.15**

For example, suppose Z is a two-dimensional array of shape [5, 7] and U and V are one-dimensional arrays of shape (3) and (4), respectively. Assume the values of U and V are:

U = [ 1, 3, 2 ]

V = [ 2, 1, 1, 3 ]

Then Z (3, V) consists of elements from the third row of Z in the order:

Z (3, 2)    Z (3, 1)    Z (3, 1)    Z (3, 3)

and Z (U, 2) consists of the column elements:

Z (1, 2)    Z (3, 2)    Z (2, 2)

and Z (U, V) consists of the elements:

Z (1, 2)    Z (1, 1)    Z (1, 1)    Z (1, 3)

Z (3, 2)    Z (3, 1)    Z (3, 1)    Z (3, 3)

Z (2, 2)    Z (2, 1)    Z (2, 1)    Z (2, 3)

Because Z (3, V) and Z (U, V) contain duplicate elements from Z, the sections Z (3, V) and Z (U, V) shall not be redefined as sections.

**6.5.4 Simply contiguous array designators**

- 1 A [section-subscript-list](#) specifies a [simply contiguous](#) section if and only if it does not have a [vector subscript](#) and
  - all but the last [subscript-triplet](#) is a colon,
  - the last [subscript-triplet](#) does not have a [stride](#), and
  - no [subscript-triplet](#) is preceded by a [section-subscript](#) that is a [subscript](#).
- 2 An array designator is [simply contiguous](#) if and only if it is
  - an [object-name](#) that has the [CONTIGUOUS](#) attribute,

- an *object-name* that is not a *pointer*, not *assumed-shape*, and not *assumed-rank*,
- a *structure-component* whose final *part-name* is an array and that either has the *CONTIGUOUS* attribute or is not a pointer, or
- an *array section*
  - that is not a *complex-part-designator*,
  - that does not have a *substring-range*,
  - whose final *part-ref* has nonzero *rank*,
  - whose rightmost *part-name* has the *CONTIGUOUS* attribute or is neither *assumed-shape* nor a pointer, and
  - which either does not have a *section-subscript-list*, or has a *section-subscript-list* which specifies a *simply contiguous* section.

- 3 An array *variable* is *simply contiguous* if and only if it is a *simply contiguous* array designator or a reference to a function that returns a pointer with the *CONTIGUOUS* attribute.

#### NOTE 6.16

Array sections that are *simply contiguous* include column, plane, cube, and hypercube subobjects of a *simply contiguous* base object, for example:

```

ARRAY1 (10:20, 3)    ! passes part of the third column of ARRAY1.
X3D (:, i:j, 2)      ! passes part of the second plane of X3D (or the whole
                    ! plane if i==LBOUND(X3D,2) and j==UBOUND(X3D,2)).
Y5D (:, :, :, :, 7) ! passes the seventh hypercube of Y5D.

```

All *simply contiguous* designators designate *contiguous* objects.

## 6.6 Image selectors

- 1 An image selector determines the *image index* for a *coindexed* object.

R624 *image-selector* is *lbracket cosubscript-list* [ , *image-selector-spec-list* ] *rbracket*

R625 *cosubscript* is *scalar-int-expr*

R626 *image-selector-spec* is STAT = *stat-variable*  
or TEAM = *team-variable*  
or TEAM\_NUMBER = *scalar-int-expr*

C628 No specifier shall appear more than once in a given *image-selector-spec-list*.

C629 A *stat-variable* in an *image-selector* shall not be a *coindexed* object.

- 2 The number of *cosubscripts* shall be equal to the *corank* of the object. The value of a *cosubscript* in an image selector shall be within the *cobounds* for its *codimension*. Taking account of the *cobounds*, the *cosubscript* list in an image selector determines the *image index* in the same way that a subscript list in an array element determines the subscript order value (6.5.3.2), taking account of the bounds.
- 3 If a *TEAM= specifier* appears in an *image-selector*, the *team* of the image selector is specified by *team-variable*, which shall be defined with a value that represents the current or an *ancestor team*; the object shall be an *established coarray* in that *team* or an *ancestor* thereof. If a *TEAM\_NUMBER= specifier* appears in an *image-selector* and the current team is not the *initial team*, the value of the *scalar-int-expr* shall be equal to the value of a *team number* for one of the teams that were formed by execution of the *FORM TEAM* statement for the *current team* and the team of the image selector is that team; the object shall be an *established coarray* in an *ancestor* of the *current team*. Otherwise, the *team* of the image selector is the *current team*.

- 1 4 An image selector shall specify an *image index* value that is not greater than the number of *images* in the *team*  
2 of the image selector, and identifies the image with that index in that *team*.
- 3 5 Execution of a statement containing an *image-selector* with a *STAT=* specifier causes the *stat-variable* to become  
4 defined. If the designator is part of an operand that is evaluated, and the object designated is on a *failed*  
5 *image*, the *stat-variable* is defined with the value *STAT\_FAILED\_IMAGE* (13.10.2.26) in the intrinsic module  
6 *ISO\_FORTRAN\_ENV*; otherwise, it is defined with the value zero.
- 7 6 The denotation of a *stat-variable* in an *image-selector* shall not depend on the evaluation of any entity in the  
8 same statement. The value of an expression shall not depend on the value of any *stat-variable* that appears in  
9 the same statement. The value of a *stat-variable* in an *image-selector* shall not be affected by the execution of  
10 any part of the statement, other than by whether the image specified by the *image-selector* has failed.

**NOTE 6.17**

For example, if there are 16 *images* and the *coarray* A is declared

```
REAL :: A(10)[5,*]
```

A(:)[1,4] is valid because it specifies *image* 16, but A(:)[2,4] is invalid because it specifies *image* 17.

## 11 6.7 Dynamic association

### 12 6.7.1 ALLOCATE statement

#### 13 6.7.1.1 Form of the ALLOCATE statement

- 14 1 The ALLOCATE statement dynamically creates pointer *targets* and *allocatable* variables.

15 R627	<i>allocate-stmt</i>	is	ALLOCATE ( [ <i>type-spec</i> :: ] <i>allocation-list</i> ■ 16 ■ [ , <i>alloc-opt-list</i> ] )
17 R628	<i>alloc-opt</i>	is	ERRMSG = <i>errmsg-variable</i> 18 or MOLD = <i>source-expr</i> 19 or SOURCE = <i>source-expr</i> 20 or STAT = <i>stat-variable</i>
21 R629	<i>stat-variable</i>	is	<i>scalar-int-variable</i>
22 R630	<i>errmsg-variable</i>	is	<i>scalar-default-char-variable</i>
23 R631	<i>source-expr</i>	is	<i>expr</i>
24 R632	<i>allocation</i>	is	<i>allocate-object</i> [ ( <i>allocate-shape-spec-list</i> ) ] ■ 25 ■ [ <i>lbracket allocate-coarray-spec rbracket</i> ]
26 R633	<i>allocate-object</i>	is	<i>variable-name</i> 27 or <i>structure-component</i>
28 R634	<i>allocate-shape-spec</i>	is	[ <i>lower-bound-expr</i> : ] <i>upper-bound-expr</i>
29 R635	<i>lower-bound-expr</i>	is	<i>scalar-int-expr</i>
30 R636	<i>upper-bound-expr</i>	is	<i>scalar-int-expr</i>
31 R637	<i>allocate-coarray-spec</i>	is	[ <i>allocate-coshape-spec-list</i> , ] [ <i>lower-bound-expr</i> : ] *
32 R638	<i>allocate-coshape-spec</i>	is	[ <i>lower-bound-expr</i> : ] <i>upper-bound-expr</i>

- 1 C630 (R633) Each *allocate-object* shall be a data pointer or an *allocatable* variable.
- 2 C631 (R627) If any *allocate-object* has a *deferred type parameter*, is *unlimited polymorphic*, or is of *abstract*  
3 *type*, either *type-spec* or *source-expr* shall appear.
- 4 C632 (R627) If *type-spec* appears, it shall specify a type with which each *allocate-object* is *type compatible*.
- 5 C633 (R627) A *type-param-value* in a *type-spec* shall be an asterisk if and only if each *allocate-object* is a *dummy*  
6 *argument* for which the corresponding type parameter is assumed.
- 7 C634 (R627) If *type-spec* appears, the kind type parameter values of each *allocate-object* shall be the same as  
8 the corresponding type parameter values of the *type-spec*.
- 9 C635 (R627) If an *allocate-object* is an array, either *allocate-shape-spec-list* shall appear in its *allocation*, or  
10 *source-expr* shall appear in the *ALLOCATE* statement and have the same *rank* as the *allocate-object*.
- 11 C636 (R632) If *allocate-object* is scalar, *allocate-shape-spec-list* shall not appear.
- 12 C637 (R632) An *allocate-coarray-spec* shall appear if and only if the *allocate-object* is a *coarray*.
- 13 C638 (R632) The number of *allocate-shape-specs* in an *allocate-shape-spec-list* shall be the same as the *rank*  
14 of the *allocate-object*. The number of *allocate-coshape-specs* in an *allocate-coarray-spec* shall be one less  
15 than the *corank* of the *allocate-object*.
- 16 C639 (R628) No *alloc-opt* shall appear more than once in a given *alloc-opt-list*.
- 17 C640 (R627) At most one of *source-expr* and *type-spec* shall appear.
- 18 C641 (R627) Each *allocate-object* shall be *type compatible* (4.3.2.3) with *source-expr*. If *SOURCE=* appears,  
19 *source-expr* shall be a scalar or have the same *rank* as each *allocate-object*.
- 20 C642 (R627) If *source-expr* appears, the *kind type parameters* of each *allocate-object* shall have the same values  
21 as the corresponding *type parameters* of *source-expr*.
- 22 C643 (R627) *type-spec* shall not specify a type that has a *coarray ultimate component*.
- 23 C644 (R627) If an *allocate-object* is a *coarray*, *type-spec* shall not specify the type *C\_PTR* or *C\_FUNPTR*  
24 from the intrinsic module *ISO\_C\_BINDING*.
- 25 C645 (R627) The *declared type* of *source-expr* shall not be *C\_PTR* or *C\_FUNPTR* from the intrinsic module  
26 *ISO\_C\_BINDING*, or *EVENT\_TYPE* or *LOCK\_TYPE* from the intrinsic module *ISO\_FORTRAN\_*  
27 *ENV*, or have a subcomponent of type *EVENT\_TYPE* or *LOCK\_TYPE*, if an *allocate-object* is a  
28 *coarray*.
- 29 C646 (R631) The *declared type* of *source-expr* shall not have a *coarray ultimate component*.
- 30 C647 (R633) An *allocate-object* shall not be a coindexed object.

**NOTE 6.18**

If a *coarray* is of a derived type that has an *allocatable* component, the component shall be allocated by its own *image*:

```
TYPE(SOMETHING), ALLOCATABLE :: T[:]
...
ALLOCATE(T[*])           ! Allowed - implies synchronization
ALLOCATE(T%AAC(N))       ! Allowed - allocated by its own image
ALLOCATE(T[Q]%AAC(N))    ! Not allowed, because it is not
                        ! necessarily executed on image Q.
```

- 1 2 An *allocate-object* or a bound or type parameter of an *allocate-object* shall not depend on the value of *stat-variable*,  
2 the value of *errmsg-variable*, or on the value, bounds, length type parameters, allocation status, or association  
3 status of any *allocate-object* in the same ALLOCATE statement.
- 4 3 *source-expr* shall not be allocated within the ALLOCATE statement in which it appears; nor shall it depend on  
5 the value, *bounds*, *deferred type parameters*, allocation status, or association status of any *allocate-object* in that  
6 statement.
- 7 4 If an *allocate-object* is a *coarray*, the ALLOCATE statement shall not have a *source-expr* with a *dynamic type*  
8 of *C\_PTR* or *C\_FUNPTR* from the intrinsic module *ISO\_C\_BINDING*, or *EVENT\_TYPE* or *LOCK\_*  
9 *TYPE* from the intrinsic module *ISO\_FORTRAN\_ENV*, or which has a *subcomponent* whose *dynamic type*  
10 is *EVENT\_TYPE* or *LOCK\_TYPE*.
- 11 5 If *type-spec* is specified, each *allocate-object* is allocated with the specified *dynamic type* and type parameter  
12 values; if *source-expr* is specified, each *allocate-object* is allocated with the *dynamic type* and type parameter  
13 values of *source-expr*; otherwise, each *allocate-object* is allocated with its *dynamic type* the same as its *declared*  
14 *type*.
- 15 6 If *type-spec* appears and the value of a type parameter it specifies differs from the value of the corresponding  
16 nondeferred type parameter specified in the declaration of any *allocate-object*, an error condition occurs. If the  
17 value of a nondeferred length type parameter of an *allocate-object* differs from the value of the corresponding type  
18 parameter of *source-expr*, an error condition occurs.
- 19 7 If a *type-param-value* in a *type-spec* in an ALLOCATE statement is an asterisk, it denotes the current value of  
20 that assumed type parameter. If it is an expression, subsequent redefinition or undefinition of any entity in the  
21 expression does not affect the type parameter value.

**NOTE 6.19**

An example of an ALLOCATE statement is:

```
ALLOCATE (X (N), B (-3 : M, 0:9), STAT = IERR_ALLOC)
```

### 22 6.7.1.2 Execution of an ALLOCATE statement

- 23 1 When an ALLOCATE statement is executed for an array for which *allocate-shape-spec-list* is specified, the values  
24 of the lower bound and upper bound expressions determine the bounds of the array. Subsequent redefinition  
25 or undefinition of any entities in the bound expressions do not affect the array bounds. If the lower bound is  
26 omitted, the default value is 1. If the upper bound is less than the lower bound, the extent in that dimension is  
27 zero and the array has zero size.
- 28 2 When an ALLOCATE statement is executed for a *coarray*, the values of the lower *cobound* and upper *cobound*  
29 expressions determine the *cobounds* of the *coarray*. Subsequent redefinition or undefinition of any entities in the  
30 *cobound* expressions do not affect the *cobounds*. If the lower *cobound* is omitted, the default value is 1. The  
31 upper *cobound* shall not be less than the lower *cobound*.
- 32 3 If an *allocation* specifies a *coarray*, its *dynamic type* and the values of corresponding type parameters shall be the  
33 same on every *active image* in the *current team*. The values of corresponding bounds and corresponding *cobounds*  
34 shall be the same on those *images*. If the *coarray* is a *dummy argument*, its *ultimate argument* (12.5.2.3) shall  
35 be the same *coarray* on those *images*.
- 36 4 When an ALLOCATE statement is executed for which an *allocate-object* is a *coarray*, there is an implicit syn-  
37 chronization of all *active images* in the *current team*. On those *images*, execution of the segment (8.6.2) following  
38 the statement is delayed until all other *active images* in the *current team* have executed the same statement the  
39 same number of times since execution last began in this team.

**NOTE 6.20**

When an [image](#) executes an ALLOCATE statement, communication is not necessarily involved apart from any required for synchronization. The [image](#) allocates its [coarray](#) and records how the corresponding [coarrays](#) on other [images](#) are to be addressed. The processor is not required to detect violations of the rule that the bounds are the same on all [images](#), nor is it responsible for detecting or resolving deadlock problems (such as two [images](#) waiting on different ALLOCATE statements).

- 1    5 If [source-expr](#) is a pointer, it shall be associated with a [target](#). If [source-expr](#) is [allocatable](#), it shall be allocated.
- 2    6 When an ALLOCATE statement is executed for an array with no [allocate-shape-spec-list](#), the bounds of [source-expr](#) determine the bounds of the array. Subsequent changes to the bounds of [source-expr](#) do not affect the array bounds.
- 3    7 If SOURCE= appears, [source-expr](#) shall be [conformable](#) with [allocation](#). If the value of a nondeferred [length type parameter](#) of [allocate-object](#) is different from the value of the corresponding [type parameter](#) of [source-expr](#), an error condition occurs. If an [allocate-object](#) is not [polymorphic](#) and the [source-expr](#) is [polymorphic](#) with a [dynamic type](#) that differs from its [declared type](#), the value provided for that [allocate-object](#) is the ancestor component of the [source-expr](#) that has the type of the [allocate-object](#); otherwise the value provided is the value of the [source-expr](#). On successful allocation, if [allocate-object](#) and [source-expr](#) have the same [rank](#) the value of [allocate-object](#) becomes the value provided, otherwise the value of each element of [allocate-object](#) becomes the value provided. The [source-expr](#) is evaluated exactly once for each execution of an ALLOCATE statement.
- 4    8 If MOLD= appears and [source-expr](#) is a variable, its value need not be defined.
- 5    9 The set of error conditions for an ALLOCATE statement is processor dependent. If an error condition occurs during execution of an ALLOCATE statement that does not contain the STAT= specifier, error termination is initiated. The STAT= specifier is described in 6.7.4. The ERRMSG= specifier is described in 6.7.5.

**6.7.1.3 Allocation of allocatable variables**

- 1    1 The allocation status of an [allocatable](#) entity is one of the following at any time.
  - 2    • The status of an [allocatable](#) variable becomes “allocated” if it is allocated by an ALLOCATE statement, if it is allocated during assignment, or if it is given that status by the intrinsic subroutine MOVE\_ALLOC (13.9.137). An [allocatable](#) variable with this status may be referenced, defined, or deallocated; allocating it causes an error condition in the ALLOCATE statement. The result of the intrinsic function ALLOCATED (13.9.11) is true for such a variable.
  - 3    • An [allocatable](#) variable has a status of “unallocated” if it is not allocated. The status of an [allocatable](#) variable becomes unallocated if it is deallocated (6.7.3) or if it is given that status by the intrinsic subroutine MOVE\_ALLOC. An [allocatable](#) variable with this status shall not be referenced or defined. It shall not be supplied as an [actual argument](#) corresponding to a nonallocatable nonoptional [dummy argument](#), except to certain intrinsic [inquiry functions](#). It may be allocated with the ALLOCATE statement. Deallocating it causes an error condition in the DEALLOCATE statement. The result of the intrinsic function ALLOCATED (13.9.11) is false for such a variable.
- 2    2 At the beginning of execution of a program, [allocatable](#) variables are unallocated.
- 3    3 When the allocation status of an [allocatable](#) variable changes, the allocation status of any associated [allocatable](#) variable changes accordingly. Allocation of an [allocatable](#) variable establishes values for the [deferred type parameters](#) of all associated [allocatable](#) variables.
- 4    4 An [unsaved allocatable local variable](#) of a procedure has a status of unallocated at the beginning of each invocation of the procedure. An [unsaved local variable](#) of a construct has a status of unallocated at the beginning of each execution of the construct.
- 5    5 When an object of derived type is created by an ALLOCATE statement, any [allocatable ultimate components](#) have an allocation status of unallocated unless the SOURCE= specifier appears and the corresponding component

of the *source-expr* is allocated.

- 6 If the evaluation of a function would change the allocation status of a variable and if a reference to the function appears in an expression in which the value of the function is not needed to determine the value of the expression, the allocation status of the variable after evaluation of the expression is processor dependent.

#### 6.7.1.4 Allocation of pointer targets

- 1 Allocation of a *pointer* creates an object that implicitly has the **TARGET attribute**. Following successful execution of an **ALLOCATE** statement for a *pointer*, the *pointer* is associated with the *target* and may be used to reference or define the *target*. Additional *pointers* may become associated with the *pointer target* or a part of the *pointer target* by *pointer assignment*. It is not an error to allocate a *pointer* that is already associated with a *target*. In this case, a new *pointer target* is created as required by the attributes of the *pointer* and any array bounds, type, and type parameters specified by the **ALLOCATE** statement. The *pointer* is then associated with this new *target*. Any previous association of the *pointer* with a *target* is broken. If the previous *target* had been created by allocation, it becomes inaccessible unless other *pointers* are associated with it. The intrinsic function **ASSOCIATED** (13.9.16) may be used to determine whether a *pointer* that does not have undefined association status is associated.
- 2 At the beginning of execution of a function whose result is a *pointer*, the association status of the result *pointer* is undefined. Before such a function returns, it shall either associate a *target* with this *pointer* or cause the association status of this *pointer* to become *disassociated*.

### 6.7.2 NULLIFY statement

- 1 The **NULLIFY** statement causes *pointers* to be *disassociated*.

R639 *nullify-stmt* is **NULLIFY** ( *pointer-object-list* )

R640 *pointer-object* is *variable-name*  
or *structure-component*  
or *proc-pointer-name*

C648 (R640) Each *pointer-object* shall have the **POINTER attribute**.

- 2 A *pointer-object* shall not depend on the value, bounds, or association status of another *pointer-object* in the same **NULLIFY** statement.

#### NOTE 6.21

When a **NULLIFY** statement is applied to a *polymorphic pointer* (4.3.2.3), its *dynamic type* becomes the *declared type*.

### 6.7.3 DEALLOCATE statement

#### 6.7.3.1 Form of the DEALLOCATE statement

- 1 The **DEALLOCATE** statement causes *allocatable* variables to be deallocated; it causes *pointer targets* to be deallocated and the *pointers* to be *disassociated*.

R641 *deallocate-stmt* is **DEALLOCATE** ( *allocate-object-list* [ , *dealloc-opt-list* ] )

R642 *dealloc-opt* is **STAT** = *stat-variable*  
or **ERRMSG** = *errmsg-variable*

C649 (R642) No *dealloc-opt* shall appear more than once in a given *dealloc-opt-list*.

- 2 An *allocate-object* shall not depend on the value, bounds, allocation status, or association status of another *allocate-object* in the same **DEALLOCATE** statement; it also shall not depend on the value of the *stat-variable*

or *errmsg-variable* in the same DEALLOCATE statement.

The set of error conditions for a DEALLOCATE statement is processor dependent. If an error condition occurs during execution of a DEALLOCATE statement that does not contain the *STAT= specifier*, *error termination* is initiated. The *STAT= specifier* is described in 6.7.4. The *ERRMSG= specifier* is described in 6.7.5.

When more than one allocated object is deallocated by execution of a DEALLOCATE statement, the order of deallocation is processor dependent.

#### NOTE 6.22

An example of a DEALLOCATE statement is:

```
DEALLOCATE (X, B)
```

### 6.7.3.2 Deallocation of allocatable variables

Deallocating an unallocated *allocatable* variable causes an error condition in the DEALLOCATE statement. Deallocating an *allocatable* variable with the *TARGET attribute* causes the *pointer association* status of any *pointer* associated with it to become *undefined*. An *allocatable* variable shall not be deallocated if it or any subobject of it is *argument associated* with a *dummy argument* or *construct associated* with an *associate name*.

When the execution of a procedure is terminated by execution of a *RETURN* or *END* statement, an *unsaved allocatable local variable* of the procedure retains its allocation and definition status if it is a *function result* or a subobject thereof; otherwise, if it is allocated it will be deallocated.

When a *BLOCK construct* terminates, any *unsaved* allocated *allocatable local variable* of the construct is deallocated.

If an executable construct references a function whose result is *allocatable* or has an *allocatable* subobject, and the function reference is executed, an *allocatable* result and any allocated *allocatable* subobject of the result is deallocated after execution of the innermost executable construct containing the reference.

If a function whose result is *allocatable* or has an *allocatable* subobject is referenced in the specification part of a *scoping unit*, and the function reference is executed, an *allocatable* result and any allocated *allocatable* subobject of the result is deallocated before execution of the executable constructs of the *scoping unit*.

When a procedure is invoked, any allocated *allocatable* object that is an *actual argument* corresponding to an *INTENT (OUT) allocatable dummy argument* is deallocated; any allocated *allocatable* object that is a subobject of an *actual argument* corresponding to an *INTENT (OUT) dummy argument* is deallocated. If a Fortran procedure that has an *INTENT (OUT) allocatable dummy argument* is invoked by a C function and the corresponding argument in the C function call is a *C descriptor* that describes an allocated *allocatable* variable, the variable is deallocated on entry to the Fortran procedure. If a C function is invoked from a Fortran procedure via an interface with an *INTENT (OUT) allocatable dummy argument* and the corresponding *actual argument* in the reference to the C function is an allocated *allocatable* variable, the variable is deallocated on invocation (before execution of the C function begins).

When an *intrinsic assignment statement* (7.2.1.3) is executed, any noncoarray allocated *allocatable* subobject of the variable is deallocated before the assignment takes place.

When a variable of derived type is deallocated, any allocated *allocatable* subobject is deallocated. If an error condition occurs during deallocation, it is *processor dependent* whether an allocated *allocatable subobject* is deallocated.

If an *allocatable* component is a subobject of a *finalizable* object, that object is finalized before the component is automatically deallocated.

When a statement that deallocates a *coarray* is executed, there is an implicit synchronization of all *active images* in the *current team*. On those *images*, execution of the segment (8.6.2) following the statement is delayed until

- 1 all other **active images** in the **current team** have executed the same statement the same number of times since  
 2 execution last began in this **team**. If an **allocate-object** is a **coarray dummy argument**, its **ultimate argument**  
 3 (12.5.2.3) shall be the same **coarray** on those **images**.
- 4 11 The effect of automatic deallocation is the same as that of a DEALLOCATE statement without a **dealloc-opt-list**.

**NOTE 6.23**

In the following example:

```
SUBROUTINE PROCESS
  REAL, ALLOCATABLE :: TEMP(:)
  REAL, ALLOCATABLE, SAVE :: X(:)
  ...
END SUBROUTINE PROCESS
```

on return from subroutine PROCESS, the allocation status of X is preserved because X has the **SAVE attribute**. TEMP does not have the **SAVE attribute**, so it will be deallocated if it was allocated. On the next invocation of PROCESS, TEMP will have an allocation status of unallocated.

**NOTE 6.24**

For example, executing a **RETURN**, **END**, or **END BLOCK** statement, or deallocating an object that has an allocatable subobject, can cause deallocation of a **coarray**, and thus an implicit synchronization of all **images**.

### 5 6.7.3.3 Deallocation of pointer targets

- 6 1 If a pointer appears in a DEALLOCATE statement, its association status shall be defined. Deallocating a pointer  
 7 that is **disassociated** or whose **target** was not created by an ALLOCATE statement causes an error condition  
 8 in the DEALLOCATE statement. If a pointer is associated with an **allocatable** entity, the pointer shall not be  
 9 deallocated. A **pointer** shall not be deallocated if its **target** or any **subobject** thereof is **argument associated** with  
 10 a **dummy argument** or **construct associated** with an **associate name**.
- 11 2 If a pointer appears in a DEALLOCATE statement, it shall be associated with the whole of an object that was  
 12 created by allocation. The pointer shall have the same **dynamic type** and type parameters as the allocated object,  
 13 and if the allocated object is an array the pointer shall be an array whose elements are the same as those of the  
 14 allocated object in array element order. Deallocating a pointer **target** causes the pointer association status of any  
 15 other pointer that is associated with the **target** or a portion of the **target** to become undefined.

### 16 6.7.4 STAT= specifier

- 17 1 The **stat-variable** shall not be allocated or deallocated within the **ALLOCATE** or **DEALLOCATE** statement  
 18 in which it appears; nor shall it depend on the value, **bounds**, **deferred type parameters**, allocation status, or  
 19 association status of any **allocate-object** in that statement.
- 20 2 If the STAT= specifier appears, successful execution of the **ALLOCATE** or **DEALLOCATE** statement causes the  
 21 **stat-variable** to become defined with a value of zero.

**Unresolved Technical Issue 016**

(DE)ALLOCATE mess re STAT=

The TS says to change the above paragraph to

If the STAT= specifier appears, successful execution of the ALLOCATE or DEALLOCATE statement on all images in the current team causes the **stat-variable** to become defined with the value zero; otherwise, if

## Unresolved Technical Issue 016 (cont.)

a failed image is detected and execution is otherwise successful, the STAT= specifier is assigned the value of the named constant STAT\_FAILED\_IMAGE in the intrinsic module ISO\_FORTRAN\_ENV (13.8.2).

This has the concept that ALLOCATE (or DEALLOCATE) might succeed on some images and fail on others. We have already rejected that concept. Either it succeeds on all or fails on all. There is no concept of “otherwise successful”. We should not make one of the error cases (failed images) so different from the other error cases that every program that tries to handle errors intelligently needs to be rewritten to take our newfangled fault tolerance into account everywhere.

I think I understand that STAT\_STOPPED\_IMAGE is supposed to take precedence over STAT\_FAILED\_IMAGE?

In any case, the existing text is **correct as is**; STAT\_FAILED\_IMAGE needs to be handled as part of the error (nonzero STAT=) case. That is, in the next paragraph, where it will

- therefore be clearer which of STOPPED and FAILED is intended to take precedence, instead of STOPPED apparently taking precedence because of sloppy wording;
- be clear that the other requirements of the error case re the status of each *allocate-object* will apply.

I have also rejected the edits to the next paragraph in the TS as being internally inconsistent. Obviously (!) if we describe ALLOCATE or DEALLOCATE as only having effect on the active images, **there is no interaction with any stopped image**, meaning STAT\_STOPPED\_IMAGE is never raised.

- 3 If an ALLOCATE or DEALLOCATE statement with a *coarray allocate-object* is executed when one or more *images* has initiated *normal termination* of execution, the *stat-variable* becomes defined with the processor-dependent positive integer value of the constant STAT\_STOPPED\_IMAGE from the intrinsic module ISO\_FORTRAN\_ENV (13.10.2). If any other error condition occurs during execution of the ALLOCATE or DEALLOCATE statement, the *stat-variable* becomes defined with a processor-dependent positive integer value different from STAT\_STOPPED\_IMAGE. In either case, each *allocate-object* has a processor-dependent status:
- each *allocate-object* that was successfully allocated shall have an allocation status of allocated or a *pointer association* status of associated;
  - each *allocate-object* that was successfully deallocated shall have an allocation status of unallocated or a *pointer association* status of disassociated;
  - each *allocate-object* that was not successfully allocated or deallocated shall retain its previous allocation status or *pointer association* status.

## NOTE 6.25

The status of objects that were not successfully allocated or deallocated can be individually checked with the intrinsic functions ALLOCATED or ASSOCIATED.

## 6.7.5 ERRMSG= specifier

- 1 The *errmsg-variable* shall not be allocated or deallocated within the ALLOCATE or DEALLOCATE statement in which it appears; nor shall it depend on the value, *bounds*, *deferred type parameters*, allocation status, or association status of any *allocate-object* in that statement.
- 2 If an error condition occurs during execution of an ALLOCATE or DEALLOCATE statement, the *errmsg-variable* is assigned an explanatory message, as if by *intrinsic assignment*. If no such condition occurs, the definition status and value of *errmsg-variable* are unchanged.



## 7 Expressions and assignment

### 7.1 Expressions

#### 7.1.1 Expression semantics

- 1 An expression represents either a data object reference or a computation, and its value is either a scalar or an array. Evaluation of an expression produces a value, which has a type, type parameters (if appropriate), and a shape (7.1.9). The *corank* of an expression that is not a variable is zero.

#### 7.1.2 Form of an expression

##### 7.1.2.1 Overall expression syntax

- 1 An expression is formed from operands, operators, and parentheses. An operand is either a scalar or an array. An operation is either intrinsic (7.1.5) or defined (7.1.6). More complicated expressions can be formed using operands which are themselves expressions.
- 2 An expression is defined in terms of several categories: primary, level-1 expression, level-2 expression, level-3 expression, level-4 expression, and level-5 expression.
- 3 These categories are related to the different operator precedence levels and, in general, are defined in terms of other categories. The simplest form of each expression category is a *primary*.

##### 7.1.2.2 Primary

R701 *primary* is *literal-constant*  
or *designator*  
or *array-constructor*  
or *structure-constructor*  
or *function-reference*  
or *type-param-inquiry*  
or *type-param-name*  
or ( *expr* )

C701 (R701) The *type-param-name* shall be the name of a type parameter.

C702 (R701) The *designator* shall not be a whole *assumed-size array*.

C703 (R701) The *expr* shall not be a function reference that returns a procedure pointer.

#### NOTE 7.1

Examples of a *primary* are:

<u>Example</u>	<u>Syntactic class</u>
1.0	<i>constant</i>
'ABCDEFGHJKLMNOPQRSTUVWXYZ' (I:I)	<i>designator</i>
[ 1.0, 2.0 ]	<i>array-constructor</i>
PERSON ('Jones', 12)	<i>structure-constructor</i>
F (X, Y)	<i>function-reference</i>
X%KIND	<i>type-param-inquiry</i>
KIND	<i>type-param-name</i>
(S + T)	( <i>expr</i> )

### 7.1.2.3 Level-1 expressions

- 1 Defined unary operators have the highest operator precedence (Table 7.1). Level-1 expressions are primaries optionally operated on by defined unary operators:

R702 *level-1-expr* is [ *defined-unary-op* ] *primary*

R703 *defined-unary-op* is . *letter* [ *letter* ] ... .

C704 (R703) A *defined-unary-op* shall not contain more than 63 letters and shall not be the same as any *intrinsic-operator* or *logical-literal-constant*.

#### NOTE 7.2

Simple examples of a level-1 expression are:

Example	Syntactic class
A	<i>primary</i> (R701)
. INVERSE. B	<i>level-1-expr</i> (R702)

A more complicated example of a level-1 expression is:

. INVERSE. (A + B)

### 7.1.2.4 Level-2 expressions

- 1 Level-2 expressions are level-1 expressions optionally involving the numeric operators *power-op*, *mult-op*, and *add-op*.

R704 *mult-operand* is *level-1-expr* [ *power-op mult-operand* ]

R705 *add-operand* is [ *add-operand mult-op* ] *mult-operand*

R706 *level-2-expr* is [ [ *level-2-expr* ] *add-op* ] *add-operand*

R707 *power-op* is \*\*

R708 *mult-op* is \*

or /

R709 *add-op* is +

or -

#### NOTE 7.3

Simple examples of a level-2 expression are:

Example	Syntactic class	Remarks
A	<i>level-1-expr</i>	A is a <i>primary</i> . (R702)
B ** C	<i>mult-operand</i>	B is a <i>level-1-expr</i> , ** is a <i>power-op</i> , and C is a <i>mult-operand</i> . (R704)
D * E	<i>add-operand</i>	D is an <i>add-operand</i> , * is a <i>mult-op</i> , and E is a <i>mult-operand</i> . (R705)
+1	<i>level-2-expr</i>	+ is an <i>add-op</i> and 1 is an <i>add-operand</i> . (R706)
F - I	<i>level-2-expr</i>	F is a <i>level-2-expr</i> , - is an <i>add-op</i> , and I is an <i>add-operand</i> . (R706)

**NOTE 7.3 (cont.)**

A more complicated example of a level-2 expression is:

- A + D \* E + B \*\* C

**7.1.2.5 Level-3 expressions**

1 Level-3 expressions are level-2 expressions optionally involving the character operator *concat-op*.

R710 *level-3-expr* is [ *level-3-expr concat-op* ] *level-2-expr*

R711 *concat-op* is //

**NOTE 7.4**

Simple examples of a level-3 expression are:

<u>Example</u>	<u>Syntactic class</u>
A	<i>level-2-expr</i> (R706)
B // C	<i>level-3-expr</i> (R710)

A more complicated example of a level-3 expression is:

X // Y // 'ABCD'

**7.1.2.6 Level-4 expressions**

1 Level-4 expressions are level-3 expressions optionally involving the relational operators *rel-op*.

R712 *level-4-expr* is [ *level-3-expr rel-op* ] *level-3-expr*

R713 *rel-op* is .EQ.  
or .NE.  
or .LT.  
or .LE.  
or .GT.  
or .GE.  
or ==  
or /=  
or <  
or <=  
or >  
or >=

**NOTE 7.5**

Simple examples of a level-4 expression are:

<u>Example</u>	<u>Syntactic class</u>
A	<i>level-3-expr</i> (R710)
B == C	<i>level-4-expr</i> (R712)
D < E	<i>level-4-expr</i> (R712)

A more complicated example of a level-4 expression is:

(A + B) /= C

### 7.1.2.7 Level-5 expressions

- 1 Level-5 expressions are level-4 expressions optionally involving the logical operators *not-op*, *and-op*, *or-op*, and *equiv-op*.

R714	<i>and-operand</i>	is	[ <i>not-op</i> ] <i>level-4-expr</i>
R715	<i>or-operand</i>	is	[ <i>or-operand and-op</i> ] <i>and-operand</i>
R716	<i>equiv-operand</i>	is	[ <i>equiv-operand or-op</i> ] <i>or-operand</i>
R717	<i>level-5-expr</i>	is	[ <i>level-5-expr equiv-op</i> ] <i>equiv-operand</i>
R718	<i>not-op</i>	is	.NOT.
R719	<i>and-op</i>	is	.AND.
R720	<i>or-op</i>	is	.OR.
R721	<i>equiv-op</i>	is	.EQV.
		or	.NEQV.

#### NOTE 7.6

Simple examples of a level-5 expression are:

Example	Syntactic class
A	<i>level-4-expr</i> (R712)
.NOT. B	<i>and-operand</i> (R714)
C .AND. D	<i>or-operand</i> (R715)
E .OR. F	<i>equiv-operand</i> (R716)
G .EQV. H	<i>level-5-expr</i> (R717)
S .NEQV. T	<i>level-5-expr</i> (R717)

A more complicated example of a level-5 expression is:

A .AND. B .EQV. .NOT. C

### 7.1.2.8 General form of an expression

- 1 Expressions are level-5 expressions optionally involving defined binary operators. Defined binary operators have the lowest operator precedence (Table 7.1).

R722	<i>expr</i>	is	[ <i>expr defined-binary-op</i> ] <i>level-5-expr</i>
R723	<i>defined-binary-op</i>	is	. letter [ letter ] ... .

- C705 (R723) A *defined-binary-op* shall not contain more than 63 letters and shall not be the same as any *intrinsic-operator* or *logical-literal-constant*.

#### NOTE 7.7

Simple examples of an expression are:

Example	Syntactic class
A	<i>level-5-expr</i> (R717)
B.UNION.C	<i>expr</i> (R722)

**NOTE 7.7 (cont.)**

More complicated examples of an expression are:

```
(B .INTERSECT. C) .UNION. (X - Y)
A + B == C * D
.INVERSE. (A + B)
A + B .AND. C * D
E // G == H (1:10)
```

**7.1.3 Precedence of operators**

- 1 There is a precedence among the intrinsic and extension operations corresponding to the form of expressions specified in 7.1.2, which determines the order in which the operands are combined unless the order is changed by the use of parentheses. This precedence order is summarized in Table 7.1.

Table 7.1: Categories of operations and relative precedence

Category of operation	Operators	Precedence
Extension	<i>defined-unary-op</i>	Highest
Numeric	**	.
Numeric	*, /	.
Numeric	unary +, -	.
Numeric	binary +, -	.
Character	//	.
Relational	.EQ., .NE., .LT., .LE., .GT., .GE., ==, /=, <, <=, >, >=	.
Logical	.NOT.	.
Logical	.AND.	.
Logical	.OR.	.
Logical	.EQV., .NEQV.	.
Extension	<i>defined-binary-op</i>	Lowest

- 2 The precedence of a *defined operation* is that of its operator.

**NOTE 7.8**

For example, in the expression

```
-A ** 2
```

the exponentiation operator (\*\*) has precedence over the negation operator (-); therefore, the operands of the exponentiation operator are combined to form an expression that is used as the operand of the negation operator. The interpretation of the above expression is the same as the interpretation of the expression

```
- (A ** 2)
```

- 3 The general form of an expression (7.1.2) also establishes a precedence among operators in the same syntactic class. This precedence determines the order in which the operands are to be combined in determining the interpretation of the expression unless the order is changed by the use of parentheses.

**NOTE 7.9**

In interpreting a *level-2-expr* containing two or more binary operators + or -, each operand (*add-operand*) is combined from left to right. Similarly, the same left-to-right interpretation for a *mult-operand* in *add-operand*, as well as for other kinds of expressions, is a consequence of the general form. However, for interpreting a *mult-operand* expression when two or more exponentiation operators \*\* combine *level-1-expr* operands, each *level-1-expr* is combined from right to left.

**NOTE 7.9 (cont.)**

For example, the expressions

```
2.1 + 3.4 + 4.9
2.1 * 3.4 * 4.9
2.1 / 3.4 / 4.9
2 ** 3 ** 4
'AB' // 'CD' // 'EF'
```

have the same interpretations as the expressions

```
(2.1 + 3.4) + 4.9
(2.1 * 3.4) * 4.9
(2.1 / 3.4) / 4.9
2 ** (3 ** 4)
('AB' // 'CD') // 'EF'
```

As a consequence of the general form (7.1.2), only the first *add-operand* of a *level-2-expr* can be preceded by the identity (+) or negation (−) operator. These formation rules do not permit expressions containing two consecutive numeric operators, such as  $A ** -B$  or  $A + -B$ . However, expressions such as  $A ** (-B)$  and  $A + (-B)$  are permitted. The rules do allow a binary operator or an intrinsic unary operator to be followed by a defined unary operator, such as:

```
A * .INVERSE. B
- .INVERSE. (B)
```

As another example, in the expression

```
A .OR. B .AND. C
```

the general form implies a higher precedence for the *.AND.* operator than for the *.OR.* operator; therefore, the interpretation of the above expression is the same as the interpretation of the expression

```
A .OR. (B .AND. C)
```

**NOTE 7.10**

An expression can contain more than one category of operator. The logical expression

```
L .OR. A + B >= C
```

where A, B, and C are of type real, and L is of type logical, contains a numeric operator, a relational operator, and a logical operator. This expression would be interpreted the same as the expression

```
L .OR. ((A + B) >= C)
```

**NOTE 7.11**

If

- the operator *\*\** is extended to type logical,
- the operator *.STARSTAR.* is defined to duplicate the function of *\*\** on type real,
- *.MINUS.* is defined to duplicate the unary operator *−*, and
- L1 and L2 are type logical and X and Y are type real,

then in precedence:  $L1 ** L2$  is higher than  $X * Y$ ;  $X * Y$  is higher than  $X .STARSTAR. Y$ ; and *.MINUS.* X is higher than  $-X$ .

## 7.1.4 Evaluation of operations

- 1 An intrinsic operation requires the values of its operands.
- 2 Execution of a function reference in the logical expression in an **IF statement** (8.1.8.4), the mask expression in a **WHERE statement** (7.2.3.1), or the *concurrent-limits* and *concurrent-steps* in a **FORALL statement** (7.2.4) is permitted to define variables in the subsidiary *action-stmt*, *where-assignment-stmt*, or *forall-assignment-stmt* respectively. Except in those cases:
  - the evaluation of a function reference shall neither affect nor be affected by the evaluation of any other entity within the statement;
  - if a function reference causes definition or undefinition of an **actual argument** of the function, that argument or any associated entities shall not appear elsewhere in the same statement.

### NOTE 7.12

For example, the statements

```
A (I) = F (I)
Y = G (X) + X
```

are prohibited if the reference to F defines or undefines I or the reference to G defines or undefines X.

However, in the statements

```
IF (F (X)) A = X
WHERE (G (X)) B = X
```

the reference to F and/or the reference to G can define X.

- 3 The appearance of an array constructor requires the evaluation of each *scalar-int-expr* of the *ac-implied-do-control* in any *ac-implied-do* it may contain.
- 4 When an **elemental** binary operation is applied to a scalar and an array or to two arrays of the same shape, the operation is performed element-by-element on corresponding array elements of the array operands.

### NOTE 7.13

For example, the array expression

```
A + B
```

produces an array of the same shape as A and B. The individual array elements of the result have the values of the first element of A added to the first element of B, the second element of A added to the second element of B, etc.

- 5 When an **elemental** unary operator operates on an array operand, the operation is performed element-by-element, and the result is the same shape as the operand. If an **elemental operation** is intrinsically pure or is implemented by a pure **elemental** function (12.8), the element operations may be performed simultaneously or in any order.

## 7.1.5 Intrinsic operations

### 7.1.5.1 Intrinsic operation classification

- 1 An intrinsic operation is either a unary or binary operation. An intrinsic unary operation is an operation of the form *intrinsic-operator*  $x_2$  where  $x_2$  is of an intrinsic type (4.4) listed in Table 7.2 for the unary intrinsic operator.
- 2 An intrinsic binary operation is an operation of the form  $x_1$  *intrinsic-operator*  $x_2$  where  $x_1$  and  $x_2$  are **conformable** and of the intrinsic types (4.4) listed in Table 7.2 for the binary intrinsic operator.

- 1 3 A numeric intrinsic operation is an intrinsic operation for which the *intrinsic-operator* is a numeric operator (+, 2 -, \*, /, or \*\*). A numeric intrinsic operator is the operator in a numeric intrinsic operation.
- 3 4 The character intrinsic operation is the intrinsic operation for which the *intrinsic-operator* is (//) and both 4 operands are of type character with the same *kind type parameter*. The character intrinsic operator is the 5 operator in a character intrinsic operation.
- 6 5 A logical intrinsic operation is an intrinsic operation for which the *intrinsic-operator* is .AND., .OR., .NOT., 7 .EQV., or .NEQV. and both operands are of type logical. A logical intrinsic operator is the operator in a logical 8 intrinsic operation.
- 9 6 A relational intrinsic operator is an *intrinsic-operator* that is .EQ., .NE., .GT., .GE., .LT., .LE., ==, /=, >, 10 >=, <, or <=. A relational intrinsic operation is an intrinsic operation for which the *intrinsic-operator* is a 11 relational intrinsic operator. A numeric relational intrinsic operation is a relational intrinsic operation for which 12 both operands are of *numeric type*. A character relational intrinsic operation is a relational intrinsic operation for 13 which both operands are of type character. The *kind type parameters* of the operands of a character relational 14 intrinsic operation shall be the same.
- 15 7 The interpretations defined in subclause 7.1.5 apply to both scalars and arrays; the interpretation for arrays is 16 obtained by applying the interpretation for scalars element by element.

**NOTE 7.14**

For example, if X is of type real, J is of type integer, and INT is the real-to-integer intrinsic conversion function, the expression INT (X + J) is an integer expression and X + J is a real expression.

Table 7.2: **Type of operands and results for intrinsic operators**

Intrinsic operator <i>op</i>	Type of $x_1$	Type of $x_2$	Type of $[x_1] \text{ op } x_2$
Unary +, -		I, R, Z	I, R, Z
Binary +, -, *, /, **	I	I, R, Z	I, R, Z
	R	I, R, Z	R, R, Z
	Z	I, R, Z	Z, Z, Z
//	C	C	C
.EQ., .NE., ==, /=	I	I, R, Z	L, L, L
	R	I, R, Z	L, L, L
	Z	I, R, Z	L, L, L
	C	C	L
.GT., .GE., .LT., .LE. >, >=, <, <=	I	I, R	L, L
	R	I, R	L, L
	C	C	L
.NOT.		L	L
.AND., .OR., .EQV., .NEQV.	L	L	L
Note: The symbols I, R, Z, C, and L stand for the types integer, real, complex, character, and logical, respectively. Where more than one type for $x_2$ is given, the type of the result of the operation is given in the same relative position in the next column.			

17 **7.1.5.2 Numeric intrinsic operations**18 **7.1.5.2.1 Interpretation of numeric intrinsic operations**

- 19 1 The two operands of numeric intrinsic binary operations may be of different *numeric types* or different *kind type* 20 *parameters*. Except for a value raised to an integer power, if the operands have different types or *kind type* 21 *parameters*, the effect is as if each operand that differs in type or *kind type parameter* from those of the result is 22 converted to the type and *kind type parameter* of the result before the operation is performed. When a value of

type real or complex is raised to an integer power, the integer operand need not be converted.

A numeric operation is used to express a numeric computation. Evaluation of a numeric operation produces a numeric value. The permitted data types for operands of the numeric intrinsic operations are specified in 7.1.5.1.

The numeric operators and their interpretation in an expression are given in Table 7.3, where  $x_1$  denotes the operand to the left of the operator and  $x_2$  denotes the operand to the right of the operator.

Table 7.3: Interpretation of the numeric intrinsic operators

Operator	Representing	Use of operator	Interpretation
**	Exponentiation	$x_1 ** x_2$	Raise $x_1$ to the power $x_2$
/	Division	$x_1 / x_2$	Divide $x_1$ by $x_2$
*	Multiplication	$x_1 * x_2$	Multiply $x_1$ by $x_2$
−	Subtraction	$x_1 - x_2$	Subtract $x_2$ from $x_1$
−	Negation	$-x_2$	Negate $x_2$
+	Addition	$x_1 + x_2$	Add $x_1$ and $x_2$
+	Identity	$+x_2$	Same as $x_2$

The interpretation of a division operation depends on the types of the operands (7.1.5.2.2).

If  $x_1$  and  $x_2$  are of type integer and  $x_2$  has a negative value, the interpretation of  $x_1 ** x_2$  is the same as the interpretation of  $1/(x_1 ** \text{ABS}(x_2))$ , which is subject to the rules of integer division (7.1.5.2.2).

#### NOTE 7.15

For example,  $2 ** (-3)$  has the value of  $1/(2 ** 3)$ , which is zero.

#### 7.1.5.2.2 Integer division

One operand of type integer may be divided by another operand of type integer. Although the mathematical quotient of two integers is not necessarily an integer, Table 7.2 specifies that an expression involving the division operator with two operands of type integer is interpreted as an expression of type integer. The result of such an operation is the integer closest to the mathematical quotient and between zero and the mathematical quotient inclusively.

#### NOTE 7.16

For example, the expression  $(-8) / 3$  has the value  $(-2)$ .

#### 7.1.5.2.3 Complex exponentiation

In the case of a complex value raised to a complex power, the value of the operation  $x_1 ** x_2$  is the principal value of  $x_1^{x_2}$ .

#### 7.1.5.2.4 Evaluation of numeric intrinsic operations

The execution of any numeric operation whose result is not defined by the arithmetic used by the processor is prohibited. Raising a negative-valued primary of type real to a real power is prohibited.

Once the interpretation of a numeric intrinsic operation is established, the processor may evaluate any mathematically equivalent expression, provided that the integrity of parentheses is not violated.

Two expressions of a numeric type are mathematically equivalent if, for all possible values of their primaries, their mathematical values are equal. However, mathematically equivalent expressions of numeric type may produce different computational results.

**NOTE 7.17**

Any difference between the values of the expressions  $(1./3.)*3.$  and  $1.$  is a computational difference, not a mathematical difference. The difference between the values of the expressions  $5/2$  and  $5./2.$  is a mathematical difference, not a computational difference.

The mathematical definition of integer division is given in [7.1.5.2.2](#).

**NOTE 7.18**

The following are examples of expressions with allowable alternative forms that can be used by the processor in the evaluation of those expressions. A, B, and C represent arbitrary real or complex operands; I and J represent arbitrary integer operands; and X, Y, and Z represent arbitrary operands of [numeric type](#).

<u>Expression</u>	<u>Allowable alternative form</u>
$X + Y$	$Y + X$
$X * Y$	$Y * X$
$-X + Y$	$Y - X$
$X + Y + Z$	$X + (Y + Z)$
$X - Y + Z$	$X - (Y - Z)$
$X * A / Z$	$X * (A / Z)$
$X * Y - X * Z$	$X * (Y - Z)$
$A / B / C$	$A / (B * C)$
$A / 5.0$	$0.2 * A$

The following are examples of expressions with forbidden alternative forms that cannot be used by a processor in the evaluation of those expressions.

<u>Expression</u>	<u>Forbidden alternative form</u>
$I / 2$	$0.5 * I$
$X * I / J$	$X * (I / J)$
$I / J / A$	$I / (J * A)$
$(X + Y) + Z$	$X + (Y + Z)$
$(X * Y) - (X * Z)$	$X * (Y - Z)$
$X * (Y - Z)$	$X * Y - X * Z$

**NOTE 7.19**

In addition to the parentheses required to establish the desired interpretation, parentheses can be included to restrict the alternative forms that can be used by the processor in the actual evaluation of the expression. This is useful for controlling the magnitude and accuracy of intermediate values developed during the evaluation of an expression.

For example, in the expression

$$A + (B - C)$$

the parenthesized expression  $(B - C)$  is evaluated and then added to A.

The inclusion of parentheses could change the mathematical value of an expression. For example, the two expressions

$$A * I / J$$

$$A * (I / J)$$

could have different mathematical values if I and J are of type integer.

**NOTE 7.20**

Each operand in a numeric intrinsic operation has a type that can depend on the order of evaluation used by the processor.

For example, in the evaluation of the expression

$$Z + R + I$$

where Z, R, and I represent data objects of complex, real, and integer type, respectively, the type of the operand that is added to I could be either complex or real, depending on which pair of operands (Z and R, R and I, or Z and I) is added first.

**7.1.5.3 Character intrinsic operation****7.1.5.3.1 Interpretation of the character intrinsic operation**

- 1 The character intrinsic operator // is used to concatenate two operands of type character with the same [kind type parameter](#). Evaluation of the character intrinsic operation produces a result of type character.
- 2 The interpretation of the character intrinsic operator // when used to form an expression is given in Table 7.4, where  $x_1$  denotes the operand to the left of the operator and  $x_2$  denotes the operand to the right of the operator.

Table 7.4: **Interpretation of the character intrinsic operator //**

Operator	Representing	Use of operator	Interpretation
//	Concatenation	$x_1 // x_2$	Concatenate $x_1$ with $x_2$

- 3 The result of the character intrinsic operation // is a character string whose value is the value of  $x_1$  concatenated on the right with the value of  $x_2$  and whose length is the sum of the lengths of  $x_1$  and  $x_2$ . Parentheses used to specify the order of evaluation have no effect on the value of a character expression.

**NOTE 7.21**

For example, the value of ('AB' // 'CDE') // 'F' is the string 'ABCDEF'. Also, the value of 'AB' // ('CDE' // 'F') is the string 'ABCDEF'.

**7.1.5.3.2 Evaluation of the character intrinsic operation**

- 1 A processor is only required to evaluate as much of the character intrinsic operation as is required by the context in which the expression appears.

**NOTE 7.22**

For example, the statements

```
CHARACTER (LEN = 2) C1, C2, C3, CF
C1 = C2 // CF (C3)
```

do not require the function CF to be evaluated, because only the value of C2 is needed to determine the value of C1 because C1 and C2 both have a length of 2.

**7.1.5.4 Logical intrinsic operations****7.1.5.4.1 Interpretation of logical intrinsic operations**

- 1 A logical operation is used to express a logical computation. Evaluation of a logical operation produces a result of type logical. The permitted types for operands of the logical intrinsic operations are specified in [7.1.5.1](#).

- 1 2 The logical operators and their interpretation when used to form an expression are given in Table 7.5, where  $x_1$   
2 denotes the operand to the left of the operator and  $x_2$  denotes the operand to the right of the operator.

Table 7.5: Interpretation of the logical intrinsic operators

Operator	Representing	Use of operator	Interpretation
.NOT.	Logical negation	.NOT. $x_2$	True if $x_2$ is false
.AND.	Logical conjunction	$x_1$ .AND. $x_2$	True if $x_1$ and $x_2$ are both true
.OR.	Logical inclusive disjunction	$x_1$ .OR. $x_2$	True if $x_1$ and/or $x_2$ is true
.EQV.	Logical equivalence	$x_1$ .EQV. $x_2$	True if both $x_1$ and $x_2$ are true or both are false
.NEQV.	Logical nonequivalence	$x_1$ .NEQV. $x_2$	True if either $x_1$ or $x_2$ is true, but not both

- 3 3 The values of the logical intrinsic operations are shown in Table 7.6.

Table 7.6: The values of operations involving logical intrinsic operators

$x_1$	$x_2$	.NOT. $x_2$	$x_1$ .AND. $x_2$	$x_1$ .OR. $x_2$	$x_1$ .EQV. $x_2$	$x_1$ .NEQV. $x_2$
true	true	false	true	true	true	false
true	false	true	false	true	false	true
false	true	false	false	true	false	true
false	false	true	false	false	true	false

#### 4 7.1.5.4.2 Evaluation of logical intrinsic operations

- 5 1 Once the interpretation of a logical intrinsic operation is established, the processor may evaluate any other  
6 expression that is logically equivalent, provided that the integrity of parentheses in any expression is not violated.

##### NOTE 7.23

For example, for the variables L1, L2, and L3 of type logical, the processor could choose to evaluate the expression

L1 .AND. L2 .AND. L3

as

L1 .AND. (L2 .AND. L3)

- 7 2 Two expressions of type logical are logically equivalent if their values are equal for all possible values of their  
8 primaries.

#### 9 7.1.5.5 Relational intrinsic operations

##### 10 7.1.5.5.1 Interpretation of relational intrinsic operations

- 11 1 A relational intrinsic operation is used to compare values of two operands using the relational intrinsic operators  
12 .LT., .LE., .GT., .GE., .EQ., .NE., <, <=, >, >=, ==, and /=. The permitted types for operands of the  
13 relational intrinsic operators are specified in 7.1.5.1.
- 14 2 The operators <, <=, >, >=, ==, and /= always have the same interpretations as the operators .LT., .LE.,  
15 .GT., .GE., .EQ., and .NE., respectively.

##### NOTE 7.24

As shown in Table 7.2, a relational intrinsic operator cannot be used to compare the value of an expression of a numeric type with one of type character or logical. Also, two operands of type logical cannot be

**NOTE 7.24 (cont.)**

compared, a complex operand can be compared with another numeric operand only when the operator is `.EQ.`, `.NE.`, `==`, or `/=`, and two character operands cannot be compared unless they have the same **kind type parameter** value.

- 1    3 Evaluation of a relational intrinsic operation produces a default logical result.
- 2    4 The interpretation of the relational intrinsic operators is given in Table 7.7, where  $x_1$  denotes the operand to the
- 3    left of the operator and  $x_2$  denotes the operand to the right of the operator.

Table 7.7: **Interpretation of the relational intrinsic operators**

Operator	Representing	Use of operator	Interpretation
<code>.LT.</code>	Less than	$x_1$ <code>.LT.</code> $x_2$	$x_1$ less than $x_2$
<code>&lt;</code>	Less than	$x_1 < x_2$	$x_1$ less than $x_2$
<code>.LE.</code>	Less than or equal to	$x_1$ <code>.LE.</code> $x_2$	$x_1$ less than or equal to $x_2$
<code>&lt;=</code>	Less than or equal to	$x_1 <= x_2$	$x_1$ less than or equal to $x_2$
<code>.GT.</code>	Greater than	$x_1$ <code>.GT.</code> $x_2$	$x_1$ greater than $x_2$
<code>&gt;</code>	Greater than	$x_1 > x_2$	$x_1$ greater than $x_2$
<code>.GE.</code>	Greater than or equal to	$x_1$ <code>.GE.</code> $x_2$	$x_1$ greater than or equal to $x_2$
<code>&gt;=</code>	Greater than or equal to	$x_1 >= x_2$	$x_1$ greater than or equal to $x_2$
<code>.EQ.</code>	Equal to	$x_1$ <code>.EQ.</code> $x_2$	$x_1$ equal to $x_2$
<code>==</code>	Equal to	$x_1 == x_2$	$x_1$ equal to $x_2$
<code>.NE.</code>	Not equal to	$x_1$ <code>.NE.</code> $x_2$	$x_1$ not equal to $x_2$
<code>/=</code>	Not equal to	$x_1 /= x_2$	$x_1$ not equal to $x_2$

- 4    5 A numeric relational intrinsic operation is interpreted as having the logical value true if and only if the values of
- 5    the operands satisfy the relation specified by the operator.
- 6    6 In the numeric relational operation
- 7     $x_1$  *rel-op*  $x_2$
- 8    7 if the types or **kind type parameters** of  $x_1$  and  $x_2$  differ, their values are converted to the type and **kind type**
- 9    **parameter** of the expression  $x_1 + x_2$  before evaluation.
- 10    8 A character relational intrinsic operation is interpreted as having the logical value true if and only if the values
- 11    of the operands satisfy the relation specified by the operator.
- 12    9 For a character relational intrinsic operation, the operands are compared one character at a time in order,
- 13    beginning with the first character of each character operand. If the operands are of unequal length, the shorter
- 14    operand is treated as if it were extended on the right with blanks to the length of the longer operand. If both
- 15     $x_1$  and  $x_2$  are of zero length,  $x_1$  is equal to  $x_2$ ; if every character of  $x_1$  is the same as the character in the
- 16    corresponding position in  $x_2$ ,  $x_1$  is equal to  $x_2$ . Otherwise, at the first position where the character operands
- 17    differ, the character operand  $x_1$  is considered to be less than  $x_2$  if the character value of  $x_1$  at this position
- 18    precedes the value of  $x_2$  in the **collating sequence** (1.3.32);  $x_1$  is greater than  $x_2$  if the character value of  $x_1$  at
- 19    this position follows the value of  $x_2$  in the **collating sequence**.

**NOTE 7.25**

The **collating sequence** depends partially on the processor; however, the result of the use of the operators `.EQ.`, `.NE.`, `==`, and `/=` does not depend on the **collating sequence**.

For nondefault character kinds, the blank padding character is processor dependent.

## 7.1.5.5.2 Evaluation of relational intrinsic operations

- 1 Once the interpretation of a relational intrinsic operation is established, the processor may evaluate any other expression that is relationally equivalent, provided that the integrity of parentheses in any expression is not violated.
- 2 Two relational intrinsic operations are relationally equivalent if their logical values are equal for all possible values of their primaries.

## 7.1.6 Defined operations

### 7.1.6.1 Definitions

- 1 A **defined operation** is either a unary operation or a binary operation. A unary **defined operation** is an operation that has the form *defined-unary-op*  $x_2$  or *intrinsic-operator*  $x_2$  and that is defined by a function and a **generic interface** (4.5.5, 12.4.3.5).
- 2 A function defines the unary operation *op*  $x_2$  if
  - (1) the function is specified with a **FUNCTION** (12.6.2.2) or **ENTRY** (12.6.2.6) statement that specifies one **dummy argument**  $d_2$ ,
  - (2) either
    - (a) a **generic interface** (12.4.3.2) provides the function with a *generic-spec* of **OPERATOR** (*op*), or
    - (b) there is a generic **binding** (4.5.5) in the **declared type** of  $x_2$  with a *generic-spec* of **OPERATOR** (*op*) and there is a corresponding **binding** to the function in the **dynamic type** of  $x_2$ ,
  - (3) the type of  $d_2$  is compatible with the **dynamic type** of  $x_2$ ,
  - (4) the type parameters, if any, of  $d_2$  match the corresponding type parameters of  $x_2$ , and
  - (5) either
    - (a) the **rank** of  $x_2$  matches that of  $d_2$  or
    - (b) the function is **elemental** and there is no other function that defines the operation.
- 3 If  $d_2$  is an array, the shape of  $x_2$  shall match the shape of  $d_2$ .
- 4 A binary **defined operation** is an operation that has the form  $x_1$  *defined-binary-op*  $x_2$  or  $x_1$  *intrinsic-operator*  $x_2$  and that is defined by a function and a **generic interface**.
- 5 A function defines the binary operation  $x_1$  *op*  $x_2$  if
  - (1) the function is specified with a **FUNCTION** (12.6.2.2) or **ENTRY** (12.6.2.6) statement that specifies two **dummy arguments**,  $d_1$  and  $d_2$ ,
  - (2) either
    - (a) a **generic interface** (12.4.3.2) provides the function with a *generic-spec* of **OPERATOR** (*op*), or
    - (b) there is a generic **binding** (4.5.5) in the **declared type** of  $x_1$  or  $x_2$  with a *generic-spec* of **OPERATOR** (*op*) and there is a corresponding **binding** to the function in the **dynamic type** of  $x_1$  or  $x_2$ , respectively,
  - (3) the types of  $d_1$  and  $d_2$  are compatible with the **dynamic types** of  $x_1$  and  $x_2$ , respectively,
  - (4) the type parameters, if any, of  $d_1$  and  $d_2$  match the corresponding type parameters of  $x_1$  and  $x_2$ , respectively, and
  - (5) either
    - (a) the **ranks** of  $x_1$  and  $x_2$  match those of  $d_1$  and  $d_2$  or
    - (b) the function is **elemental**,  $x_1$  and  $x_2$  are **conformable**, and there is no other function that defines the operation.

- 1 6 If  $d_1$  or  $d_2$  is an array, the shapes of  $x_1$  and  $x_2$  shall match the shapes of  $d_1$  and  $d_2$ , respectively.

**NOTE 7.26**

An intrinsic operator can be used as the operator in a [defined operation](#). In such a case, the generic properties of the operator are extended.

2 **7.1.6.2 Interpretation of a defined operation**

- 3 1 The interpretation of a [defined operation](#) is provided by the function that defines the operation.
- 4 2 The operators  $<$ ,  $<=$ ,  $>$ ,  $>=$ ,  $=$ , and  $\neq$  always have the same interpretations as the operators [.LT.](#), [.LE.](#),  
5 [.GT.](#), [.GE.](#), [.EQ.](#), and [.NE.](#), respectively.

6 **7.1.6.3 Evaluation of a defined operation**

- 7 1 Once the interpretation of a [defined operation](#) is established, the processor may evaluate any other expression  
8 that is equivalent, provided that the integrity of parentheses is not violated.
- 9 2 Two expressions of derived type are equivalent if their values are equal for all possible values of their primaries.

10 **7.1.7 Evaluation of operands**

- 11 1 It is not necessary for a processor to evaluate all of the operands of an expression, or to evaluate entirely each  
12 operand, if the value of the expression can be determined otherwise.

**NOTE 7.27**

This principle is most often applicable to logical expressions, zero-sized arrays, and zero-length strings, but it applies to all expressions.

For example, in evaluating the expression

$$X > Y \text{ .OR. } L(Z)$$

where  $X$ ,  $Y$ , and  $Z$  are real and  $L$  is a function of type logical, the function reference  $L(Z)$  need not be evaluated if  $X$  is greater than  $Y$ . Similarly, in the array expression

$$W(Z) + A$$

where  $A$  is of size zero and  $W$  is a function, the function reference  $W(Z)$  need not be evaluated.

- 13 2 If a statement contains a function reference in a part of an expression that need not be evaluated, all entities that  
14 would have become defined in the execution of that reference become undefined at the completion of evaluation  
15 of the expression containing the function reference.

**NOTE 7.28**

In the examples in Note [7.27](#), if  $L$  or  $W$  defines its argument, evaluation of the expressions under the specified conditions causes  $Z$  to become undefined, no matter whether or not  $L(Z)$  or  $W(Z)$  is evaluated.

- 16 3 If a statement contains a function reference in a part of an expression that need not be evaluated, no invocation  
17 of that function in that part of the expression shall execute an [image control statement](#) other than [CRITICAL](#)  
18 or [END CRITICAL](#).

**NOTE 7.29**

This restriction is intended to avoid inadvertent deadlock caused by optimization.

## 7.1.8 Integrity of parentheses

- 1 The rules for evaluation specified in subclause 7.1.5 state certain conditions under which a processor may evaluate an expression that is different from the one specified by applying the rules given in 7.1.2 and rules for interpretation specified in subclause 7.1.5. However, any expression in parentheses shall be treated as a data entity.

### NOTE 7.30

For example, in evaluating the expression  $A + (B - C)$  where A, B, and C are of **numeric types**, the difference of B and C shall be evaluated before the addition operation is performed; the processor shall not evaluate the mathematically equivalent expression  $(A + B) - C$ .

## 7.1.9 Type, type parameters, and shape of an expression

### 7.1.9.1 General

- 1 The type, type parameters, and shape of an expression depend on the operators and on the types, type parameters, and shapes of the primaries used in the expression, and are determined recursively from the syntactic form of the expression. The type of an expression is one of the intrinsic types (4.4) or a derived type (4.5).
- 2 If an expression is a **polymorphic** primary or **defined operation**, the type parameters and the **declared** and **dynamic** types of the expression are the same as those of the primary or **defined operation**. Otherwise the type parameters and **dynamic type** of the expression are the same as its **declared type** and type parameters; they are referred to simply as the type and type parameters of the expression.

R724 *logical-expr* is *expr*

C706 (R724) *logical-expr* shall be of type logical.

R725 *default-char-expr* is *expr*

C707 (R725) *default-char-expr* shall be default character.

R726 *int-expr* is *expr*

C708 (R726) *int-expr* shall be of type integer.

R727 *numeric-expr* is *expr*

C709 (R727) *numeric-expr* shall be of type integer, real, or complex.

### 7.1.9.2 Type, type parameters, and shape of a primary

- 1 The type, **type parameters**, and shape of a primary are determined according to whether the primary is a **literal constant**, **designator**, array constructor, **structure constructor**, function reference, **type parameter inquiry**, type parameter name, or parenthesized expression. If a primary is a **literal constant**, its type, **type parameters**, and shape are those of the **literal constant**. If it is a **structure constructor**, it is scalar and its type and **type parameters** are as described in 4.5.10. If it is an array constructor, its type, **type parameters**, and shape are as described in 4.8. If it is a **designator** or function reference, its type, **type parameters**, and shape are those of the **designator** (5.2, 5.5) or the function reference (12.5.3), respectively. If the function reference is generic (12.4.3.2, 13.7) then its type, **type parameters**, and shape are those of the specific function referenced, which is determined by the types, **type parameters**, and **ranks** of its **actual arguments** as specified in 12.5.5.2. If it is a **type parameter inquiry** or type parameter name, it is a scalar integer with the kind of the type parameter.
- 2 If a primary is a parenthesized expression, its type, **type parameters**, and shape are those of the expression.
- 3 The associated **target** object is referenced if a pointer appears as
- a primary in an intrinsic or **defined operation**,
  - the *expr* of a parenthesized primary, or

- the only primary on the right-hand side of an [intrinsic assignment statement](#).

4 The type, type parameters, and shape of the primary are those of the [target](#). If the pointer is not associated with a [target](#), it may appear as a primary only as an [actual argument](#) in a reference to a procedure whose corresponding [dummy argument](#) is declared to be a pointer, or as the target in a [pointer assignment statement](#).

5 A [disassociated array pointer](#) or an unallocated [allocatable](#) array has no shape but does have [rank](#). The type, type parameters, and [rank](#) of the result of the intrinsic function [NULL](#) (13.9.144) depend on context.

### 7.1.9.3 Type, type parameters, and shape of the result of an operation

1 The type of the result of an intrinsic operation  $[x_1] \text{ op } x_2$  is specified by Table 7.2. The shape of the result of an intrinsic operation is the shape of  $x_2$  if  $\text{op}$  is unary or if  $x_1$  is scalar, and is the shape of  $x_1$  otherwise.

2 The type, type parameters, and shape of the result of a defined operation  $[x_1] \text{ op } x_2$  are specified by the function defining the operation (7.1.6).

3 An expression of an intrinsic type has a [kind type parameter](#). An expression of type character also has a character length parameter.

4 The type parameters of the result of an intrinsic operation are as follows.

- For an expression  $x_1 // x_2$  where  $//$  is the character intrinsic operator and  $x_1$  and  $x_2$  are of type character, the character length parameter is the sum of the lengths of the operands and the [kind type parameter](#) is the [kind type parameter](#) of  $x_1$ , which shall be the same as the [kind type parameter](#) of  $x_2$ .
- For an expression  $\text{op } x_2$  where  $\text{op}$  is an intrinsic unary operator and  $x_2$  is of type integer, real, complex, or logical, the [kind type parameter](#) of the expression is that of the operand.
- For an expression  $x_1 \text{ op } x_2$  where  $\text{op}$  is a numeric intrinsic binary operator with one operand of type integer and the other of type real or complex, the [kind type parameter](#) of the expression is that of the real or complex operand.
- For an expression  $x_1 \text{ op } x_2$  where  $\text{op}$  is a numeric intrinsic binary operator with both operands of the same type and [kind type parameters](#), or with one real and one complex with the same [kind type parameters](#), the [kind type parameter](#) of the expression is identical to that of each operand. In the case where both operands are integer with different [kind type parameters](#), the [kind type parameter](#) of the expression is that of the operand with the greater decimal exponent range if the decimal exponent ranges are different; if the decimal exponent ranges are the same, the [kind type parameter](#) of the expression is processor dependent, but it is the same as that of one of the operands. In the case where both operands are any of type real or complex with different [kind type parameters](#), the [kind type parameter](#) of the expression is that of the operand with the greater decimal precision if the decimal precisions are different; if the decimal precisions are the same, the [kind type parameter](#) of the expression is processor dependent, but it is the same as that of one of the operands.
- For an expression  $x_1 \text{ op } x_2$  where  $\text{op}$  is a logical intrinsic binary operator with both operands of the same [kind type parameter](#), the [kind type parameter](#) of the expression is identical to that of each operand. In the case where both operands are of type logical with different [kind type parameters](#), the [kind type parameter](#) of the expression is processor dependent, but it is the same as that of one of the operands.
- For an expression  $x_1 \text{ op } x_2$  where  $\text{op}$  is a relational intrinsic operator, the expression has the default logical [kind type parameter](#).

### 7.1.10 Conformability rules for elemental operations

1 An [elemental operation](#) is an intrinsic operation or a [defined operation](#) for which the function is [elemental](#) (12.8).

2 For all [elemental](#) binary operations, the two operands shall be conformable. In the case where one is a scalar and the other an array, the scalar is treated as if it were an array of the same shape as the array operand with every element, if any, of the array equal to the value of the scalar.

### 7.1.11 Specification expression

1 A **specification expression** is an expression with limitations that make it suitable for use in specifications such as length type parameters (C404) and array bounds (R517, R518). A *specification-expr* shall be a **constant expression** unless it is in an **interface body** (12.4.3.2), the specification part of a subprogram or **BLOCK construct**, a derived type definition, or the *declaration-type-spec* of a **FUNCTION statement** (12.6.2.2).

R728 *specification-expr*                      is    *scalar-int-expr*

C710 (R728) The *scalar-int-expr* shall be a restricted expression.

2 A restricted expression is an expression in which each operation is intrinsic or defined by a specification function and each primary is

- (1) a constant or subobject of a constant,
- (2) an **object designator** with a **base object** that is a **dummy argument** that has neither the **OPTIONAL** nor the **INTENT (OUT)** attribute,
- (3) an **object designator** with a **base object** that is in a **common block**,
- (4) an **object designator** with a **base object** that is made accessible by use or **host** association,
- (5) an **object designator** with a **base object** that is a **local variable** of the procedure containing the **BLOCK construct** in which the restricted expression appears,
- (6) an **object designator** with a **base object** that is a **local variable** of an outer **BLOCK construct** containing the **BLOCK construct** in which the restricted expression appears,
- (7) an array constructor where each element and each *scalar-int-expr* of each *ac-implied-do-control* is a restricted expression,
- (8) a **structure constructor** where each component is a restricted expression,
- (9) a specification inquiry where each **designator** or argument is
  - (a) a restricted expression or
  - (b) a variable that is not an optional dummy argument, and whose properties inquired about are not
    - (i) dependent on the upper bound of the last dimension of an **assumed-size array**,
    - (ii) deferred, or
    - (iii) defined by an expression that is not a restricted expression,
- (10) a specification inquiry that is a constant expression,
- (11) a reference to the intrinsic function **PRESENT**,
- (12) a reference to any other standard intrinsic function where each argument is a restricted expression,
- (13) a reference to a specification function where each argument is a restricted expression,
- (14) a type parameter of the derived type being defined,
- (15) an *ac-do-variable* within an array constructor where each *scalar-int-expr* of the corresponding *ac-implied-do-control* is a restricted expression, or
- (16) a restricted expression enclosed in parentheses,

3 where each subscript, section subscript, substring starting point, substring ending point, and type parameter value is a restricted expression, and where any **final subroutine** that is invoked is pure.

4 A specification inquiry is a reference to

- (1) an intrinsic **inquiry function** other than **PRESENT**,
- (2) a **type parameter inquiry** (6.4.5),
- (3) an **inquiry function** from the intrinsic modules **IEEE\_ARITHMETIC** and **IEEE\_EXCEPTIONS** (14.10),
- (4) the function **C\_SIZEOF** from the intrinsic module **ISO\_C\_BINDING** (15.2.3.7), or
- (5) the **COMPILER\_VERSION** or **COMPILER\_OPTIONS** function from the intrinsic module **ISO\_FORTRAN\_ENV** (13.10.2.6, 13.10.2.7).

- 1 5 A function is a specification function if it is a pure function, is not a standard intrinsic function, is not an internal  
 2 function, is not a statement function, and does not have a **dummy procedure** argument.
- 3 6 Evaluation of a **specification expression** shall not directly or indirectly cause a procedure defined by the subpro-  
 4 gram in which it appears to be invoked.

**NOTE 7.31**

Specification functions are nonintrinsic functions that can be used in specification expressions to determine the attributes of data objects. The requirement that they be pure ensures that they cannot have side effects that could affect other objects being declared in the same *specification-part*. The requirement that they not be internal ensures that they cannot inquire, via **host association**, about other objects being declared in the same *specification-part*. The prohibition against recursion avoids the creation of a new instance of a procedure while construction of one is in progress.

- 5 7 A variable in a **specification expression** shall have its type and type parameters, if any, specified by a previous  
 6 declaration in the same **scoping unit**, by the implicit typing rules in effect for the **scoping unit**, or by **host** or **use**  
 7 association. If a variable in a **specification expression** is typed by the implicit typing rules, its appearance in any  
 8 subsequent **type declaration statement** shall confirm the implied type and type parameters.
- 9 8 If a **specification expression** includes a specification inquiry that depends on a type parameter or an array bound  
 10 of an entity specified in the same *specification-part*, the type parameter or array bound shall be specified in a prior  
 11 specification of the *specification-part*. The prior specification may be to the left of the specification inquiry in the  
 12 same statement, but shall not be within the same *entity-decl*. If a **specification expression** includes a reference to  
 13 the value of an element of an array specified in the same *specification-part*, the array shall be completely specified  
 14 in prior declarations.
- 15 9 A generic entity referenced in a **specification expression** in the *specification-part* of a **scoping unit** shall have no  
 16 specific procedures defined in the **scoping unit**, or its **host scoping unit**, subsequent to the **specification expression**.

**NOTE 7.32**

The following are examples of **specification expressions**:

```
LBOUND (B, 1) + 5 ! B is an assumed-shape dummy array
M + LEN (C)       ! M and C are dummy arguments
2 * PRECISION (A) ! A is a real variable made accessible
                  ! by a USE statement
```

## 17 7.1.12 Constant expression

- 18 1 A **constant expression** is an expression with limitations that make it suitable for use as a **kind type parameter**,  
 19 initializer, or **named constant**. It is an expression in which each operation is intrinsic, and each primary is
- 20 (1) a constant or subobject of a constant,
  - 21 (2) an array constructor where each element and each *scalar-int-expr* of each *ac-implied-do-control* is a  
 22 **constant expression**,
  - 23 (3) a **structure constructor** where each *component-spec* corresponding to
    - 24 (a) an **allocatable** component is a reference to the intrinsic function **NULL**,
    - 25 (b) a pointer component is an initialization target or a reference to the intrinsic function **NULL**,  
 26 and
    - 27 (c) any other component is a **constant expression**,
  - 28 (4) a specification inquiry where each **designator** or argument is
    - 29 (a) a **constant expression** or
    - 30 (b) a variable whose properties inquired about are not

- (i) assumed,
- (ii) deferred, or
- (iii) defined by an expression that is not a *constant expression*,

- (5) a reference to an *elemental* standard intrinsic function, where each argument is a *constant expression*,
- (6) a reference to a *transformational* standard intrinsic function other than *COMMAND\_ARGUMENT\_COUNT*, *NULL*, *NUM\_IMAGES*, *THIS\_IMAGE*, or *TRANSFER*, where each argument is a *constant expression*,
- (7) a reference to the intrinsic function *NULL* that does not have an argument with a type parameter that is assumed or is defined by an expression that is not a *constant expression*,
- (8) a reference to the intrinsic function *TRANSFER* where each argument is a *constant expression* and each *ultimate pointer* component of the *SOURCE* argument is *disassociated*,
- (9) a reference to the *transformational* function *IEEE\_SELECTED\_REAL\_KIND* from the intrinsic module *IEEE\_ARITHMETIC*(14), where each argument is a *constant expression*,
- (10) a previously declared *kind type parameter* of the derived type being defined,
- (11) a *data-i-do-variable* within a *data-implied-do*,
- (12) an *ac-do-variable* within an array constructor where each *scalar-int-expr* of the corresponding *ac-implied-do-control* is a *constant expression*, or
- (13) a *constant expression* enclosed in parentheses,

and where each subscript, section subscript, substring starting point, substring ending point, and type parameter value is a *constant expression*.

R729 *constant-expr* is *expr*

C711 (R729) *constant-expr* shall be a *constant expression*.

R730 *default-char-constant-expr* is *default-char-expr*

C712 (R730) *default-char-constant-expr* shall be a *constant expression*.

R731 *int-constant-expr* is *int-expr*

C713 (R731) *int-constant-expr* shall be a *constant expression*.

- 2 If a *constant expression* includes a specification inquiry that depends on a type parameter or an array bound of an entity specified in the same *specification-part*, the type parameter or array bound shall be specified in a prior specification of the *specification-part*. The prior specification may be to the left of the specification inquiry in the same statement, but shall not be within the same *entity-decl* unless the specification inquiry appears within an *initialization*.
- 3 A generic entity referenced in a *constant expression* in the *specification-part* of a *scoping unit* shall have no specific procedures defined in that *scoping unit*, or its *host scoping unit*, subsequent to the *constant expression*.

#### NOTE 7.33

The following are examples of *constant expressions*:

```

3
-3 + 4
'AB'
'AB' // 'CD'
('AB' // 'CD') // 'EF'
SIZE (A)
DIGITS (X) + 4
4.0 * atan(1.0)
ceiling(number_of_decimal_digits / log10(radix(0.0)))

```

**NOTE 7.33 (cont.)**

where A is an [explicit-shape array](#) with constant bounds and X is default real.

**7.2 Assignment****7.2.1 Assignment statement****7.2.1.1 General form**

R732 *assignment-stmt* is *variable* = *expr*

C714 (R732) The *variable* shall not be a whole [assumed-size array](#).

**NOTE 7.34**

Examples of an assignment statement are:

```
A = 3.5 + X * Y
I = INT (A)
```

- 1 An [assignment-stmt](#) shall meet the requirements of either a [defined assignment](#) statement or an intrinsic assignment statement.

**7.2.1.2 Intrinsic assignment statement**

- 1 An intrinsic assignment statement is an assignment statement that is not a [defined assignment](#) statement (7.2.1.4). In an intrinsic assignment statement,

- (1) if the variable is [polymorphic](#) it shall be [allocatable](#) and not a [coarray](#),
- (2) if *expr* is an array then the variable shall also be an array,
- (3) the variable and *expr* shall be [conformable](#) unless the variable is an [allocatable](#) array that has the same [rank](#) as *expr* and is not a [coarray](#),
- (4) if the variable is [polymorphic](#) it shall be [type compatible](#) with *expr*; otherwise the [declared types](#) of the variable and *expr* shall conform as specified in Table 7.8,
- (5) if the variable is of type character and of [ISO 10646](#), [ASCII](#), or default character kind, *expr* shall be of [ISO 10646](#), [ASCII](#), or default character kind,
- (6) otherwise if the variable is of type character *expr* shall have the same [kind type parameter](#),
- (7) if the variable is of derived type each [kind type parameter](#) of the variable shall have the same value as the corresponding [kind type parameter](#) of *expr*, and
- (8) if the variable is of derived type each length type parameter of the variable shall have the same value as the corresponding type parameter of *expr* unless the variable is [allocatable](#), is not a [coarray](#), and its corresponding [type parameter](#) is [deferred](#).

Table 7.8: **Type conformance for the intrinsic assignment statement**

Type of the variable	Type of <i>expr</i>
integer	integer, real, complex
real	integer, real, complex
complex	integer, real, complex
character	character
logical	logical
derived type	same derived type as the variable

- 2 If the variable in an intrinsic assignment statement is a [coindexed object](#),

- the variable shall not be *polymorphic*,
- the variable shall not have an *allocatable* ultimate component,
- the variable shall be *conformable* with *expr*, and
- each *deferred length type parameter* of the variable shall have the same value as the corresponding type parameter of *expr*.

- 3 If the variable is a pointer, it shall be associated with a *definable target* such that the type, type parameters, and shape of the *target* and *expr* conform.

### 7.2.1.3 Interpretation of intrinsic assignments

- 1 Execution of an intrinsic assignment causes, in effect, the evaluation of the expression *expr* and all expressions within *variable* (7.1), the possible conversion of *expr* to the type and type parameters of the variable (Table 7.9), and the definition of the variable with the resulting value. The execution of the assignment shall have the same effect as if the evaluation of *expr* and the evaluation of all expressions in *variable* occurred before any portion of the variable is defined by the assignment. The evaluation of expressions within *variable* shall neither affect nor be affected by the evaluation of *expr*.
- 2 If the variable is a pointer, the value of *expr* is assigned to the *target* of the variable.
- 3 If the variable is an unallocated *allocatable* array, *expr* shall have the same *rank*. If the variable is an allocated *allocatable* variable, it is deallocated if *expr* is an array of different shape, any of the corresponding length type parameter values of the variable and *expr* differ, or the variable is *polymorphic* and the *dynamic type* of the variable and *expr* differ. If the variable is or becomes an unallocated *allocatable* variable, it is then allocated with
- the same *dynamic type* as *expr* if the variable is *polymorphic*,
  - each *deferred type parameter* equal to the corresponding type parameter of *expr*,
  - the same bounds as before if the variable is an array and *expr* is scalar, and
  - the shape of *expr* with each lower bound equal to the corresponding element of *LBOUND* (*expr*) if *expr* is an array.

#### NOTE 7.35

For example, given the declaration

```
CHARACTER(:),ALLOCATABLE :: NAME
```

then after the assignment statement

```
NAME = 'Dr. '//FIRST_NAME//' '//SURNAME
```

NAME will have the length LEN(FIRST\_NAME)+LEN(SURNAME)+5, even if it had previously been unallocated, or allocated with a different length. However, for the assignment statement

```
NAME(:) = 'Dr. '//FIRST_NAME//' '//SURNAME
```

NAME must already be allocated at the time of the assignment; the assigned value is truncated or blank padded to the previously allocated length of NAME.

- 4 Both *variable* and *expr* may contain references to any portion of the variable.

#### NOTE 7.36

For example, in the character intrinsic assignment statement:

```
STRING (2:5) = STRING (1:4)
```

**NOTE 7.36 (cont.)**

the assignment of the first character of STRING to the second character does not affect the evaluation of STRING (1:4). If the value of STRING prior to the assignment was 'ABCDEF', the value following the assignment is 'AABCDF'.

- 1 5 If *expr* is a scalar and the variable is an array, the *expr* is treated as if it were an array of the same shape as the  
 2 variable with every element of the array equal to the scalar value of *expr*.
- 3 6 If the variable is an array, the assignment is performed element-by-element on corresponding array elements of  
 4 the variable and *expr*.

**NOTE 7.37**

For example, if A and B are arrays of the same shape, the array intrinsic assignment

A = B

assigns the corresponding elements of B to those of A; that is, the first element of B is assigned to the first element of A, the second element of B is assigned to the second element of A, etc.

If C is an allocatable array of rank 1, then

C = PACK (ARRAY, ARRAY > 0)

will cause C to contain all the positive elements of ARRAY in array element order; if C is not allocated or is allocated with the wrong size, it will be re-allocated to be of the correct size to hold the result of PACK.

- 5 7 The processor may perform the element-by-element assignment in any order.

**NOTE 7.38**

For example, the following program segment results in the values of the elements of array X being reversed:

```
REAL X (10)
...
X (1:10) = X (10:1:-1)
```

- 6 8 For an intrinsic assignment statement where the variable is of numeric type, the *expr* can have a different numeric  
 7 type or kind type parameter, in which case the value of *expr* is converted to the type and kind type parameter  
 8 of the variable according to the rules of Table 7.9.

Table 7.9: **Numeric conversion and the assignment statement**

Type of the variable	Value Assigned
integer	INT ( <i>expr</i> , KIND = KIND ( <i>variable</i> ))
real	REAL ( <i>expr</i> , KIND = KIND ( <i>variable</i> ))
complex	CMPLX ( <i>expr</i> , KIND = KIND ( <i>variable</i> ))
Note: INT, REAL, CMPLX, and KIND are the generic names of functions defined in 13.9.	

- 9 9 For an intrinsic assignment statement where the variable is of type logical, the *expr* can have a different kind  
 10 type parameter, in which case the value of *expr* is converted to the kind type parameter of the variable.
- 11 10 For an intrinsic assignment statement where the variable is of type character, the *expr* can have a different  
 12 character length parameter in which case the conversion of *expr* to the length of the variable is as follows.
- 13 (1) If the length of the variable is less than that of *expr*, the value of *expr* is truncated from the right  
 14 until it is the same length as the variable.

- (2) If the length of the variable is greater than that of *expr*, the value of *expr* is extended on the right with blanks until it is the same length as the variable.

- 11 For an intrinsic assignment statement where the variable is of type character, if *expr* has a different **kind type parameter**, each character *c* in *expr* is converted to the **kind type parameter** of the variable by **ACHAR** (**IACHAR**(*c*), **KIND** (*variable*)).

#### NOTE 7.39

For nondefault character kinds, the blank padding character is processor dependent. When assigning a character expression to a variable of a different kind, each character of the expression that is not representable in the kind of the variable is replaced by a processor-dependent character.

- 12 For an intrinsic assignment of the type **C\_PTR** or **C\_FUNPTR** from the intrinsic module **ISO\_C\_BINDING**, the variable becomes undefined if the variable and *expr* are not on the same **image**.

#### NOTE 7.40

An intrinsic assignment statement for a variable of type **C\_PTR** or **C\_FUNPTR** is not permitted to involve a **coindexed object**, see C614, which prevents inappropriate copying from one **image** to another. However, such copying can occur as an intrinsic assignment for a component in a derived-type assignment.

- 13 An intrinsic assignment where the variable is of derived type is performed as if each component of the variable were assigned from the corresponding component of *expr* using **pointer assignment** (7.2.2) for each pointer component, **defined assignment** for each nonpointer nonallocatable component of a type that has a **type-bound defined assignment** consistent with the component, intrinsic assignment for each other nonpointer nonallocatable component, and intrinsic assignment for each allocated **coarray** component. For unallocated **coarray** components, the corresponding component of the variable shall be unallocated. For a noncoarray **allocatable** component the following sequence of operations is applied.

- (1) If the component of the variable is allocated, it is deallocated.
- (2) If the component of the value of *expr* is allocated, the corresponding component of the variable is allocated with the same **dynamic type** and type parameters as the component of the value of *expr*. If it is an array, it is allocated with the same bounds. The value of the component of the value of *expr* is then assigned to the corresponding component of the variable using **defined assignment** if the **declared type** of the component has a **type-bound defined assignment** consistent with the component, and intrinsic assignment for the **dynamic type** of that component otherwise.

- 14 The processor may perform the component-by-component assignment in any order or by any means that has the same effect.

#### NOTE 7.41

For an example of a derived-type intrinsic assignment statement, if C and D are of the same derived type with a pointer component P and nonpointer components S, T, U, and V of type integer, logical, character, and another derived type, respectively, the intrinsic

C = D

pointer assigns D%P to C%P. It assigns D%S to C%S, D%T to C%T, and D%U to C%U using intrinsic assignment. It assigns D%V to C%V using **defined assignment** if objects of that type have a compatible **type-bound defined assignment**, and intrinsic assignment otherwise.

#### NOTE 7.42

If an **allocatable** component of *expr* is unallocated, the corresponding component of the variable has an allocation status of unallocated after execution of the assignment.

#### 7.2.1.4 Defined assignment statement

- 1 A **defined assignment** statement is an assignment statement that is defined by a subroutine and a **generic interface** (4.5.5, 12.4.3.5.3) that specifies **ASSIGNMENT** (=).
- 2 A subroutine defines the **defined assignment**  $x_1 = x_2$  if
  - (1) the subroutine is specified with a **SUBROUTINE** (12.6.2.3) or **ENTRY** (12.6.2.6) statement that specifies two **dummy arguments**,  $d_1$  and  $d_2$ ,
  - (2) either
    - (a) a **generic interface** (12.4.3.2) provides the subroutine with a *generic-spec* of **ASSIGNMENT** (=), or
    - (b) there is a generic **binding** (4.5.5) in the **declared type** of  $x_1$  or  $x_2$  with a *generic-spec* of **ASSIGNMENT** (=) and there is a corresponding **binding** to the subroutine in the **dynamic type** of  $x_1$  or  $x_2$ , respectively,
  - (3) the types of  $d_1$  and  $d_2$  are compatible with the **dynamic types** of  $x_1$  and  $x_2$ , respectively,
  - (4) the type parameters, if any, of  $d_1$  and  $d_2$  match the corresponding type parameters of  $x_1$  and  $x_2$ , respectively, and
  - (5) either
    - (a) the **ranks** of  $x_1$  and  $x_2$  match those of  $d_1$  and  $d_2$  or
    - (b) the subroutine is **elemental**,  $x_1$  and  $x_2$  are **conformable**, and there is no other subroutine that defines the assignment.
- 3 If  $d_1$  or  $d_2$  is an array, the shapes of  $x_1$  and  $x_2$  shall match the shapes of  $d_1$  and  $d_2$ , respectively.

#### 7.2.1.5 Interpretation of defined assignment statements

- 1 The interpretation of a **defined assignment** is provided by the subroutine that defines it.
- 2 If the **defined assignment** is an **elemental assignment** and the variable in the assignment is an array, the **defined assignment** is performed element-by-element, on corresponding elements of the variable and *expr*. If *expr* is a scalar, it is treated as if it were an array of the same shape as the variable with every element of the array equal to the scalar value of *expr*.

##### NOTE 7.43

The rules of **defined assignment** (12.4.3.5.3), procedure references (12.5), subroutine references (12.5.4), and **elemental** subroutine arguments (12.8.3) ensure that the **defined assignment** has the same effect as if the evaluation of all operations in  $x_2$  and  $x_1$  occurs before any portion of  $x_1$  is defined. If an **elemental assignment** is defined by a pure **elemental** subroutine, the element assignments can be performed simultaneously or in any order.

### 7.2.2 Pointer assignment

#### 7.2.2.1 General

- 1 **Pointer assignment** causes a pointer to become associated with a **target** or causes its **pointer association** status to become **disassociated** or undefined. Any previous association between the pointer and a **target** is broken.
- 2 **Pointer assignment** for a pointer component of a structure may also take place by execution of a derived-type **intrinsic assignment statement** (7.2.1.3).

#### 7.2.2.2 Syntax of the pointer assignment statement

R733 *pointer-assignment-stmt*      **is** *data-pointer-object* [ (*bounds-spec-list*) ] => *data-target*  
    **or** *data-pointer-object* (*bounds-remapping-list*) => *data-target*  
    **or** *proc-pointer-object* => *proc-target*

- 1 R734 *data-pointer-object* is *variable-name*  
 2 or *scalar-variable* % *data-pointer-component-name*
- 3 C715 (R733) If *data-target* is not unlimited polymorphic, *data-pointer-object* shall be type compatible (4.3.2.3)  
 4 with it and the corresponding kind type parameters shall be equal.
- 5 C716 (R733) If *data-target* is unlimited polymorphic, *data-pointer-object* shall be unlimited polymorphic, or of  
 6 a type with the BIND attribute or the SEQUENCE attribute.
- 7 C717 (R733) If *bounds-spec-list* is specified, the number of *bounds-specs* shall equal the rank of *data-pointer-*  
 8 *object*.
- 9 C718 (R733) If *bounds-remapping-list* is specified, the number of *bounds-remappings* shall equal the rank of  
 10 *data-pointer-object*.
- 11 C719 (R733) If *bounds-remapping-list* is not specified, the ranks of *data-pointer-object* and *data-target* shall be  
 12 the same.
- 13 C720 (R733) A coarray *data-target* shall have the VOLATILE attribute if and only if the *data-pointer-object*  
 14 has the VOLATILE attribute.
- 15 C721 (R734) A *variable-name* shall have the POINTER attribute.
- 16 C722 (R734) A *scalar-variable* shall be a *data-ref*.
- 17 C723 (R734) A *data-pointer-component-name* shall be the name of a component of *scalar-variable* that is a  
 18 data pointer.
- 19 C724 (R734) A *data-pointer-object* shall not be a coindexed object.
- 20 R735 *bounds-spec* is *lower-bound-expr* :  
 21 R736 *bounds-remapping* is *lower-bound-expr* : *upper-bound-expr*
- 22 R737 *data-target* is *expr*
- 23 C725 (R737) The *expr* shall be a *designator* that designates a variable with either the TARGET or POINTER  
 24 attribute and is not an array section with a vector subscript, or it shall be a reference to a function that  
 25 returns a data pointer.
- 26 C726 (R737) A *data-target* shall not be a coindexed object.

**NOTE 7.44**

A data pointer and its target are always on the same image. A coarray can be of a derived type with pointer or allocatable subcomponents. For example, if PTR is a pointer component, Z[P]%PTR is a reference to the target of component PTR of Z on image P. This target is on image P and its association with Z[P]%PTR must have been established by the execution of an ALLOCATE statement or a pointer assignment on image P.

- 27 R738 *proc-pointer-object* is *proc-pointer-name*  
 28 or *proc-component-ref*
- 29 R739 *proc-component-ref* is *scalar-variable* % *procedure-component-name*
- 30 C727 (R739) The *scalar-variable* shall be a *data-ref* that is not a coindexed object.
- 31 C728 (R739) The *procedure-component-name* shall be the name of a procedure pointer component of the de-  
 32 clared type of *scalar-variable*.
- 33 R740 *proc-target* is *expr*

or *procedure-name*  
or *proc-component-ref*

C729 (R740) An *expr* shall be a reference to a function whose result is a procedure pointer.

C730 (R740) A *procedure-name* shall be the name of an *internal*, *module*, or *dummy* procedure, a *procedure pointer*, a *specific* intrinsic function listed in Table 13.2, or an *external procedure* that is accessed by *use* or *host* association, referenced in the *scoping unit* as a procedure, or that has the *EXTERNAL* attribute.

C731 (R740) The *proc-target* shall not be a nonintrinsic *elemental procedure*.

1 In a pointer assignment statement, *data-pointer-object* or *proc-pointer-object* denotes the pointer object and *data-target* or *proc-target* denotes the pointer target.

2 For pointer assignment performed by a derived-type *intrinsic assignment statement*, the pointer object is the pointer component of the variable and the pointer target is the corresponding component of *expr*.

### 7.2.2.3 Data pointer assignment

1 If the pointer object is not polymorphic (4.3.2.3) and the pointer target is *polymorphic* with *dynamic type* that differs from its *declared type*, the assignment target is the ancestor component of the pointer target that has the type of the pointer object. Otherwise, the assignment target is the pointer target.

2 If the pointer target is not a pointer, the pointer object becomes *pointer associated* with the assignment target; if the pointer target is a pointer with a target that is not on the same *image*, the *pointer association* status of the pointer object becomes undefined. Otherwise, the *pointer association* status of the pointer object becomes that of the pointer target; if the pointer target is associated with an object, the pointer object becomes associated with the assignment target. If the pointer target is *allocatable*, it shall be allocated.

#### NOTE 7.45

A *pointer assignment statement* is not permitted to involve a *coindexed* pointer or target, see C724 and C726. This prevents a pointer assignment statement from associating a pointer with a target on another *image*. If such an association would otherwise be implied, the association status of the pointer becomes undefined. For example, a derived-type intrinsic assignment where the variable and *expr* are on different images and the variable has an ultimate pointer component.

3 If the pointer object is *polymorphic*, it assumes the *dynamic type* of the pointer target. If the pointer object is of a type with the *BIND* attribute or the *SEQUENCE* attribute, the *dynamic type* of the pointer target shall be that type.

4 If the pointer target is a *disassociated* pointer, all nondeferred type parameters of the *declared type* of the pointer object that correspond to nondeferred type parameters of the pointer target shall have the same values as the corresponding type parameters of the pointer target.

5 Otherwise, all nondeferred type parameters of the *declared type* of the pointer object shall have the same values as the corresponding type parameters of the pointer target.

6 If the pointer object has nondeferred type parameters that correspond to *deferred type parameters* of the pointer target, the pointer target shall not be a pointer with undefined association status.

7 If the pointer object has the *CONTIGUOUS* attribute, the pointer target shall be *contiguous*.

8 If the target of a pointer is a *coarray*, the pointer shall have the *VOLATILE* attribute if and only if the *coarray* has the *VOLATILE* attribute.

9 If *bounds-remapping-list* appears, it specifies the upper and lower *bounds* of each dimension of the pointer, and thus the *extents*; the pointer target shall be *simply contiguous* (6.5.4) or of *rank* one, and shall not be a *disassociated* or *undefined* pointer. The number of elements of the pointer target shall not be less than the

1 number implied by the *bounds-remapping-list*. The elements of the pointer object are associated with those of  
 2 the pointer *target*, in array element order; if the pointer target has more elements than specified for the pointer  
 3 object, the remaining elements are not associated with the pointer object.

4 10 If no *bounds-remapping-list* appears, the extent of a dimension of the pointer object is the extent of the corres-  
 5 ponding dimension of the pointer target. If *bounds-spec-list* appears, it specifies the lower bounds; otherwise,  
 6 the lower bound of each dimension is the result of the intrinsic function **LBOUND** (13.9.109) applied to the  
 7 corresponding dimension of the pointer target. The upper bound of each dimension is one less than the sum of  
 8 the lower bound and the extent.

#### 9 7.2.2.4 Procedure pointer assignment

10 1 If the pointer target is not a pointer, the pointer object becomes *pointer associated* with the pointer target.  
 11 Otherwise, the *pointer association* status of the pointer object becomes that of the pointer target; if the pointer  
 12 target is associated with a procedure, the pointer object becomes associated with the same procedure.

13 2 The *host instance* (12.6.2.4) of an associated *procedure pointer* is the *host instance* of its target.

14 3 If the pointer object has an *explicit interface*, its *characteristics* shall be the same as the pointer target except  
 15 that the pointer target may be pure even if the pointer object is not pure and the pointer target may be an  
 16 *elemental* intrinsic procedure even if the pointer object is not *elemental*.

17 4 If the *characteristics* of the pointer object or the pointer target are such that an *explicit interface* is required,  
 18 both the pointer object and the pointer target shall have an *explicit interface*.

19 5 If the pointer object has an *implicit interface* and is explicitly typed or referenced as a function, the pointer target  
 20 shall be a function. If the pointer object has an *implicit interface* and is referenced as a subroutine, the pointer  
 21 target shall be a subroutine.

22 6 If the pointer object is a function with an *implicit interface*, the pointer target shall be a function with the same  
 23 type; corresponding type parameters shall have the same value.

24 7 If *procedure-name* is a specific procedure name that is also a generic name, only the specific procedure is associated  
 25 with the pointer object.

#### 26 7.2.2.5 Examples

##### NOTE 7.46

The following are examples of pointer assignment statements. (See Note 12.18 for declarations of P and BESSEL.)

```
NEW_NODE % LEFT => CURRENT_NODE
SIMPLE_NAME => TARGET_STRUCTURE % SUBSTRUCT % COMPONENT
PTR => NULL ( )
ROW => MAT2D (N, :)
WINDOW => MAT2D (I-1:I+1, J-1:J+1)
POINTER_OBJECT => POINTER_FUNCTION (ARG_1, ARG_2)
EVERY_OTHER => VECTOR (1:N:2)
WINDOW2 (0:, 0:) => MAT2D (ML:MU, NL:NU)
! P is a procedure pointer and BESSEL is a procedure with a
! compatible interface.
P => BESSEL

! Likewise for a structure component.
STRUCT % COMPONENT => BESSEL
```

**NOTE 7.47**

It is possible to obtain different-rank views of parts of an object by specifying upper bounds in pointer assignment statements. This requires that the object be either [rank](#) one or [contiguous](#). Consider the following example, in which a matrix is under consideration. The matrix is stored as a rank-one object in MYDATA because its diagonal is needed for some reason – the diagonal cannot be gotten as a single object from a rank-two representation. The matrix is represented as a rank-two view of MYDATA.

```
real, target :: MYDATA ( NR*NC )      ! An automatic array
real, pointer :: MATRIX ( :, : )      ! A rank-two view of MYDATA
real, pointer :: VIEW_DIAG ( : )
MATRIX( 1:NR, 1:NC ) => MYDATA        ! The MATRIX view of the data
VIEW_DIAG => MYDATA( 1::NR+1 )        ! The diagonal of MATRIX
```

Rows, columns, or blocks of the matrix can be accessed as sections of MATRIX.

[Rank](#) remapping can be applied to CONTIGUOUS arrays, for example:

```
REAL, CONTIGUOUS, POINTER :: A(:)
REAL, CONTIGUOUS, TARGET :: B(:, :) ! Dummy argument
A(1:SIZE(B)) => B                    ! Linear view of a rank-2 array
```

**7.2.3 Masked array assignment – WHERE****7.2.3.1 General form of the masked array assignment**

- 1 A [masked array assignment](#) is either a WHERE statement or a WHERE construct. It is used to mask the evaluation of expressions and assignment of values in array assignment statements, according to the value of a logical array expression.

R741 *where-stmt* is WHERE ( *mask-expr* ) *where-assignment-stmt*

R742 *where-construct* is *where-construct-stmt*  
                                   [ *where-body-construct* ] ...  
                                   [ *masked-elsewhere-stmt*  
                                   [ *where-body-construct* ] ... ] ...  
                                   [ *elsewhere-stmt*  
                                   [ *where-body-construct* ] ... ]  
                                   *end-where-stmt*

R743 *where-construct-stmt* is [ *where-construct-name* ] WHERE ( *mask-expr* )

R744 *where-body-construct* is *where-assignment-stmt*  
                                   or *where-stmt*  
                                   or *where-construct*

R745 *where-assignment-stmt* is *assignment-stmt*

R746 *mask-expr* is *logical-expr*

R747 *masked-elsewhere-stmt* is ELSEWHERE ( *mask-expr* ) [ *where-construct-name* ]

R748 *elsewhere-stmt* is ELSEWHERE [ *where-construct-name* ]

R749 *end-where-stmt* is END WHERE [ *where-construct-name* ]

C732 (R745) A *where-assignment-stmt* that is a [defined assignment](#) shall be [elemental](#).

C733 (R742) If the *where-construct-stmt* is identified by a *where-construct-name*, the corresponding *end-where-*

*stmt* shall specify the same *where-construct-name*. If the *where-construct-stmt* is not identified by a *where-construct-name*, the corresponding *end-where-stmt* shall not specify a *where-construct-name*. If an *elsewhere-stmt* or a *masked-elsewhere-stmt* is identified by a *where-construct-name*, the corresponding *where-construct-stmt* shall specify the same *where-construct-name*.

C734 (R744) A statement that is part of a *where-body-construct* shall not be a *branch target statement*.

- 2 If a *where-construct* contains a *where-stmt*, a *masked-elsewhere-stmt*, or another *where-construct* then each *mask-expr* within the *where-construct* shall have the same shape. In each *where-assignment-stmt*, the *mask-expr* and the variable being defined shall be arrays of the same shape.

#### NOTE 7.48

Examples of a masked array assignment are:

```
WHERE (TEMP > 100.0) TEMP = TEMP - REDUCE_TEMP
WHERE (PRESSURE <= 1.0)
    PRESSURE = PRESSURE + INC_PRESSURE
    TEMP = TEMP - 5.0
ELSEWHERE
    RAINING = .TRUE.
END WHERE
```

#### 7.2.3.2 Interpretation of masked array assignments

- 1 When a WHERE statement or a *where-construct-stmt* is executed, a control mask is established. In addition, when a WHERE construct statement is executed, a pending control mask is established. If the statement does not appear as part of a *where-body-construct*, the *mask-expr* of the statement is evaluated, and the control mask is established to be the value of *mask-expr*. The pending control mask is established to have the value *.NOT. mask-expr* upon execution of a WHERE construct statement that does not appear as part of a *where-body-construct*. The *mask-expr* is evaluated only once.
- 2 Each statement in a WHERE construct is executed in sequence.
- 3 Upon execution of a *masked-elsewhere-stmt*, the following actions take place in sequence.
  - (1) The control mask  $m_c$  is established to have the value of the pending control mask.
  - (2) The pending control mask is established to have the value  $m_c$  *.AND. (.NOT. mask-expr)*.
  - (3) The control mask  $m_c$  is established to have the value  $m_c$  *.AND. mask-expr*.
- 4 The *mask-expr* is evaluated at most once.
- 5 Upon execution of an ELSEWHERE statement, the control mask is established to have the value of the pending control mask. No new pending control mask value is established.
- 6 Upon execution of an ENDWHERE statement, the control mask and pending control mask are established to have the values they had prior to the execution of the corresponding WHERE construct statement. Following the execution of a WHERE statement that appears as a *where-body-construct*, the control mask is established to have the value it had prior to the execution of the WHERE statement.

#### NOTE 7.49

The establishment of control masks and the pending control mask is illustrated with the following example:

```
WHERE(cond1)          ! Statement 1
. . .
ELSEWHERE(cond2)      ! Statement 2
. . .
ELSEWHERE             ! Statement 3
```

## NOTE 7.49 (cont.)

```

      . . .
      END WHERE

```

Following execution of statement 1, the control mask has the value `cond1` and the pending control mask has the value `.NOT. cond1`. Following execution of statement 2, the control mask has the value `(.NOT. cond1) .AND. cond2` and the pending control mask has the value `(.NOT. cond1) .AND. (.NOT. cond2)`. Following execution of statement 3, the control mask has the value `(.NOT. cond1) .AND. (.NOT. cond2)`. The false condition values are propagated through the execution of the masked ELSEWHERE statement.

- 1    7    Upon execution of a WHERE construct statement that is part of a *where-body-construct*, the pending control
- 2        mask is established to have the value  $m_c$  `.AND. (.NOT. mask-expr)`. The control mask is then established to
- 3        have the value  $m_c$  `.AND. mask-expr`. The *mask-expr* is evaluated at most once.
  
- 4    8    Upon execution of a WHERE statement that is part of a *where-body-construct*, the control mask is established
- 5        to have the value  $m_c$  `.AND. mask-expr`. The pending control mask is not altered.
  
- 6    9    If a nonelemental function reference occurs in the *expr* or *variable* of a *where-assignment-stmt* or in a *mask-expr*,
- 7        the function is evaluated without any masked control; that is, all of its argument expressions are fully evaluated
- 8        and the function is fully evaluated. If the result is an array and the reference is not within the argument list
- 9        of a nonelemental function, elements corresponding to true values in the control mask are selected for use in
- 10       evaluating the *expr*, variable or *mask-expr*.
  
- 11 10    If an elemental operation or function reference occurs in the *expr* or *variable* of a *where-assignment-stmt* or in a
- 12       *mask-expr*, and is not within the argument list of a nonelemental function reference, the operation is performed
- 13       or the function is evaluated only for the elements corresponding to true values of the control mask.
  
- 14 11    If an array constructor appears in a *where-assignment-stmt* or in a *mask-expr*, the array constructor is evaluated
- 15       without any masked control and then the *where-assignment-stmt* is executed or the *mask-expr* is evaluated.
  
- 16 12    When a *where-assignment-stmt* is executed, the values of *expr* that correspond to true values of the control mask
- 17       are assigned to the corresponding elements of the variable.
  
- 18 13    The value of the control mask is established by the execution of a WHERE statement, a WHERE construct
- 19       statement, an ELSEWHERE statement, a masked ELSEWHERE statement, or an ENDWHERE statement.
- 20       Subsequent changes to the value of entities in a *mask-expr* have no effect on the value of the control mask. The
- 21       execution of a function reference in the mask expression of a WHERE statement is permitted to affect entities in
- 22       the assignment statement.

## NOTE 7.50

Examples of function references in masked array assignments are:

```

WHERE (A > 0.0)
  A = LOG (A)           ! LOG is invoked only for positive elements.
  A = A / SUM (LOG (A)) ! LOG is invoked for all elements
                        ! because SUM is transformational
END WHERE

```

## 23    7.2.4    FORALL

## 24    7.2.4.1    Form of the FORALL Construct

- 25    1    The FORALL construct allows multiple assignments, masked array (WHERE) assignments, and nested FORALL constructs and
- 26        statements to be controlled by a single *concurrent-control-list* and *scalar-mask-expr*.

27    R750    *forall-construct*                                    is    *forall-construct-stmt*

1 *[forall-body-construct ] ...*  
 2 *end-forall-stmt*

3 R751 *forall-construct-stmt* is *[forall-construct-name :] FORALL concurrent-header*

4 R752 *forall-body-construct* is *forall-assignment-stmt*  
 5 or *where-stmt*  
 6 or *where-construct*  
 7 or *forall-construct*  
 8 or *forall-stmt*

9 R753 *forall-assignment-stmt* is *assignment-stmt*  
 10 or *pointer-assignment-stmt*

11 R754 *end-forall-stmt* is END FORALL *[forall-construct-name ]*

12 C735 (R754) If the *forall-construct-stmt* has a *forall-construct-name*, the *end-forall-stmt* shall have the same *forall-construct-*  
 13 *name*. If the *end-forall-stmt* has a *forall-construct-name*, the *forall-construct-stmt* shall have the same *forall-construct-*  
 14 *name*.

15 C736 (R752) A statement in a *forall-body-construct* shall not define an *index-name* of the *forall-construct*.

16 C737 (R752) Any procedure referenced in a *forall-body-construct*, including one referenced by a defined operation, assignment,  
 17 or *finalization*, shall be a *pure procedure*.

18 C738 (R752) A *forall-body-construct* shall not be a branch target.

19 2 The scope and attributes of an *index-name* in a *concurrent-header* in a FORALL construct or statement are described in 16.4.

## 20 7.2.4.2 Execution of the FORALL construct

### 21 7.2.4.2.1 Execution stages

- 22 1 There are three stages in the execution of a FORALL construct:
- 23 (1) determination of the values for *index-name* variables,
  - 24 (2) evaluation of the *scalar-mask-expr*, and
  - 25 (3) execution of the FORALL body constructs.

### 26 7.2.4.2.2 Determination of the values for index variables

- 27 1 The values of the index variables are determined as they are for the DO CONCURRENT statement (8.1.7.4.2).

### 28 7.2.4.2.3 Evaluation of the mask expression

- 29 1 The mask expression is evaluated as it is for the DO CONCURRENT statement (8.1.7.4.2).

### 30 7.2.4.2.4 Execution of the FORALL body constructs

- 31 1 The *forall-body-constructs* are executed in the order in which they appear. Each construct is executed for all active combinations of  
 32 the *index-name* values with the following interpretation:
- 33 2 Execution of a *forall-assignment-stmt* that is an *assignment-stmt* causes the evaluation of *expr* and all expressions within *variable*  
 34 for all active combinations of *index-name* values. These evaluations may be done in any order. After all these evaluations have been  
 35 performed, each *expr* value is assigned to the corresponding *variable*. The assignments may occur in any order.
  - 36 3 Execution of a *forall-assignment-stmt* that is a *pointer-assignment-stmt* causes the evaluation of all expressions within *data-target*  
 37 and *data-pointer-object* or *proc-target* and *proc-pointer-object*, the determination of any pointers within *data-pointer-object* or *proc-*  
 38 *pointer-object*, and the determination of the *target* for all active combinations of *index-name* values. These evaluations may be done  
 39 in any order. After all these evaluations have been performed, each *data-pointer-object* or *proc-pointer-object* is associated with the  
 40 corresponding *target*. These associations may occur in any order.

- 1 4 In a *forall-assignment-stmt*, a *defined assignment* subroutine shall not reference any *variable* that becomes defined by the statement.

#### NOTE 7.51

If a variable defined in an assignment statement within a FORALL construct is referenced in a later statement in that construct, the later statement uses the value(s) computed in the preceding assignment statement, not the value(s) the variable had prior to execution of the FORALL.

- 2 5 Each statement in a *where-construct* (7.2.3) within a *forall-construct* is executed in sequence. When a *where-stmt*, *where-construct-stmt* or *masked-elsewhere-stmt* is executed, the statement's *mask-expr* is evaluated for all active combinations of *index-name* values as determined by the outer *forall-constructs*, masked by any control mask corresponding to outer *where-constructs*. Any *where-assignment-stmt* is executed for all active combinations of *index-name* values, masked by the control mask in effect for the *where-assignment-stmt*.
- 3 6 Execution of a *forall-stmt* or *forall-construct* causes the evaluation of the *concurrent-limit* and *concurrent-step* expressions in the *concurrent-control-list* for all active combinations of the *index-name* values of the outer FORALL construct. The set of combinations of *index-name* values for the inner FORALL is the union of the sets defined by these limits and steps for each active combination of the outer *index-name* values; it also includes the outer *index-name* values. The *scalar-mask-expr* is then evaluated for all combinations of the *index-name* values of the inner construct to produce a set of active combinations for the inner construct. If there is no *scalar-mask-expr*, it is as if it appeared with the value true. Each statement in the inner FORALL is then executed for each active combination of the *index-name* values.

#### 7.2.4.3 The FORALL statement

- 15 1 The FORALL statement allows a single assignment statement or *pointer assignment statement* to be controlled by a set of index values and an optional mask expression.

17 R755 *forall-stmt* is FORALL *concurrent-header forall-assignment-stmt*

- 18 2 A FORALL statement is equivalent to a FORALL construct containing a single *forall-body-construct* that is a *forall-assignment-stmt*.

- 19 3 The scope of an *index-name* in a *forall-stmt* is the statement itself (16.4).

#### 7.2.4.4 Restrictions on FORALL constructs and statements

- 21 1 A many-to-one assignment is more than one assignment to the same object, or association of more than one *target* with the same pointer, whether the object is referenced directly or indirectly through a pointer. A many-to-one assignment shall not occur within a single statement in a FORALL construct or statement. It is possible to assign or pointer-assign to the same object in different *assignment* or *pointer assignment* statements in a FORALL construct.

#### NOTE 7.52

The appearance of each *index-name* in the identification of the left-hand side of an assignment statement is helpful in eliminating many-to-one assignments, but it is not sufficient to guarantee there will be none. For example, the following is allowed

```
FORALL (I = 1:10)
  A (INDEX (I)) = B(I)
END FORALL
```

if and only if INDEX(1:10) contains no repeated values.

- 25 2 Within the scope of a FORALL construct, a nested FORALL statement or FORALL construct shall not have the same *index-name*.  
26 The *concurrent-header* expressions within a nested FORALL may depend on the values of outer *index-name* variables.

1 (Blank page)

## 8 Execution control

### 8.1 Executable constructs containing blocks

#### 8.1.1 Blocks

1 The following are executable constructs that contain blocks:

- [ASSOCIATE construct](#);
- [BLOCK construct](#);
- [CHANGE TEAM construct](#);
- [CRITICAL construct](#);
- [DO construct](#);
- [IF construct](#);
- [SELECT CASE construct](#);
- [SELECT RANK construct](#);
- [SELECT TYPE construct](#).

R801 *block* is [ [execution-part-construct](#) ] ...

2 Executable constructs may be used to control which blocks of a program are executed or how many times a block is executed. Blocks are always bounded by statements that are particular to the construct in which they are embedded.

#### NOTE 8.1

An example of a construct containing a block is:

```
IF (A > 0.0) THEN
  B = SQRT (A)  ! These two statements
  C = LOG (A)   ! form a block.
END IF
```

#### 8.1.2 Rules governing blocks

##### 8.1.2.1 Control flow in blocks

- 1 Transfer of control to the interior of a block from outside the block is prohibited, except for the return from a procedure invoked within the block. Transfers within a block and transfers from the interior of a block to outside the block may occur.
- 2 Subroutine and function references ([12.5.3](#), [12.5.4](#)) may appear in a block.

##### 8.1.2.2 Execution of a block

- 1 Execution of a block begins with the execution of the first executable construct in the block.
- 2 Execution of the block is completed when
  - execution of the last executable construct in the block completes without branching to a statement within the block,
  - a branch ([8.2](#)) within the block that has a branch target outside the block occurs,

- a **RETURN statement** within the block is executed, or
- an **EXIT statement** or **CYCLE statement** statement that belongs to a construct that contains the block is executed.

**NOTE 8.2**

The action that takes place at the terminal boundary depends on the particular construct and on the block within that construct.

**8.1.3 ASSOCIATE construct****8.1.3.1 Purpose and form of the ASSOCIATE construct**

- 1 The ASSOCIATE construct associates named entities with expressions or variables during the execution of its block. These named **construct entities** (16.4) are **associating entities** (16.5.1.6). The names are **associate names**.

R802 *associate-construct*            **is**    *associate-stmt*  
    *block*  
    *end-associate-stmt*

R803 *associate-stmt*                            **is**    [ *associate-construct-name* : ] ASSOCIATE ■  
    ■ (*association-list* )

R804 *association*                                **is**    *associate-name* => *selector*

R805 *selector*                                    **is**    *expr*  
    **or**    *variable*

C801 (R804) If *selector* is not a *variable* or is a *variable* that has a **vector subscript**, *associate-name* shall not appear in a variable definition context (16.6.7).

C802 (R804) An *associate-name* shall not be the same as another *associate-name* in the same *associate-stmt*.

C803 (R805) *variable* shall not be a **coindexed object**.

C804 (R805) *expr* shall not be a variable.

C805 (R805) *expr* shall not be a **designator** of a procedure pointer or a function reference that returns a procedure pointer.

R806 *end-associate-stmt*                            **is**    END ASSOCIATE [ *associate-construct-name* ]

C806 (R806) If the *associate-stmt* of an *associate-construct* specifies an *associate-construct-name*, the corresponding *end-associate-stmt* shall specify the same *associate-construct-name*. If the *associate-stmt* of an *associate-construct* does not specify an *associate-construct-name*, the corresponding *end-associate-stmt* shall not specify an *associate-construct-name*.

**8.1.3.2 Execution of the ASSOCIATE construct**

- 1 Execution of an ASSOCIATE construct causes evaluation of every expression within every *selector* that is a variable designator and evaluation of every other *selector*, followed by execution of its block. During execution of that block each **associate name** identifies an entity which is associated (16.5.1.6) with the corresponding selector. The **associating entity** assumes the **declared type** and type parameters of the selector. If and only if the selector is **polymorphic**, the **associating entity** is **polymorphic**.
- 2 The other attributes of the **associating entity** are described in 8.1.3.3.
- 3 It is permissible to branch to an *end-associate-stmt* only from within its ASSOCIATE construct.

### 8.1.3.3 Other attributes of associate names

- 1 Within an ASSOCIATE or SELECT TYPE construct, each **associating entity** has the same **rank** as its associated selector. The lower bound of each dimension is the result of the intrinsic function **LBOUND** (13.9.109) applied to the corresponding dimension of **selector**. The upper bound of each dimension is one less than the sum of the lower bound and the extent. The **associating entity** has the **ASYNCHRONOUS** or **VOLATILE** attribute if and only if the selector is a variable and has the attribute. The **associating entity** does not have the **ALLOCATABLE** or **POINTER** attributes; it has the **TARGET** attribute if and only if the selector is a variable and has either the **TARGET** or **POINTER** attribute.
- 2 Within an ASSOCIATE, SELECT RANK, or SELECT TYPE construct, each **associating entity** has the same **corank** as its associated selector. If the selector is a **coarray**, the **cobounds** of each **codimension** of the **associating entity** are the same as those of the selector. If the **associating entity** is **polymorphic**, it assumes the **dynamic type** and **type parameter** values of the selector. If the selector has the **OPTIONAL** attribute, it shall be present. The **associating entity** is **contiguous** if and only if the selector is **contiguous**.
- 3 The **associating entity** itself is a variable, but if the selector is not a **definable** variable, the **associating entity** is not **definable** and shall not be **defined** or become **undefined**. If the selector is not permitted to appear in a variable definition context (16.6.7), the **associate name** shall not appear in a variable definition context.

### 8.1.3.4 Examples of the ASSOCIATE construct

#### NOTE 8.3

The following example illustrates an association with an expression.

```
ASSOCIATE ( Z => EXP(-(X**2+Y**2)) * COS(THETA) )
  PRINT *, A+Z, A-Z
END ASSOCIATE
```

The following example illustrates an association with a derived-type variable.

```
ASSOCIATE ( XC => AX%B(I,J)%C )
  XC%DV = XC%DV + PRODUCT(XC%EV(1:N))
END ASSOCIATE
```

The following example illustrates association with an **array section**.

```
ASSOCIATE ( ARRAY => AX%B(I,:)%C )
  ARRAY(N)%EV = ARRAY(N-1)%EV
END ASSOCIATE
```

The following example illustrates multiple associations.

```
ASSOCIATE ( W => RESULT(I,J)%W, ZX => AX%B(I,J)%D, ZY => AY%B(I,J)%D )
  W = ZX*X + ZY*Y
END ASSOCIATE
```

### 8.1.4 BLOCK construct

- 1 The BLOCK construct is an executable construct that may contain declarations.

R807	<i>block-construct</i>	is	<i>block-stmt</i> [ <i>block-specification-part</i> ] <i>block</i> <i>end-block-stmt</i>
R808	<i>block-stmt</i>	is	[ <i>block-construct-name</i> : ] BLOCK

- 1 R809 *block-specification-part* is [ *use-stmt* ]...  
 2 [ *import-stmt* ] ...  
 3 [ *implicit-part* ]  
 4 [ [ *declaration-construct* ] ...  
 5 *specification-construct* ]
- 6 R810 *end-block-stmt* is END BLOCK [ *block-construct-name* ]
- 7 C807 (R807) A *block-specification-part* shall not contain a COMMON, EQUIVALENCE, INTENT, NAMELIST,  
 8 OPTIONAL, statement function, or VALUE statement.
- 9 C808 (R807) A SAVE statement in a BLOCK construct shall contain a *saved-entity-list* that does not specify a  
 10 *common-block-name*.
- 11 C809 (R807) If the *block-stmt* of a *block-construct* specifies a *block-construct-name*, the corresponding *end-block-*  
 12 *stmt* shall specify the same *block-construct-name*. If the *block-stmt* does not specify a *block-construct-*  
 13 *name*, the corresponding *end-block-stmt* shall not specify a *block-construct-name*.
- 14 2 Except for the ASYNCHRONOUS and VOLATILE statements, specifications in a BLOCK construct declare  
 15 construct entities whose scope is that of the BLOCK construct (16.4). The appearance of the name of an object  
 16 that is not a construct entity in an ASYNCHRONOUS or VOLATILE statement in a BLOCK construct specifies  
 17 that the object has the attribute within the construct even if it does not have the attribute outside the construct.
- 18 3 Execution of a BLOCK construct causes evaluation of the specification expressions within its specification part  
 19 in a processor-dependent order, followed by execution of its block.

**NOTE 8.4**

The following is an example of a BLOCK construct.

```

IF (swapxy) THEN
  BLOCK
    REAL(KIND(x)) tmp
    tmp = x
    x = y
    y = tmp
  END BLOCK
END IF

```

Actions on a variable local to a BLOCK construct do not affect any variable of the same name outside the construct. For example,

```

F = 254E-2
BLOCK
  REAL F
  F = 39.37
END BLOCK
! F is still equal to 254E-2.

```

A SAVE statement outside a BLOCK construct does not affect variables local to the BLOCK construct, because a SAVE statement affects variables in its scoping unit rather than in its inclusive scope. For example,

```

SUBROUTINE S
  ...
  SAVE
  ...
  BLOCK

```

## NOTE 8.4 (cont.)

```

      REAL X                ! Not saved.
      REAL,SAVE :: Y(100)   ! SAVE attribute is allowed.
      Z = 3                  ! Implicitly declared in S, thus saved.
      ...
    END BLOCK
      ...
END SUBROUTINE

```

## 8.1.5 CHANGE TEAM construct

1 The CHANGE TEAM construct changes the [current team](#).

R811 *change-team-construct* is *change-team-stmt*  
*block*  
*end-change-team-stmt*

R812 *change-team-stmt* is [ *team-construct-name*: ] CHANGE TEAM ( *team-variable* ■  
 ■ [ *coarray-association-list*] [, *sync-stat-list*] )

R813 *coarray-association* is *codimension-decl* => *coselector-name*

R814 *end-change-team-stmt* is END TEAM [ ( [ *sync-stat-list*] ) ] [ *team-construct-name* ]

R815 *team-variable* is *scalar-variable*

C810 A branch (8.2) within a CHANGE TEAM construct shall not have a branch target that is outside the construct.

C811 A RETURN statement shall not appear within a CHANGE TEAM construct.

C812 An *exit-stmt* or *cycle-stmt* within a CHANGE TEAM construct shall not belong to an outer construct.

C813 If the *change-team-stmt* of a *change-team-construct* specifies a *team-construct-name*, the corresponding *end-change-team-stmt* shall specify the same *team-construct-name*. If the *change-team-stmt* of a *change-team-construct* does not specify a *team-construct-name*, the corresponding *end-change-team-stmt* shall not specify a *team-construct-name*.

C814 In a *change-team-stmt*, a *coarray-name* in a *codimension-decl* shall not be the same as a *coselector-name*, or another *coarray-name*, in that statement.

C815 A *team-variable* shall be of type `TEAM_TYPE` from the intrinsic module `ISO_FORTRAN_ENV`.

C816 No *coselector-name* shall appear more than once in a given *change-team-stmt*.

C817 (R813) A *coselector-name* shall be a coarray.

2 The values of the *team-variables* on the *active images* of the *team* shall be those of team variables defined by the same execution of the same FORM TEAM statement on all *active images* of the *current team*. The current team for the statements of the CHANGE TEAM *block* is the team specified by the value of the *team-variable*. The current team is not changed by a redefinition of the *team variable* during execution of the CHANGE TEAM construct. A CHANGE TEAM construct completes execution by executing its END TEAM statement.

## Unresolved Technical Issue 017

Same execution is not sufficiently well-defined

That is, image 3 executing a FORM TEAM statement is ipso facto **NOT** the same “execution” as image

**Unresolved Technical Issue 017 (cont.)**

177 executing a FORM TEAM statement. It needs to be something like “corresponding”, but then we need to say what makes them correspond, which would I guess be something like “synchronized executions”.

Or, say that the team variables defined by the “corresponding”/“synchronized”/whatever execution of FORM TEAM are “corresponding”/“compatible”/“synchronized”/weasel-like, and use that.

I would probably recommend saying something about synchronized FORM TEAM executions producing synchronized values and just talking about such values, as being on the whole simpler.

- 1 3 A coselector name identifies a coarray. The coarray shall be established when the CHANGE TEAM statement  
 2 begins execution. A *codimension-decl* in a *coarray-association* associates a coarray with an established coarray  
 3 during the execution of the block. This coarray is an associating entity (8.1.3.2, 8.1.3.3, 16.5.1.6). Its name is  
 4 an associate name that has the scope of the construct. It has the declared type, dynamic type, type parameters,  
 5 rank, and bounds of the established coarray. Its corank and cobounds are those specified in the *codimension-decl*.

**Unresolved Technical Issue 019****Coarray association incompletely specified.**

We need to spell coarray association out properly, not just put in a list of cross-references. See how we wrote the text in ASSOCIATE construct para 1... and the cross refs are wrong and do not say what is intended, at least not without modification, since a **CHANGE TEAM construct is not an ASSOCIATE, SELECT RANK or SELECT TYPE construct**. The subclause about “Execution of the ASSOCIATE construct” is also obviously irrelevant to this.

It should not need saying that 16.5.1.6 needs careful checking for applicability. Calling the selector a “coselector” automatically makes p4 wrong.

- 6 4 An allocatable coarray that was allocated immediately before executing a CHANGE TEAM statement shall not  
 7 be deallocated during execution of the construct. An allocatable coarray that was unallocated immediately before  
 8 executing a CHANGE TEAM statement, and which is allocated immediately before executing the corresponding  
 9 END TEAM statement, is deallocated by the execution of the END TEAM statement.

**Unresolved Technical Issue 020****Can expressions in CHANGE TEAM do image control?**

That would be a bad idea, but does not seem to be disallowed apart from the general “we did not say what it means” rule. We should probably require all function references in CHANGE TEAM to be pure.

Also if we don’t do that, there is no specification I could see of whether the functions are referenced before the synchronisation or afterwards.

- 10 5 All *active images* of the *current team* shall execute the same CHANGE TEAM statement. When a CHANGE  
 11 TEAM statement is executed, there is an implicit synchronization of all *active images* of the team containing  
 12 the executing image that is identified by *team-variable*. On each *active image* of the *team*, execution of the  
 13 segment following the statement is delayed until all other *active images* of the *team* have executed the same  
 14 statement the same number of times since execution last began in the *team that was current* before execution of  
 15 the CHANGE TEAM statement. When a CHANGE TEAM construct completes execution, there is an implicit  
 16 synchronization of all *active images* in the current team. On each *active image* of the team, execution of the  
 17 segment following the END TEAM statement is delayed until all other *active images* of the team have executed  
 18 the same construct the same number of times since execution last began in the team that was current before  
 19 execution of the corresponding CHANGE TEAM statement.

## Unresolved Technical Issue 021

## Implicit synchronisation badly specified here

1. What is “all active images of the team containing the executing image that is identified by *team-variable*” supposed to mean? A team variable does not identify an image.

Perhaps simply “all active images of the team specified by *team-variable*”?

2. “completes execution, there is an implicit synchronization of all active images in the current team”. And which “current team” is that? Hint: The “current team” on execution of the END TEAM statement is probably not the one you mean, but it is definitely the one this specifies.

Probably should be rewritten so that we call the “current team” immediately before executing CHANGE TEAM the “original team”, and then we can use that description whenever that is what we mean.

Might improve the horrible “last began in the team that was current before execution of the CHANGE TEAM statement” slightly too.

Also more wordsmithing required in the paragraph as a whole...

## NOTE 8.5

Deallocation of an allocatable coarray that was not allocated at the beginning of a CHANGE TEAM construct, but is allocated at the end of execution of the construct, occurs even for allocatable coarrays with the SAVE attribute.

## 8.1.6 CRITICAL construct

1 A CRITICAL construct limits execution of a block to one image at a time.

R816 *critical-construct* is *critical-stmt*  
*block*  
*end-critical-stmt*

R817 *critical-stmt* is [ *critical-construct-name* : ] CRITICAL [ ( [ *sync-stat-list* ] ) ]

R818 *end-critical-stmt* is END CRITICAL [ *critical-construct-name* ]

C818 (R816) If the *critical-stmt* of a *critical-construct* specifies a *critical-construct-name*, the corresponding *end-critical-stmt* shall specify the same *critical-construct-name*. If the *critical-stmt* of a *critical-construct* does not specify a *critical-construct-name*, the corresponding *end-critical-stmt* shall not specify a *critical-construct-name*.

C819 (R816) The *block* of a *critical-construct* shall not contain a RETURN statement or an image control statement.

C820 A branch (8.2) within a CRITICAL construct shall not have a branch target that is outside the construct.

2 Execution of the CRITICAL construct is completed when execution of its block is completed, or the executing image fails (2.3.6). A procedure invoked, directly or indirectly, from a CRITICAL construct shall not execute an image control statement.

3 The processor shall ensure that once an image has commenced executing *block*, no other image shall commence executing *block* until this image has completed executing *block*. The image shall not execute an image control statement during the execution of *block*. The sequence of executed statements is therefore a segment (8.6.2). If image T is the next to execute the construct after image M, the segment on image M precedes the segment on image T.

## Unresolved Technical Issue 018

**CRITICAL semantics do not permit image failure.**

Read the above paragraph.

- 1     4   The effect of a **STAT=** or **ERRMSG=** specifier in a **CRITICAL** statement is specified in subclause 8.6.11.

**NOTE 8.6**

If more than one **image** executes the block of a **CRITICAL** construct, its execution by one **image** always either precedes or succeeds its execution by another **image**. Typically no other statement ordering is needed. Consider the following example:

```
CRITICAL
  GLOBAL_COUNTER[1] = GLOBAL_COUNTER[1] + 1
END CRITICAL
```

The definition of `GLOBAL_COUNTER [1]` by a particular **image** will always precede the reference to the same variable by the next **image** to execute the block.

**NOTE 8.7**

The following example permits a large number of jobs to be shared among the **images**:

```
INTEGER :: NUM_JOBS[*], JOB

IF (THIS_IMAGE() == 1) READ(*,*) NUM_JOBS
SYNC ALL
DO
  CRITICAL
    JOB = NUM_JOBS[1]
    NUM_JOBS[1] = JOB - 1
  END CRITICAL
  IF (JOB > 0) THEN
    ! Work on JOB
  ELSE
    EXIT
  END IF
END DO
SYNC ALL
```

2     **8.1.7 DO construct**3     **8.1.7.1 Purpose and form of the DO construct**

- 4     1   The DO construct specifies the repeated execution of a sequence of executable constructs. Such a repeated  
5     sequence is called a loop.
- 6     2   The number of iterations of a loop can be determined at the beginning of execution of the DO construct, or can  
7     be left indefinite (“DO forever” or DO WHILE). The execution order of the iterations can be left indeterminate  
8     (DO CONCURRENT); except in this case, the loop can be terminated immediately (8.1.7.4.5). An iteration of  
9     the loop can be curtailed by executing a **CYCLE statement** (8.1.7.4.4).
- 10    3   There are three phases in the execution of a DO construct: initiation of the loop, execution of each iteration of  
11    the loop, and termination of the loop.

1 4 The scope and attributes of an *index-name* in a *concurrent-header* (DO CONCURRENT) are described in 16.4.

## 2 8.1.7.2 Form of the DO construct

3 R819 *do-construct* is *do-stmt*  
 4 *block*  
 5 *end-do*

6 R820 *do-stmt* is *nonlabel-do-stmt*  
 7 or *label-do-stmt*

8 R821 *label-do-stmt* is [ *do-construct-name* : ] DO *label* [ *loop-control* ]

9 R822 *nonlabel-do-stmt* is [ *do-construct-name* : ] DO [ *loop-control* ]

10 R823 *loop-control* is [ , ] *do-variable* = *scalar-int-expr*, *scalar-int-expr* ■  
 11 ■ [ , *scalar-int-expr* ]  
 12 or [ , ] WHILE ( *scalar-logical-expr* )  
 13 or [ , ] CONCURRENT *concurrent-header* *concurrent-locality*

14 R824 *do-variable* is *scalar-int-variable-name*

15 C821 (R824) The *do-variable* shall be a variable of type integer.

16 R825 *concurrent-header* is ( [ *integer-type-spec* :: ] *concurrent-control-list* [ , *scalar-mask-expr* ] )

17 R826 *concurrent-control* is *index-name* = *concurrent-limit* : *concurrent-limit* [ : *concurrent-step* ]

18 R827 *concurrent-limit* is *scalar-int-expr*

19 R828 *concurrent-step* is *scalar-int-expr*

20 R829 *concurrent-locality* is [ *locality-spec* ]...

21 R830 *locality-spec* is LOCAL ( *variable-name-list* )  
 22 or LOCAL\_INIT ( *variable-name-list* )  
 23 or SHARED ( *variable-name-list* )  
 24 or DEFAULT ( NONE )

25 C822 (R825) Any procedure referenced in the *scalar-mask-expr*, including one referenced by a *defined operation*,  
 26 shall be a *pure procedure* (12.7).

27 C823 (R826) The *index-name* shall be a named scalar variable of type integer.

28 C824 (R826) A *concurrent-limit* or *concurrent-step* in a *concurrent-control* shall not contain a reference to any  
 29 *index-name* in the *concurrent-control-list* in which it appears.

30 C825 A *variable-name* in a *locality-spec* shall be the name of a variable in the innermost executable construct  
 31 or *scoping unit* that includes the DO CONCURRENT statement.

32 C826 A *variable-name* in a *locality-spec* shall not be the same as an *index-name* in the *concurrent-header* of  
 33 the same DO CONCURRENT statement.

34 C827 The name of a variable shall not appear in more than one *variable-name-list*, or more than once in a  
 35 *variable-name-list*, in a *locality-spec*.

36 C828 A *variable-name* that appears in a LOCAL or LOCAL\_INIT *locality-spec* shall not have the *ALLOC-*  
 37 *ATABLE*, *INTENT (IN)*, or *OPTIONAL* attributes, shall not be of finalizable type, and shall not be a  
 38 *coarray* or an *assumed-size array*. A *variable-name* that is not permitted to appear in a variable definition  
 39 context shall not appear in a LOCAL or LOCAL\_INIT *locality-spec*.

1 C829 A variable that is referenced by the *scalar-mask-expr* of a *concurrent-header* or by any *concurrent-limit*  
 2 or *concurrent-step* in that *concurrent-header* shall not appear in a LOCAL *locality-spec* in the same DO  
 3 CONCURRENT statement.

4 C830 If the *locality-spec* DEFAULT ( NONE ) appears in a DO CONCURRENT statement, a variable that  
 5 appears in the *block* of the construct and is not an *index-name* of that construct shall have its locality  
 6 explicitly specified.

7 R831 *end-do* is *end-do-stmt*  
 8 or *continue-stmt*

9 R832 *end-do-stmt* is END DO [ *do-construct-name* ]

10 C831 (R819) If the *do-stmt* of a *do-construct* specifies a *do-construct-name*, the corresponding *end-do* shall be  
 11 an *end-do-stmt* specifying the same *do-construct-name*. If the *do-stmt* of a *do-construct* does not specify  
 12 a *do-construct-name*, the corresponding *end-do* shall not specify a *do-construct-name*.

13 C832 (R819) If the *do-stmt* is a *nonlabel-do-stmt*, the corresponding *end-do* shall be an *end-do-stmt*.

14 C833 (R819) If the *do-stmt* is a *label-do-stmt*, the corresponding *end-do* shall be identified with the same *label*.

15 1 It is permissible to branch to an *end-do* only from within its DO construct.

### 16 8.1.7.3 Active and inactive DO constructs

17 1 A DO construct is either active or inactive. Initially inactive, a DO construct becomes active only when its DO  
 18 *statement* is executed.

19 2 Once active, the DO construct becomes inactive only when it terminates (8.1.7.4.5).

### 20 8.1.7.4 Execution of a DO construct

#### 21 8.1.7.4.1 Loop initiation

22 1 When the DO *statement* is executed, the DO construct becomes active. If *loop-control* is

23 [ , ] *do-variable* = *scalar-int-expr*<sub>1</sub> , *scalar-int-expr*<sub>2</sub> [ , *scalar-int-expr*<sub>3</sub> ]

24 the following steps are performed in sequence.

25 (1) The initial parameter  $m_1$ , the terminal parameter  $m_2$ , and the incrementation parameter  $m_3$  are  
 26 of type integer with the same *kind type parameter* as the *do-variable*. Their values are established  
 27 by evaluating *scalar-int-expr*<sub>1</sub>, *scalar-int-expr*<sub>2</sub>, and *scalar-int-expr*<sub>3</sub>, respectively, including, if ne-  
 28 cessary, conversion to the *kind type parameter* of the *do-variable* according to the rules for numeric  
 29 conversion (Table 7.9). If *scalar-int-expr*<sub>3</sub> does not appear,  $m_3$  has the value 1. The value of  $m_3$   
 30 shall not be zero.

31 (2) The DO variable becomes defined with the value of the initial parameter  $m_1$ .

32 (3) The iteration count is established and is the value of the expression  $(m_2 - m_1 + m_3)/m_3$ , unless that  
 33 value is negative, in which case the iteration count is 0.

#### NOTE 8.8

The iteration count is zero whenever:

$m_1 > m_2$  and  $m_3 > 0$ , or  
 $m_1 < m_2$  and  $m_3 < 0$ .

34 2 If *loop-control* is omitted, no iteration count is calculated. The effect is as if a large positive iteration count,  
 35 impossible to decrement to zero, were established. If *loop-control* is [ , ] WHILE (*scalar-logical-expr*), the effect  
 36 is as if *loop-control* were omitted and the following statement inserted as the first statement of the *block*:

1 IF (.NOT. (*scalar-logical-expr*)) EXIT

2 3 For a DO CONCURRENT construct, the values of the index variables for the iterations of the construct are  
3 determined by the rules in 8.1.7.4.2.

4 4 At the completion of the execution of the DO statement, the execution cycle begins.

#### 5 8.1.7.4.2 DO CONCURRENT loop control

6 1 The *concurrent-limit* and *concurrent-step* expressions in the *concurrent-control-list* are evaluated. These ex-  
7 pressions may be evaluated in any order. The set of values that a particular *index-name* variable assumes is  
8 determined as follows.

- 9 (1) The lower bound  $m_1$ , the upper bound  $m_2$ , and the step  $m_3$  are of type integer with the same *kind type*  
10 *parameter* as the *index-name*. Their values are established by evaluating the first *concurrent-limit*,  
11 the second *concurrent-limit*, and the *concurrent-step* expressions, respectively, including, if necessary,  
12 conversion to the *kind type parameter* of the *index-name* according to the rules for numeric conversion  
13 (Table 7.9). If *concurrent-step* does not appear,  $m_3$  has the value 1. The value  $m_3$  shall not be zero.
- 14 (2) Let the value of  $max$  be  $(m_2 - m_1 + m_3)/m_3$ . If  $max \leq 0$  for some *index-name*, the execution of the  
15 construct is complete. Otherwise, the set of values for the *index-name* is

$$m_1 + (k - 1) \times m_3 \quad \text{where } k = 1, 2, \dots, max.$$

17 2 The set of combinations of *index-name* values is the Cartesian product of the sets defined by each triplet specific-  
18 ation. An *index-name* becomes defined when this set is evaluated.

19 3 The *scalar-mask-expr*, if any, is evaluated for each combination of *index-name* values. If there is no *scalar-*  
20 *mask-expr*, it is as if it appeared with the value true. The *index-name* variables may be primaries in the  
21 *scalar-mask-expr*.

22 4 The set of active combinations of *index-name* values is the subset of all possible combinations for which the  
23 *scalar-mask-expr* has the value true.

#### NOTE 8.9

The *index-name* variables can appear in the mask, for example

```
DO CONCURRENT (I=1:10, J=1:10, A(I) > 0.0 .AND. B(J) < 1.0)
. . .
```

#### 24 8.1.7.4.3 The execution cycle

25 1 The execution cycle of a DO construct that is not a DO CONCURRENT construct consists of the following steps  
26 performed in sequence repeatedly until termination.

- 27 (1) The iteration count, if any, is tested. If it is zero, the loop terminates and the DO construct becomes  
28 inactive. If *loop-control* is [ , ] WHILE (*scalar-logical-expr*), the *scalar-logical-expr* is evaluated; if  
29 the value of this expression is false, the loop terminates and the DO construct becomes inactive.
- 30 (2) The *block* of the loop is executed.
- 31 (3) The iteration count, if any, is decremented by one. The DO variable, if any, is incremented by the  
32 value of the incrementation parameter  $m_3$ .

33 2 Except for the incrementation of the DO variable that occurs in step (3), the DO variable shall neither be redefined  
34 nor become undefined while the DO construct is active.

35 3 The *block* of a DO CONCURRENT construct is executed for every active combination of the *index-name* values.  
36 Each execution of the *block* is an iteration. The executions may occur in any order.

#### 8.1.7.4.4 CYCLE statement

1 Execution of a loop iteration can be curtailed by executing a CYCLE statement that belongs to the construct.

R833 *cycle-stmt* is CYCLE [ *do-construct-name* ]

C834 If a *do-construct-name* appears on a CYCLE statement, the CYCLE statement shall be within that *do-construct*; otherwise, it shall be within at least one *do-construct*.

C835 A *cycle-stmt* shall not appear within a CHANGE TEAM, CRITICAL, or DO CONCURRENT construct if it belongs to an outer construct.

2 A CYCLE statement belongs to a particular DO construct. If the CYCLE statement contains a DO construct name, it belongs to that DO construct; otherwise, it belongs to the innermost DO construct in which it appears.

3 Execution of a CYCLE statement that belongs to a DO construct that is not a DO CONCURRENT construct causes immediate progression to step (3) of the execution cycle of the DO construct to which it belongs.

4 Execution of a CYCLE statement that belongs to a DO CONCURRENT construct completes execution of that iteration of the construct.

5 In a DO construct, a transfer of control to the *end-do* has the same effect as execution of a CYCLE statement belonging to that construct.

#### 8.1.7.4.5 Loop termination

1 For a DO construct that is not a DO CONCURRENT construct, the loop terminates, and the DO construct becomes inactive, when any of the following occurs.

- The iteration count is determined to be zero or the *scalar-logical-expr* is false, when tested during step (1) of the above execution cycle.
- An EXIT statement that belongs to the DO construct is executed.
- An EXIT or CYCLE statement that belongs to an outer construct and is within the DO construct is executed.
- A branch occurs within the DO construct and the *branch target statement* is outside the construct.
- A RETURN statement within the DO construct is executed.

2 For a DO CONCURRENT construct, the loop terminates, and the DO construct becomes inactive when all of the iterations have completed execution.

3 When a DO construct becomes inactive, the DO variable, if any, of the DO construct retains its last defined value.

#### 8.1.7.5 Additional semantics for DO CONCURRENT constructs

C836 A RETURN statement shall not appear within a DO CONCURRENT construct.

C837 An *image control statement* shall not appear within a DO CONCURRENT construct.

C838 A branch (8.2) within a DO CONCURRENT construct shall not have a branch target that is outside the construct.

C839 A reference to an impure procedure shall not appear within a DO CONCURRENT construct.

C840 A reference to the procedure IEEE\_GET\_FLAG, IEEE\_SET\_HALTING\_MODE, or IEEE\_GET\_-HALTING\_MODE from the intrinsic module IEEE\_EXCEPTIONS, shall not appear within a DO CONCURRENT construct.

- 1 1 The locality of a variable that appears in a DO CONCURRENT construct is **LOCAL**, **LOCAL\_INIT**, **SHARED**,  
2 or unspecified.
- 3 2 A variable that has **LOCAL** or **LOCAL\_INIT** locality is a construct entity with the same type, type parameters,  
4 rank as the variable with the same name in the innermost executable construct or **scoping unit** that includes the  
5 DO CONCURRENT construct, and the outside variable is inaccessible by that name within the construct. The  
6 construct entity has the **ASYNCHRONOUS**, **CONTIGUOUS**, **POINTER**, **TARGET**, or **VOLATILE** attribute if  
7 and only if the outside variable has that attribute; it does not have the **BIND**, **INTENT**, **PROTECTED**, **SAVE**  
8 or **VALUE** attribute, even if the outside variable has that attribute. If it is not a pointer, it has the same bounds  
9 as the outside variable. At the beginning of execution of each iteration,  
10 • if a variable with **LOCAL** locality is a pointer it has undefined pointer association status, and otherwise it  
11 is undefined except for any subobjects that are default-initialized;  
12 • a variable with **LOCAL\_INIT** locality has the pointer association status and definition status of the out-  
13 side variable with that name; the outside variable shall not be an undefined pointer or a nonallocatable  
14 nonpointer variable that is undefined.
- 15 If a variable with **LOCAL** or **LOCAL\_INIT** locality becomes an affector of a pending input/output operation,  
16 the operation shall have completed before the end of the iteration. If a variable with **LOCAL** or **LOCAL\_INIT**  
17 locality has the **TARGET attribute**, a pointer associated with it during an iteration becomes undefined when  
18 execution of that iteration completes.
- 19 3 If a variable has **SHARED** locality, appearances of the variable within the DO CONCURRENT construct refer  
20 to the variable in the innermost executable construct or **scoping unit** that includes the DO CONCURRENT  
21 construct. If it is defined or becomes undefined during any iteration, it shall not be referenced, defined, or  
22 become undefined during any other iteration. If it is allocated or deallocated during an iteration it shall not  
23 have its allocation or association status, dynamic type, array bounds, shape, or a deferred type parameter value  
24 inquired about in any other iteration.
- 25 4 If a variable has unspecified locality,  
26 • if it is referenced in an iteration it shall either be previously defined during that iteration, or shall not be  
27 defined or become undefined during any other iteration; if it is defined or becomes undefined by more than  
28 one iteration it becomes undefined when the loop terminates;  
29 • if it is a pointer and is used in an iteration other than as the pointer in pointer assignment, allocation,  
30 or nullification, it shall either be previously **pointer associated** during that iteration or shall not have its  
31 **pointer association** changed during any iteration;  
32 • if it is a pointer whose **pointer association** is changed in more than one iteration, it has an association status  
33 of undefined when the construct terminates;  
34 • if it is **allocatable** and is allocated in more than one iteration, it shall have an allocation status of unallocated  
35 at the end of every iteration;  
36 • if it is **allocatable** and is referenced, defined, deallocated, or has its allocation status, dynamic type, or a  
37 deferred type parameter value inquired about, in any iteration, it shall either be previously allocated in  
38 that iteration or shall not be allocated or deallocated in any other iteration.
- 39 5 If data are written to a file record or position in one iteration, that record or position in that file shall not be  
40 read from or written to in a different iteration. If records are written to a file connected for sequential access by  
41 more than one iteration, the ordering of records written by different iterations is processor dependent.

**NOTE 8.10**

The restrictions on referencing variables defined in an iteration of a DO CONCURRENT construct apply to any procedure invoked within the loop.

**NOTE 8.11**

The restrictions on the statements in a DO CONCURRENT construct are designed to ensure there are no data dependencies between iterations of the loop. This permits code optimizations that might otherwise be difficult or impossible because they would depend on properties of the program not visible to the compiler.

## 1 8.1.7.6 Examples of DO constructs

**NOTE 8.12**

The following program fragment computes a tensor product of two arrays:

```
DO I = 1, M
  DO J = 1, N
    C (I, J) = DOT_PRODUCT (A (I, J, :), B(:, I, J))
  END DO
END DO
```

**NOTE 8.13**

The following program fragment contains a DO construct that uses the **WHILE** form of *loop-control*. The loop will continue to execute until an end-of-file or input/output error is encountered, at which point the **DO statement** terminates the loop. When a negative value of X is read, the program skips immediately to the next **READ statement**, bypassing most of the *block* of the loop.

```
READ (IUN, '(1X, G14.7)', IOSTAT = IOS) X
DO WHILE (IOS == 0)
  IF (X >= 0.) THEN
    CALL SUBA (X)
    CALL SUBB (X)
    ...
    CALL SUBZ (X)
  ENDIF
  READ (IUN, '(1X, G14.7)', IOSTAT = IOS) X
END DO
```

**NOTE 8.14**

The following example behaves exactly the same as the one in Note 8.13. However, the **READ statement** has been moved to the interior of the loop, so that only one **READ statement** is needed. Also, a **CYCLE statement** has been used to avoid an extra level of **IF** nesting.

```
DO      ! A "DO WHILE + 1/2" loop
  READ (IUN, '(1X, G14.7)', IOSTAT = IOS) X
  IF (IOS /= 0) EXIT
  IF (X < 0.) CYCLE
  CALL SUBA (X)
  CALL SUBB (X)
  . . .
  CALL SUBZ (X)
END DO
```

**NOTE 8.15**

The following example illustrates a case in which the user knows that there are no repeated values in the index array IND. The DO CONCURRENT construct makes it easier for the processor to generate vector gather/scatter code, unroll the loop, or parallelize the code for this loop, potentially improving performance.

```
INTEGER :: A(N), IND(N)

DO CONCURRENT (I=1:M)
  A(IND(I)) = I
END DO
```

**NOTE 8.16**

The following code demonstrates the use of the **LOCAL** clause so that the X inside the **DO CONCURRENT construct** is a temporary variable, and will not affect the X outside the construct.

```

X = 1.0
DO CONCURRENT (I=1:10) LOCAL (X)
  IF (A (I) > 0) THEN
    X = SQRT (A (I))
    A (I) = A (I) - X**2
  END IF
  B (I) = B (I) - A (I)
END DO
PRINT *, X
! Always prints 1.0.

```

**NOTE 8.17**

Additional examples of DO constructs are in [C.5.3](#).

**8.1.8 IF construct and statement****8.1.8.1 Purpose and form of the IF construct**

- 1 The IF construct selects for execution at most one of its constituent blocks. The selection is based on a sequence of logical expressions.

```

R834  if-construct           is  if-then-stmt
                                block
                                [ else-if-stmt
                                block ] ...
                                [ else-stmt
                                block ]
                                end-if-stmt

R835  if-then-stmt           is  [ if-construct-name : ] IF ( scalar-logical-expr ) THEN

R836  else-if-stmt           is  ELSE IF ( scalar-logical-expr ) THEN [ if-construct-name ]

R837  else-stmt              is  ELSE [ if-construct-name ]

R838  end-if-stmt            is  END IF [ if-construct-name ]

```

C841 (R834) If the *if-then-stmt* of an *if-construct* specifies an *if-construct-name*, the corresponding *end-if-stmt* shall specify the same *if-construct-name*. If the *if-then-stmt* of an *if-construct* does not specify an *if-construct-name*, the corresponding *end-if-stmt* shall not specify an *if-construct-name*. If an *else-if-stmt* or *else-stmt* specifies an *if-construct-name*, the corresponding *if-then-stmt* shall specify the same *if-construct-name*.

**8.1.8.2 Execution of an IF construct**

- 1 At most one of the blocks in the IF construct is executed. If there is an ELSE statement in the construct, exactly one of the blocks in the construct is executed. The scalar logical expressions are evaluated in the order of their appearance in the construct until a true value is found or an ELSE statement or END IF statement is encountered. If a true value or an ELSE statement is found, the block immediately following is executed and this completes the execution of the construct. The scalar logical expressions in any remaining ELSE IF statements of the IF construct are not evaluated. If none of the evaluated expressions is true and there is no ELSE statement, the execution of the construct is completed without the execution of any block within the construct.

- 1 2 It is permissible to branch to an END IF statement only from within its IF construct. Execution of an END IF  
2 statement has no effect.

### 3 8.1.8.3 Examples of IF constructs

#### NOTE 8.18

```

IF (CVAR == 'RESET') THEN
  I = 0; J = 0; K = 0
END IF
PROOF_DONE: IF (PROP) THEN
  WRITE (3, '( 'QED' )')
  STOP
ELSE
  PROP = NEXTPROP
END IF PROOF_DONE
IF (A > 0) THEN
  B = C/A
  IF (B > 0) THEN
    D = 1.0
  END IF
ELSE IF (C > 0) THEN
  B = A/C
  D = -1.0
ELSE
  B = ABS (MAX (A, C))
  D = 0
END IF

```

### 4 8.1.8.4 IF statement

- 5 1 The IF statement controls the execution of a single action statement based on a single logical expression.

6 R839 *if-stmt* is IF ( *scalar-logical-expr* ) *action-stmt*

7 C842 (R839) The *action-stmt* in the *if-stmt* shall not be an *end-function-stmt*, *end-mp-subprogram-stmt*, *end-*  
8 *program-stmt*, *end-subroutine-stmt*, or *if-stmt*.

- 9 2 Execution of an IF statement causes evaluation of the scalar logical expression. If the value of the expression is  
10 true, the action statement is executed. If the value is false, the action statement is not executed.

- 11 3 The execution of a function reference in the scalar logical expression may affect entities in the action statement.

#### NOTE 8.19

An example of an IF statement is:

```

IF (A > 0.0) A = LOG (A)

```

## 12 8.1.9 SELECT CASE construct

### 13 8.1.9.1 Purpose and form of the SELECT CASE construct

- 14 1 The SELECT CASE construct selects for execution at most one of its constituent blocks. The selection is based  
15 on the value of an expression.

16 R840 *case-construct* is *select-case-stmt*  
17 [ *case-stmt*

1 *block* ] ...  
 2 *end-select-stmt*

3 R841 *select-case-stmt* is [ *case-construct-name* : ] SELECT CASE ( *case-expr* )

4 R842 *case-stmt* is CASE *case-selector* [*case-construct-name*]

5 R843 *end-select-stmt* is END SELECT [ *case-construct-name* ]

6 C843 (R840) If the *select-case-stmt* of a *case-construct* specifies a *case-construct-name*, the corresponding *end-*  
 7 *select-stmt* shall specify the same *case-construct-name*. If the *select-case-stmt* of a *case-construct* does  
 8 not specify a *case-construct-name*, the corresponding *end-select-stmt* shall not specify a *case-construct-*  
 9 *name*. If a *case-stmt* specifies a *case-construct-name*, the corresponding *select-case-stmt* shall specify the  
 10 same *case-construct-name*.

11 R844 *case-expr* is *scalar-expr*

12 C844 *case-expr* shall be of type character, integer, or logical.

13 R845 *case-selector* is ( *case-value-range-list* )  
 14 or DEFAULT

15 C845 (R840) No more than one of the selectors of one of the CASE statements shall be DEFAULT.

16 R846 *case-value-range* is *case-value*  
 17 or *case-value* :  
 18 or : *case-value*  
 19 or *case-value* : *case-value*

20 R847 *case-value* is *scalar-constant-expr*

21 C846 (R840) For a given *case-construct*, each *case-value* shall be of the same type as *case-expr*. For character  
 22 type, the *kind type parameters* shall be the same; character length differences are allowed.

23 C847 (R840) A *case-value-range* using a colon shall not be used if *case-expr* is of type logical.

24 C848 (R840) For a given *case-construct*, there shall be no possible value of the *case-expr* that matches more  
 25 than one *case-value-range*.

## 26 8.1.9.2 Execution of a SELECT CASE construct

27 1 The execution of the SELECT CASE statement causes the case expression to be evaluated. For a case value  
 28 range list, a match occurs if the case expression value matches any of the case value ranges in the list. For a case  
 29 expression with a value of *c*, a match is determined as follows.

- 30 (1) If the case value range contains a single value *v* without a colon, a match occurs for type logical if  
 31 the expression *c* .EQV. *v* is true, and a match occurs for type integer or character if the expression  
 32 *c* == *v* is true.
- 33 (2) If the case value range is of the form *low* : *high*, a match occurs if the expression *low* <= *c* .AND.  
 34 *c* <= *high* is true.
- 35 (3) If the case value range is of the form *low* :, a match occurs if the expression *low* <= *c* is true.
- 36 (4) If the case value range is of the form : *high*, a match occurs if the expression *c* <= *high* is true.
- 37 (5) If no other selector matches and a DEFAULT selector appears, it matches the case index.
- 38 (6) If no other selector matches and the DEFAULT selector does not appear, there is no match.

39 2 The block following the CASE statement containing the matching selector, if any, is executed. This completes  
 40 execution of the construct.

41 3 It is permissible to branch to an *end-select-stmt* only from within its SELECT CASE construct.

## 1 8.1.9.3 Examples of SELECT CASE constructs

**NOTE 8.20**

An integer signum function:

```

INTEGER FUNCTION SIGNUM (N)
SELECT CASE (N)
CASE (:-1)
    SIGNUM = -1
CASE (0)
    SIGNUM = 0
CASE (1:)
    SIGNUM = 1
END SELECT
END

```

**NOTE 8.21**

A code fragment to check for balanced parentheses:

```

CHARACTER (80) :: LINE
...
LEVEL = 0
SCAN_LINE: DO I = 1, 80
    CHECK_PARENS: SELECT CASE (LINE (I:I))
    CASE ('(')
        LEVEL = LEVEL + 1
    CASE (')')
        LEVEL = LEVEL - 1
        IF (LEVEL < 0) THEN
            PRINT *, 'UNEXPECTED RIGHT PARENTHESIS'
            EXIT SCAN_LINE
        END IF
    CASE DEFAULT
        ! Ignore all other characters
    END SELECT CHECK_PARENS
END DO SCAN_LINE
IF (LEVEL > 0) THEN
    PRINT *, 'MISSING RIGHT PARENTHESIS'
END IF

```

**NOTE 8.22**

The following three fragments are equivalent:

```

IF (SILLY == 1) THEN
    CALL THIS
ELSE
    CALL THAT
END IF
SELECT CASE (SILLY == 1)
CASE (.TRUE.)
    CALL THIS
CASE (.FALSE.)
    CALL THAT
END SELECT

```

**NOTE 8.22 (cont.)**

```

SELECT CASE (SILLY)
CASE DEFAULT
    CALL THAT
CASE (1)
    CALL THIS
END SELECT

```

**NOTE 8.23**

A code fragment showing several selections of one block:

```

SELECT CASE (N)
CASE (1, 3:5, 8) ! Selects 1, 3, 4, 5, 8
    CALL SUB
CASE DEFAULT
    CALL OTHER
END SELECT

```

**8.1.10 SELECT RANK construct****8.1.10.1 Purpose and form of the SELECT RANK construct**

- 1 The SELECT RANK construct selects for execution at most one of its constituent blocks. The selection is based on the rank of an [assumed-rank variable](#). A name is associated with the variable (16.4, 16.5.1.6), in the same way as for the [ASSOCIATE construct](#).

R848 *select-rank-construct*      **is** *select-rank-stmt*  
    [ *select-rank-case-stmt*  
    *block* ]...  
    *end-select-rank-stmt*

R849 *select-rank-stmt*      **is** [ *select-construct-name* : ] SELECT RANK ■  
    ■ ( [ *associate-name* => ] *selector* )

C849 The *selector* in a *select-rank-stmt* shall be the name of an [assumed-rank](#) array.

R850 *select-rank-case-stmt*      **is** RANK ( *scalar-int-constant-expr* ) [ *select-construct-name* ]  
    **or** RANK ( \* ) [ *select-construct-name* ]  
    **or** RANK DEFAULT [ *select-construct-name* ]

C850 A *scalar-int-constant-expr* in a *select-rank-case-stmt* shall be nonnegative and less than or equal to the maximum rank supported by the processor.

C851 For a given *select-rank-construct*, the same rank value shall not be specified in more than one *select-rank-case-stmt*.

C852 For a given *select-rank-construct*, there shall be at most one RANK ( \* ) *select-rank-case-stmt* and at most one RANK DEFAULT *select-rank-case-stmt*.

C853 If *select-construct-name* appears on a *select-rank-case-stmt* the corresponding *select-rank-stmt* shall specify the same *select-construct-name*.

C854 A SELECT RANK construct shall not have a *select-rank-case-stmt* that is RANK ( \* ) if the selector has the [ALLOCATABLE](#) or [POINTER](#) attribute.

R851 *end-select-rank-stmt*      **is** END SELECT [ *select-construct-name* ]

C855 If the *select-rank-stmt* of a *select-rank-construct* specifies a *select-construct-name*, the corresponding *end-select-rank-stmt* shall specify the same *select-construct-name*. If the *select-rank-stmt* of a *select-rank-construct* does not specify a *select-construct-name*, the corresponding *end-select-rank-stmt* shall not specify a *select-construct-name*.

2 The associate name of a SELECT RANK construct is the *associate-name* if specified; otherwise it is the name that constitutes the selector.

#### 8.1.10.2 Execution of the SELECT RANK construct

1 A SELECT RANK construct selects at most one block to be executed. During execution of that block, the associate name identifies an entity which is associated (16.5.1.6) with the selector. A RANK ( \* ) statement matches the selector if the selector is argument associated with an *assumed-size array*. A RANK ( *scalar-int-constant-expr* ) statement matches the selector if the selector has that rank and is not argument associated with an *assumed-size array*. A RANK DEFAULT statement matches the selector if no other *select-rank-case-stmt* of the construct matches the selector. If a *select-rank-case-stmt* matches the selector, the block following that statement is executed; otherwise, control is transferred to the *end-select-rank-stmt*.

2 It is permissible to branch to an *end-select-rank-stmt* only from within its SELECT RANK construct.

#### 8.1.10.3 Attributes of a SELECT RANK associate name

1 Within the block following a RANK DEFAULT statement, the associating entity (16.5.5) is *assumed-rank* and has exactly the same attributes as the selector. Within the block following a RANK ( \* ) statement, the associating entity has rank 1 and is *assumed-size*, as if it were declared with DIMENSION(1:\*). Within the block following a RANK ( *scalar-int-constant-expr* ) statement, the associating entity has the specified rank; the lower bound of each dimension is the result of the intrinsic function LBOUND (13.9.109) applied to the corresponding dimension of the selector, and the upper bound of each dimension is the result of the intrinsic function UBOUND (13.9.196) applied to the corresponding dimension of the selector.

2 If the selector has the *ALLOCATABLE attribute*, the associating entity has that attribute. If the selector has the *POINTER attribute*, the associating entity has that attribute (and not the *TARGET attribute*). The other attributes of the associating entity are described in 8.1.3.3.

#### 8.1.10.4 Examples of the SELECT RANK construct

1 This example shows how to use a SELECT RANK construct to process scalars and rank-2 arrays; anything else will be rejected as an error.

```

SUBROUTINE process(x)
  REAL x(..)
  !
  SELECT RANK(x)
  RANK (0)
    x = 0
  RANK (2)
    IF (SIZE(x,2)>=2) x(:,2) = 2
  RANK DEFAULT
    Print *, 'I did not expect rank', RANK(y), 'shape', SHAPE(y)
    ERROR STOP 'process bad arg'
  END SELECT

```

2 The following example shows how to process *assumed-size arrays*, including how to use *sequence association* if you want to do multi-dimensional processing of an *assumed-size array*.

```

SELECT RANK (y => x)
RANK (*)

```

```

1      IF (RANK(x)==2) THEN
2          ! Special code for the rank two case.
3          CALL sequence_assoc_2(y, LBOUND(x,1), UBOUND(x,1), LBOUND(x,2))
4      ELSE
5          ! We just do all the other ranks in array element order.
6          i = 1
7          DO
8              IF (y(i)==0) Exit
9              y(i) = -y(i)
10             i = i + 1
11         END DO
12     END IF
13 END SELECT
14 ...
15 CONTAINS
16 ...
17 SUBROUTINE sequence_assoc_2(a, lb1, ub1, lb2)
18     REAL a(lb1:ub1,lb2:*)
19     j = lb2
20 outer: DO
21     DO i=lb1,ub1
22         IF (a(i,j)==0) EXIT outer
23         a(i,j) = a(i,j)**2
24     END DO
25     j = j + 1
26     IF (ANY(a(:,j))==0) EXIT
27     j = j + 1
28 END DO outer
29 END SUBROUTINE

```

## 30 8.1.11 SELECT TYPE construct

### 31 8.1.11.1 Purpose and form of the SELECT TYPE construct

32 1 The SELECT TYPE construct selects for execution at most one of its constituent blocks. The selection is based  
33 on the [dynamic type](#) of an expression. A [name](#) is associated with the expression or variable ([16.4](#), [16.5.1.6](#)), in  
34 the same way as for the ASSOCIATE construct.

35 R852 *select-type-construct* is *select-type-stmt*  
36 [ *type-guard-stmt*  
37 *block* ] ...  
38 *end-select-type-stmt*

39 R853 *select-type-stmt* is [ *select-construct-name* : ] SELECT TYPE ■  
40 ■ ( [ *associate-name* => ] *selector* )

41 C856 (R853) If *selector* is not a named *variable*, *associate-name* => shall appear.

42 C857 (R853) If *selector* is not a *variable* or is a *variable* that has a *vector subscript*, *associate-name* shall not  
43 appear in a variable definition context ([16.6.7](#)).

44 C858 (R853) The *selector* in a *select-type-stmt* shall be *polymorphic*.

45 R854 *type-guard-stmt* is TYPE IS ( *type-spec* ) [ *select-construct-name* ]  
46 or CLASS IS ( *derived-type-spec* ) [ *select-construct-name* ]  
47 or CLASS DEFAULT [ *select-construct-name* ]

- 1 C859 (R854) The *type-spec* or *derived-type-spec* shall specify that each length type parameter is assumed.
- 2 C860 (R854) The *type-spec* or *derived-type-spec* shall not specify a type with the *BIND attribute* or the *SEQUENCE attribute*.
- 3
- 4 C861 (R852) If *selector* is not *unlimited polymorphic*, each TYPE IS or CLASS IS *type-guard-stmt* shall specify
- 5 an *extension* of the *declared type* of *selector*.
- 6 C862 (R852) For a given *select-type-construct*, the same type and *kind type parameter* values shall not be
- 7 specified in more than one TYPE IS *type-guard-stmt* and shall not be specified in more than one CLASS
- 8 IS *type-guard-stmt*.
- 9 C863 (R852) For a given *select-type-construct*, there shall be at most one CLASS DEFAULT *type-guard-stmt*.
- 10 R855 *end-select-type-stmt* is END SELECT [ *select-construct-name* ]
- 11 C864 (R852) If the *select-type-stmt* of a *select-type-construct* specifies a *select-construct-name*, the correspond-
- 12 ing *end-select-type-stmt* shall specify the same *select-construct-name*. If the *select-type-stmt* of a *select-*
- 13 *type-construct* does not specify a *select-construct-name*, the corresponding *end-select-type-stmt* shall not
- 14 specify a *select-construct-name*. If a *type-guard-stmt* specifies a *select-construct-name*, the corresponding
- 15 *select-type-stmt* shall specify the same *select-construct-name*.
- 16 2 The *associate name* of a SELECT TYPE construct is the *associate-name* if specified; otherwise it is the *name*
- 17 that constitutes the *selector*.

#### 8.1.11.2 Execution of the SELECT TYPE construct

- 19 1 Execution of a SELECT TYPE construct causes evaluation of every expression within a selector that is a variable
- 20 designator, or evaluation of a selector that is not a variable designator.
- 21 2 A SELECT TYPE construct selects at most one block to be executed. During execution of that block, the
- 22 *associate name* identifies an entity which is associated (16.5.1.6) with the selector.
- 23 3 A TYPE IS type guard statement matches the selector if the *dynamic type* and *kind type parameter* values of
- 24 the selector are the same as those specified by the statement. A CLASS IS type guard statement matches the
- 25 selector if the *dynamic type* of the selector is an *extension* of the type specified by the statement and the *kind*
- 26 *type parameter* values specified by the statement are the same as the corresponding type parameter values of the
- 27 *dynamic type* of the selector.
- 28 4 The block to be executed is selected as follows.
- 29 (1) If a TYPE IS type guard statement matches the selector, the block following that statement is
- 30 executed.
- 31 (2) Otherwise, if exactly one CLASS IS type guard statement matches the selector, the block following
- 32 that statement is executed.
- 33 (3) Otherwise, if several CLASS IS type guard statements match the selector, one of these statements
- 34 must specify a type that is an *extension* of all the types specified in the others; the block following
- 35 that statement is executed.
- 36 (4) Otherwise, if there is a CLASS DEFAULT type guard statement, the block following that statement
- 37 is executed.
- 38 (5) Otherwise, no block is executed.

#### NOTE 8.24

This algorithm does not examine the type guard statements in source text order when it looks for a match; it selects the most particular type guard when there are several potential matches.

- 39 5 Within the block following a TYPE IS type guard statement, the *associating entity* (16.5.5) is not polymorphic
- 40 (4.3.2.3), has the type named in the type guard statement, and has the type parameter values of the selector.

1 6 Within the block following a CLASS IS type guard statement, the *associating entity* is *polymorphic* and has the  
 2 *declared type* named in the type guard statement. The type parameter values of the *associating entity* are the  
 3 corresponding type parameter values of the selector.

4 7 Within the block following a CLASS DEFAULT type guard statement, the *associating entity* is *polymorphic* and  
 5 has the same *declared type* as the selector. The type parameter values of the *associating entity* are those of the  
 6 *declared type* of the selector.

#### NOTE 8.25

If the *declared type* of the *selector* is T, specifying CLASS DEFAULT has the same effect as specifying CLASS IS (T).

7 8 The other attributes of the *associating entity* are described in 8.1.3.3.

8 9 It is permissible to branch to an *end-select-type-stmt* only from within its SELECT TYPE construct.

### 8.1.11.3 Examples of the SELECT TYPE construct

#### NOTE 8.26

```
TYPE POINT
  REAL :: X, Y
END TYPE POINT
TYPE, EXTENDS(POINT) :: POINT_3D
  REAL :: Z
END TYPE POINT_3D
TYPE, EXTENDS(POINT) :: COLOR_POINT
  INTEGER :: COLOR
END TYPE COLOR_POINT

TYPE(POINT), TARGET :: P
TYPE(POINT_3D), TARGET :: P3
TYPE(COLOR_POINT), TARGET :: C
CLASS(POINT), POINTER :: P_OR_C
P_OR_C => C
SELECT TYPE ( A => P_OR_C )
CLASS IS ( POINT )
  ! "CLASS ( POINT ) :: A" implied here
  PRINT *, A%X, A%Y ! This block gets executed
TYPE IS ( POINT_3D )
  ! "TYPE ( POINT_3D ) :: A" implied here
  PRINT *, A%X, A%Y, A%Z
END SELECT
```

#### NOTE 8.27

The following example illustrates the omission of *associate-name*. It uses the declarations from Note 8.26.

```
P_OR_C => P3
SELECT TYPE ( P_OR_C )
CLASS IS ( POINT )
  ! "CLASS ( POINT ) :: P_OR_C" implied here
  PRINT *, P_OR_C%X, P_OR_C%Y
TYPE IS ( POINT_3D )
  ! "TYPE ( POINT_3D ) :: P_OR_C" implied here
  PRINT *, P_OR_C%X, P_OR_C%Y, P_OR_C%Z ! This block gets executed
```

## NOTE 8.27 (cont.)

END SELECT

## 8.1.12 EXIT statement

1 The EXIT statement provides one way of terminating a loop, or completing execution of another construct.

R856 *exit-stmt* is EXIT [ *construct-name* ]

C865 If a *construct-name* appears on an EXIT statement, the EXIT statement shall be within that construct; otherwise, it shall be within at least one *do-construct*.

2 An EXIT statement belongs to a particular construct. If a construct name appears, the EXIT statement belongs to that construct; otherwise, it belongs to the innermost *DO construct* in which it appears.

C866 An *exit-stmt* shall not appear within a **CHANGE TEAM**, **CRITICAL**, or **DO CONCURRENT** construct if it belongs to that construct or an outer construct.

3 When an EXIT statement that belongs to a *DO construct* is executed, it terminates the loop (8.1.7.4.5) and any active loops contained within the terminated loop. When an EXIT statement that belongs to a non-DO construct is executed, it terminates any active loops contained within that construct, and completes execution of that construct.

## 8.2 Branching

## 8.2.1 Branch concepts

1 Branching is used to alter the normal execution sequence. A branch causes a transfer of control from one statement to a labeled *branch target statement* in the same *inclusive scope*. Branching may be caused by a **GO TO statement**, a *computed GO TO statement*, a **CALL statement** that has an *alt-return-spec*, or an input/output statement that has an **END=**, **EOR=**, or **ERR=** specifier. Although procedure references and control constructs can cause transfer of control, they are not branches. A *branch target statement* is an *action-stmt*, *associate-stmt*, *end-associate-stmt*, *if-then-stmt*, *end-if-stmt*, *select-case-stmt*, *end-select-stmt*, *select-rank-stmt*, *end-select-rank-stmt*, *select-type-stmt*, *end-select-type-stmt*, *do-stmt*, *end-do-stmt*, *block-stmt*, *end-block-stmt*, *critical-stmt*, *end-critical-stmt*, *forall-construct-stmt*, *forall-stmt*, or *where-construct-stmt*.

## 8.2.2 GO TO statement

R857 *goto-stmt* is GO TO *label*

C867 (R857) The *label* shall be the statement label of a *branch target statement* that appears in the same *inclusive scope* as the *goto-stmt*.

1 Execution of a GO TO statement causes a branch to the *branch target statement* identified by the label.

## 8.2.3 Computed GO TO statement

R858 *computed-goto-stmt* is GO TO ( *label-list* ) [ , ] *scalar-int-expr*

C868 (R858) Each *label* in *label-list* shall be the statement label of a *branch target statement* that appears in the same *inclusive scope* as the *computed-goto-stmt*.

1 Execution of a computed GO TO statement causes evaluation of the scalar integer expression. If this value is  $i$  such that  $1 \leq i \leq n$  where  $n$  is the number of labels in *label-list*, a branch occurs to the *branch target statement* identified by the  $i^{th}$  label in the list of labels. If  $i$  is less than 1 or greater than  $n$ , the execution sequence continues as though a **CONTINUE statement** were executed.

### 8.3 CONTINUE statement

- 1 Execution of a CONTINUE statement has no effect.

R859 *continue-stmt* is CONTINUE

### 8.4 STOP and ERROR STOP statements

R860 *stop-stmt* is STOP [ *stop-code* ] [, QUIET = *scalar-logical-expr*]

R861 *error-stop-stmt* is ERROR STOP [ *stop-code* ] [, QUIET = *scalar-logical-expr*]

R862 *stop-code* is *scalar-default-char-expr*  
or *scalar-int-expr*

C869 (R862) The *scalar-int-expr* shall be of default kind.

- 1 Execution of a STOP statement initiates **normal termination** of execution. Execution of an ERROR STOP statement initiates **error termination** of execution.
- 2 When an **image** is terminated by a STOP or ERROR STOP statement, its stop code, if any, is made available in a processor-dependent manner. If the *stop-code* is an integer, it is recommended that the value be used as the process exit status, if the processor supports that concept. If the *stop-code* in a STOP statement is of type character or does not appear, or if an *end-program-stmt* is executed, it is recommended that the value zero be supplied as the process exit status, if the processor supports that concept. If the *stop-code* in an ERROR STOP statement is of type character or does not appear, or if an *end-program-stmt* is executed, it is recommended that a processor-dependent nonzero value be supplied as the process exit status, if the processor supports that concept.
- 3 If QUIET= is omitted or the *scalar-logical-expr* has the value false:
  - if any exception (14) is signaling on that **image**, the processor shall issue a warning indicating which exceptions are signaling, and this warning shall be on the **unit** identified by the named constant **ERROR\_UNIT** from the intrinsic module **ISO\_FORTRAN\_ENV** (13.10.2.9);
  - if a stop code is specified, it is recommended that it be made available by formatted output to the same **unit**.
- 4 If QUIET= appears and the *scalar-logical-expr* has the value true, no output of signaling exceptions or stop code shall be produced.

#### NOTE 8.28

When **normal termination** occurs on more than one **image**, it is expected that a processor-dependent summary of any stop codes and signaling exceptions will be made available.

#### NOTE 8.29

If the integer *stop-code* is used as the process exit status, the processor might be able to interpret only values within a limited range, or only a limited portion of the integer value (for example, only the least-significant 8 bits).

### 8.5 FAIL IMAGE statement

R863 *fail-image-stmt* is FAIL IMAGE

- 1 Execution of a FAIL IMAGE statement causes the executing image to cease participating in program execution without initiating termination. No further statements are executed by that image.

**NOTE 8.30**

The FAIL IMAGE statement allows a program to test a recovery algorithm without needing an actual failure.

On a processor that does not have the ability to detect that an image has failed, execution of a FAIL IMAGE statement might provide a simulated failure environment that provides debug information.

In a piece of code that executes about once a second, invoking this subroutine on an image

```
SUBROUTINE FAIL
  REAL :: X
  CALL RANDOM_NUMBER(X)
  IF (X<0.001) FAIL IMAGE
END SUBROUTINE FAIL
```

will cause that image to have an independent 1/1000 chance of failure every second if the random number generators on different images are independent.

Note that FAIL IMAGE is not an [image control statement](#).

## 8.6 Image execution control

### 8.6.1 Image control statements

1 The execution sequence on each [image](#) is specified in [2.3.5](#).

2 Execution of an [image control statement](#) divides the execution sequence on an [image](#) into segments. Each of the following is an [image control statement](#):

- [SYNC ALL](#) statement;
- [SYNC IMAGES](#) statement;
- [SYNC MEMORY](#) statement;
- [SYNC TEAM](#) statement;
- [ALLOCATE](#) or [DEALLOCATE](#) statement that has a [coarray allocate-object](#);
- [CHANGE TEAM](#) or [END TEAM](#) statement ([8.1.5](#));
- [CRITICAL](#) or [END CRITICAL](#) statement ([8.1.6](#));
- [EVENT POST](#) or [EVENT WAIT](#) statement;
- [FORM TEAM](#) statement;
- [LOCK](#) or [UNLOCK](#) statement;
- any statement that completes execution of a block or procedure and which results in the implicit deallocation of a [coarray](#);
- a [CALL](#) statement that references the intrinsic subroutine [MOVE\\_ALLOC](#) with [coarray](#) arguments;
- [STOP](#) statement;
- [END](#) statement of a main program.

3 All image control statements except [CRITICAL](#), [END CRITICAL](#), [EVENT POST](#), [EVENT WAIT](#), [FORM TEAM](#), [LOCK](#), and [UNLOCK](#) include the effect of executing a [SYNC MEMORY](#) statement ([8.6.5](#)).

4 During an execution of a statement that invokes more than one procedure, at most one invocation shall cause execution of an [image control statement](#) other than [CRITICAL](#) or [END CRITICAL](#).

## 8.6.2 Segments

- 1 On each **image**, the sequence of statements executed before the first execution of an **image control statement**, between the execution of two **image control statements**, or after the last execution of an **image control statement** is a segment. The segment executed immediately before the execution of an **image control statement** includes the evaluation of all expressions within the statement.
- By execution of **image control statements** or user-defined ordering (8.6.5), the program can ensure that the execution of the  $i^{th}$  segment on **image** P,  $P_i$ , either precedes or succeeds the execution of the  $j^{th}$  segment on another **image** Q,  $Q_j$ . If the program does not ensure this, segments  $P_i$  and  $Q_j$  are unordered; depending on the relative execution speeds of the **images**, some or all of the execution of the segment  $P_i$  may take place at the same time as some or all of the execution of the segment  $Q_j$ .
- A **coarray** may be referenced or defined by execution of an **atomic subroutine** during the execution of a segment that is unordered relative to the execution of a segment in which the **coarray** is referenced or defined by execution of an **atomic subroutine**. An **event variable** may be referenced or defined during the execution of a segment that is unordered relative to the execution of another segment in which that **event variable** is defined. Otherwise,
  - if a variable is defined on an **image** in a segment, it shall not be referenced, defined, or become undefined in a segment on another **image** unless the segments are ordered,
  - if the allocation of an **allocatable** subobject of a **coarray** or the **pointer association** of a pointer subobject of a **coarray** is changed on an **image** in a segment, that subobject shall not be referenced or defined in a segment on another **image** unless the segments are ordered, and
  - if a procedure invocation on **image** P is in execution in segments  $P_i, P_{i+1}, \dots, P_k$  and defines a noncoarray **dummy argument**, the **effective argument** shall not be referenced, defined, or become undefined on another **image** Q in a segment  $Q_j$  unless  $Q_j$  precedes  $P_i$  or succeeds  $P_k$ .

### NOTE 8.31

The set of all segments on all **images** is partially ordered: the segment  $P_i$  precedes segment  $Q_j$  if and only if there is a sequence of segments starting with  $P_i$  and ending with  $Q_j$  such that each segment of the sequence precedes the next either because they are consecutive segments on the same **image** or because of the execution of **image control statements**.

### NOTE 8.32

If the segments  $S_1, S_2, \dots, S_k$  on the distinct **images**  $P_1, P_2, \dots, P_k$  are all unordered with respect to each other, it is expected that the processor will ensure that each of these **images** is provided with an equitable share of resources for executing its segment.

### NOTE 8.33

Because of the restrictions on references and definitions in unordered segments, the processor can apply code motion optimizations within a segment as if it were the only **image** in execution, provided calls to **atomic subroutines** are not involved.

### NOTE 8.34

The model upon which the interpretation of a program is based is that there is a permanent memory location for each **coarray** and that all **images** on which it is established can access it.

In practice, apart from executions of **atomic subroutines**, the processor could make a copy of a nonvolatile **coarray** on an **image** (in cache or a register, for example) and, as an optimization, defer copying a changed value back to the permanent memory location while it is still being used. Since the variable is not volatile, it is safe to defer this transfer until the end of the segment and thereafter to reload from permanent memory any **coarray** that was not defined within the segment. It might not be safe to defer these actions beyond the end of the segment since another **image** might reference the variable then.

**NOTE 8.34 (cont.)**

The value of the ATOM argument of an [atomic subroutine](#) might be accessed or modified by another concurrently executing [image](#). Therefore, execution of an [atomic subroutine](#) that references the ATOM argument cannot rely on a local copy, but instead always gets its value from its permanent memory location. Execution of an [atomic subroutine](#) that defines the ATOM argument does not complete until the value of its ATOM argument has been sent to its permanent memory location.

**NOTE 8.35**

The incorrect sequencing of [image control statements](#) can suspend execution indefinitely. For example, one [image](#) might be executing a [SYNC ALL statement](#) while another is executing an [ALLOCATE statement](#) for a [coarray](#).

**8.6.3 SYNC ALL statement**

R864 *sync-all-stmt*                      **is** SYNC ALL [ ( [ *sync-stat-list* ] ) ]

R865 *sync-stat*                      **is** STAT = *stat-variable*  
    **or** ERRMSG = *errmsg-variable*

C870 No specifier shall appear more than once in a given *sync-stat-list*.

1 The **STAT=** and **ERRMSG=** specifiers for [image control statements](#) are described in 8.6.11.

2 Execution of a SYNC ALL statement performs a synchronization of all [images](#) in the [current team](#). Execution on an [image](#), M, of the segment following the SYNC ALL statement is delayed until each other [image](#) in the [current team](#) has executed a SYNC ALL statement as many times as has [image](#) M since execution last began in this [team](#). The segments that executed before the SYNC ALL statement on an [image](#) precede the segments that execute after the SYNC ALL statement on another [image](#).

**Unresolved Technical Issue 022****Incompatible change of semantics.**

According to F2008, SYNC ALL operates on all images, and if there is a stopped image it returns STAT\_STOPPED\_IMAGE, and does not do any syncing, just doing the effect of SYNC MEMORY.

According to the TS, SYNC ALL only operates on active images. Which means it does not do any “interaction with an image that has initiated termination”. Which means that (assuming no other errors) it should assign zero to the *stat-variable* and do the syncing.

Oh yes, and it means that it ought not to detect failed images either.

We could keep the ignoring of failed images by inserting “nonfailed” above, but that likewise seems problematic, and is unlikely to be our intent.

Therefore, I did not insert “active” in the above paragraph. If this is correct, little more needs to be done (but see last para of this UTI).

Otherwise, further edits will be required, and if it is supposed to be incompatible, edits to the compatibility subclauses.

And if we want SYNC ALL to detect stopped/failed but to synchronise the active images even when stopped (or failed) images exist, the contradiction between the wording here and in the STAT= subclause needs to be resolved, as well as adding to the incompatibility list.

Also, in any case the paragraph should begin with the word “Successful”.

**NOTE 8.36**

The processor might have special hardware or employ an optimized algorithm to make the SYNC ALL statement execute efficiently.

Here is a simple example of its use. [Image 1](#) reads data and broadcasts it to other [images](#):

```
REAL :: P[*]
...
SYNC ALL
IF (THIS_IMAGE()==1) THEN
  READ (*,*) P
  DO I = 2, NUM_IMAGES()
    P[I] = P
  END DO
END IF
SYNC ALL
```

**8.6.4 SYNC IMAGES statement**

R866 *sync-images-stmt* is SYNC IMAGES ( *image-set* [ , *sync-stat-list* ] )

R867 *image-set* is *int-expr*  
or \*

C871 An *image-set* that is an *int-expr* shall be scalar or of [rank](#) one.

1 If *image-set* is an array expression, the value of each element shall be positive and not greater than the number of [images](#) in the [current team](#), and there shall be no repeated values.

2 If *image-set* is a scalar expression, its value shall be positive and not greater than the number of [images](#) in the [current team](#).

3 An *image-set* that is an asterisk specifies all [images](#) in the [current team](#).

4 Execution of a SYNC IMAGES statement performs a synchronization of the [image](#) with each of the other [images](#) in the *image-set*. Executions of SYNC IMAGES statements on [images](#) M and T correspond if the number of times [image](#) M has executed a SYNC IMAGES statement in the [current team](#) with T in its [image](#) set is the same as the number of times [image](#) T has executed a SYNC IMAGES statement with M in its [image](#) set since execution last began in this [team](#). The segments that executed before the SYNC IMAGES statement on either [image](#) precede the segments that execute after the corresponding SYNC IMAGES statement on the other [image](#).

**NOTE 8.37**

A SYNC IMAGES statement that specifies the single [image index](#) value `THIS_IMAGE ( )` in its [image](#) set is allowed. This simplifies writing programs for an arbitrary number of [images](#) by allowing correct execution in the limiting case of the number of [images](#) being equal to one.

**NOTE 8.38**

In a program that uses SYNC ALL as its only synchronization mechanism, every SYNC ALL statement could be replaced by a SYNC IMAGES (\*) statement, but SYNC ALL might give better performance.

SYNC IMAGES statements are not required to specify the entire [image](#) set, or even the same [image](#) set, on all [images](#) participating in the synchronization. In the following example, [image 1](#) will wait for each of the other [images](#) to complete its use of the data. The other [images](#) wait for [image 1](#) to set up the data, but do not wait on any other [image](#).

## NOTE 8.38 (cont.)

**Unresolved Technical Issue 023****Glitches in SYNC IMAGES**

1. I declined to insert “active” since that would have been an incompatible change in semantics, pace UTI 022. However that is resolved, this needs to be resolved the same way.
2. “last began” really needs to go; either it is unnecessary or it needs to be firmly nailed down. Also seems to be missing “in the current team” after “T has executed a SYNC IMAGES statement”.

```

IF (THIS_IMAGE() == 1) then
  ! Set up coarray data needed by all other images
  SYNC IMAGES(*)
ELSE
  SYNC IMAGES(1)
  ! Use the data set up by image 1
END IF

```

When the following example runs on five or more `images`, each `image` synchronizes with both of its neighbors, in a circular fashion.

```

INTEGER :: up, down
...
IF (NUM_IMAGES()>1) THEN
  up = THIS_IMAGE()+1; IF (up>NUM_IMAGES()) up = 1
  down = THIS_IMAGE()-1; IF (down==0) down = NUM_IMAGES()
  SYNC IMAGES ( (/ up, down /) )
END IF

```

This might appear to have the same effect as `SYNC ALL` but there is no ordering between the preceding and succeeding segments on non-adjacent `images`. For example, the segment preceding the `SYNC IMAGES statement` on `image 3` will be ordered before those succeeding it on `images 2` and `4`, but not those on `images 1` and `5`.

## NOTE 8.39

In the following example, each `image` synchronizes with its neighbor.

```

INTEGER :: ME, NE, STEP, NSTEPS
NE = NUM_IMAGES()
ME = THIS_IMAGE()
  ! Initial calculation
SYNC ALL
DO STEP = 1, NSTEPS
  IF (ME > 1) SYNC IMAGES(ME-1)
    ! Perform calculation
  IF (ME < NE) SYNC IMAGES(ME+1)
END DO
SYNC ALL

```

The calculation starts on `image 1` since all the others will be waiting on `SYNC IMAGES (ME-1)`. When this is done, `image 2` can start and `image 1` can perform its second calculation. This continues until they are all executing different steps at the same time. Eventually, `image 1` will finish and then the others will finish one by one.

## 8.6.5 SYNC MEMORY statement

- 1 Execution of a SYNC MEMORY statement ends one segment and begins another; those two segments can be ordered by a user-defined way with respect to segments on other **images**.

R868 *sync-memory-stmt* is SYNC MEMORY [ ( [ *sync-stat-list* ] ) ]

- 2 If, by execution of statements on **image** P,

- a variable X on **image** Q is defined, referenced, becomes undefined, or has its allocation status, **pointer association** status, array bounds, **dynamic type**, or type parameters changed or inquired about by execution of a statement,
- that statement precedes a successful execution of a SYNC MEMORY statement, and
- a variable Y on **image** Q is defined, referenced, becomes undefined, or has its allocation status, **pointer association** status, array bounds, **dynamic type**, or type parameters changed or inquired about by execution of a statement that succeeds execution of that SYNC MEMORY statement,

then the action regarding X on **image** Q precedes the action regarding Y on **image** Q.

- 3 User-defined ordering of segment  $P_i$  on **image** P to precede segment  $Q_j$  on **image** Q occurs when

- **image** P executes an **image control statement** that ends segment  $P_i$ , and then executes statements that initiate a cooperative synchronization between **images** P and Q, and
- **image** Q executes statements that complete the cooperative synchronization between **images** P and Q and then executes an **image control statement** that begins segment  $Q_j$ .

- 4 Execution of the cooperative synchronization between **images** P and Q shall include a dependency that forces execution on **image** P of the statements that initiate the synchronization to precede the execution on **image** Q of the statements that complete the synchronization. The mechanisms available for creating such a dependency are processor dependent.

### NOTE 8.40

SYNC MEMORY usually suppresses compiler optimizations that might reorder memory operations across the segment boundary defined by the SYNC MEMORY statement and ensures that all memory operations initiated in the preceding segments in its **image** complete before any memory operations in the subsequent segment in its **image** are initiated. It needs to do this unless it can establish that failure to do so could not alter processing on another **image**.

### NOTE 8.41

SYNC MEMORY can be used to implement specialized schemes for segment ordering, such as the spin-wait loop. For example:

```
USE, INTRINSIC :: ISO_FORTRAN_ENV
LOGICAL(ATOMICAL_KIND), SAVE :: LOCKED[*] = .TRUE.
LOGICAL :: VAL
INTEGER :: IAM, P, Q
...
IAM = THIS_IMAGE()
IF (IAM == P) THEN
    SYNC MEMORY
    CALL ATOMIC_DEFINE (LOCKED[Q], .FALSE.)
ELSE IF (IAM == Q) THEN
    VAL = .TRUE.
    DO WHILE (VAL)
    CALL ATOMIC_REF (VAL, LOCKED)
```

! Segment  $P_i$   
! A  
! Segment  $P_{i+1}$   
! Segment  $Q_{j-1}$

**NOTE 8.41 (cont.)**

```

END DO
SYNC MEMORY                                ! B
                                           ! Segment  $Q_j$ 
END IF

```

The **DO WHILE** loop does not complete until VAL is defined with the value false. This is the cooperative synchronization that provides the dependency that **image** Q does not complete segment  $Q_{j-1}$  until the **CALL statement** in segment  $P_{i+1}$  completes. This ensures that the execution of segment  $P_i$  on **image** P precedes execution of segment  $Q_j$  on **image** Q.

The first SYNC MEMORY statement (A) ensures that the compiler does not reorder the following statement (segment  $P_{i+1}$ ) with the previous statements, since the lock is supposed to be freed only after the work in segment  $P_i$  has been completed.

The second SYNC MEMORY statement (B) marks the beginning of a new segment, informing the compiler that the values of **coarrays** referenced in that segment might have been changed by other **images** in preceding segments, so need to be loaded from memory.

**NOTE 8.42**

As a second example, the user might have access to an **external procedure** that performs synchronization between **images**. That library procedure might not be aware of the mechanisms used by the processor to manage remote data references and definitions, and therefore not, by itself, be able to ensure the correct memory state before and after its reference. The SYNC MEMORY statement provides the needed memory ordering that enables the safe use of the external synchronization routine. For example:

```

INTEGER :: IAM
REAL    :: X[*]

IAM = THIS_IMAGE()
IF (IAM == 1) X = 1.0
SYNC MEMORY
CALL EXTERNAL_SYNC()
SYNC MEMORY
IF (IAM == 2) WRITE(*,*) X[1]

```

where executing the subroutine EXTERNAL\_SYNC has an **image** synchronization effect similar to executing a **SYNC ALL statement**.

**8.6.6 SYNC TEAM statement**

R869 *sync-team-stmt* **is** SYNC TEAM ( *team-variable* [, *sync-stat-list*] )

- 1 The *team-variable* shall be defined with a value that represents an **ancestor team**, the **current team**, or a **team** whose **parent** is the **current team**.
- 2 Successful execution of a SYNC TEAM statement performs a synchronization of the team specified by *team-variable*. Execution on an image, M, of the segment following the SYNC TEAM statement is delayed until each other image of the specified team has executed a SYNC TEAM statement specifying the same team as many times as has image M since execution last began in this team. The segments that executed before the SYNC TEAM statement on an image precede the segments that execute after the SYNC TEAM statement on another image.

## Unresolved Technical Issue 024

## SYNC TEAM shares flaws with SYNC ALL and SYNC IMAGES

1. It tried to synchronise all the active images, but according to the STAT= subclause, SYNC TEAM does no syncing (acts like SYNC MEMORY) on error conditions (which includes stopped/failed images).
2. “last began” needs to be sorted out.
3. The last sentence is ambiguous (lacking well-defined antecedents for the second “the SYNC TEAM statement”).

## NOTE 8.43

A SYNC TEAM statement synchronizes a particular team whereas a SYNC ALL statement synchronizes the current team.

## 8.6.7 EVENT POST statement

1 The EVENT POST statement posts an event.

R870 *event-post-stmt* is EVENT POST ( *event-variable* [ , *sync-stat-list* ] )

R871 *event-variable* is scalar-variable

C872 (R871) An *event-variable* shall be of type EVENT\_TYPE from the intrinsic module ISO\_FORTRAN\_ENV(13.10.2).

2 Successful execution of an EVENT POST statement atomically increments the count of the *event variable* by one. If an error condition occurs during execution of an EVENT POST statement, the value of the count of the *event variable* is processor dependent. The completion of an EVENT POST statement does not depend on the execution of a corresponding EVENT WAIT statement.

## 8.6.8 EVENT WAIT statement

1 The EVENT WAIT statement waits until an event is posted.

R872 *event-wait-stmt* is EVENT WAIT ( *event-variable* [ , *event-wait-spec-list* ] )

R873 *event-wait-spec* is *until-spec*  
or *sync-stat*

R874 *until-spec* is UNTIL\_COUNT = *scalar-int-expr*

C873 (R872) The *event-variable* in an *event-wait-stmt* shall not be coindexed.

C874 (R872) No specifier shall appear more than once in an *event-wait-spec-list*.

2 Execution of an EVENT WAIT statement consists of the following sequence of actions:

1. if the UNTIL\_COUNT= specifier does not appear, the threshold value is set to one; otherwise, the threshold value is set to the maximum of the value of the *scalar-int-expr* and one;
2. the executing image waits until the count of the *event variable* is greater than or equal to the threshold value or an error condition occurs;
3. if no error condition occurs, the count of the *event variable* is atomically decremented by the threshold value.

1 3 If an error occurs during execution of an EVENT WAIT statement, the value of the count of its *event variable*  
2 is processor dependent.

3 4 An *EVENT POST statement* execution is initially unsatisfied. Successful execution of an EVENT WAIT state-  
4 ment with a threshold of  $k$  satisfies the first  $k$  unsatisfied *EVENT POST statement* executions for that *event*  
5 *variable*. This EVENT WAIT statement execution causes the segment following the EVENT WAIT statement  
6 execution to succeed the segments preceding those  $k$  *EVENT POST statement* executions.

## 7 8.6.9 FORM TEAM statement

8 1 The FORM TEAM statement creates new teams whose *parent team* is the *current team*.

9 R875 *form-team-stmt* is FORM TEAM ( *team-number*, *team-variable* ■  
10 ■ [, *form-team-spec-list* ] )

11 R876 *team-number* is *scalar-int-expr*

12 R877 *form-team-spec* is NEW\_INDEX = *scalar-int-expr*  
13 or *sync-stat*

14 C875 No specifier shall appear more than once in given a *form-team-spec-list*.

15 2 Successful execution of a FORM TEAM statement creates a new team for each unique *team-number* value specified  
16 by the *active images* of the *current team*. The value of *team-number* shall be positive. Each executing image will  
17 belong to the team whose *team number* is equal to the value of *team-number* on that image, and the *team-variable*  
18 becomes defined with a value that specifies that team.

19 3 The value of the *scalar-int-expr* in a NEW\_INDEX= specifier specifies the *image index* that the executing image  
20 will have in its new *team*. It shall be positive, less than or equal to the number of *images* in the *team*, and  
21 different from the value specified by every other *image* that belongs to that *team*.

22 4 If the NEW\_INDEX= specifier does not appear, the *image index* of the executing image in the new *team* is  
23 processor dependent. This *image index* will be positive, less than or equal to the number of *images* in the *team*,  
24 and different from that of every other *image* in the *team*.

25 5 If the FORM TEAM statement is executed on one *image*, the same statement shall be executed on all *active*  
26 *images* of the *current team*. When a FORM TEAM statement is executed, there is an implicit synchronization  
27 of all *active images* in the *current team*. On those *images*, execution of the segment following the statement is  
28 delayed until all other *active images* in the *current team* have executed the same statement the same number of  
29 times since execution last began in this team. If an error condition other than detection of a failed image occurs,  
30 the team variable becomes undefined.

31 6 If execution of a FORM TEAM statement assigns the value *STAT\_FAILED\_IMAGE* to the *stat-variable*, the  
32 effect is the same as for the successful execution of FORM TEAM except for the value assigned to *stat-variable*.

### NOTE 8.44

Executing the statement

```
FORM TEAM ( 2-MOD(THIS_IMAGE(),2), ODD_EVEN )
```

will create two subteams of the *current team*, with *images* whose image index is odd being in the team with number 1, and those with an even image index being in the team with number 2.

### NOTE 8.45

If the current team consists of  $P^2$  images, with corresponding coarrays on each image representing parts of a larger array spread over a  $P \times P$  square, the following code will establish teams for the rows with image indices equal to the column indices.

## NOTE 8.45 (cont.)

```

USE, INTRINSIC :: ISO_FORTRAN_ENV
TYPE(Team_Type) :: Row
REAL :: A [P, *]
INTEGER :: ME (2)
ME (:) = THIS_IMAGE (A)
FORM Team (ME(1), Row, NEW_INDEX=ME(2))

```

## 8.6.10 LOCK and UNLOCK statements

R878 *lock-stmt* is LOCK ( *lock-variable* [ , *lock-stat-list* ] )

R879 *lock-stat* is ACQUIRED\_LOCK = *scalar-logical-variable*  
or *sync-stat*

C876 No specifier shall appear more than once in a given *lock-stat-list*.

R880 *unlock-stmt* is UNLOCK ( *lock-variable* [ , *sync-stat-list* ] )

R881 *lock-variable* is *scalar-variable*

C877 (R881) A *lock-variable* shall be of type LOCK\_TYPE from the intrinsic module ISO\_FORTRAN\_ENV (13.10.2.19).

- 1 A lock variable is unlocked if and only if the value of each component is the same as its default value. If it has any other value, it is locked. A lock variable is locked by an *image* if it was locked by execution of a LOCK statement on that *image*, has not been subsequently unlocked by execution of an UNLOCK statement on the same *image*, and that *image* has not failed.
- 2 Successful execution of a LOCK statement without an ACQUIRED\_LOCK= specifier causes the lock variable to become locked by that *image*. If the lock variable is already locked by another *image*, that LOCK statement causes the lock variable to become locked after the other *image* causes the lock variable to become unlocked.
- 3 If the lock variable is unlocked, successful execution of a LOCK statement with an ACQUIRED\_LOCK= specifier causes the lock variable to become locked by that *image* and the scalar logical variable to become defined with the value true. If the lock variable is already locked by a different *image*, successful execution of a LOCK statement with an ACQUIRED\_LOCK= specifier leaves the lock variable unchanged and causes the scalar logical variable to become defined with the value false.
- 4 Successful execution of an UNLOCK statement causes the lock variable to become unlocked. Failure of an *image* causes all lock variables that are locked by that *image* to become unlocked.
- 5 During execution of the program, the value of a lock variable changes through a sequence of locked and unlocked states due to the execution of LOCK and UNLOCK statements. If a lock variable becomes unlocked by execution of an UNLOCK statement on *image* M and next becomes locked by execution of a LOCK statement on *image* T, the segments preceding the UNLOCK statement on *image* M precede the segments following the LOCK statement on *image* T. Execution of a LOCK statement that does not cause the lock variable to become locked does not affect segment ordering.
- 6 An error condition occurs if the lock variable in a LOCK statement is already locked by the executing *image*. An error condition occurs if the lock variable in an UNLOCK statement is not already locked by the executing *image*. If an error condition occurs during execution of a LOCK or UNLOCK statement, the value of the lock variable is not changed and the value of the ACQUIRED\_LOCK variable, if any, is not changed.

## NOTE 8.46

A lock variable is effectively defined atomically by a LOCK or UNLOCK statement. If LOCK statements on two *images* both attempt to acquire a lock, one will succeed and the other will either fail if an ACQUIRED\_ -

**NOTE 8.46 (cont.)**

LOCK= specifier appears, or will wait until the lock is later released if an ACQUIRED\_LOCK= specifier does not appear.

**NOTE 8.47**

An [image](#) might wait for a LOCK statement to successfully complete for a long period of time if other [images](#) frequently lock and unlock the same lock variable. This situation might result from executing LOCK statements with ACQUIRED\_LOCK= specifiers inside a spin loop.

**NOTE 8.48**

The following example illustrates the use of LOCK and UNLOCK statements to manage a work queue:

```
USE, INTRINSIC :: ISO_FORTRAN_ENV

TYPE(LOCK_TYPE) :: queue_lock[*] ! Lock on each image to manage its work queue
INTEGER :: work_queue_size[*]
TYPE(Task) :: work_queue(100)[*] ! List of tasks to perform

TYPE(Task) :: job ! Current task working on
INTEGER :: me

me = THIS\_IMAGE()
DO
  ! Process the next item in your work queue

  LOCK (queue_lock) ! New segment A starts
  ! This segment A is ordered with respect to
  ! segment B executed by image me-1 below because of lock exclusion
  IF (work_queue_size>0) THEN
    ! Fetch the next job from the queue
    job = work_queue(work_queue_size)
    work_queue_size = work_queue_size-1
  END IF
  UNLOCK (queue_lock) ! Segment ends
  ... ! Actually process the task

  ! Add a new task on neighbors queue:
  LOCK(queue_lock[me+1]) ! Starts segment B
  ! This segment B is ordered with respect to
  ! segment A executed by image me+1 above because of lock exclusion
  IF (work_queue_size[me+1]<SIZE(work_queue)) THEN
    work_queue_size[me+1] = work_queue_size[me+1]+1
    work_queue(work_queue_size[me+1])[me+1] = job
  END IF
  UNLOCK (queue_lock[me+1]) ! Ends segment B

END DO
```

## 1 8.6.11 STAT= and ERRMSG= specifiers in image control statements

### Unresolved Technical Issue 025

STAT= with a positive value needs to be an error indication

## Unresolved Technical Issue 025 (cont.)

None of the TS changes have been applied to this subclause, as

- the basic idea that we have “successful” execution returning an error indication is completely unacceptable;
- the text is massively over-complicated and internally inconsistent (this is likely why no-one complained about the former issue).

I have no problem if we want to say that “if *stat-variable* is assigned the value STAT\_FAILED\_IMAGE, the teams involving the other images are still formed; otherwise, the effect is that of SYNC MEMORY”. I would prefer not to have STAT\_FAILED\_IMAGE require the synchronization to take place for SYNC ALL et al, for consistency with STAT\_STOPPED\_IMAGE and so as not to hamper more primitive error recovery schemes. But I draw the line at pretending that detected image failure is not an error, or conversely that successful execution sets *stat-variable* to anything other than zero.

- 1 If the STAT= specifier appears, successful execution of the LOCK, SYNC ALL, SYNC IMAGES, SYNC MEMORY, or UNLOCK statement causes the specified variable to become defined with the value zero.
- 2 If the STAT= specifier appears in a SYNC ALL or SYNC IMAGES statement and execution of one of these statements involves synchronization with an image that has initiated normal termination, the variable becomes defined with the value of the constant STAT\_STOPPED\_IMAGE (13.10.2.29) in the intrinsic module ISO\_FORTRAN\_ENV(13.10.2), and the effect of executing the statement is otherwise the same as that of executing the SYNC MEMORY statement. If any other error condition occurs during execution of one of these statements, the variable becomes defined with a processor-dependent positive integer value that is different from the value of STAT\_STOPPED\_IMAGE.
- 3 If the STAT= specifier appears in a LOCK statement and the lock variable is locked by the executing image, the specified variable becomes defined with the value of STAT\_LOCKED (13.10.2.27). If the STAT= specifier appears in an UNLOCK statement and the lock variable has the value unlocked, the variable specified by the STAT= specifier becomes defined with the value of STAT\_UNLOCKED (13.10.2.30). If the STAT= specifier appears in an UNLOCK statement and the lock variable is locked by a different image, the specified variable becomes defined with the value STAT\_LOCKED\_OTHER\_IMAGE (13.10.2.28). The named constants STAT\_LOCKED, STAT\_UNLOCKED, and STAT\_LOCKED\_OTHER\_IMAGE are defined in the intrinsic module ISO\_FORTRAN\_ENV. If any other error condition occurs during execution of a LOCK or UNLOCK statement, the specified variable becomes defined with a positive integer value that is different from STAT\_LOCKED, STAT\_UNLOCKED, and STAT\_LOCKED\_OTHER\_IMAGE.
- 4 If an error condition occurs during execution of an EVENT POST, EVENT WAIT, LOCK, SYNC ALL, SYNC IMAGES, SYNC MEMORY, or UNLOCK statement that does not contain the STAT= specifier, error termination is initiated.
- 5 If an ERRMSG= specifier appears in an EVENT POST, EVENT WAIT, LOCK, SYNC ALL, SYNC IMAGES, SYNC MEMORY, or UNLOCK statement, and an error condition occurs during execution of that statement, the specified variable is assigned an explanatory message as if by intrinsic assignment. If no such condition occurs, the definition status and value of the specified variable are unchanged.
- 6 The set of error conditions that can occur in an image control statement is processor dependent.

## NOTE 8.49

A processor might detect communication failure between images and treat it as an error condition. A processor might also treat an invalid set of images in a SYNC IMAGES statement as an error condition.

1 (Blank page)

## 9 Input/output statements

### 9.1 Input/output concepts

- 1 **Input statements** provide the means of transferring data from external media to internal storage or from an **internal file** to internal storage. This process is called reading. **Output statements** provide the means of transferring data from internal storage to external media or from internal storage to an **internal file**. This process is called writing. Some **input/output statements** specify that editing of the data is to be performed.
- 2 In addition to the statements that transfer data, there are auxiliary input/output statements to manipulate the external medium, or to describe or inquire about the properties of the connection to the external medium.
- 3 The input/output statements are the **BACKSPACE**, **CLOSE**, **ENDFILE**, **FLUSH**, **INQUIRE**, **OPEN**, **PRINT**, **READ**, **REWIND**, **WAIT**, and **WRITE** statements.
- 4 A file is composed of either a sequence of **file storage units** (9.3.5) or a sequence of records, which provide an extra level of organization to the file. A file composed of records is called a **record file**. A file composed of **file storage units** is called a **stream file**. A processor may allow a file to be viewed both as a **record file** and as a **stream file**; in this case the relationship between the **file storage units** when viewed as a **stream file** and the records when viewed as a **record file** is processor dependent.
- 5 A file is either an **external file** (9.3) or an **internal file** (9.4).

### 9.2 Records

#### 9.2.1 Definition of a record

- 1 A **record** is a sequence of values or a sequence of characters. For example, a line on a terminal is usually considered to be a record. However, a record does not necessarily correspond to a physical entity. There are three kinds of records:
  - (1) formatted;
  - (2) unformatted;
  - (3) endfile.

#### NOTE 9.1

What is called a “record” in Fortran is commonly called a “logical record”. There is no concept in Fortran of a “physical record.”

#### 9.2.2 Formatted record

- 1 A formatted record consists of a sequence of characters that are representable in the processor; however, a processor may prohibit some control characters (3.1.1) from appearing in a formatted record. The length of a formatted record is measured in characters and depends primarily on the number of characters put into the record when it is written. However, it may depend on the processor and the external medium. The length may be zero. Formatted records shall be read or written only by **formatted input/output statements**.

#### 9.2.3 Unformatted record

- 1 An unformatted record consists of a sequence of values in a processor-dependent form and may contain data of any type or may contain no data. The length of an unformatted record is measured in **file storage units**

(9.3.5) and depends on the output list (9.6.3) used when it is written, as well as on the processor and the external medium. The length may be zero. Unformatted records may be read or written only by [unformatted input/output statements](#).

## 9.2.4 Endfile record

- 1 An endfile record is written explicitly by the [ENDFILE statement](#); the file shall be [connected](#) for sequential access. An endfile record is written implicitly to a file [connected](#) for sequential access when the most recent [data transfer statement](#) referring to the file is an [output statement](#), no intervening [file positioning statement](#) referring to the file has been executed, and

- a [REWIND](#) or [BACKSPACE](#) statement references the [unit](#) to which the file is [connected](#), or
- the [unit](#) is closed, either explicitly by a [CLOSE statement](#), implicitly by [normal termination](#), or implicitly by another [OPEN statement](#) for the same [unit](#).

- 2 An endfile record may occur only as the last record of a file. An endfile record does not have a length property.

### NOTE 9.2

An endfile record does not necessarily have any physical embodiment. The processor can use a record count or any other means to register the position of the file at the time an [ENDFILE statement](#) is executed, so that it can take appropriate action when that position is reached again during a read operation. The endfile record, however it is implemented, is considered to exist for the [BACKSPACE statement](#) (9.8.2).

## 9.3 External files

### 9.3.1 External file concepts

- 1 An [external file](#) is any file that exists in a medium external to the program.
- 2 At any given time, there is a processor-dependent set of allowed access methods, a processor-dependent set of allowed forms, a processor-dependent set of allowed actions, and a processor-dependent set of allowed record lengths for a file.

### NOTE 9.3

For example, the processor-dependent set of allowed actions for a printer would likely include the write action, but not the read action.

- 3 A file may have a name; a file that has a name is called a named file. The name of a named file is represented by a character string value. The set of allowable names for a file is processor dependent. Whether a named file on one [image](#) is the same as a file with the same name on another [image](#) is processor dependent.

### NOTE 9.4

If different files are needed on each [image](#), using a different file name on each [image](#) will improve portability of the code. One technique is to incorporate the [image index](#) as part of the name.

- 4 An [external file](#) that is [connected](#) to a [unit](#) has a position property (9.3.4).

### NOTE 9.5

For more explanatory information on [external files](#), see C.6.1.

### 9.3.2 File existence

- 1 At any given time, there is a processor-dependent set of [external files](#) that exist for a program. A file may be known to the processor, yet not exist for a program at a particular time.

1 2 To create a file means to cause a file to exist that did not exist previously. To delete a file means to terminate  
2 the existence of the file.

3 3 All input/output statements may refer to files that exist. A [CLOSE](#), [ENDFILE](#), [FLUSH](#), [INQUIRE](#), [OPEN](#),  
4 [PRINT](#), [REWIND](#), or [WRITE](#) statement is permitted to refer to a file that does not exist. No other input/output  
5 statement shall refer to a file that does not exist. Execution of a [WRITE](#), [PRINT](#), or [ENDFILE](#) statement  
6 referring to a [preconnected](#) file that does not exist creates the file. This file is a different file from one [preconnected](#)  
7 on any other [image](#).

### 8 9.3.3 File access

#### 9 9.3.3.1 File access methods

10 1 There are three methods of accessing the data of an [external file](#): sequential, direct, and stream. Some files may  
11 have more than one allowed access method; other files may be restricted to one access method.

#### NOTE 9.6

For example, a processor might provide only sequential access to a file on magnetic tape. Thus, the set of allowed access methods depends on the file and the processor.

12 2 The method of accessing a file is determined when the file is [connected](#) to a [unit](#) (9.5.4) or when the file is created  
13 if the file is [preconnected](#) (9.5.5).

#### 14 9.3.3.2 Sequential access

15 1 Sequential access is a method of accessing the records of an external [record file](#) in order.

16 2 While [connected](#) for sequential access, an [external file](#) has the following properties.

- 17 • The order of the records is the order in which they were written if the direct access method is not a member  
18 of the set of allowed access methods for the file. If the direct access method is also a member of the set of  
19 allowed access methods for the file, the order of the records is the same as that specified for direct access.  
20 In this case, the first record accessible by sequential access is the record whose record number is 1 for direct  
21 access. The second record accessible by sequential access is the record whose record number is 2 for direct  
22 access, etc. A record that has not been written since the file was created shall not be read.
- 23 • The records of the file are either all formatted or all unformatted, except that the last record of the file may  
24 be an endfile record. Unless the previous reference to the file was an [output statement](#), the last record, if  
25 any, of the file shall be an endfile record.
- 26 • The records of the file shall be read or written only by sequential access [data transfer statements](#).

#### 27 9.3.3.3 Direct access

28 1 Direct access is a method of accessing the records of an external [record file](#) in arbitrary order.

29 2 While [connected](#) for direct access, an [external file](#) has the following properties.

- 30 • Each record of the file is uniquely identified by a positive integer called the record number. The record  
31 number of a record is specified when the record is written. Once established, the record number of a record  
32 can never be changed. The order of the records is the order of their record numbers.
- 33 • The records of the file are either all formatted or all unformatted. If the sequential access method is also a  
34 member of the set of allowed access methods for the file, its endfile record, if any, is not considered to be  
35 part of the file while it is [connected](#) for direct access. If the sequential access method is not a member of  
36 the set of allowed access methods for the file, the file shall not contain an endfile record.
- 37 • The records of the file shall be read or written only by direct access [data transfer statements](#).
- 38 • All records of the file have the same length.

- Records need not be read or written in the order of their record numbers. Any record may be written into the file while it is [connected](#) to a [unit](#). For example, it is permissible to write record 3, even though records 1 and 2 have not been written. Any record may be read from the file while it is [connected](#) to a [unit](#), provided that the record has been written since the file was created, and if a [READ statement](#) for this connection is permitted.
- The records of the file shall not be read or written using list-directed formatting ([10.10](#)), namelist formatting ([10.11](#)), or a nonadvancing [data transfer statement](#) ([9.3.4.2](#)).

**NOTE 9.7**

A record cannot be deleted; however, a record can be rewritten.

**9.3.3.4 Stream access**

- 1 Stream access is a method of accessing the [file storage units](#) ([9.3.5](#)) of an external [stream file](#).
- 2 The properties of an [external file connected](#) for stream access depend on whether the connection is for unformatted or formatted access. While connected for stream access, the [file storage units](#) of the file shall be read or written only by stream access [data transfer statements](#).
- 3 While [connected](#) for unformatted stream access, an [external file](#) has the following properties.
  - Each [file storage unit](#) in the file is uniquely identified by a positive integer called the position. The first [file storage unit](#) in the file is at position 1. The position of each subsequent [file storage unit](#) is one greater than that of its preceding [file storage unit](#).
  - If it is possible to position the file, the [file storage units](#) need not be read or written in order of their position. For example, it might be permissible to write the [file storage unit](#) at position 3, even though the [file storage units](#) at positions 1 and 2 have not been written. Any [file storage unit](#) may be read from the file while it is [connected](#) to a [unit](#), provided that the [file storage unit](#) has been written since the file was created, and if a [READ statement](#) for this connection is permitted.
- 4 While [connected](#) for formatted stream access, an [external file](#) has the following properties.
  - Some [file storage units](#) of the file may contain record markers; this imposes a record structure on the file in addition to its stream structure. There might or might not be a record marker at the end of the file. If there is no record marker at the end of the file, the final record is incomplete.
  - No maximum length ([9.5.6.15](#)) is applicable to these records.
  - Writing an empty record with no record marker has no effect.
  - Each [file storage unit](#) in the file is uniquely identified by a positive integer called the position. The first [file storage unit](#) in the file is at position 1. The relationship between positions of successive [file storage units](#) is processor dependent; not all positive integers need correspond to valid positions.
  - If it is possible to position the file, the file position can be set to a position that was previously identified by the [POS= specifier](#) in an [INQUIRE statement](#).
  - A processor may prohibit some control characters ([3.1.1](#)) from appearing in a formatted [stream file](#).

**NOTE 9.8**

Because the record structure is determined from the record markers that are stored in the file itself, an incomplete record at the end of the file is necessarily not empty.

**NOTE 9.9**

There might be some character positions in the file that do not correspond to characters written; this is because on some processors a record marker could be written to the file as a carriage-return/line-feed or other sequence. The means of determining the position in a file [connected](#) for stream access is via the [POS= specifier](#) in an [INQUIRE statement](#) ([9.10.2.22](#)).

### 9.3.4 File position

#### 9.3.4.1 General

- 1 Execution of certain input/output statements affects the position of an [external file](#). Certain circumstances can cause the position of a file to become indeterminate.
- 2 The initial point of a file is the position just before the first record or [file storage unit](#). The terminal point is the position just after the last record or [file storage unit](#). If there are no records or [file storage units](#) in the file, the initial point and the terminal point are the same position.
- 3 If a [record file](#) is positioned within a record, that record is the current record; otherwise, there is no current record.
- 4 Let  $n$  be the number of records in the file. If  $1 < i \leq n$  and a file is positioned within the  $i$ th record or between the  $(i - 1)$ th record and the  $i$ th record, the  $(i - 1)$ th record is the preceding record. If  $n \geq 1$  and the file is positioned at its terminal point, the preceding record is the  $n$ th and last record. If  $n = 0$  or if a file is positioned at its initial point or within the first record, there is no preceding record.
- 5 If  $1 \leq i < n$  and a file is positioned within the  $i$ th record or between the  $i$ th and  $(i + 1)$ th record, the  $(i + 1)$ th record is the next record. If  $n \geq 1$  and the file is positioned at its initial point, the first record is the next record. If  $n = 0$  or if a file is positioned at its terminal point or within the  $n$ th (last) record, there is no next record.
- 6 For a file [connected](#) for stream access, the file position is either between two [file storage units](#), at the initial point of the file, at the terminal point of the file, or undefined.

#### 9.3.4.2 Advancing and nonadvancing input/output

- 1 An advancing input/output statement always positions a [record file](#) after the last record read or written, unless there is an error condition.
- 2 A nonadvancing input/output statement may position a [record file](#) at a character position within the current record, or a subsequent record ([10.8.2](#)). Using nonadvancing input/output, it is possible to read or write a record of the file by a sequence of [data transfer statements](#), each accessing a portion of the record. It is also possible to read variable-length records and be notified of their lengths. If a nonadvancing [output statement](#) leaves a file positioned within a current record and no further [output statement](#) is executed for the file before it is closed or a [BACKSPACE](#), [ENDFILE](#), or [REWIND](#) statement is executed for it, the effect is as if the [output statement](#) were the corresponding advancing [output statement](#).

#### 9.3.4.3 File position prior to data transfer

- 1 The positioning of the file prior to data transfer depends on the method of access: sequential, direct, or stream.
- 2 For sequential access on input, if there is a current record, the file position is not changed. Otherwise, the file is positioned at the beginning of the next record and this record becomes the current record. Input shall not occur if there is no next record or if there is a current record and the last [data transfer statement](#) accessing the file performed output.
- 3 If the file contains an endfile record, the file shall not be positioned after the endfile record prior to data transfer. However, a [REWIND](#) or [BACKSPACE](#) statement may be used to reposition the file.
- 4 For sequential access on output, if there is a current record, the file position is not changed and the current record becomes the last record of the file. Otherwise, a new record is created as the next record of the file; this new record becomes the last and current record of the file and the file is positioned at the beginning of this record.
- 5 For direct access, the file is positioned at the beginning of the record specified by the [REC= specifier](#). This record becomes the current record.
- 6 For stream access, the file is positioned immediately before the [file storage unit](#) specified by the [POS= specifier](#);

1 if there is no **POS= specifier**, the file position is not changed.

2 7 File positioning for child **data transfer statements** is described in 9.6.4.8.

### 3 9.3.4.4 File position after data transfer

4 1 If an error condition (9.11) occurred, the position of the file is indeterminate. If no error condition occurred, but  
5 an end-of-file condition (9.11) occurred as a result of reading an endfile record, the file is positioned after the  
6 endfile record.

7 2 For unformatted stream input/output, if no error condition occurred, the file position is not changed. For  
8 unformatted stream output, if the file position exceeds the previous terminal point of the file, the terminal point  
9 is set to the file position.

#### NOTE 9.10

An unformatted stream **output statement** with a **POS= specifier** and an empty output list can have the effect of extending the terminal point of a file without actually writing any data.

10 3 For formatted stream input, if an end-of-file condition occurred, the file position is not changed.

11 4 For nonadvancing input, if no error condition or end-of-file condition occurred, but an end-of-record condition  
12 (9.11) occurred, the file is positioned after the record just read. If no error condition, end-of-file condition, or  
13 end-of-record condition occurred in a nonadvancing **input statement**, the file position is not changed. If no error  
14 condition occurred in a nonadvancing **output statement**, the file position is not changed.

15 5 In all other cases, the file is positioned after the record just read or written and that record becomes the preceding  
16 record.

17 6 For a formatted stream **output statement**, if no error condition occurred, the terminal point of the file is set to  
18 the next position after the highest-numbered position to which a datum was transferred by the statement.

#### NOTE 9.11

The highest-numbered position might not be the current one if the output involved a T, TL, TR, or X edit descriptor (10.8.1) and the statement is a nonadvancing **output statement**.

### 19 9.3.5 File storage units

20 1 A **file storage unit** is the basic unit of storage in a **stream file** or an unformatted **record file**. It is the unit of file  
21 position for stream access, the unit of record length for unformatted files, and the unit of file size for all **external**  
22 **files**.

23 2 Every value in a **stream file** or an unformatted **record file** shall occupy an integer number of **file storage units**; if  
24 the **stream** or **record** file is unformatted, this number shall be the same for all scalar values of the same type and  
25 type parameters. The number of **file storage units** required for an item of a given type and type parameters may  
26 be determined using the **IOLength= specifier** of the **INQUIRE statement** (9.10.3).

27 3 For a file **connected** for unformatted stream access, the processor shall not have alignment restrictions that prevent  
28 a value of any type from being stored at any positive integer file position.

29 4 The number of bits in a **file storage unit** is given by the constant **FILE\_STORAGE\_SIZE** (13.10.2.11) defined  
30 in the intrinsic module **ISO\_FORTRAN\_ENV**. It is recommended that the **file storage unit** be an 8-bit octet  
31 where this choice is practical.

#### NOTE 9.12

The requirement that every data value occupy an integer number of **file storage units** implies that data items inherently smaller than a **file storage unit** will require padding. This suggests that the **file storage unit** be small to avoid wasted space. Ideally, the file storage unit would be chosen such that padding is

**NOTE 9.12** (cont.)

never required. A [file storage unit](#) of one bit would always meet this goal, but would likely be impractical because of the alignment requirements.

The prohibition on alignment restrictions prohibits the processor from requiring data alignments larger than the [file storage unit](#).

The 8-bit octet is recommended as a good compromise that is small enough to accommodate the requirements of many applications, yet not so small that the data alignment requirements are likely to cause significant performance problems.

## 9.4 Internal files

1 **Internal files** provide a means of transferring and converting data from internal storage to internal storage.

2 An **internal file** is a **record file** with the following properties.

- The file is a variable of default, [ASCII](#), or [ISO 10646 character](#) that is not an [array section](#) with a [vector subscript](#).
- A record of an [internal file](#) is a scalar character variable.
- If the file is a scalar character variable, it consists of a single record whose length is the same as the length of the scalar character variable. If the file is a character array, it is treated as a sequence of character array elements. Each array element, if any, is a record of the file. The ordering of the records of the file is the same as the ordering of the array elements in the array ([6.5.3.2](#)) or the [array section](#) ([6.5.3.3](#)). Every record of the file has the same length, which is the length of an array element in the array.
- A record of the [internal file](#) becomes defined by writing the record. If the number of characters written in a record is less than the length of the record, the remaining portion of the record is filled with blanks. The number of characters to be written shall not exceed the length of the record.
- A record may be read only if the record is defined.
- A record of an [internal file](#) may become defined (or undefined) by means other than an [output statement](#). For example, the character variable may become defined by a character [assignment statement](#).
- An [internal file](#) is always positioned at the beginning of the first record prior to data transfer, except for child [data transfer statements](#) ([9.6.4.8](#)). This record becomes the current record.
- The initial value of a connection mode ([9.5.2](#)) is the value that would be implied by an initial [OPEN statement](#) without the corresponding keyword.
- Reading and writing records shall be accomplished only by sequential access formatted [data transfer statements](#).
- An [internal file](#) shall not be specified as the [unit](#) in a [CLOSE](#), [INQUIRE](#), or [OPEN](#) statement.

## 9.5 File connection

### 9.5.1 Referring to a file

1 A [unit](#), specified by an *io-unit*, provides a means for referring to a file.

```
R901      io-unit          is file-unit-number  
            or             *
```

or internal-file-variable

R902 *file-unit-number* is scalar-*int-expr*

R903 *internal-file-variable* is *char-variable*

C901 (R903) The *char-variable* shall not be an array section with a vector subscript.

C902 (R903) The *char-variable* shall be default character, *ASCII character*, or *ISO 10646 character*.

A *unit* is either an *external unit* or an *internal unit*. An *external unit* is used to refer to an *external file* and is specified by an asterisk or a *file-unit-number*. The value of *file-unit-number* shall be nonnegative, equal to one of the *named constants* *INPUT\_UNIT*, *OUTPUT\_UNIT*, or *ERROR\_UNIT* of the intrinsic module *ISO\_FORTRAN\_ENV* (13.10.2), the *unit* argument of an active *defined input/output* procedure (9.6.4.8), or a *NEWUNIT* value (9.5.6.12). An *internal unit* is used to refer to an *internal file* and is specified by an *internal-file-variable* or a *file-unit-number* whose value is equal to the *unit* argument of an active *defined input/output* procedure. The value of a *file-unit-number* shall identify a valid *unit*.

On an *image*, the *external unit* identified by a particular value of a *scalar-int-expr* is the same *external unit* in all *program units*.

#### NOTE 9.13

In the example:

```
SUBROUTINE A
  READ (6) X
  ...
SUBROUTINE B
  N = 6
  REWIND N
```

the value 6 used in both *program units* identifies the same *external unit*.

In a *READ statement*, an *io-unit* that is an asterisk identifies an *external unit* that is *preconnected* for sequential formatted input on *image 1* in the *initial team* only (9.6.4.3). This *unit* is also identified by the value of the *named constant* *INPUT\_UNIT* of the intrinsic module *ISO\_FORTRAN\_ENV* (13.10.2.13). In a *WRITE statement*, an *io-unit* that is an asterisk identifies an *external unit* that is *preconnected* for sequential formatted output. This *unit* is also identified by the value of the *named constant* *OUTPUT\_UNIT* of the intrinsic module *ISO\_FORTRAN\_ENV* (13.10.2.22).

This part of ISO/IEC 1539 identifies a processor-dependent *external unit* for the purpose of error reporting. This *unit* shall be *preconnected* for sequential formatted output. The processor may define this to be the same as the output *unit* identified by an asterisk. This *unit* is also identified by a unit number defined by the *named constant* *ERROR\_UNIT* of the intrinsic module *ISO\_FORTRAN\_ENV*.

#### NOTE 9.14

Even though *OUTPUT\_UNIT* is connected to a separate file on each *image*, it is expected that the processor could merge the sequences of records from these files into a single sequence of records that is sent to the physical device associated with this *unit*, such as the user's terminal. If *ERROR\_UNIT* is associated with the same physical device, the sequences of records from files connected to *ERROR\_UNIT* on each of the *images* could be merged into the same sequence generated from the *OUTPUT\_UNIT* files. Otherwise, it is expected that the sequence of records in the files connected to *ERROR\_UNIT* on each *image* could be merged into a single sequence of records that is sent to the physical device associated with *ERROR\_UNIT*.

### 9.5.2 Connection modes

A connection for formatted input/output has several changeable modes: these are the blank interpretation mode (10.8.6), delimiter mode (10.10.4, 10.11.4.2), sign mode (10.8.4), decimal edit mode (10.8.8), input/output rounding mode (10.7.2.3.8), pad mode (9.6.4.5.3), and scale factor (10.8.5). A connection for unformatted input/output has no changeable modes.

Values for the modes of a connection are established when the connection is initiated. If the connection is initiated by an *OPEN statement*, the values are as specified, either explicitly or implicitly, by the *OPEN statement*. If the connection is initiated other than by an *OPEN statement* (that is, if the file is an *internal file* or *preconnected file*)

1 the values established are those that would be implied by an initial **OPEN statement** without the corresponding  
2 keywords.

3 3 The scale factor cannot be explicitly specified in an **OPEN statement**; it is implicitly 0.

4 4 The modes of a connection to an **external file** may be changed by a subsequent **OPEN statement** that modifies  
5 the connection.

6 5 The modes of a connection may be temporarily changed by a corresponding keyword specifier in a **data transfer**  
7 **statement** or by an edit descriptor. Keyword specifiers take effect at the beginning of execution of the **data**  
8 **transfer statement**. Edit descriptors take effect when they are encountered in format processing. When a **data**  
9 **transfer statement** terminates, the values for the modes are reset to the values in effect immediately before the  
10 **data transfer statement** was executed.

### 11 9.5.3 Unit existence

12 1 At any given time, there is a processor-dependent set of **external units** that exist for an **image**.

13 2 All input/output statements are permitted to refer to **units** that exist. The **CLOSE**, **INQUIRE**, and **WAIT**  
14 statements are also permitted to refer to **units** that do not exist. No other input/output statement shall refer to  
15 a **unit** that does not exist.

### 16 9.5.4 Connection of a file to a unit

17 1 An **external unit** has a property of being **connected** or not connected. If **connected**, it refers to an **external file**. An  
18 **external unit** may become **connected** by **preconnection** or by the execution of an **OPEN statement**. The property  
19 of connection is symmetric; the unit is **connected** to a file if and only if the file is **connected** to the **unit**.

20 2 Every input/output statement except an **OPEN**, **CLOSE**, **INQUIRE**, or **WAIT** statement shall refer to a **unit**  
21 that is **connected** to a file and thereby make use of or affect that file.

22 3 A file may be **connected** and not exist (9.3.2).

#### NOTE 9.15

An example is a **preconnected external file** that has not yet been written.

23 4 A **unit** shall not be **connected** to more than one file at the same time. However, means are provided to change  
24 the status of an **external unit** and to connect a **unit** to a different file. It is processor dependent whether a file  
25 can be **connected** to more than one **unit** at the same time.

26 5 This part of ISO/IEC 1539 defines means of portable interoperation with C. C streams are described in 7.21.2 of  
27 ISO/IEC 9899:2011. Whether a **unit** can be **connected** to a file that is also **connected** to a C stream is processor  
28 dependent. If a **unit** is **connected** to a file that is also **connected** to a C stream, the results of performing  
29 input/output operations on such a file are processor dependent. It is processor dependent whether the files  
30 **connected** to the **units** **INPUT\_UNIT**, **OUTPUT\_UNIT**, and **ERROR\_UNIT** correspond to the predefined C  
31 text streams standard input, standard output, and standard error. If a main program or procedure defined by  
32 means of Fortran and a main program or procedure defined by means other than Fortran perform input/output  
33 operations on the same **external file**, the results are processor dependent. A main program or procedure defined by  
34 means of Fortran and a main program or procedure defined by means other than Fortran can perform input/output  
35 operations on different **external files** without interference.

36 6 If input/output operations are performed on more than one unit while they are connected to the same **external**  
37 **file**, the results are processor dependent.

38 7 After an **external unit** has been disconnected by the execution of a **CLOSE statement**, it may be **connected** again  
39 within the same program to the same file or to a different file. After an **external file** has been disconnected by  
40 the execution of a **CLOSE statement**, it may be **connected** again within the same program to the same **unit** or

to a different **unit**.

#### NOTE 9.16

The only means of referencing a file that has been disconnected is by the appearance of its name in an **OPEN** or **INQUIRE** statement. There might be no means of reconnecting an unnamed file once it is disconnected.

8 An **internal unit** is always **connected** to the **internal file** designated by the variable that identifies the **unit**.

#### NOTE 9.17

For more explanatory information on file connection properties, see C.6.4.

## 9.5.5 Preconnection

1 Preconnection means that the **unit** is **connected** to a file at the beginning of execution of the program and therefore it may be specified in input/output statements without the prior execution of an **OPEN statement**.

## 9.5.6 OPEN statement

### 9.5.6.1 General

1 An **OPEN statement** initiates or modifies the connection between an **external file** and a specified **unit**. The **OPEN statement** may be used to connect an existing file to a **unit**, create a file that is **preconnected**, create a file and connect it to a **unit**, or change certain modes of a connection between a file and a **unit**.

2 An **external unit** may be **connected** by an **OPEN statement** in the main program or any subprogram.

3 If the file to be **connected** to the **unit** does not exist but is the same as the file to which the **unit** is **preconnected**, the modes specified by an **OPEN statement** become a part of the connection.

4 If the file to be **connected** to the **unit** is not the same as the file to which the **unit** is **connected**, the effect is as if a **CLOSE statement** without a **STATUS= specifier** had been executed for the **unit** immediately prior to the execution of an **OPEN statement**.

5 If a **unit** is **connected** to a file that exists, execution of an **OPEN statement** for that **unit** is permitted. If the **FILE= specifier** is not included in such an **OPEN statement**, the file to be **connected** to the **unit** is the same as the file to which the **unit** is already **connected**.

6 If the file to be **connected** to the **unit** is the same as the file to which the **unit** is **connected**, a new connection is not established and values for any changeable modes (9.5.2) specified come into effect for the established connection; the current file position is unaffected. Before any effect on changeable modes, a wait operation is performed for any pending asynchronous data transfer operations for the specified **unit**. If the **POSITION= specifier** appears in such an **OPEN statement**, the value specified shall not disagree with the current position of the file. If the **STATUS= specifier** is included in such an **OPEN statement**, it shall be specified with the value **OLD**. Other than **ERR=**, **IOSTAT=**, and **IOMSG=**, and the changeable modes, the values of all other specifiers in such an **OPEN statement** shall not differ from those in effect for the established connection.

7 A **STATUS= specifier** with a value of **OLD** is always allowed when the file to be **connected** to the **unit** is the same as the file to which the **unit** is **connected**. In this case, if the status of the file was **SCRATCH** before execution of the **OPEN statement**, the file will still be deleted when the **unit** is closed, and the file is still considered to have a status of **SCRATCH**.

### 9.5.6.2 Syntax

R904	<i>open-stmt</i>	is	<b>OPEN</b> ( <i>connect-spec-list</i> )
R905	<i>connect-spec</i>	is	[ <b>UNIT</b> = ] <i>file-unit-number</i>
		or	<b>ACCESS</b> = <i>scalar-default-char-expr</i>

- 1                   or ACTION = *scalar-default-char-expr*  
2                   or ASYNCHRONOUS = *scalar-default-char-expr*  
3                   or BLANK = *scalar-default-char-expr*  
4                   or DECIMAL = *scalar-default-char-expr*  
5                   or DELIM = *scalar-default-char-expr*  
6                   or ENCODING = *scalar-default-char-expr*  
7                   or ERR = *label*  
8                   or FILE = *file-name-expr*  
9                   or FORM = *scalar-default-char-expr*  
10                  or IOMSG = *iomsg-variable*  
11                  or IOSTAT = *scalar-int-variable*  
12                  or NEWUNIT = *scalar-int-variable*  
13                  or PAD = *scalar-default-char-expr*  
14                  or POSITION = *scalar-default-char-expr*  
15                  or RECL = *scalar-int-expr*  
16                  or ROUND = *scalar-default-char-expr*  
17                  or SIGN = *scalar-default-char-expr*  
18                  or STATUS = *scalar-default-char-expr*
- 19    R906    *file-name-expr*                   is   *scalar-default-char-expr*
- 20    R907    *iomsg-variable*                   is   *scalar-default-char-variable*
- 21    C903    No specifier shall appear more than once in a given *connect-spec-list*.
- 22    C904    (R904) If the NEWUNIT= specifier does not appear, a *file-unit-number* shall be specified; if the optional  
23            characters UNIT= are omitted, the *file-unit-number* shall be the first item in the *connect-spec-list*.
- 24    C905    (R904) If a NEWUNIT= specifier appears, a *file-unit-number* shall not appear.
- 25    C906    (R904) The *label* used in the ERR= specifier shall be the statement label of a *branch target statement*  
26            that appears in the same *inclusive scope* as the OPEN statement.
- 27    1 If the STATUS= specifier has the value NEW or REPLACE, the FILE= specifier shall appear. If the STATUS=  
28        specifier has the value SCRATCH, the FILE= specifier shall not appear. If the STATUS= specifier has the value  
29        OLD, the FILE= specifier shall appear unless the unit is *connected* and the file *connected* to the unit exists.
- 30    2 If the NEWUNIT= specifier appears in an OPEN statement, either the FILE= specifier shall appear, or the  
31        STATUS= specifier shall appear with a value of SCRATCH. The unit identified by a NEWUNIT value shall not  
32        be *preconnected*.
- 33    3 A specifier that requires a *scalar-default-char-expr* may have a limited list of character values. These values are  
34        listed for each such specifier. Any trailing blanks are ignored. The value specified is without regard to case. Some  
35        specifiers have a default value if the specifier is omitted.
- 36    4 The IOSTAT=, ERR=, and IOMSG= specifiers are described in 9.11.

**NOTE 9.18**

An example of an OPEN statement is:

```
OPEN (10, FILE = 'employee.names', ACTION = 'READ', PAD = 'YES')
```

**NOTE 9.19**

For more explanatory information on the OPEN statement, see C.6.3.

### 9.5.6.3 ACCESS= specifier in the OPEN statement

- 1 The *scalar-default-char-expr* shall evaluate to SEQUENTIAL, DIRECT, or STREAM. The ACCESS= specifier specifies the access method for the connection of the file as being sequential, direct, or stream. If this specifier is omitted, the default value is SEQUENTIAL. For an existing file, the specified access method shall be included in the set of allowed access methods for the file. For a new file, the processor creates the file with a set of allowed access methods that includes the specified method.

### 9.5.6.4 ACTION= specifier in the OPEN statement

- 1 The *scalar-default-char-expr* shall evaluate to READ, WRITE, or READWRITE. READ specifies that the WRITE, PRINT, and ENDFILE statements shall not refer to this connection. WRITE specifies that READ statements shall not refer to this connection. READWRITE permits any input/output statements to refer to this connection. If this specifier is omitted, the default value is processor dependent. If READWRITE is included in the set of allowable actions for a file, both READ and WRITE also shall be included in the set of allowed actions for that file. For an existing file, the specified action shall be included in the set of allowed actions for the file. For a new file, the processor creates the file with a set of allowed actions that includes the specified action.

### 9.5.6.5 ASYNCHRONOUS= specifier in the OPEN statement

- 1 The *scalar-default-char-expr* shall evaluate to YES or NO. If YES is specified, asynchronous input/output on the *unit* is allowed. If NO is specified, asynchronous input/output on the *unit* is not allowed. If this specifier is omitted, the default value is NO.

### 9.5.6.6 BLANK= specifier in the OPEN statement

- 1 The *scalar-default-char-expr* shall evaluate to NULL or ZERO. The BLANK= specifier is permitted only for a connection for formatted input/output. It specifies the blank interpretation mode (10.8.6, 9.6.2.6) for input for this connection. This mode has no effect on output. It is a changeable mode (9.5.2). If this specifier is omitted in an OPEN statement that initiates a connection, the default value is NULL.

### 9.5.6.7 DECIMAL= specifier in the OPEN statement

- 1 The *scalar-default-char-expr* shall evaluate to COMMA or POINT. The DECIMAL= specifier is permitted only for a connection for formatted input/output. It specifies the decimal edit mode (10.6, 10.8.8, 9.6.2.7) for this connection. This is a changeable mode (9.5.2). If this specifier is omitted in an OPEN statement that initiates a connection, the default value is POINT.

### 9.5.6.8 DELIM= specifier in the OPEN statement

- 1 The *scalar-default-char-expr* shall evaluate to APOSTROPHE, QUOTE, or NONE. The DELIM= specifier is permitted only for a connection for formatted input/output. It specifies the delimiter mode (9.6.2.8) for list-directed (10.10.4) and namelist (10.11.4.2) output for the connection. This mode has no effect on input. It is a changeable mode (9.5.2). If this specifier is omitted in an OPEN statement that initiates a connection, the default value is NONE.

### 9.5.6.9 ENCODING= specifier in the OPEN statement

- 1 The *scalar-default-char-expr* shall evaluate to UTF-8 or DEFAULT. The ENCODING= specifier is permitted only for a connection for formatted input/output. The value UTF-8 specifies that the encoding form of the file is UTF-8 as specified in ISO/IEC 10646. Such a file is called a Unicode file, and all characters therein are of ISO 10646 character kind. The value UTF-8 shall not be specified if the processor does not support the ISO 10646 character kind. The value DEFAULT specifies that the encoding form of the file is processor dependent. If this specifier is omitted in an OPEN statement that initiates a connection, the default value is DEFAULT.

#### 9.5.6.10 FILE= specifier in the OPEN statement

- 1 The value of the FILE= specifier is the name of the file to be [connected](#) to the specified [unit](#). Any trailing blanks are ignored. The [file-name-expr](#) shall be a name that is allowed by the processor. If this specifier is omitted and the [unit](#) is not connected to a file, the STATUS= specifier shall be specified with a value of SCRATCH; in this case, the connection is made to a processor-dependent file. The interpretation of case is processor dependent.

#### 9.5.6.11 FORM= specifier in the OPEN statement

- 1 The [scalar-default-char-expr](#) shall evaluate to FORMATTED or UNFORMATTED. The FORM= specifier determines whether the file is being connected for formatted or unformatted input/output. If this specifier is omitted, the default value is UNFORMATTED if the file is being connected for direct access or stream access, and the default value is FORMATTED if the file is being connected for sequential access. For an existing file, the specified form shall be included in the set of allowed forms for the file. For a new file, the processor creates the file with a set of allowed forms that includes the specified form.

#### 9.5.6.12 NEWUNIT= specifier in the OPEN statement

- 1 The variable is defined with a processor determined NEWUNIT value if no error occurs during the execution of the OPEN statement. If an error occurs, the processor shall not change the value of the variable.
- 2 A NEWUNIT value is a negative number, and shall not be equal to  $-1$ , any of the [named constants](#) [ERROR\\_UNIT](#), [INPUT\\_UNIT](#), or [OUTPUT\\_UNIT](#) from the intrinsic module [ISO\\_FORTRAN\\_ENV](#) (13.10.2), any value used by the processor for the [unit](#) argument to a [defined input/output](#) procedure, nor any previous NEWUNIT value that identifies a file that is [connected](#).

#### 9.5.6.13 PAD= specifier in the OPEN statement

- 1 The [scalar-default-char-expr](#) shall evaluate to YES or NO. The PAD= specifier is permitted only for a connection for formatted input/output. It specifies the pad mode ([9.6.4.5.3](#), [9.6.2.10](#)) for input for this connection. This mode has no effect on output. It is a changeable mode ([9.5.2](#)). If this specifier is omitted in an OPEN statement that initiates a connection, the default value is YES.

#### 9.5.6.14 POSITION= specifier in the OPEN statement

- 1 The [scalar-default-char-expr](#) shall evaluate to ASIS, REWIND, or APPEND. The connection shall be for sequential or stream access. A new file is positioned at its initial point. REWIND positions an existing file at its initial point. APPEND positions an existing file such that the endfile record is the next record, if it has one. If an existing file does not have an endfile record, APPEND positions the file at its terminal point. ASIS leaves the position unchanged if the file exists and already is [connected](#). If the file exists but is not connected, the position resulting from ASIS is processor dependent. If this specifier is omitted, the default value is ASIS.

#### 9.5.6.15 RECL= specifier in the OPEN statement

- 1 The value of the RECL= specifier shall be positive. It specifies the length of each record in a file being connected for direct access, or specifies the maximum length of a record in a file being connected for sequential access. This specifier shall not appear when a file is being connected for stream access. This specifier shall appear when a file is being connected for direct access. If this specifier is omitted when a file is being connected for sequential access, the default value is processor dependent. If the file is being connected for formatted input/output, the length is the number of characters for all records that contain only characters of default kind. When a record contains any nondefault characters, the effect of the RECL= specifier is processor dependent. If the file is being connected for unformatted input/output, the length is measured in [file storage units](#). For an existing file, the value of the RECL= specifier shall be included in the set of allowed record lengths for the file. For a new file, the processor creates the file with a set of allowed record lengths that includes the specified value.

### 9.5.6.16 ROUND= specifier in the OPEN statement

- 1 The *scalar-default-char-expr* shall evaluate to one of UP, DOWN, ZERO, NEAREST, COMPATIBLE, or PROCESSOR\_DEFINED. The ROUND= specifier is permitted only for a connection for formatted input/output. It specifies the input/output rounding mode (10.7.2.3.8, 9.6.2.13) for this connection. This is a changeable mode (9.5.2). If this specifier is omitted in an OPEN statement that initiates a connection, the input/output rounding mode is processor dependent; it shall be one of the above modes.

#### NOTE 9.20

A processor is free to select any input/output rounding mode for the default mode. The mode might correspond to UP, DOWN, ZERO, NEAREST, or COMPATIBLE; or it might be a completely different input/output rounding mode.

### 9.5.6.17 SIGN= specifier in the OPEN statement

- 1 The *scalar-default-char-expr* shall evaluate to one of PLUS, SUPPRESS, or PROCESSOR\_DEFINED. The SIGN= specifier is permitted only for a connection for formatted input/output. It specifies the sign mode (10.8.4, 9.6.2.14) for this connection. This is a changeable mode (9.5.2). If this specifier is omitted in an OPEN statement that initiates a connection, the default value is PROCESSOR\_DEFINED.

### 9.5.6.18 STATUS= specifier in the OPEN statement

- 1 The *scalar-default-char-expr* shall evaluate to OLD, NEW, SCRATCH, REPLACE, or UNKNOWN. If OLD is specified, the file shall exist. If NEW is specified, the file shall not exist.
- 2 Successful execution of an OPEN statement with NEW specified creates the file and changes the status to OLD. If REPLACE is specified and the file does not already exist, the file is created and the status is changed to OLD. If REPLACE is specified and the file does exist, the file is deleted, a new file is created with the same name, and the status is changed to OLD. If SCRATCH is specified, the file is created and connected to the specified unit for use by the program but is deleted at the execution of a CLOSE statement referring to the same unit or at the normal termination of the program.

#### NOTE 9.21

SCRATCH shall not be specified with a named file.

- 3 If UNKNOWN is specified, the status is processor dependent. If this specifier is omitted, the default value is UNKNOWN.

## 9.5.7 CLOSE statement

### 9.5.7.1 General

- 1 The CLOSE statement is used to terminate the connection of a specified unit to an external file.
- 2 Execution of a CLOSE statement for a unit may occur in any program unit of a program and need not occur in the same program unit as the execution of an OPEN statement referring to that unit.
- 3 Execution of a CLOSE statement performs a wait operation for any pending asynchronous data transfer operations for the specified unit.
- 4 Execution of a CLOSE statement specifying a unit that does not exist, exists but is connected to a file that does not exist, or has no file connected to it, is permitted and affects no file or unit.
- 5 After a unit has been disconnected by execution of a CLOSE statement, it may be connected again within the same program, either to the same file or to a different file. After a named file has been disconnected by execution of a CLOSE statement, it may be connected again within the same program, either to the same unit or to a different unit, provided that the file still exists.

- 6 During the completion step (2.3.7) of **normal termination**, all **units** that are **connected** are closed. Each **unit** is closed with status KEEP unless the file status prior to **termination of execution** was SCRATCH, in which case the **unit** is closed with status DELETE.

**NOTE 9.22**

The effect is as though a CLOSE statement without a STATUS= specifier were executed on each **connected unit**.

#### 9.5.7.2 Syntax

- R908 *close-stmt* is CLOSE ( *close-spec-list* )
- R909 *close-spec* is [ UNIT = ] *file-unit-number*  
or IOSTAT = *scalar-int-variable*  
or IOMSG = *iomsg-variable*  
or ERR = *label*  
or STATUS = *scalar-default-char-expr*
- C907 No specifier shall appear more than once in a given *close-spec-list*.
- C908 A *file-unit-number* shall be specified in a *close-spec-list*; if the optional characters UNIT= are omitted, the *file-unit-number* shall be the first item in the *close-spec-list*.
- C909 (R909) The *label* used in the ERR= specifier shall be the statement label of a **branch target statement** that appears in the same **inclusive scope** as the CLOSE statement.
- 1 The *scalar-default-char-expr* has a limited list of character values. Any trailing blanks are ignored. The value specified is without regard to case.
- 2 The IOSTAT=, ERR=, and IOMSG= specifiers are described in 9.11.

**NOTE 9.23**

An example of a CLOSE statement is:

```
CLOSE (10, STATUS = 'KEEP')
```

#### 9.5.7.3 STATUS= specifier in the CLOSE statement

- 1 The *scalar-default-char-expr* shall evaluate to KEEP or DELETE. The STATUS= specifier determines the disposition of the file that is **connected** to the specified **unit**. KEEP shall not be specified for a file whose status prior to execution of a CLOSE statement is SCRATCH. If KEEP is specified for a file that exists, the file continues to exist after the execution of a CLOSE statement. If DELETE is specified, the file will not exist after the execution of a CLOSE statement. If this specifier is omitted, the default value is KEEP, unless the file status prior to execution of the CLOSE statement is SCRATCH, in which case the default value is DELETE.

## 9.6 Data transfer statements

### 9.6.1 Form of input and output statements

- 1 The READ statement is the data transfer input statement. The WRITE statement and the PRINT statement are the data transfer output statements.

- R910 *read-stmt* is READ ( *io-control-spec-list* ) [ *input-item-list* ]  
or READ *format* [ , *input-item-list* ]
- R911 *write-stmt* is WRITE ( *io-control-spec-list* ) [ *output-item-list* ]

1 R912 *print-stmt* is PRINT *format* [ , *output-item-list* ]

#### NOTE 9.24

Examples of data transfer statements are:

```

READ (6, *) SIZE
READ 10, A, B
WRITE (6, 10) A, S, J
PRINT 10, A, S, J
10 FORMAT (2E16.3, I5)

```

## 2 9.6.2 Control information list

### 3 9.6.2.1 Syntax

4 1 A control information list is an *io-control-spec-list*. It governs data transfer.

5 R913 *io-control-spec* is [ UNIT = ] *io-unit*  
6 or [ FMT = ] *format*  
7 or [ NML = ] *namelist-group-name*  
8 or ADVANCE = *scalar-default-char-expr*  
9 or ASYNCHRONOUS = *scalar-default-char-constant-expr*  
10 or BLANK = *scalar-default-char-expr*  
11 or DECIMAL = *scalar-default-char-expr*  
12 or DELIM = *scalar-default-char-expr*  
13 or END = *label*  
14 or EOR = *label*  
15 or ERR = *label*  
16 or ID = *id-variable*  
17 or IOMSG = *iomsg-variable*  
18 or IOSTAT = *scalar-int-variable*  
19 or PAD = *scalar-default-char-expr*  
20 or POS = *scalar-int-expr*  
21 or REC = *scalar-int-expr*  
22 or ROUND = *scalar-default-char-expr*  
23 or SIGN = *scalar-default-char-expr*  
24 or SIZE = *scalar-int-variable*

25 R914 *id-variable* is *scalar-int-variable*

26 C910 No specifier shall appear more than once in a given *io-control-spec-list*.

27 C911 An *io-unit* shall be specified in an *io-control-spec-list*; if the optional characters UNIT= are omitted, the  
28 *io-unit* shall be the first item in the *io-control-spec-list*.

29 C912 (R913) A DELIM= or SIGN= specifier shall not appear in a *read-stmt*.

30 C913 (R913) A BLANK=, PAD=, END=, EOR=, or SIZE= specifier shall not appear in a *write-stmt*.

31 C914 (R913) The *label* in the ERR=, EOR=, or END= specifier shall be the statement label of a *branch target*  
32 *statement* that appears in the same *inclusive scope* as the data transfer statement.

33 C915 (R913) A *namelist-group-name* shall be the name of a namelist group.

34 C916 (R913) A *namelist-group-name* shall not appear if a REC= specifier, *format*, *input-item-list*, or an *output-*  
35 *item-list* appears in the data transfer statement.

36 C917 (R913) If *format* appears without a preceding FMT=, it shall be the second item in the *io-control-spec-list*

- 1 and the first item shall be *io-unit*.
- 2 C918 (R913) If *namelist-group-name* appears without a preceding NML=, it shall be the second item in the  
3 *io-control-spec-list* and the first item shall be *io-unit*.
- 4 C919 (R913) If *io-unit* is not a *file-unit-number*, the *io-control-spec-list* shall not contain a REC= specifier or  
5 a POS= specifier.
- 6 C920 (R913) If the REC= specifier appears, an END= specifier shall not appear, and the *format*, if any, shall  
7 not be an asterisk.
- 8 C921 (R913) An ADVANCE= specifier may appear only in a formatted sequential or stream data transfer  
9 statement with explicit format specification (10.2) whose *io-control-spec-list* does not contain an *internal-*  
10 *file-variable* as the *io-unit*.
- 11 C922 (R913) If an EOR= specifier appears, an ADVANCE= specifier also shall appear.
- 12 C923 (R913) The *scalar-default-char-constant-expr* in an ASYNCHRONOUS= specifier shall have the value  
13 YES or NO.
- 14 C924 (R913) An ASYNCHRONOUS= specifier with a value YES shall not appear unless *io-unit* is a *file-unit-*  
15 *number*.
- 16 C925 (R913) If an ID= specifier appears, an ASYNCHRONOUS= specifier with the value YES shall also  
17 appear.
- 18 C926 (R913) If a POS= specifier appears, the *io-control-spec-list* shall not contain a REC= specifier.
- 19 C927 (R913) If a DECIMAL=, BLANK=, PAD=, SIGN=, or ROUND= specifier appears, a *format* or  
20 *namelist-group-name* shall also appear.
- 21 C928 (R913) If a DELIM= specifier appears, either *format* shall be an asterisk or *namelist-group-name* shall  
22 appear.
- 23 C929 (R914) The *scalar-int-variable* shall have a decimal exponent range no smaller than that of default integer.
- 24 2 If an EOR= specifier appears, an ADVANCE= specifier with the value NO shall also appear.
- 25 3 If the data transfer statement contains a *format* or *namelist-group-name*, the statement is a formatted in-  
26 put/output statement; otherwise, it is an unformatted input/output statement.
- 27 4 The ADVANCE=, ASYNCHRONOUS=, DECIMAL=, BLANK=, DELIM=, PAD=, SIGN=, and ROUND=   
28 specifiers have a limited list of character values. Any trailing blanks are ignored. The values specified are without  
29 regard to case.
- 30 5 The IOSTAT=, ERR=, EOR=, END=, and IOMSG= specifiers are described in 9.11.

**NOTE 9.25**

An example of a READ statement is:

```
READ (IOSTAT = IOS, UNIT = 6, FMT = '(10F8.2)') A, B
```

**9.6.2.2 Format specification in a data transfer statement**

- 32 1 The *format* specifier supplies a format specification or specifies list-directed formatting for a formatted in-  
33 put/output statement.
- 34 R915 *format* is *default-char-expr*  
35 or *label*

or \*

C930 (R915) The *label* shall be the label of a **FORMAT statement** that appears in the same **inclusive scope** as the statement containing the **FMT=** specifier.

The *default-char-expr* shall evaluate to a valid format specification (10.2.1 and 10.2.2).

If *default-char-expr* is an array, it is treated as if all of the elements of the array were specified in array element order and were concatenated.

If *format* is \*, the statement is a list-directed input/output statement.

#### NOTE 9.26

An example in which the format is a character expression is:

```
READ (6, FMT = "(" // CHAR_FMT // ")") X, Y, Z
```

where CHAR\_FMT is a default character variable.

#### 9.6.2.3 NML= specifier in a data transfer statement

The **NML=** specifier supplies the *namelist-group-name* (5.8). This name identifies a particular collection of data objects on which transfer is to be performed.

If a *namelist-group-name* appears, the statement is a namelist input/output statement.

#### 9.6.2.4 ADVANCE= specifier in a data transfer statement

The *scalar-default-char-expr* shall evaluate to YES or NO. The **ADVANCE=** specifier determines whether advancing input/output occurs for a nonchild data transfer statement. If YES is specified for a nonchild data transfer statement, advancing input/output occurs. If NO is specified, nonadvancing input/output occurs (9.3.4.2). If this specifier is omitted from a nonchild data transfer statement that allows the specifier, the default value is YES. A formatted child data transfer statement is a nonadvancing input/output statement, and any **ADVANCE=** specifier is ignored.

#### 9.6.2.5 ASYNCHRONOUS= specifier in a data transfer statement

The **ASYNCHRONOUS=** specifier determines whether this data transfer statement is synchronous or asynchronous. If YES is specified, the statement and the input/output operation are asynchronous. If NO is specified or if the specifier is omitted, the statement and the input/output operation are synchronous.

Asynchronous input/output is permitted only for **external files** opened with an **ASYNCHRONOUS=** specifier with the value YES in the **OPEN statement**.

#### NOTE 9.27

Both synchronous and asynchronous input/output are allowed for files opened with an **ASYNCHRONOUS=** specifier of YES. For other files, only synchronous input/output is allowed; this includes files opened with an **ASYNCHRONOUS=** specifier of NO, files opened without an **ASYNCHRONOUS=** specifier, **preconnected** files accessed without an **OPEN statement**, and **internal files**.

The **ASYNCHRONOUS=** specifier value in a data transfer statement is a **constant expression** because it effects compiler optimizations and, therefore, needs to be known at compile time.

The processor may perform an asynchronous data transfer operation asynchronously, but it is not required to do so. For each **external file**, records and **file storage units** read or written by asynchronous data transfer statements are read, written, and processed in the same order as they would have been if the data transfer statements were synchronous. The documentation of the Fortran processor should describe when input/output will be performed asynchronously.

1 4 If a variable is used in an asynchronous data transfer statement as

- 2 • an item in an input/output list,
- 3 • a group object in a namelist, or
- 4 • a SIZE= specifier,

5 the [base object](#) of the [data-ref](#) is implicitly given the [ASYNCHRONOUS](#) attribute in the [scoping unit](#) of the data  
6 transfer statement. This attribute may be confirmed by explicit declaration.

7 5 When an asynchronous input/output statement is executed, the set of [storage units](#) specified by the item list or  
8 NML= specifier, plus the [storage units](#) specified by the SIZE= specifier, is defined to be the pending input/output  
9 storage sequence for the data transfer operation.

#### NOTE 9.28

A pending input/output storage sequence is not necessarily a [contiguous](#) set of [storage units](#).

10 6 A pending input/output storage sequence affector is a variable of which any part is associated with a [storage unit](#)  
11 in a pending input/output storage sequence.

#### 12 9.6.2.6 BLANK= specifier in a data transfer statement

13 1 The [scalar-default-char-expr](#) shall evaluate to NULL or ZERO. The BLANK= specifier temporarily changes  
14 (9.5.2) the blank interpretation mode (10.8.6, 9.5.6.6) for the connection. If the specifier is omitted, the mode is  
15 not changed.

#### 16 9.6.2.7 DECIMAL= specifier in a data transfer statement

17 1 The [scalar-default-char-expr](#) shall evaluate to COMMA or POINT. The DECIMAL= specifier temporarily changes  
18 (9.5.2) the decimal edit mode (10.6, 10.8.8, 9.5.6.7) for the connection. If the specifier is omitted, the mode is  
19 not changed.

#### 20 9.6.2.8 DELIM= specifier in a data transfer statement

21 1 The [scalar-default-char-expr](#) shall evaluate to APOSTROPHE, QUOTE, or NONE. The DELIM= specifier tem-  
22 porarily changes (9.5.2) the delimiter mode (10.10.4, 10.11.4.2, 9.5.6.8) for the connection. If the specifier is  
23 omitted, the mode is not changed.

#### 24 9.6.2.9 ID= specifier in a data transfer statement

25 1 Successful execution of an asynchronous data transfer statement containing an ID= specifier causes the variable  
26 specified in the ID= specifier to become defined with a processor determined value. If this value is zero, the  
27 data transfer operation has been completed. A nonzero value is referred to as the identifier of the data transfer  
28 operation. This identifier is different from the identifier of any other pending data transfer operation for this [unit](#).  
29 It can be used in a subsequent [WAIT](#) or [INQUIRE](#) statement to identify the particular data transfer operation.

30 2 If an error occurs during the execution of a data transfer statement containing an ID= specifier, the variable  
31 specified in the ID= specifier becomes undefined.

32 3 A child data transfer statement shall not specify the ID= specifier.

#### 33 9.6.2.10 PAD= specifier in a data transfer statement

34 1 The [scalar-default-char-expr](#) shall evaluate to YES or NO. The PAD= specifier temporarily changes (9.5.2) the  
35 pad mode (9.6.4.5.3, 9.5.6.13) for the connection. If the specifier is omitted, the mode is not changed.

### 9.6.2.11 POS= specifier in a data transfer statement

- 1 The POS= specifier specifies the file position in [file storage units](#). This specifier may appear in a data transfer statement only if the statement specifies a [unit connected](#) for stream access. A child data transfer statement shall not specify this specifier.
- 2 A processor may prohibit the use of POS= with particular files that do not have the properties necessary to support random positioning. A processor may also prohibit positioning a particular file to any position prior to its current file position if the file does not have the properties necessary to support such positioning.

#### NOTE 9.29

A [unit](#) that is [connected](#) to a device or data stream might not be positionable.

- 3 If the file is [connected](#) for formatted stream access, the file position specified by POS= shall be equal to either 1 (the beginning of the file) or a value previously returned by a [POS= specifier](#) in an [INQUIRE statement](#) for the file.

### 9.6.2.12 REC= specifier in a data transfer statement

- 1 The REC= specifier specifies the number of the record that is to be read or written. This specifier may appear only in a data transfer statement that specifies a [unit connected](#) for direct access; it shall not appear in a child data transfer statement. If the [io-control-spec-list](#) contains a REC= specifier, the statement is a direct access data transfer statement. A child data transfer statement is a direct access data transfer statement if the parent is a direct access data transfer statement. Any other data transfer statement is a sequential access data transfer statement or a stream access data transfer statement, depending on whether the file connection is for sequential access or stream access.

### 9.6.2.13 ROUND= specifier in a data transfer statement

- 1 The [scalar-default-char-expr](#) shall evaluate to UP, DOWN, ZERO, NEAREST, COMPATIBLE or PROCESSOR\_DEFINED. The ROUND= specifier temporarily changes (9.5.2) the input/output rounding mode (10.7.2.3.8, 9.5.6.16) for the connection. If the specifier is omitted, the mode is not changed.

### 9.6.2.14 SIGN= specifier in a data transfer statement

- 1 The [scalar-default-char-expr](#) shall evaluate to PLUS, SUPPRESS, or PROCESSOR\_DEFINED. The SIGN= specifier temporarily changes (9.5.2) the sign mode (10.8.4, 9.5.6.17) for the connection. If the specifier is omitted, the mode is not changed.

### 9.6.2.15 SIZE= specifier in a data transfer statement

- 1 The SIZE= specifier in an [input statement](#) causes the variable specified to become [defined](#) with the count of the characters transferred from the file by data edit descriptors during the input operation. Blanks inserted as padding are not counted.
- 2 For a synchronous [input statement](#), this definition occurs when execution of the statement completes. For an asynchronous [input statement](#), this definition occurs when the corresponding wait operation is performed.

## 9.6.3 Data transfer input/output list

- 1 An input/output list specifies the entities whose values are transferred by a data transfer statement.

R916 *input-item* is [variable](#)  
or [io-implied-do](#)

R917 *output-item* is [expr](#)  
or [io-implied-do](#)

- 1 R918 *io-implied-do* is ( *io-implied-do-object-list* , *io-implied-do-control* )
- 2 R919 *io-implied-do-object* is *input-item*
- 3 or *output-item*
- 4 R920 *io-implied-do-control* is *do-variable* = *scalar-int-expr* , ■
- 5 ■ *scalar-int-expr* [ , *scalar-int-expr* ]
- 6 C931 (R916) A variable that is an *input-item* shall not be a whole *assumed-size* array.
- 7 C932 (R919) In an *input-item-list*, an *io-implied-do-object* shall be an *input-item*. In an *output-item-list*, an
- 8 *io-implied-do-object* shall be an *output-item*.
- 9 C933 (R917) An expression that is an *output-item* shall not have a value that is a procedure pointer.
- 10 2 An *input-item* shall not appear as, nor be associated with, the *do-variable* of any *io-implied-do* that contains the
- 11 *input-item*.

**NOTE 9.30**

A constant, an expression involving operators or function references that does not have a pointer result, or an expression enclosed in parentheses shall not appear as an input list item.

- 12 3 If an input item is a pointer, it shall be associated with a *definable target* and data are transferred from the file to
- 13 the associated *target*. If an output item is a pointer, it shall be associated with a *target* and data are transferred
- 14 from the *target* to the file.

**NOTE 9.31**

Data transfers always involve the movement of values between a file and internal storage. A pointer as such cannot be read or written. Therefore, a pointer shall not appear as an item in an input/output list unless it is associated with a *target* that can receive a value (input) or can deliver a value (output).

- 15 4 If an input item or an output item is *allocatable*, it shall be allocated.
- 16 5 A list item shall not be *polymorphic* unless it is processed by a *defined input/output* procedure (9.6.4.8).
- 17 6 The *do-variable* of an *io-implied-do* that is in another *io-implied-do* shall not appear as, nor be associated with,
- 18 the *do-variable* of the containing *io-implied-do*.
- 19 7 The following rules describing whether to expand an input/output list item are re-applied to each expanded list
- 20 item until none of the rules apply.
- 21 • If an array appears as an input/output list item, it is treated as if the elements, if any, were specified in
- 22 array element order (6.5.3.2). However, no element of that array may affect the value of any expression in
- 23 the *input-item*, nor may any element appear more than once in an *input-item*.

**NOTE 9.32**

For example:

```

INTEGER A (100), J (100)
...
READ *, A (A)                ! Not allowed
READ *, A (LBOUND (A, 1) : UBOUND (A, 1)) ! Allowed
READ *, A (J)                ! Allowed if no two elements
                                !   of J have the same value
A(1) = 1; A(10) = 10
READ *, A (A (1) : A (10))    ! Not allowed

```

- If a list item of derived type in an unformatted input/output statement is not processed by a [defined input/output](#) procedure (9.6.4.8), and if any subobject of that list item would be processed by a [defined input/output](#) procedure, the list item is treated as if all of the components of the object were specified in the list in [component order](#) (4.5.4.7); those components shall be accessible in the [scoping unit](#) containing the data transfer statement and shall not be pointers or [allocatable](#).
- An effective item of derived type in an unformatted input/output statement is treated as a single value in a processor-dependent form unless the list item or a subobject thereof is processed by a [defined input/output](#) procedure (9.6.4.8).

**NOTE 9.33**

The appearance of a derived-type object as an input/output list item in an unformatted input/output statement is not equivalent to the list of its components.

Unformatted input/output involving derived-type list items forms the single exception to the rule that the appearance of an aggregate list item (such as an array) is equivalent to the appearance of its expanded list of component parts. This exception permits the processor greater latitude in improving efficiency or in matching the processor-dependent sequence of values for a derived-type object to similar sequences for aggregate objects used by means other than Fortran. However, formatted input/output of all list items and unformatted input/output of list items other than those of derived types adhere to the above rule.

- If a list item of derived type in a formatted input/output statement is not processed by a [defined input/output](#) procedure, that list item is treated as if all of the components of the list item were specified in the list in [component order](#); those components shall be accessible in the [scoping unit](#) containing the input/output statement and shall not be pointers or [allocatable](#).
- If a derived-type list item is not processed by a [defined input/output](#) procedure and is not treated as a list of its individual components, all the subcomponents of that list item shall be accessible in the [scoping unit](#) containing the data transfer statement and shall not be pointers or [allocatable](#).
- For an *io-implied-do*, the loop initialization and execution are the same as for a [DO construct](#) (8.1.7.4).

**NOTE 9.34**

An example of an output list with an implied DO is:

```
WRITE (LP, FMT = '(10F8.2)') (LOG (A (I)), I = 1, N + 9, K), G
```

- 8 The scalar objects resulting when a data transfer statement's list items are expanded according to the rules in this subclause for handling array and derived-type list items are called [effective items](#). Zero-sized arrays and *io-implied-dos* with an iteration count of zero do not contribute to the list of [effective items](#). A scalar character item of zero length is an [effective item](#).

**NOTE 9.35**

In a formatted input/output statement, edit descriptors are associated with [effective items](#), which are always scalar. The rules in 9.6.3 determine the set of [effective items](#) corresponding to each actual list item in the statement. These rules might have to be applied repetitively until all of the [effective items](#) are scalar items.

- 9 An input/output list shall not contain an effective item of nondefault character kind if the data transfer statement specifies an [internal file](#) of default character kind. An input/output list shall not contain an effective item that is nondefault character except for [ISO 10646](#) or [ASCII character](#) if the data transfer statement specifies an [internal file](#) of [ISO 10646 character](#) kind. An input/output list shall not contain an effective item of type character of any kind other than [ASCII](#) if the data transfer statement specifies an [ASCII character internal file](#).

## 9.6.4 Execution of a data transfer input/output statement

### 9.6.4.1 Data transfer sequence of operations

- 1 Execution of a WRITE or PRINT statement for a file that does not exist creates the file unless an error condition occurs.
- 2 The effect of executing a synchronous data transfer statement shall be as if the following operations were performed in the order specified.
  - (1) Determine the direction of data transfer.
  - (2) Identify the [unit](#).
  - (3) Perform a wait operation for all pending input/output operations for the [unit](#). If an error, end-of-file, or end-of-record condition occurs during any of the wait operations, steps 4 through 8 are skipped.
  - (4) Establish the format if one is specified.
  - (5) If the statement is not a child data transfer statement ([9.6.4.8](#)),
    - (a) position the file prior to data transfer ([9.3.4.3](#)), and
    - (b) for formatted data transfer, set the left tab limit ([10.8.1.2](#)).
  - (6) Transfer data between the file and the entities specified by the input/output list (if any) or namelist, possibly mediated by [defined input/output](#) procedures ([9.6.4.8](#)).
  - (7) Determine whether an error, end-of-file, or end-of-record condition has occurred.
  - (8) Position the file after data transfer ([9.3.4.4](#)) unless the statement is a child data transfer statement ([9.6.4.8](#)).
  - (9) Cause any variable specified in a SIZE= specifier to become defined.
  - (10) If an error, end-of-file, or end-of-record condition occurred, processing continues as specified in [9.11](#); otherwise, any variable specified in an IOSTAT= specifier is assigned the value zero.
- 3 The effect of executing an asynchronous data transfer statement shall be as if the following operations were performed in the order specified.
  - (1) Determine the direction of data transfer.
  - (2) Identify the [unit](#).
  - (3) Optionally, perform wait operations for one or more pending input/output operations for the [unit](#). If an error, end-of-file, or end-of-record condition occurs during any of the wait operations, steps 4 through 9 are skipped.
  - (4) Establish the format if one is specified.
  - (5) Position the file prior to data transfer ([9.3.4.3](#)) and, for formatted data transfer, set the left tab limit ([10.8.1.2](#)).
  - (6) Establish the set of [storage units](#) identified by the input/output list. For an [input statement](#), this might require some or all of the data in the file to be read if an input variable is used as a *scalar-int-expr* in an *io-implied-do-control* in the input/output list, as a *subscript*, *substring-range*, *stride*, or is otherwise referenced.
  - (7) Initiate an asynchronous data transfer between the file and the entities specified by the input/output list (if any) or namelist. The asynchronous data transfer may complete (and an error, end-of-file, or end-of-record condition may occur) during the execution of this data transfer statement or during a later wait operation.
  - (8) Determine whether an error, end-of-file, or end-of-record condition has occurred. The conditions may occur during the execution of this data transfer statement or during the corresponding wait operation, but not both.
  - (9) Position the file as if the data transfer had finished ([9.3.4.4](#)).
  - (10) Cause any variable specified in a SIZE= specifier to become undefined.
  - (11) If an error, end-of-file, or end-of-record condition occurred, processing continues as specified in [9.11](#); otherwise, any variable specified in an IOSTAT= specifier is assigned the value zero.

- 1 4 For an asynchronous data transfer statement, the data transfers may occur during execution of the statement,  
2 during execution of the corresponding wait operation, or anywhere between. The data transfer operation is  
3 considered to be pending until a corresponding wait operation is performed.
- 4 5 For asynchronous output, a pending input/output storage sequence affector (9.6.2.5) shall not be redefined,  
5 become undefined, or have its [pointer association](#) status changed.
- 6 6 For asynchronous input, a pending input/output storage sequence affector shall not be referenced, become defined,  
7 become undefined, become associated with a [dummy argument](#) that has the [VALUE attribute](#), or have its [pointer  
association](#) status changed.
- 9 7 Error, end-of-file, and end-of-record conditions in an asynchronous data transfer operation may occur during  
10 execution of either the data transfer statement or the corresponding wait operation. If an ID= specifier does not  
11 appear in the initiating data transfer statement, the conditions may occur during the execution of any subsequent  
12 data transfer or wait operation for the same [unit](#). When a condition occurs for a previously executed asynchronous  
13 data transfer statement, a wait operation is performed for all pending data transfer operations on that [unit](#). When  
14 a condition occurs during a subsequent statement, any actions specified by [IOSTAT=](#), [IOMSG=](#), [ERR=](#), [END=](#),  
15 and [EOR=](#) specifiers for that statement are taken.
- 16 8 If execution of the program is terminated during execution of an [output statement](#), the contents of the file become  
17 undefined.

**NOTE 9.36**

Because end-of-file and error conditions for asynchronous data transfer statements without an ID= specifier can be reported by the processor during the execution of a subsequent data transfer statement, it might be impossible for the user to determine which data transfer statement caused the condition. Reliably detecting which [input statement](#) caused an end-of-file condition requires that all asynchronous [input statements](#) for the [unit](#) include an ID= specifier.

#### 18 9.6.4.2 Direction of data transfer

- 19 1 Execution of a READ statement causes values to be transferred from a file to the entities specified by the input  
20 list, if any, or specified within the file itself for namelist input. Execution of a WRITE or PRINT statement  
21 causes values to be transferred to a file from the entities specified by the output list and format specification, if  
22 any, or by the *namelist-group-name* for namelist output.

#### 23 9.6.4.3 Identifying a unit

- 24 1 A data transfer statement that contains an input/output control list includes a UNIT= specifier that identifies  
25 an [external](#) or [internal](#) unit. A READ statement that does not contain an input/output control list specifies a  
26 particular processor-dependent [unit](#), which is the same as the [unit](#) identified by \* in a [READ statement](#) that  
27 contains an input/output control list (9.5.1) and is the same as the [unit](#) identified by the value of the [named  
constant](#) INPUT\_UNIT of the intrinsic module ISO\_FORTRAN\_ENV (13.10.2.13). The [PRINT statement](#)  
28 specifies some other processor-dependent [unit](#), which is the same as the [unit](#) identified by \* in a [WRITE statement](#)  
29 and is the same as the [unit](#) identified by the value of the [named constant](#) OUTPUT\_UNIT of the intrinsic module  
30 ISO\_FORTRAN\_ENV (13.10.2.22). Thus, each data transfer statement identifies an [external](#) or [internal](#) unit.  
31
- 32 2 The [unit](#) identified by an unformatted data transfer statement shall be an [external unit](#).
- 33 3 The [unit](#) identified by a data transfer statement shall be [connected](#) to a file when execution of the statement  
34 begins.

**NOTE 9.37**

The [unit](#) could be [preconnected](#).

#### 9.6.4.4 Establishing a format

- 1 If the input/output control list contains \* as a format, list-directed formatting is established. If *namelist-group-name* appears, namelist formatting is established. If no *format* or *namelist-group-name* is specified, unformatted data transfer is established. Otherwise, the format specified by *format* is established.
- 2 For output to an *internal file*, a format specification that is in the file or is associated with the file shall not be specified.
- 3 An input list item, or an entity associated with it, shall not contain any portion of an established format specification.

#### 9.6.4.5 Data transfer

##### 9.6.4.5.1 General

- 1 Data are transferred between the file and the entities specified by the input/output list or namelist. The list items are processed in the order of the input/output list for all data transfer statements except namelist data transfer statements. The list items for a namelist *input statement* are processed in the order of the entities specified within the input records. The list items for a namelist *output statement* are processed in the order in which the variables are specified in the *namelist-group-object-list*. *Effective items* are derived from the input/output list items as described in 9.6.3.
- 2 All values needed to determine which entities are specified by an input/output list item are determined at the beginning of the processing of that item.
- 3 All values are transmitted to or from the entities specified by a list item prior to the processing of any succeeding list item for all data transfer statements.

#### NOTE 9.38

In the example

```
READ (N) N, X (N)
```

the old value of N identifies the *unit*, but the new value of N is the subscript of X.

- 4 All values following the *name=* part of the namelist entity (10.11) within the input records are transmitted to the matching entity specified in the *namelist-group-object-list* prior to processing any succeeding entity within the input record for namelist *input statements*. If an entity is specified more than once within the input record during a namelist *input statement*, the last occurrence of the entity specifies the value or values to be used for that entity.
- 5 If the input/output item is a pointer, data are transferred between the file and the associated *target*.
- 6 If an *internal file* has been specified, an input/output list item shall not be in the file or associated with the file.
- 7 During the execution of an *output statement* that specifies an *internal file*, no part of that *internal file* shall be referenced, defined, or become undefined as the result of evaluating any output list item.
- 8 During the execution of an *input statement* that specifies an *internal file*, no part of that *internal file* shall be defined or become undefined as the result of transferring a value to any input list item.
- 9 A DO variable becomes defined and its iteration count established at the beginning of processing of the *io-implied-do-object-list* an *io-implied-do*.
- 10 On output, every entity whose value is to be transferred shall be defined.

#### 9.6.4.5.2 Unformatted data transfer

- 1 If the file is not [connected](#) for unformatted input/output, unformatted data transfer is prohibited.
- 2 During unformatted data transfer, data are transferred without editing between the file and the entities specified by the input/output list. If the file is [connected](#) for sequential or direct access, exactly one record is read or written.
- 3 A value in the file is stored in a [contiguous](#) sequence of [file storage units](#), beginning with the [file storage unit](#) immediately following the current file position.
- 4 After each value is transferred, the current file position is moved to a point immediately after the last [file storage unit](#) of the value.
- 5 On input from a file [connected](#) for sequential or direct access, the number of [file storage units](#) required by the input list shall be less than or equal to the number of [file storage units](#) in the record.
- 6 On input, if the [file storage units](#) transferred do not contain a value with the same type and type parameters as the input list entity, then the resulting value of the entity is processor dependent except in the following cases.
  - A complex entity may correspond to two real values with the same [kind type parameter](#) as the complex entity.
  - A default character list entity of length *n* may correspond to *n* default characters stored in the file, regardless of the length parameters of the entities that were written to these [storage units](#) of the file. If the file is [connected](#) for stream input, the characters may have been written by formatted stream output.
- 7 On output to a file [connected](#) for unformatted direct access, the output list shall not specify more values than can fit into the record. If the file is [connected](#) for direct access and the values specified by the output list do not fill the record, the remainder of the record is undefined.
- 8 If the file is [connected](#) for unformatted sequential access, the record is created with a length sufficient to hold the values from the output list. This length shall be one of the set of allowed record lengths for the file and shall not exceed the value specified in the [RECL= specifier](#), if any, of the [OPEN statement](#) that established the connection.

#### 9.6.4.5.3 Formatted data transfer

- 1 If the file is not [connected](#) for formatted input/output, formatted data transfer is prohibited.
- 2 During formatted data transfer, data are transferred with editing between the file and the entities specified by the input/output list or by the *namelist-group-name*. Format control is initiated and editing is performed as described in [Clause 10](#).
- 3 The current record and possibly additional records are read or written.
- 4 During advancing input when the pad mode has the value NO, the input list and format specification shall not require more characters from the record than the record contains.
- 5 During advancing input when the pad mode has the value YES, blank characters are supplied by the processor if the input list and format specification require more characters from the record than the record contains.
- 6 During nonadvancing input when the pad mode has the value NO, an end-of-record condition ([9.11](#)) occurs if the input list and format specification require more characters from the record than the record contains, and the record is complete ([9.3.3.4](#)). If the record is incomplete, an end-of-file condition occurs instead of an end-of-record condition.
- 7 During nonadvancing input when the pad mode has the value YES, blank characters are supplied by the processor if an [effective item](#) and its corresponding data edit descriptors require more characters from the record than the record contains. If the record is incomplete, an end-of-file condition occurs; otherwise, an end-of-record condition

1 occurs.

2 8 If the file is [connected](#) for direct access, the record number is increased by one as each succeeding record is read  
3 or written.

4 9 On output, if the file is [connected](#) for direct access or is an internal file and the characters specified by the output  
5 list and format do not fill a record, blank characters are added to fill the record.

6 10 On output, the output list and format specification shall not specify more characters for a record than have been  
7 specified by a [RECL= specifier](#) in the [OPEN statement](#) or the record length of an [internal file](#).

#### 8 **9.6.4.6 List-directed formatting**

9 1 If list-directed formatting has been established, editing is performed as described in [10.10](#).

#### 10 **9.6.4.7 Namelist formatting**

11 1 If namelist formatting has been established, editing is performed as described in [10.11](#).

12 2 Every [allocatable namelist-group-object](#) in the namelist group shall be allocated and every [namelist-group-object](#)  
13 that is a pointer shall be associated with a [target](#). If a [namelist-group-object](#) is [polymorphic](#) or has an [ultimate](#)  
14 [component](#) that is [allocatable](#) or a pointer, that object shall be processed by a [defined input/output](#) procedure  
15 ([9.6.4.8](#)).

#### 16 **9.6.4.8 Defined input/output**

##### 17 **9.6.4.8.1 General**

18 1 [Defined input/output](#) allows a program to override the default handling of derived-type objects and values in  
19 data transfer statements described in [9.6.3](#).

20 2 A [defined input/output](#) procedure is a procedure accessible by a [defined-io-generic-spec](#) ([12.4.3.2](#)). A particular  
21 [defined input/output](#) procedure is selected as described in [9.6.4.8.4](#).

##### 22 **9.6.4.8.2 Executing defined input/output data transfers**

23 1 If a [defined input/output](#) procedure is selected for an effective item as specified in [9.6.4.8.4](#), the processor shall  
24 call the selected [defined input/output](#) procedure for that item. The [defined input/output](#) procedure controls the  
25 actual data transfer operations for the derived-type list item.

26 2 A data transfer statement that includes a derived-type list item and that causes a [defined input/output](#) procedure  
27 to be invoked is called a parent data transfer statement. A data transfer statement that is executed while a parent  
28 data transfer statement is being processed and that specifies the [unit](#) passed into a [defined input/output](#) procedure  
29 is called a child data transfer statement.

#### **NOTE 9.39**

A [defined input/output](#) procedure will usually contain child data transfer statements that read values from or write values to the current record or at the current file position. The effect of executing the [defined input/output](#) procedure is similar to that of substituting the list items from any child data transfer statements into the parent data transfer statement's list items, along with similar substitutions in the format specification.

#### **NOTE 9.40**

A particular execution of a READ, WRITE or PRINT statement can be both a parent and a child data transfer statement. A [defined input/output](#) procedure can indirectly call itself or another [defined input/output](#) procedure by executing a child data transfer statement containing a list item of derived type,

**NOTE 9.40 (cont.)**

where a matching [interface](#) is accessible for that derived type. If a [defined input/output](#) procedure calls itself indirectly in this manner, it cannot be declared `NON_RECURSIVE`.

1 3 A child data transfer statement is processed differently from a nonchild data transfer statement in the following  
2 ways.

- 3 • Executing a child data transfer statement does not position the file prior to data transfer.
- 4 • An unformatted child data transfer statement does not position the file after data transfer is complete.
- 5 • Any [ADVANCE= specifier](#) in a child input/output statement is ignored.

### 6 9.6.4.8.3 Defined input/output procedures

7 1 For a particular derived type and a particular set of [kind type parameter](#) values, there are four possible sets of  
8 [characteristics](#) for [defined input/output](#) procedures; one each for formatted input, formatted output, unformatted  
9 input, and unformatted output. The user need not supply all four procedures. The procedures are specified  
10 to be used for derived-type input/output by [interface blocks](#) (12.4.3.2) or by generic [bindings](#) (4.5.5), with a  
11 [defined-io-generic-spec](#) (R1209). The [defined-io-generic-specs](#) for these procedures are [READ \(FORMATTED\)](#),  
12 [READ \(UNFORMATTED\)](#), [WRITE \(FORMATTED\)](#), and [WRITE \(UNFORMATTED\)](#), for formatted input,  
13 unformatted input, formatted output, and unformatted output respectively.

14 2 In the four [interfaces](#), which specify the [characteristics](#) of [defined input/output](#) procedures, the following syntax  
15 term is used:

16 R921 *dtv-type-spec*                                **is** TYPE( *derived-type-spec* )  
17    **or** CLASS( *derived-type-spec* )

18 C934 (R921) If *derived-type-spec* specifies an [extensible type](#), the **CLASS** keyword shall be used; otherwise, the  
19 **TYPE** keyword shall be used.

20 C935 (R921) All length type parameters of *derived-type-spec* shall be assumed.

21 3 If the [defined-io-generic-spec](#) is [READ \(FORMATTED\)](#), the [characteristics](#) shall be the same as those specified  
22 by the following [interface](#):

```

23 4      SUBROUTINE my_read_routine_formatted                                &
24                                     (dtv,                                &
25                                     unit,                                &
26                                     iotype, v_list,                      &
27                                     iostat, iomsg)
28
29      ! the derived-type variable
30      dtv-type-spec, INTENT(INOUT) :: dtv
31      INTEGER, INTENT(IN) :: unit ! unit number
32      ! the edit descriptor string
33      CHARACTER (LEN=*), INTENT(IN) :: iotype
34      INTEGER, INTENT(IN) :: v_list(:)
35      INTEGER, INTENT(OUT) :: iostat
36      CHARACTER (LEN=*), INTENT(INOUT) :: iomsg
37      END

```

37 5 If the [defined-io-generic-spec](#) is [READ \(UNFORMATTED\)](#), the [characteristics](#) shall be the same as those specified  
38 by the following [interface](#):

```

39 6      SUBROUTINE my_read_routine_unformatted                            &
40                                     (dtv,                                &

```

```

1          unit,                                &
2          iostat, iomsg)
3      ! the derived-type variable
4      dtv-type-spec, INTENT(INOUT) :: dtv
5      INTEGER, INTENT(IN) :: unit
6      INTEGER, INTENT(OUT) :: iostat
7      CHARACTER (LEN=*), INTENT(INOUT) :: iomsg
8      END

```

9 7 If the *defined-io-generic-spec* is **WRITE (FORMATTED)**, the *characteristics* shall be the same as those specified  
10 by the following *interface*:

```

11 8      SUBROUTINE my_write_routine_formatted                                &
12          (dtv,                                                                &
13          unit,                                                                &
14          iotype, v_list,                                                    &
15          iostat, iomsg)
16      ! the derived-type value/variable
17      dtv-type-spec, INTENT(IN) :: dtv
18      INTEGER, INTENT(IN) :: unit
19      ! the edit descriptor string
20      CHARACTER (LEN=*), INTENT(IN) :: iotype
21      INTEGER, INTENT(IN) :: v_list(:)
22      INTEGER, INTENT(OUT) :: iostat
23      CHARACTER (LEN=*), INTENT(INOUT) :: iomsg
24      END

```

25 9 If the *defined-io-generic-spec* is **WRITE (UNFORMATTED)**, the *characteristics* shall be the same as those  
26 specified by the following *interface*:

```

27 10     SUBROUTINE my_write_routine_unformatted                                &
28         (dtv,                                                                &
29         unit,                                                                &
30         iostat, iomsg)
31     ! the derived-type value/variable
32     dtv-type-spec, INTENT(IN) :: dtv
33     INTEGER, INTENT(IN) :: unit
34     INTEGER, INTENT(OUT) :: iostat
35     CHARACTER (LEN=*), INTENT(INOUT) :: iomsg
36     END

```

37 11 The actual specific procedure names (the *my\_...\_routine\_...* procedure names above) are not significant. In  
38 the discussion here and elsewhere, the *dummy arguments* in these *interfaces* are referred to by the names given  
39 above; the names are, however, arbitrary.

40 12 When a *defined input/output* procedure is invoked, the processor shall pass a *unit* argument that has a value as  
41 follows.

- 42 • If the parent data transfer statement uses a *file-unit-number*, the value of the *unit* argument shall be that  
43 of the *file-unit-number*.
- 44 • If the parent data transfer statement is a **WRITE statement** with an asterisk *unit* or a **PRINT statement**,  
45 the *unit* argument shall have the same value as the *named constant OUTPUT\_UNIT* of the intrinsic  
46 module **ISO\_FORTRAN\_ENV(13.10.2)**.

- If the parent data transfer statement is a **READ statement** with an asterisk **unit** or a **READ statement** without an *io-control-spec-list*, the **unit** argument shall have the same value as the **INPUT\_UNIT** named constant of the intrinsic module **ISO\_FORTRAN\_ENV**(13.10.2).
- Otherwise the parent data transfer statement must access an **internal file**, in which case the **unit** argument shall have a processor-dependent negative value.

**NOTE 9.41**

The **unit** argument passed to a **defined input/output** procedure will be negative when the parent data transfer statement specified an **internal unit**, or specified an **external unit** that is a **NEWUNIT** value. When an **internal unit** is used with the **INQUIRE statement**, an error condition will occur, and any variable specified in an **IOSTAT= specifier** will be assigned the value **IOSTAT\_INQUIRE\_INTERNAL\_UNIT** from the intrinsic module **ISO\_FORTRAN\_ENV**(13.10.2).

- 13 For formatted data transfer, the processor shall pass an **iotype** argument that has the value
- “LISTDIRECTED” if the parent data transfer statement specified list directed formatting,
  - “NAMELIST” if the parent data transfer statement specified namelist formatting, or
  - “DT” concatenated with the *char-literal-constant*, if any, of the DT edit descriptor in the format specification of the parent data transfer statement.
- 14 If the parent data transfer statement is an **input statement**, the **dtv dummy argument** is argument associated with the **effective item** that caused the **defined input** procedure to be invoked, as if the **effective item** were an **actual argument** in this **procedure reference** (2.4.5).
- 15 If the parent data transfer statement is an **output statement**, the processor shall provide the value of the **effective item** in the **dtv dummy argument**.
- 16 If the *v-list* of the edit descriptor appears in the parent data transfer statement, the processor shall provide the values from it in the **v\_list dummy argument**, with the same number of elements in the same order as *v-list*. If there is no *v-list* in the edit descriptor or if the data transfer statement specifies list-directed or namelist formatting, the processor shall provide **v\_list** as a zero-sized array.

**NOTE 9.42**

The user’s procedure might choose to interpret an element of the **v\_list** argument as a field width, but this is not required. If it does, it would be appropriate to fill an output field with “\*”s if the width is too small.

- 17 The **iostat** argument is used to report whether an error, end-of-record, or end-of-file condition (9.11) occurs. If an error condition occurs, the **defined input/output** procedure shall assign a positive value to the **iostat** argument. Otherwise, if an end-of-file condition occurs, the **defined input** procedure shall assign the value of the **named constant IOSTAT\_END** (13.10.2.16) to the **iostat** argument. Otherwise, if an end-of-record condition occurs, the **defined input** procedure shall assign the value of the **named constant IOSTAT\_EOR** (13.10.2.17) to **iostat**. Otherwise, the **defined input/output** procedure shall assign the value zero to the **iostat** argument.
- 18 If the **defined input/output** procedure returns a nonzero value for the **iostat** argument, the procedure shall also return an explanatory message in the **iomsg** argument. Otherwise, the procedure shall not change the value of the **iomsg** argument.

**NOTE 9.43**

The values of the **iostat** and **iomsg** arguments set in a **defined input/output** procedure need not be passed to all of the parent data transfer statements.

- 19 If the **iostat** argument of the **defined input/output** procedure has a nonzero value when that procedure returns, and the processor therefore **terminates execution** of the program as described in 9.11, the processor shall make the value of the **iomsg** argument available in a processor-dependent manner.

- 1 20 While a parent **READ statement** is active, an input/output statement shall not read from any **external unit** other  
 2 than the one specified by the **unit dummy argument** and shall not perform output to any **external unit**.
- 3 21 While a parent **WRITE** or **PRINT** statement is active, an input/output statement shall not perform output to  
 4 any **external unit** other than the one specified by the **unit dummy argument** and shall not read from any **external**  
 5 **unit**.
- 6 22 While a parent data transfer statement is active, a data transfer statement that specifies an **internal file** is  
 7 permitted.
- 8 23 **OPEN**, **CLOSE**, **BACKSPACE**, **ENDFILE**, and **REWIND** statements shall not be executed while a parent data  
 9 transfer statement is active.
- 10 24 A **defined input/output** procedure may use a format specification with a **DT edit descriptor** for handling a  
 11 component of the derived type that is itself of a derived type. A child data transfer statement that is a list  
 12 directed or namelist input/output statement may contain a list item of derived type.
- 13 25 Because a child data transfer statement does not position the file prior to data transfer, the child data transfer  
 14 statement starts transferring data from where the file was positioned by the parent data transfer statement's  
 15 most recently processed **effective item** or edit descriptor. This is not necessarily at the beginning of a record.
- 16 26 The edit descriptors **T** and **TL** used on **unit** by a child data transfer statement shall not cause the file to be  
 17 positioned before the file position at the time the **defined input/output** procedure was invoked.

**NOTE 9.44**

A **defined input/output** procedure could use **INQUIRE** to determine the settings of **BLANK=**, **PAD=**, **ROUND=**, **DECIMAL=**, and **DELIM=** for an **external unit**. The **INQUIRE statement** provides values as specified in 9.10.

- 18 27 Neither a parent nor child data transfer statement shall be asynchronous.
- 19 28 A **defined input/output** procedure, and any procedures invoked therefrom, shall not define, nor cause to become  
 20 undefined, any **storage unit** referenced by any input/output list item, the corresponding format, or any specifier  
 21 in any active parent data transfer statement, except through the **dtv** argument.

**NOTE 9.45**

A child data transfer statement shall not specify the **ID=**, **POS=**, or **REC=** specifiers in an input/output control list.

**NOTE 9.46**

A simple example of derived type formatted output follows. The derived type variable **chairman** has two components. The type and an associated write formatted procedure are defined in a module so as to be accessible from wherever they might be needed. It would also be possible to check that **iotype** indeed has the value 'DT' and to set **iostat** and **iomsg** accordingly.

```

MODULE p

  TYPE :: person
    CHARACTER (LEN=20) :: name
    INTEGER :: age
  CONTAINS
    PROCEDURE,PRIVATE :: pwf
    GENERIC             :: WRITE(FORMATTED) => pwf
  END TYPE person

CONTAINS

```

## NOTE 9.46 (cont.)

```

SUBROUTINE pwf (dtv,unit,iotype,vlist,iostat,iomsg)
! argument definitions
  CLASS(person), INTENT(IN) :: dtv
  INTEGER, INTENT(IN) :: unit
  CHARACTER (LEN=*), INTENT(IN) :: iotype
  INTEGER, INTENT(IN) :: vlist(:)
  INTEGER, INTENT(OUT) :: iostat
  CHARACTER (LEN=*), INTENT(INOUT) :: iomsg
! local variable
  CHARACTER (LEN=9) :: pfmt

!  vlist(1) and (2) are to be used as the field widths of the two
!  components of the derived type variable. First set up the format to
!  be used for output.
  WRITE(pfmt,'(A,I2,A,I2,A)' ) '(A', vlist(1), ',I', vlist(2), ' )'

!  now the basic output statement
  WRITE(unit, FMT=pfmt, IOSTAT=iostat) dtv%name, dtv%age

END SUBROUTINE pwf

END MODULE p

PROGRAM committee
  USE p
  INTEGER id, members
  TYPE (person) :: chairman
  ...
  WRITE(6, FMT="(I2, DT (15,6), I5)" ) id, chairman, members
! this writes a record with four fields, with lengths 2, 15, 6, 5
! respectively

END PROGRAM

```

## NOTE 9.47

In the following example, the variables of the derived type `node` form a linked list, with a single value at each node. The subroutine `pwf` is used to write the values in the list, one per line.

```

MODULE p

  TYPE node
    INTEGER :: value = 0
    TYPE (NODE), POINTER :: next_node => NULL ( )
  CONTAINS
    PROCEDURE,PRIVATE :: pwf
    GENERIC :: WRITE(FORMATTED) => pwf
  END TYPE node

CONTAINS

  SUBROUTINE pwf (dtv,unit,iotype,vlist,iostat,iomsg)
! Write the chain of values, each on a separate line in I9 format.
    CLASS(node), INTENT(IN) :: dtv

```

## NOTE 9.47 (cont.)

```

    INTEGER, INTENT(IN) :: unit
    CHARACTER (LEN=*), INTENT(IN) :: iotype
    INTEGER, INTENT(IN) :: vlist(:)
    INTEGER, INTENT(OUT) :: iostat
    CHARACTER (LEN=*), INTENT(INOUT) :: iomsg

    WRITE(unit,'(i9 /)', IOSTAT = iostat) dtv%value
    IF(iostat/=0) RETURN
    IF(ASSOCIATED(dtv%next_node)) WRITE(unit,'(dt)', IOSTAT=iostat) dtv%next_node
END SUBROUTINE pwf

END MODULE p

```

## 9.6.4.8.4 Resolving defined input/output procedure references

- 1 A suitable [generic interface](#) for [defined input/output](#) of an [effective item](#) is one that has a [defined-io-generic-spec](#) that is appropriate to the direction (read or write) and form (formatted or unformatted) of the data transfer as specified in [9.6.4.8.3](#), and has a [specific interface](#) whose `dtv` argument is compatible with the [effective item](#) according to the rules for argument association in [12.5.2.4](#).
- 2 When an [effective item](#) ([9.6.3](#)) that is of derived type is encountered during a data transfer, [defined input/output](#) occurs if both of the following conditions are true.
  - (1) The circumstances of the input/output are such that [defined input/output](#) is permitted; that is, either
    - (a) the transfer was initiated by a list-directed, namelist, or unformatted input/output statement, or
    - (b) a format specification is supplied for the data transfer statement, and the edit descriptor corresponding to the [effective item](#) is a DT edit descriptor.
  - (2) A suitable [defined input/output](#) procedure is available; that is, either
    - (a) the [declared type](#) of the [effective item](#) has a suitable generic [type-bound procedure](#), or
    - (b) a suitable [generic interface](#) is accessible.
- 3 If (2a) is true, the procedure referenced is determined as for explicit [type-bound procedure](#) references ([12.5](#)); that is, the [binding](#) with the appropriate [specific interface](#) is located in the [declared type](#) of the [effective item](#), and the corresponding [binding](#) in the [dynamic type](#) of the [effective item](#) is selected.
- 4 If (2a) is false and (2b) is true, the reference is to the procedure identified by the appropriate [specific interface](#) in the [interface block](#).

## 9.6.5 Termination of data transfer statements

- 1 Termination of a data transfer statement occurs when
  - format processing encounters a colon or data edit descriptor and there are no remaining elements in the [input-item-list](#) or [output-item-list](#),
  - unformatted or list-directed data transfer exhausts the [input-item-list](#) or [output-item-list](#),
  - namelist output exhausts the [namelist-group-object-list](#),
  - an error condition occurs,
  - an end-of-file condition occurs,
  - a slash (/) is encountered as a value separator ([10.10](#), [10.11](#)) in the record being read during list-directed or namelist input, or
  - an end-of-record condition occurs during execution of a nonadvancing input statement ([9.11](#)).

## 9.7 Waiting on pending data transfer

### 9.7.1 Wait operation

- 1 Execution of an asynchronous [data transfer statement](#) in which neither an error, end-of-record, nor end-of-file condition occurs initiates a pending data transfer operation. There may be multiple pending data transfer operations for the same or multiple [units](#) simultaneously. A pending data transfer operation remains pending until a corresponding wait operation is performed. A wait operation may be performed by a [BACKSPACE](#), [CLOSE](#), [ENDFILE](#), [FLUSH](#), [INQUIRE](#), [PRINT](#), [READ](#), [REWIND](#), [WAIT](#), or [WRITE](#) statement.
- 2 A wait operation completes the processing of a pending data transfer operation. Each wait operation completes only a single data transfer operation, although a single statement may perform multiple wait operations.
- 3 If the actual data transfer is not yet complete, the wait operation first waits for its completion. If the data transfer operation is an input operation that completed without error, the [storage units](#) of the input/output storage sequence then become defined with the values as described in [9.6.2.15](#) and [9.6.4.5](#).
- 4 If any error, end-of-file, or end-of-record conditions occur, the applicable actions specified by the [IOSTAT=](#), [IOMSG=](#), [ERR=](#), [END=](#), and [EOR=](#) specifiers of the statement that performs the wait operation are taken.
- 5 If an error or end-of-file condition occurs during a wait operation for a [unit](#), the processor performs a wait operation for all pending data transfer operations for that [unit](#).

#### NOTE 9.48

Error, end-of-file, and end-of-record conditions can be raised either during the [data transfer statement](#) that initiates asynchronous input/output, a subsequent asynchronous [data transfer statement](#) for the same [unit](#), or during the wait operation. If such conditions are raised during a [data transfer statement](#), they trigger actions according to the [IOSTAT=](#), [ERR=](#), [END=](#), and [EOR=](#) specifiers of that statement; if they are raised during the wait operation, the actions are in accordance with the specifiers of the statement that performs the wait operation.

- 6 After completion of the wait operation, the data transfer operation and its input/output storage sequence are no longer considered to be pending.

### 9.7.2 WAIT statement

- 1 A WAIT statement performs a wait operation for specified pending asynchronous data transfer operations.

R922 *wait-stmt*                                    **is** WAIT (*wait-spec-list*)

R923 *wait-spec*                                    **is** [ UNIT = ] *file-unit-number*  
    **or** END = *label*  
    **or** EOR = *label*  
    **or** ERR = *label*  
    **or** ID = *scalar-int-expr*  
    **or** IOMSG = *iomsg-variable*  
    **or** IOSTAT = *scalar-int-variable*

C936 No specifier shall appear more than once in a given *wait-spec-list*.

C937 A *file-unit-number* shall be specified in a *wait-spec-list*; if the optional characters UNIT= are omitted, the *file-unit-number* shall be the first item in the *wait-spec-list*.

C938 (R923) The *label* in the [ERR=](#), [EOR=](#), or [END=](#) specifier shall be the statement label of a [branch target statement](#) that appears in the same [inclusive scope](#) as the WAIT statement.

- 2 The [IOSTAT=](#), [ERR=](#), [EOR=](#), [END=](#), and [IOMSG=](#) specifiers are described in [9.11](#).

- 1 3 The value of the expression specified in the ID= specifier shall be zero or the identifier of a pending data transfer  
 2 operation for the specified **unit**. If the ID= specifier appears, a wait operation for the specified data transfer  
 3 operation, if any, is performed. If the ID= specifier is omitted, wait operations for all pending data transfers for  
 4 the specified **unit** are performed.
- 5 4 Execution of a WAIT statement specifying a **unit** that does not exist, has no file connected to it, or is not open  
 6 for asynchronous input/output is permitted, provided that the WAIT statement has no ID= specifier; such a  
 7 WAIT statement does not cause an error or end-of-file condition to occur.

**NOTE 9.49**

An **EOR= specifier** has no effect if the pending data transfer operation is not a nonadvancing read. An **END= specifier** has no effect if the pending data transfer operation is not a READ.

## 8 **9.8 File positioning statements**

### 9 **9.8.1 Syntax**

10 R924 *backspace-stmt*                    **is** BACKSPACE *file-unit-number*  
 11    **or** BACKSPACE ( *position-spec-list* )

12 R925 *endfile-stmt*                    **is** ENDFILE *file-unit-number*  
 13    **or** ENDFILE ( *position-spec-list* )

14 R926 *rewind-stmt*                    **is** REWIND *file-unit-number*  
 15    **or** REWIND ( *position-spec-list* )

- 16 1 A **unit** that is connected for direct access shall not be referred to by a BACKSPACE, ENDFILE, or REWIND  
 17 statement. A **unit** that is connected for unformatted stream access shall not be referred to by a BACKSPACE  
 18 statement. A **unit** that is connected with an ACTION= specifier having the value READ shall not be referred  
 19 to by an ENDFILE statement.

20 R927 *position-spec*                    **is** [ UNIT = ] *file-unit-number*  
 21    **or** IOMSG = *iomsg-variable*  
 22    **or** IOSTAT = *scalar-int-variable*  
 23    **or** ERR = *label*

24 C939 No specifier shall appear more than once in a given *position-spec-list*.

25 C940 A *file-unit-number* shall be specified in a *position-spec-list*; if the optional characters UNIT= are omitted,  
 26 the *file-unit-number* shall be the first item in the *position-spec-list*.

27 C941 (R927) The *label* in the **ERR= specifier** shall be the statement label of a **branch target statement** that  
 28 appears in the same **inclusive scope** as the file positioning statement.

29 2 The **IOSTAT=**, **ERR=**, and **IOMSG=** specifiers are described in 9.11.

30 3 Execution of a file positioning statement performs a wait operation for all pending asynchronous data transfer  
 31 operations for the specified **unit**.

### 32 **9.8.2 BACKSPACE statement**

- 33 1 Execution of a BACKSPACE statement causes the file connected to the specified **unit** to be positioned before  
 34 the current record if there is a current record, or before the preceding record if there is no current record. If the  
 35 file is at its initial point, the position of the file is not changed.

**NOTE 9.50**

If the preceding record is an endfile record, the file is positioned before the endfile record.

- 1 2 If a BACKSPACE statement causes the implicit writing of an endfile record, the file is positioned before the
- 2 record that precedes the endfile record.
- 3 3 Backspacing a file that is connected but does not exist is prohibited.
- 4 4 Backspacing over records written using list-directed or namelist formatting is prohibited.

**NOTE 9.51**

An example of a BACKSPACE statement is:

```
BACKSPACE (10, IOSTAT = N)
```

**5 9.8.3 ENDFILE statement**

- 6 1 Execution of an ENDFILE statement for a file connected for sequential access writes an endfile record as the next
- 7 record of the file. The file is then positioned after the endfile record, which becomes the last record of the file.
- 8 If the file also may be connected for direct access, only those records before the endfile record are considered
- 9 have been written. Thus, only those records may be read during subsequent direct access connections to the file.
- 10 2 After execution of an ENDFILE statement for a file connected for sequential access, a BACKSPACE or REWIND
- 11 statement shall be used to reposition the file prior to execution of any data transfer input/output statement or
- 12 ENDFILE statement.
- 13 3 Execution of an ENDFILE statement for a file connected for stream access causes the terminal point of the file
- 14 to become equal to the current file position. Only file storage units before the current position are considered
- 15 to have been written; thus only those file storage units may be subsequently read. Subsequent stream output
- 16 statements may be used to write further data to the file.
- 17 4 Execution of an ENDFILE statement for a file that is connected but does not exist creates the file; if the file is
- 18 connected for sequential access, it is created prior to writing the endfile record.

**NOTE 9.52**

An example of an ENDFILE statement is:

```
ENDFILE K
```

**19 9.8.4 REWIND statement**

- 20 1 Execution of a REWIND statement causes the specified file to be positioned at its initial point.

**NOTE 9.53**

If the file is already positioned at its initial point, execution of this statement has no effect on the position of the file.

- 21 2 Execution of a REWIND statement for a file that is connected but does not exist is permitted and has no effect
- 22 on any file.

**NOTE 9.54**

An example of a REWIND statement is:

```
REWIND 10
```

## 9.9 FLUSH statement

R928 *flush-stmt* is FLUSH *file-unit-number*  
or FLUSH ( *flush-spec-list* )

R929 *flush-spec* is [UNIT =] *file-unit-number*  
or IOSTAT = *scalar-int-variable*  
or IOMSG = *iomsg-variable*  
or ERR = *label*

C942 No specifier shall appear more than once in a given *flush-spec-list*.

C943 A *file-unit-number* shall be specified in a *flush-spec-list*; if the optional characters UNIT= are omitted from the unit specifier, the *file-unit-number* shall be the first item in the *flush-spec-list*.

C944 (R929) The *label* in the ERR= specifier shall be the statement label of a *branch target statement* that appears in the same *inclusive scope* as the FLUSH statement.

1 The IOSTAT=, IOMSG= and ERR= specifiers are described in 9.11.

2 Execution of a FLUSH statement causes data written to an *external file* to be available to other processes, or causes data placed in an *external file* by means other than Fortran to be available to a *READ statement*. These actions are processor dependent.

3 Execution of a FLUSH statement for a file that is connected but does not exist is permitted and has no effect on any file. A FLUSH statement has no effect on file position.

4 Execution of a FLUSH statement performs a wait operation for all pending asynchronous data transfer operations for the specified *unit*.

### NOTE 9.55

Because this part of ISO/IEC 1539 does not specify the mechanism of file storage, the exact meaning of the flush operation is not precisely defined. It is expected that the flush operation will make all data written to a file available to other processes or devices, or make data recently added to a file by other processes or devices available to the program via a subsequent read operation. This is commonly called “flushing input/output buffers”.

### NOTE 9.56

An example of a FLUSH statement is:

```
FLUSH (10, IOSTAT = N)
```

## 9.10 File inquiry statement

### 9.10.1 Forms of the INQUIRE statement

1 The INQUIRE statement can be used to inquire about properties of a particular named file, of the connection to a particular *unit*, or the number of file storage units required for an output list. There are three forms of the INQUIRE statement: inquire by file, which uses the FILE= specifier, inquire by *unit*, which uses the UNIT= specifier, and inquire by output list, which uses only the *IOLength= specifier*. Assignments to specifier variables are performed as if by *intrinsic assignment*.

2 For inquiry by *unit*, the *unit* specified need not exist or be connected to a file. If it is connected to a file, the inquiry is being made about the connection and about the file connected.

3 An INQUIRE statement may be executed before, while, or after a file is connected to a *unit*. All values assigned

by an INQUIRE statement are those that are current at the time the statement is executed.

R930 *inquire-stmt* is INQUIRE ( *inquire-spec-list* )  
 or INQUIRE ( IOLENGTH = *scalar-int-variable* ) ■  
 ■ *output-item-list*

#### NOTE 9.57

Examples of INQUIRE statements are:

```
INQUIRE (IOLENGTH = IOL) A (1:N)
INQUIRE (UNIT = JOAN, OPENED = LOG_01, NAMED = LOG_02, &
  FORM = CHAR_VAR, IOSTAT = IOS)
```

## 9.10.2 Inquiry specifiers

### 9.10.2.1 Syntax

- 1 Unless constrained, the following inquiry specifiers may be used in either of the inquire by file or inquire by unit forms of the INQUIRE statement.

R931 *inquire-spec* is [ UNIT = ] *file-unit-number*  
 or FILE = *file-name-expr*  
 or ACCESS = *scalar-default-char-variable*  
 or ACTION = *scalar-default-char-variable*  
 or ASYNCHRONOUS = *scalar-default-char-variable*  
 or BLANK = *scalar-default-char-variable*  
 or DECIMAL = *scalar-default-char-variable*  
 or DELIM = *scalar-default-char-variable*  
 or DIRECT = *scalar-default-char-variable*  
 or ENCODING = *scalar-default-char-variable*  
 or ERR = *label*  
 or EXIST = *scalar-logical-variable*  
 or FORM = *scalar-default-char-variable*  
 or FORMATTED = *scalar-default-char-variable*  
 or ID = *scalar-int-expr*  
 or IOMSG = *iomsg-variable*  
 or IOSTAT = *scalar-int-variable*  
 or NAME = *scalar-default-char-variable*  
 or NAMED = *scalar-logical-variable*  
 or NEXTREC = *scalar-int-variable*  
 or NUMBER = *scalar-int-variable*  
 or OPENED = *scalar-logical-variable*  
 or PAD = *scalar-default-char-variable*  
 or PENDING = *scalar-logical-variable*  
 or POS = *scalar-int-variable*  
 or POSITION = *scalar-default-char-variable*  
 or READ = *scalar-default-char-variable*  
 or READWRITE = *scalar-default-char-variable*  
 or RECL = *scalar-int-variable*  
 or ROUND = *scalar-default-char-variable*  
 or SEQUENTIAL = *scalar-default-char-variable*  
 or SIGN = *scalar-default-char-variable*  
 or SIZE = *scalar-int-variable*  
 or STREAM = *scalar-default-char-variable*  
 or UNFORMATTED = *scalar-default-char-variable*  
 or WRITE = *scalar-default-char-variable*

- 1 C945 No specifier shall appear more than once in a given *inquire-spec-list*.
- 2 C946 An *inquire-spec-list* shall contain one **FILE= specifier** or one *file-unit-number*, but not both.
- 3 C947 In the inquire by unit form of the INQUIRE statement, if the optional characters UNIT= are omitted,  
4 the *file-unit-number* shall be the first item in the *inquire-spec-list*.
- 5 C948 If an **ID= specifier** appears in an *inquire-spec-list*, a **PENDING= specifier** shall also appear.
- 6 C949 (R929) The *label* in the **ERR= specifier** shall be the statement label of a **branch target statement** that  
7 appears in the same **inclusive scope** as the INQUIRE statement.
- 8 2 If *file-unit-number* identifies an **internal unit** (9.6.4.8.3), an error condition occurs.
- 9 3 When a returned value of a specifier other than the **NAME= specifier** is of type character, the value returned is  
10 in upper case.
- 11 4 If an error condition occurs during execution of an INQUIRE statement, all of the inquiry specifier variables  
12 become undefined, except for variables in the **IOSTAT=** and **IOMSG=** specifiers (if any).
- 13 5 The **IOSTAT=**, **ERR=**, and **IOMSG=** specifiers are described in 9.11.

#### 14 9.10.2.2 FILE= specifier in the INQUIRE statement

- 15 1 The value of the *file-name-expr* in the **FILE=** specifier specifies the name of the file being inquired about. The  
16 named file need not exist or be connected to a **unit**. The value of the *file-name-expr* shall be of a form acceptable  
17 to the processor as a file name. Any trailing blanks are ignored. The interpretation of case is processor dependent.

#### 18 9.10.2.3 ACCESS= specifier in the INQUIRE statement

- 19 1 The *scalar-default-char-variable* in the **ACCESS=** specifier is assigned the value SEQUENTIAL if the connection  
20 is for sequential access, DIRECT if the connection is for direct access, or STREAM if the connection is for stream  
21 access. If there is no connection, it is assigned the value UNDEFINED.

#### 22 9.10.2.4 ACTION= specifier in the INQUIRE statement

- 23 1 The *scalar-default-char-variable* in the **ACTION=** specifier is assigned the value READ if the connection is for  
24 input only, WRITE if the connection is for output only, and READWRITE if the connection is for both input  
25 and output. If there is no connection, the *scalar-default-char-variable* is assigned the value UNDEFINED.

#### 26 9.10.2.5 ASYNCHRONOUS= specifier in the INQUIRE statement

- 27 1 The *scalar-default-char-variable* in the **ASYNCHRONOUS=** specifier is assigned the value YES if the connection  
28 allows asynchronous input/output; it is assigned the value NO if the connection does not allow asynchronous  
29 input/output. If there is no connection, the *scalar-default-char-variable* is assigned the value UNDEFINED.

#### 30 9.10.2.6 BLANK= specifier in the INQUIRE statement

- 31 1 The *scalar-default-char-variable* in the **BLANK=** specifier is assigned the value ZERO or NULL, corresponding  
32 to the blank interpretation mode in effect for a connection for formatted input/output. If there is no connection,  
33 or if the connection is not for formatted input/output, the *scalar-default-char-variable* is assigned the value  
34 UNDEFINED.

#### 35 9.10.2.7 DECIMAL= specifier in the INQUIRE statement

- 36 1 The *scalar-default-char-variable* in the **DECIMAL=** specifier is assigned the value COMMA or POINT, corres-  
37 ponding to the decimal edit mode in effect for a connection for formatted input/output. If there is no connection,  
38 or if the connection is not for formatted input/output, the *scalar-default-char-variable* is assigned the value  
39 UNDEFINED.

### 9.10.2.8 DELIM= specifier in the INQUIRE statement

- 1 The *scalar-default-char-variable* in the DELIM= specifier is assigned the value APOSTROPHE, QUOTE, or NONE, corresponding to the delimiter mode in effect for a connection for formatted input/output. If there is no connection or if the connection is not for formatted input/output, the *scalar-default-char-variable* is assigned the value UNDEFINED.

### 9.10.2.9 DIRECT= specifier in the INQUIRE statement

- 1 The *scalar-default-char-variable* in the DIRECT= specifier is assigned the value YES if DIRECT is included in the set of allowed access methods for the file, NO if DIRECT is not included in the set of allowed access methods for the file, and UNKNOWN if the processor is unable to determine whether DIRECT is included in the set of allowed access methods for the file or if the *unit* identified by *file-unit-number* is not connected to a file.

### 9.10.2.10 ENCODING= specifier in the INQUIRE statement

- 1 The *scalar-default-char-variable* in the ENCODING= specifier is assigned the value UTF-8 if the connection is for formatted input/output with an encoding form of UTF-8, and is assigned the value UNDEFINED if the connection is for unformatted input/output. If there is no connection, it is assigned the value UTF-8 if the processor is able to determine that the encoding form of the file is UTF-8; if the processor is unable to determine the encoding form of the file or if the *unit* identified by *file-unit-number* is not connected to a file, the variable is assigned the value UNKNOWN.

#### NOTE 9.58

The value assigned could be something other than UTF-8, UNDEFINED, or UNKNOWN if the processor supports other specific encoding forms (e.g. UTF-16BE).

### 9.10.2.11 EXIST= specifier in the INQUIRE statement

- 1 Execution of an INQUIRE by file statement causes the *scalar-logical-variable* in the EXIST= specifier to be assigned the value true if there exists a file with the specified name; otherwise, false is assigned. Execution of an INQUIRE by unit statement causes true to be assigned if the specified *unit* exists; otherwise, false is assigned.

### 9.10.2.12 FORM= specifier in the INQUIRE statement

- 1 The *scalar-default-char-variable* in the FORM= specifier is assigned the value FORMATTED if the connection is for formatted input/output, and is assigned the value UNFORMATTED if the connection is for unformatted input/output. If there is no connection, it is assigned the value UNDEFINED.

### 9.10.2.13 FORMATTED= specifier in the INQUIRE statement

- 1 The *scalar-default-char-variable* in the FORMATTED= specifier is assigned the value YES if FORMATTED is included in the set of allowed forms for the file, NO if FORMATTED is not included in the set of allowed forms for the file, and UNKNOWN if the processor is unable to determine whether FORMATTED is included in the set of allowed forms for the file or if the *unit* identified by *file-unit-number* is not connected to a file.

### 9.10.2.14 ID= specifier in the INQUIRE statement

- 1 The value of the expression specified in the ID= specifier shall be the identifier of a pending data transfer operation for the specified *unit*. This specifier interacts with the PENDING= specifier (9.10.2.21).

### 9.10.2.15 NAME= specifier in the INQUIRE statement

- 1 The *scalar-default-char-variable* in the NAME= specifier is assigned the value of the name of the file if the file has a name; otherwise, it becomes undefined. The value assigned shall be suitable for use as the value of the *file-name-expr* in the FILE= specifier in an OPEN statement.

**NOTE 9.59**

If this specifier appears in an INQUIRE by file statement, its value is not necessarily the same as the name given in the FILE= specifier.

The processor could assign a file name qualified by a user identification, device, directory, or other relevant information.

2 The case of the characters assigned to *scalar-default-char-variable* is processor dependent.

#### 2 **9.10.2.16 NAMED= specifier in the INQUIRE statement**

3 1 The *scalar-logical-variable* in the NAMED= specifier is assigned the value true if the file has a name; otherwise,  
4 it is assigned the value false.

#### 5 **9.10.2.17 NEXTREC= specifier in the INQUIRE statement**

6 1 The *scalar-int-variable* in the NEXTREC= specifier is assigned the value  $n + 1$ , where  $n$  is the record number of  
7 the last record read from or written to the connection for direct access. If there is a connection but no records have  
8 been read or written since the connection, the *scalar-int-variable* is assigned the value 1. If there is no connection,  
9 the connection is not for direct access, or the position is indeterminate because of a previous error condition, the  
10 *scalar-int-variable* becomes undefined. If there are pending data transfer operations for the specified *unit*, the  
11 value assigned is computed as if all the pending data transfers had already completed.

#### 12 **9.10.2.18 NUMBER= specifier in the INQUIRE statement**

13 1 Execution of an INQUIRE by file statement causes the *scalar-int-variable* in the NUMBER= specifier to be  
14 assigned the value of the *external unit* number of the unit that is connected to the file. If more than one unit  
15 on an *image* is connected to the file, which of the connected external unit numbers is assigned to the *scalar-int-*  
16 *variable* is processor dependent. If there is no *unit* connected to the file, the value  $-1$  is assigned. Execution of  
17 an INQUIRE by *unit* statement causes the *scalar-int-variable* to be assigned the value of *file-unit-number*.

#### 18 **9.10.2.19 OPENED= specifier in the INQUIRE statement**

19 1 Execution of an INQUIRE by file statement causes the *scalar-logical-variable* in the OPENED= specifier to be  
20 assigned the value true if the file specified is connected to a *unit*; otherwise, false is assigned. Execution of an  
21 INQUIRE by unit statement causes the *scalar-logical-variable* to be assigned the value true if the specified *unit*  
22 is connected to a file; otherwise, false is assigned.

#### 23 **9.10.2.20 PAD= specifier in the INQUIRE statement**

24 1 The *scalar-default-char-variable* in the PAD= specifier is assigned the value YES or NO, corresponding to the  
25 pad mode in effect for a connection for formatted input/output. If there is no connection or if the connection is  
26 not for formatted input/output, the *scalar-default-char-variable* is assigned the value UNDEFINED.

#### 27 **9.10.2.21 PENDING= specifier in the INQUIRE statement**

28 1 The PENDING= specifier is used to determine whether previously pending asynchronous data transfers are  
29 complete. A data transfer operation is previously pending if it is pending at the beginning of execution of the  
30 INQUIRE statement.

31 2 If an ID= specifier appears and the specified data transfer operation is complete, then the variable specified in  
32 the PENDING= specifier is assigned the value false and the INQUIRE statement performs the wait operation  
33 for the specified data transfer.

34 3 If the ID= specifier is omitted and all previously pending data transfer operations for the specified *unit* are  
35 complete, then the variable specified in the PENDING= specifier is assigned the value false and the INQUIRE  
36 statement performs wait operations for all previously pending data transfers for the specified *unit*.

- 1 4 In all other cases, the variable specified in the PENDING= specifier is assigned the value true, no wait operations  
 2 are performed, and the previously pending data transfers remain pending after the execution of the INQUIRE  
 3 statement.

#### NOTE 9.60

The processor has considerable flexibility in defining when it considers a transfer to be complete. Any of the following approaches could be used:

- The INQUIRE statement could consider an asynchronous data transfer to be incomplete until after the corresponding wait operation. In this case PENDING= would always return true unless there were no previously pending data transfers for the [unit](#).
- The INQUIRE statement could wait for all specified data transfers to complete and then always return false for PENDING=.
- The INQUIRE statement could actually test the state of the specified data transfer operations.

#### 4 9.10.2.22 POS= specifier in the INQUIRE statement

- 5 1 The *scalar-int-variable* in the POS= specifier is assigned the number of the [file storage unit](#) immediately following  
 6 the current position of a file connected for stream access. If the file is positioned at its terminal position, the  
 7 variable is assigned a value one greater than the number of the highest-numbered [file storage unit](#) in the file.  
 8 If there are pending data transfer operations for the specified unit, the value assigned is computed as if all the  
 9 pending data transfers had already completed. If there is no connection, the file is not connected for stream  
 10 access, or if the position of the file is indeterminate because of previous error conditions, the variable becomes  
 11 undefined.

#### 12 9.10.2.23 POSITION= specifier in the INQUIRE statement

- 13 1 The *scalar-default-char-variable* in the POSITION= specifier is assigned the value REWIND if the connection  
 14 was opened for positioning at its initial point, APPEND if the connection was opened for positioning before its  
 15 endfile record or at its terminal point, and ASIS if the connection was opened without changing its position.  
 16 If there is no connection or if the file is connected for direct access, the *scalar-default-char-variable* is assigned  
 17 the value UNDEFINED. If the file has been repositioned since the connection, the *scalar-default-char-variable*  
 18 is assigned a processor-dependent value, which shall not be REWIND unless the file is positioned at its initial  
 19 point and shall not be APPEND unless the file is positioned so that its endfile record is the next record or at its  
 20 terminal point if it has no endfile record.

#### 21 9.10.2.24 READ= specifier in the INQUIRE statement

- 22 1 The *scalar-default-char-variable* in the READ= specifier is assigned the value YES if READ is included in the  
 23 set of allowed actions for the file, NO if READ is not included in the set of allowed actions for the file, and  
 24 UNKNOWN if the processor is unable to determine whether READ is included in the set of allowed actions for  
 25 the file or if the [unit](#) identified by [file-unit-number](#) is not connected to a file.

#### 26 9.10.2.25 READWRITE= specifier in the INQUIRE statement

- 27 1 The *scalar-default-char-variable* in the READWRITE= specifier is assigned the value YES if READWRITE is  
 28 included in the set of allowed actions for the file, NO if READWRITE is not included in the set of allowed actions  
 29 for the file, and UNKNOWN if the processor is unable to determine whether READWRITE is included in the  
 30 set of allowed actions for the file or if the [unit](#) identified by [file-unit-number](#) is not connected to a file.

#### 31 9.10.2.26 RECL= specifier in the INQUIRE statement

- 32 1 The *scalar-int-variable* in the RECL= specifier is assigned the value of the record length of a connection for direct  
 33 access, or the value of the maximum record length of a connection for sequential access. If the connection is for  
 34 formatted input/output, the length is the number of characters for all records that contain only characters of

default kind. If the connection is for unformatted input/output, the length is measured in [file storage units](#). If there is no connection, the *scalar-int-variable* is assigned the value  $-1$ , and if the connection is for stream access, the *scalar-int-variable* is assigned the value  $-2$ .

#### 9.10.2.27 ROUND= specifier in the INQUIRE statement

The *scalar-default-char-variable* in the ROUND= specifier is assigned the value UP, DOWN, ZERO, NEAREST, COMPATIBLE, or PROCESSOR\_DEFINED, corresponding to the input/output rounding mode in effect for a connection for formatted input/output. If there is no connection or if the connection is not for formatted input/output, the *scalar-default-char-variable* is assigned the value UNDEFINED. The processor shall return the value PROCESSOR\_DEFINED only if the behavior of the input/output rounding mode is different from that of the UP, DOWN, ZERO, NEAREST, and COMPATIBLE modes.

#### 9.10.2.28 SEQUENTIAL= specifier in the INQUIRE statement

The *scalar-default-char-variable* in the SEQUENTIAL= specifier is assigned the value YES if SEQUENTIAL is included in the set of allowed access methods for the file, NO if SEQUENTIAL is not included in the set of allowed access methods for the file, and UNKNOWN if the processor is unable to determine whether SEQUENTIAL is included in the set of allowed access methods for the file or if the [unit](#) identified by *file-unit-number* is not connected to a file.

#### 9.10.2.29 SIGN= specifier in the INQUIRE statement

The *scalar-default-char-variable* in the SIGN= specifier is assigned the value PLUS, SUPPRESS, or PROCESSOR\_DEFINED, corresponding to the sign mode in effect for a connection for formatted input/output. If there is no connection, or if the connection is not for formatted input/output, the *scalar-default-char-variable* is assigned the value UNDEFINED.

#### 9.10.2.30 SIZE= specifier in the INQUIRE statement

- 1 The *scalar-int-variable* in the SIZE= specifier is assigned the size of the file in [file storage units](#). If the file size cannot be determined or if the [unit](#) identified by *file-unit-number* is not connected to a file, the variable is assigned the value  $-1$ .
- 2 For a file that may be connected for stream access, the file size is the number of the highest-numbered [file storage unit](#) in the file.
- 3 For a file that may be connected for sequential or direct access, the file size may be different from the number of [storage units](#) implied by the data in the records; the exact relationship is processor dependent.
- 4 If there are pending data transfer operations for the specified unit, the value assigned is computed as if all the pending data transfers had already completed.

#### 9.10.2.31 STREAM= specifier in the INQUIRE statement

The *scalar-default-char-variable* in the STREAM= specifier is assigned the value YES if STREAM is included in the set of allowed access methods for the file, NO if STREAM is not included in the set of allowed access methods for the file, and UNKNOWN if the processor is unable to determine whether STREAM is included in the set of allowed access methods for the file or if the [unit](#) identified by *file-unit-number* is not connected to a file.

#### 9.10.2.32 UNFORMATTED= specifier in the INQUIRE statement

The *scalar-default-char-variable* in the UNFORMATTED= specifier is assigned the value YES if UNFORMATTED is included in the set of allowed forms for the file, NO if UNFORMATTED is not included in the set of allowed forms for the file, and UNKNOWN if the processor is unable to determine whether UNFORMATTED is included in the set of allowed forms for the file or if the [unit](#) identified by *file-unit-number* is not connected to a file.

### 9.10.2.33 WRITE= specifier in the INQUIRE statement

- 1 The *scalar-default-char-variable* in the WRITE= specifier is assigned the value YES if WRITE is included in the set of allowed actions for the file, NO if WRITE is not included in the set of allowed actions for the file, and UNKNOWN if the processor is unable to determine whether WRITE is included in the set of allowed actions for the file or if the *unit* identified by *file-unit-number* is not connected to a file.

### 9.10.3 Inquire by output list

- 1 The *scalar-int-variable* in the IOLENGTH= specifier is assigned the processor-dependent number of *file storage units* that would be required to store the data of the output list in an unformatted file. The value shall be suitable as a RECL= specifier in an OPEN statement that connects a file for unformatted direct access if data will be read from or written to the file using data transfer statements with an input/output list that specifies transfer of a sequence of objects having the same types, type parameters, and extents, in the same order as the output list in the INQUIRE statement.
- 2 The output list in an INQUIRE statement shall not contain any derived-type list items that require a *defined input/output* procedure as described in subclause 9.6.3. If a derived-type list item appears in the output list, the value returned for the IOLENGTH= specifier assumes that no *defined input/output* procedure will be invoked.

## 9.11 Error, end-of-record, and end-of-file conditions

### 9.11.1 Occurrence of input/output conditions

- 1 The set of input/output error conditions is processor dependent. Except as otherwise specified, when an error condition occurs or is detected is processor dependent.
- 2 An end-of-record condition occurs when a nonadvancing *input statement* attempts to transfer data from a position beyond the end of the current record, unless the file is a *stream file* and the current record is at the end of the file (an end-of-file condition occurs instead).
- 3 An end-of-file condition occurs when
  - an endfile record is encountered during the reading of a file connected for sequential access,
  - an attempt is made to read a record beyond the end of an *internal file*, or
  - an attempt is made to read beyond the end of a *stream file*.
- 4 An end-of-file condition may occur at the beginning of execution of an *input statement*. An end-of-file condition also may occur during execution of a formatted *input statement* when more than one record is required by the interaction of the input list and the format. An end-of-file condition also may occur during execution of a stream *input statement*.

### 9.11.2 Error conditions and the ERR= specifier

- 1 If an error condition occurs during execution of an input/output statement, the position of the file becomes indeterminate.
- 2 If an error condition occurs during execution of an input/output statement that contains neither an ERR= nor IOSTAT= specifier, *error termination* is initiated. If an error condition occurs during execution of an input/output statement that contains either an ERR= specifier or an IOSTAT= specifier then:
  - (1) processing of the input/output list, if any, terminates;
  - (2) if the statement is a *data transfer statement* or the error occurs during a wait operation, all *do-variables* in the statement that initiated the transfer become undefined;
  - (3) if an IOSTAT= specifier appears, the *scalar-int-variable* in the IOSTAT= specifier becomes defined as specified in 9.11.5;

- (4) if an **IOMSG= specifier** appears, the *iomsg-variable* becomes defined as specified in 9.11.6;
- (5) if the statement is a **READ statement** and it contains a **SIZE= specifier**, the *scalar-int-variable* in the **SIZE= specifier** becomes defined as specified in 9.6.2.15;
- (6) if the statement is a **READ statement** or the error condition occurs in a wait operation for a transfer initiated by a **READ statement**, all input items or namelist group objects in the statement that initiated the transfer become undefined;
- (7) if an **ERR= specifier** appears, a branch to the statement labeled by the *label* in the **ERR= specifier** occurs.

### 9.11.3 End-of-file condition and the **END= specifier**

- 1 If an end-of-file condition occurs during execution of an input/output statement that contains neither an **END= specifier** nor an **IOSTAT= specifier**, **error termination** is initiated. If an end-of-file condition occurs during execution of an input/output statement that contains either an **END= specifier** or an **IOSTAT= specifier**, and an error condition does not occur then:

- (1) processing of the input list, if any, terminates;
- (2) if the statement is a **data transfer statement** or the end-of-file condition occurs during a wait operation, all *do-variables* in the statement that initiated the transfer become undefined;
- (3) if the statement is an **input statement** or the end-of-file condition occurs during a wait operation for a transfer initiated by an **input statement**, all input list items or namelist group objects in the statement that initiated the transfer become undefined;
- (4) if the file specified in the **input statement** is an external record file, it is positioned after the endfile record;
- (5) if an **IOSTAT= specifier** appears, the *scalar-int-variable* in the **IOSTAT= specifier** becomes defined as specified in 9.11.5;
- (6) if an **IOMSG= specifier** appears, the *iomsg-variable* becomes defined as specified in 9.11.6;
- (7) if an **END= specifier** appears, a branch to the statement labeled by the *label* in the **END= specifier** occurs.

### 9.11.4 End-of-record condition and the **EOR= specifier**

- 1 If an end-of-record condition occurs during execution of an input/output statement that contains neither an **EOR= specifier** nor an **IOSTAT= specifier**, **error termination** is initiated. If an end-of-record condition occurs during execution of an input/output statement that contains either an **EOR= specifier** or an **IOSTAT= specifier**, and an error condition does not occur then:

- (1) if the pad mode has the value
  - (a) YES, the record is padded with blanks to satisfy the **effective item** (9.6.4.5.3) and corresponding data edit descriptors that require more characters than the record contains,
  - (b) NO, the input list item becomes undefined;
- (2) processing of the input list, if any, terminates;
- (3) if the statement is a **data transfer statement** or the end-of-record condition occurs during a wait operation, all *do-variables* in the statement that initiated the transfer become undefined;
- (4) the file specified in the input statement is positioned after the current record;
- (5) if an **IOSTAT= specifier** appears, the *scalar-int-variable* in the **IOSTAT= specifier** becomes defined as specified in 9.11.5;
- (6) if an **IOMSG= specifier** appears, the *iomsg-variable* becomes defined as specified in 9.11.6;
- (7) if a **SIZE= specifier** appears, the *scalar-int-variable* in the **SIZE= specifier** becomes defined as specified in (9.6.2.15);
- (8) if an **EOR= specifier** appears, a branch to the statement labeled by the *label* in the **EOR= specifier** occurs.

### 9.11.5 IOSTAT= specifier

1 Execution of an input/output statement containing the IOSTAT= specifier causes the *scalar-int-variable* in the IOSTAT= specifier to become defined with

- a zero value if neither an error condition, an end-of-file condition, nor an end-of-record condition occurs,
- the processor-dependent positive integer value of the constant `IOSTAT_INQUIRE_INTERNAL_UNIT` from the intrinsic module `ISO_FORTRAN_ENV`(13.10.2) if a `unit` number in an `INQUIRE` statement identifies an *internal file*,
- a processor-dependent positive integer value different from `IOSTAT_INQUIRE_INTERNAL_UNIT` if any other error condition occurs,
- the processor-dependent negative integer value of the constant `IOSTAT_END` (13.10.2.16) from the intrinsic module `ISO_FORTRAN_ENV` if an end-of-file condition occurs and no error condition occurs,
- the processor-dependent negative integer value of the constant `IOSTAT_EOR` (13.10.2.17) from the intrinsic module `ISO_FORTRAN_ENV` if an end-of-record condition occurs and no error condition or end-of-file condition occurs, or
- a processor-dependent negative integer value different from `IOSTAT_EOR` and `IOSTAT_END`, if the IO-STAT= specifier appears in a `FLUSH` statement and the processor does not support the flush operation for the specified unit.

#### NOTE 9.61

An end-of-file condition can occur only for sequential or stream input and an end-of-record condition can occur only for nonadvancing input.

For example,

```

READ (FMT = "(E8.3)", UNIT = 3, IOSTAT = IOSS) X
IF (IOSS < 0) THEN
    ! Perform end-of-file processing on the file connected to unit 3.
    CALL END_PROCESSING
ELSE IF (IOSS > 0) THEN
    ! Perform error processing
    CALL ERROR_PROCESSING
END IF

```

### 9.11.6 IOMSG= specifier

1 If an error, end-of-file, or end-of-record condition occurs during execution of an input/output statement, *iomsg-variable* is assigned an explanatory message as if by *intrinsic assignment*. If no such condition occurs, the definition status and value of *iomsg-variable* are unchanged.

## 9.12 Restrictions on input/output statements

- 1 If a *unit*, or a file connected to a *unit*, does not have all of the properties required for the execution of certain input/output statements, those statements shall not refer to the *unit*.
- 2 An input/output statement that is executed while another input/output statement is being executed is a recursive input/output statement. A recursive input/output statement shall not identify an *external unit* that is identified by another input/output statement being executed except that a child *data transfer statement* may identify its parent *data transfer statement external unit*.
- 3 An input/output statement shall not cause the value of any established format specification to be modified.
- 4 A recursive input/output statement shall not modify the value of any *internal unit* except that a recursive `WRITE` statement may modify the *internal unit* identified by that recursive `WRITE` statement.

- 1 5 The value of a specifier in an input/output statement shall not depend on the definition or evaluation of any other  
2 specifier in the *io-control-spec-list* or *inquire-spec-list* in that statement. The value of an *internal-file-variable* or  
3 of a **FMT=**, **ID=**, **IOMSG=**, **IOSTAT=**, or **SIZE=** specifier shall not depend on the value of any *input-item* or  
4 *io-implied-do do-variable* in the same statement.
- 5 6 The value of any subscript or substring bound of a variable that appears in a specifier in an input/output  
6 statement shall not depend on any *input-item*, *io-implied-do do-variable*, or on the definition or evaluation of any  
7 other specifier in the *io-control-spec-list* or *inquire-spec-list* in that statement.
- 8 7 In a *data transfer statement*, the variable specified in an **IOSTAT=**, **IOMSG=**, or **SIZE=** specifier, if any, shall  
9 not be associated with any entity in the data transfer input/output list (9.6.3) or *namelist-group-object-list*, nor  
10 with a *do-variable* of an *io-implied-do* in the data transfer input/output list.
- 11 8 In a *data transfer statement*, if a variable specified in an **IOSTAT=**, **IOMSG=**, or **SIZE=** specifier is an array  
12 element reference, its subscript values shall not be affected by the data transfer, the *io-implied-do* processing, or  
13 the definition or evaluation of any other specifier in the *io-control-spec-list*.
- 14 9 A variable that may become defined or undefined as a result of its use in a specifier in an **INQUIRE statement**,  
15 or any associated entity, shall not appear in another specifier in the same **INQUIRE statement**.

**NOTE 9.62**

Restrictions on the evaluation of expressions (7.1.4) prohibit certain side effects.
--



## 10 Input/output editing

### 10.1 Format specifications

- 1 A format used in conjunction with a [data transfer statement](#) provides information that directs the editing between the internal representation of data and the characters of a sequence of formatted records.
- 2 A *format* ([9.6.2.2](#)) in a [data transfer statement](#) may refer to a [FORMAT statement](#) or to a character expression that contains a format specification. A format specification provides explicit editing information. The *format* alternatively may be an asterisk (\*), which indicates list-directed formatting ([10.10](#)). Namelist formatting ([10.11](#)) may be indicated by specifying a *namelist-group-name* instead of a *format*.

### 10.2 Explicit format specification methods

#### 10.2.1 FORMAT statement

R1001 *format-stmt* is FORMAT *format-specification*

R1002 *format-specification* is ( [ *format-items* ] )  
or ( [ *format-items*, ] *unlimited-format-item* )

C1001 (R1001) The *format-stmt* shall be labeled.

- 1 Blank characters may precede the initial left parenthesis of the format specification. Additional blank characters may appear at any point within the format specification, with no effect on the interpretation of the format specification, except within a character string edit descriptor ([10.9](#)).

#### NOTE 10.1

Examples of FORMAT statements are:

```
5      FORMAT (1PE12.4, I10)
9      FORMAT (I12, /, ' Dates: ', 2 (2I3, I5))
```

#### 10.2.2 Character format specification

- 1 A character expression used as a *format* in a formatted input/output statement shall evaluate to a character string whose leading part is a valid format specification.

#### NOTE 10.2

The format specification begins with a left parenthesis and ends with a right parenthesis.

- 2 All character positions up to and including the final right parenthesis of the format specification shall be defined at the time the [data transfer statement](#) is executed, and shall not become redefined or undefined during the execution of the statement. Character positions, if any, following the right parenthesis that ends the format specification need not be defined and may contain any character data with no effect on the interpretation of the format specification.
- 3 If the *format* is a character array, it is treated as if all of the elements of the array were specified in array element order and were concatenated. However, if a *format* is a character array element, the format specification shall be entirely within that array element.

**NOTE 10.3**

If a character constant is used as a *format* in *data transfer statement*, care shall be taken that the value of the character constant is a valid format specification. In particular, if a format specification delimited by apostrophes contains a character constant edit descriptor delimited with apostrophes, two apostrophes shall be written to delimit the edit descriptor and four apostrophes shall be written for each apostrophe that occurs within the edit descriptor. For example, the text:

```
2 ISN'T 3
```

can be written by various combinations of *output statements* and format specifications:

```
WRITE (6, 100) 2, 3
100 FORMAT (1X, I1, 1X, 'ISN'T', 1X, I1)
WRITE (6, '(1X, I1, 1X, ''ISN''''T'', 1X, I1)') 2, 3
WRITE (6, '(A)') ' 2 ISN'T 3'
```

Doubling of internal apostrophes usually can be avoided by using quotation marks to delimit the format specification and doubling of internal quotation marks usually can be avoided by using apostrophes as delimiters.

## 10.3 Form of a format item list

### 10.3.1 Syntax

R1003 *format-items* is *format-item* [ [ , ] *format-item* ] ...

R1004 *format-item* is [ *r* ] *data-edit-desc*  
or *control-edit-desc*  
or *char-string-edit-desc*  
or [ *r* ] ( *format-items* )

R1005 *unlimited-format-item* is \* ( *format-items* )

R1006 *r* is *int-literal-constant*

C1002 (R1003) The optional comma shall not be omitted except

- between a P edit descriptor and an immediately following F, E, EN, ES, D, or G edit descriptor (10.8.5), possibly preceded by a repeat specification,
- before a slash edit descriptor when the optional repeat specification does not appear (10.8.2),
- after a slash edit descriptor, or
- before or after a colon edit descriptor (10.8.3)

C1003 (R1005) An *unlimited-format-item* shall contain at least one data edit descriptor.

C1004 (R1006) *r* shall be positive.

C1005 (R1006) A kind parameter shall not be specified for *r*.

1 The integer literal constant *r* is called a repeat specification.

### 10.3.2 Edit descriptors

1 An edit descriptor is a data edit descriptor (*data-edit-desc*), control edit descriptor (*control-edit-desc*), or character string edit descriptor (*char-string-edit-desc*).

1	R1007	<i>data-edit-desc</i>	<b>is</b> I <i>w</i> [ . <i>m</i> ]
2			<b>or</b> B <i>w</i> [ . <i>m</i> ]
3			<b>or</b> O <i>w</i> [ . <i>m</i> ]
4			<b>or</b> Z <i>w</i> [ . <i>m</i> ]
5			<b>or</b> F <i>w</i> . <i>d</i>
6			<b>or</b> E <i>w</i> . <i>d</i> [ E <i>e</i> ]
7			<b>or</b> EN <i>w</i> . <i>d</i> [ E <i>e</i> ]
8			<b>or</b> ES <i>w</i> . <i>d</i> [ E <i>e</i> ]
9			<b>or</b> EX <i>w</i> . <i>d</i> [ E <i>e</i> ]
10			<b>or</b> G <i>w</i> [ . <i>d</i> [ E <i>e</i> ] ]
11			<b>or</b> L <i>w</i>
12			<b>or</b> A [ <i>w</i> ]
13			<b>or</b> D <i>w</i> . <i>d</i>
14			<b>or</b> DT [ <i>char-literal-constant</i> ] [ ( <i>v-list</i> ) ]
15	R1008	<i>w</i>	<b>is</b> <i>int-literal-constant</i>
16	R1009	<i>m</i>	<b>is</b> <i>int-literal-constant</i>
17	R1010	<i>d</i>	<b>is</b> <i>int-literal-constant</i>
18	R1011	<i>e</i>	<b>is</b> <i>int-literal-constant</i>
19	R1012	<i>v</i>	<b>is</b> <i>signed-int-literal-constant</i>
20	C1006	(R1008) <i>w</i> shall be zero or positive for the I, B, O, Z, D, E, EN, ES, EX, F, and G edit descriptors. <i>w</i>	
21		shall be positive for all other edit descriptors.	
22	C1007	(R1007) For the G edit descriptor, <i>d</i> shall be specified if <i>w</i> is not zero.	
23	C1008	(R1007) For the G edit descriptor, <i>e</i> shall not be specified if <i>w</i> is zero.	
24	C1009	(R1007) A kind parameter shall not be specified for the <i>char-literal-constant</i> in the DT edit descriptor,	
25		or for <i>w</i> , <i>m</i> , <i>d</i> , <i>e</i> , and <i>v</i> .	
26	2	An I, B, O, Z, F, E, EN, ES, EX, G, L, A, D, or DT edit descriptor indicates the manner of editing.	
27	R1013	<i>control-edit-desc</i>	<b>is</b> <i>position-edit-desc</i>
28			<b>or</b> [ <i>r</i> ] /
29			<b>or</b> :
30			<b>or</b> <i>sign-edit-desc</i>
31			<b>or</b> <i>k</i> P
32			<b>or</b> <i>blank-interp-edit-desc</i>
33			<b>or</b> <i>round-edit-desc</i>
34			<b>or</b> <i>decimal-edit-desc</i>
35	R1014	<i>k</i>	<b>is</b> <i>signed-int-literal-constant</i>
36	C1010	(R1014) A kind parameter shall not be specified for <i>k</i> .	
37	3	In <i>k</i> P, <i>k</i> is called the scale factor.	
38	R1015	<i>position-edit-desc</i>	<b>is</b> T <i>n</i>
39			<b>or</b> TL <i>n</i>
40			<b>or</b> TR <i>n</i>
41			<b>or</b> <i>n</i> X
42	R1016	<i>n</i>	<b>is</b> <i>int-literal-constant</i>
43	C1011	(R1016) <i>n</i> shall be positive.	

1 C1012 (R1016) A kind parameter shall not be specified for *n*.

2 R1017 *sign-edit-desc* is SS  
3 or SP  
4 or S

5 R1018 *blank-interp-edit-desc* is BN  
6 or BZ

7 R1019 *round-edit-desc* is RU  
8 or RD  
9 or RZ  
10 or RN  
11 or RC  
12 or RP

13 R1020 *decimal-edit-desc* is DC  
14 or DP

15 4 A T, TL, TR, X, slash, colon, SS, SP, S, P, BN, BZ, RU, RD, RZ, RN, RC, RP, DC, or DP edit descriptor  
16 indicates the manner of editing.

17 R1021 *char-string-edit-desc* is *char-literal-constant*

18 C1013 (R1021) A kind parameter shall not be specified for the *char-literal-constant*.

19 5 Each *rep-char* in a character string edit descriptor shall be capable of representation by the processor.

20 6 A character string edit descriptor provides constant data to be output, and is not valid for input.

21 7 The edit descriptors are without regard to case except within a character string edit descriptor.

### 22 10.3.3 Fields

23 1 A field is a part of a record that is read on input or written on output when format control encounters a data  
24 edit descriptor or a character string edit descriptor. The field width is the size in characters of the field.

## 25 10.4 Interaction between input/output list and format

26 1 The start of formatted data transfer using a format specification initiates format control (9.6.4.5.3). Each action  
27 of format control depends on information jointly provided by the next edit descriptor in the format specification  
28 and the next *effective item* in the input/output list, if one exists.

29 2 If an input/output list specifies at least one *effective item*, at least one data edit descriptor shall exist in the  
30 format specification.

#### NOTE 10.4

An empty format specification of the form ( ) can be used only if the input/output list has no *effective item* (9.6.4.5). A zero length character item is an *effective item*, but a zero sized array and an implied DO list with an iteration count of zero is not.

31 3 A format specification is interpreted from left to right. The exceptions are format items preceded by a repeat  
32 specification *r*, and format reversion (described below).

33 4 A format item preceded by a repeat specification is processed as a list of *r* items, each identical to the format  
34 item but without the repeat specification and separated by commas.

**NOTE 10.5**

An omitted repeat specification is treated in the same way as a repeat specification whose value is one.

- 1    5 To each data edit descriptor interpreted in a format specification, there corresponds one [effective item](#) specified by  
 2    the input/output list ([9.6.3](#)), except that an input/output list item of type complex requires the interpretation of  
 3    two F, E, EN, ES, D, or G edit descriptors. For each control edit descriptor or character edit descriptor, there is  
 4    no corresponding item specified by the input/output list, and format control communicates information directly  
 5    with the record.
- 6    6 Whenever format control encounters a data edit descriptor in a format specification, it determines whether  
 7    there is a corresponding [effective item](#) specified by the input/output list. If there is such an item, it transmits  
 8    appropriately edited information between the item and the record, and then format control proceeds. If there is  
 9    no such item, format control terminates.
- 10   7 If format control encounters a colon edit descriptor in a format specification and another effective item is not  
 11   specified, format control terminates.
- 12   8 If format control encounters the rightmost parenthesis of an unlimited format item, control reverts to the leftmost  
 13   parenthesis of that unlimited format item. This reversion of format control has no effect on the changeable modes  
 14   ([9.5.2](#)).
- 15   9 If format control encounters the rightmost parenthesis of a complete format specification and another effective  
 16   item is not specified, format control terminates. However, if another effective item is specified, format control  
 17   then reverts to the beginning of the format item terminated by the last preceding right parenthesis that is not  
 18   part of a DT edit descriptor. If there is no such preceding right parenthesis, format control reverts to the first  
 19   left parenthesis of the format specification. If any reversion occurs, the reused portion of the format specification  
 20   shall contain at least one data edit descriptor. If format control reverts to a parenthesis that is preceded by a  
 21   repeat specification, the repeat specification is reused. Reversion of format control, of itself, has no effect on  
 22   the [changeable modes](#). The file is positioned in a manner identical to the way it is positioned when a slash edit  
 23   descriptor is processed ([10.8.2](#)).

**NOTE 10.6**

Example: The format specification:

```
10 FORMAT (1X, 2(F10.3, I5))
```

with an output list of

```
WRITE (10,10) 10.1, 3, 4.7, 1, 12.4, 5, 5.2, 6
```

produces the same output as the format specification:

```
10 FORMAT (1X, F10.3, I5, F10.3, I5/F10.3, I5, F10.3, I5)
```

**NOTE 10.7**

The effect of an [unlimited-format-item](#) is as if its enclosed list were preceded by a very large repeat count. There is no file positioning implied by [unlimited-format-item](#) reversion. This can be used to write what is commonly called a comma separated value record.

For example,

```
WRITE( 10, '( "IARRAY =", *( I0, :, ",") )' ) IARRAY
```

produces a single record with a header and a comma separated list of integer values.

## 10.5 Positioning by format control

- 1 After each data edit descriptor or character string edit descriptor is processed, the file is positioned after the last character read or written in the current record.
- 2 After each T, TL, TR, or X edit descriptor is processed, the file is positioned as described in 10.8.1.1. After each slash edit descriptor is processed, the file is positioned as described in 10.8.2.
- 3 During formatted stream output, processing of an A edit descriptor can cause file positioning to occur (10.7.4).
- 4 If format control reverts as described in 10.4, the file is positioned in a manner identical to the way it is positioned when a slash edit descriptor is processed (10.8.2).
- 5 During a read operation, any unprocessed characters of the current record are skipped whenever the next record is read.

## 10.6 Decimal symbol

- 1 The [decimal symbol](#) is the character that separates the whole and fractional parts in the decimal representation of a real number in an internal or [external](#) file. When the decimal edit mode is POINT, the [decimal symbol](#) is a decimal point. When the decimal edit mode is COMMA, the [decimal symbol](#) is a comma.
- 2 If the decimal edit mode is COMMA during list-directed input/output, the character used as a value separator is a semicolon in place of a comma.

## 10.7 Data edit descriptors

### 10.7.1 Purpose of data edit descriptors

- 1 Data edit descriptors cause the conversion of data to or from its internal representation; during formatted stream output, the A data edit descriptor may also cause file positioning. On input, the specified variable becomes defined unless an error condition, an end-of-file condition, or an end-of-record condition occurs. On output, the specified expression is evaluated.
- 2 During input from a Unicode file,
  - characters in the record that correspond to an [ASCII character](#) variable shall have a position in the [ISO 10646 character collating sequence](#) of 127 or less, and
  - characters in the record that correspond to a default character variable shall be representable as default characters.
- 3 During input from a non-Unicode file,
  - characters in the record that correspond to a character variable shall have the kind of the character variable, and
  - characters in the record that correspond to a numeric or logical variable shall be default characters.
- 4 During output to a Unicode file, all characters transmitted to the record are of [ISO 10646 character](#) kind. If a character input/output list item or character string edit descriptor contains a character that is not representable as an [ISO 10646 character](#), the result is processor dependent.
- 5 During output to a non-Unicode file, characters transmitted to the record as a result of processing a character string edit descriptor or as a result of evaluating a numeric, logical, or default character data entity, are of default kind.

## 10.7.2 Numeric editing

### 10.7.2.1 General rules

1 The I, B, O, Z, F, E, EN, ES, EX, D, and G edit descriptors may be used to specify the input/output of integer, real, and complex data. The following general rules apply.

- (1) On input, leading blanks are not significant. When the input field is not an IEEE exceptional specification or hexadecimal-significand number (10.7.2.3.2), the interpretation of blanks, other than leading blanks, is determined by the blank interpretation mode (10.8.6). Plus signs may be omitted. A field containing only blanks is considered to be zero.
- (2) On input, with F, E, EN, ES, EX, D, and G editing, a decimal symbol appearing in the input field overrides the portion of an edit descriptor that specifies the decimal symbol location. The input field may have more digits than the processor uses to approximate the value of the datum.
- (3) On output with I, F, E, EN, ES, EX, D, and G editing, the representation of a positive or zero internal value in the field may be prefixed with a plus sign, as controlled by the S, SP, and SS edit descriptors or the processor. The representation of a negative internal value in the field shall be prefixed with a minus sign.
- (4) On output, the representation is right justified in the field. If the number of characters produced by the editing is smaller than the field width, leading blanks are inserted in the field.
- (5) On output, if an exponent exceeds its specified or implied width using the E, EN, ES, EX, D, or G edit descriptor, or the number of characters produced exceeds the field width, the processor shall fill the entire field of width *w* with asterisks. However, the processor shall not produce asterisks if the field width is not exceeded when optional characters are omitted.

#### NOTE 10.8

When the sign mode is PLUS, a plus sign is not optional.

- (6) On output, with I, B, O, Z, D, E, EN, ES, EX, F, and G editing, the specified value of the field width *w* may be zero. In such cases, the processor selects the smallest positive actual field width that does not result in a field filled with asterisks. The specified value of *w* shall not be zero on input.
- (7) On output of a real zero value, the digits in the exponent field shall all be zero.

### 10.7.2.2 Integer editing

1 The *Iw* and *Iw.m* edit descriptors indicate that the field to be edited occupies *w* positions, except when *w* is zero. When *w* is zero, the processor selects the field width. On input, *w* shall not be zero. The specified input/output list item shall be of type integer. The G, B, O, and Z edit descriptor also may be used to edit integer data (10.7.5.2.1, 10.7.2.4).

2 On input, *m* has no effect.

3 In the standard form of the input field for the I edit descriptor, the character string is a *signed-digit-string* (R410), except for the interpretation of blanks. If the input field does not have the standard form and is not acceptable to the processor, an error condition occurs.

4 The output field for the *Iw* edit descriptor consists of zero or more leading blanks followed by a minus sign if the internal value is negative, or an optional plus sign otherwise, followed by the magnitude of the internal value as a *digit-string* without leading zeros.

#### NOTE 10.9

A *digit-string* always consists of at least one digit.

5 The output field for the *Iw.m* edit descriptor is the same as for the *Iw* edit descriptor, except that the *digit-string* consists of at least *m* digits. If necessary, sufficient leading zeros are included to achieve the minimum of *m* digits. The value of *m* shall not exceed the value of *w*, except when *w* is zero. If *m* is zero and the internal value is zero,

the output field consists of only blank characters, regardless of the sign control in effect. When *m* and *w* are both zero, and the internal value is zero, one blank character is produced.

### 10.7.2.3 Real and complex editing

#### 10.7.2.3.1 General

The F, E, EN, ES, and D edit descriptors specify the editing of real and complex data. An input/output list item corresponding to an F, E, EN, ES, or D edit descriptor shall be real or complex. The G, B, O, and Z edit descriptors also may be used to edit real and complex data (10.7.5.2.2, 10.7.2.4).

#### 10.7.2.3.2 F editing

The F*w.d* edit descriptor indicates that the field occupies *w* positions, except when *w* is zero in which case the processor selects the field width. The fractional part of the field consists of *d* digits. On input, *w* shall not be zero.

A lower-case letter is equivalent to the corresponding upper-case letter in an IEEE exceptional specification or the exponent in a numeric input field.

The standard form of the input field is an IEEE exceptional specification, a hexadecimal-significant number, or consists of a mantissa optionally followed by an exponent. The form of the mantissa is an optional sign, followed by a string of one or more digits optionally containing a decimal symbol, including any blanks interpreted as zeros. The *d* has no effect on input if the input field contains a decimal symbol. If the decimal symbol is omitted, the rightmost *d* digits of the string, with leading zeros assumed if necessary, are interpreted as the fractional part of the value represented. The string of digits may contain more digits than a processor uses to approximate the value. The form of the exponent is one of the following:

- a sign followed by a digit-string;
- the letter E followed by zero or more blanks, followed by a signed-digit-string;
- the letter D followed by zero or more blanks, followed by a signed-digit-string.

An exponent containing a D is processed identically to an exponent containing an E.

#### NOTE 10.10

If the input field does not contain an exponent, the effect is as if the basic form were followed by an exponent with a value of  $-k$ , where *k* is the established scale factor (10.8.5).

An input field that is an IEEE exceptional specification consists of optional blanks, followed by either

- an optional sign, followed by the string 'INF' or the string 'INFINITY', or
- an optional sign, followed by the string 'NAN', optionally followed by zero or more alphanumeric characters enclosed in parentheses,

optionally followed by blanks.

The value specified by 'INF' or 'INFINITY' is an IEEE infinity; this form shall not be used if the processor does not support IEEE infinities for the input variable. The value specified by 'NAN' is an IEEE NaN; this form shall not be used if the processor does not support IEEE NaNs for the input variable. The NaN value is a quiet NaN if the only nonblank characters in the field are 'NAN' or 'NAN()'; otherwise, the NaN value is processor dependent. The interpretation of a sign in a NaN input field is processor dependent.

An input field that is a hexadecimal-significant number consists of an optional sign, followed by the hexadecimal indicator which is the digit 0 immediately followed by the letter X, followed by a hexadecimal significant followed by a hexadecimal exponent. A hexadecimal significant is a string of one or more hexadecimal characters optionally containing a decimal symbol. The decimal symbol indicates the position of the hexadecimal point; if no decimal symbol appears, the hexadecimal point implicitly follows the last hexadecimal symbol. A hexadecimal exponent

- 1 is the letter P followed by a (decimal) *signed-digit-string*. Embedded blanks are not permitted in a hexadecimal-  
 2 significand number. The value is equal to the significand multiplied by two raised to the power of the exponent,  
 3 negated if the optional sign is minus.
- 4 8 If the input field does not have one of the standard forms, and is not acceptable to the processor, an error  
 5 condition occurs.
- 6 9 For an internal value that is an IEEE infinity, the output field consists of blanks, if necessary, followed by a minus  
 7 sign for negative infinity or an optional plus sign otherwise, followed by the letters 'Inf' or 'Infinity', right justified  
 8 within the field. The minimum field width required for output of the form 'Inf' is 3 if no sign is produced, and  
 9 4 otherwise. The minimum field width required for output of the form 'Infinity' is 8 if no sign is produced, and  
 10 9 otherwise. If *w* is greater than or equal to the minimum required for the form 'Infinity', the form 'Infinity' is  
 11 output. If *w* is zero or *w* is less than the minimum required for the form 'Infinity' and greater than or equal to  
 12 the minimum required for the form 'Inf', the form 'Inf' is output. Otherwise (*w* is greater than zero but less than  
 13 the minimum required for any form), the field is filled with asterisks.
- 14 10 For an internal value that is an IEEE NaN, the output field consists of blanks, if necessary, followed by the  
 15 letters 'NaN' and optionally followed by one to *w*−5 alphanumeric processor-dependent characters enclosed in  
 16 parentheses, right justified within the field. If *w* is greater than zero and less than 3, the field is filled with  
 17 asterisks. If *w* is zero, the output field is 'NaN'.

**NOTE 10.11**

The processor-dependent characters following 'NaN' might convey additional information about that particular NaN.

- 18 11 For an internal value that is neither an IEEE infinity nor a NaN, the output field consists of blanks, if necessary,  
 19 followed by a minus sign if the internal value is negative, or an optional plus sign otherwise, followed by a string  
 20 of digits that contains a *decimal symbol* and represents the magnitude of the internal value, as modified by the  
 21 established scale factor and rounded (10.7.2.3.8) to *d* fractional digits. Leading zeros are not permitted except  
 22 for an optional zero immediately to the left of the *decimal symbol* if the magnitude of the value in the output  
 23 field is less than one. The optional zero shall appear if there would otherwise be no digits in the output field.

**10.7.2.3.3 E and D editing**

- 25 1 The E*w.d*, D*w.d*, and E*w.d* E*e* edit descriptors indicate that the external field occupies *w* positions, except when  
 26 *w* is zero in which case the processor selects the field width. The fractional part of the field contains *d* digits,  
 27 unless a scale factor greater than one is in effect. If *e* is positive the exponent part contains *e* digits, otherwise it  
 28 contains the minimum number of digits required to represent the exponent value. The *e* has no effect on input.
- 29 2 The form and interpretation of the input field is the same as for F*w.d* editing (10.7.2.3.2).
- 30 3 For an internal value that is an IEEE infinity or NaN, the form of the output field is the same as for F*w.d*.
- 31 4 For an internal value that is neither an IEEE infinity nor a NaN, the form of the output field for a scale factor  
 32 of zero is  
 33 
$$[ \pm ] [0].x_1x_2 \dots x_d exp$$
  
 34 where:
- 35 •  $\pm$  signifies a plus sign or a minus sign;
  - 36 • . signifies a *decimal symbol* (10.6);
  - 37 •  $x_1x_2 \dots x_d$  are the *d* most significant digits of the internal value after rounding (10.7.2.3.8);
  - 38 • *exp* is a decimal exponent having one of the forms specified in table 10.1.

Table 10.1: **E and D exponent forms**

Edit Descriptor	Absolute Value of Exponent	Form of Exponent <sup>1</sup>
$Ew.d$	$ exp  \leq 99$	$E\pm z_1 z_2$ or $\pm 0 z_1 z_2$
	$99 <  exp  \leq 999$	$\pm z_1 z_2 z_3$
$Ew.d Ee$ with $e > 0$	$ exp  \leq 10^e - 1$	$E\pm z_1 z_2 \dots z_e$
$Ew.d E0$	any	$E\pm z_1 z_2 \dots z_s$
$Dw.d$	$ exp  \leq 99$	$D\pm z_1 z_2$ or $E\pm z_1 z_2$ or $\pm 0 z_1 z_2$
	$99 <  exp  \leq 999$	$\pm z_1 z_2 z_3$
(1) where each $z$ is a digit, and $s$ is the minimum number of digits required to represent the exponent.		

- 1    5    The sign in the exponent is produced. A plus sign is produced if the exponent value is zero.
- 2    6    The scale factor  $k$  controls the decimal normalization (10.3.2, 10.8.5). If  $-d < k \leq 0$ , the output field contains  
3    exactly  $|k|$  leading zeros and  $d - |k|$  significant digits after the decimal symbol. If  $0 < k < d + 2$ , the output field  
4    contains exactly  $k$  significant digits to the left of the decimal symbol and  $d - k + 1$  significant digits to the right  
5    of the decimal symbol. Other values of  $k$  are not permitted.
- 6    **10.7.2.3.4    EN editing**
- 7    1    The EN edit descriptor produces an output field in the form of a real number in engineering notation such that  
8    the decimal exponent is divisible by three and the absolute value of the significand (R415) is greater than or  
9    equal to 1 and less than 1000, except when the output value is zero. The scale factor has no effect on output.
- 10   2    The forms of the edit descriptor are  $ENw.d$  and  $ENw.d Ee$  indicating that the external field occupies  $w$  positions,  
11   except when  $w$  is zero in which case the processor selects the field width. The fractional part of the field contains  
12    $d$  digits. If  $e$  is positive the exponent part contains  $e$  digits, otherwise it contains the minimum number of digits  
13   required to represent the exponent value.
- 14   3    The form and interpretation of the input field is the same as for  $Fw.d$  editing (10.7.2.3.2).
- 15   4    For an internal value that is an IEEE infinity or NaN, the form of the output field is the same as for  $Fw.d$ .
- 16   5    For an internal value that is neither an IEEE infinity nor a NaN, the form of the output field is  
17    $[\pm] yyy . x_1 x_2 \dots x_d exp$   
18   where:
- 19   •  $\pm$  signifies a plus sign or a minus sign;
  - 20   •  $yyy$  are the 1 to 3 decimal digits representative of the most significant digits of the internal value after  
21   rounding (10.7.2.3.8);
  - 22   •  $yyy$  is an integer such that  $1 \leq yyy < 1000$  or, if the output value is zero,  $yyy = 0$ ;
  - 23   •  $.$  signifies a decimal symbol (10.6);
  - 24   •  $x_1 x_2 \dots x_d$  are the  $d$  next most significant digits of the internal value after rounding;
  - 25   •  $exp$  is a decimal exponent, divisible by three, having one of the forms specified in table 10.2.

Table 10.2: **EN exponent forms**

Edit Descriptor	Absolute Value of Exponent	Form of Exponent <sup>1</sup>
EN $w.d$	$ exp  \leq 99$	$E\pm z_1 z_2$ or $\pm 0 z_1 z_2$
	$99 <  exp  \leq 999$	$\pm z_1 z_2 z_3$
EN $w.d$ E $e$ with $e > 0$	$ exp  \leq 10^e - 1$	$E\pm z_1 z_2 \dots z_e$
EN $w.d$ E0	any	$E\pm z_1 z_2 \dots z_s$
(1) where each $z$ is a digit, and $s$ is the minimum number of digits required to represent the exponent.		

- 1    6    The sign in the exponent is produced. A plus sign is produced if the exponent value is zero.

**NOTE 10.12**

Examples:

Internal Value	Output field Using SS, EN12.3
6.421	6.421E+00
-.5	-500.000E-03
.00217	2.170E-03
4721.3	4.721E+03

2    **10.7.2.3.5 ES editing**

- 3    1    The ES edit descriptor produces an output field in the form of a real number in scientific notation such that the  
4    absolute value of the significand (R415) is greater than or equal to 1 and less than 10, except when the output  
5    value is zero. The scale factor has no effect on output.
- 6    2    The forms of the edit descriptor are ES $w.d$  and ES $w.d$  E $e$  indicating that the external field occupies  $w$  positions,  
7    except when  $w$  is zero in which case the processor selects the field width. The fractional part of the field contains  
8     $d$  digits. If  $e$  is positive the exponent part contains  $e$  digits, otherwise it contains the minimum number of digits  
9    required to represent the exponent value.
- 10    3    The form and interpretation of the input field is the same as for F $w.d$  editing (10.7.2.3.2).
- 11    4    For an internal value that is an IEEE infinity or NaN, the form of the output field is the same as for F $w.d$ .
- 12    5    For an internal value that is neither an IEEE infinity nor a NaN, the form of the output field is  
13     $[\pm] y . x_1 x_2 \dots x_d exp$   
14    where:
- 15    •  $\pm$  signifies a plus sign or a minus sign;
  - 16    •  $y$  is a decimal digit representative of the most significant digit of the internal value after rounding (10.7.2.3.8);
  - 17    •  $.$  signifies a decimal symbol (10.6);
  - 18    •  $x_1 x_2 \dots x_d$  are the  $d$  next most significant digits of the internal value after rounding;
  - 19    •  $exp$  is a decimal exponent having one of the forms specified in table 10.3.

Table 10.3: **ES exponent forms**

Edit Descriptor	Absolute Value of Exponent	Form of Exponent <sup>1</sup>
ES $w.d$	$ exp  \leq 99$	$E\pm z_1 z_2$ or $\pm 0 z_1 z_2$
	$99 <  exp  \leq 999$	$\pm z_1 z_2 z_3$
ES $w.d$ E $e$ with $e > 0$	$ exp  \leq 10^e - 1$	$E\pm z_1 z_2 \dots z_e$
ES $w.d$ E0	any	$E\pm z_1 z_2 \dots z_s$
(1) where each $z$ is a digit, and $s$ is the minimum number of digits required to represent the exponent.		

- 1 6 The sign in the exponent is produced. A plus sign is produced if the exponent value is zero.

#### NOTE 10.13

Examples:

Internal Value	Output field Using SS, ES12.3
6.421	6.421E+00
-.5	-5.000E-01
.00217	2.170E-03
4721.3	4.721E+03

#### 2 10.7.2.3.6 EX editing

- 3 1 The EX edit descriptor produces an output field in the form of a hexadecimal-significand number.
- 4 2 The EX $w.d$  and EX $w.dEe$  edit descriptors indicate that the external field occupies  $w$  positions, except when  $w$  is zero in which case the processor selects the field width. The fractional part of the field contains  $d$  hexadecimal digits, except when  $d$  is zero in which case the processor selects the number of hexadecimal digits to be the minimum required so that the output field is equal to the internal value;  $d$  shall not be zero if the radix of the internal value is not a power of two. The hexadecimal point, represented by a decimal symbol, appears after the first hexadecimal digit. For the form EX $w.d$ , and for EX $w.dE0$ , the exponent part contains the minimum number of digits needed to represent the exponent; otherwise the exponent contains  $e$  digits. The  $e$  has no effect on input. The scale factor has no effect on output.
- 12 3 The form and interpretation of the input field is the same as for F $w.d$  editing (10.7.2.3.2).
- 13 4 For an internal value that is an IEEE infinity or NaN, the form of the output field is the same as for F $w.d$ .
- 14 5 For an internal value that is neither an IEEE infinity nor a NaN, the form of the output field is
- 15  $[ \pm ] x_0 . x_1 x_2 \dots exp$
- 16 where:
- 17 •  $\pm$  signifies a plus sign or a minus sign;
  - 18 •  $.$  signifies a decimal symbol (10.6);
  - 19 •  $x_0 x_1 x_2 \dots$  are the most significant hexadecimal digits of the internal value, after rounding if  $d$  is not zero (10.7.2.3.8);
  - 21 •  $exp$  is a binary exponent expressed as a decimal integer; for EX $w.d$  and EX $w.dE0$ , the form is  $P \pm z_1 \dots z_n$ , where  $n$  is the minimum number of digits needed to represent  $exp$ , and for EX $w.dEe$  with  $e$  greater than zero the form is  $P \pm z_1 \dots z_e$ .
- 24 6 The sign in the exponent is produced. A plus sign is produced if the exponent value is zero.

#### NOTE 10.14

Examples:

Internal value	Edit descriptor	Possible output with SS in effect
1.375	EX0.1	0X1.6P+0
-15.625	EX14.4E3	-0X1.F400P+003
1048580.0	EX0.0	0X1.00003P+20

#### 25 10.7.2.3.7 Complex editing

- 26 1 A complex datum consists of a pair of separate real data. The editing of a scalar datum of complex type is specified by two edit descriptors each of which specifies the editing of real data. The first edit descriptor specifies the editing for the real part; the second specifies it for the imaginary part. The two edit descriptors may be different. Control and character string edit descriptors may be processed between the edit descriptor for the real part and the edit descriptor for the imaginary part.

### 10.7.2.3.8 Input/output rounding mode

- 1 The input/output rounding mode can be specified by an **OPEN statement** (9.5.2), a **data transfer statement** (9.6.2.13), or an edit descriptor (10.8.7).
- 2 In what follows, the term “decimal value” means the exact decimal number as given by the character string, while the term “internal value” means the number actually stored in the processor. For example, in dealing with the decimal constant 0.1, the decimal value is the mathematical quantity 1/10, which has no exact representation in binary form. Formatted output of real data involves conversion from an internal value to a decimal value; formatted input involves conversion from a decimal value to an internal value.
- 3 When the input/output rounding mode is UP, the value resulting from conversion shall be the smallest representable value that is greater than or equal to the original value. When the input/output rounding mode is DOWN, the value resulting from conversion shall be the largest representable value that is less than or equal to the original value. When the input/output rounding mode is ZERO, the value resulting from conversion shall be the value closest to the original value and no greater in magnitude than the original value. When the input/output rounding mode is NEAREST, the value resulting from conversion shall be the closer of the two nearest representable values if one is closer than the other. If the two nearest representable values are equidistant from the original value, it is processor dependent which one of them is chosen. When the input/output rounding mode is COMPATIBLE, the value resulting from conversion shall be the closer of the two nearest representable values or the value away from zero if halfway between them. When the input/output rounding mode is PROCESSOR\_DEFINED, rounding during conversion shall be a processor-dependent default mode, which may correspond to one of the other modes.
- 4 On processors that support IEEE rounding on conversions (14.4), NEAREST shall correspond to round to nearest, as specified in ISO/IEC/IEEE 60559:2011.

#### NOTE 10.15

On processors that support IEEE rounding on conversions, the input/output rounding modes COMPATIBLE and NEAREST will produce the same results except when the datum is halfway between the two representable values. In that case, NEAREST will pick the even value, but COMPATIBLE will pick the value away from zero. The input/output rounding modes UP, DOWN, and ZERO have the same effect as those specified in ISO/IEC/IEEE 60559:2011 for round toward  $+\infty$ , round toward  $-\infty$ , and round toward 0, respectively.

### 10.7.2.4 B, O, and Z editing

- 1 The **B $w$** , **B $w.m$** , **O $w$** , **O $w.m$** , **Z $w$** , and **Z $w.m$**  edit descriptors indicate that the field to be edited occupies  $w$  positions, except when  $w$  is zero. When  $w$  is zero, the processor selects the field width. On input,  $w$  shall not be zero. The corresponding input/output list item shall be of type integer, real, or complex.
- 2 On input,  $m$  has no effect.
- 3 In the standard form of the input field for the B, O, and Z edit descriptors the character string consists of binary, octal, or hexadecimal digits (as in R465, R466, R467) in the respective input field. The lower-case hexadecimal digits a through f in a hexadecimal input field are equivalent to the corresponding upper-case hexadecimal digits. If the input field does not have the standard form, and is not acceptable to the processor, an error condition occurs.
- 4 The value is INT (X) if the input list item is of type integer and REAL (X) if the input list item is of type real or complex, where X is a **boz-literal-constant** that specifies the same bit sequence as the digits of the input field.
- 5 The output field for the **B $w$** , **O $w$** , and **Z $w$**  descriptors consists of zero or more leading blanks followed by the internal value in a form identical to the digits of a binary, octal, or hexadecimal constant, respectively, that specifies the same bit sequence but without leading zero bits.

**NOTE 10.16**

A binary, octal, or hexadecimal constant always consists of at least one digit or hexadecimal digit.

1 R1022 *hex-digit-string* is *hex-digit* [ *hex-digit* ] ...

2 6 The output field for the B*w.m*, O*w.m*, and Z*w.m* edit descriptor is the same as for the B*w*, O*w*, and Z*w* edit  
3 descriptor, except that the *digit-string* or *hex-digit-string* consists of at least *m* digits. If necessary, sufficient  
4 leading zeros are included to achieve the minimum of *m* digits. The value of *m* shall not exceed the value of *w*,  
5 except when *w* is zero. If *m* is zero and the internal value consists of all zero bits, the output field consists of  
6 only blank characters. When *m* and *w* are both zero, and the internal value consists of all zero bits, one blank  
7 character is produced.

### 8 10.7.3 Logical editing

- 9 1 The L*w* edit descriptor indicates that the field occupies *w* positions. The specified input/output list item shall  
10 be of type logical. The G edit descriptor also may be used to edit logical data (10.7.5.3).
- 11 2 The standard form of the input field consists of optional blanks, optionally followed by a period, followed by a T  
12 for true or F for false. The T or F may be followed by additional characters in the field, which are ignored. If the  
13 input field does not have the standard form, and is not acceptable to the processor, an error condition occurs.
- 14 3 A lower-case letter is equivalent to the corresponding upper-case letter in a logical input field.

**NOTE 10.17**

The logical constants *.TRUE.* and *.FALSE.* are acceptable input forms.

- 15 4 The output field consists of *w*−1 blanks followed by a T or F, depending on whether the internal value is true or  
16 false, respectively.

### 17 10.7.4 Character editing

- 18 1 The A[*w*] edit descriptor is used with an input/output list item of type character. The G edit descriptor also may  
19 be used to edit character data (10.7.5.4). The kind type parameter of all characters transferred and converted  
20 under control of one A or G edit descriptor is implied by the kind of the corresponding list item.
- 21 2 If a field width *w* is specified with the A edit descriptor, the field consists of *w* characters. If a field width *w* is  
22 not specified with the A edit descriptor, the number of characters in the field is the length of the corresponding  
23 list item, regardless of the value of the kind type parameter.
- 24 3 Let *len* be the length of the input/output list item. If the specified field width *w* for an A edit descriptor  
25 corresponding to an input item is greater than or equal to *len*, the rightmost *len* characters will be taken from the  
26 input field. If the specified field width *w* is less than *len*, the *w* characters will appear left justified with *len*−*w*  
27 trailing blanks in the internal value.
- 28 4 If the specified field width *w* for an A edit descriptor corresponding to an output item is greater than *len*, the  
29 output field will consist of *w*−*len* blanks followed by the *len* characters from the internal value. If the specified  
30 field width *w* is less than or equal to *len*, the output field will consist of the leftmost *w* characters from the internal  
31 value.

**NOTE 10.18**

For nondefault character kinds, the blank padding character is processor dependent.

- 32 5 If the file is connected for stream access, the output may be split across more than one record if it contains  
33 newline characters. A newline character is a nonblank character returned by the intrinsic function *NEW\_LINE*.  
34 Beginning with the first character of the output field, each character that is not a newline is written to the current  
35 record in successive positions; each newline character causes file positioning at that point as if by slash editing

(the current record is terminated at that point, a new empty record is created following the current record, this new record becomes the last and current record of the file, and the file is positioned at the beginning of this new record).

#### NOTE 10.19

If the intrinsic function `NEW_LINE` returns a blank character for a particular character kind, then the processor does not support using a character of that kind to cause record termination in a formatted stream file.

## 10.7.5 Generalized editing

### 10.7.5.1 Overview

The `Gw`, `Gw.d` and `Gw.d Ee` edit descriptors are used with an input/output list item of any intrinsic type. When `w` is nonzero, these edit descriptors indicate that the external field occupies `w` positions. For real or complex data the fractional part consists of a maximum of `d` digits and the exponent part consists of `e` digits. When these edit descriptors are used to specify the input/output of integer, logical, or character data, `d` and `e` have no effect. When `w` is zero the processor selects the field width. On input, `w` shall not be zero.

### 10.7.5.2 Generalized numeric editing

When used to specify the input/output of integer, real, and complex data, the `Gw`, `Gw.d` and `Gw.d Ee` edit descriptors follow the general rules for numeric editing (10.7.2).

#### NOTE 10.20

The `Gw.d Ee` edit descriptor follows any additional rules for the `Ew.d Ee` edit descriptor.

#### 10.7.5.2.1 Generalized integer editing

When used to specify the input/output of integer data, the `Gw.d` and `Gw.d Ee` edit descriptors follow the rules for the `Iw` edit descriptor (10.7.2.2), except that `w` shall not be zero. When used to specify the output of integer data, the `G0` and `G0.d` edit descriptors follow the rules for the `I0` edit descriptor.

#### 10.7.5.2.2 Generalized real and complex editing

The form and interpretation of the input field is the same as for `Fw.d` editing (10.7.2.3.2).

If `d` is zero, `kPEw.0` or `kPEw.0Ee` editing is used for `Gw.0` editing or `Gw.0Ee` editing respectively.

When used to specify the output of real or complex data that is not an IEEE infinity or NaN, the `G0` and `G0.d` edit descriptors follow the rules for the `Gw.dEe` edit descriptor, except that any leading or trailing blanks are removed. Reasonable processor-dependent values of `w`, `d` (if not specified), and `e` are used with each output value.

For an internal value that is an IEEE infinity or NaN, the form of the output field for the `Gw.d` and `Gw.d Ee` edit descriptors is the same as for `Fw.d`, and the form of the output field for the `G0` and `G0.d` edit descriptors is the same as for `F0.0`.

Otherwise, the method of representation in the output field depends on the magnitude of the internal value being edited. If the internal value is zero, let  $s$  be one. If the internal value is a number other than zero, let  $N$  be the decimal value that is the result of converting the internal value to `d` significant digits according to the input/output rounding mode and let  $s$  be the integer such that  $10^{s-1} \leq |N| < 10^s$ . If  $s < 0$  or  $s > d$ , `kPEw.d` or `kPEw.dEe` editing is used for `Gw.d` editing or `Gw.dEe` editing respectively, where  $k$  is the scale factor (10.8.5). If  $0 \leq s \leq d$ , the scale factor has no effect and `F(w-n).(d-s),n('b')` editing is used where  $b$  is a blank and  $n$  is 4 for `Gw.d` editing,  $e+2$  for `Gw.dEe` editing if  $e > 0$ , and 4 for `Gw.dE0` editing.

The value of  $w-n$  shall be positive.

**NOTE 10.21**

The scale factor has no effect on output unless the magnitude of the datum to be edited is outside the range that permits effective use of F editing.

**10.7.5.3 Generalized logical editing**

- 1 When used to specify the input/output of logical data, the G*w.d* and G*w.d Ee* edit descriptors follow the rules for the L*w* edit descriptor (10.7.3). When used to specify the output of logical data, the G0 and G0.*d* edit descriptors follow the rules for the L1 edit descriptor.

**10.7.5.4 Generalized character editing**

- 1 When used to specify the input/output of character data, the G*w.d* and G*w.d Ee* edit descriptors follow the rules for the A*w* edit descriptor (10.7.4). When used to specify the output of character data, the G0 and G0.*d* edit descriptors follow the rules for the A edit descriptor with no field width.

**10.7.6 User-defined derived-type editing**

- 1 The DT edit descriptor specifies that a user-provided procedure shall be used instead of the processor's default input/output formatting for processing a list item of derived type.
- 2 The DT edit descriptor may include a character literal constant. The character value "DT" concatenated with the character literal constant is passed to the [defined input/output](#) procedure as the `iotype` argument (9.6.4.8). The *v* values of the edit descriptor are passed to the [defined input/output](#) procedure as the `v_list` array argument.

**NOTE 10.22**

For the edit descriptor DT'Link List'(10, 4, 2), `iotype` is "DTLink List" and `v_list` is [10, 4, 2].

- 3 If a derived-type variable or value corresponds to a DT edit descriptor, there shall be an accessible [interface](#) to a corresponding [defined input/output](#) procedure for that derived type (9.6.4.8). A DT edit descriptor shall not correspond to a list item that is not of a derived type.

**10.8 Control edit descriptors****10.8.1 Position edit descriptors****10.8.1.1 Position editing**

- 1 The position edit descriptors T, TL, TR, and X, specify the position at which the next character will be transmitted to or from the record. If any character skipped by a position edit descriptor is of type nondefault character, and the `unit` is a default character [internal file](#) or an [external](#) non-Unicode file, the result of that position editing is processor dependent.
- 2 On input, if the position specified by a position edit descriptor is before the current position, portions of a record can be processed more than once, possibly with different editing.
- 3 On input, a position beyond the last character of the record may be specified if no characters are transmitted from such positions.
- 4 On output, a position edit descriptor does not by itself cause characters to be transmitted and therefore does not by itself affect the length of the record. If characters are transmitted to positions at or after the position specified by a position edit descriptor, positions skipped and not previously filled are filled with blanks. The result is as if the entire record were initially filled with blanks.
- 5 On output, a character in the record can be replaced. A position edit descriptor never directly causes a character

already placed in the record to be replaced, but it might result in positioning such that subsequent editing causes a replacement.

### 10.8.1.2 T, TL, and TR editing

1 The left tab limit affects file positioning by the T and TL edit descriptors. Immediately prior to nonchild data transfer (9.6.4.8.2), the left tab limit becomes defined as the character position of the current record or the current position of the stream file. If, during data transfer, the file is positioned to another record, the left tab limit becomes defined as character position one of that record.

2 The T $n$  edit descriptor indicates that the transmission of the next character to or from a record is to occur at the  $n$ th character position of the record, relative to the left tab limit. This position can be in either direction from the current position.

3 The TL $n$  edit descriptor indicates that the transmission of the next character to or from the record is to occur at the character position  $n$  characters backward from the current position. However, if  $n$  is greater than the difference between the current position and the left tab limit, the TL $n$  edit descriptor indicates that the transmission of the next character to or from the record is to occur at the left tab limit.

4 The TR $n$  edit descriptor indicates that the transmission of the next character to or from the record is to occur at the character position  $n$  characters forward from the current position.

### 10.8.1.3 X editing

1 The  $n$ X edit descriptor indicates that the transmission of the next character to or from a record is to occur at the character position  $n$  characters forward from the current position.

#### NOTE 10.23

An  $n$ X edit descriptor has the same effect as a TR $n$  edit descriptor.

## 10.8.2 Slash editing

1 The slash edit descriptor indicates the end of data transfer to or from the current record.

2 On input from a file connected for sequential or stream access, the remaining portion of the current record is skipped and the file is positioned at the beginning of the next record. This record becomes the current record. On output to a file connected for sequential or stream access, a new empty record is created following the current record; this new record then becomes the last and current record of the file and the file is positioned at the beginning of this new record.

3 For a file connected for direct access, the record number is increased by one and the file is positioned at the beginning of the record that has that record number, if there is such a record, and this record becomes the current record.

#### NOTE 10.24

A record that contains no characters can be written on output. If the file is an [internal file](#) or a file connected for direct access, the record is filled with blank characters.

An entire record can be skipped on input.

4 The repeat specification is optional in the slash edit descriptor. If it is not specified, the default value is one.

## 10.8.3 Colon editing

1 The colon edit descriptor terminates format control if there are no more [effective items](#) in the input/output list (9.6.3). The colon edit descriptor has no effect if there are more [effective items](#) in the input/output list.

## 10.8.4 SS, SP, and S editing

- 1 The SS, SP, and S edit descriptors temporarily change (9.5.2) the sign mode (9.5.6.17, 9.6.2.14) for the connection. The edit descriptors SS, SP, and S set the sign mode corresponding to the SIGN= specifier values SUPPRESS, PLUS, and PROCESSOR\_DEFINED, respectively.
- 2 The sign mode controls optional plus characters in numeric output fields. When the sign mode is PLUS, the processor shall produce a plus sign in any position that normally contains an optional plus sign. When the sign mode is SUPPRESS, the processor shall not produce a plus sign in such positions. When the sign mode is PROCESSOR\_DEFINED, the processor has the option of producing a plus sign or not in such positions, subject to 10.7.2(5).
- 3 The SS, SP, and S edit descriptors affect only I, F, E, EN, ES, D, and G editing during the execution of an output statement. The SS, SP, and S edit descriptors have no effect during the execution of an input statement.

## 10.8.5 P editing

- 1 The *k*P edit descriptor temporarily changes (9.5.2) the scale factor for the connection to *k*. The scale factor affects the editing done by the F, E, EN, ES, D, and G edit descriptors for numeric quantities.
- 2 The scale factor *k* affects the appropriate editing in the following manner.
  - On input, with F, E, EN, ES, D, and G editing (provided that no exponent exists in the field), the effect is that the externally represented number equals the internally represented number multiplied by  $10^k$ ; the scale factor is applied to the external decimal value and then this is converted using the input/output rounding mode.
  - On input, with F, E, EN, ES, D, and G editing, the scale factor has no effect if there is an exponent in the field.
  - On output, with F output editing, the effect is that the externally represented number equals the internally represented number multiplied by  $10^k$ ; the internal value is converted using the input/output rounding mode and then the scale factor is applied to the converted decimal value.
  - On output, with E and D editing, the effect is that the significand (R415) part of the quantity to be produced is multiplied by  $10^k$  and the exponent is reduced by *k*.
  - On output, with G editing, the effect is suspended unless the magnitude of the datum to be edited is outside the range that permits the use of F editing. If the use of E editing is required, the scale factor has the same effect as with E output editing.
  - On output, with EN and ES editing, the scale factor has no effect.

## 10.8.6 BN and BZ editing

- 1 The BN and BZ edit descriptors temporarily change (9.5.2) the blank interpretation mode (9.5.6.6, 9.6.2.6) for the connection. The edit descriptors BN and BZ set the blank interpretation mode corresponding to the BLANK= specifier values NULL and ZERO, respectively.
- 2 The blank interpretation mode controls the interpretation of nonleading blanks in numeric input fields. Such blank characters are interpreted as zeros when the blank interpretation mode has the value ZERO; they are ignored when the blank interpretation mode has the value NULL. The effect of ignoring blanks is to treat the input field as if blanks had been removed, the remaining portion of the field right justified, and the blanks replaced as leading blanks. However, a field containing only blanks has the value zero.
- 3 The blank interpretation mode affects only numeric editing (10.7.2) and generalized numeric editing (10.7.5.2) on input. It has no effect on output.

## 10.8.7 RU, RD, RZ, RN, RC, and RP editing

- 1 The round edit descriptors temporarily change (9.5.2) the connection's input/output rounding mode (9.5.6.16, 9.6.2.13, 10.7.2.3.8). The round edit descriptors RU, RD, RZ, RN, RC, and RP set the input/output rounding

mode corresponding to the `ROUND= specifier` values UP, DOWN, ZERO, NEAREST, COMPATIBLE, and PROCESSOR\_DEFINED, respectively. The input/output rounding mode affects the conversion of real and complex values in formatted input/output. It affects only D, E, EN, ES, F, and G editing.

### 10.8.8 DC and DP editing

- 1 The decimal edit descriptors temporarily change (9.5.2) the decimal edit mode (9.5.6.7, 9.6.2.7, 10.6) for the connection. The edit descriptors DC and DP set the decimal edit mode corresponding to the `DECIMAL= specifier` values COMMA and POINT, respectively.
- 2 The decimal edit mode controls the representation of the decimal symbol (10.6) during conversion of real and complex values in formatted input/output. The decimal edit mode affects only D, E, EN, ES, F, and G editing.

## 10.9 Character string edit descriptors

- 1 A character string edit descriptor shall not be used on input.
- 2 The character string edit descriptor causes characters to be written from the enclosed characters of the edit descriptor itself, including blanks. For a character string edit descriptor, the width of the field is the number of characters between the delimiting characters. Within the field, two consecutive delimiting characters are counted as a single character.

#### NOTE 10.25

A delimiter for a character string edit descriptor is either an apostrophe or quote.

## 10.10 List-directed formatting

### 10.10.1 Purpose of list-directed formatting

- 1 List-directed input/output allows data editing according to the type of the list item instead of by a format specification. It also allows data to be free-field, that is, separated by commas (or semicolons) or blanks.

### 10.10.2 Values and value separators

- 1 The characters in one or more list-directed records constitute a sequence of values and value separators. The end of a record has the same effect as a blank character, unless it is within a character constant. Any sequence of two or more consecutive blanks is treated as a single blank, unless it is within a character constant.
- 2 Each value is either a null value,  $c$ ,  $r^*c$ , or  $r^*$ , where  $c$  is a literal constant, optionally signed if integer or real, or an undelimited character constant and  $r$  is an unsigned, nonzero, integer literal constant. Neither  $c$  nor  $r$  shall have kind type parameters specified. The constant  $c$  is interpreted as though it had the same kind type parameter as the corresponding list item. The  $r^*c$  form is equivalent to  $r$  successive appearances of the constant  $c$ , and the  $r^*$  form is equivalent to  $r$  successive appearances of the null value. Neither of these forms may contain embedded blanks, except where permitted within the constant  $c$ .
- 3 A value separator is
  - a comma optionally preceded by one or more contiguous blanks and optionally followed by one or more contiguous blanks, unless the decimal edit mode is COMMA, in which case a semicolon is used in place of the comma,
  - a slash optionally preceded by one or more contiguous blanks and optionally followed by one or more contiguous blanks, or
  - one or more contiguous blanks between two nonblank values or following the last nonblank value, where a nonblank value is a constant, an  $r^*c$  form, or an  $r^*$  form.

**NOTE 10.26**

Although a slash encountered in an input record is referred to as a separator, it actually causes termination of list-directed and namelist [input statements](#); it does not actually separate two values.

**NOTE 10.27**

If no list items are specified in a list-directed input/output statement, one input record is skipped or one empty output record is written.

**10.10.3 List-directed input****10.10.3.1 List-directed input forms**

- 1 Input forms acceptable to edit descriptors for a given type are acceptable for list-directed formatting, except as noted below. If the form of the input value is not acceptable to the processor for the type of the next [effective item](#) in the list, an error condition occurs. Blanks are never used as zeros, and embedded blanks are not permitted in constants, except within character constants and complex constants as specified below.
- 2 For the  $r^*c$  form of an input value, the constant  $c$  is interpreted as an unlimited character constant if the first list item corresponding to this value is default, [ASCII](#), or [ISO 10646 character](#), there is a nonblank character immediately after  $r^*$ , and that character is not an apostrophe or a quotation mark; otherwise,  $c$  is interpreted as a literal constant.

**NOTE 10.28**

The end of a record has the effect of a blank, except when it appears within a character constant.

- 3 When the next [effective item](#) is of type integer, the value in the input record is interpreted as if an *Iw* edit descriptor with a suitable value of *w* were used.
- 4 When the next [effective item](#) is of type real, the input form is that of a numeric input field. A numeric input field is a field suitable for F editing ([10.7.2.3.2](#)) that is assumed to have no fractional digits unless a [decimal symbol](#) appears within the field.
- 5 When the next [effective item](#) is of type complex, the input form consists of a left parenthesis followed by an ordered pair of numeric input fields separated by a comma (if the decimal edit mode is POINT) or semicolon (if the decimal edit mode is COMMA), and followed by a right parenthesis. The first numeric input field is the real part of the complex constant and the second is the imaginary part. Each of the numeric input fields may be preceded or followed by any number of blanks and ends of records. The end of a record may occur after the real part or before the imaginary part.
- 6 When the next [effective item](#) is of type logical, the input form shall not include value separators among the optional characters permitted for L editing.
- 7 When the next [effective item](#) is of type character, the input form consists of a possibly delimited sequence of zero or more *rep-chars* whose kind type parameter is implied by the kind of the [effective item](#). Character sequences may be continued from the end of one record to the beginning of the next record, but the end of record shall not occur between a doubled apostrophe in an apostrophe-delimited character sequence, nor between a doubled quote in a quote-delimited character sequence. The end of the record does not cause a blank or any other character to become part of the character sequence. The character sequence may be continued on as many records as needed. The characters blank, comma, semicolon, and slash may appear in default, [ASCII](#), or [ISO 10646 character](#) sequences.
- 8 If the next [effective item](#) is default, [ASCII](#), or [ISO 10646 character](#) and
  - the character sequence does not contain value separators,
  - the character sequence does not cross a record boundary,
  - the first nonblank character is not a quotation mark or an apostrophe,

- the leading characters are not *digits* followed by an asterisk, and
- the character sequence contains at least one character,

the delimiting apostrophes or quotation marks are not required. If the delimiters are omitted, the character sequence is terminated by the first blank, comma (if the decimal edit mode is POINT), semicolon (if the decimal edit mode is COMMA), slash, or end of record; in this case apostrophes and quotation marks within the datum are not to be doubled.

- 9 Let *len* be the length of the next [effective item](#), and let *w* be the length of the character sequence. If *len* is less than or equal to *w*, the leftmost *len* characters of the sequence are transmitted to the next [effective item](#). If *len* is greater than *w*, the sequence is transmitted to the leftmost *w* characters of the next [effective item](#) and the remaining *len*−*w* characters of the next [effective item](#) are filled with blanks. The effect is as though the sequence were assigned to the next [effective item](#) in an [intrinsic assignment statement](#) (7.2.1.3).

### 10.10.3.2 Null values

- 1 A null value is specified by

- the *r*\* form,
- no characters between consecutive value separators, or
- no characters before the first value separator in the first record read by each execution of a list-directed [input statement](#).

#### NOTE 10.29

The end of a record following any other value separator, with or without separating blanks, does not specify a null value in list-directed input.

- 2 A null value has no effect on the definition status of the next [effective item](#). A null value shall not be used for either the real or imaginary part of a complex constant, but a single null value may represent an entire complex constant.
- 3 A slash encountered as a value separator during execution of a list-directed [input statement](#) causes termination of execution of that [input statement](#) after the transference of the previous value. Any characters remaining in the current record are ignored. If there are additional items in the input list, the effect is as if null values had been supplied for them. Any *do-variable* in the input list becomes defined as if enough null values had been supplied for any remaining input list items.

#### NOTE 10.30

All blanks in a list-directed input record are considered to be part of some value separator except for

- blanks embedded in a character sequence,
- embedded blanks surrounding the real or imaginary part of a complex constant, and
- leading blanks in the first record read by each execution of a list-directed [input statement](#), unless immediately followed by a slash or comma.

#### NOTE 10.31

List-directed input example:

```
INTEGER I; REAL X (8); CHARACTER (11) P;
COMPLEX Z; LOGICAL G
...
READ *, I, X, P, Z, G
...
```

The input data records are:

**NOTE 10.31 (cont.)**

```
12345,12345,,2*1.5,4*
ISN'T_BOB'S,(123,0),.TEXAS$
```

The results are:

Variable	Value
I	12345
X (1)	12345.0
X (2)	unchanged
X (3)	1.5
X (4)	1.5
X (5) – X (8)	unchanged
P	ISN'T_BOB'S
Z	(123.0,0.0)
G	true

**10.10.4 List-directed output**

1 The form of the values produced is the same as that required for input, except as noted otherwise. With the exception of adjacent undelimited character sequences, the values are separated by one or more blanks or by a comma, or a semicolon if the decimal edit mode is COMMA, optionally preceded by one or more blanks and optionally followed by one or more blanks. Two undelimited character sequences are considered adjacent when both were written using list-directed input/output, no intervening [data transfer](#) or [file positioning](#) operations on that [unit](#) occurred, and both were written either by a single [data transfer statement](#), or during the execution of a parent [data transfer statement](#) along with its child [data transfer statements](#). The form of the values produced by [defined output \(9.6.4.8\)](#) is determined by the [defined output](#) procedure; this form need not be compatible with list-directed input.

2 The processor may begin new records as necessary, but the end of record shall not occur within a constant except as specified for complex constants and character sequences. The processor shall not insert blanks within character sequences or within constants, except as specified for complex constants.

3 Logical output values are T for the value true and F for the value false.

4 Integer output constants are produced with the effect of an [Iw](#) edit descriptor.

5 Real constants are produced with the effect of either an F edit descriptor or an E edit descriptor, depending on the magnitude  $x$  of the value and a range  $10^{d_1} \leq x < 10^{d_2}$ , where  $d_1$  and  $d_2$  are processor-dependent integers. If the magnitude  $x$  is within this range or is zero, the constant is produced using [0PFw.d](#); otherwise, [1PEw.d Ee](#) is used.

6 For numeric output, reasonable processor-dependent values of [w](#), [d](#), and [e](#) are used for each of the numeric constants output.

7 Complex constants are enclosed in parentheses with a separator between the real and imaginary parts, each produced as defined above for real constants. The separator is a comma if the decimal edit mode is POINT; it is a semicolon if the decimal edit mode is COMMA. The end of a record may occur between the separator and the imaginary part only if the entire constant is as long as, or longer than, an entire record. The only embedded blanks permitted within a complex constant are between the separator and the end of a record and one blank at the beginning of the next record.

8 Character sequences produced when the delimiter mode has a value of NONE

- are not delimited by apostrophes or quotation marks,
- are not separated from each other by value separators,

- have each internal apostrophe or quotation mark represented externally by one apostrophe or quotation mark, and
- have a blank character inserted by the processor at the beginning of any record that begins with the continuation of a character sequence from the preceding record.

Character sequences produced when the delimiter mode has a value of QUOTE are delimited by quotes, are preceded and followed by a value separator, and have each internal quote represented on the external medium by two contiguous quotes.

Character sequences produced when the delimiter mode has a value of APOSTROPHE are delimited by apostrophes, are preceded and followed by a value separator, and have each internal apostrophe represented on the external medium by two contiguous apostrophes.

If two or more successive values in an output record have identical values, the processor has the option of producing a repeated constant of the form  $r*c$  instead of the sequence of identical values.

Slashes, as value separators, and null values are not produced as output by list-directed formatting.

Except for new records created by explicit formatting within a [defined output](#) procedure or by continuation of delimited character sequences, each output record begins with a blank character.

#### NOTE 10.32

The length of the output records is not specified and is processor dependent.

## 10.11 Namelist formatting

### 10.11.1 Purpose of namelist formatting

Namelist input/output allows data editing with name-value subsequences. This facilitates documentation of input and output files and more flexibility on input.

### 10.11.2 Name-value subsequences

The characters in one or more namelist records constitute a sequence of name-value subsequences, each of which consists of an [object designator](#) followed by an equals and followed by one or more values and value separators. The equals may optionally be preceded or followed by one or more contiguous blanks. The end of a record has the same effect as a blank character, unless it is within a character constant. Any sequence of two or more consecutive blanks is treated as a single blank, unless it is within a character constant.

Each object designator shall begin with a name from the [namelist-group-object-list](#) (5.8) and shall follow the syntax of [designator](#) (R601). It shall not contain a [vector subscript](#) or an [image-selector](#) and shall not designate a zero-sized array, a zero-sized array section, or a zero-length character string. Each subscript, stride, and substring range expression shall be an optionally signed integer literal constant with no [kind type parameter](#) specified. If a section subscript list appears, the number of section subscripts shall be equal to the rank of the object. If the namelist group object is of derived type, the designator in the input record may be either the name of the variable or the designator of one of its components, indicated by qualifying the variable name with the appropriate component name. Successive qualifications may be applied as appropriate to the shape and type of the variable represented. Each designator may be preceded and followed by one or more optional blanks but shall not contain embedded blanks.

A value separator for namelist formatting is the same as for list-directed formatting (10.10.2), or one or more contiguous blanks between a nonblank value and the following [object designator](#) or namelist comment (10.11.3.6).

### 10.11.3 Namelist input

#### 10.11.3.1 Overall syntax

1 Input for a namelist **input statement** consists of

- (1) optional blanks and namelist comments,
- (2) the character & followed immediately by the *namelist-group-name* as specified in the **NAMELIST statement**,
- (3) one or more blanks,
- (4) a sequence of zero or more name-value subsequences separated by value separators, and
- (5) a slash to terminate the namelist input.

#### NOTE 10.33

A slash encountered in a namelist input record causes the **input statement** to terminate. A slash cannot be used to separate two values in a namelist **input statement**.

2 The order of the name-value subsequences in the input records need not match the order of the *namelist-group-object-list*. The input records need not specify all objects in the *namelist-group-object-list*. They may specify a part of an object more than once.

3 A group name or object name is without regard to case.

#### 10.11.3.2 Namelist input processing

1 The name-value subsequences are evaluated serially, in left-to-right order. A namelist group object **designator** may appear in more than one name-value subsequence. The definition status of an object that is not a subobject of a **designator** in any name-value subsequence remains unchanged.

2 When the **designator** in the input record represents an array variable or a variable of derived type, the effect is as if the variable represented were expanded into a sequence of scalar list items, in the same way that formatted input/output list items are expanded (9.6.3). The number of values following the equals shall not exceed the number of list items in the expanded sequence, but may be less; in the latter case, the effect is as if sufficient null values had been appended to match any remaining list items in the expanded sequence. Except as noted elsewhere in this subclause, if an input value is not acceptable to the processor for the type of the list item in the corresponding position in the expanded sequence, an error condition occurs.

#### NOTE 10.34

For example, if the designator in the input record designates an integer array of size 100, at most 100 values, each of which is either a digit string or a null value, can follow the equals; these values would then be assigned to the elements of the array in array element order.

3 A slash encountered as a value separator during the execution of a namelist **input statement** causes termination of execution of that **input statement** after transference of the previous value. If there are additional items in the namelist group object being transferred, the effect is as if null values had been supplied for them.

4 Successive namelist records are read by namelist input until a slash is encountered; the remainder of the record is ignored.

5 A namelist comment may appear after any value separator except a slash (which terminates namelist input). A namelist comment is also permitted to start in the first nonblank position of an input record except within a character literal constant.

#### 10.11.3.3 Namelist input values

1 Each value is either a null value (10.11.3.4),  $c$ ,  $r^*c$ , or  $r^*$ , where  $c$  is a literal constant, optionally signed if integer or real, and  $r$  is an unsigned, nonzero, integer literal constant. A kind type parameter shall not be specified for  $c$

or  $r$ . The constant  $c$  is interpreted as though it had the same kind type parameter as the corresponding [effective item](#). The  $r^*c$  form is equivalent to  $r$  successive appearances of the constant  $c$ , and the  $r^*$  form is equivalent to  $r$  successive null values. Neither of these forms may contain embedded blanks, except where permitted within the constant  $c$ .

2 The datum  $c$  ([10.11](#)) is any input value acceptable to format specifications for a given type, except for a restriction on the form of input values corresponding to list items of types logical, integer, and character as specified in this subclause. The form of a real or complex value is dependent on the decimal edit mode in effect ([10.6](#)). The form of an input value shall be acceptable for the type of the namelist group object list item. The number and forms of the input values that may follow the equals in a name-value subsequence depend on the shape and type of the object represented by the name in the input record. When the name in the input record is that of a scalar variable of an intrinsic type, the equals shall not be followed by more than one value. Blanks are never used as zeros, and embedded blanks are not permitted in constants except within character constants and complex constants as specified in this subclause.

3 When the next [effective item](#) is of type real, the input form of the input value is that of a numeric input field. A numeric input field is a field suitable for F editing ([10.7.2.3.2](#)) that is assumed to have no fractional digits unless a decimal symbol appears within the field.

4 When the next [effective item](#) is of type complex, the input form of the input value consists of a left parenthesis followed by an ordered pair of numeric input fields separated by a comma (if the decimal edit mode is POINT) or a semicolon (if the decimal edit mode is COMMA), and followed by a right parenthesis. The first numeric input field is the real part of the complex constant and the second field is the imaginary part. Each of the numeric input fields may be preceded or followed by any number of blanks and ends of records. The end of a record may occur between the real part and the comma or semicolon, or between the comma or semicolon and the imaginary part.

5 When the next [effective item](#) is of type logical, the input form of the input value shall not include equals or value separators among the optional characters permitted for L editing ([10.7.3](#)).

6 When the next [effective item](#) is of type integer, the value in the input record is interpreted as if an  $Iw$  edit descriptor with a suitable value of  $w$  were used.

7 When the next [effective item](#) is of type character, the input form consists of a delimited sequence of zero or more *rep-chars* whose kind type parameter is implied by the kind of the corresponding list item. Such a sequence may be continued from the end of one record to the beginning of the next record, but the end of record shall not occur between a doubled apostrophe in an apostrophe-delimited sequence, nor between a doubled quote in a quote-delimited sequence. The end of the record does not cause a blank or any other character to become part of the sequence. The sequence may be continued on as many records as needed. The characters blank, comma, semicolon, and slash may appear in such character sequences.

#### NOTE 10.35

A character sequence corresponding to a namelist input item of character type shall be delimited either with apostrophes or with quotes. The delimiter is required to avoid ambiguity between unlimited character sequences and object names. The value of the `DELIM= specifier`, if any, in the `OPEN statement` for an `external file` is ignored during namelist input ([9.5.6.8](#)).

8 Let  $len$  be the length of the next [effective item](#), and let  $w$  be the length of the character sequence. If  $len$  is less than or equal to  $w$ , the leftmost  $len$  characters of the sequence are transmitted to the next [effective item](#). If  $len$  is greater than  $w$ , the constant is transmitted to the leftmost  $w$  characters of the next [effective item](#) and the remaining  $len-w$  characters of the next [effective item](#) are filled with blanks. The effect is as though the sequence were assigned to the next [effective item](#) in an [intrinsic assignment statement](#) ([7.2.1.3](#)).

#### 10.11.3.4 Null values

- 1 A null value is specified by
- the  $r^*$  form,

- blanks between two consecutive nonblank value separators following an equals,
- zero or more blanks preceding the first value separator and following an equals, or
- two consecutive nonblank value separators.

2 A null value has no effect on the definition status of the corresponding input list item. If the namelist group object list item is defined, it retains its previous value; if it is undefined, it remains undefined. A null value shall not be used as either the real or imaginary part of a complex constant, but a single null value may represent an entire complex constant.

#### NOTE 10.36

The end of a record following a value separator, with or without intervening blanks, does not specify a null value in namelist input.

### 10.11.3.5 Blanks

- 1 All blanks in a namelist input record are considered to be part of some value separator except for
- blanks embedded in a character constant,
  - embedded blanks surrounding the real or imaginary part of a complex constant,
  - leading blanks following the equals unless followed immediately by a slash or comma, or a semicolon if the decimal edit mode is COMMA, and
  - blanks between a name and the following equals.

### 10.11.3.6 Namelist comments

- 1 Except within a character literal constant, a “!” character after a value separator or in the first nonblank position of a namelist input record initiates a comment. The comment extends to the end of the record and may contain any graphic character in the processor-dependent character set. The comment is ignored. A slash within the namelist comment does not terminate execution of the namelist [input statement](#). Namelist comments are not allowed in stream input because comments depend on record structure.

#### NOTE 10.37

Namelist input example:

```
INTEGER I; REAL X (8); CHARACTER (11) P; COMPLEX Z; LOGICAL G
NAMELIST / TODAY / G, I, P, Z, X
READ (*, NML = TODAY)
```

The input data records are:

```
&TODAY I = 12345, X(1) = 12345, X(3:4) = 2*1.5, I=6, ! This is a comment.
P = 'ISN'T_BOB'S', Z = (123,0)/
```

The results stored are:

Variable	Value
I	6
X (1)	12345.0
X (2)	unchanged
X (3)	1.5
X (4)	1.5
X (5) – X (8)	unchanged
P	ISN'T_BOB'S
Z	(123.0,0.0)
G	unchanged

## 10.11.4 Namelist output

### 10.11.4.1 Form of namelist output

1 The form of the output produced by intrinsic namelist output shall be suitable for input, except for character output. The names in the output are in upper case. With the exception of adjacent undelimited character values, the values are separated by one or more blanks or by a comma, or a semicolon if the decimal edit mode is COMMA, optionally preceded by one or more blanks and optionally followed by one or more blanks. The form of the output produced by [defined output](#) (9.6.4.8) is determined by the [defined output](#) procedure; this form need not be compatible with namelist input.

2 Namelist output shall not include namelist comments.

3 The processor may begin new records as necessary. However, except for complex constants and character values, the end of a record shall not occur within a constant, character value, or name, and blanks shall not appear within a constant, character value, or name.

#### NOTE 10.38

The length of the output records is not specified exactly and is processor dependent.

### 10.11.4.2 Namelist output editing

1 Values in namelist output records are edited as for list-directed output ([10.10.4](#)).

#### NOTE 10.39

Namelist output records produced with a [DELIM= specifier](#) with a value of NONE and which contain a character sequence might not be acceptable as namelist input records.

### 10.11.4.3 Namelist output records

1 If two or more successive values for the same namelist group item in an output record produced have identical values, the processor has the option of producing a repeated constant of the form  $r*c$  instead of the sequence of identical values.

2 The name of each namelist group object list item is placed in the output record followed by an equals and a list of values of the namelist group object list item.

3 An ampersand character followed immediately by a *namelist-group-name* will be produced by namelist formatting at the start of the first output record to indicate which particular group of data objects is being output. A slash is produced by namelist formatting to indicate the end of the namelist formatting.

4 A null value is not produced by namelist formatting.

5 Except for new records created by explicit formatting within a [defined output](#) procedure or by continuation of delimited character sequences, each output record begins with a blank character.

1 (Blank page)

## 11 Program units

## 11.1 Main program

1 A Fortran main program is a **program unit** that does not contain a **SUBROUTINE**, **FUNCTION**, **MODULE**, **SUBMODULE**, or **BLOCK DATA** statement as its first statement.

```
R1101  main-program           is  [ program-stmt ]
                                   [ specification-part ]
                                   [ execution-part   ]
                                   [ internal-subprogram-part ]
                                   end-program-stmt
```

R1102   *program-stmt*                      is   PROGRAM *program-name*

R1103 *end-program-stmt*            is    END [ PROGRAM [ *program-name* ] ]

C1101 (R1101) The *program-name* may be included in the *end-program-stmt* only if the optional *program-stmt* is used and, if included, shall be identical to the *program-name* specified in the *program-stmt*.

### NOTE 11.1

The program name is global to the program (16.2). For explanatory information about uses for the program name, see subclause C.8.1.

### NOTE 11.2

An example of a main program is:

```

PROGRAM ANALYZE
    REAL A, B, C (10,10)      ! Specification part
    CALL FIND                  ! Execution part
CONTAINS
    SUBROUTINE FIND            ! Internal subprogram
        ...
    END SUBROUTINE FIND
END PROGRAM ANALYZE

```

2 The main program may be defined by means other than Fortran; in that case, the program shall not contain a *main-program* program unit.

3 A reference to a Fortran *main-program* shall not appear in any **program unit** in the program, including itself.

## 11.2 Modules

### 11.2.1 Module syntax and semantics

1 A **module** contains declarations, specifications, and definitions. Public identifiers of module entities are accessible to other **program units** by **use association** as specified in 11.2.2. A module that is provided as an inherent part of the processor is an intrinsic module. A nonintrinsic module is defined by a module **program unit** or a means other than Fortran.

2 Procedures and types defined in an intrinsic module are not themselves intrinsic.

1 R1104 *module* is *module-stmt*  
 2 [ *specification-part* ]  
 3 [ *module-subprogram-part* ]  
 4 *end-module-stmt*

5 R1105 *module-stmt* is MODULE *module-name*

6 R1106 *end-module-stmt* is END [ MODULE [ *module-name* ] ]

7 R1107 *module-subprogram-part* is *contains-stmt*  
 8 [ *module-subprogram* ] ...

9 R1108 *module-subprogram* is *function-subprogram*  
 10 or *subroutine-subprogram*  
 11 or *separate-module-subprogram*

12 C1102 (R1104) If the *module-name* is specified in the *end-module-stmt*, it shall be identical to the *module-name*  
 13 specified in the *module-stmt*.

14 C1103 (R1104) A module *specification-part* shall not contain a *stmt-function-stmt*, an *entry-stmt*, or a *format-stmt*.

15 3 If a procedure declared in the *scoping unit* of a module has an *implicit interface*, it shall be given the **EXTERNAL**  
 16 **attribute** in that *scoping unit*; if it is a function, its type and type parameters shall be explicitly declared in a  
 17 *type declaration statement* in that *scoping unit*.

18 4 If an intrinsic procedure is declared in the *scoping unit* of a module, it shall explicitly be given the **INTRINSIC**  
 19 **attribute** in that *scoping unit* or be used as an intrinsic procedure in that *scoping unit*.

#### NOTE 11.3

The module name is global to the program (16.2).

#### NOTE 11.4

Although *statement function definitions*, **ENTRY** statements, and **FORMAT** statements cannot appear in the specification part of a module, they can appear in the specification part of a module subprogram in the module.

#### NOTE 11.5

For a discussion of the impact of modules on dependent compilation, see subclause C.8.2.

#### NOTE 11.6

For examples of the use of modules, see subclause C.8.3.

## 20 11.2.2 The USE statement and use association

21 1 The USE statement specifies *use association*. A USE statement is a *reference* to the module it specifies. At the  
 22 time a USE statement is processed, the public portions of the specified module shall be available. A module shall  
 23 not reference itself, either directly or indirectly.

24 2 The USE statement provides the means by which a *scoping unit* accesses named data objects, derived types,  
 25 procedures, *abstract interfaces*, *generic identifiers*, and namelist groups in a module. The entities in the *scoping*  
 26 *unit* are use associated with the entities in the module. The accessed entities have the attributes specified in the  
 27 module, except that a local entity may have a different *accessibility attribute*, it may have the **ASYNCHRONOUS**  
 28 **attribute** even if the associated module entity does not, and if it is not a *coarray* it may have the **VOLATILE**  
 29 **attribute** even if the associated module entity does not. The entities made accessible are identified by the names  
 30 or *generic identifiers* used to identify them in the module. By default, the local entities are identified by the  
 31 same identifiers in the *scoping unit* containing the USE statement, but it is possible to specify that different local

identifiers are used. A use-associated entity is considered to have been previously defined.

### NOTE 11.7

The accessibility of module entities can be controlled by accessibility attributes (4.5.2.2, 5.5.2), and the ONLY option of the USE statement. Definability of module entities can be controlled by the **PROTECTED attribute** (5.5.15).

```
R1109  use-stmt      is  USE [ [ , module-nature ] :: ] module-name [ , rename-list ]
      or  USE [ [ , module-nature ] :: ] module-name , ■
      ■ ONLY : [ only-list ]
```

[illegible]

```
R1111  rename          is local-name => use-name  
      or OPERATOR (local-defined-operator) => ■  
      ■ OPERATOR (use-defined-operator)
```

R1112    *only*                  is    *generic-spec*  
                                     or    *only-use-name*  
                                     or    *rename*

R1113    *only-use-name*                      **is**    *use-name*

C1104 (R1109) If *module-nature* is INTRINSIC, *module-name* shall be the name of an intrinsic module.

C1105 (R1109) If *module-nature* is NON\_INTRINSIC, *module-name* shall be the name of a nonintrinsic module.

C1106 (R1109) A **scoping unit** shall not access an intrinsic module and a nonintrinsic module of the same name.

C1107 (R1111) OPERATOR (*use-defined-operator*) shall not identify a type-bound generic interface.

C1108 (R1112) The *generic-spec* shall not identify a type-bound generic interface.

### NOTE 11.8

Constraints C1107 and C1108 do not prevent accessing a *generic-spec* that is declared by an *interface block*, even if a *type-bound generic interface* has the same *generic-spec*.

C1109 (R1112) Each *generic-spec* shall be a public entity in the module.

C1110 (R1113) Each *use-name* shall be the name of a public entity in the module.

R1114 *local-defined-operator* is *defined-unary-op*  
or *defined-binary-op*

R1115	<i>use-defined-operator</i>	is	<i>defined-unary-op</i>
		or	<i>defined-binary-op</i>

C1111 (R1115) Each *use-defined-operator* shall be a public entity in the module.

3 A *use-stmt* without a *module-nature* provides access either to an intrinsic or to a nonintrinsic module. If the *module-name* is the name of both an intrinsic and a nonintrinsic module, the nonintrinsic module is accessed.

4 The USE statement without the **ONLY** option provides access to all public entities in the specified module.

5 A USE statement with the **ONLY** option provides access only to those entities that appear as *generic-specs*, *use-names*, or *use-defined-operators* in the *only-list*.

6 More than one USE statement for a given module may appear in a specification part. If one of the USE statements

is without an **ONLY** option, all public entities in the module are accessible. If all the USE statements have **ONLY** options, only those entities in one or more of the *only-lists* are accessible.

7 An accessible entity in the referenced module has one or more local identifiers. These identifiers are

- the identifier of the entity in the referenced module if that identifier appears as an *only-use-name* or as the *defined-operator* of a *generic-spec* in any *only* for that module,
- each of the *local-names* or *local-defined-operators* that the entity is given in any *rename* for that module, and
- the identifier of the entity in the referenced module if that identifier does not appear as a *use-name* or *use-defined-operator* in any *rename* for that module.

8 Two or more accessible entities, other than **generic interfaces** or defined operators, may have the same local identifier only if the identifier is not used. **Generic interfaces** and defined operators are handled as described in 12.4.3.5. Except for these cases, the local identifier of any entity given accessibility by a USE statement shall differ from the local identifiers of all other entities accessible to the **scoping unit**.

#### NOTE 11.9

There is no prohibition against a *use-name* or *use-defined-operator* appearing multiple times in one USE statement or in multiple USE statements involving the same module. As a result, it is possible for one use-associated entity to be accessible by more than one local identifier.

9 The local identifier of an entity made accessible by a USE statement shall not appear in any other nonexecutable statement that would cause any **attribute** (5.5) of the entity to be specified in the **scoping unit** that contains the USE statement, except that it may appear in a **PUBLIC** or **PRIVATE** statement in the **scoping unit** of a module and it may be given the **ASYNCHRONOUS** or **VOLATILE** attribute.

10 The appearance of such a local identifier in a **PUBLIC statement** in a module causes the entity accessible by the USE statement to be a public entity of that module. If the identifier appears in a **PRIVATE statement** in a module, the entity is not a public entity of that module. If the local identifier does not appear in either a **PUBLIC** or **PRIVATE** statement, it assumes the default accessibility attribute (5.6.1) of that **scoping unit**.

#### NOTE 11.10

The constraints in subclauses 5.9.1, 5.9.2, and 5.8 prohibit the *local-name* from appearing as a *common-block-object* in a **COMMON statement**, an *equivalence-object* in an **EQUIVALENCE statement**, or a *namelist-group-name* in a **NAMELIST statement**, respectively. There is no prohibition against the *local-name* appearing as a *common-block-name* or a *namelist-group-object*.

#### NOTE 11.11

For a discussion of the impact of the **ONLY** option and renaming on dependent compilation, see subclause C.8.2.1.

#### NOTE 11.12

Examples:

```
USE STATS_LIB
```

provides access to all public entities in the module STATS\_LIB.

```
USE MATH_LIB; USE STATS_LIB, SPROD => PROD
```

makes all public entities in both MATH\_LIB and STATS\_LIB accessible. If MATH\_LIB contains an entity called PROD, it is accessible by its own name while the entity PROD of STATS\_LIB is accessible by the name SPROD.

**NOTE 11.12 (cont.)**

```
USE STATS_LIB, ONLY: YPROD; USE STATS_LIB, ONLY : PROD
```

makes public entities YPROD and PROD in STATS\_LIB accessible.

```
USE STATS_LIB, ONLY : YPROD; USE STATS_LIB
```

makes all public entities in STATS\_LIB accessible.

**11.2.3 Submodules**

1 A **submodule** is a **program unit** that extends a module or another submodule. The **program unit** that it extends is its **host**, and is specified by the **parent-identifier** in the **submodule-stmt**.

2 A module or submodule is an ancestor program unit of all of its descendants, which are its submodules and their descendants. The submodule identifier is the ordered pair whose first element is the ancestor module name and whose second element is the submodule name; the submodule name by itself is not a local or global identifier.

**NOTE 11.13**

A module and its submodules stand in a tree-like relationship one to another, with the module at the root. Therefore, a submodule has exactly one ancestor module and can have one or more ancestor submodules.

3 A submodule may provide implementations for separate module procedures (12.6.2.5), each of which is declared (12.4.3.2) within that submodule or one of its ancestors, and declarations and definitions of other entities that are accessible by **host association** in its **descendants**.

```
R1116  submodule                is  submodule-stmt
                                     [ specification-part ]
                                     [ module-subprogram-part ]
                                     end-submodule-stmt
```

```
R1117  submodule-stmt          is  SUBMODULE ( parent-identifier ) submodule-name
```

```
R1118  parent-identifier        is  ancestor-module-name [ : parent-submodule-name ]
```

```
R1119  end-submodule-stmt      is  END [ SUBMODULE [ submodule-name ] ]
```

C1112 (R1116) A submodule **specification-part** shall not contain a **format-stmt**, **entry-stmt**, or **stmt-function-stmt**.

C1113 (R1118) The **ancestor-module-name** shall be the name of a nonintrinsic module; the **parent-submodule-name** shall be the name of a **descendant** of that module.

C1114 (R1116) If a **submodule-name** appears in the **end-submodule-stmt**, it shall be identical to the one in the **submodule-stmt**.

**11.3 Block data program units**

1 A block data program unit is used to provide initial values for data objects in named **common blocks**.

```
R1120  block-data              is  block-data-stmt
                                     [ specification-part ]
                                     end-block-data-stmt
```

```
R1121  block-data-stmt         is  BLOCK DATA [ block-data-name ]
```

```
R1122  end-block-data-stmt     is  END [ BLOCK DATA [ block-data-name ] ]
```

C1115 (R1120) The **block-data-name** shall be included in the **end-block-data-stmt** only if it was provided in the **block-data-stmt** and, if included, shall be identical to the **block-data-name** in the **block-data-stmt**.

- 1 C1116 (R1120) A *block-data specification-part* shall contain only definitions of derived-type definitions and ASYNCHRONOUS,  
2 BIND, COMMON, DATA, DIMENSION, EQUIVALENCE, IMPLICIT, INTRINSIC, PARAMETER, POINTER, SAVE,  
3 TARGET, USE, VOLATILE, and type declaration statements.
- 4 C1117 (R1120) A type declaration statement in a *block-data specification-part* shall not contain ALLOCATABLE, EXTERNAL,  
5 or BIND attribute specifiers.
- 6 2 If an object in a named common block is initially defined, all storage units in the common block storage sequence shall be specified  
7 even if they are not all initially defined. More than one named common block may have objects initially defined in a single block  
8 data program unit.
- 9 3 Only an object in a named common block may be initially defined in a block data program unit.
- 10 4 The same named common block shall not be specified in more than one block data program unit in a program.
- 11 5 There shall not be more than one unnamed block data program unit in a program.

## 12 Procedures

### 12.1 Concepts

- 1 The concept of a procedure was introduced in 2.2.3. This clause contains a complete description of procedures. The actions specified by a procedure are performed when the procedure is invoked by execution of a reference to it.
- 2 The sequence of actions encapsulated by a procedure has access to entities in the procedure reference by way of [argument association](#) (12.5.2). A name that appears as a *dummy-arg-name* in the SUBROUTINE, FUNCTION, or ENTRY statement in the declaration of a procedure (R1237) is a [dummy argument](#). [Dummy arguments](#) are also specified for intrinsic procedures and procedures in intrinsic modules in Clauses 13, 14, and 15.

### 12.2 Procedure classifications

#### 12.2.1 Procedure classification by reference

- 1 The definition of a procedure specifies it to be a function or a subroutine. A reference to a function either appears explicitly as a primary within an expression, or is implied by a [defined operation](#) (7.1.6) within an expression. A reference to a subroutine is a [CALL statement](#), a [defined assignment](#) statement (7.2.1.4), the appearance of an object processed by [defined input/output](#) (9.6.4.8) in an input/output list, or [finalization](#) (4.5.6).
- 2 A procedure is classified as [elemental](#) if it is a procedure that may be referenced [elementally](#) (12.8).

#### 12.2.2 Procedure classification by means of definition

##### 12.2.2.1 Intrinsic procedures

- 1 A procedure that is provided as an inherent part of the processor is an intrinsic procedure.

##### 12.2.2.2 External, internal, and module procedures

- 1 An [external procedure](#) is a procedure that is defined by an external subprogram or by a means other than Fortran.
- 2 An [internal procedure](#) is a procedure that is defined by an internal subprogram. Internal subprograms may appear in the main program, in an external subprogram, or in a module subprogram. Internal subprograms shall not appear in other internal subprograms. Internal subprograms are the same as external subprograms except that the name of the [internal procedure](#) is not a global identifier, an internal subprogram shall not contain an [ENTRY statement](#), and the internal subprogram has access to host entities by [host association](#).
- 3 A [module procedure](#) is a procedure that is defined by a module subprogram.
- 4 A subprogram defines a procedure for the SUBROUTINE or FUNCTION statement. If the subprogram has one or more [ENTRY statements](#), it also defines a procedure for each of them.

##### 12.2.2.3 Dummy procedures

- 1 A [dummy argument](#) that is specified to be a procedure or appears as the procedure designator in a [procedure reference](#) is a [dummy procedure](#). A [dummy procedure](#) with the [POINTER attribute](#) is a [dummy procedure pointer](#).

#### 12.2.2.4 Procedure pointers

- 1 A [procedure pointer](#) is a procedure that has the [EXTERNAL](#) and [POINTER](#) attributes; it may be [pointer associated](#) with an [external procedure](#), an [internal procedure](#), an intrinsic procedure, a [module procedure](#), or a [dummy procedure](#) that is not a [procedure pointer](#).

#### 12.2.2.5 Statement functions

- 1 A function that is defined by a single statement is a statement function ([12.6.4](#)).

## 12.3 Characteristics

### 12.3.1 Characteristics of procedures

- 1 The [characteristics](#) of a procedure are the classification of the procedure as a function or subroutine, whether it is pure, whether it is [elemental](#), whether it has the [BIND attribute](#), the [characteristics](#) of its [dummy arguments](#), and the [characteristics](#) of its [function result](#) if it is a function.

### 12.3.2 Characteristics of dummy arguments

#### 12.3.2.1 General

- 1 Each [dummy argument](#) has the [characteristic](#) that it is a [dummy data object](#), a [dummy procedure](#), or an asterisk (alternate return indicator).

#### 12.3.2.2 Characteristics of dummy data objects

- 1 The [characteristics](#) of a [dummy data object](#) are its type, its [type parameters](#) (if any), its shape (unless it is [assumed-rank](#)), its [corank](#), its [codimensions](#), its intent ([5.5.10](#), [5.6.9](#)), whether it is optional ([5.5.12](#), [5.6.10](#)), whether it is [allocatable](#) ([5.5.3](#)), whether it has the [ASYNCHRONOUS](#) ([5.5.4](#)), [CONTIGUOUS](#) ([5.5.7](#)), [VALUE](#) ([5.5.18](#)), or [VOLATILE](#) ([5.5.19](#)) attributes, whether it is [polymorphic](#), and whether it is a [pointer](#) ([5.5.14](#), [5.6.12](#)) or a target ([5.5.17](#), [5.6.15](#)). If a [type parameter](#) of an object or a bound of an array is not a [constant expression](#), the exact dependence on the entities in the expression is a characteristic. If a rank, shape, size, type, or type parameter is assumed or [deferred](#), it is a characteristic.

#### 12.3.2.3 Characteristics of dummy procedures

- 1 The [characteristics](#) of a [dummy procedure](#) are the explicitness of its [interface](#) ([12.4.2](#)), its [characteristics](#) as a procedure if the [interface](#) is [explicit](#), whether it is a pointer, and whether it is optional ([5.5.12](#), [5.6.10](#)).

#### 12.3.2.4 Characteristics of asterisk dummy arguments

- 1 An asterisk as a [dummy argument](#) has no [characteristics](#).

### 12.3.3 Characteristics of function results

- 1 The [characteristics](#) of a function result are its type, type parameters (if any), [rank](#), whether it is [polymorphic](#), whether it is [allocatable](#), whether it is a [pointer](#), whether it has the [CONTIGUOUS attribute](#), and whether it is a [procedure pointer](#). If a function result is an array that is not [allocatable](#) or a pointer, its shape is a characteristic. If a type parameter of a function result or a bound of a function result array is not a [constant expression](#), the exact dependence on the entities in the expression is a characteristic. If type parameters of a function result are [deferred](#), which parameters are [deferred](#) is a characteristic. Whether the length of a character function result is assumed is a characteristic.

## 12.4 Procedure interface

### 12.4.1 Interface and abstract interface

- 1 The **interface** of a procedure determines the forms of reference through which it may be invoked. The procedure's interface consists of its name, **binding label**, generic identifiers, characteristics, and the names of its dummy arguments. The **characteristics** and **binding label** of a procedure are fixed, but the remainder of the interface may differ in differing contexts, except that for a separate module procedure body (12.6.2.5), the **dummy argument** names and whether it is recursive shall be the same as in its corresponding module procedure interface body (12.4.3.2).
- 2 An **abstract interface** is a set of procedure **characteristics** with the **dummy argument** names.

### 12.4.2 Implicit and explicit interfaces

#### 12.4.2.1 Interfaces and scopes

- 1 The interface of a procedure is either explicit or implicit. It is explicit if it is
  - an **internal procedure**, **module procedure**, or **intrinsic** procedure,
  - a subroutine, or a function with a separate result name, within the **scoping unit** that defines it, or
  - a procedure declared by a **procedure declaration statement** that specifies an explicit interface, or by an interface body.
- Otherwise, the interface of the identifier is implicit. The interface of a statement function is always implicit.

#### NOTE 12.1

For example, the subroutine LLS of C.8.3.4 has an **explicit interface**.

#### 12.4.2.2 Explicit interface

- 1 Within the scope of a procedure identifier, the procedure shall have an **explicit interface** if it is not a statement function and
  - (1) a reference to the procedure appears
    - (a) with an **argument keyword** (12.5.2), or
    - (b) in a context that requires it to be pure,
  - (2) the procedure has a **dummy argument** that
    - (a) has the **ALLOCATABLE**, **ASYNCHRONOUS**, **OPTIONAL**, **POINTER**, **TARGET**, **VALUE**, or **VOLATILE** attribute,
    - (b) is an **assumed-shape array**,
    - (c) is **assumed-rank**,
    - (d) is a **coarray**,
    - (e) is of a parameterized derived type, or
    - (f) is **polymorphic**,
  - (3) the procedure has a result that
    - (a) is an array,
    - (b) is a **pointer** or is **allocatable**, or
    - (c) has a nonassumed type parameter value that is not a **constant expression**,
  - (4) the procedure is **elemental**, or
  - (5) the procedure has the **BIND attribute**.

## 12.4.3 Specification of the procedure interface

### 12.4.3.1 General

- 1 The interface for an [internal](#), [external](#), [module](#), or [dummy procedure](#) is specified by a [FUNCTION](#), [SUBROUTINE](#), or [ENTRY](#) statement and by specification statements for the [dummy arguments](#) and the result of a function. These statements may appear in the procedure definition, in an [interface body](#), or both, except that the [ENTRY statement](#) shall not appear in an [interface body](#).

#### NOTE 12.2

An [interface body](#) cannot be used to describe the interface of an [internal procedure](#), a [module procedure](#) that is not a separate [module procedure](#), or an intrinsic procedure because the interfaces of such procedures are already [explicit](#). However, the name of a procedure can appear in a [PROCEDURE statement](#) in an [interface block](#) (12.4.3.2).

### 12.4.3.2 Interface block

- |       |                                |                      |  |
|-------|--------------------------------|----------------------|--|
| R1201 | <i>interface-block</i>         | is                   | <i>interface-stmt</i><br>[ <i>interface-specification</i> ] ...<br><i>end-interface-stmt</i>   |
| R1202 | <i>interface-specification</i> | is<br>or             | <i>interface-body</i><br><i>procedure-stmt</i>   |
| R1203 | <i>interface-stmt</i>          | is<br>or             | INTERFACE [ <i>generic-spec</i> ]<br>ABSTRACT INTERFACE  |
| R1204 | <i>end-interface-stmt</i>      | is                   | END INTERFACE [ <i>generic-spec</i> ]  |
| R1205 | <i>interface-body</i>          | is<br>or             | <i>function-stmt</i><br>[ <i>specification-part</i> ]<br><i>end-function-stmt</i><br><i>subroutine-stmt</i><br>[ <i>specification-part</i> ]<br><i>end-subroutine-stmt</i> |
| R1206 | <i>procedure-stmt</i>          | is                   | [ MODULE ] PROCEDURE [ :: ] <i>specific-procedure-list</i>   |
| R1207 | <i>specific-procedure</i>      | is                   | <i>procedure-name</i>  |
| R1208 | <i>generic-spec</i>            | is<br>or<br>or<br>or | <i>generic-name</i><br>OPERATOR ( <i>defined-operator</i> )<br>ASSIGNMENT ( = )<br><i>defined-io-generic-spec</i>  |
| R1209 | <i>defined-io-generic-spec</i> | is<br>or<br>or<br>or | READ (FORMATTED)<br>READ (UNFORMATTED)<br>WRITE (FORMATTED)<br>WRITE (UNFORMATTED)   |
- C1201 (R1201) An [interface-block](#) in a subprogram shall not contain an [interface-body](#) for a procedure defined by that subprogram.
- C1202 (R1201) If the [end-interface-stmt](#) includes *generic-name*, the [interface-stmt](#) shall specify the same *generic-name*. If the [end-interface-stmt](#) includes [ASSIGNMENT\(=\)](#), the [interface-stmt](#) shall specify [ASSIGNMENT\(=\)](#). If the [end-interface-stmt](#) includes *defined-io-generic-spec*, the [interface-stmt](#) shall specify the same *defined-io-generic-spec*. If the [end-interface-stmt](#) includes [OPERATOR\(defined-operator\)](#), the [interface-stmt](#) shall specify the same *defined-operator*. If one *defined-operator* is [.LT.](#), [.LE.](#), [.GT.](#), [.GE.](#), [.EQ.](#), or [.NE.](#), the other is permitted to be the corresponding operator [<](#), [<=](#), [>](#), [>=](#), [==](#), or [/=](#).

- 1 C1203 (R1203) If the *interface-stmt* is ABSTRACT INTERFACE, then the *function-name* in the *function-stmt*  
 2 or the *subroutine-name* in the *subroutine-stmt* shall not be the same as a keyword that specifies an  
 3 intrinsic type.
- 4 C1204 (R1202) A *procedure-stmt* is allowed only in an interface block that has a *generic-spec*.
- 5 C1205 (R1205) An *interface-body* of a pure procedure shall specify the intents of all dummy arguments except  
 6 alternate return indicators, dummy procedures, and arguments with the POINTER or VALUE attribute.
- 7 C1206 (R1205) An *interface-body* shall not contain a *data-stmt*, *format-stmt*, *entry-stmt*, or *stmt-function-stmt*.
- 8 C1207 (R1206) If MODULE appears in a *procedure-stmt*, each *procedure-name* in that statement shall denote a  
 9 module procedure.
- 10 C1208 (R1207) A *procedure-name* shall denote a nonintrinsic procedure that has an explicit interface.
- 11 C1209 (R1201) An *interface-specification* in a generic interface block shall not specify a procedure that was  
 12 specified previously in any accessible interface with the same generic identifier.
- 13 1 An external or module subprogram specifies a specific interface for each procedure defined in that subprogram.
- 14 2 An interface block introduced by ABSTRACT INTERFACE is an abstract interface block. An interface body  
 15 in an abstract interface block specifies an abstract interface. An interface block with a generic specification is  
 16 a generic interface block. An interface block with neither ABSTRACT nor a generic specification is a specific  
 17 interface block.
- 18 3 The name of the entity declared by an interface body is the *function-name* in the *function-stmt* or the *subroutine-*  
 19 *name* in the *subroutine-stmt* that begins the interface body.
- 20 4 A module procedure interface body is an interface body whose initial statement contains the keyword MODULE.  
 21 It specifies the interface for a separate module procedure (12.6.2.5). A separate module procedure is accessible  
 22 by use association if and only if its interface body is declared in the specification part of a module and is public.  
 23 If a corresponding (12.6.2.5) separate module procedure is not defined, the interface may be used to specify an  
 24 explicit specific interface but the procedure shall not be used in any other way.
- 25 5 An interface body in a generic or specific interface block specifies the EXTERNAL attribute and an explicit  
 26 specific interface for an external procedure or a dummy procedure. If the name of the declared procedure is that  
 27 of a dummy argument in the subprogram containing the interface body, the procedure is a dummy procedure;  
 28 otherwise, it is an external procedure.
- 29 6 An interface body specifies all of the characteristics of the explicit specific interface or abstract interface. The  
 30 specification part of an interface body may specify attributes or define values for data entities that do not  
 31 determine characteristics of the procedure. Such specifications have no effect.
- 32 7 If an explicit specific interface for an external procedure is specified by an interface body or a procedure declaration  
 33 statement (12.4.3.7), the characteristics shall be consistent with those specified in the procedure definition, except  
 34 that the interface may specify a procedure that is not pure even if the procedure is defined to be pure. An interface  
 35 for a procedure defined by an ENTRY statement may be specified by using the entry name as the procedure name in the interface body.  
 36 If an external procedure does not exist in the program, an interface body for it may be used to specify an explicit  
 37 specific interface but the procedure shall not be used in any other way. A procedure shall not have more than  
 38 one explicit specific interface in a given scoping unit, except that if the interface is accessed by use association,  
 39 there may be more than one local name for the procedure. If a procedure is accessed by use association, each  
 40 access shall be to the same procedure declaration or definition.

**NOTE 12.3**

The dummy argument names in an interface body can be different from the corresponding dummy argument names in the procedure definition because the name of a dummy argument is not a characteristic.

**NOTE 12.4**

An example of a specific interface block is:

```

INTERFACE
  SUBROUTINE EXT1 (X, Y, Z)
    REAL, DIMENSION (100, 100) :: X, Y, Z
  END SUBROUTINE EXT1
  SUBROUTINE EXT2 (X, Z)
    REAL X
    COMPLEX (KIND = 4) Z (2000)
  END SUBROUTINE EXT2
  FUNCTION EXT3 (P, Q)
    LOGICAL EXT3
    INTEGER P (1000)
    LOGICAL Q (1000)
  END FUNCTION EXT3
END INTERFACE

```

This interface block specifies [explicit interfaces](#) for the three [external procedures](#) EXT1, EXT2, and EXT3. Invocations of these procedures can use [argument keywords](#) (12.5.2); for example:

```
PRINT *, EXT3 (Q = P_MASK (N+1 : N+1000), P = ACTUAL_P)
```

**12.4.3.3 GENERIC statement**

- 1 A **GENERIC** statement specifies a generic identifier for one or more specific procedures, in the same way as a generic interface block that does not contain interface bodies.

R1210 *generic-stmt*                                    **is** **GENERIC** [ , *access-spec* ] :: *generic-spec* => *specific-procedure-list*

- 2 If *access-spec* appears, it specifies the accessibility (5.5.2) of *generic-spec*.

**12.4.3.4 IMPORT statement**

R1211 *import-stmt*                                    **is** **IMPORT** [[ :: ] *import-name-list* ]  
**or** **IMPORT, ONLY** : *import-name-list*  
**or** **IMPORT, NONE**  
**or** **IMPORT, ALL**

C1210 (R1211) An **IMPORT** statement shall not appear in the *specification-part* of a *main-program*, *external-subprogram*, *module*, or *block-data*.

C1211 (R1211) Each *import-name* shall be the name of an entity in the *host scoping unit*.

C1212 If any **IMPORT** statement in a *scoping unit* has an **ONLY** specifier, all **IMPORT** statements in that *scoping unit* shall have an **ONLY** specifier.

C1213 **IMPORT, NONE** shall not appear in a submodule.

C1214 If an **IMPORT, NONE** or **IMPORT, ALL** statement appears in a *scoping unit*, no other **IMPORT** statement shall appear in that *scoping unit*.

C1215 Within an interface body, an entity that is accessed by host association shall be accessible by host or use association within the *host scoping unit*, or explicitly declared prior to the interface body.

C1216 An entity whose name appears as an *import-name* or which is made accessible by an **IMPORT, ALL** statement shall not appear in any context described in 16.5.1.4 that would cause the host entity of that name to be inaccessible.

- 1 1 If the **ONLY** specifier appears on an IMPORT statement in a **scoping unit**, an entity is only accessible by host association if its name appears as an *import-name* in that **scoping unit**.
- 2 2 An IMPORT, NONE statement in a **scoping unit** specifies that no entities in the **host scoping unit** are accessible by host association in that **scoping unit**. This is the default for an interface body for an **external** or **dummy** procedure.
- 3 3 An IMPORT, ALL statement in a **scoping unit** specifies that all entities from the **host scoping unit** are accessible by host association in that **scoping unit**.
- 4 4 If an IMPORT statement with no *import-name-list* appears in a **scoping unit**, every entity in the **host scoping unit** is accessible unless its name appears in a context described in 16.5.1.4 that causes it to be inaccessible. This is the default for a nested scoping unit other than an interface body for an **external** or **dummy** procedure.
- 5 5 If an IMPORT statement with an *import-name-list* appears in a **scoping unit**, each named entity from the **host scoping unit** is accessible by host association.

**NOTE 12.5**

The IMPORT, NONE statement can be used to prevent accidental host association:

```

SUBROUTINE s(x,n)
  IMPLICIT NONE
  IMPORT, NONE
  ...
  DO i=1,n ! Forces I to be locally declared.

```

**NOTE 12.6**

The IMPORT, ALL statement can be used to confirm the default rules and prevent accidental “shadowing” of host entities:

```

SUBROUTINE outer
  REAL x
  ...
CONTAINS
  SUBROUTINE inner
    IMPORT, ALL
    ...
    x = x + 1 ! X was prevented from being locally declared...
              ! so must be the host X.

```

**NOTE 12.7**

The IMPORT, **ONLY** statement can be used to document deliberate access via host association whilst blocking accidental access:

```

SUBROUTINE sub
  IMPORT,ONLY :: x, y
  ...
  x = y + z ! X and Y imported, Z must be local.

```

**NOTE 12.8**

The IMPORT statement can be used to allow **module procedures** to have **dummy arguments** that are procedures with **assumed-shape** arguments of an opaque type. For example:

```

MODULE M

```

## NOTE 12.8 (cont.)

```

TYPE T
  PRIVATE      ! T is an opaque type
  ...
END TYPE
CONTAINS
SUBROUTINE PROCESS(X,Y,RESULT,MONITOR)
  TYPE(T),INTENT(IN) :: X(:,,:),Y(:,,:)
  TYPE(T),INTENT(OUT) :: RESULT(:,,:)
  INTERFACE
    SUBROUTINE MONITOR(ITERATION_NUMBER,CURRENT_ESTIMATE)
      IMPORT T
      INTEGER,INTENT(IN) :: ITERATION_NUMBER
      TYPE(T),INTENT(IN) :: CURRENT_ESTIMATE(:,,:)
    END SUBROUTINE
  END INTERFACE
  ...
END SUBROUTINE
END MODULE

```

The MONITOR dummy procedure requires an explicit interface because it has an assumed-shape array argument, but TYPE(T) would not be available inside the interface body without the IMPORT statement.

## 12.4.3.5 Generic interfaces

## 12.4.3.5.1 Generic identifiers

1 A generic interface block specifies a generic interface for each of the procedures in the interface block. The PROCEDURE statement lists procedure pointers, external procedures, dummy procedures, or module procedures that have this generic interface. A GENERIC statement specifies a generic interface for each of the procedures named in its specific-procedure-list. A generic interface is always explicit.

2 The generic-spec in an interface-stmt is a generic identifier for all the procedures in the interface block. The generic-spec in a GENERIC statement is a generic identifier for all of the procedures named in its specific-procedure-list. The rules specifying how any two procedures with the same generic identifier shall differ are given in 12.4.3.5.5. They ensure that any generic invocation applies to at most one specific procedure. If a specific procedure in a generic interface has a function dummy argument, that argument shall have its type and type parameters explicitly declared in the specific interface.

3 A generic name is a generic identifier that refers to all of the procedure names in the interface block. A generic name may be the same as any one of the procedure names in the interface block, or the same as any accessible generic name.

4 A generic name may be the same as a derived-type name, in which case all of the procedures in the interface block shall be functions.

5 An interface-stmt having a defined-io-generic-spec is an interface for a defined input/output procedure (9.6.4.8).

## NOTE 12.9

An example of a generic procedure interface is:

```

INTERFACE SWITCH
  SUBROUTINE INT_SWITCH (X, Y)
    INTEGER, INTENT (INOUT) :: X, Y
  END SUBROUTINE INT_SWITCH
  SUBROUTINE REAL_SWITCH (X, Y)

```

**NOTE 12.9 (cont.)**

```

    REAL, INTENT (INOUT) :: X, Y
  END SUBROUTINE REAL_SWITCH
  SUBROUTINE COMPLEX_SWITCH (X, Y)
    COMPLEX, INTENT (INOUT) :: X, Y
  END SUBROUTINE COMPLEX_SWITCH
END INTERFACE SWITCH

```

Any of these three subroutines (INT\_SWITCH, REAL\_SWITCH, COMPLEX\_SWITCH) can be referenced with the generic name SWITCH, as well as by its [specific name](#). For example, a reference to INT\_SWITCH could take the form:

```
CALL SWITCH (MAX_VAL, LOC_VAL) ! MAX_VAL and LOC_VAL are of type INTEGER
```

**NOTE 12.10**

A [type-bound-generic-stmt](#) within a derived-type definition (4.5.5) specifies a [generic identifier](#) for a set of [type-bound procedures](#).

**12.4.3.5.2 Defined operations**

- 1 If [OPERATOR](#) is specified in a generic specification, all of the procedures specified in the [generic interface](#) shall be functions that may be referenced as [defined operations](#) (7.1.6, 12.5). In the case of functions of two arguments, infix binary operator notation is implied. In the case of functions of one argument, prefix operator notation is implied. [OPERATOR](#) shall not be specified for functions with no arguments or for functions with more than two arguments. The [dummy arguments](#) shall be nonoptional [dummy data objects](#) and shall have the [INTENT \(IN\)](#) or [VALUE](#) attribute. The function result shall not have assumed character length. If the operator is an *intrinsic-operator* (R308), the number of [dummy arguments](#) shall be consistent with the intrinsic uses of that operator, and the types, kind type parameters, or [ranks](#) of the [dummy arguments](#) shall differ from those required for the intrinsic operation (7.1.5).
- 2 A [defined operation](#) is treated as a reference to the function. For a unary [defined operation](#), the operand corresponds to the function's [dummy argument](#); for a binary operation, the left-hand operand corresponds to the first [dummy argument](#) of the function and the right-hand operand corresponds to the second [dummy argument](#). All restrictions and constraints that apply to [actual arguments](#) in a [reference](#) to the function also apply to the corresponding operands in the expression as if they were used as [actual arguments](#).
- 3 A given defined operator may, as with generic names, apply to more than one function, in which case it is generic in exact analogy to generic procedure names. For intrinsic operator symbols, the generic properties include the intrinsic operations they represent. Because both forms of each relational operator have the same interpretation (7.1.6.2), extending one form (such as <=) has the effect of defining both forms (<= and [.LE.](#)).

**NOTE 12.11**

An example of the use of the [OPERATOR](#) generic specification is:

```

INTERFACE OPERATOR ( * )
  FUNCTION BOOLEAN_AND (B1, B2)
    LOGICAL, INTENT (IN) :: B1 (:), B2 (SIZE (B1))
    LOGICAL :: BOOLEAN_AND (SIZE (B1))
  END FUNCTION BOOLEAN_AND
END INTERFACE OPERATOR ( * )

```

This allows, for example

```
SENSOR (1:N) * ACTION (1:N)
```

**NOTE 12.11 (cont.)**

as an alternative to the function call

```
BOOLEAN_AND (SENSOR (1:N), ACTION (1:N))    ! SENSOR and ACTION are
                                              ! of type LOGICAL
```

**12.4.3.5.3 Defined assignments**

- 1 If **ASSIGNMENT** ( = ) is specified in a generic specification, all the procedures in the **generic interface** shall be subroutines that may be referenced as **defined assignments** (7.2.1.4, 7.2.1.5). **Defined assignment** may, as with generic names, apply to more than one subroutine, in which case it is generic in exact analogy to generic procedure names.
- 2 Each of these subroutines shall have exactly two **dummy arguments**. The **dummy arguments** shall be nonoptional **dummy data objects**. The first argument shall have **INTENT (OUT)** or **INTENT (INOUT)** and the second argument shall have the **INTENT (IN)** or **VALUE** attribute. Either the second argument shall be an array whose **rank** differs from that of the first argument, the **declared types** and **kind type parameters** of the arguments shall not conform as specified in Table 7.8, or the first argument shall be of derived type. A **defined assignment** is treated as a reference to the subroutine, with the left-hand side as the first argument and the right-hand side enclosed in parentheses as the second argument. All restrictions and constraints that apply to **actual arguments** in a reference to the subroutine also apply to the left-hand-side and to the right-hand-side enclosed in parentheses as if they were used as **actual arguments**. The **ASSIGNMENT** generic specification specifies that assignment is extended or redefined.

**NOTE 12.12**

An example of the use of the **ASSIGNMENT** generic specification is:

```
INTERFACE ASSIGNMENT ( = )
```

```
  SUBROUTINE LOGICAL_TO_NUMERIC (N, B)
```

```
    INTEGER, INTENT (OUT) :: N
```

```
    LOGICAL, INTENT (IN)  :: B
```

```
  END SUBROUTINE LOGICAL_TO_NUMERIC
```

```
  SUBROUTINE CHAR_TO_STRING (S, C)
```

```
    USE STRING_MODULE      ! Contains definition of type STRING
```

```
    TYPE (STRING), INTENT (OUT) :: S ! A variable-length string
```

```
    CHARACTER (*), INTENT (IN)  :: C
```

```
  END SUBROUTINE CHAR_TO_STRING
```

```
END INTERFACE ASSIGNMENT ( = )
```

Example assignments are:

```
KOUNT = SENSOR (J)    ! CALL LOGICAL_TO_NUMERIC (KOUNT, (SENSOR (J)))
```

```
NOTE  = '89AB'        ! CALL CHAR_TO_STRING (NOTE, ('89AB'))
```

**NOTE 12.13**

A procedure whose second **dummy argument** has the **ALLOCATABLE** or **POINTER** attribute cannot be accessed via **defined assignment**, even if it given the **ASSIGNMENT** ( = ) generic identifier. This is because the **actual argument** associated with that **dummy argument** is the right-hand side of the assignment enclosed in parentheses, which makes the **actual argument** an expression that does not have the **ALLOCATABLE**, **POINTER**, or **TARGET** attribute.

#### 12.4.3.5.4 Defined input/output procedure interfaces

- 1 All of the procedures specified in an interface block for a **defined input/output** procedure shall be subroutines that have interfaces as described in 9.6.4.8.3.

#### 12.4.3.5.5 Restrictions on generic declarations

- 1 This subclause contains the rules that shall be satisfied by every pair of specific procedures that have the same **generic identifier** within the scope of the identifier. If a generic procedure is accessed from a module, the rules apply to all the specific versions even if some of them are inaccessible by their **specific names**.

##### NOTE 12.14

In most **scoping units**, the possible sources of procedures with a particular **generic identifier** are the accessible interface blocks and the generic **bindings** other than names for the accessible objects in that **scoping unit**. In a type definition, they are the generic **bindings**, including those from a **parent type**.

- 2 A **dummy argument** is type, kind, and **rank** compatible, or TKR compatible, with another **dummy argument** if the first is **type compatible** with the second, the kind type parameters of the first have the same values as the corresponding kind type parameters of the second, and both have the same **rank** or either is **assumed-rank**.

- 3 Two **dummy arguments** are distinguishable if

- one is a procedure and the other is a data object,
- they are both data objects or known to be functions, and neither is TKR compatible with the other,
- one has the **ALLOCATABLE attribute** and the other has the **POINTER attribute** and not the **INTENT (IN) attribute**, or
- one is a function with nonzero **rank** and the other is not known to be a function.

C1217 Within the scope of a generic operator, if two procedures with that identifier have the same number of arguments, one shall have a **dummy argument** that corresponds by position in the argument list to a **dummy argument** of the other that is distinguishable from it.

C1218 Within the scope of the generic **ASSIGNMENT (=)** identifier, if two procedures have that identifier, one shall have a **dummy argument** that corresponds by position in the argument list to a **dummy argument** of the other that is distinguishable from it.

C1219 Within the scope of a **defined-io-generic-spec**, if two procedures have that **generic identifier**, their **dtv arguments** (9.6.4.8.3) shall be distinguishable.

C1220 Within the scope of a generic name, each pair of procedures identified by that name shall both be subroutines or both be functions, and

- (1) there is a non-passed-object **dummy data object** in one or the other of them such that
  - (a) the number of **dummy data objects** in one that are nonoptional, are not passed-object, and with which that **dummy data object** is TKR compatible, possibly including that **dummy data object** itself, exceeds
  - (b) the number of non-passed-object **dummy data objects**, both optional and nonoptional, in the other that are not distinguishable from that **dummy data object**,
- (2) the number of nonoptional **dummy procedures** in one of them exceeds the number of **dummy procedures** in the other,
- (3) both have **passed-object dummy arguments** and the **passed-object dummy arguments** are distinguishable, or
- (4) at least one of them shall have both

- (a) a nonoptional non-passed-object [dummy argument](#) at an effective position such that either the other procedure has no [dummy argument](#) at that effective position or the [dummy argument](#) at that position is distinguishable from it, and
- (b) a nonoptional non-passed-object [dummy argument](#) whose name is such that either the other procedure has no [dummy argument](#) with that name or the [dummy argument](#) with that name is distinguishable from it,
- and the [dummy argument](#) that disambiguates by position shall either be the same as or occur earlier in the argument list than the one that disambiguates by name.

4 The effective position of a [dummy argument](#) is its position in the argument list after any [passed-object dummy argument](#) has been removed.

5 Within the scope of a generic name that is the same as the generic name of an intrinsic procedure, the intrinsic procedure is not accessible by its generic name if the procedures in the interface and the intrinsic procedure are not all functions or not all subroutines. If a generic invocation is consistent with both a specific procedure from an interface and an accessible intrinsic procedure, it is the specific procedure from the interface that is referenced.

#### NOTE 12.15

An extensive explanation of the application of these rules is in [C.9.6](#).

### 12.4.3.6 EXTERNAL statement

1 An EXTERNAL statement specifies the [EXTERNAL attribute](#) (5.5.9) for a list of names.

R1212 *external-stmt* is EXTERNAL [ :: ] *external-name-list*

2 The appearance of the name of a [block data program unit](#) in an EXTERNAL statement confirms that the [block data program unit](#) is a part of the program.

#### NOTE 12.16

For explanatory information on potential portability problems with [external procedures](#), see subclause [C.9.1](#).

#### NOTE 12.17

An example of an EXTERNAL statement is:

EXTERNAL FOCUS

### 12.4.3.7 Procedure declaration statement

1 A procedure declaration statement declares [procedure pointers](#), [dummy procedures](#), and [external procedures](#). It specifies the [EXTERNAL attribute](#) (5.5.9) for all entities in the *proc-decl-list*.

R1213 *procedure-declaration-stmt* is PROCEDURE ( [ *proc-interface* ] ) ■  
■ [ [ , *proc-attr-spec* ] ... :: ] *proc-decl-list*

R1214 *proc-interface* is *interface-name*  
or *declaration-type-spec*

R1215 *proc-attr-spec* is *access-spec*  
or *proc-language-binding-spec*  
or INTENT ( *intent-spec* )  
or OPTIONAL  
or POINTER  
or PROTECTED  
or SAVE

1 R1216 *proc-decl* is *procedure-entity-name* [ => *proc-pointer-init* ]

2 R1217 *interface-name* is *name*

3 R1218 *proc-pointer-init* is *null-init*  
4 or *initial-proc-target*

5 R1219 *initial-proc-target* is *procedure-name*  
6

7 C1221 (R1217) The *name* shall be the name of an *abstract interface* or of a procedure that has an *explicit*  
8 *interface*. If *name* is declared by a *procedure-declaration-stmt* it shall be previously declared. If *name*  
9 denotes an intrinsic procedure it shall be one that is listed in Table 13.2.

10 C1222 (R1217) The *name* shall not be the same as a keyword that specifies an intrinsic type.

11 C1223 (R1213) If a *proc-interface* describes an *elemental procedure*, each *procedure-entity-name* shall specify an  
12 *external procedure*.

13 C1224 (R1216) If => appears in *proc-decl*, the procedure entity shall have the *POINTER* attribute.

14 C1225 (R1219) The *procedure-name* shall be the name of a *nonelemental external* or *module procedure*, or a  
15 specific intrinsic function listed in Table 13.2.

16 C1226 (R1213) If *proc-language-binding-spec* with a *NAME=* is specified, then *proc-decl-list* shall contain exactly  
17 one *proc-decl*, which shall neither have the *POINTER* attribute nor be a *dummy procedure*.

18 C1227 (R1213) If *proc-language-binding-spec* is specified, the *proc-interface* shall appear, it shall be an *interface-*  
19 *name*, and *interface-name* shall be declared with a *proc-language-binding-spec*.

20 2 If *proc-interface* appears and consists of *interface-name*, it specifies an explicit *specific interface* (12.4.3.2) for the  
21 declared procedure entities. The *abstract interface* (12.4) is that specified by the interface named by *interface-*  
22 *name*. The interface specified by *interface-name* shall not depend on any characteristic of a procedure identified  
23 by a *procedure-entity-name* in the *proc-decl-list* of the same procedure declaration statement.

24 3 If *proc-interface* appears and consists of *declaration-type-spec*, it specifies that the declared procedure entities are  
25 functions having *implicit interfaces* and the specified result type. If a type is specified for an external function,  
26 its function definition (12.6.2.2) shall specify the same result type and type parameters.

27 4 If *proc-interface* does not appear, the procedure declaration statement does not specify whether the declared  
28 procedure entities are subroutines or functions.

29 5 If a *proc-attr-spec* other than a *proc-language-binding-spec* appears, it specifies that the declared procedure entities  
30 have that attribute. These attributes are described in 5.5. If a *proc-language-binding-spec* with *NAME=* appears,  
31 it specifies a *binding label* or its absence, as described in 15.10.2. A *proc-language-binding-spec* without a *NAME=*  
32 is allowed, but is redundant with the *proc-interface* required by C1227.

33 6 If => appears in a *proc-decl* in a *procedure-declaration-stmt* it specifies the initial association status of the  
34 corresponding procedure entity, and implies the *SAVE attribute*, which may be confirmed by explicit specification.  
35 If => *null-init* appears, the procedure entity is initially *disassociated*. If => *initial-proc-target* appears, the  
36 procedure entity is initially associated with the *target*.

37 7 If *procedure-entity-name* has an *explicit interface*, its *characteristics* shall be the same as *initial-proc-target* except  
38 that *initial-proc-target* may be pure even if *procedure-entity-name* is not pure and *initial-proc-target* may be an *elemental*  
39 *intrinsic procedure*.

40 8 If the *characteristics* of *procedure-entity-name* or *initial-proc-target* are such that an *explicit interface* is required,  
41 both *procedure-entity-name* and *initial-proc-target* shall have an *explicit interface*.

42 9 If *procedure-entity-name* has an *implicit interface* and is explicitly typed or referenced as a function, *initial-proc-*

- 1 *target* shall be a function. If *procedure-entity-name* has an **implicit interface** and is referenced as a subroutine,  
 2 *initial-proc-target* shall be a subroutine.
- 3 10 If *initial-proc-target* and *procedure-entity-name* are functions, their results shall have the same **characteristics**.

**NOTE 12.18**

The following code illustrates procedure declaration statements. Note 7.46 illustrates the use of the P and BESSEL defined by this code.

```

ABSTRACT INTERFACE
  FUNCTION REAL_FUNC (X)
    REAL, INTENT (IN) :: X
    REAL :: REAL_FUNC
  END FUNCTION REAL_FUNC
END INTERFACE

INTERFACE
  SUBROUTINE SUB (X)
    REAL, INTENT (IN) :: X
  END SUBROUTINE SUB
END INTERFACE

!-- Some external or dummy procedures with explicit interface.
PROCEDURE (REAL_FUNC) :: BESSEL, GFUN
PROCEDURE (SUB) :: PRINT_REAL
!-- Some procedure pointers with explicit interface,
!-- one initialized to NULL().
PROCEDURE (REAL_FUNC), POINTER :: P, R => NULL()
PROCEDURE (REAL_FUNC), POINTER :: PTR_TO_GFUN
!-- A derived type with a procedure pointer component ...
TYPE STRUCT_TYPE
  PROCEDURE (REAL_FUNC), POINTER, NOPASS :: COMPONENT
END TYPE STRUCT_TYPE
!-- ... and a variable of that type.
TYPE(STRUCT_TYPE) :: STRUCT
!-- An external or dummy function with implicit interface
PROCEDURE (REAL) :: PSI

```

#### 4 12.4.3.8 INTRINSIC statement

- 5 1 An INTRINSIC statement specifies the **INTRINSIC attribute** (5.5.11) for a list of names.

6 R1220 *intrinsic-stmt* is INTRINSIC [ :: ] *intrinsic-procedure-name-list*

7 C1228 (R1220) Each *intrinsic-procedure-name* shall be the name of an intrinsic procedure.

#### 8 12.4.3.9 Implicit interface specification

- 9 1 If the interface of a function is **implicit**, the type and type parameters of the function result are specified by an  
 10 implicit or explicit type specification of the function name. The type, type parameters, and shape of the **dummy**  
 11 **arguments** of a procedure invoked from where the interface of the procedure is **implicit** shall be such that each  
 12 **actual argument** is consistent with the **characteristics** of the corresponding **dummy argument**.

## 12.5 Procedure reference

### 12.5.1 Syntax of a procedure reference

- 1 The form of a procedure reference is dependent on the [interface](#) of the procedure or [procedure pointer](#), but is independent of the means by which the procedure is defined. The forms of procedure references are as follows.

R1221 *function-reference* is *procedure-designator* ( [ *actual-arg-spec-list* ] )

C1229 (R1221) The *procedure-designator* shall designate a function.

C1230 (R1221) The *actual-arg-spec-list* shall not contain an *alt-return-spec*.

R1222 *call-stmt* is CALL *procedure-designator* [ ( [ *actual-arg-spec-list* ] ) ]

C1231 (R1222) The *procedure-designator* shall designate a subroutine.

R1223 *procedure-designator* is *procedure-name*  
or *proc-component-ref*  
or *data-ref* % *binding-name*

C1232 (R1223) A *procedure-name* shall be a generic name or the name of a procedure.

C1233 (R1223) A *binding-name* shall be a [binding name](#) (4.5.5) of the [declared type](#) of *data-ref*.

C1234 (R1223) A *data-ref* shall not be a [polymorphic](#) subobject of a [coindexed object](#).

C1235 (R1223) If *data-ref* is an array, the referenced [type-bound procedure](#) shall have the [PASS attribute](#).

- 2 The *data-ref* in a *procedure-designator* shall not be an unallocated [allocatable](#) variable or a pointer that is not associated.

- 3 Resolving references to [type-bound procedures](#) is described in 12.5.6.

- 4 A function may also be referenced as a [defined operation](#) (7.1.6). A subroutine may also be referenced as a [defined assignment](#) (7.2.1.4, 7.2.1.5), by [defined input/output](#) (9.6.4.8), or by [finalization](#) (4.5.6).

#### NOTE 12.19

When resolving [type-bound procedure](#) references, constraints on the use of [coindexed objects](#) ensure that the [coindexed object](#) (on the remote [image](#)) has the same [dynamic type](#) as the corresponding object on the local [image](#). Thus a processor can resolve the [type-bound procedure](#) using the [coarray](#) variable on its own [image](#) and pass the [coindexed object](#) as the [actual argument](#).

R1224 *actual-arg-spec* is [ *keyword* = ] *actual-arg*

R1225 *actual-arg* is *expr*  
or *variable*  
or *procedure-name*  
or *proc-component-ref*  
or *alt-return-spec*

R1226 *alt-return-spec* is \* *label*

C1236 (R1224) The *keyword* = shall not appear if the [interface](#) of the procedure is [implicit](#).

C1237 (R1224) The *keyword* = shall not be omitted from an *actual-arg-spec* unless it has been omitted from each preceding *actual-arg-spec* in the argument list.

C1238 (R1224) Each *keyword* shall be the name of a [dummy argument](#) in the [explicit interface](#) of the procedure.

- 1 C1239 (R1225) A nonintrinsic [elemental procedure](#) shall not be used as an [actual argument](#).
- 2 C1240 (R1225) A *procedure-name* shall be the name of an [external](#), [internal](#), [module](#), or [dummy](#) procedure, a  
3 specific intrinsic function listed in Table 13.2, or a [procedure pointer](#).
- 4 C1241 (R1225) *expr* shall not be a variable.
- 5 C1242 (R1226) The *label* shall be the [statement label](#) of a [branch target statement](#) that appears in the same [inclusive scope](#) as the  
6 [call-stmt](#).
- 7 C1243 An actual argument that is a [coindexed object](#) shall not have a pointer ultimate component.

**NOTE 12.20**

Examples of procedure reference using [procedure pointers](#):

```
P => BESSEL
WRITE (*, *) P(2.5)      !-- BESSEL(2.5)

S => PRINT_REAL
CALL S(3.14)
```

**NOTE 12.21**

An [internal procedure](#) cannot be invoked using a [procedure pointer](#) from either Fortran or C after the [host instance](#) completes execution, because the pointer is then undefined. While the [host instance](#) is active, however, if an [internal procedure](#) was passed as an [actual argument](#) or is the [target](#) of a [procedure pointer](#), it could be invoked from outside of the host subprogram.

Assume there is a procedure with the following [interface](#) that calculates  $\int_a^b f(x) dx$ .

```
INTERFACE
  FUNCTION INTEGRATE(F, A, B) RESULT(INTEGRAL) BIND(C)
    USE ISO_C_BINDING
    INTERFACE
      FUNCTION F(X) BIND(C) ! Integrand
        USE ISO_C_BINDING
        REAL(C_FLOAT), VALUE :: X
        REAL(C_FLOAT) :: F
      END FUNCTION
    END INTERFACE
    REAL(C_FLOAT), VALUE :: A, B ! Bounds
    REAL(C_FLOAT) :: INTEGRAL
  END FUNCTION INTEGRATE
END INTERFACE
```

This procedure can be called from Fortran or C, and could be written in either Fortran or C. The argument F representing the mathematical function  $f(x)$  can be written as an [internal procedure](#); this [internal procedure](#) will have access to any [host instance local variables](#) necessary to actually calculate  $f(x)$ . For example:

```
REAL FUNCTION MY_INTEGRATION(N, A, B) RESULT(INTEGRAL)
  ! Integrate f(x)=x^n over [a,b]
  USE ISO_C_BINDING
  INTEGER, INTENT(IN) :: N
  REAL, INTENT(IN) :: A, B

  INTEGRAL = INTEGRATE(MY_F, REAL(A, C_FLOAT), REAL(B, C_FLOAT))
  ! This will call the internal function MY_F to calculate f(x).
```

**NOTE 12.21 (cont.)**

! The above interface of INTEGRATE must be explicit and available.

CONTAINS

```
REAL(C_FLOAT) FUNCTION MY_F(X) BIND(C) ! Integrand
  REAL(C_FLOAT), VALUE :: X
  MY_F = X**N ! N is taken from the host instance of MY_INTEGRATION.
END FUNCTION
```

END FUNCTION MY\_INTEGRATION

The function INTEGRATE cannot retain a function pointer to MY\_F and use it after INTEGRATE has finished execution, because the [host instance](#) of MY\_F might no longer exist, making the pointer undefined. If such a pointer is retained, then it can only be used to invoke MY\_F during the execution of the [host instance](#) of MY\_INTEGRATION called from INTEGRATE.

## 12.5.2 Actual arguments, dummy arguments, and argument association

### 12.5.2.1 Argument correspondence

- 1 In either a subroutine reference or a function reference, the [actual argument](#) list identifies the correspondence between the [actual arguments](#) and the [dummy arguments](#) of the procedure. This correspondence may be established either by keyword or by position. If an [argument keyword](#) appears, the [actual argument](#) corresponds to the [dummy argument](#) whose name is the same as the [argument keyword](#) (using the [dummy argument](#) names from the [interface](#) accessible by the [procedure reference](#)). In the absence of an [argument keyword](#), an [actual argument](#) corresponds to the [dummy argument](#) occupying the corresponding position in the reduced [dummy argument](#) list; that is, the first [actual argument](#) corresponds to the first [dummy argument](#) in the reduced list, the second [actual argument](#) corresponds to the second [dummy argument](#) in the reduced list, etc. The reduced [dummy argument](#) list is either the full [dummy argument](#) list or, if there is a [passed-object dummy argument](#) (4.5.4.5), the [dummy argument](#) list with the [passed-object dummy argument](#) omitted. Exactly one [actual argument](#) shall correspond to each nonoptional [dummy argument](#). At most one [actual argument](#) shall correspond to each optional [dummy argument](#). Each [actual argument](#) shall correspond to a [dummy argument](#).

**NOTE 12.22**

For example, the procedure defined by

```
SUBROUTINE SOLVE (FUNCT, SOLUTION, METHOD, STRATEGY, PRINT)
  INTERFACE
    FUNCTION FUNCT (X)
      REAL FUNCT, X
    END FUNCTION FUNCT
  END INTERFACE
  REAL SOLUTION
  INTEGER, OPTIONAL :: METHOD, STRATEGY, PRINT
  ...
```

can be invoked with

```
CALL SOLVE (FUN, SOL, PRINT = 6)
```

provided its [interface](#) is [explicit](#), and if the [interface](#) is specified by an [interface body](#), the name of the last argument is PRINT.

### 12.5.2.2 The passed-object dummy argument and argument correspondence

- 1 In a reference to a [type-bound procedure](#), or a [procedure pointer](#) component, that has a [passed-object dummy argument](#) (4.5.4.5), the [data-ref](#) of the [function-reference](#) or [call-stmt](#) corresponds, as an [actual argument](#), with the [passed-object dummy argument](#).

### 12.5.2.3 Argument association

- 1 Except in references to intrinsic [inquiry functions](#), a pointer [actual argument](#) that corresponds to a nonoptional nonpointer [dummy argument](#) shall be [pointer associated](#) with a [target](#).
- 2 If a nonpointer [dummy argument](#) without the [VALUE attribute](#) corresponds to a pointer [actual argument](#) that is [pointer associated](#) with a [target](#),
- if the [dummy argument](#) is [polymorphic](#), it becomes argument associated with that [target](#);
  - if the [dummy argument](#) is nonpolymorphic, it becomes argument associated with the [declared type](#) part of that [target](#).
- 3 If a present nonpointer [dummy argument](#) without the [VALUE attribute](#) corresponds to a nonpointer [actual argument](#),
- if the [dummy argument](#) is [polymorphic](#), it becomes argument associated with that [actual argument](#);
  - if the [dummy argument](#) is nonpolymorphic, it becomes argument associated with the [declared type](#) part of that [actual argument](#).
- 4 A present [dummy argument](#) with the [VALUE attribute](#) becomes argument associated with a [definable](#) anonymous data object whose initial value is the value of the [actual argument](#).
- 5 A present pointer [dummy argument](#) that corresponds to a pointer [actual argument](#) becomes argument associated with that [actual argument](#). A present pointer [dummy argument](#) that does not correspond to a pointer [actual argument](#) is not argument associated.
- 6 The entity that is argument associated with a [dummy argument](#) is called its [effective argument](#).
- 7 The [ultimate argument](#) is the [effective argument](#) if the [effective argument](#) is not a [dummy argument](#) or a subobject of a [dummy argument](#). If the [effective argument](#) is a dummy argument, the [ultimate argument](#) is the [ultimate argument](#) of that dummy argument. If the [effective argument](#) is a subobject of a dummy argument, the [ultimate argument](#) is the corresponding subobject of the [ultimate argument](#) of that dummy argument.

#### NOTE 12.23

For the sequence of subroutine calls

```

INTEGER :: X(100)
CALL SUBA (X)
...
SUBROUTINE SUBA(A)
  INTEGER :: A(:)
  CALL SUBB (A(1:5), A(5:1:-1))
  ...
SUBROUTINE SUBB(B, C)
  INTEGER :: B(:), C(:)

```

the [ultimate argument](#) of B is X(1:5). The [ultimate argument](#) of C is X(5:1:-1) and this is not the same object as the [ultimate argument](#) of B.

#### NOTE 12.24

Fortran [argument association](#) is usually similar to call by reference and call by value-result. If the [VALUE attribute](#) is specified, the effect is as if the [actual argument](#) is assigned to a temporary, and the temporary is

**NOTE 12.24 (cont.)**

then argument associated with the dummy argument. Subsequent changes to the value or definition status of the dummy argument do not affect the **actual argument**. The actual mechanism by which this happens is determined by the processor.

**12.5.2.4 Ordinary dummy variables**

- 1 The requirements in this subclause apply to **actual arguments** that correspond to nonallocatable nonpointer dummy data objects.
- 2 The **dummy argument** shall be **type compatible** with the **actual argument**. If the **actual argument** is a **polymorphic coindexed object**, the dummy argument shall not be **polymorphic**. If the **actual argument** is a **polymorphic assumed-size array**, the **dummy argument** shall be **polymorphic**. If the **actual argument** is of a derived type that has **type parameters**, **type-bound procedures**, or **final subroutines**, the **dummy argument** shall not be **assumed-type**.
- 3 The **kind type parameter** values of the **actual argument** shall agree with the corresponding ones of the **dummy argument**. The **length type parameter** values of a present **actual argument** shall agree with the corresponding ones of the **dummy argument** that are not assumed, except for the case of the character length parameter of an **actual argument** of type character with default kind or C character kind (15.2.2) associated with a **dummy argument** that is not **assumed-shape** or **assumed-rank**.
- 4 If a present scalar **dummy argument** is of type character with default kind or C character kind, the length *len* of the **dummy argument** shall be less than or equal to the length of the **actual argument**. The **dummy argument** becomes associated with the leftmost *len* characters of the **actual argument**. If a present array **dummy argument** is of type character with default kind or C character kind and is not **assumed-shape** or **assumed-rank**, it becomes associated with the leftmost characters of the **actual argument** element sequence (12.5.2.11).
- 5 The values of **assumed type parameters** of a **dummy argument** are assumed from the corresponding type parameters of its **effective argument**.
- 6 If the **actual argument** is a **coindexed object** with an **allocatable ultimate component**, the dummy argument shall have the **INTENT (IN)** or the **VALUE** attribute.

**NOTE 12.25**

If the **actual argument** is a **coindexed object**, a processor that uses distributed memory might create a copy on the executing **image** of the **actual argument**, including copies of any allocated **allocatable subobjects**, and associate the **dummy argument** with that copy. If necessary, on return from the procedure, the value of the copy would be copied back to the **actual argument**.

- 7 Except in references to intrinsic **inquiry functions**, if the dummy argument is nonoptional and the **actual argument** is **allocatable**, the corresponding **actual argument** shall be allocated.
- 8 If the dummy argument does not have the **TARGET attribute**, any pointers associated with the **effective argument** do not become associated with the corresponding dummy argument on invocation of the procedure. If such a dummy argument is used as an **actual argument** that corresponds to a dummy argument with the **TARGET attribute**, whether any pointers associated with the original **effective argument** become associated with the dummy argument with the **TARGET attribute** is processor dependent.
- 9 If the dummy argument has the **TARGET attribute**, does not have the **VALUE attribute**, and either the **effective argument** is **simply contiguous** or the **dummy argument** is scalar, **assumed-rank**, or **assumed-shape**, and does not have the **CONTIGUOUS attribute**, and the **effective argument** has the **TARGET attribute** but is not a **coindexed object** or an **array section** with a **vector subscript** then
  - any pointers associated with the **effective argument** become associated with the corresponding dummy argument on invocation of the procedure, and

- when execution of the procedure completes, any pointers that do not become undefined (16.5.2.5) and are associated with the dummy argument remain associated with the **effective argument**.

10 If the dummy argument has the **TARGET attribute** and is an **explicit-shape array**, an **assumed-shape array** with the **CONTIGUOUS attribute**, an **assumed-rank object** with the **CONTIGUOUS attribute**, or an **assumed-size array**, and the **effective argument** has the **TARGET attribute** but is not **simply contiguous** and is not an **array section** with a **vector subscript** then

- on invocation of the procedure, whether any pointers associated with the **effective argument** become associated with the corresponding dummy argument is processor dependent, and
- when execution of the procedure completes, the **pointer association** status of any pointer that is **pointer associated** with the **dummy argument** is processor dependent.

11 If the dummy argument has the **TARGET attribute** and the **effective argument** does not have the **TARGET attribute** or is an **array section** with a **vector subscript**, any pointers associated with the dummy argument become undefined when execution of the procedure completes.

12 If the **dummy argument** has the **TARGET attribute** and the **VALUE attribute**, any pointers associated with the **dummy argument** become undefined when execution of the procedure completes.

13 If the **actual argument** is a **coindexed scalar**, the corresponding **dummy argument** shall be scalar. If the **actual argument** is a noncoindexed scalar, the corresponding **dummy argument** shall be scalar unless the **actual argument** is default character, of type character with the C character kind (15.2.2), or is an element or substring of an element of an array that is not an **assumed-shape**, **pointer**, or **polymorphic array**. If the procedure is nonelemental and is referenced by a generic name or as a defined operator or **defined assignment**, the **ranks** of the **actual arguments** and corresponding **dummy arguments** shall agree.

14 If a dummy argument is an **assumed-shape array**, the **rank** of the **actual argument** shall be the same as the **rank** of the dummy argument; the **actual argument** shall not be an **assumed-size array** (including an array element **designator** or an array element substring **designator**).

15 An **actual argument** of any **rank** may correspond to an **assumed-rank** dummy argument. The **rank** and **shape** of the dummy argument are the **rank** and **shape** of the corresponding **actual argument**. If the rank is nonzero, the lower and upper bounds of the dummy argument are those that would be given by the intrinsic functions **LBOUND** and **UBOUND** respectively if applied to the **actual argument**, except that when the **actual argument** is **assumed-size**, the upper bound of the last dimension of the dummy argument is 2 less than the lower bound of that dimension.

16 Except when a procedure reference is **elemental** (12.8), each element of an array **actual argument** or of a sequence in a sequence association (12.5.2.11) is associated with the element of the dummy array that has the same position in array element order (6.5.3.2).

#### NOTE 12.26

For default character sequence associations, the interpretation of element is provided in 12.5.2.11.

17 A scalar dummy argument of a nonelemental procedure shall correspond only to a scalar actual argument.

18 If a dummy argument has **INTENT (OUT)** or **INTENT (INOUT)**, the **actual argument** shall be **definable**. If a dummy argument has **INTENT (OUT)**, the **effective argument** becomes undefined at the time the association is established, except for **direct components** of an object of derived type for which **default initialization** has been specified.

19 If the procedure is nonelemental and the **actual argument** is an **array section** having a **vector subscript**, the dummy argument is not **definable** and shall not have the **ASYNCHRONOUS**, **INTENT (OUT)**, **INTENT (INOUT)**, or **VOLATILE** attributes.

**NOTE 12.27**

Argument intent specifications serve several purposes. See Note 5.17.

**NOTE 12.28**

For more explanatory information on [targets](#) as dummy arguments, see subclause C.9.4.

- 1 C1244 An [actual argument](#) that is a [coindexed object](#) with the [ASYNCHRONOUS](#) or [VOLATILE](#) attribute shall  
2 not correspond to a dummy argument that has either the [ASYNCHRONOUS](#) or [VOLATILE](#) attribute.
- 3 C1245 (R1225) If an [actual argument](#) is a nonpointer array that has the [ASYNCHRONOUS](#) or [VOLATILE](#)  
4 attribute but is not [simply contiguous](#) (6.5.4), and the corresponding dummy argument has either the  
5 [ASYNCHRONOUS](#) or [VOLATILE](#) attribute, that dummy argument shall be [assumed-shape](#) or [assumed-](#)  
6 [rank](#) and shall not have the [CONTIGUOUS](#) attribute.
- 7 C1246 (R1225) If an [actual argument](#) is an array pointer that has the [ASYNCHRONOUS](#) or [VOLATILE](#)  
8 attribute but does not have the [CONTIGUOUS](#) attribute, and the corresponding dummy argument has  
9 either the [ASYNCHRONOUS](#) or [VOLATILE](#) attribute, that dummy argument shall be an array pointer,  
10 an [assumed-shape](#) array without the [CONTIGUOUS](#) attribute, or an [assumed-rank entity](#) without the  
11 [CONTIGUOUS](#) attribute.

**NOTE 12.29**

The constraints on an [actual argument](#) with the [ASYNCHRONOUS](#) or [VOLATILE](#) attribute that corresponds to a dummy argument with either the [ASYNCHRONOUS](#) or [VOLATILE](#) attribute are designed to avoid forcing a processor to use the so-called copy-in/copy-out argument passing mechanism. Making a copy of an [actual argument](#) whose value is likely to change due to an asynchronous input/output operation completing or in some unpredictable manner will cause the new value to be lost when a called procedure returns and the copy-out overwrites the [actual argument](#).

**12.5.2.5 Allocatable and pointer dummy variables**

- 1 The requirements in this subclause apply to an [actual argument](#) with the [ALLOCATABLE](#) or [POINTER](#) attribute  
2 that corresponds to a [dummy argument](#) with the same attribute.
- 3 The [actual argument](#) shall be [polymorphic](#) if and only if the associated [dummy argument](#) is [polymorphic](#), and  
4 either both the [actual](#) and [dummy](#) arguments shall be [unlimited polymorphic](#), or the [declared type](#) of the [actual](#)  
5 [argument](#) shall be the same as the [declared type](#) of the [dummy argument](#).

**NOTE 12.30**

The [dynamic type](#) of a [polymorphic allocatable](#) or pointer dummy argument can change as a result of execution of an [ALLOCATE](#) statement or [pointer assignment](#) in the subprogram. Because of this the corresponding [actual argument](#) needs to be [polymorphic](#) and have a [declared type](#) that is the same as the [declared type](#) of the dummy argument or an [extension](#) of that type. However, type compatibility requires that the [declared type](#) of the dummy argument be the same as, or an [extension](#) of, the type of the [actual argument](#). Therefore, the dummy and [actual](#) arguments need to have the same [declared type](#).

[Dynamic type](#) information is not maintained for a nonpolymorphic [allocatable](#) or pointer dummy argument. However, allocating or pointer-assigning such a dummy argument would require maintenance of this information if the corresponding [actual argument](#) is [polymorphic](#). Therefore, the corresponding [actual argument](#) needs to be nonpolymorphic.

- 3 The [rank](#) of the [actual argument](#) shall be the same as that of the dummy argument, unless the dummy argument  
4 is [assumed-rank](#). The type parameter values of the [actual argument](#) shall agree with the corresponding ones of  
5 the dummy argument that are not [assumed](#) or [deferred](#).
- 6 The [actual argument](#) shall have [deferred](#) the same type parameters as the dummy argument.

1 5 If the **actual argument** is a **coindexed object**, the dummy argument shall have the **INTENT (IN)** attribute.

#### 2 12.5.2.6 Allocatable dummy variables

3 1 The requirements in this subclause apply to **actual arguments** that correspond to **allocatable** dummy data objects.

4 2 The **actual argument** shall be **allocatable**. It is permissible for the **actual argument** to have an allocation status  
5 of unallocated.

6 3 The **corank** of the **actual argument** shall be the same as that of the **dummy argument**.

7 4 The values of **assumed type parameters** of a **dummy argument** are assumed from the corresponding type para-  
8 meters of its **effective argument**.

9 5 If the **dummy argument** does not have the **TARGET attribute**, any pointers associated with the **actual argument**  
10 do not become associated with the corresponding **dummy argument** on invocation of the procedure. If such a  
11 **dummy argument** is used as an **actual argument** that is associated with a **dummy argument** with the **TARGET**  
12 **attribute**, whether any pointers associated with the original **actual argument** become associated with the **dummy**  
13 **argument** with the **TARGET attribute** is processor dependent.

14 6 If the dummy argument has the **TARGET attribute**, does not have the **INTENT (OUT)** or **VALUE** attribute,  
15 and the corresponding **actual argument** has the **TARGET attribute** then

- 16 • any pointers associated with the **actual argument** become associated with the corresponding dummy argu-  
17 ment on invocation of the procedure, and
- 18 • when execution of the procedure completes, any pointers that do not become undefined (16.5.2.5) and are  
19 associated with the dummy argument remain associated with the **actual argument**.

20 7 If a dummy argument has **INTENT (OUT)** or **INTENT (INOUT)**, the **actual argument** shall be **definable**. If a  
21 dummy argument has **INTENT (OUT)** and its associated **actual argument** is allocated, the **actual argument** is  
22 deallocated on procedure invocation (6.7.3.2).

#### 23 12.5.2.7 Pointer dummy variables

24 1 The requirements in this subclause apply to **actual arguments** that correspond to dummy data pointers.

25 C1247 The **actual argument** corresponding to a dummy pointer with the **CONTIGUOUS attribute** shall be  
26 **simply contiguous** (6.5.4).

27 C1248 The **actual argument** corresponding to a dummy pointer shall not be a **coindexed object**.

28 2 If the dummy argument does not have **INTENT (IN)**, the **actual argument** shall be a pointer. Otherwise, the  
29 **actual argument** shall be a pointer or a valid **target** for the dummy pointer in a **pointer assignment statement**. If  
30 the **actual argument** is not a pointer, the dummy pointer becomes **pointer associated** with the **actual argument**.

31 3 The **rank** of the **actual argument** shall be the same as that of the dummy argument. The type parameter values  
32 of the **actual argument** shall agree with the corresponding ones of the dummy argument that are not **assumed**  
33 or **deferred**. The values of **assumed type parameters** of a dummy argument are assumed from the corresponding  
34 type parameters of its **effective argument**.

35 4 If the dummy argument has **INTENT (OUT)**, the **pointer association** status of the **actual argument** becomes  
36 undefined on invocation of the procedure.

37 5 If the dummy argument is nonoptional and the **actual argument** is **allocatable**, the **actual argument** shall be  
38 allocated.

#### NOTE 12.31

For more explanatory information on pointers as dummy arguments, see subclause C.9.4.

### 12.5.2.8 Coarray dummy variables

- 1 If the **dummy argument** is a **coarray**, the corresponding **actual argument** shall be a **coarray** and shall have the **VOLATILE attribute** if and only if the **dummy argument** has the **VOLATILE attribute**.
- 2 If the dummy argument is an array **coarray** that has the **CONTIGUOUS attribute** or is not of **assumed shape**, the corresponding **actual argument** shall be **simply contiguous** or an element of a **simply contiguous** array.

#### NOTE 12.32

Consider the invocation of a procedure on a particular **image**. Each dummy **coarray** is associated with its **ultimate argument** on the **image**. In addition, during this execution of the procedure, this **image** can access the **coarray** corresponding to the **ultimate argument** on any other **image**. For example, consider

```
INTERFACE
  SUBROUTINE SUB(X)
    REAL :: X[*]
  END SUBROUTINE SUB
END INTERFACE
...
REAL :: A(1000)[*]
...
CALL SUB(A(10))
```

During execution of this invocation of SUB, the executing **image** has access through the syntax X[P] to A(10) on **image** P.

#### NOTE 12.33

Each invocation of a procedure with a nonallocatable **coarray dummy argument** establishes a dummy **coarray** for the **image** with its own bounds and **cobounds**. During this execution of the procedure, this **image** can use its own bounds and **cobounds** to access the **coarray** corresponding to the **ultimate argument** on any other **image**. For example, consider

```
INTERFACE
  SUBROUTINE SUB(X,N)
    INTEGER :: N
    REAL :: X(N,N)[N,*]
  END SUBROUTINE SUB
END INTERFACE
...
REAL :: A(1000)[*]
...
CALL SUB(A,10)
```

During execution of this invocation of SUB, the executing **image** has access through the syntax X(1,2)[3,4] to A(11) on the **image** with **image index** 33.

#### NOTE 12.34

The requirements on an **actual argument** that corresponds to a dummy **coarray** that is not of **assumed-shape** or has the **CONTIGUOUS attribute** are designed to avoid forcing a processor to use the so-called copy-in/copy-out argument passing mechanism.

### 12.5.2.9 Actual arguments associated with dummy procedure entities

- 1 If the **interface** of a **dummy procedure** is **explicit**, its **characteristics** as a procedure (12.3.1) shall be the same as those of its **effective argument**, except that a pure **effective argument** may be associated with a dummy argument that is not pure and an **elemental** intrinsic **actual** procedure may be associated with a **dummy procedure** (which cannot be **elemental**).
- 2 If the **interface** of a **dummy procedure** is **implicit** and either the dummy argument is explicitly typed or referenced as a function, it shall not be referenced as a subroutine and any corresponding **actual argument** shall be a function, function **procedure pointer**, or **dummy procedure**. If both the **actual argument** and **dummy argument** are known to be functions, they shall have the same type and type parameters. If only the **dummy argument** is known to be a function, the function that would be invoked by a reference to the **dummy argument** shall have the same type and type parameters, except that an external function with assumed character length may be associated with a dummy argument with explicit character length.
- 3 If the **interface** of a **dummy procedure** is **implicit** and a reference to it appears as a subroutine reference, any corresponding **actual argument** shall be a subroutine, subroutine **procedure pointer**, or **dummy procedure**.
- 4 If a dummy argument is a **dummy procedure** without the **POINTER** attribute, its **effective argument** shall be an **external**, **internal**, **module**, or **dummy** procedure, or a specific intrinsic procedure listed in Table 13.2. If the **specific name** is also a generic name, only the specific procedure is associated with the dummy argument.
- 5 If a dummy argument is a **procedure pointer**, the corresponding **actual argument** shall be a **procedure pointer**, a reference to a function that returns a **procedure pointer**, a reference to the intrinsic function **NULL**, or a valid **target** for the dummy pointer in a **pointer assignment statement**. If the **actual argument** is not a pointer, the dummy argument shall have **INTENT (IN)** and becomes **pointer associated** with the **actual argument**.
- 6 When the **actual argument** is a procedure, the **host instance** of the **dummy argument** is the **host instance** of the **actual argument** (12.6.2.4).
- 7 If an **external procedure** or a **dummy procedure** is used as an **actual argument**, its **interface** shall be **explicit** or it shall be explicitly declared to have the **EXTERNAL** attribute.

### 12.5.2.10 Actual arguments and alternate return indicators

- 1 If a dummy argument is an asterisk (12.6.2.3), the corresponding **actual argument** shall be an alternate return specifier (R1226).

### 12.5.2.11 Sequence association

- 1 An **actual argument** represents an element sequence if it is an array expression, an array element **designator**, a default character scalar, or a scalar of type character with the C character kind (15.2.2). If the **actual argument** is an array expression, the element sequence consists of the elements in array element order. If the **actual argument** is an array element **designator**, the element sequence consists of that array element and each element that follows it in array element order.
- 2 If the **actual argument** is default character or of type character with the C character kind, and is an array expression, array element, or array element substring **designator**, the element sequence consists of the **storage units** beginning with the first **storage unit** of the **actual argument** and continuing to the end of the array. The **storage units** of an array element substring **designator** are viewed as array elements consisting of consecutive groups of **storage units** having the character length of the dummy array.
- 3 If the **actual argument** is default character or of type character with the C character kind, and is a scalar that is not an array element or array element substring **designator**, the element sequence consists of the **storage units** of the **actual argument**.

#### NOTE 12.35

Some of the elements in the element sequence might consist of **storage units** from different elements of the original array.

4 An **actual argument** that represents an element sequence and corresponds to a dummy argument that is an array is sequence associated with the dummy argument if the dummy argument is an **explicit-shape** or **assumed-size** array. The **rank** and shape of the **actual argument** need not agree with the **rank** and shape of the dummy argument, but the number of elements in the dummy argument shall not exceed the number of elements in the element sequence of the **actual argument**. If the dummy argument is **assumed-size**, the number of elements in the dummy argument is exactly the number of elements in the element sequence.

#### 12.5.2.12 Argument presence and restrictions on arguments not present

1 A **dummy argument** or an entity that is **host associated** with a **dummy argument** is not present if the **dummy argument**

- does not correspond to an **actual argument**,
- corresponds to an **actual argument** that is not present, or
- does not have the **ALLOCATABLE** or **POINTER** attribute, and corresponds to an **actual argument** that
  - has the **ALLOCATABLE** attribute and is not allocated, or
  - has the **POINTER** attribute and is **disassociated**;

otherwise, it is present.

2 A nonoptional **dummy argument** shall be present. If an optional nonpointer **dummy argument** corresponds to a present pointer **actual argument**, the **pointer association** status of the **actual argument** shall not be undefined.

3 An optional dummy argument that is not present is subject to the following restrictions.

- (1) If it is a data object, it shall not be referenced or be defined. If it is of a type that has default initialization, the initialization has no effect.
- (2) It shall not be used as the **data-target** or **proc-target** of a **pointer assignment**.
- (3) If it is a procedure or **procedure pointer**, it shall not be invoked.
- (4) It shall not be supplied as an **actual argument** corresponding to a nonoptional dummy argument other than as the argument of the intrinsic function **PRESENT** or as an argument of a function reference that is a **constant expression**.
- (5) A **designator** with it as the **base object** and with one or more subobject selectors shall not be supplied as an **actual argument**.
- (6) If it is an array, it shall not be supplied as an **actual argument** to an **elemental procedure** unless an array of the same **rank** is supplied as an **actual argument** corresponding to a nonoptional dummy argument of that **elemental procedure**.
- (7) If it is a pointer, it shall not be allocated, deallocated, nullified, pointer-assigned, or supplied as an **actual argument** corresponding to an optional nonpointer dummy argument.
- (8) If it is **allocatable**, it shall not be allocated, deallocated, or supplied as an **actual argument** corresponding to an optional nonallocatable **dummy argument**.
- (9) If it has length type parameters, they shall not be the subject of an inquiry.
- (10) It shall not be used as the **selector** in an **ASSOCIATE**, **SELECT RANK**, or **SELECT TYPE** construct.
- (11) It shall not be supplied as the **data-ref** in a **procedure-designator**.
- (12) It shall not be supplied as the **scalar-variable** in a **proc-component-ref**.

4 Except as noted in the list above, it may be supplied as an **actual argument** corresponding to an optional dummy argument, which is then also considered not to be present.

#### 12.5.2.13 Restrictions on entities associated with dummy arguments

1 While an entity is associated with a dummy argument, the following restrictions hold.

- (1) Action that affects the allocation status of the entity or a subobject thereof shall be taken through the dummy argument.

- (2) If the allocation status of the entity or a subobject thereof is affected through the dummy argument, then at any time during the invocation and execution of the procedure, either before or after the allocation or deallocation, it shall be referenced only through the dummy argument.
- (3) Action that affects the value of the entity or any subobject of it shall be taken only through the dummy argument unless
  - (a) the dummy argument has the **POINTER attribute**,
  - (b) the dummy argument is a scalar, **assumed-shape**, or **assumed-rank** object, and has the **TARGET attribute** but not the **INTENT (IN)** or **CONTIGUOUS** attributes, and the **actual argument** is a target other than an **array section** with a **vector subscript**, or
  - (c) the dummy argument is an **assumed-rank** object with the **TARGET attribute** and not the **INTENT (IN) attribute**, and the **actual argument** is a scalar target.
- (4) If the value of the entity or any subobject of it is affected through the dummy argument, then at any time during the invocation and execution of the procedure, either before or after the definition, it may be referenced only through that dummy argument unless
  - (a) the dummy argument has the **POINTER attribute**,
  - (b) the dummy argument is a scalar, **assumed-shape**, or **assumed-rank** object, and has the **TARGET attribute** but not the **INTENT (IN)** or **CONTIGUOUS** attributes, and the **actual argument** is a target other than an **array section** with a **vector subscript**, or
  - (c) the dummy argument is an **assumed-rank** object with the **TARGET attribute** and not the **INTENT (IN) attribute**, and the **actual argument** is a scalar target.

**NOTE 12.36**

In

```

SUBROUTINE OUTER
  REAL, POINTER :: A (:)
  ...
  ALLOCATE (A (1:N))
  ...
  CALL INNER (A)
  ...
CONTAINS
  SUBROUTINE INNER (B)
    REAL :: B (:)
    ...
  END SUBROUTINE INNER
  SUBROUTINE SET (C, D)
    REAL, INTENT (OUT) :: C
    REAL, INTENT (IN) :: D
    C = D
  END SUBROUTINE SET
END SUBROUTINE OUTER

```

an **assignment statement** such as

```
A (1) = 1.0
```

would not be permitted during the execution of INNER because this would be changing A without using B, but statements such as

```
B (1) = 1.0
```

or

**NOTE 12.36 (cont.)**

```
CALL SET (B (1), 1.0)
```

would be allowed. Similarly,

```
DEALLOCATE (A)
```

would not be allowed because this affects the allocation of B without using B. In this case,

```
DEALLOCATE (B)
```

also would not be permitted. If B were declared with the [POINTER attribute](#), either of the statements

```
DEALLOCATE (A)
```

and

```
DEALLOCATE (B)
```

would be permitted, but not both.

**NOTE 12.37**

If there is a partial or complete overlap between the [effective arguments](#) of two different dummy arguments of the same procedure and the dummy arguments have neither the [POINTER](#) nor [TARGET](#) attribute, the overlapped portions shall not be defined, redefined, or become undefined during the execution of the procedure. For example, in

```
CALL SUB (A (1:5), A (3:9))
```

A (3:5) shall not be defined, redefined, or become undefined through the first dummy argument because it is part of the argument associated with the second dummy argument and shall not be defined, redefined, or become undefined through the second dummy argument because it is part of the argument associated with the first dummy argument. A (1:2) remains [definable](#) through the first dummy argument and A (6:9) remains [definable](#) through the second dummy argument.

**NOTE 12.38**

This restriction applies equally to pointer targets. In

```
REAL, DIMENSION (10), TARGET :: A
```

```
REAL, DIMENSION (:), POINTER :: B, C
```

```
B => A (1:5)
```

```
C => A (3:9)
```

```
CALL SUB (B, C) ! The dummy arguments of SUB are neither pointers nor targets.
```

B (3:5) cannot be defined because it is part of the argument associated with the second dummy argument. C (1:3) cannot be defined because it is part of the argument associated with the first dummy argument. A (1:2) [which is B (1:2)] remains [definable](#) through the first dummy argument and A (6:9) [which is C (4:7)] remains [definable](#) through the second dummy argument.

**NOTE 12.39**

In

```
MODULE DATA
```

```
  REAL :: W, X, Y, Z
```

```
END MODULE DATA
```

**NOTE 12.39 (cont.)**

```

PROGRAM MAIN
  USE DATA
  ...
  CALL INIT (X)
  ...
END PROGRAM MAIN
SUBROUTINE INIT (V)
  USE DATA
  ...
  READ (*, *) V
  ...
END SUBROUTINE INIT

```

variable X cannot be directly referenced at any time during the execution of INIT because it is being defined through the dummy argument V. X can be (indirectly) referenced through V. W, Y, and Z can be directly referenced. X can, of course, be directly referenced once execution of INIT is complete.

**NOTE 12.40**

The restrictions on entities associated with dummy arguments are intended to facilitate a variety of optimizations in the translation of the subprogram, including implementations of [argument association](#) in which the value of an [actual argument](#) that is neither a pointer nor a target is maintained in a register or in local storage.

**12.5.3 Function reference**

- 1 A function is invoked during expression evaluation by a [function-reference](#) or by a defined operation (7.1.6). When it is invoked, all [actual argument](#) expressions are evaluated, then the arguments are associated, and then the function is executed. When execution of the function is complete, the value of the function result is available for use in the expression that caused the function to be invoked. The [characteristics](#) of the function result (12.3.3) are determined by the [interface](#) of the function. If a reference to an [elemental](#) function (12.8) is an [elemental reference](#), all array arguments shall have the same shape.

**12.5.4 Subroutine reference**

- 1 A subroutine is invoked by execution of a [CALL statement](#), execution of a [defined assignment statement](#) (7.2.1.4), [defined input/output](#) (9.6.4.8.2), or [finalization](#)(4.5.6). When a subroutine is invoked, all [actual argument](#) expressions are evaluated, then the arguments are associated, and then the subroutine is executed. When the actions specified by the subroutine are completed, the execution of the [CALL statement](#), the execution of the [defined assignment statement](#), the processing of an input or output list item, or [finalization](#) of an object is also completed. If a [CALL statement](#) includes one or more alternate return specifiers among its arguments, a branch to one of the statements indicated might occur, depending on the action specified by the subroutine. If a reference to an [elemental](#) subroutine (12.8) is an [elemental reference](#), at least one [actual argument](#) shall correspond to an [INTENT \(OUT\)](#) or [INTENT \(INOUT\)](#) dummy argument, all such [actual arguments](#) shall be arrays, and all [actual arguments](#) shall be [conformable](#).

**12.5.5 Resolving named procedure references****12.5.5.1 Establishment of procedure names**

- 1 The rules for interpreting a procedure reference depend on whether the procedure name in the reference is established by the available declarations and specifications to be generic in the [scoping unit](#) containing the reference, is established to be only specific in the [scoping unit](#) containing the reference, or is not established.
- 2 A procedure name is established to be generic in a [scoping unit](#)

- (1) if that **scoping unit** contains an **interface block** with that name;
- (2) if that **scoping unit** contains a **GENERIC statement** with a *generic-spec* that is that name;
- (3) if that **scoping unit** contains an **INTRINSIC attribute** specification for that name and it is the generic name of an intrinsic procedure;
- (4) if that **scoping unit** contains a **USE statement** that makes that procedure name accessible and the corresponding name in the module is established to be generic; or
- (5) if that **scoping unit** contains no declarations of that name, that **scoping unit** has a **host scoping unit**, and that name is established to be generic in the **host scoping unit**.

3 A procedure name is established to be only specific in a **scoping unit** if it is established to be specific and not established to be generic. It is established to be specific

- (1) if that **scoping unit** contains a **module subprogram**, **internal subprogram**, or **statement function statement** that defines a procedure with that name;
- (2) if that **scoping unit** is of a subprogram that defines a procedure with that name;
- (3) if that **scoping unit** contains an **INTRINSIC attribute** specification for that name and it is the name of a specific intrinsic procedure;
- (4) if that **scoping unit** contains an explicit **EXTERNAL attribute** specification for that name;
- (5) if that **scoping unit** contains a **USE statement** that makes that procedure name accessible and the corresponding name in the module is established to be specific; or
- (6) if that **scoping unit** contains no declarations of that name, that **scoping unit** has a **host scoping unit**, and that name is established to be specific in the **host scoping unit**.

4 A procedure name is not established in a **scoping unit** if it is neither established to be generic nor established to be specific.

#### 12.5.5.2 Resolving procedure references to names established to be generic

- 1 If the reference is consistent with a nonelemental reference to one of the **specific interfaces** of a **generic interface** that has that name and either is defined in the **scoping unit** in which the reference appears or is made accessible by a **USE statement** in the **scoping unit**, the reference is to the specific procedure in the **interface block** that provides that **interface**. The rules in 12.4.3.5.5 ensure that there can be at most one such specific procedure.
- 2 Otherwise, if the reference is consistent with an **elemental reference** to one of the **specific interfaces** of a **generic interface** that has that name and either is defined in the **scoping unit** in which the reference appears or is made accessible by a **USE statement** in the **scoping unit**, the reference is to the specific **elemental procedure** in the **interface block** that provides that **interface**. The rules in 12.4.3.5.5 ensure that there can be at most one such specific **elemental procedure**.
- 3 Otherwise, if the **scoping unit** contains either an **INTRINSIC attribute** specification for that name or a **USE statement** that makes that name accessible from a module in which the corresponding name is specified to have the **INTRINSIC attribute**, and if the reference is consistent with the **interface** of that **intrinsic** procedure, the reference is to that **intrinsic** procedure.
- 4 Otherwise, if the **scoping unit** has a **host scoping unit**, the name is established to be generic in that **host scoping unit**, and there is agreement between the **scoping unit** and the **host scoping unit** as to whether the name is a function name or a subroutine name, the name is resolved by applying the rules in this subclause to the **host scoping unit** as if the reference appeared there.
- 5 Otherwise, if the name is that of an intrinsic procedure and the reference is consistent with that intrinsic procedure, the reference is to that intrinsic procedure.

**NOTE 12.41**

These rules allow particular specific procedures with the same [generic identifier](#) to be used for particular array [ranks](#) and a general [elemental](#) version to be used for other [ranks](#). For example, given an [interface block](#) such as:

```
INTERFACE RANF
  ELEMENTAL FUNCTION SCALAR_RANF(X)
    REAL, INTENT(IN) :: X
  END FUNCTION SCALAR_RANF
  FUNCTION VECTOR_RANDOM(X)
    REAL X(:)
    REAL VECTOR_RANDOM(SIZE(X))
  END FUNCTION VECTOR_RANDOM
END INTERFACE RANF
```

and a declaration such as:

```
REAL A(10,10), AA(10,10)
```

then the statement

```
A = RANF(AA)
```

is an [elemental reference](#) to SCALAR\_RANF. The statement

```
A(6:10,2) = RANF(AA(6:10,2))
```

is a [nonelemental reference](#) to VECTOR\_RANDOM.

**NOTE 12.42**

In the [USE statement](#) case, it is possible, because of the renaming facility, for the name in the reference to be different from the name of the intrinsic procedure.

### 1 12.5.5.3 Resolving procedure references to names established to be only specific

- 2 1 If the name has the [EXTERNAL attribute](#),
  - 3 • if it is a [procedure pointer](#), the reference is to its target;
  - 4 • if it is a [dummy procedure](#) that is not a [procedure pointer](#), the reference is to the [effective argument](#) corresponding to that name;
  - 5 • otherwise, the reference is to the [external procedure](#) with that name.
- 7 2 If the name is that of an accessible [external procedure](#), [internal procedure](#), [module procedure](#), [intrinsic procedure](#),
  - 8 or statement function, the reference is to that procedure.

**NOTE 12.43**

Because of the renaming facility of the [USE statement](#), the name in the reference can be different from the original name of the procedure.

### 9 12.5.5.4 Resolving procedure references to names not established

- 10 1 If the name is the name of a dummy argument of the [scoping unit](#), the dummy argument is a [dummy procedure](#)
 and the reference is to that [dummy procedure](#). That is, the procedure invoked by executing that reference is the
 [effective argument](#) corresponding to that [dummy procedure](#).
  - 11
  - 12

- 1 2 Otherwise, if the name is the name of an intrinsic procedure, and if there is agreement between the reference and  
 2 the status of the intrinsic procedure as being a function or subroutine, the reference is to that intrinsic procedure.
- 3 3 Otherwise, the reference is to an **external procedure** with that name.

## 4 12.5.6 Resolving type-bound procedure references

- 5 1 If the *binding-name* in a *procedure-designator* (R1223) is that of a specific **type-bound procedure**, the procedure  
 6 referenced is the one bound to that name in the **dynamic type** of the *data-ref*.
- 7 2 If the *binding-name* in a *procedure-designator* is that of a generic type bound procedure, the generic **binding** with  
 8 that name in the **declared type** of the *data-ref* is used to select a specific **binding** using the following criteria.
- 9 • If the reference is consistent with one of the specific **bindings** of that generic **binding**, that specific **binding**  
 10 is selected.
  - 11 • Otherwise, the reference shall be consistent with an **elemental reference** to one of the specific **bindings** of  
 12 that generic **binding**; that specific **binding** is selected.
- 13 3 The reference is to the procedure bound to the same name as the selected specific **binding** in the **dynamic type**  
 14 of the *data-ref*.

## 15 12.6 Procedure definition

### 16 12.6.1 Intrinsic procedure definition

- 17 1 Intrinsic procedures are defined as an inherent part of the processor. A standard-conforming processor shall  
 18 include the intrinsic procedures described in Clause 13, but may include others. However, a standard-conforming  
 19 program shall not make use of intrinsic procedures other than those described in Clause 13.

### 20 12.6.2 Procedures defined by subprograms

#### 21 12.6.2.1 General

- 22 1 A procedure is defined by the initial **SUBROUTINE** or **FUNCTION** statement of a subprogram, and each **ENTRY**  
 23 **statement** defines an additional procedure (12.6.2.6).
- 24 2 A subprogram is specified to be **elemental** (12.8), **pure** (12.7), **recursive**, or a separate module subprogram  
 25 (12.6.2.5) by a *prefix* in its initial **SUBROUTINE** or **FUNCTION** statement.

26 R1227 *prefix* is *prefix-spec* [ *prefix-spec* ] ...

27 R1228 *prefix-spec* is *declaration-type-spec*  
 28 or ELEMENTAL  
 29 or IMPURE  
 30 or MODULE  
 31 or NON\_RECURSIVE  
 32 or PURE  
 33 or RECURSIVE

34 C1249 (R1227) A *prefix* shall contain at most one of each *prefix-spec*.

35 C1250 (R1227) A *prefix* shall not specify both PURE and IMPURE.

36 C1251 (R1227) A *prefix* shall not specify both NON\_RECURSIVE and RECURSIVE.

37 C1252 An **elemental procedure** shall not have the **BIND attribute**.

- 1 C1253 (R1227) **MODULE** shall appear only in the *function-stmt* or *subroutine-stmt* of a module subprogram or  
2 of a nonabstract *interface body* that is declared in the *scoping unit* of a module or submodule.
- 3 C1254 (R1227) If **MODULE** appears in the *prefix* of a module subprogram, it shall have been declared to be a  
4 separate *module procedure* in the containing program unit or an ancestor of that program unit.
- 5 C1255 (R1227) If **MODULE** appears in the *prefix* of a module subprogram, the subprogram shall specify the  
6 same *characteristics* and dummy argument names as its corresponding module procedure *interface body*.
- 7 C1256 (R1227) If **MODULE** appears in the *prefix* of a module subprogram and a *binding label* is specified, it  
8 shall be the same as the *binding label* specified in the corresponding module procedure *interface body*.
- 9 C1257 (R1227) If **MODULE** appears in the *prefix* of a module subprogram, **NON\_RECURSIVE** shall appear  
10 if and only if **NON\_RECURSIVE** appears in the *prefix* in the corresponding module procedure *interface*  
11 *body*.
- 12 3 The **NON\_RECURSIVE** *prefix-spec* shall not appear if any procedure defined by the subprogram directly or  
13 indirectly invokes itself or any other procedure defined by the subprogram.
- 14 4 If the *prefix-spec* **PURE** appears, or the *prefix-spec* **ELEMENTAL** appears and **IMPURE** does not appear, the  
15 subprogram is a pure subprogram and shall meet the additional constraints of 12.7.
- 16 5 If the *prefix-spec* **ELEMENTAL** appears, the subprogram is an *elemental subprogram* and shall meet the additional  
17 constraints of 12.8.1.
- 18 **12.6.2.2 Function subprogram**
- 19 1 A function subprogram is a subprogram that has a **FUNCTION** statement as its first statement.
- 20 R1229 *function-subprogram* is *function-stmt*  
21 [ *specification-part* ]  
22 [ *execution-part* ]  
23 [ *internal-subprogram-part* ]  
24 *end-function-stmt*
- 25 R1230 *function-stmt* is [ *prefix* ] **FUNCTION** *function-name* ■  
26 ■ ( [ *dummy-arg-name-list* ] ) [ *suffix* ]
- 27 C1258 (R1230) If **RESULT** appears, *result-name* shall not be the same as *function-name* and shall not be the same  
28 as the *entry-name* in any **ENTRY** statement in the subprogram.
- 29 C1259 (R1230) If **RESULT** appears, the *function-name* shall not appear in any specification statement in the  
30 *scoping unit* of the function subprogram.
- 31 R1231 *proc-language-binding-spec* is *language-binding-spec*
- 32 C1260 (R1231) A *proc-language-binding-spec* with a **NAME=** specifier shall not be specified in the *function-stmt*  
33 or *subroutine-stmt* of an *internal procedure*, or of an *interface body* for an *abstract interface* or a *dummy*  
34 *procedure*.
- 35 C1261 If *proc-language-binding-spec* is specified for a procedure, each of its dummy arguments shall be an *inter-*  
36 *operable* procedure (15.3.7) or a variable that is *interoperable* (15.3.5, 15.3.6), *assumed-shape*, *assumed-*  
37 *rank*, *assumed-type*, of type **CHARACTER** with assumed length, or that has the **ALLOCATABLE** or  
38 **POINTER** attribute.
- 39 C1262 If *proc-language-binding-spec* is specified for a function, the function result shall be an *interoperable* scalar  
40 variable.
- 41 C1263 A variable that is a dummy argument of a procedure that has a *proc-language-binding-spec* shall be  
42 assumed-type or of *interoperable* type.

- 1 R1232 *dummy-arg-name* is *name*
- 2 C1264 (R1232) A *dummy-arg-name* shall be the name of a dummy argument.
- 3 R1233 *suffix* is *proc-language-binding-spec* [ RESULT ( *result-name* ) ]  
 4 or RESULT ( *result-name* ) [ *proc-language-binding-spec* ]
- 5 R1234 *end-function-stmt* is END [ FUNCTION [ *function-name* ] ]
- 6 C1265 (R1229) An internal function subprogram shall not contain an *internal-subprogram-part*.
- 7 C1266 (R1234) If a *function-name* appears in the *end-function-stmt*, it shall be identical to the *function-name*  
 8 specified in the *function-stmt*.
- 9 2 The name of the function is *function-name*.
- 10 3 The type and type parameters (if any) of the result of the function defined by a function subprogram may be  
 11 specified by a type specification in the FUNCTION statement or by the name of the *function result* appearing  
 12 in a *type declaration statement* in the specification part of the function subprogram. They shall not be specified  
 13 both ways. If they are not specified either way, they are determined by the implicit typing rules in effect within  
 14 the function subprogram. If the *function result* is an array, *allocatable*, or a pointer, this shall be specified by  
 15 specifications of the name of the *function result* within the function body. The specifications of the *function*  
 16 *result* attributes, the specification of dummy argument attributes, and the information in the procedure heading  
 17 collectively define the *characteristics* of the function (12.3.1).
- 18 4 If RESULT appears, the name of the *function result* of the function is *result-name* and all occurrences of the  
 19 function name in *execution-part* statements in its scope refer to the function itself. If RESULT does not appear,  
 20 the name of the *function result* is *function-name* and all occurrences of the function name in *execution-part*  
 21 statements in its scope are references to the *function result*. On completion of execution of the function, the value  
 22 returned is that of its *function result*. If the *function result* is a data pointer, the shape of the value returned by  
 23 the function is determined by the shape of the *function result* when the execution of the function is completed. If  
 24 the *function result* is not a pointer, its value shall be defined by the function. If the *function result* is a pointer,  
 25 on return the *pointer association* status of the *function result* shall not be undefined.

**NOTE 12.44**

The *function result* is similar to any other entity (variable or procedure pointer) local to a function subprogram. Its existence begins when execution of the function is initiated and ends when execution of the function is terminated. However, because the final value of this entity is used subsequently in the evaluation of the expression that invoked the function, an implementation might defer releasing the storage occupied by that entity until after its value has been used in expression evaluation.

**NOTE 12.45**

The following is an example of the declaration of an *interface body* with the *BIND attribute*, and a reference to the procedure declared.

```
USE, INTRINSIC :: ISO_C_BINDING

INTERFACE
  FUNCTION JOE (I, J, R) BIND(C,NAME="FrEd")
    USE, INTRINSIC :: ISO_C_BINDING
    INTEGER(C_INT) :: JOE
    INTEGER(C_INT), VALUE :: I, J
    REAL(C_FLOAT), VALUE :: R
  END FUNCTION JOE
END INTERFACE

INT = JOE(1_C_INT, 3_C_INT, 4.0_C_FLOAT)
```

**NOTE 12.45 (cont.)**

END PROGRAM

The invocation of the function JOE results in a reference to a function with a [binding label](#) "FrEd". FrEd could be a C function described by the C prototype

```
int FrEd(int n, int m, float x);
```

**12.6.2.3 Subroutine subprogram**

1 A subroutine subprogram is a subprogram that has a SUBROUTINE statement as its first statement.

R1235 *subroutine-subprogram* is *subroutine-stmt*  
   [ *specification-part* ]  
   [ *execution-part* ]  
   [ *internal-subprogram-part* ]  
   *end-subroutine-stmt*

R1236 *subroutine-stmt* is [ *prefix* ] SUBROUTINE *subroutine-name* ■  
   ■ [ ( [ *dummy-arg-list* ] ) [ *proc-language-binding-spec* ] ]

C1267 (R1236) The *prefix* of a *subroutine-stmt* shall not contain a *declaration-type-spec*.

R1237 *dummy-arg* is *dummy-arg-name*  
   or \*

R1238 *end-subroutine-stmt* is END [ SUBROUTINE [ *subroutine-name* ] ]

C1268 (R1235) An internal subroutine subprogram shall not contain an *internal-subprogram-part*.

C1269 (R1238) If a *subroutine-name* appears in the *end-subroutine-stmt*, it shall be identical to the *subroutine-name* specified in the *subroutine-stmt*.

2 The name of the subroutine is *subroutine-name*.

**12.6.2.4 Instances of a subprogram**

1 When a procedure defined by a subprogram is invoked, an instance of that subprogram is created. Each instance has an independent sequence of execution and an independent set of dummy arguments, [unsaved local variables](#), and [unsaved local procedure pointers](#). Saved local entities are shared by all instances of the subprogram.

2 When a statement function is invoked, an instance of that statement function is created.

3 When execution of an instance completes it ceases to exist.

4 The caller of an instance of a procedure is the instance of the main program, subprogram, or statement function that invoked it. The call sequence of an instance of a procedure is its caller, followed by the call sequence of its caller. The call sequence of the main program is empty. The [host instance](#) of an instance of a statement function or an [internal procedure](#) that is invoked by its name is the first element of the call sequence that is an instance of the host of the statement function or internal subprogram. The [host instance](#) of an [internal procedure](#) that is invoked via a [dummy procedure](#) or [procedure pointer](#) is the [host instance](#) of the [associating entity](#) from when the [argument association](#) or [pointer association](#) was established (16.5.5). The [host instance](#) of a [module procedure](#) is the module or submodule in which it is defined. A main program or external subprogram has no [host instance](#).

**12.6.2.5 Separate module procedures**

1 A separate module procedure is a [module procedure](#) defined by a *separate-module-subprogram*, by a *function-subprogram* whose initial statement contains the keyword MODULE, or by a *subroutine-subprogram* whose initial

statement contains the keyword MODULE.

R1239 *separate-module-subprogram* **is** *mp-subprogram-stmt*  
   [ *specification-part* ]  
   [ *execution-part* ]  
   [ *internal-subprogram-part* ]  
   *end-mp-subprogram-stmt*

R1240 *mp-subprogram-stmt*               **is** MODULE PROCEDURE *procedure-name*

R1241 *end-mp-subprogram-stmt*       **is** END [PROCEDURE [*procedure-name*]]

C1270 (R1239) The *procedure-name* shall have been declared to be a separate **module procedure** in the containing **program unit** or an ancestor of that **program unit**.

C1271 (R1241) If a *procedure-name* appears in the *end-mp-subprogram-stmt*, it shall be identical to the *procedure-name* in the *mp-subprogram-stmt*.

2 A separate **module procedure** shall not be defined more than once.

3 The **interface** of a procedure defined by a *separate-module-subprogram* is explicitly declared by the *mp-subprogram-stmt* to be the same as its module procedure **interface body**. It is recursive if and only if it is declared to be recursive by the **interface body**, and if it is a function its result name is determined by the **FUNCTION statement** in the **interface body**.

#### NOTE 12.46

A separate **module procedure** can be accessed by **use association** only if its **interface body** is declared in the specification part of a module and is public.

#### 12.6.2.6 ENTRY statement

1 An ENTRY statement permits a procedure reference to begin with a particular executable statement within the function or subroutine subprogram in which the ENTRY statement appears.

R1242 *entry-stmt*                               **is** ENTRY *entry-name* [ ( [ *dummy-arg-list* ] ) [ *suffix* ] ]

C1272 (R1242) If **RESULT** appears, the *entry-name* shall not appear in any specification or **type declaration statement** in the **scoping unit** of the function program.

C1273 (R1242) An *entry-stmt* shall appear only in an *external-subprogram* or a *module-subprogram* that does not define a separate **module procedure**. An *entry-stmt* shall not appear within an *executable-construct*.

C1274 (R1242) **RESULT** shall appear only if the *entry-stmt* is in a function subprogram.

C1275 (R1242) A *dummy-arg* shall not be an alternate return indicator if the ENTRY statement is in a function subprogram.

C1276 (R1242) If **RESULT** appears, *result-name* shall not be the same as the *function-name* in the **FUNCTION statement** and shall not be the same as the *entry-name* in any ENTRY statement in the subprogram.

2 Optionally, a subprogram may have one or more ENTRY statements.

3 If the ENTRY statement is in a function subprogram, an additional function is defined by that subprogram. The name of the function is *entry-name* and the name of its **result** is *result-name* or is *entry-name* if no *result-name* is provided. The **dummy arguments** of the function are those specified in the ENTRY statement. If the **characteristics** of the result of the function named in the ENTRY statement are the same as the **characteristics** of the result of the function named in the **FUNCTION statement**, their **result** names identify the same entity, although their names need not be the same. Otherwise, they are **storage associated** and shall all be nonpointer, nonallocatable scalar variables that are default integer, default real, double precision real, default complex, or default logical.

4 If the ENTRY statement is in a subroutine subprogram, an additional subroutine is defined by that subprogram. The name of the subroutine is *entry-name*. The dummy arguments of the subroutine are those specified in the ENTRY statement.

- 1 5 The order, number, types, kind type parameters, and names of the dummy arguments in an ENTRY statement may differ from the  
 2 order, number, types, kind type parameters, and names of the dummy arguments in the **FUNCTION** or **SUBROUTINE** statement  
 3 in the containing subprogram.
- 4 6 Because an ENTRY statement defines an additional function or an additional subroutine, it is referenced in the same manner as any  
 5 other function or subroutine (12.5).
- 6 7 In a subprogram, a dummy argument specified in an ENTRY statement shall not appear in an executable statement preceding that  
 7 ENTRY statement, unless it also appears in a **FUNCTION**, **SUBROUTINE**, or ENTRY statement that precedes the executable  
 8 statement. A function result specified by a *result-name* in an ENTRY statement shall not appear in any executable statement that  
 9 precedes the first **RESULT** clause with that name.
- 10 8 In a subprogram, a **dummy argument** specified in an ENTRY statement shall not appear in the expression of a **statement function**  
 11 that precedes the first *dummy-arg* with that name in the subprogram. A function result specified by a *result-name* in an ENTRY  
 12 statement shall not appear in the expression of a **statement function** that precedes the first **RESULT** clause with that name.
- 13 9 If a **dummy argument** appears in an executable statement, the execution of the executable statement is permitted during the  
 14 execution of a reference to the function or subroutine only if the **dummy argument** appears in the **dummy argument** list of the  
 15 referenced procedure.
- 16 10 If a **dummy argument** is used in a **specification expression** to specify an array bound or character length of an object, the appearance  
 17 of the object in a statement that is executed during a procedure reference is permitted only if the **dummy argument** appears in the  
 18 **dummy argument** list of the referenced procedure and it is present (12.5.2.12).
- 19 11 The NON\_RECURSIVE and RECURSIVE keywords are not used in an ENTRY statement. Instead, the presence or absence of  
 20 NON\_RECURSIVE in the initial **SUBROUTINE** or **FUNCTION** statement controls whether the procedure defined by an ENTRY  
 21 statement is permitted to reference itself or another procedure defined by the subprogram.
- 22 12 The keywords **PURE** and **IMPURE** are not used in an ENTRY statement. Instead, the procedure defined by an ENTRY statement  
 23 is pure if and only if the subprogram is a pure subprogram.
- 24 13 The keyword **ELEMENTAL** is not used in an ENTRY statement. Instead, the procedure defined by an ENTRY statement is **elemental**  
 25 if and only if **ELEMENTAL** is specified in the **SUBROUTINE** or **FUNCTION** statement.

### 12.6.2.7 RETURN statement

R1243 *return-stmt* is RETURN [ *scalar-int-expr* ]

C1277 (R1243) The *return-stmt* shall be in the **inclusive scope** of a function or subroutine subprogram.

C1278 (R1243) The *scalar-int-expr* is allowed only in the **inclusive scope** of a subroutine subprogram.

- 1 Execution of the RETURN statement completes execution of the instance of the subprogram in which it appears.  
 If the expression appears and has a value  $n$  between 1 and the number of asterisks in the dummy argument list, the **CALL statement**  
 that invoked the subroutine branches (8.2) to the **branch target statement** identified by the  $n^{th}$  alternate return specifier in the **actual**  
**argument** list of the referenced procedure. If the expression is omitted or has a value outside the required range, there is no transfer  
 of control to an alternate return.
- 2 Execution of an *end-function-stmt*, *end-mp-subprogram-stmt*, or *end-subroutine-stmt* is equivalent to execution  
 of a RETURN statement with no expression.

### 12.6.2.8 CONTAINS statement

R1244 *contains-stmt* is CONTAINS

- 1 The CONTAINS statement separates the body of a main program, module, submodule, or subprogram from any  
 internal or module subprograms it might contain, or it introduces the type-bound procedure part of a derived-type  
 definition (4.5.5). The CONTAINS statement is not executable.

### 12.6.3 Definition and invocation of procedures by means other than Fortran

- 1 A procedure may be defined by means other than Fortran. The [interface](#) of a procedure defined by means other than Fortran may be specified by an [interface body](#) or [procedure declaration statement](#). A reference to such a procedure is made as though it were defined by an external subprogram.
- 2 A procedure defined by means other than Fortran that is invoked by a Fortran procedure and does not cause [termination of execution](#) shall return to its caller.

#### NOTE 12.47

Examples of code that might cause a transfer of control that bypasses the normal return mechanism of a Fortran procedure are `setjmp` and `longjmp` in C and exception handling in other languages. No such behavior is permitted by this part of ISO/IEC 1539.

- 3 If the [interface](#) of a procedure has a [proc-language-binding-spec](#), the procedure is [interoperable](#) (15.10).
- 4 Interoperation with C functions is described in 15.10.

#### NOTE 12.48

For explanatory information on definition of procedures by means other than Fortran, see subclause C.9.2.

### 12.6.4 Statement function

- 1 A statement function is a function defined by a single statement.

R1245 *stmt-function-stmt*                      is    *function-name* ( [ *dummy-arg-name-list* ] ) = *scalar-expr*

(R1245) Each [primary](#) in *scalar-expr* shall be a constant (literal or named), a reference to a variable, a reference to a function, or an expression in parentheses. Each operation shall be intrinsic. If *scalar-expr* contains a reference to a function, the reference shall not require an [explicit interface](#), the function shall not require an [explicit interface](#) unless it is an intrinsic function, the function shall not be a [transformational](#) intrinsic, and the result shall be scalar. If an argument to a function is an array, it shall be an array name. If a reference to a statement function appears in *scalar-expr*, its definition shall have been provided earlier in the [scoping unit](#) and shall not be the name of the statement function being defined.

(R1245) [Named constants](#) in *scalar-expr* shall have been declared earlier in the [scoping unit](#) or made accessible by use or [host](#) association. If array elements appear in *scalar-expr*, the array shall have been declared as an array earlier in the [scoping unit](#) or made accessible by use or [host](#) association.

(R1245) If a *dummy-arg-name*, variable, function reference, or dummy function reference is typed by the implicit typing rules, its appearance in any subsequent [type declaration statement](#) shall confirm this implied type and the values of any implied type parameters.

(R1245) The *function-name* and each *dummy-arg-name* shall be specified, explicitly or implicitly, to be scalar.

(R1245) A given *dummy-arg-name* shall not appear more than once in any *dummy-arg-name-list*.

- 2 The definition of a statement function with the same name as an accessible entity from the [host](#) shall be preceded by the declaration of its type in a [type declaration statement](#).
- 3 The dummy arguments have a scope of the statement function statement. Each dummy argument has the same type and type parameters as the entity of the same name in the [scoping unit](#) containing the statement function statement.
- 4 A statement function shall not be supplied as an actual argument.
- 5 Execution of a statement function consists of evaluating the expression using the values of the [actual arguments](#) for the values of the corresponding dummy arguments and, if necessary, converting the result to the [declared type](#) and type parameters of the function.
- 6 A function reference in the scalar expression shall not cause a dummy argument of the statement function to become redefined or undefined.

## 12.7 Pure procedures

1 A **pure procedure** is

- a pure intrinsic procedure (13.1),
- a **module procedure** in an intrinsic module, if it is specified to be pure,
- defined by a pure subprogram,
- a **dummy procedure** that has been specified to be **PURE**, or
- a statement function that references only pure functions.

2 A pure subprogram is a subprogram that has the *prefix-spec* **PURE** or that has the *prefix-spec* **ELEMENTAL** and does not have the *prefix-spec* **IMPURE**. The following additional constraints apply to pure subprograms.

C1284 The *specification-part* of a pure function subprogram shall specify that all its nonpointer dummy data objects have the **INTENT (IN)** or the **VALUE** attribute.

C1285 The function result of a pure function shall not be such that **finalization** of a reference to the function would reference an impure procedure.

C1286 A pure function shall not have a **polymorphic** allocatable result.

C1287 The *specification-part* of a pure subroutine subprogram shall specify the intents of all its nonpointer dummy data objects that do not have the **VALUE** attribute.

C1288 An **INTENT (OUT)** **dummy argument** of a pure procedure shall not be such that **finalization** of the actual argument would reference an impure procedure.

C1289 An **INTENT (OUT)** **dummy argument** of a pure procedure shall not be **polymorphic**.

C1290 A local variable of a pure subprogram, or of a **BLOCK construct** within a pure subprogram, shall not have the **SAVE** attribute.

### NOTE 12.49

Variable initialization in a *type-declaration-stmt* or a *data-stmt* implies the **SAVE** attribute; therefore, such initialization is also disallowed.

C1291 The *specification-part* of a pure subprogram shall specify that all its **dummy procedures** are pure.

C1292 If a procedure that is neither an intrinsic procedure nor a statement function is used in a context that requires it to be pure, then its **interface** shall be **explicit** in the scope of that use. The **interface** shall specify that the procedure is pure.

C1293 All internal subprograms in a pure subprogram shall be pure.

C1294 A **designator** of a variable with the **VOLATILE** attribute shall not appear in a pure subprogram.

C1295 In a pure subprogram any **designator** with a **base object** that is in common or accessed by **host** or **use** association, is a **pointer dummy argument** of a pure function, is a **dummy argument** with the **INTENT (IN)** attribute, is a **coindexed object**, or an object that is **storage associated** with any such variable, shall not be used

- (1) in a variable definition context (16.6.7),
- (2) in a pointer association context (16.6.8),
- (3) as the *data-target* in a *pointer-assignment-stmt*,
- (4) as the *expr* corresponding to a component with the **POINTER** attribute in a *structure-constructor*,
- (5) as the *expr* of an **intrinsic assignment statement** in which the variable is of a derived type if the derived type has a pointer component at any level of component selection,

- (6) as the *source-expr* in a **SOURCE= specifier** if the designator is of a derived type that has a pointer component at any level of component selection,
- (7) as an **actual argument** corresponding to a dummy argument with the **POINTER attribute**, or
- (8) as the **actual argument** to the function **C\_LOC** from the intrinsic module **ISO\_C\_BINDING**.

**NOTE 12.50**

Item 5 requires that processors be able to determine if entities with the **PRIVATE attribute** or with private components have a pointer component.

- C1296 Any procedure referenced in a pure subprogram, including one referenced via a **defined operation**, **defined assignment**, **defined input/output**, or **finalization**, shall be pure.
- C1297 A statement that might result in the deallocation of a **polymorphic** entity is not permitted in a pure procedure.

**NOTE 12.51**

This includes **intrinsic assignment** to a variable that has a **potential subobject component** that is **polymorphic** and **allocatable**.

- C1298 A pure subprogram shall not contain a *print-stmt*, *open-stmt*, *close-stmt*, *backspace-stmt*, *endfile-stmt*, *rewind-stmt*, *flush-stmt*, *wait-stmt*, or *inquire-stmt*.
- C1299 A pure subprogram shall not contain a *read-stmt* or *write-stmt* whose *io-unit* is a *file-unit-number* or \*.
- C12100 A pure subprogram shall not contain an **image control statement** (8.6.1).

**NOTE 12.52**

The above constraints are designed to guarantee that a **pure procedure** is free from side effects (modifications of data visible outside the procedure), which means that it is safe to reference it in constructs such as **DO CONCURRENT** and **FORALL**, where there is no explicit order of evaluation.

The constraints on pure subprograms appear to be complicated, but it is not necessary for a programmer to be intimately familiar with them. From the programmer's point of view, these constraints can be summarized as follows: a pure subprogram shall not contain any operation that could conceivably result in an assignment or **pointer assignment** to a common variable, a variable accessed by use or **host** association, or an **INTENT (IN)** dummy argument; nor shall a pure subprogram contain any operation that could conceivably perform any **external file** input/output or **STOP** operation. Note the use of the word conceivably; it is not sufficient for a pure subprogram merely to be side-effect free in practice. For example, a function that contains an assignment to a global variable but in a block that is not executed in any invocation of the function is nevertheless not a pure function. The exclusion of functions of this nature is required if strict compile-time checking is to be used.

It is expected that most library procedures will conform to the constraints required of **pure procedures**, and so can be declared pure and referenced in **DO CONCURRENT constructs**, **FORALL** statements and constructs, and within user-defined **pure procedures**.

**NOTE 12.53**

Pure subroutines are included to allow subroutine calls from **pure procedures** in a safe way, and to allow *forall-assignment-stmts* to be **defined assignments**. The constraints for pure subroutines are based on the same principles as for pure functions, except that side effects to **INTENT (OUT)**, **INTENT (INOUT)**, and pointer dummy arguments are permitted.

## 12.8 Elemental procedures

### 12.8.1 Elemental procedure declaration and interface

- 1 An [elemental procedure](#) is an elemental intrinsic procedure or a procedure that is defined by an [elemental subprogram](#). An [elemental procedure](#) has only [scalar dummy arguments](#), but may have array [actual arguments](#).
- 2 An [elemental subprogram](#) has the *prefix-spec* **ELEMENTAL**. An [elemental subprogram](#) is a pure subprogram unless it has the *prefix-spec* **IMPURE**. The following additional constraints apply to [elemental subprograms](#).
  - C12101 All [dummy arguments](#) of an [elemental procedure](#) shall be scalar noncoarray dummy data objects and shall not have the **POINTER** or **ALLOCATABLE** attribute.
  - C12102 The [result](#) of an [elemental](#) function shall be scalar, and shall not have the **POINTER** or **ALLOCATABLE** attribute.
  - C12103 The *specification-part* of an [elemental subprogram](#) shall specify the intents of all of its [dummy arguments](#) that do not have the **VALUE** attribute.
  - C12104 In the *specification-expr* that specifies a [type parameter](#) value of the result of an elemental function, an [object designator](#) with a [dummy argument](#) of the function as the [base object](#) shall appear only as the subject of a specification inquiry (7.1.11), and that specification inquiry shall not depend on a property that is deferred.
- 3 In a reference to an elemental procedure, if any argument is an array, each [actual argument](#) that corresponds to an **INTENT (OUT)** or **INTENT (INOUT)** [dummy argument](#) shall be an array. All [actual arguments](#) shall be [conformable](#). An array [actual argument](#) is considered to be associated with the scalar [dummy arguments](#) of the procedure throughout the entire execution of the elemental reference; thus, the restrictions on actions specified in 12.5.2.13 apply to the entirety of the [actual array argument](#).

### 12.8.2 Elemental function actual arguments and results

- 1 If a generic name or a [specific name](#) is used to reference an [elemental](#) function, the shape of the result is the same as the shape of the [actual argument](#) with the greatest [rank](#). If there are no [actual arguments](#) or the [actual arguments](#) are all scalar, the result is scalar. In the array case, the values of the elements, if any, of the result are the same as would have been obtained if the scalar function had been applied separately, in array element order, to corresponding elements of each array [actual argument](#).

#### NOTE 12.54

An example of an [elemental reference](#) to the intrinsic function **MAX**:

if X and Y are arrays of shape (M, N),

**MAX** (X, 0.0, Y)

is an array expression of shape (M, N) whose elements have values

**MAX** (X(I, J), 0.0, Y(I, J)), I = 1, 2, ..., M, J = 1, 2, ..., N

### 12.8.3 Elemental subroutine actual arguments

- 1 In a reference to an elemental subroutine, if the [actual arguments](#) corresponding to **INTENT (OUT)** and **INTENT (INOUT)** [dummy arguments](#) are arrays, the values of the elements, if any, of the results are the same as would be obtained if the subroutine had been applied separately, in array element order, to corresponding elements of each array [actual argument](#).

## 13 Intrinsic procedures and modules

### 13.1 Classes of intrinsic procedures

- 1 Intrinsic procedures are divided into eight classes: [inquiry functions](#), [elemental functions](#), [transformational functions](#), [elemental subroutines](#), [pure subroutines](#), [atomic subroutines](#), [collective subroutines](#), and (impure) subroutines.
- 2 An intrinsic [inquiry function](#) is one whose result depends on the properties of one or more of its arguments instead of their values; in fact, these argument values may be undefined. Unless the description of an intrinsic [inquiry function](#) states otherwise, these arguments are permitted to be unallocated [allocatable](#) variables or pointers that are undefined or [disassociated](#). An [elemental](#) intrinsic function is one that is specified for scalar arguments, but may be applied to array arguments as described in [12.8](#). All other intrinsic functions are [transformational functions](#); they almost all have one or more array arguments or an array result. All standard intrinsic functions are pure.
- 3 An [atomic subroutine](#) is an intrinsic subroutine that performs an atomic action. The semantics of atomic actions are described in [13.5](#).
- 4 A [collective subroutine](#) is an intrinsic subroutine that performs a calculation on a [team](#) of [images](#) without requiring synchronization. The semantics of [collective subroutines](#) are described in [13.6](#).
- 5 The subroutine [MOVE\\_ALLOC](#) and the [elemental](#) subroutine [MVBITS](#) are [pure](#). No other standard intrinsic subroutine is pure.
- 6 [Generic names](#) of standard intrinsic procedures are listed in [13.7](#). In most cases, generic functions accept arguments of more than one type and the type of the result is the same as the type of the arguments. [Specific names](#) of standard intrinsic functions with corresponding generic names are listed in [13.8](#).
- 7 If an intrinsic procedure is used as an [actual argument](#) to a procedure, its [specific name](#) shall be used and it may be referenced in the called procedure only with scalar arguments. If an intrinsic procedure does not have a [specific name](#), it shall not be used as an [actual argument](#) ([12.5.2.9](#)).
- 8 [Elemental](#) intrinsic procedures behave as described in [12.8](#).

### 13.2 Arguments to intrinsic procedures

#### 13.2.1 General rules

- 1 All intrinsic procedures may be invoked with either positional arguments or [argument keywords](#) ([12.5](#)). The descriptions in [13.7](#) through [13.9](#) give the [argument keyword](#) names and positional sequence for standard intrinsic procedures.
- 2 Many of the intrinsic procedures have optional arguments. These arguments are identified by the notation “optional” in the argument descriptions. In addition, the names of the optional arguments are enclosed in square brackets in description headings and in lists of procedures. The valid forms of reference for procedures with optional arguments are described in [12.5.2](#).

#### NOTE 13.1

The text `CMPLX (X [, Y, KIND])` indicates that Y and KIND are both optional arguments. Valid reference forms include `CMPLX(x)`, `CMPLX(x, y)`, `CMPLX(x, KIND=kind)`, `CMPLX(x, y, kind)`, and `CMPLX(KIND=kind, X=x, Y=y)`.

**NOTE 13.2**

Some intrinsic procedures impose additional requirements on their optional arguments. For example, [SELECTED\\_REAL\\_KIND](#) requires that at least one of its optional arguments be present, and [RANDOM\\_SEED](#) requires that at most one of its optional arguments be present.

- 1    3    The dummy arguments of the specific intrinsic procedures in [13.8](#) have [INTENT \(IN\)](#). The dummy arguments of the intrinsic  
2    procedures in [13.9](#) have [INTENT \(IN\)](#) if the intent is not stated explicitly.
- 3    4    The [actual argument](#) corresponding to an intrinsic function dummy argument named KIND shall be a scalar  
4    integer [constant expression](#) and its value shall specify a representation method for the function result that exists  
5    on the processor.
- 6    5    Intrinsic subroutines that assign values to arguments of type character do so in accordance with the rules of  
7    intrinsic assignment ([7.2.1.3](#)).
- 8    6    In a reference to the intrinsic subroutine [MVBITS](#), the [actual arguments](#) corresponding to the TO and FROM  
9    dummy arguments may be the same variable and may be associated scalar variables or associated array variables  
10   all of whose corresponding elements are associated. Apart from this, the [actual arguments](#) in a reference to an  
11   [intrinsic](#) subroutine shall be such that the execution of the intrinsic subroutine would satisfy the restrictions of  
12   [12.5.2.13](#).
- 13   7    An argument to an intrinsic procedure other than [ASSOCIATED](#), [NULL](#), or [PRESENT](#) shall be a data object.

## 14    **13.2.2    The shape of array arguments**

- 15   1    Unless otherwise specified, the intrinsic [inquiry functions](#) accept array arguments for which the shape need not  
16   be defined. The shape of array arguments to [transformational](#) and [elemental](#) intrinsic functions shall be defined.

## 17    **13.2.3    Mask arguments**

- 18   1    Some array intrinsic functions have an optional MASK argument of type logical that is used by the function to  
19   select the elements of one or more arguments to be operated on by the function. Any element not selected by the  
20   mask need not be defined at the time the function is invoked.
- 21   2    The MASK affects only the value of the function, and does not affect the evaluation, prior to invoking the  
22   function, of arguments that are array expressions.

## 23    **13.2.4    DIM arguments and reduction functions**

- 24   1    Some array intrinsic functions are “reduction” functions; that is, they reduce the rank of an array by collapsing  
25   one dimension (or all dimensions, usually producing a scalar result). These functions have a DIM argument that  
26   can specify the dimension to be reduced.
- 27   2    The process of reducing a dimension usually combines the selected elements with a simple operation such as  
28   addition or an intrinsic function such as [MAX](#), but more sophisticated reductions are also provided, e.g. by  
29   [COUNT](#) and [MAXLOC](#).

## 30    **13.3    Bit model**

### 31    **13.3.1    General**

- 32   1    The bit manipulation procedures are described in terms of a model for the representation and behavior of bits  
33   on a processor.
- 2    2    For the purposes of these procedures, a bit is defined to be a binary digit  $w$  located at position  $k$  of a nonnegative

integer scalar object based on a model nonnegative integer defined by

$$j = \sum_{k=0}^{z-1} w_k \times 2^k$$

and for which  $w_k$  may have the value 0 or 1. This defines a sequence of bits  $w_{z-1} \dots w_0$ , with  $w_{z-1}$  the leftmost bit and  $w_0$  the rightmost bit. The positions of bits in the sequence are numbered from right to left, with the position of the rightmost bit being zero. The length of a sequence of bits is  $z$ . An example of a model number compatible with the examples used in 13.4 would have  $z = 32$ , thereby defining a 32-bit integer.

3 The interpretation of a negative integer as a sequence of bits is processor dependent.

4 The inquiry function `BIT_SIZE` provides the value of the parameter  $z$  of the model.

5 Effectively, this model defines an integer object to consist of  $z$  bits in sequence numbered from right to left from 0 to  $z - 1$ . This model is valid only in the context of the use of such an object as the argument or result of an intrinsic procedure that interprets that object as a sequence of bits. In all other contexts, the model defined for an integer in 13.4 applies. In particular, whereas the models are identical for  $r = 2$  and  $w_{z-1} = 0$ , they do not correspond for  $r \neq 2$  or  $w_{z-1} = 1$  and the interpretation of bits in such objects is processor dependent.

### 13.3.2 Bit sequence comparisons

1 When bit sequences of unequal length are compared, the shorter sequence is considered to be extended to the length of the longer sequence by padding with zero bits on the left.

2 Bit sequences are compared from left to right, one bit at a time, until unequal bits are found or all bits have been compared and found to be equal. If unequal bits are found, the sequence with zero in the unequal position is considered to be less than the sequence with one in the unequal position. Otherwise the sequences are considered to be equal.

### 13.3.3 Bit sequences as arguments to INT and REAL

1 When a *boz-literal-constant* is the argument A of the intrinsic function `INT` or `REAL`,

- if the length of the sequence of bits specified by A is less than the size in bits of a scalar variable of the same type and kind type parameter as the result, the *boz-literal-constant* is treated as if it were extended to a length equal to the size in bits of the result by padding on the left with zero bits, and
- if the length of the sequence of bits specified by A is greater than the size in bits of a scalar variable of the same type and kind type parameter as the result, the *boz-literal-constant* is treated as if it were truncated from the left to a length equal to the size in bits of the result.

C1301 If a *boz-literal-constant* is truncated as an argument to the intrinsic function `REAL`, the discarded bits shall all be zero.

#### NOTE 13.3

The result values of the intrinsic functions `CMPLX` and `DBLE` are defined by references to the intrinsic function `REAL` with the same arguments. Therefore, the padding and truncation of *boz-literal-constant* arguments to those intrinsic functions is the same as for the intrinsic function `REAL`.

## 13.4 Numeric models

1 The numeric manipulation and inquiry functions are described in terms of a model for the representation and behavior of numbers on a processor. The model has parameters that are determined so as to make the model best fit the machine on which the program is executed.

- 2 The model set for integer  $i$  is defined by

$$i = s \times \sum_{k=0}^{q-1} w_k \times r^k$$

1 where  $r$  is an integer exceeding one,  $q$  is a positive integer, each  $w_k$  is a nonnegative integer less than  $r$ , and  $s$  is  
2  $+1$  or  $-1$ .

- 3 The model set for real  $x$  is defined by

$$x = \begin{cases} 0 \text{ or} \\ s \times b^e \times \sum_{k=1}^p f_k \times b^{-k} \end{cases},$$

3 where  $b$  and  $p$  are integers exceeding one; each  $f_k$  is a nonnegative integer less than  $b$ , with  $f_1$  nonzero;  $s$  is  $+1$  or  
4  $-1$ ; and  $e$  is an integer that lies between some integer maximum  $e_{\max}$  and some integer minimum  $e_{\min}$  inclusively.  
5 For  $x = 0$ , its exponent  $e$  and digits  $f_k$  are defined to be zero. The integer parameters  $r$  and  $q$  determine the  
6 set of model integers and the integer parameters  $b$ ,  $p$ ,  $e_{\min}$ , and  $e_{\max}$  determine the set of model floating-point  
7 numbers.

- 8 4 The parameters of the integer and real models are available for each representation method of the integer and  
9 real types. The parameters characterize the set of available numbers in the definition of the model. Intrinsic  
10 functions provide the values of some parameters and other values related to the models.
- 11 5 There is also an extended model set for each kind of real  $x$ ; this extended model is the same as the ordinary  
12 model except that there are no limits on the range of the exponent  $e$ .

#### NOTE 13.4

Some of the function descriptions use the models

$$i = s \times \sum_{k=0}^{30} w_k \times 2^k$$

and

$$x = 0 \text{ or } s \times 2^e \times \left( \frac{1}{2} + \sum_{k=2}^{24} f_k \times 2^{-k} \right), \quad -126 \leq e \leq 127$$

## 13.5 Atomic subroutines

- 14 1 An [atomic subroutine](#) is an intrinsic subroutine that performs an action on its ATOM argument, and if it has an  
15 OLD argument, determination of the value to be assigned to that argument, atomically. Definition or evaluation  
16 of any argument other than ATOM is not performed atomically.
- 17 2 For any two executions in unordered segments of atomic subroutines whose ATOM argument is the same object,  
18 the effect is as if one of the executions is performed completely before the other execution begins. Which execution  
19 is performed first is processor dependent. The sequence of [atomic](#) actions within ordered segments is specified in  
20 [2.3.5](#). If successive atomic subroutine invocations on image  $P$  redefine a variable atomically in segments  $P_i$  and  
21  $P_j$ , atomic references to that variable from image  $Q$  in a segment  $Q_k$  that is unordered relative to  $P_i$  and  $P_j$  may  
22 observe the changes in the value of that variable in any order.
- 23 3 Atomic operations shall make asynchronous progress. If a variable  $X$  on image  $P$  is defined by an [atomic](#)  
24 [subroutine](#) on image  $Q$ , image  $R$  repeatedly references  $X[P]$  by an atomic subroutine in an unordered segment,

- 1 and no other image defines  $X[P]$  in an unordered segment, image  $R$  shall eventually receive the value assigned  
 2 by image  $Q$ , even if none of the images  $P$ ,  $Q$ , or  $R$  execute an **image control statement** until after the definition  
 3 of  $X[P]$  by image  $Q$  and the reception of that value by image  $R$ .
- 4 4 If the STAT argument is present in an invocation of an atomic subroutine and no error condition occurs, it is  
 5 assigned the value zero.
- 6 5 If the STAT argument is present in an invocation of an atomic subroutine and an error condition occurs, it  
 7 is assigned a processor-dependent positive value, and any other argument that is not **INTENT (IN)** becomes  
 8 undefined.
- 9 6 If the STAT argument is not present in an invocation of an atomic subroutine and an error condition occurs,  
 10 **error termination** is initiated.

**NOTE 13.5**

The properties of atomic subroutines are intended to support custom synchronization mechanisms. The program will need to handle all possible orderings of sequences of atomic subroutine executions that can arise as a consequence of the above rules; note that the orderings can appear to be different on different images even in the same program execution.

## 11 13.6 Collective subroutines

- 12 1 A **collective subroutine** is one that is invoked on each active image of the current team to perform a calculation  
 13 on those images and that assigns the computed value on one or all of them.
- 14 2 If a **collective subroutine** is invoked by one **image**, it shall be invoked by the same statement on all **active**  
 15 **images** of its **current team** in segments that are not ordered with respect to each other; corresponding references  
 16 (invocations) participate in the same collective computation. Before execution of the first **CHANGE TEAM**  
 17 **statement** on an **image**, in between executions of **CHANGE TEAM** and **END TEAM** statements, and after the  
 18 last execution of an **END TEAM** statement, the sequence of invocations of **collective subroutines** shall be the  
 19 same on all **active images** of a **team**. A **collective subroutine** shall not be referenced when an **image control**  
 20 **statement** is not permitted to be executed (for example, in a procedure invoked from a **CRITICAL construct**).
- 21 C1302 A reference to a **collective subroutine** shall not appear in a context where an **image control statement** is  
 22 not permitted to appear.
- 23 3 If the A argument in a reference to a **collective subroutine** is a **coarray**, the corresponding **ultimate arguments** on  
 24 all **active images** of the **current team** shall be corresponding **coarrays** as described in 2.4.7.

**Unresolved Technical Issue 026**

**There is no such thing as a “whole coarray”.**

And since putting in an image selector makes it *not a coarray*, there is no obvious meaning. Or perhaps you mean “coarray and either scalar or a whole array”? If so, that is what needs to be said.

- 25 4 If the optional STAT argument is present in a reference to a **collective subroutine** on one image, it shall be present  
 26 on all the corresponding references.
- 27 5 If the STAT argument is present in an invocation of a **collective subroutine** and its execution is successful, it is  
 28 assigned the value zero.
- 29 6 If the STAT argument is present in an invocation of a **collective subroutine** and an error condition occurs, it is  
 30 assigned a nonzero value and the A argument becomes undefined.
- 31 7 If the STAT argument is present and the **current team** contains **stopped images**, an error condition occurs and  
 32 the value **STAT\_STOPPED\_IMAGE** from the intrinsic module **ISO\_FORTRAN\_ENV** is assigned to the STAT

argument. Otherwise, if the `current team` contained `failed images`, an error condition occurs and the value `STAT_FAILED_IMAGE` from the intrinsic module `ISO_FORTRAN_ENV` is assigned to the `STAT` argument. If any other error condition occurs, the `STAT` argument is assigned a processor-dependent positive value that is different from the value of `STAT_STOPPED_IMAGE` or `STAT_FAILED_IMAGE`.

#### Unresolved Technical Issue 027

##### Involves synchronization with a stopped image?

The TS said that `STAT_STOPPED_IMAGE` only arose if there was synchronization with a stopped image. This is obviously nonsense since (a) we keep being told repeatedly that collectives do not involve synchronization, and more importantly (b) the description above says it applies only to the active images.

I took a guess that if any image in the current team is stopped, it should return `STAT_STOPPED_IMAGE`, just as it returns `STAT_FAILED_IMAGE` if any image is failed. If that guess is wrong, edits are required to unravel just exactly what it is that would cause `STAT_STOPPED_IMAGE` to occur.

- 8 If the `STAT` argument is not present in an invocation of a `collective subroutine` and an error condition occurs, error termination is initiated.
- 9 If the `ERRMSG` argument is present in an invocation of a `collective subroutine` and an error condition occurs, it is assigned an explanatory message, truncated or padded with blanks if required. If no error condition occurs, the definition status and value of `ERRMSG` are unchanged.

#### NOTE 13.6

The argument `A` becomes undefined if an error condition occurs during execution of a `collective subroutine` because it is intended to allow the processor to use `A` for intermediate values during calculation.

## 13.7 Standard generic intrinsic procedures

- 1 For all of the standard intrinsic procedures, the arguments shown are the names that shall be used for `argument keywords` if the keyword form is used for `actual arguments`.

#### NOTE 13.7

For example, a reference to `CMPLX` can be written in the form `CMPLX (A, B, M)` or in the form `CMPLX (Y = B, KIND = M, X = A)`.

#### NOTE 13.8

Many of the `argument keywords` have names that are indicative of their usage. For example:

<code>KIND</code>	Describes the kind type parameter of the result
<code>STRING, STRING_A</code>	An arbitrary character string
<code>BACK</code>	Controls the direction of string scan (forward or backward)
<code>MASK</code>	A mask to be applied to the arguments
<code>DIM</code>	A selected dimension of an array argument

- 2 In the Class column of Table 13.1,
  - `A` indicates that the procedure is an `atomic subroutine`,
  - `C` indicates that the procedure is a `collective subroutine`,
  - `E` indicates that the procedure is an `elemental function`,
  - `ES` indicates that the procedure is an `elemental subroutine`,
  - `I` indicates that the procedure is an `inquiry function`,

PS indicates that the procedure is a [pure](#) subroutine,  
 S indicates that the procedure is an impure subroutine, and  
 T indicates that the procedure is a [transformational function](#).

Table 13.1: **Standard generic intrinsic procedure summary**

Procedure	Arguments	Class	Description
ABS	(A)	E	Absolute value.
ACHAR	(I [, KIND])	E	Character from ASCII code value.
ACOS	(X)	E	Arccosine (inverse cosine) function.
ACOSH	(X)	E	Inverse hyperbolic cosine function.
ADJUSTL	(STRING)	E	Left-justified string value.
ADJUSTR	(STRING)	E	Right-justified string value.
AIMAG	(Z)	E	Imaginary part of a complex number.
AINIT	(A [, KIND])	E	Truncation toward 0 to a whole number.
ALL	(MASK) or (MASK, DIM)	T	Array reduced by <a href="#">.AND.</a> operator.
ALLOCATED	(ARRAY) or (SCALAR)	I	Allocation status of allocatable variable.
ANINT	(A [, KIND])	E	Nearest whole number.
ANY	(MASK) or (MASK, DIM)	T	Array reduced by <a href="#">.OR.</a> operator.
ASIN	(X)	E	Arcsine (inverse sine) function.
ASINH	(X)	E	Inverse hyperbolic sine function.
ASSOCIATED	(POINTER [, TARGET])	I	<a href="#">Pointer association</a> status inquiry.
ATAN	(X) or (Y, X)	E	Arctangent (inverse tangent) function.
ATAN2	(Y, X)	E	Arctangent (inverse tangent) function.
ATANH	(X)	E	Inverse hyperbolic tangent function.
ATOMIC_ADD	(ATOM, VALUE [, STAT])	A	Atomic addition.
ATOMIC_AND	(ATOM, VALUE [, STAT])	A	Atomic bitwise AND.
ATOMIC_CAS	(ATOM, OLD, COMPARE, VALUE [, STAT])	A	Atomic compare and swap.
ATOMIC_DEFINE	(ATOM, VALUE [, STAT])	A	Define a variable atomically.
ATOMIC_FETCH_ - ADD	(ATOM, VALUE, OLD [, STAT])	A	Atomic fetch and add.
ATOMIC_FETCH_ - AND	(ATOM, VALUE, OLD [, STAT])	A	Atomic fetch and bitwise AND.
ATOMIC_FETCH_ - OR	(ATOM, VALUE, OLD [, STAT])	A	Atomic fetch and bitwise OR.
ATOMIC_FETCH_ - XOR	(ATOM, VALUE, OLD [, STAT])	A	Atomic fetch and bitwise exclusive OR.
ATOMIC_OR	(ATOM, VALUE [, STAT])	A	Atomic bitwise OR.
ATOMIC_REF	(VALUE, ATOM [, STAT])	A	Reference a variable atomically.
ATOMIC_XOR	(ATOM, VALUE [, STAT])	A	Atomic bitwise exclusive OR.
BESSEL_J0	(X)	E	Bessel function of the 1 <sup>st</sup> kind, order 0.
BESSEL_J1	(X)	E	Bessel function of the 1 <sup>st</sup> kind, order 1.
BESSEL_JN	(N, X)	E	Bessel function of the 1 <sup>st</sup> kind, order N.
BESSEL_JN	(N1, N2, X)	T	Bessel functions of the 1 <sup>st</sup> kind.
BESSEL_Y0	(X)	E	Bessel function of the 2 <sup>nd</sup> kind, order 0.
BESSEL_Y1	(X)	E	Bessel function of the 2 <sup>nd</sup> kind, order 1.
BESSEL_YN	(N, X)	E	Bessel function of the 2 <sup>nd</sup> kind, order N.
BESSEL_YN	(N1, N2, X)	T	Bessel functions of the 2 <sup>nd</sup> kind.
BGE	(I, J)	E	Bitwise greater than or equal to.
BGT	(I, J)	E	Bitwise greater than.
BIT_SIZE	(I)	I	Number of bits in integer model <a href="#">13.3</a> .
BLE	(I, J)	E	Bitwise less than or equal to.
BLT	(I, J)	E	Bitwise less than.
BTEST	(I, POS)	E	Test single bit in an integer.
CEILING	(A [, KIND])	E	Least integer greater than or equal to A.

Table 13.1: Standard generic intrinsic procedure summary (cont.)

Procedure	Arguments	Class	Description
CHAR	(I [, KIND])	E	Character from code value.
CMPLX	(X [, KIND]) or (X [, Y, KIND])	E	Conversion to complex type.
CO_BROADCAST	(A, SOURCE_IMAGE [, STAT, ERRMSG])	C	Broadcast value to images.
CO_MAX	(A [, RESULT_IMAGE, STAT, ERRMSG])	C	Compute maximum value across images.
CO_MIN	(A [, RESULT_IMAGE, STAT, ERRMSG])	C	Compute minimum value across images.
CO_REDUCE	(A, OPERATOR [, RESULT_IMAGE, STAT, ERRMSG])	C	Generalized reduction across images.
CO_SUM	(A [, RESULT_IMAGE, STAT, ERRMSG])	C	Compute sum across images.
COMMAND_ARGUMENT_COUNT	( )	T	Number of command arguments.
CONJG	(Z)	E	Conjugate of a complex number.
COS	(X)	E	Cosine function.
COSH	(X)	E	Hyperbolic cosine function.
COSHAPE	(COARRAY [, KIND])	I	Sizes of codimensions of a coarray.
COUNT	(MASK [, DIM, KIND])	T	Logical array reduced by counting true values.
CPU_TIME	(TIME)	S	Processor time used.
CSHIFT	(ARRAY, SHIFT [, DIM])	T	Circular shift of an array.
DATE_AND_TIME	([DATE, TIME, ZONE, VALUES])	S	Date and time.
DBLE	(A)	E	Conversion to double precision real.
DIGITS	(X)	I	Significant digits in numeric model.
DIM	(X, Y)	E	Maximum of X – Y and zero.
DOT_PRODUCT	(VECTOR_A, VECTOR_B)	T	Dot product of two vectors.
DPROD	(X, Y)	E	Double precision real product.
DSHIFTL	(I, J, SHIFT)	E	Combined left shift.
DSHIFTR	(I, J, SHIFT)	E	Combined right shift.
EOSHIFT	(ARRAY, SHIFT [, BOUNDARY, DIM])	T	End-off shift of the elements of an array.
EPSILON	(X)	I	Model number that is small compared to 1.
ERF	(X)	E	Error function.
ERFC	(X)	E	Complementary error function.
ERFC_SCALED	(X)	E	Scaled complementary error function.
EVENT_QUERY	(EVENT, COUNT [, STAT])	S	Query event count.
EXECUTE_COMMAND_LINE	(COMMAND [, WAIT, EXITSTAT, CMDSTAT, CMDMSG])	S	Execute a command line.
EXP	(X)	E	Exponential function.
EXPONENT	(X)	E	Exponent of floating-point number.
EXTENDS_TYPE_OF	(A, MOLD)	I	Dynamic type extension inquiry.
FAILED_IMAGES	([TEAM, KIND])	T	Indices of failed images.
FINDLOC	(ARRAY, VALUE, DIM [, MASK, KIND, BACK]) or (ARRAY, VALUE [, MASK, KIND, BACK])	T	Location(s) of a specified value.

Table 13.1: Standard generic intrinsic procedure summary (cont.)

Procedure	Arguments	Class	Description
FLOOR	(A [, KIND])	E	Greatest integer less than or equal to A.
FRACTION	(X)	E	Fractional part of number.
GAMMA	(X)	E	Gamma function.
GET_COMMAND	([COMMAND, LENGTH, STATUS, ERRMSG])	S	Get program invocation command.
GET_COMMAND_ARGUMENT	(NUMBER [, VALUE, LENGTH, STATUS, ERRMSG])	S	Get program invocation argument.
GET_ENVIRONMENT_VARIABLE	(NAME [, VALUE, LENGTH, STATUS, TRIM_NAME, ERRMSG])	S	Get environment variable.
GET_TEAM	([LEVEL])	T	Team.
HUGE	(X)	I	Largest model number.
HYPOT	(X, Y)	E	Euclidean distance function.
IACHAR	(C [, KIND])	E	ASCII code value for character.
IALL	(ARRAY, DIM [, MASK]) or (ARRAY [, MASK])	T	Array reduced by IAND function.
IAND	(I, J)	E	Bitwise AND.
IANY	(ARRAY, DIM [, MASK]) or (ARRAY [, MASK])	T	Array reduced by IOR function.
IBCLR	(I, POS)	E	I with bit POS replaced by zero.
IBITS	(I, POS, LEN)	E	Specified sequence of bits.
IBSET	(I, POS)	E	I with bit POS replaced by one.
ICHAR	(C [, KIND])	E	Code value for character.
IEOR	(I, J)	E	Bitwise exclusive OR.
IMAGE_INDEX	(COARRAY, SUB) or (COARRAY, SUB, TEAM) or (COARRAY, SUB, TEAM_NUMBER)	I	<a href="#">Image index</a> from <a href="#">cosubscripts</a> .
IMAGE_STATUS	(IMAGE [, TEAM])	T	Image execution state.
INDEX	(STRING, SUBSTRING [, BACK, KIND])	E	Character string search.
INT	(A [, KIND])	E	Conversion to integer type.
IOR	(I, J)	E	Bitwise inclusive OR.
IPARITY	(ARRAY, DIM [, MASK]) or (ARRAY [, MASK])	T	Array reduced by IEOB function.
ISHFT	(I, SHIFT)	E	Logical shift.
ISHFTC	(I, SHIFT [, SIZE])	E	Circular shift of the rightmost bits.
IS_CONTIGUOUS	(ARRAY)	I	Array contiguity test ( <a href="#">5.5.7</a> ).
IS_IOSTAT_END	(I)	E	IOSTAT value test for end of file.
IS_IOSTAT_EOR	(I)	E	IOSTAT value test for end of record.
KIND	(X)	I	Value of the kind type parameter of X.
LBOUND	(ARRAY [, DIM, KIND])	I	Lower bound(s).
LCBOUND	(COARRAY [, DIM, KIND])	I	Lower <a href="#">cobound</a> (s) of a <a href="#">coarray</a> .
LEADZ	(I)	E	Number of leading zero bits.
LEN	(STRING [, KIND])	I	Length of a character entity.
LEN_TRIM	(STRING [, KIND])	E	Length without trailing blanks.
LGE	(STRING_A, STRING_B)	E	ASCII greater than or equal.
LGT	(STRING_A, STRING_B)	E	ASCII greater than.
LLE	(STRING_A, STRING_B)	E	ASCII less than or equal.
LLT	(STRING_A, STRING_B)	E	ASCII less than.
LOG	(X)	E	Natural logarithm.
LOG_GAMMA	(X)	E	Logarithm of the absolute value of the gamma function.

Table 13.1: Standard generic intrinsic procedure summary (cont.)

Procedure	Arguments	Class	Description
LOG10	(X)	E	Common logarithm.
LOGICAL	(L [, KIND])	E	Conversion between kinds of logical.
MASKL	(I [, KIND])	E	Left justified mask.
MASKR	(I [, KIND])	E	Right justified mask.
MATMUL	(MATRIX_A, MATRIX_B)	T	Matrix multiplication.
MAX	(A1, A2 [, A3, ...])	E	Maximum value.
MAXEXPONENT	(X)	I	Maximum exponent of a real model.
MAXLOC	(ARRAY, DIM [, MASK, KIND, BACK]) or (ARRAY [, MASK, KIND, BACK])	T	Location(s) of maximum value.
MAXVAL	(ARRAY, DIM [, MASK]) or (ARRAY [, MASK])	T	Maximum value(s) of array.
MERGE	(TSOURCE, FSOURCE, MASK)	E	Expression value selection.
MERGE_BITS	(I, J, MASK)	E	Merge of bits under mask.
MIN	(A1, A2 [, A3, ...])	E	Minimum value.
MINEXPONENT	(X)	I	Minimum exponent of a real model.
MINLOC	(ARRAY, DIM [, MASK, KIND, BACK]) or (ARRAY [, MASK, KIND, BACK])	T	Location(s) of minimum value.
MINVAL	(ARRAY, DIM [, MASK]) or (ARRAY [, MASK])	T	Minimum value(s) of array.
MOD	(A, P)	E	Remainder function.
MODULO	(A, P)	E	Modulo function.
MOVE_ALLOC	(FROM, TO [, STAT, ERRMSG])	PS	Move an allocation.
MVBITS	(FROM, FROMPOS, LEN, TO, TOPOS)	ES	Copy a sequence of bits.
NEAREST	(X, S)	E	Adjacent machine number.
NEW_LINE	(A)	I	Newline character.
NINT	(A [, KIND])	E	Nearest integer.
NORM2	(X) or (X, DIM)	T	$L_2$ norm of an array.
NOT	(I)	E	Bitwise complement.
NULL	([MOLD])	T	<a href="#">Disassociated</a> pointer or unallocated allocatable entity.
NUM_IMAGES	( ) or (TEAM) or (TEAM_NUMBER)	T	Number of <a href="#">images</a> .
OUT_OF_RANGE	(X, MOLD [, ROUND])	E	Whether a value cannot be converted safely.
PACK	(ARRAY, MASK [, VECTOR])	T	Array packed into a vector.
PARITY	(MASK) or (MASK, DIM)	T	Array reduced by <a href="#">.NEQV.</a> operator.
POPCNT	(I)	E	Number of one bits.
POPPAR	(I)	E	Parity expressed as 0 or 1.
PRECISION	(X)	I	Decimal precision of a real model.
PRESENT	(A)	I	Presence of optional argument.
PRODUCT	(ARRAY, DIM [, MASK]) or (ARRAY [, MASK])	T	Array reduced by multiplication.
RADIX	(X)	I	Base of a numeric model.
RANDOM_INIT	(REPEATABLE, IMAGE_DISTINCT)	S	Initialise the pseudorandom number generator.
RANDOM_NUMBER	(HARVEST)	S	Generate pseudorandom number(s).
RANDOM_SEED	([SIZE, PUT, GET])	S	Restart or query the pseudorandom number generator.

Table 13.1: Standard generic intrinsic procedure summary (cont.)

Procedure	Arguments	Class	Description
RANGE	(X)	I	Decimal exponent range of a numeric model (13.4).
RANK	(A)	I	Rank of a data object.
REAL	(A [, KIND])	E	Conversion to real type.
REDUCE	(ARRAY, OPERATION, DIM [, MASK, IDENTITY, ORDERED]) or (ARRAY, OPERATION [, MASK, IDENTITY, ORDERED])	T	General reduction of array
REPEAT	(STRING, NCOPIES)	T	Repetitive string concatenation.
RESHAPE	(SOURCE, SHAPE [, PAD, ORDER])	T	Arbitrary shape array construction.
RRSPACING	(X)	E	Reciprocal of relative spacing of model numbers.
SAME_TYPE_AS	(A, B)	I	<a href="#">Dynamic type</a> equality test.
SCALE	(X, I)	E	Real number scaled by radix power.
SCAN	(STRING, SET [, BACK, KIND])	E	Character set membership search.
SELECTED_CHAR_KIND	(NAME)	T	Character kind selection.
SELECTED_INT_KIND	(R)	T	Integer kind selection.
SELECTED_REAL_KIND	([P, R, RADIX])	T	Real kind selection.
SET_EXPONENT	(X, I)	E	Real value with specified exponent.
SHAPE	(SOURCE [, KIND])	I	Shape of an array or a scalar.
SHIFTA	(I, SHIFT)	E	Right shift with fill.
SHIFTL	(I, SHIFT)	E	Left shift.
SHIFTR	(I, SHIFT)	E	Right shift.
SIGN	(A, B)	E	Magnitude of A with the sign of B.
SIN	(X)	E	Sine function.
SINH	(X)	E	Hyperbolic sine function.
SIZE	(ARRAY [, DIM, KIND])	I	Size of an array or one extent.
SPACING	(X)	E	Spacing of model numbers.
SPREAD	(SOURCE, DIM, NCOPIES)	T	Value replicated in a new dimension.
SQRT	(X)	E	Square root.
STOPPED_IMAGES	([TEAM, KIND])	T	Indices of stopped images.
STORAGE_SIZE	(A [, KIND])	I	Storage size in bits.
SUM	(ARRAY, DIM [, MASK]) or (ARRAY [, MASK])	T	Array reduced by addition.
SYSTEM_CLOCK	([COUNT, COUNT_RATE, COUNT_MAX])	S	Query system clock.
TAN	(X)	E	Tangent function.
TANH	(X)	E	Hyperbolic tangent function.
TEAM_NUMBER	([TEAM])	T	Team number.
THIS_IMAGE	([TEAM])	T	<a href="#">Index</a> of the invoking <a href="#">image</a> .
THIS_IMAGE	(COARRAY [, TEAM]) or (COARRAY, DIM [, TEAM])	T	<a href="#">Cosubscript(s)</a> for this <a href="#">image</a> .
TINY	(X)	I	Smallest positive model number.
TRAILZ	(I)	E	Number of trailing zero bits.
TRANSFER	(SOURCE, MOLD [, SIZE])	T	Transfer physical representation.
TRANSPOSE	(MATRIX)	T	Transpose of an array of <a href="#">rank</a> two.

Table 13.1: Standard generic intrinsic procedure summary (cont.)

Procedure	Arguments	Class	Description
TRIM	(STRING)	T	String without trailing blanks.
UBOUND	(ARRAY [, DIM, KIND])	I	Upper bound(s).
UCOBOUND	(COARRAY [, DIM, KIND])	I	Upper <a href="#">cobound</a> (s) of a <a href="#">coarray</a> .
UNPACK	(VECTOR, MASK, FIELD)	T	Vector unpacked into an array.
VERIFY	(STRING, SET [, BACK, KIND])	E	Character set non-membership search.

- 1 3 The effect of calling [EXECUTE\\_COMMAND\\_LINE](#) on any [image](#) other than [image](#) 1 in the [initial team](#) is  
2 processor dependent.
- 3 4 If [RANDOM\\_INIT](#) or [RANDOM\\_SEED](#) is called in a segment A, and [RANDOM\\_INIT](#), [RANDOM\\_SEED](#), or  
4 [RANDOM\\_NUMBER](#) is called in segment B, then segments A and B shall be ordered. It is processor dependent  
5 whether each [image](#) uses a separate random number generator, or if some or all [images](#) use common random  
6 number generators.
- 7 5 The use of all other standard intrinsic procedures in unordered segments is subject only to their argument use  
8 following the rules in [8.6.2](#).

## 9 13.8 Specific names for standard intrinsic functions

- 10 1 Except for AMAX0, AMIN0, MAX1, and MIN1, the result type of the specific function is the same that the result type of the  
11 corresponding generic function reference would be if it were invoked with the same arguments as the specific function.
- 12 2 A function listed in Table 13.3 is not permitted to be used as an [actual argument](#) ([12.5.1](#), C1240), as a target in a procedure  
13 [pointer assignment statement](#) ([7.2.2.2](#), C730), as an initial target in a [procedure declaration statement](#) ([12.4.3.7](#), C1225), or to specify  
14 an interface ([12.4.3.7](#), C1221).

Table 13.2: Unrestricted specific intrinsic functions

Specific name	Generic name	Argument type and kind
ABS	ABS	default real
ACOS	ACOS	default real
AIMAG	AIMAG	default complex
AINT	AINT	default real
ALOG	LOG	default real
ALOG10	LOG10	default real
AMOD	MOD	default real
ANINT	ANINT	default real
ASIN	ASIN	default real
ATAN	ATAN (X)	default real
ATAN2	ATAN2	default real
CABS	ABS	default complex
CCOS	COS	default complex
CEXP	EXP	default complex
CLOG	LOG	default complex
CONJG	CONJG	default complex
COS	COS	default real
COSH	COSH	default real
CSIN	SIN	default complex
CSQRT	SQRT	default complex
DABS	ABS	double precision real
DACOS	ACOS	double precision real
DASIN	ASIN	double precision real
DATAN	ATAN	double precision real
DATAN2	ATAN2	double precision real
DCOS	COS	double precision real
DCOSH	COSH	double precision real
DDIM	DIM	double precision real
DEXP	EXP	double precision real
DIM	DIM	default real
DINT	AINT	double precision real

Unrestricted specific intrinsic functions (cont.)

Specific name	Generic name	Argument type and kind
DLOG	LOG	double precision real
DLOG10	LOG10	double precision real
DMOD	MOD	double precision real
DNINT	ANINT	double precision real
DPROD	DPROD	default real
DSIGN	SIGN	double precision real
DSIN	SIN	double precision real
DSINH	SINH	double precision real
DSQRT	SQRT	double precision real
DTAN	TAN	double precision real
DTANH	TANH	double precision real
EXP	EXP	default real
IABS	ABS	default integer
IDIM	DIM	default integer
IDNINT	NINT	double precision real
INDEX	INDEX	default character
ISIGN	SIGN	default integer
LEN	LEN	default character
MOD	MOD	default integer
NINT	NINT	default real
SIGN	SIGN	default real
SIN	SIN	default real
SINH	SINH	default real
SQRT	SQRT	default real
TAN	TAN	default real
TANH	TANH	default real

Table 13.3: Restricted specific intrinsic functions

Specific name	Generic name	Argument type and kind
AMAX0 (...)	REAL (MAX (...))	default integer
AMAX1	MAX	default real
AMIN0 (...)	REAL (MIN (...))	default integer
AMIN1	MIN	default real
CHAR	CHAR	default integer
DMAX1	MAX	double precision real
DMIN1	MIN	double precision real
FLOAT	REAL	default integer
ICHAR	ICHAR	default character
IDINT	INT	double precision real
IFIX	INT	default real
INT	INT	default real
LGE	LGE	default character
LGT	LGT	default character
LLE	LLE	default character
LLT	LLT	default character
MAX0	MAX	default integer
MAX1 (...)	INT (MAX (...))	default real
MIN0	MIN	default integer
MIN1 (...)	INT (MIN (...))	default real
REAL	REAL	default integer
SNGL	REAL	double precision real

## 13.9 Specifications of the standard intrinsic procedures

### 13.9.1 General

- 1 Detailed specifications of the standard generic intrinsic procedures are provided in [13.9](#) in alphabetical order.
- 2 The types and type parameters of standard intrinsic procedure arguments and function results are determined by these specifications. The “Argument(s)” paragraphs specify requirements on the [actual arguments](#) of the

procedures. The result [characteristics](#) are sometimes specified in terms of the [characteristics](#) of dummy arguments. A program shall not invoke an intrinsic procedure under circumstances where a value to be assigned to a subroutine argument or returned as a function result is not representable by objects of the specified type and type parameters.

3 If an IEEE infinity is assigned or returned by an intrinsic procedure, the intrinsic module [IEEE\\_ARITHMETIC](#) is accessible, and the actual arguments were finite numbers, the flag IEEE\_OVERFLOW or IEEE\_DIVIDE\_BY\_ZERO shall signal. If an IEEE NaN is assigned or returned, the actual arguments were finite numbers, the intrinsic module [IEEE\\_ARITHMETIC](#) is accessible, and the exception IEEE\_INVALID is supported, the flag IEEE\_INVALID shall signal. If no IEEE infinity or NaN is assigned or returned, these flags shall have the same status as when the intrinsic procedure was invoked.

## 13.9.2 ABS (A)

**1 Description.** Absolute value.

**2 Class.** [Elemental](#) function.

**3 Argument.** A shall be of type integer, real, or complex.

**4 Result Characteristics.** The same as A except that if A is complex, the result is real.

**5 Result Value.** If A is of type integer or real, the value of the result is  $|A|$ ; if A is complex with value  $(x, y)$ , the result is equal to a processor-dependent approximation to  $\sqrt{x^2 + y^2}$  computed without undue overflow or underflow.

**6 Example.** ABS ((3.0, 4.0)) has the value 5.0 (approximately).

## 13.9.3 ACHAR (I [, KIND])

**1 Description.** Character from ASCII code value.

**2 Class.** [Elemental](#) function.

**3 Arguments.**

I shall be of type integer.

KIND (optional) shall be a scalar integer [constant expression](#).

**4 Result Characteristics.** Character of length one. If KIND is present, the kind type parameter is that specified by the value of KIND; otherwise, the kind type parameter is that of default character.

**5 Result Value.** If I has a value in the range  $0 \leq I \leq 127$ , the result is the character in position I of the [ASCII collating sequence](#), provided the processor is capable of representing that character in the character kind of the result; otherwise, the result is processor dependent. ACHAR (IACHAR (C)) shall have the value C for any character C capable of representation as a default character.

**6 Example.** ACHAR (88) has the value 'X'.

## 13.9.4 ACOS (X)

**1 Description.** Arccosine (inverse cosine) function.

**2 Class.** [Elemental](#) function.

**3 Argument.** X shall be of type real with a value that satisfies the inequality  $|X| \leq 1$ , or of type complex.

**4 Result Characteristics.** Same as X.

**5 Result Value.** The result has a value equal to a processor-dependent approximation to  $\arccos(X)$ . If it is real it is expressed in radians and lies in the range  $0 \leq \text{ACOS}(X) \leq \pi$ . If it is complex the real part is expressed in

1 radians and lies in the range  $0 \leq \text{REAL}(\text{ACOS}(X)) \leq \pi$ .

2 6 **Example.** ACOS (0.54030231) has the value 1.0 (approximately).

### 3 13.9.5 ACOSH (X)

4 1 **Description.** Inverse hyperbolic cosine function.

5 2 **Class.** [Elemental](#) function.

6 3 **Argument.** X shall be of type real or complex.

7 4 **Result Characteristics.** Same as X.

8 5 **Result Value.** The result has a value equal to a processor-dependent approximation to the inverse hyperbolic  
9 cosine function of X. If the result is complex the imaginary part is expressed in radians and lies in the range  
10  $0 \leq \text{AIMAG}(\text{ACOSH}(X)) \leq \pi$

11 6 **Example.** ACOSH (1.5430806) has the value 1.0 (approximately).

### 12 13.9.6 ADJUSTL (STRING)

13 1 **Description.** Left-justified string value.

14 2 **Class.** [Elemental](#) function.

15 3 **Argument.** STRING shall be of type character.

16 4 **Result Characteristics.** Character of the same length and kind type parameter as STRING.

17 5 **Result Value.** The value of the result is the same as STRING except that any leading blanks have been deleted  
18 and the same number of trailing blanks have been inserted.

19 6 **Example.** ADJUSTL (' WORD') has the value 'WORD '.

### 20 13.9.7 ADJUSTR (STRING)

21 1 **Description.** Right-justified string value.

22 2 **Class.** [Elemental](#) function.

23 3 **Argument.** STRING shall be of type character.

24 4 **Result Characteristics.** Character of the same length and kind type parameter as STRING.

25 5 **Result Value.** The value of the result is the same as STRING except that any trailing blanks have been deleted  
26 and the same number of leading blanks have been inserted.

27 6 **Example.** ADJUSTR ('WORD ') has the value ' WORD'.

### 28 13.9.8 AIMAG (Z)

29 1 **Description.** Imaginary part of a complex number.

30 2 **Class.** [Elemental](#) function.

31 3 **Argument.** Z shall be of type complex.

32 4 **Result Characteristics.** Real with the same kind type parameter as Z.

33 5 **Result Value.** If Z has the value  $(x, y)$ , the result has the value  $y$ .

1 6 **Example.** AIMAG ((2.0, 3.0)) has the value 3.0.

## 2 13.9.9 AINT (A [, KIND])

3 1 **Description.** Truncation toward 0 to a whole number.

4 2 **Class.** [Elemental](#) function.

5 3 **Arguments.**

6 A shall be of type real.

7 KIND (optional) shall be a scalar integer [constant expression](#).

8 4 **Result Characteristics.** The result is of type real. If KIND is present, the kind type parameter is that specified by the value of KIND; otherwise, the kind type parameter is that of A.

10 5 **Result Value.** If  $|A| < 1$ , AINT (A) has the value 0; if  $|A| \geq 1$ , AINT (A) has a value equal to the integer whose magnitude is the largest integer that does not exceed the magnitude of A and whose sign is the same as the sign of A.

13 6 **Examples.** AINT (2.783) has the value 2.0. AINT (−2.783) has the value −2.0.

## 14 13.9.10 ALL (MASK) or ALL (MASK, DIM)

15 1 **Description.** Array reduced by [.AND.](#) operator.

16 2 **Class.** [Transformational](#) function.

17 3 **Arguments.**

18 MASK shall be a logical array.

19 DIM shall be an integer scalar with value in the range  $1 \leq \text{DIM} \leq n$ , where  $n$  is the [rank](#) of MASK.

20 4 **Result Characteristics.** The result is of type logical with the same kind type parameter as MASK. It is scalar if DIM does not appear or  $n = 1$ ; otherwise, the result has [rank](#)  $n - 1$  and shape  $[d_1, d_2, \dots, d_{\text{DIM}-1}, d_{\text{DIM}+1}, \dots, d_n]$  where  $[d_1, d_2, \dots, d_n]$  is the shape of MASK.

23 5 **Result Value.**

24 *Case (i):* The result of ALL (MASK) has the value true if all elements of MASK are true or if MASK has size zero, and the result has value false if any element of MASK is false.

26 *Case (ii):* If MASK has [rank](#) one, ALL (MASK, DIM) is equal to ALL (MASK). Otherwise, the value of element  $(s_1, s_2, \dots, s_{\text{DIM}-1}, s_{\text{DIM}+1}, \dots, s_n)$  of ALL (MASK, DIM) is equal to ALL (MASK  $(s_1, s_2, \dots, s_{\text{DIM}-1}, :, s_{\text{DIM}+1}, \dots, s_n)$ ).

29 6 **Examples.**

30 *Case (i):* The value of ALL ([[.TRUE.](#), [.FALSE.](#), [.TRUE.](#)]) is false.

31 *Case (ii):* If B is the array  $\begin{bmatrix} 1 & 3 & 5 \\ 2 & 4 & 6 \end{bmatrix}$  and C is the array  $\begin{bmatrix} 0 & 3 & 5 \\ 7 & 4 & 8 \end{bmatrix}$  then ALL (B /= C, DIM = 1) is [true, false, false] and ALL (B /= C, DIM = 2) is [false, false].

## 33 13.9.11 ALLOCATED (ARRAY) or ALLOCATED (SCALAR)

34 1 **Description.** Allocation status of allocatable variable.

35 2 **Class.** [Inquiry](#) function.

36 3 **Arguments.**

37 ARRAY shall be an [allocatable](#) array.

1 SCALAR shall be an [allocatable](#) scalar.

2 4 **Result Characteristics.** Default logical scalar.

3 5 **Result Value.** The result has the value true if the argument (ARRAY or SCALAR) is allocated and has the  
4 value false if the argument is unallocated.

## 5 13.9.12 ANINT (A [, KIND])

6 1 **Description.** Nearest whole number.

7 2 **Class.** [Elemental](#) function.

8 3 **Arguments.**

9 A shall be of type real.

10 KIND (optional) shall be a scalar integer [constant expression](#).

11 4 **Result Characteristics.** The result is of type real. If KIND is present, the kind type parameter is that specified  
12 by the value of KIND; otherwise, the kind type parameter is that of A.

13 5 **Result Value.** The result is the integer nearest A, or if there are two integers equally near A, the result is  
14 whichever such integer has the greater magnitude.

15 6 **Examples.** ANINT (2.783) has the value 3.0. ANINT (−2.783) has the value −3.0.

## 16 13.9.13 ANY (MASK) or ANY (MASK, DIM)

17 1 **Description.** Array reduced by [.OR.](#) operator.

18 2 **Class.** [Transformational](#) function.

19 3 **Arguments.**

20 MASK shall a logical array.

21 DIM shall be an integer scalar with a value in the range  $1 \leq \text{DIM} \leq n$ , where  $n$  is the [rank](#) of MASK.

22 4 **Result Characteristics.** The result is of type logical with the same kind type parameter as MASK. It is scalar  
23 if DIM does not appear or  $n = 1$ ; otherwise, the result has [rank](#)  $n - 1$  and shape  $[d_1, d_2, \dots, d_{\text{DIM}-1}, d_{\text{DIM}+1},$   
24  $\dots, d_n]$  where  $[d_1, d_2, \dots, d_n]$  is the shape of MASK.

25 5 **Result Value.**

26 *Case (i):* The result of ANY (MASK) has the value true if any element of MASK is true and has the value  
27 false if no elements are true or if MASK has size zero.

28 *Case (ii):* If MASK has [rank](#) one, ANY (MASK, DIM) is equal to ANY (MASK). Otherwise, the value of  
29 element  $(s_1, s_2, \dots, s_{\text{DIM}-1}, s_{\text{DIM}+1}, \dots, s_n)$  of ANY (MASK, DIM) is equal to ANY (MASK  $(s_1,$   
30  $s_2, \dots, s_{\text{DIM}-1}, :, s_{\text{DIM}+1}, \dots, s_n)$ ).

31 6 **Examples.**

32 *Case (i):* The value of ANY ([[.TRUE.](#), [.FALSE.](#), [.TRUE.](#)]) is true.

33 *Case (ii):* If B is the array  $\begin{bmatrix} 1 & 3 & 5 \\ 2 & 4 & 6 \end{bmatrix}$  and C is the array  $\begin{bmatrix} 0 & 3 & 5 \\ 7 & 4 & 8 \end{bmatrix}$  then ANY (B /= C, DIM = 1) is  
34 [true, false, true] and ANY (B /= C, DIM = 2) is [true, true].

### 13.9.14 ASIN (X)

**Description.** Arcsine (inverse sine) function.

**Class.** [Elemental](#) function.

**Argument.** X shall be of type real with a value that satisfies the inequality  $|X| \leq 1$ , or of type complex.

**Result Characteristics.** Same as X.

**Result Value.** The result has a value equal to a processor-dependent approximation to  $\arcsin(X)$ . If it is real it is expressed in radians and lies in the range  $-\pi/2 \leq \text{ASIN}(X) \leq \pi/2$ . If it is complex the real part is expressed in radians and lies in the range  $-\pi/2 \leq \text{REAL}(\text{ASIN}(X)) \leq \pi/2$ .

**Example.**  $\text{ASIN}(0.84147098)$  has the value 1.0 (approximately).

### 13.9.15 ASINH (X)

**Description.** Inverse hyperbolic sine function.

**Class.** [Elemental](#) function.

**Argument.** X shall be of type real or complex.

**Result Characteristics.** Same as X.

**Result Value.** The result has a value equal to a processor-dependent approximation to the inverse hyperbolic sine function of X. If the result is complex the imaginary part is expressed in radians and lies in the range  $-\pi/2 \leq \text{AIMAG}(\text{ASINH}(X)) \leq \pi/2$ .

**Example.**  $\text{ASINH}(1.1752012)$  has the value 1.0 (approximately).

### 13.9.16 ASSOCIATED (POINTER [, TARGET])

**Description.** [Pointer association](#) status inquiry.

**Class.** [Inquiry function](#).

**Arguments.**

POINTER shall be a pointer. It may be of any type or may be a procedure pointer. Its [pointer association](#) status shall not be undefined.

TARGET (optional) shall be allowable as the [data-target](#) or [proc-target](#) in a [pointer assignment statement](#) (7.2.2) in which POINTER is [data-pointer-object](#) or [proc-pointer-object](#). If TARGET is a pointer then its [pointer association](#) status shall not be undefined.

**Result Characteristics.** Default logical scalar.

**Result Value.**

*Case (i):* If TARGET is absent, the result is true if and only if POINTER is associated with a target.

*Case (ii):* If TARGET is present and is a procedure, the result is true if and only if POINTER is associated with TARGET.

*Case (iii):* If TARGET is present and is a procedure pointer, the result is true if and only if POINTER and TARGET are associated with the same procedure.

*Case (iv):* If TARGET is present and is a scalar target, the result is true if and only if TARGET is not a zero-sized [storage sequence](#) and POINTER is associated with a target that occupies the same [storage units](#) as TARGET.

*Case (v):* If TARGET is present and is an array target, the result is true if and only if POINTER is associated with a target that has the same shape as TARGET, is neither of size zero nor an array whose elements

are zero-sized [storage sequences](#), and occupies the same [storage units](#) as TARGET in array element order.

*Case (vi):* If TARGET is present and is a scalar pointer, the result is true if and only if POINTER and TARGET are associated, the targets are not zero-sized [storage sequences](#), and they occupy the same [storage units](#).

*Case (vii):* If TARGET is present and is an [array pointer](#), the result is true if and only if POINTER and TARGET are both associated, have the same shape, are neither of size zero nor arrays whose elements are zero-sized [storage sequences](#), and occupy the same [storage units](#) in array element order.

#### NOTE 13.9

The references to TARGET in the above cases are referring to properties that might be possessed by the actual argument, so the case of TARGET being a disassociated pointer will be covered by case [\(iii\)](#), [\(vi\)](#), or [\(vii\)](#).

**6 Examples.** ASSOCIATED (CURRENT, HEAD) is true if CURRENT is associated with the target HEAD. After the execution of  
`A_PART => A (:N)`  
 ASSOCIATED (A\_PART, A) is true if N is equal to UBOUND (A, DIM = 1). After the execution of  
`NULLIFY (CUR); NULLIFY (TOP)`  
 ASSOCIATED (CUR, TOP) is false.

### 13.9.17 ATAN (X) or ATAN (Y, X)

**1 Description.** Arctangent (inverse tangent) function.

**2 Class.** [Elemental](#) function.

**3 Arguments.**

Y shall be of type real.

X If Y appears, X shall be of type real with the same kind type parameter as Y. If Y has the value zero, X shall not have the value zero. If Y does not appear, X shall be of type real or complex.

**4 Result Characteristics.** Same as X.

**5 Result Value.** If Y appears, the result is the same as the result of [ATAN2](#) (Y,X). If Y does not appear, the result has a value equal to a processor-dependent approximation to  $\arctan(X)$  whose real part is expressed in radians and lies in the range  $-\pi/2 \leq \text{ATAN}(X) \leq \pi/2$ .

**6 Example.** ATAN (1.5574077) has the value 1.0 (approximately).

### 13.9.18 ATAN2 (Y, X)

**1 Description.** Arctangent (inverse tangent) function.

**2 Class.** [Elemental](#) function.

**3 Arguments.**

Y shall be of type real.

X shall be of the same type and kind type parameter as Y. If Y has the value zero, X shall not have the value zero.

**4 Result Characteristics.** Same as X.

**5 Result Value.** The result has a value equal to a processor-dependent approximation to the principal value of the argument of the complex number (X, Y), expressed in radians. It lies in the range  $-\pi \leq \text{ATAN2}(Y,X)$

$\leq \pi$  and is equal to a processor-dependent approximation to a value of  $\arctan(Y/X)$  if  $X \neq 0$ . If  $Y > 0$ , the result is positive. If  $Y = 0$  and  $X > 0$ , the result is  $Y$ . If  $Y = 0$  and  $X < 0$ , then the result is approximately  $\pi$  if  $Y$  is positive real zero or the processor does not distinguish between positive and negative real zero, and approximately  $-\pi$  if  $Y$  is negative real zero. If  $Y < 0$ , the result is negative. If  $X = 0$ , the absolute value of the result is approximately  $\pi/2$ .

**Examples.**  $\text{ATAN2}(1.5574077, 1.0)$  has the value 1.0 (approximately). If  $Y$  has the value  $\begin{bmatrix} 1 & 1 \\ -1 & -1 \end{bmatrix}$  and  $X$  has the value  $\begin{bmatrix} -1 & 1 \\ -1 & 1 \end{bmatrix}$ , the value of  $\text{ATAN2}(Y, X)$  is approximately  $\begin{bmatrix} 3\pi/4 & \pi/4 \\ -3\pi/4 & -\pi/4 \end{bmatrix}$ .

### 13.9.19 ATANH (X)

**Description.** Inverse hyperbolic tangent function.

**Class.** [Elemental](#) function.

**Argument.**  $X$  shall be of type real or complex.

**Result Characteristics.** Same as  $X$ .

**Result Value.** The result has a value equal to a processor-dependent approximation to the inverse hyperbolic tangent function of  $X$ . If the result is complex the imaginary part is expressed in radians and lies in the range  $-\pi/2 \leq \text{AIMAG}(\text{ATANH}(X)) \leq \pi/2$ .

**Example.**  $\text{ATANH}(0.76159416)$  has the value 1.0 (approximately).

### 13.9.20 ATOMIC\_ADD (ATOM, VALUE [, STAT])

**Description.** Atomic addition.

**Class.** [Atomic subroutine](#).

**Arguments.**

**ATOM** shall be a scalar [coarray](#) or [coindexed object](#). It shall be of type integer with kind [ATOMIC\\_INT\\_KIND](#) from the intrinsic module [ISO\\_FORTRAN\\_ENV](#). It is an [INTENT \(INOUT\)](#) argument. If an error condition occurs, **ATOM** becomes undefined; otherwise, it becomes defined with the value of  $\text{ATOM} + \text{VALUE}$ .

**VALUE** shall be an integer scalar. It is an [INTENT \(IN\)](#) argument. The values of **VALUE** and **ATOM** + **VALUE** shall be representable in [ATOMIC\\_INT\\_KIND](#).

**STAT** (optional) shall be a noncoindexed integer scalar with a decimal range of at least four. It is an [INTENT \(OUT\)](#) argument. It is assigned a value as specified in [13.5](#). If an error condition occurs and **STAT** is not present, [error termination](#) is initiated.

**Example.**  $\text{CALL ATOMIC\_ADD}(I[3], 42)$  will cause  $I$  on [image 3](#) to become defined with the value 46 if the value of  $I[3]$  is 4 when the atomic operation is executed.

### 13.9.21 ATOMIC\_AND (ATOM, VALUE [, STAT])

**Description.** Atomic bitwise AND.

**Class.** [Atomic subroutine](#).

**Arguments.**

**ATOM** shall be a scalar [coarray](#) or [coindexed object](#). It shall be of type integer with kind [ATOMIC\\_INT\\_KIND](#) from the intrinsic module [ISO\\_FORTRAN\\_ENV](#). It is an [INTENT \(INOUT\)](#) argument. If an error condition occurs, **ATOM** becomes undefined; otherwise, it becomes defined with the value

of **IA**ND (ATOM, INT (VALUE, **ATOMIC\_INT\_KIND**)).

VALUE shall be an integer scalar. It is an **INTENT (IN)** argument. The value of VALUE shall be representable in **ATOMIC\_INT\_KIND**.

STAT (optional) shall be a noncoindexed integer scalar with a decimal range of at least four. It is an **INTENT (OUT)** argument. It is assigned a value as specified in 13.5. If an error condition occurs and STAT is not present, **error termination** is initiated.

4 **Example.** CALL ATOMIC\_AND (I [3], 6) will cause I on **image** 3 to become defined with the value 4 if the value of I [3] is 5 when the atomic operation is executed.

### 13.9.22 ATOMIC\_CAS (ATOM, OLD, COMPARE, NEW [, STAT])

1 **Description.** Atomic compare and swap.

2 **Class.** **Atomic subroutine.**

3 **Arguments.**

ATOM shall be a scalar **coarray** or **coindexed object**. It shall be of type integer with kind **ATOMIC\_INT\_KIND** from the intrinsic module **ISO\_FORTRAN\_ENV**, or of type logical with kind **ATOMIC\_LOGICAL\_KIND** from the intrinsic module **ISO\_FORTRAN\_ENV**. It is an **INTENT (INOUT)** argument. If an error condition occurs, ATOM becomes undefined; otherwise, if ATOM is of type integer and equal to COMPARE, or of type logical and equivalent to COMPARE, it becomes defined with the value of NEW.

OLD shall be scalar and of the same type and kind as ATOM. It is an **INTENT (OUT)** argument. If an error condition occurs, it becomes undefined; otherwise, it becomes defined with the value that ATOM had at the start of the atomic operation.

COMPARE shall be scalar and of the same type and kind as ATOM. It is an **INTENT (IN)** argument.

NEW shall be scalar and of the same type and kind as ATOM. It is an **INTENT (IN)** argument.

STAT (optional) shall be a noncoindexed integer scalar with a decimal range of at least four. It is an **INTENT (OUT)** argument. It is assigned a value as specified in 13.5. If an error condition occurs and STAT is not present, **error termination** is initiated.

4 **Example.** If the value of I on **image** 3 is equal to 13 at the beginning of the atomic operation performed by CALL ATOMIC\_CAS (I [3], OLD, 0, 1), the value of I on **image** 3 will be unchanged, and OLD will become defined with the value 13. If the value of I on **image** 3 is equal to 0 at the beginning of the atomic operation performed by CALL ATOMIC\_CAS (I [3], OLD, 0, 1), I on **image** 3 will become defined with the value 1, and OLD will become defined with the value 0.

### 13.9.23 ATOMIC\_DEFINE (ATOM, VALUE [, STAT])

1 **Description.** Define a variable atomically.

2 **Class.** **Atomic subroutine.**

3 **Arguments.**

ATOM shall be a scalar **coarray** or **coindexed object**. It shall be of type integer with kind **ATOMIC\_INT\_KIND** from the intrinsic module **ISO\_FORTRAN\_ENV**, or of type logical with kind **ATOMIC\_LOGICAL\_KIND** from the intrinsic module **ISO\_FORTRAN\_ENV**. It is an **INTENT (OUT)** argument. On successful execution, it becomes defined with the value of VALUE. If an error condition occurs, it becomes undefined.

VALUE shall be scalar and of the same type as ATOM. It is an **INTENT (IN)** argument.

STAT (optional) shall be a noncoindexed integer scalar with a decimal range of at least four. It is an **INTENT (OUT)** argument. It is assigned a value as specified in 13.5. If an error condition occurs and STAT is not present, **error termination** is initiated.

1 4 **Example.** CALL ATOMIC\_DEFINE (I [3], 4) causes I on [image 3](#) to become defined with the value 4.

## 2 13.9.24 ATOMIC\_FETCH\_ADD (ATOM, VALUE, OLD [, STAT])

3 1 **Description.** Atomic fetch and add.

4 2 **Class.** [Atomic subroutine](#).

5 3 **Arguments.**

6 ATOM shall be a scalar [coarray](#) or [coindexed object](#). It shall be of type integer with kind [ATOMIC\\_INT\\_KIND](#) from the intrinsic module [ISO\\_FORTRAN\\_ENV](#). It is an [INTENT \(INOUT\)](#) argument. If  
7 an error condition occurs, ATOM becomes undefined; otherwise, it becomes defined with the value  
8 of ATOM + VALUE.

10 VALUE shall be an integer scalar. It is an [INTENT \(IN\)](#) argument. The values of VALUE and ATOM +  
11 VALUE shall be representable in [ATOMIC\\_INT\\_KIND](#).

12 OLD shall be scalar and of the same type and kind as ATOM. It is an [INTENT \(OUT\)](#) argument. If  
13 an error condition occurs, it becomes undefined; otherwise, it becomes defined with the value that  
14 ATOM had at the start of the atomic operation.

15 STAT (optional) shall be a noncoindexed integer scalar with a decimal range of at least four. It is an [INTENT](#)  
16 [\(OUT\)](#) argument. It is assigned a value as specified in [13.5](#). If an error condition occurs and STAT  
17 is not present, [error termination](#) is initiated.

18 4 **Example.** CALL ATOMIC\_FETCH\_ADD (I [3], 7, J) will cause I on [image 3](#) to become defined with the value  
19 12, and J to become defined with the value 5, if the value of I [3] is 5 when the atomic operation is executed.

## 20 13.9.25 ATOMIC\_FETCH\_AND (ATOM, VALUE, OLD [, STAT])

21 1 **Description.** Atomic fetch and bitwise AND.

22 2 **Class.** [Atomic subroutine](#).

23 3 **Arguments.**

24 ATOM shall be a scalar [coarray](#) or [coindexed object](#). It shall be of type integer with kind [ATOMIC\\_INT\\_-](#)  
25 [KIND](#) from the intrinsic module [ISO\\_FORTRAN\\_ENV](#). It is an [INTENT \(INOUT\)](#) argument. If  
26 an error condition occurs, ATOM becomes undefined; otherwise, it becomes defined with the value  
27 of [IAND](#) (ATOM, [INT](#) (VALUE, [ATOMIC\\_INT\\_KIND](#))).

28 VALUE shall be an integer scalar. It is an [INTENT \(IN\)](#) argument. The value of VALUE shall be repres-  
29 entable in [ATOMIC\\_INT\\_KIND](#).

30 OLD shall be scalar and of the same type and kind as ATOM. It is an [INTENT \(OUT\)](#) argument. If  
31 an error condition occurs, it becomes undefined; otherwise, it becomes defined with the value that  
32 ATOM had at the start of the atomic operation.

33 STAT (optional) shall be a noncoindexed integer scalar with a decimal range of at least four. It is an [INTENT](#)  
34 [\(OUT\)](#) argument. It is assigned a value as specified in [13.5](#). If an error condition occurs and STAT  
35 is not present, [error termination](#) is initiated.

36 4 **Example.** CALL ATOMIC\_FETCH\_AND (I [3], 6, J) will cause I on [image 3](#) to become defined with the value  
37 4, and J to become defined with the value 5, if the value of I [3] is 5 when the atomic operation is executed.

## 38 13.9.26 ATOMIC\_FETCH\_OR (ATOM, VALUE, OLD [, STAT])

39 1 **Description.** Atomic fetch and bitwise OR.

40 2 **Class.** [Atomic subroutine](#).

41 3 **Arguments.**

1     ATOM       shall be a scalar [coarray](#) or [coindexed object](#). It shall be of type integer with kind [ATOMIC\\_INT\\_KIND](#) from the intrinsic module [ISO\\_FORTRAN\\_ENV](#). It is an [INTENT \(INOUT\)](#) argument. If  
2     an error condition occurs, ATOM becomes undefined; otherwise, it becomes defined with the value  
3     of [IOR](#) (ATOM, [INT](#) (VALUE, [ATOMIC\\_INT\\_KIND](#))).  
4  
5     VALUE       shall be an integer scalar. It is an [INTENT \(IN\)](#) argument. The value of VALUE shall be repres-  
6     entable in [ATOMIC\\_INT\\_KIND](#).  
7  
8     OLD         shall be scalar and of the same type and kind as ATOM. It is an [INTENT \(OUT\)](#) argument. If  
9     an error condition occurs, it becomes undefined; otherwise, it becomes defined with the value that  
10    ATOM had at the start of the atomic operation.  
11  
12    STAT (optional) shall be a noncoindexed integer scalar with a decimal range of at least four. It is an [INTENT \(OUT\)](#) argument. It is assigned a value as specified in [13.5](#). If an error condition occurs and STAT is not present, [error termination](#) is initiated.

13   4 **Example.** CALL ATOMIC\_FETCH\_OR (I [3], 1, J) will cause I on [image 3](#) to become defined with the value  
14   3, and J to become defined with the value 2, if the value of I [3] is 2 when the atomic operation is executed.

### 15   13.9.27    ATOMIC\_FETCH\_XOR (ATOM, VALUE, OLD [, STAT])

16   1 **Description.** Atomic fetch and bitwise exclusive OR.

17   2 **Class.** [Atomic subroutine](#).

18   3 **Arguments.**

19    ATOM       shall be a scalar [coarray](#) or [coindexed object](#). It shall be of type integer with kind [ATOMIC\\_INT\\_KIND](#) from the intrinsic module [ISO\\_FORTRAN\\_ENV](#). It is an [INTENT \(INOUT\)](#) argument. If  
20    an error condition occurs, ATOM becomes undefined; otherwise, it becomes defined with the value  
21    of [IEOR](#) (ATOM, [INT](#) (VALUE, [ATOMIC\\_INT\\_KIND](#))).  
22  
23    VALUE       shall be an integer scalar. It is an [INTENT \(IN\)](#) argument. The value of VALUE shall be repres-  
24    entable in [ATOMIC\\_INT\\_KIND](#).  
25  
26    OLD         shall be scalar and of the same type and kind as ATOM. It is an [INTENT \(OUT\)](#) argument. If  
27    an error condition occurs, it becomes undefined; otherwise, it becomes defined with the value that  
28    ATOM had at the start of the atomic operation.  
29  
30    STAT (optional) shall be a noncoindexed integer scalar with a decimal range of at least four. It is an [INTENT \(OUT\)](#) argument. It is assigned a value as specified in [13.5](#). If an error condition occurs and STAT is not present, [error termination](#) is initiated.

31   4 **Example.** CALL ATOMIC\_FETCH\_XOR (I [3], 1, J) will cause I on [image 3](#) to become defined with the value  
32   2, and J to become defined with the value 3, if the value of I [3] is 3 when the atomic operation is executed.

### 33   13.9.28    ATOMIC\_OR (ATOM, VALUE [, STAT])

34   1 **Description.** Atomic bitwise OR.

35   2 **Class.** [Atomic subroutine](#).

36   3 **Arguments.**

37    ATOM       shall be a scalar [coarray](#) or [coindexed object](#). It shall be of type integer with kind [ATOMIC\\_INT\\_KIND](#) from the intrinsic module [ISO\\_FORTRAN\\_ENV](#). It is an [INTENT \(INOUT\)](#) argument. If  
38    an error condition occurs, ATOM becomes undefined; otherwise, it becomes defined with the value  
39    of [IOR](#) (ATOM, [INT](#) (VALUE, [ATOMIC\\_INT\\_KIND](#))).  
40  
41    VALUE       shall be an integer scalar. It is an [INTENT \(IN\)](#) argument. The value of VALUE shall be repres-  
42    entable in [ATOMIC\\_INT\\_KIND](#).  
43  
44    STAT (optional) shall be a noncoindexed integer scalar with a decimal range of at least four. It is an [INTENT \(OUT\)](#) argument. It is assigned a value as specified in [13.5](#). If an error condition occurs and STAT

is not present, [error termination](#) is initiated.

**Example.** CALL ATOMIC\_OR (I [3], 1) will cause I on [image 3](#) to become defined with the value 3 if the value of I [3] is 2 when the atomic operation is executed.

### 13.9.29 ATOMIC\_REF (VALUE, ATOM [, STAT])

**Description.** Reference a variable atomically.

**Class.** [Atomic subroutine](#).

**Arguments.**

VALUE shall be scalar and of the same type as ATOM. It is an [INTENT \(OUT\)](#) argument. On successful execution, it becomes defined with the value of ATOM. If an error condition occurs, it becomes undefined.

ATOM shall be a scalar [coarray](#) or [coindexed object](#). It shall be of type integer with kind [ATOMIC\\_INT\\_KIND](#) from the intrinsic module [ISO\\_FORTRAN\\_ENV](#), or of type logical with kind [ATOMIC\\_LOGICAL\\_KIND](#) from the intrinsic module [ISO\\_FORTRAN\\_ENV](#). It is an [INTENT \(IN\)](#) argument.

STAT (optional) shall be a noncoindexed integer scalar with a decimal range of at least four. It is an [INTENT \(OUT\)](#) argument. It is assigned a value as specified in [13.5](#). If an error condition occurs and STAT is not present, [error termination](#) is initiated.

**Example.** CALL ATOMIC\_REF (VAL, I [3]) causes VAL to become defined with the value of I on [image 3](#).

### 13.9.30 ATOMIC\_XOR (ATOM, VALUE [, STAT])

**Description.** Atomic bitwise exclusive OR.

**Class.** [Atomic subroutine](#).

**Arguments.**

ATOM shall be a scalar [coarray](#) or [coindexed object](#). It shall be of type integer with kind [ATOMIC\\_INT\\_KIND](#) from the intrinsic module [ISO\\_FORTRAN\\_ENV](#). It is an [INTENT \(INOUT\)](#) argument. If an error condition occurs, ATOM becomes undefined; otherwise, it becomes defined with the value of [IEOR](#) (ATOM, [INT](#) (VALUE, [ATOMIC\\_INT\\_KIND](#))).

VALUE shall be an integer scalar. It is an [INTENT \(IN\)](#) argument. The value of VALUE shall be representable in [ATOMIC\\_INT\\_KIND](#).

STAT (optional) shall be a noncoindexed integer scalar with a decimal range of at least four. It is an [INTENT \(OUT\)](#) argument. It is assigned a value as specified in [13.5](#). If an error condition occurs and STAT is not present, [error termination](#) is initiated.

**Example.** CALL ATOMIC\_XOR (I [3], 1) will cause I on [image 3](#) to become defined with the value 2 if the value of I [3] is 3 when the atomic operation is executed.

### 13.9.31 BESSEL\_J0 (X)

**Description.** Bessel function of the 1<sup>st</sup> kind, order 0.

**Class.** [Elemental function](#).

**Argument.** X shall be of type real.

**Result Characteristics.** Same as X.

**Result Value.** The result has a value equal to a processor-dependent approximation to the Bessel function of the first kind and order zero of X.

1 6 **Example.** BESSEL\_J0 (1.0) has the value 0.765 (approximately).

## 2 13.9.32 BESSEL\_J1 (X)

3 1 **Description.** Bessel function of the 1<sup>st</sup> kind, order 1.

4 2 **Class.** [Elemental](#) function.

5 3 **Argument.** X shall be of type real.

6 4 **Result Characteristics.** Same as X.

7 5 **Result Value.** The result has a value equal to a processor-dependent approximation to the Bessel function of the first kind and order one of X.

9 6 **Example.** BESSEL\_J1 (1.0) has the value 0.440 (approximately).

## 10 13.9.33 BESSEL\_JN (N, X) or BESSEL\_JN (N1, N2, X)

11 1 **Description.** Bessel functions of the 1<sup>st</sup> kind.

12 2 **Class.**

13 *Case (i):* BESSEL\_JN (N,X) is an [elemental](#) function.

14 *Case (ii):* BESSEL\_JN (N1,N2,X) is a [transformational function](#).

15 3 **Arguments.**

16 N shall be of type integer and nonnegative.

17 N1 shall be an integer scalar with a nonnegative value.

18 N2 shall be an integer scalar with a nonnegative value.

19 X shall be of type real; if the function is transformational, X shall be scalar.

20 4 **Result Characteristics.** Same type and kind as X. The result of BESSEL\_JN (N1, N2, X) is a rank-one array with extent [MAX](#) (N2–N1+1, 0).

22 5 **Result Value.**

23 *Case (i):* The result value of BESSEL\_JN (N, X) is a processor-dependent approximation to the Bessel function of the first kind and order N of X.

25 *Case (ii):* Element *i* of the result value of BESSEL\_JN (N1, N2, X) is a processor-dependent approximation to the Bessel function of the first kind and order N1+*i* – 1 of X.

27 6 **Example.** BESSEL\_JN (2, 1.0) has the value 0.115 (approximately).

## 28 13.9.34 BESSEL\_Y0 (X)

29 1 **Description.** Bessel function of the 2<sup>nd</sup> kind, order 0.

30 2 **Class.** [Elemental](#) function.

31 3 **Argument.** X shall be of type real. Its value shall be greater than zero.

32 4 **Result Characteristics.** Same as X.

33 5 **Result Value.** The result has a value equal to a processor-dependent approximation to the Bessel function of the second kind and order zero of X.

35 6 **Example.** BESSEL\_Y0 (1.0) has the value 0.088 (approximately).

### 13.9.35 BESSEL\_Y1 (X)

**Description.** Bessel function of the 2<sup>nd</sup> kind, order 1.

**Class.** [Elemental](#) function.

**Argument.** X shall be of type real. Its value shall be greater than zero.

**Result Characteristics.** Same as X.

**Result Value.** The result has a value equal to a processor-dependent approximation to the Bessel function of the second kind and order one of X.

**Example.** BESSEL\_Y1 (1.0) has the value  $-0.781$  (approximately).

### 13.9.36 BESSEL\_YN (N, X) or BESSEL\_YN (N1, N2, X)

**Description.** Bessel functions of the 2<sup>nd</sup> kind.

**Class.**

*Case (i):* BESSEL\_YN (N, X) is an [elemental](#) function.

*Case (ii):* BESSEL\_YN (N1, N2, X) is a [transformational function](#).

**Arguments.**

N shall be of type integer and nonnegative.

N1 shall be an integer scalar with a nonnegative value.

N2 shall be an integer scalar with a nonnegative value.

X shall be of type real; if the function is transformational, X shall be scalar. Its value shall be greater than zero.

**Result Characteristics.** Same type and kind as X. The result of BESSEL\_YN (N1, N2, X) is a rank-one array with extent [MAX](#) (N2–N1+1, 0).

**Result Value.**

*Case (i):* The result value of BESSEL\_YN (N, X) is a processor-dependent approximation to the Bessel function of the second kind and order N of X.

*Case (ii):* Element *i* of the result value of BESSEL\_YN (N1, N2, X) is a processor-dependent approximation to the Bessel function of the second kind and order  $N1+i-1$  of X.

**Example.** BESSEL\_YN (2, 1.0) has the value  $-1.651$  (approximately).

### 13.9.37 BGE (I, J)

**Description.** Bitwise greater than or equal to.

**Class.** [Elemental](#) function.

**Arguments.**

I shall be of type integer or a [boz-literal-constant](#).

J shall be of type integer or a [boz-literal-constant](#).

**Result Characteristics.** Default logical.

**Result Value.** The result is true if the sequence of bits represented by I is greater than or equal to the sequence of bits represented by J, according to the method of bit sequence comparison in [13.3.2](#); otherwise the result is false.

6 The interpretation of a *boz-literal-constant* as a sequence of bits is described in 4.7. The interpretation of an integer value as a sequence of bits is described in 13.3.

7 **Example.** If BIT\_SIZE (J) has the value 8, BGE (Z'FF', J) has the value true for any value of J. BGE (0, -1) has the value false.

### 13.9.38 BGT (I, J)

1 **Description.** Bitwise greater than.

2 **Class.** *Elemental* function.

3 **Arguments.**

I shall be of type integer or a *boz-literal-constant*.

J shall be of type integer or a *boz-literal-constant*.

4 **Result Characteristics.** Default logical.

5 **Result Value.** The result is true if the sequence of bits represented by I is greater than the sequence of bits represented by J, according to the method of bit sequence comparison in 13.3.2; otherwise the result is false.

6 The interpretation of a *boz-literal-constant* as a sequence of bits is described in 4.7. The interpretation of an integer value as a sequence of bits is described in 13.3.

7 **Example.** BGT (Z'FF', Z'FC') has the value true. BGT (0, -1) has the value false.

### 13.9.39 BIT\_SIZE (I)

1 **Description.** Number of bits in integer model 13.3.

2 **Class.** *Inquiry* function.

3 **Argument.** I shall be of type integer. It may be a scalar or an array.

4 **Result Characteristics.** Scalar integer with the same kind type parameter as I.

5 **Result Value.** The result has the value of the number of bits *z* of the model integer defined for bit manipulation contexts in 13.3.

6 **Example.** BIT\_SIZE (1) has the value 32 if *z* of the model is 32.

### 13.9.40 BLE (I, J)

1 **Description.** Bitwise less than or equal to.

2 **Class.** *Elemental* function.

3 **Arguments.**

I shall be of type integer or a *boz-literal-constant*.

J shall be of type integer or a *boz-literal-constant*.

4 **Result Characteristics.** Default logical.

5 **Result Value.** The result is true if the sequence of bits represented by I is less than or equal to the sequence of bits represented by J, according to the method of bit sequence comparison in 13.3.2; otherwise the result is false.

6 The interpretation of a *boz-literal-constant* as a sequence of bits is described in 4.7. The interpretation of an integer value as a sequence of bits is described in 13.3.

1 7 **Example.** BLE (0, J) has the value true for any value of J. BLE (−1, 0) has the value false.

## 2 13.9.41 BLT (I, J)

3 1 **Description.** Bitwise less than.

4 2 **Class.** [Elemental](#) function.

5 3 **Arguments.**

6 I shall be of type integer or a *boz-literal-constant*.

7 J shall be of type integer or a *boz-literal-constant*.

8 4 **Result Characteristics.** Default logical.

9 5 **Result Value.** The result is true if the sequence of bits represented by I is less than the sequence of bits  
10 represented by J, according to the method of bit sequence comparison in [13.3.2](#); otherwise the result is false.

11 6 The interpretation of a *boz-literal-constant* as a sequence of bits is described in [4.7](#). The interpretation of an  
12 integer value as a sequence of bits is described in [13.3](#).

13 7 **Example.** BLT (0, −1) has the value true. BLT (Z'FF', Z'FC') has the value false.

## 14 13.9.42 BTEST (I, POS)

15 1 **Description.** Test single bit in an integer.

16 2 **Class.** [Elemental](#) function.

17 3 **Arguments.**

18 I shall be of type integer.

19 POS shall be of type integer. It shall be nonnegative and be less than BIT\_SIZE (I).

20 4 **Result Characteristics.** Default logical.

21 5 **Result Value.** The result has the value true if bit POS of I has the value 1 and has the value false if bit POS  
22 of I has the value 0. The model for the interpretation of an integer value as a sequence of bits is in [13.3](#).

23 6 **Examples.** BTEST (8, 3) has the value true. If A has the value  $\begin{bmatrix} 1 & 2 \\ 3 & 4 \end{bmatrix}$ , the value of BTEST (A, 2) is  
24  $\begin{bmatrix} \text{false} & \text{false} \\ \text{false} & \text{true} \end{bmatrix}$  and the value of BTEST (2, A) is  $\begin{bmatrix} \text{true} & \text{false} \\ \text{false} & \text{false} \end{bmatrix}$ .

## 25 13.9.43 CEILING (A [, KIND])

26 1 **Description.** Least integer greater than or equal to A.

27 2 **Class.** [Elemental](#) function.

28 3 **Arguments.**

29 A shall be of type real.

30 KIND (optional) shall be a scalar integer [constant expression](#).

31 4 **Result Characteristics.** Integer. If KIND is present, the kind type parameter is that specified by the value of  
32 KIND; otherwise, the kind type parameter is that of default integer type.

33 5 **Result Value.** The result has a value equal to the least integer greater than or equal to A.

1 6 **Examples.** CEILING (3.7) has the value 4. CEILING (−3.7) has the value −3.

## 2 13.9.44 CHAR (I [, KIND])

3 1 **Description.** Character from code value.

4 2 **Class.** [Elemental](#) function.

5 3 **Arguments.**

6 I shall be of type integer with a value in the range  $0 \leq I \leq n - 1$ , where  $n$  is the number of characters  
7 in the [collating sequence](#) associated with the specified kind type parameter.

8 KIND (optional) shall be a scalar integer [constant expression](#).

9 4 **Result Characteristics.** Character of length one. If KIND is present, the kind type parameter is that specified  
10 by the value of KIND; otherwise, the kind type parameter is that of default character.

11 5 **Result Value.** The result is the character in position I of the [collating sequence](#) associated with the spe-  
12 cified kind type parameter. ICHAR (CHAR (I, KIND (C))) shall have the value I for  $0 \leq I \leq n - 1$  and  
13 CHAR (ICHAR (C), KIND (C)) shall have the value C for any character C capable of representation in the  
14 processor.

15 6 **Example.** CHAR (88) has the value 'X' on a processor using the [ASCII collating sequence](#) for default characters.

## 16 13.9.45 CMPLX (X [, KIND]) or CMPLX (X [, Y, KIND])

17 1 **Description.** Conversion to complex type.

18 2 **Class.** [Elemental](#) function.

19 3 **Arguments for CMPLX(X [, KIND]).**

20 X shall be of type complex.

21 KIND (optional) shall be a scalar integer [constant expression](#).

22 4 **Arguments for CMPLX(X [, Y, KIND]).**

23 X shall be of type integer or real, or a [boz-literal-constant](#).

24 Y (optional) shall be of type integer or real, or a [boz-literal-constant](#).

25 KIND (optional) shall be a scalar integer [constant expression](#).

26 5 **Result Characteristics.** The result is of type complex. If KIND is present, the kind type parameter is that  
27 specified by the value of KIND; otherwise, the kind type parameter is that of default real kind.

28 6 **Result Value.** If Y is absent and X is not complex, it is as if Y were present with the value zero. If KIND is  
29 absent, it is as if KIND were present with the value [KIND](#) (0.0). If X is complex, the result is the same as that  
30 of CMPLX ([REAL](#) (X), [AIMAG](#) (X), KIND). The result of CMPLX (X, Y, KIND) has the complex value whose  
31 real part is [REAL](#) (X, KIND) and whose imaginary part is [REAL](#) (Y, KIND).

32 7 **Example.** CMPLX (−3) has the value (−3.0, 0.0).

## 33 13.9.46 CO\_BROADCAST(A, SOURCE\_IMAGE [, STAT, ERRMSG])

34 1 **Description.** Broadcast value to images.

35 2 **Class.** [Collective subroutine](#).

36 3 **Arguments.**

37 A shall have the same shape, dynamic type, and type parameter values, in corresponding references  
38 on all [active images](#) in the [current team](#). It shall not be a coindexed object. It is an [INTENT](#)

(**INOUT**) argument. If no error condition occurs, A becomes defined, as if by intrinsic assignment, on all images in the **current team** with the value of A on image `SOURCE_IMAGE`.

`SOURCE_IMAGE` shall be an integer scalar. It is an **INTENT (IN)** argument. Its value shall be that of an image index of an image in the current team. The value shall be the same in all corresponding references.

`STAT` (optional) shall be a noncoindexed integer scalar with a decimal exponent range of at least four. It is an **INTENT (OUT)** argument.

`ERRMSG` (optional) shall be a noncoindexed default character scalar. It is an **INTENT (INOUT)** argument.

4 The semantics of `STAT` and `ERRMSG` are described in 13.6.

5 **Example.** If A is the array [1, 5, 3] on image one, after execution of

```
CALL CO_BROADCAST (A, 1)
```

the value of A on all images of the current team is [1, 5, 3].

### 13.9.47 CO\_MAX(A [, RESULT\_IMAGE, STAT, ERRMSG])

1 **Description.** Compute maximum value across images.

2 **Class.** Collective subroutine.

3 **Arguments.**

A shall be of type integer, real, or character. It shall have the same shape, type, and type parameter values, in corresponding references on all **active images** in the **current team**. It shall not be a coindexed object. It is an **INTENT (INOUT)** argument. If it is scalar, the computed value is equal to the maximum value of A in all corresponding references. If it is an array each element of the computed value is equal to the maximum value of all corresponding elements of A in all corresponding references.

The computed value is assigned to A if no error condition occurs, and either `RESULT_IMAGE` is absent, or the executing image is the one identified by `RESULT_IMAGE`. Otherwise, A becomes undefined.

`RESULT_IMAGE` (optional) shall be an integer scalar. It is an **INTENT (IN)** argument. Its presence, and value if present, shall be the same in all corresponding references. If it is present, its value shall be that of an image index in the current team.

`STAT` (optional) shall be a noncoindexed integer scalar with a decimal exponent range of at least four. It is an **INTENT (OUT)** argument.

`ERRMSG` (optional) shall be a noncoindexed default character scalar. It is an **INTENT (INOUT)** argument.

4 The semantics of `STAT` and `ERRMSG` are described in 13.6.

5 **Example.** If the number of images in the current team is two and A is the array [1, 5, 3] on one image and [4, 1, 6] on the other image, the value of A after executing the statement `CALL CO_MAX (A)` is [4, 5, 6] on both images.

### 13.9.48 CO\_MIN(A [, RESULT\_IMAGE, STAT, ERRMSG])

1 **Description.** Compute minimum value across images.

2 **Class.** Collective subroutine.

3 **Arguments.**

A shall be of type integer, real, or character. It shall have the same shape, type, and type parameter values, in corresponding references on all **active images** in the **current team**. It shall not be a coindexed object. It is an **INTENT (INOUT)** argument. If it is scalar, the computed value is

equal to the minimum value of A in all corresponding references. If it is an array each element of the computed value is equal to the minimum value of all corresponding elements of A in all corresponding references.

The computed value is assigned to A if no error condition occurs, and either RESULT\_IMAGE is absent, or the executing image is the one identified by RESULT\_IMAGE. Otherwise, A becomes undefined.

RESULT\_IMAGE (optional) shall be an integer scalar. It is an **INTENT (IN)** argument. Its presence, and value if present, shall be the same in all corresponding references. If it is present, its value shall be that of an image index in the current team.

STAT (optional) shall be a noncoindexed integer scalar with a decimal exponent range of at least four. It is an **INTENT (OUT)** argument.

ERRMSG (optional) shall be a noncoindexed default character scalar. It is an **INTENT (INOUT)** argument.

4 The semantics of STAT and ERRMSG are described in 13.6.

5 **Example.** If the number of images in the current team is two and A is the array [1, 5, 3] on one image and [4, 1, 6] on the other image, the value of A after executing the statement CALL CO\_MIN (A) is [1, 1, 3] on both images.

### 13.9.49 CO\_REDUCE(A, OPERATOR [, RESULT\_IMAGE, STAT, ERRMSG])

1 **Description.** Generalized reduction across images.

2 **Class.** Collective subroutine.

3 **Arguments.**

A shall not be polymorphic. It shall have the same shape, type, and type parameter values, in corresponding references. It shall not be a coindexed object. It is an **INTENT (INOUT)** argument. If A is scalar, the computed value is the result of the reduction operation of applying OPERATOR to the values of A in all corresponding references. If A is an array, each element of the computed value is equal to the result of the reduction operation of applying OPERATOR to corresponding elements of A in all corresponding references.

The computed value is assigned to A if no error condition occurs, and either RESULT\_IMAGE is absent, or the executing image is the one identified by RESULT\_IMAGE. Otherwise, A becomes undefined.

OPERATOR shall be a pure function with two scalar, nonallocatable, nonpointer, nonpolymorphic, nonoptional arguments of the same type and type parameters as A. Its result shall have the same type and type parameters as A. If one argument has the **ASYNCHRONOUS**, **TARGET**, or **VALUE** attribute, the other shall have that attribute. OPERATOR shall implement a mathematically commutative and associative operation. OPERATOR shall be the same function on all images in corresponding references.

The computed value of a reduction operation over a set of values is the result of an iterative process. Each iteration involves the evaluation of OPERATOR ( $x, y$ ) for  $x$  and  $y$  in the set, the removal of  $x$  and  $y$  from the set, and the addition of the value of OPERATOR ( $x, y$ ) to the set. The process terminates when the set has only one element; this is the computed value.

RESULT\_IMAGE (optional) shall be an integer scalar. It is an **INTENT (IN)** argument. Its presence, and value if present, shall be the same in all corresponding references. If it is present, its value shall be that of an image index in the current team.

STAT (optional) shall be a noncoindexed integer scalar with a decimal exponent range of at least four. It is an **INTENT (OUT)** argument.

ERRMSG (optional) shall be a noncoindexed default character scalar. It is an **INTENT (INOUT)** argument.

4 The semantics of STAT and ERRMSG are described in 13.6.

1 5 **Example.** The subroutine below demonstrates how to use CO\_REDUCE to create a collective counterpart to  
2 the intrinsic function ALL:

```
3      SUBROUTINE co_all(boolean)
4          LOGICAL, INTENT(INOUT) :: boolean
5          CALL CO_REDUCE(boolean,both)
6      CONTAINS
7          PURE FUNCTION both(lhs,rhs) RESULT(lhs_and_rhs)
8              LOGICAL, INTENT(IN) :: lhs,rhs
9              LOGICAL :: lhs_and_rhs
10             lhs_and_rhs = lhs .AND. rhs
11         END FUNCTION both
12     END SUBROUTINE co_all
```

### 13 13.9.50 CO\_SUM(A [, RESULT\_IMAGE, STAT, ERRMSG])

14 1 **Description.** Compute sum across images.

15 2 **Class.** Collective subroutine.

16 3 **Arguments.**

17 A shall be of numeric type. It shall have the same shape, type, and type parameter values, in cor-  
18 responding references. It shall not be a coindexed object. It is an **INTENT (INOUT)** argument.  
19 If it is scalar, the computed value is equal to a processor-dependent approximation to the sum of  
20 the values of A in corresponding references. If it is an array, each element of the computed value  
21 is equal to a processor-dependent approximation to the sum of all corresponding elements of A in  
22 corresponding.

23 The computed value is assigned to A if no error condition occurs, and either RESULT\_IMAGE is  
24 absent, or the executing image is the one identified by RESULT\_IMAGE. Otherwise, A becomes  
25 undefined.

26 RESULT\_IMAGE (optional) shall be an integer scalar. It is an **INTENT (IN)** argument. Its presence, and value  
27 if present, shall be the same in all corresponding references. If it is present, its value shall be that  
28 of an image index in the current team.

29 STAT (optional) shall be a noncoindexed integer scalar with a decimal exponent range of at least four. It is an  
30 **INTENT (OUT)** argument.

31 ERRMSG (optional) shall be a noncoindexed default character scalar. It is an **INTENT (INOUT)** argument.

32 4 The semantics of STAT and ERRMSG are described in 13.6.

33 5 **Example.** If the number of images in the current team is two and A is the array [1, 5, 3] on one image and [4,  
34 1, 6] on the other image, the value of A after executing the statement CALL CO\_SUM(A) is [5, 6, 9] on both  
35 images.

### 36 13.9.51 COMMAND\_ARGUMENT\_COUNT ( )

37 1 **Description.** Number of command arguments.

38 2 **Class.** Transformational function.

39 3 **Argument.** None.

40 4 **Result Characteristics.** Default integer scalar.

41 5 **Result Value.** The result value is equal to the number of command arguments available. If there are no

1 command arguments available or if the processor does not support command arguments, then the result has the  
 2 value zero. If the processor has a concept of a command name, the command name does not count as one of the  
 3 command arguments.

4 6 **Example.** See [13.9.83](#).

## 5 **13.9.52 CONJG (Z)**

6 1 **Description.** Conjugate of a complex number.

7 2 **Class.** [Elemental](#) function.

8 3 **Argument.** Z shall be of type complex.

9 4 **Result Characteristics.** Same as Z.

10 5 **Result Value.** If Z has the value  $(x, y)$ , the result has the value  $(x, -y)$ .

11 6 **Example.** CONJG ((2.0, 3.0)) has the value (2.0, -3.0).

## 12 **13.9.53 COS (X)**

13 1 **Description.** Cosine function.

14 2 **Class.** [Elemental](#) function.

15 3 **Argument.** X shall be of type real or complex.

16 4 **Result Characteristics.** Same as X.

17 5 **Result Value.** The result has a value equal to a processor-dependent approximation to  $\cos(X)$ . If X is of type  
 18 real, it is regarded as a value in radians. If X is of type complex, its real part is regarded as a value in radians.

19 6 **Example.** COS (1.0) has the value 0.54030231 (approximately).

## 20 **13.9.54 COSH (X)**

21 1 **Description.** Hyperbolic cosine function.

22 2 **Class.** [Elemental](#) function.

23 3 **Argument.** X shall be of type real or complex.

24 4 **Result Characteristics.** Same as X.

25 5 **Result Value.** The result has a value equal to a processor-dependent approximation to  $\cosh(X)$ . If X is of type  
 26 complex its imaginary part is regarded as a value in radians.

27 6 **Example.** COSH (1.0) has the value 1.5430806 (approximately).

## 28 **13.9.55 COSHAPE (COARRAY [, KIND])**

29 1 **Description.** Sizes of [codimensions](#) of a [coarray](#).

30 2 **Class.** [Inquiry function](#).

31 3 **Arguments.**

32 COARRAY shall be a [coarray](#) of any type. It shall not be an unallocated [allocatable coarray](#).

33 4 **Result Characteristics.** Integer. If KIND is present, the kind type parameter is that specified by the value

of KIND; otherwise the kind type parameter is that of default integer type. The result is an array of **rank** one whose size is equal to the **corank** of COARRAY.

**5 Result Value.** The result has a value whose  $i^{th}$  element is equal to the size of the  $i^{th}$  **codimension** of COARRAY, as given by  $UCOBOUND(COARRAY,i) - LCOBOUND(COARRAY,i) + 1$ .

**6 Example.**

The following code allocates the coarray D with the same size in each codimension as that of the coarray C, with the lower cobound 1.

```
REAL, ALLOCATABLE :: C[:,:], D[:,:]
INTEGER, ALLOCATABLE :: COSHAPE_C(:)
...
COSHAPE_C = COSHAPE(C)
ALLOCATE ( D[COSHAPE_C(1),*] )
```

### 13.9.56 COUNT (MASK [, DIM, KIND])

**1 Description.** Logical array reduced by counting true values.

**2 Class.** Transformational function.

**3 Arguments.**

**MASK** shall be a logical array.

**DIM** (optional) shall be an integer scalar with a value in the range  $1 \leq \text{DIM} \leq n$ , where  $n$  is the **rank** of MASK. The corresponding **actual argument** shall not be an optional dummy argument, a disassociated pointer, or an unallocated allocatable.

**KIND** (optional) shall be a scalar integer **constant expression**.

**4 Result Characteristics.** Integer. If KIND is present, the kind type parameter is that specified by the value of KIND; otherwise the kind type parameter is that of default integer type. The result is scalar if DIM is absent or  $n = 1$ ; otherwise, the result has **rank**  $n - 1$  and shape  $[d_1, d_2, \dots, d_{\text{DIM}-1}, d_{\text{DIM}+1}, \dots, d_n]$  where  $[d_1, d_2, \dots, d_n]$  is the shape of MASK.

**5 Result Value.**

*Case (i):* If DIM is absent or MASK has rank one, the result has a value equal to the number of true elements of MASK or has the value zero if MASK has size zero.

*Case (ii):* If DIM is present and MASK has rank  $n > 1$ , the value of element  $(s_1, s_2, \dots, s_{\text{DIM}-1}, s_{\text{DIM}+1}, \dots, s_n)$  of the result is equal to the number of true elements of MASK  $(s_1, s_2, \dots, s_{\text{DIM}-1}, :, s_{\text{DIM}+1}, \dots, s_n)$ .

**6 Examples.**

*Case (i):* The value of COUNT ( [.TRUE., .FALSE., .TRUE.] ) is 2.

*Case (ii):* If B is the array  $\begin{bmatrix} 1 & 3 & 5 \\ 2 & 4 & 6 \end{bmatrix}$  and C is the array  $\begin{bmatrix} 0 & 3 & 5 \\ 7 & 4 & 8 \end{bmatrix}$ , COUNT ( B /= C, DIM = 1 ) is [2, 0, 1] and COUNT ( B /= C, DIM = 2 ) is [1, 2].

### 13.9.57 CPU\_TIME (TIME)

**1 Description.** Processor time used.

**2 Class.** Subroutine.

**3 Argument.** TIME shall be a real scalar. It is an **INTENT (OUT)** argument. If the processor cannot provide a meaningful value for the time, it is assigned a processor-dependent negative value; otherwise, it is assigned a

1 processor-dependent approximation to the processor time in seconds. Whether the value assigned is an approx-  
 2 imation to the amount of time used by the invoking [image](#), or the amount of time used by the whole program, is  
 3 processor dependent.

#### 4 **Example.**

```
5      REAL T1, T2
6      ...
7      CALL CPU_TIME(T1)
8      ... ! Code to be timed.
9      CALL CPU_TIME(T2)
10     WRITE (*,*) 'Time taken by code was ', T2-T1, ' seconds'
```

11 writes the processor time taken by a piece of code.

#### **NOTE 13.10**

A processor for which a single result is inadequate (for example, a parallel processor) might choose to provide an additional version for which time is an array.

The exact definition of time is left imprecise because of the variability in what different processors are able to provide. The primary purpose is to compare different algorithms on the same processor or discover which parts of a calculation are the most expensive.

The start time is left imprecise because the purpose is to time sections of code, as in the example.

Most computer systems have multiple concepts of time. One common concept is that of time expended by the processor for a given program. This might or might not include system overhead, and has no obvious connection to elapsed “wall clock” time.

## 12 **13.9.58 CSHIFT (ARRAY, SHIFT [, DIM])**

13 1 **Description.** Circular shift of an array.

14 2 **Class.** [Transformational function](#).

15 3 **Arguments.**

16 ARRAY may be of any type. It shall be an array.

17 SHIFT shall be of type integer and shall be scalar if ARRAY has [rank](#) one; otherwise, it shall be scalar or  
 18 of [rank](#)  $n - 1$  and of shape  $[d_1, d_2, \dots, d_{\text{DIM}-1}, d_{\text{DIM}+1}, \dots, d_n]$  where  $[d_1, d_2, \dots, d_n]$  is the shape  
 19 of ARRAY.

20 DIM (optional) shall be an integer scalar with a value in the range  $1 \leq \text{DIM} \leq n$ , where  $n$  is the [rank](#) of ARRAY.  
 21 If DIM is absent, it is as if it were present with the value 1.

22 4 **Result Characteristics.** The result is of the type and type parameters of ARRAY, and has the shape of  
 23 ARRAY.

24 5 **Result Value.**

25 *Case (i):* If ARRAY has [rank](#) one, element  $i$  of the result is  $\text{ARRAY}(1 + \text{MODULO}(i + \text{SHIFT} - 1, \text{SIZE}(\text{ARRAY})))$ .  
 26

27 *Case (ii):* If ARRAY has [rank](#) greater than one, section  $(s_1, s_2, \dots, s_{\text{DIM}-1}, :, s_{\text{DIM}+1}, \dots, s_n)$  of the result  
 28 has a value equal to  $\text{CSHIFT}(\text{ARRAY}(s_1, s_2, \dots, s_{\text{DIM}-1}, :, s_{\text{DIM}+1}, \dots, s_n), sh, 1)$ , where  $sh$  is  
 29 SHIFT or SHIFT  $(s_1, s_2, \dots, s_{\text{DIM}-1}, s_{\text{DIM}+1}, \dots, s_n)$ .

30 6 **Examples.**

*Case (i):* If  $V$  is the array  $[1, 2, 3, 4, 5, 6]$ , the effect of shifting  $V$  circularly to the left by two positions is achieved by  $\text{CSHIFT}(V, \text{SHIFT} = 2)$  which has the value  $[3, 4, 5, 6, 1, 2]$ ;  $\text{CSHIFT}(V, \text{SHIFT} = -2)$  achieves a circular shift to the right by two positions and has the value  $[5, 6, 1, 2, 3, 4]$ .

*Case (ii):* The rows of an array of [rank](#) two may all be shifted by the same amount or by different amounts.

If  $M$  is the array  $\begin{bmatrix} 1 & 2 & 3 \\ 4 & 5 & 6 \\ 7 & 8 & 9 \end{bmatrix}$ , the value of

$\text{CSHIFT}(M, \text{SHIFT} = -1, \text{DIM} = 2)$  is  $\begin{bmatrix} 3 & 1 & 2 \\ 6 & 4 & 5 \\ 9 & 7 & 8 \end{bmatrix}$ , and the value of

$\text{CSHIFT}(M, \text{SHIFT} = [-1, 1, 0], \text{DIM} = 2)$  is  $\begin{bmatrix} 3 & 1 & 2 \\ 5 & 6 & 4 \\ 7 & 8 & 9 \end{bmatrix}$ .

### 13.9.59 DATE\_AND\_TIME ([DATE, TIME, ZONE, VALUES])

1 **Description.** Date and time.

2 **Class.** Subroutine.

3 **Arguments.**

DATE (optional) shall be a default character scalar. It is an [INTENT \(OUT\)](#) argument. It is assigned a value of the form *YYYYMMDD*, where *YYYY* is the year in the Gregorian calendar, *MM* is the month within the year, and *DD* is the day within the month. The characters of this value shall all be decimal digits. If there is no date available, DATE is assigned all blanks.

TIME (optional) shall be a default character scalar. It is an [INTENT \(OUT\)](#) argument. It is assigned a value of the form *hhmmss.sss*, where *hh* is the hour of the day, *mm* is the minutes of the hour, and *ss.sss* is the seconds and milliseconds of the minute. The characters of this value shall all be decimal digits. If there is no clock available, TIME is assigned all blanks.

ZONE (optional) shall be a default character scalar. It is an [INTENT \(OUT\)](#) argument. It is assigned a value of the form *+hhmm* or *-hhmm*, where *hh* and *mm* are the time difference with respect to Coordinated Universal Time (UTC) in hours and minutes, respectively. The characters of this value following the sign character shall all be decimal digits. If this information is not available, ZONE is assigned all blanks.

VALUES (optional) shall be a rank-one array of type integer with a decimal exponent range of at least four. It is an [INTENT \(OUT\)](#) argument. Its size shall be at least 8. The values assigned to VALUES are as follows:

VALUES (1) the year, including the century (for example, 2008), or *-HUGE (VALUES)* if there is no date available;

VALUES (2) the month of the year, or *-HUGE (VALUES)* if there is no date available;

VALUES (3) the day of the month, or *-HUGE (VALUES)* if there is no date available;

VALUES (4) the time difference with respect to Coordinated Universal Time (UTC) in minutes, or *-HUGE (VALUES)* if this information is not available;

VALUES (5) the hour of the day, in the range of 0 to 23, or *-HUGE (VALUES)* if there is no clock;

VALUES (6) the minutes of the hour, in the range 0 to 59, or *-HUGE (VALUES)* if there is no clock;

VALUES (7) the seconds of the minute, in the range 0 to 60, or *-HUGE (VALUES)* if there is no clock;

VALUES (8) the milliseconds of the second, in the range 0 to 999, or *-HUGE (VALUES)* if there is no clock.

4 The date, clock, and time zone information might be available on some [images](#) and not others. If the date, clock, or time zone information is available on more than one [image](#), it is processor dependent whether or not those [images](#) share the same information.

5 **Example.** If run in Geneva, Switzerland on April 12, 2008 at 15:27:35.5 with a system configured for the

1 local time zone, this example would have assigned the value 20080412 to BIG\_BEN (1), the value 152735.500 to  
 2 BIG\_BEN (2), the value +0100 to BIG\_BEN (3), and the value [2008, 4, 12, 60, 15, 27, 35, 500] to DATE\_TIME.

3 INTEGER DATE\_TIME (8)

4 CHARACTER (LEN = 10) BIG\_BEN (3)

5 CALL DATE\_AND\_TIME (BIG\_BEN (1), BIG\_BEN (2), BIG\_BEN (3), DATE\_TIME)

#### NOTE 13.11

These forms are compatible with the representations defined in ISO 8601:2004. UTC is established by the International Bureau of Weights and Measures (BIPM, i.e. Bureau International des Poids et Mesures) and the International Earth Rotation Service (IERS).

### 6 13.9.60 DBLE (A)

7 1 **Description.** Conversion to double precision real.

8 2 **Class.** [Elemental](#) function.

9 3 **Argument.** A shall be of type integer, real, complex, or a *boz-literal-constant*.

10 4 **Result Characteristics.** Double precision real.

11 5 **Result Value.** The result has the value REAL (A, KIND (0.0D0)).

12 6 **Example.** DBLE (−3) has the value −3.0D0.

### 13 13.9.61 DIGITS (X)

14 1 **Description.** Significant digits in numeric model.

15 2 **Class.** [Inquiry](#) function.

16 3 **Argument.** X shall be of type integer or real. It may be a scalar or an array.

17 4 **Result Characteristics.** Default integer scalar.

18 5 **Result Value.** The result has the value  $q$  if X is of type integer and  $p$  if X is of type real, where  $q$  and  $p$  are as  
 19 defined in [13.4](#) for the model representing numbers of the same type and kind type parameter as X.

20 6 **Example.** DIGITS (X) has the value 24 for real X whose model is as in Note [13.4](#).

### 21 13.9.62 DIM (X, Y)

22 1 **Description.** Maximum of  $X - Y$  and zero.

23 2 **Class.** [Elemental](#) function.

24 3 **Arguments.**

25 X shall be of type integer or real.

26 Y shall be of the same type and kind type parameter as X.

27 4 **Result Characteristics.** Same as X.

28 5 **Result Value.** The value of the result is the maximum of  $X - Y$  and zero.

29 6 **Example.** DIM (−3.0, 2.0) has the value 0.0.

### 13.9.63 DOT\_PRODUCT (VECTOR\_A, VECTOR\_B)

**1 Description.** Dot product of two vectors.

**2 Class.** Transformational function.

**3 Arguments.**

VECTOR\_A shall be of **numeric type** (integer, real, or complex) or of logical type. It shall be a rank-one array.

VECTOR\_B shall be of **numeric type** if VECTOR\_A is of **numeric type** or of type logical if VECTOR\_A is of type logical. It shall be a rank-one array. It shall be of the same size as VECTOR\_A.

**4 Result Characteristics.** If the arguments are of **numeric type**, the type and **kind type parameter** of the result are those of the expression VECTOR\_A \* VECTOR\_B determined by the types and kinds of the arguments according to 7.1.9.3. If the arguments are of type logical, the result is of type logical with the kind type parameter of the expression VECTOR\_A .AND. VECTOR\_B according to 7.1.9.3. The result is scalar.

**5 Result Value.**

*Case (i):* If VECTOR\_A is of type integer or real, the result has the value SUM (VECTOR\_A\*VECTOR\_B). If the vectors have size zero, the result has the value zero.

*Case (ii):* If VECTOR\_A is of type complex, the result has the value SUM (CONJG (VECTOR\_A)\*VECTOR\_B). If the vectors have size zero, the result has the value zero.

*Case (iii):* If VECTOR\_A is of type logical, the result has the value ANY (VECTOR\_A .AND. VECTOR\_B). If the vectors have size zero, the result has the value false.

**6 Example.** DOT\_PRODUCT ([1, 2, 3], [2, 3, 4]) has the value 20.

### 13.9.64 DPROD (X, Y)

**1 Description.** Double precision real product.

**2 Class.** Elemental function.

**3 Arguments.**

X shall be default real.

Y shall be default real.

**4 Result Characteristics.** Double precision real.

**5 Result Value.** The result has a value equal to a processor-dependent approximation to the product of X and Y. DPROD (X, Y) should have the same value as DBLE (X) \* DBLE (Y).

**6 Example.** DPROD (−3.0, 2.0) has the value −6.0D0.

### 13.9.65 DSHIFTL (I, J, SHIFT)

**1 Description.** Combined left shift.

**2 Class.** Elemental function.

**3 Arguments.**

I shall be of type integer or a *boz-literal-constant*.

J shall be of type integer or a *boz-literal-constant*. If both I and J are of type integer, they shall have the same kind type parameter. I and J shall not both be *boz-literal-constants*.

SHIFT shall be of type integer. It shall be nonnegative and less than or equal to BIT\_SIZE (I) if I is of type integer; otherwise, it shall be less than or equal to BIT\_SIZE (J).

1 4 **Result Characteristics.** Same as I if I is of type integer; otherwise, same as J.

2 5 **Result Value.** If either I or J is a *boz-literal-constant*, it is first converted as if by the intrinsic function `INT` to  
3 type integer with the kind type parameter of the other. The rightmost SHIFT bits of the result value are the same  
4 as the leftmost bits of J, and the remaining bits of the result value are the same as the rightmost bits of I. This  
5 is equal to `IOR (SHIFTL (I, SHIFT), SHIFTR (J, BIT_SIZE (J)–SHIFT))`. The model for the interpretation of  
6 an integer value as a sequence of bits is in 13.3.

7 6 **Examples.** `DSHIFTL (1, 2**30, 2)` has the value 5 if default integer has 32 bits. `DSHIFTL (I, I, SHIFT)` has  
8 the same result value as `ISHFTC (I, SHIFT)`.

### 9 13.9.66 DSHIFTR (I, J, SHIFT)

10 1 **Description.** Combined right shift.

11 2 **Class.** *Elemental* function.

12 3 **Arguments.**

13 I shall be of type integer or a *boz-literal-constant*.

14 J shall be of type integer or a *boz-literal-constant*. If both I and J are of type integer, they shall have  
15 the same kind type parameter. I and J shall not both be *boz-literal-constants*.

16 SHIFT shall be of type integer. It shall be nonnegative and less than or equal to `BIT_SIZE (I)` if I is of  
17 type integer; otherwise, it shall be less than or equal to `BIT_SIZE (J)`.

18 4 **Result Characteristics.** Same as I if I is of type integer; otherwise, same as J.

19 5 **Result Value.** If either I or J is a *boz-literal-constant*, it is first converted as if by the intrinsic function `INT` to  
20 type integer with the kind type parameter of the other. The leftmost SHIFT bits of the result value are the same  
21 as the rightmost bits of I, and the remaining bits of the result value are the same as the leftmost bits of J. This  
22 is equal to `IOR (SHIFTL (I, BIT_SIZE (I)–SHIFT), SHIFTR (J, SHIFT))`. The model for the interpretation of  
23 an integer value as a sequence of bits is in 13.3.

24 6 **Examples.** `DSHIFTR (1, 16, 3)` has the value  $2^{29} + 2$  if default integer has 32 bits. `DSHIFTR (I, I, SHIFT)` has  
25 the same result value as `ISHFTC (I, –SHIFT)`.

### 26 13.9.67 EOSHIFT (ARRAY, SHIFT [, BOUNDARY, DIM])

27 1 **Description.** End-off shift of the elements of an array.

28 2 **Class.** *Transformational* function.

29 3 **Arguments.**

30 ARRAY shall be an array be of any type.

31 SHIFT shall be of type integer and shall be scalar if ARRAY has *rank* one; otherwise, it shall be scalar or  
32 of *rank*  $n - 1$  and of shape  $[d_1, d_2, \dots, d_{\text{DIM}-1}, d_{\text{DIM}+1}, \dots, d_n]$  where  $[d_1, d_2, \dots, d_n]$  is the shape  
33 of ARRAY.

34 BOUNDARY (optional) shall be of the same type and type parameters as ARRAY and shall be scalar if ARRAY  
35 has *rank* one; otherwise, it shall be either scalar or of *rank*  $n - 1$  and of shape  $[d_1, d_2, \dots, d_{\text{DIM}-1},$   
36  $d_{\text{DIM}+1}, \dots, d_n]$ . BOUNDARY is permitted to be absent only for the types in Table 13.4, and in  
37 this case it is as if it were present with the scalar value shown, converted if necessary to the kind  
38 type parameter value of ARRAY.

Table 13.4: Default BOUNDARY values for EOSHIFT

Type of ARRAY	Value of BOUNDARY
Integer	0
Real	0.0

## Default BOUNDARY values for EOSHIFT.

Type of ARRAY	Value of BOUNDARY
Complex	(0.0, 0.0)
Logical	.FALSE.
Character ( <i>len</i> )	<i>len</i> blanks

DIM (optional) shall be an integer scalar with a value in the range  $1 \leq \text{DIM} \leq n$ , where  $n$  is the **rank** of ARRAY.  
If DIM is absent, it is as if it were present with the value 1.

**Result Characteristics.** The result has the type, type parameters, and shape of ARRAY.

**Result Value.** Element  $(s_1, s_2, \dots, s_n)$  of the result has the value  $\text{ARRAY}(s_1, s_2, \dots, s_{\text{DIM}-1}, s_{\text{DIM}} + sh, s_{\text{DIM}+1}, \dots, s_n)$  where  $sh$  is SHIFT or SHIFT  $(s_1, s_2, \dots, s_{\text{DIM}-1}, s_{\text{DIM}+1}, \dots, s_n)$  provided the inequality  $\text{LBOUND}(\text{ARRAY}, \text{DIM}) \leq s_{\text{DIM}} + sh \leq \text{UBOUND}(\text{ARRAY}, \text{DIM})$  holds and is otherwise BOUNDARY or BOUNDARY  $(s_1, s_2, \dots, s_{\text{DIM}-1}, s_{\text{DIM}+1}, \dots, s_n)$ .

#### 6 Examples.

*Case (i):* If V is the array [1, 2, 3, 4, 5, 6], the effect of shifting V end-off to the left by 3 positions is achieved by EOSHIFT (V, SHIFT = 3), which has the value [4, 5, 6, 0, 0, 0]; EOSHIFT (V, SHIFT = -2, BOUNDARY = 99) achieves an end-off shift to the right by 2 positions with the boundary value of 99 and has the value [99, 99, 1, 2, 3, 4].

*Case (ii):* The rows of an array of **rank** two may all be shifted by the same amount or by different amounts

and the boundary elements can be the same or different. If M is the array  $\begin{bmatrix} A & B & C \\ D & E & F \\ G & H & I \end{bmatrix}$ , then the

value of EOSHIFT (M, SHIFT = -1, BOUNDARY = '\*', DIM = 2) is  $\begin{bmatrix} * & A & B \\ * & D & E \\ * & G & H \end{bmatrix}$ , and the value

of EOSHIFT (M, SHIFT = [-1, 1, 0], BOUNDARY = ['\*', '/', '?'], DIM = 2) is  $\begin{bmatrix} * & A & B \\ E & F & / \\ G & H & I \end{bmatrix}$ .

### 13.9.68 EPSILON (X)

**Description.** Model number that is small compared to 1.

**Class.** Inquiry function.

**Argument.** X shall be of type real. It may be a scalar or an array.

**Result Characteristics.** Scalar of the same type and kind type parameter as X.

**Result Value.** The result has the value  $b^{1-p}$  where  $b$  and  $p$  are as defined in 13.4 for the model representing numbers of the same type and kind type parameter as X.

**Example.** EPSILON (X) has the value  $2^{-23}$  for real X whose model is as in Note 13.4.

### 13.9.69 ERF (X)

**Description.** Error function.

**Class.** Elemental function.

**Argument.** X shall be of type real.

**Result Characteristics.** Same as X.

1 5 **Result Value.** The result has a value equal to a processor-dependent approximation to the error function of  $X$ ,  
 2  $\frac{2}{\sqrt{\pi}} \int_0^X \exp(-t^2) dt$ .

3 6 **Example.** `ERF (1.0)` has the value 0.843 (approximately).

## 4 13.9.70 ERFC (X)

5 1 **Description.** Complementary error function.

6 2 **Class.** [Elemental](#) function.

7 3 **Argument.**  $X$  shall be of type real.

8 4 **Result Characteristics.** Same as  $X$ .

9 5 **Result Value.** The result has a value equal to a processor-dependent approximation to the complementary error  
 10 function of  $X$ ,  $1 - \text{ERF}(X)$ ; this is equivalent to  $\frac{2}{\sqrt{\pi}} \int_X^\infty \exp(-t^2) dt$ .

11 6 **Example.** `ERFC (1.0)` has the value 0.157 (approximately).

## 12 13.9.71 ERFC\_SCALED (X)

13 1 **Description.** Scaled complementary error function.

14 2 **Class.** [Elemental](#) function.

15 3 **Argument.**  $X$  shall be of type real.

16 4 **Result Characteristics.** Same as  $X$ .

17 5 **Result Value.** The result has a value equal to a processor-dependent approximation to the exponentially-scaled  
 18 complementary error function of  $X$ ,  $\exp(X^2) \frac{2}{\sqrt{\pi}} \int_X^\infty \exp(-t^2) dt$ .

19 6 **Example.** `ERFC_SCALED (20.0)` has the value 0.02817434874 (approximately).

### NOTE 13.12

The complementary error function is asymptotic to  $\exp(-X^2)/(X\sqrt{\pi})$ . As such it underflows for  $X > \approx 9$  when using ISO/IEC/IEEE 60559:2011 single precision arithmetic. The exponentially-scaled complementary error function is asymptotic to  $1/(X\sqrt{\pi})$ . As such it does not underflow until  $X > \text{HUGE}(X)/\sqrt{\pi}$ .

## 20 13.9.72 EVENT\_QUERY (EVENT, COUNT [, STAT])

21 1 **Description.** Query event count.

22 2 **Class.** Subroutine.

23 3 **Arguments.**

24 **EVENT** shall be an [event variable](#) (13.10.2.10). It shall not be [coindexed](#). It is an [INTENT \(IN\)](#) argument.  
 25 The **EVENT** argument is accessed atomically with respect to the execution of [EVENT POST](#)  
 26 [statements](#) in unordered segments, in exact analogy to [atomic subroutines](#).

27 **COUNT** shall be an integer scalar with a decimal range no smaller than that of default integer. It is an  
 28 [INTENT \(OUT\)](#) argument. If no error condition occurs, **COUNT** is assigned the value of the count  
 29 of **EVENT**; otherwise, it is assigned the value  $-1$ .

30 **STAT** (optional) shall be a noncoindexed integer scalar with a decimal exponent range of at least four. It is an  
 31 [INTENT \(OUT\)](#) argument. If the **STAT** argument is present, it is assigned a processor-dependent  
 32 positive value if an error condition occurs; otherwise it is assigned the value zero. If the **STAT**  
 33 argument is not present and an error condition occurs, [error termination](#) is initiated.

1 4 **Example.** If EVENT is an [event variable](#) for which there have been no successful posts or waits in preceding  
 2 segments, and for which there are no posts or waits in an unordered segment, after execution of

3 CALL EVENT\_QUERY (EVENT, COUNT)

4 the integer variable COUNT will have the value zero. If there have been ten successful posts to EVENT and two  
 5 successful waits without an [UNTIL\\_COUNT= specifier](#) in preceding segments, and for which there are no posts  
 6 or waits in an unordered segment, after execution of

7 CALL EVENT\_QUERY (EVENT, COUNT)

8 the variable COUNT will have the value eight.

#### NOTE 13.13

Execution of EVENT\_QUERY does not imply any synchronization.

### 9 13.9.73 EXECUTE\_COMMAND\_LINE (COMMAND [, WAIT, EXITSTAT, CMDSTAT, CMDMSG ])

10 1 **Description.** Execute a command line.

11 2 **Class.** Subroutine.

12 3 **Arguments.**

13 COMMAND shall be a default character scalar. It is an [INTENT \(IN\)](#) argument. Its value is the command line  
 14 to be executed. The interpretation is processor dependent.

15 WAIT (optional) shall be a logical scalar. It is an [INTENT \(IN\)](#) argument. If WAIT is present with the value  
 16 false, and the processor supports asynchronous execution of the command, the command is executed  
 17 asynchronously; otherwise it is executed synchronously.

18 EXITSTAT (optional) shall be a scalar of type integer with a decimal exponent range of at least nine. It is an  
 19 [INTENT \(INOUT\)](#) argument. If the command is executed synchronously, it is assigned the value  
 20 of the processor-dependent exit status. Otherwise, the value of EXITSTAT is unchanged.

21 CMDSTAT (optional) shall be a scalar of type integer with a decimal exponent range of at least four. It is an  
 22 [INTENT \(OUT\)](#) argument. It is assigned the value  $-1$  if the processor does not support command  
 23 line execution, a processor-dependent positive value if an error condition occurs, or the value  $-2$   
 24 if no error condition occurs but WAIT is present with the value false and the processor does not  
 25 support asynchronous execution. Otherwise it is assigned the value 0.

26 CMDMSG (optional) shall be a default character scalar. It is an [INTENT \(INOUT\)](#) argument. If an error condi-  
 27 tion occurs, it is assigned a processor-dependent explanatory message. Otherwise, it is unchanged.

28 4 If the processor supports command line execution, it shall support synchronous and may support asynchronous  
 29 execution of the command line.

30 5 When the command is executed synchronously, EXECUTE\_COMMAND\_LINE returns after the command line  
 31 has completed execution. Otherwise, EXECUTE\_COMMAND\_LINE returns without waiting.

32 6 If a condition occurs that would assign a nonzero value to CMDSTAT but the CMDSTAT variable is not present,  
 33 [error termination](#) is initiated.

### 34 13.9.74 EXP (X)

35 1 **Description.** Exponential function.

36 2 **Class.** [Elemental](#) function.

37 3 **Argument.** X shall be of type real or complex.

1 4 **Result Characteristics.** Same as X.

2 5 **Result Value.** The result has a value equal to a processor-dependent approximation to  $e^X$ . If X is of type  
3 complex, its imaginary part is regarded as a value in radians.

4 6 **Example.** EXP (1.0) has the value 2.7182818 (approximately).

### 5 13.9.75 EXPONENT (X)

6 1 **Description.** Exponent of floating-point number.

7 2 **Class.** [Elemental](#) function.

8 3 **Argument.** X shall be of type real.

9 4 **Result Characteristics.** Default integer.

10 5 **Result Value.** The result has a value equal to the exponent  $e$  of the representation for the value of X in the  
11 extended real model for the kind of X ([13.4](#)), provided X is nonzero and  $e$  is within the range for default integers.  
12 If X has the value zero, the result has the value zero. If X is an IEEE infinity or NaN, the result has the value  
13 HUGE (0).

14 6 **Examples.** EXPONENT (1.0) has the value 1 and EXPONENT (4.1) has the value 3 for reals whose model is  
15 as in Note [13.4](#).

### 16 13.9.76 EXTENDS\_TYPE\_OF (A, MOLD)

17 1 **Description.** [Dynamic type](#) extension inquiry.

18 2 **Class.** [Inquiry function](#).

19 3 **Arguments.**

20 A shall be an object of [extensible declared](#) type or [unlimited polymorphic](#). If it is a [polymorphic](#)  
21 pointer, it shall not have an undefined association status.

22 MOLD shall be an object of [extensible declared](#) type or [unlimited polymorphic](#). If it is a [polymorphic](#)  
23 pointer, it shall not have an undefined association status.

24 4 **Result Characteristics.** Default logical scalar.

25 5 **Result Value.** If MOLD is [unlimited polymorphic](#) and is either a [disassociated](#) pointer or unallocated [allocatable](#)  
26 variable, the result is true; otherwise if A is [unlimited polymorphic](#) and is either a [disassociated](#) pointer or  
27 unallocated [allocatable](#) variable, the result is false; otherwise if the [dynamic type](#) of A or MOLD is [extensible](#), the  
28 result is true if and only if the [dynamic type](#) of A is an [extension type](#) of the [dynamic type](#) of MOLD; otherwise  
29 the result is processor dependent.

#### NOTE 13.14

The [dynamic type](#) of a [disassociated](#) pointer or unallocated [allocatable](#) variable is its [declared type](#).

30 6 **Example.** Given the declarations and assignments

```
31      TYPE T1
32      REAL C
33      END TYPE
34      TYPE, EXTENDS(T1) :: T2
35      END TYPE
36      CLASS(T1), POINTER :: P, Q
```

1           ALLOCATE (P)

2           ALLOCATE (T2 :: Q)

3       the result of EXTENDS\_TYPE\_OF (P, Q) will be false, and the result of EXTENDS\_TYPE\_OF (Q, P) will  
4       be true.

### 5       **13.9.77    FAILED\_IMAGES([TEAM, KIND])**

6   1 **Description.** Indices of failed images.

7   2 **Class.** [Transformational function](#).

8   3 **Arguments.**

9       TEAM (optional) shall be a scalar of type [TEAM\\_TYPE](#) from the intrinsic module [ISO\\_FORTRAN\\_ENV](#). Its  
10       value shall be that of the [current](#) or an [ancestor](#) team. If TEAM is absent, the team specified is the  
11       [current team](#).

12       KIND (optional) shall be a scalar integer [constant expression](#). Its value shall specify a representation method  
13       whose decimal range is at least as large as that of default integer.

#### Unresolved Technical Issue 028

##### Why should FAILED\_IMAGES nanny the user?

The requirement for KIND not to specify a small integer kind is unnecessary except on big machines and supercomputers, and we don't require it for UCBOUND or LCOBOUND so it is pointless. We do not and cannot forbid assigning the result to a small integer variable anyway. We do not even require it for UBOUND and LBOUND, which are far more likely to give bad answers if assigned to a small integer.

This applies doubly to STOPPED\_IMAGES.

14 4 **Result Characteristics.** Integer. If KIND is present, the kind type parameter is that specified by the value  
15 of KIND; otherwise, the kind type parameter is that of default integer type. The result is an array of rank one  
16 whose size is equal to the number of images in the specified team that are known by the invoking image to have  
17 failed.

18 5 **Result Value.** The elements of the result are the values of the image indices of the known failed images in the  
19 specified team, in numerically increasing order. If the executing image has previously executed an [image control](#)  
20 [statement](#) whose STAT= specifier assigned the value [STAT\\_FAILED\\_IMAGE](#) from the intrinsic module [ISO\\_](#)  
21 [FORTRAN\\_ENV](#), or referenced a [collective subroutine](#) whose STAT argument was set to [STAT\\_FAILED\\_](#)  
22 [IMAGE](#), at least one image in the set of images participating in that [image control statement](#) or [collective](#)  
23 [subroutine](#) reference shall be known to have failed.

24 6 **Examples.** If image 3 is the only image in the [current team](#) that is known by the invoking image to have failed,  
25 FAILED\_IMAGES() will have the value [3]. If there are no images in the current team that are known by the  
26 invoking image to have failed, the value of FAILED\_IMAGES() will be a zero-sized array.

### 27       **13.9.78    FINDLOC (ARRAY, VALUE, DIM [, MASK, KIND, BACK]) or**          **FINDLOC (ARRAY, VALUE [, MASK, KIND, BACK])**

28 1 **Description.** Location(s) of a specified value.

29 2 **Class.** [Transformational function](#).

30 3 **Arguments.**

31       ARRAY       shall be an array of intrinsic type.

32       VALUE       shall be scalar and in type conformance with ARRAY, as specified in Table 7.2 for the operator ==  
33       or the operator [.EQV.](#).

1 DIM shall be an integer scalar with a value in the range  $1 \leq \text{DIM} \leq n$ , where  $n$  is the **rank** of ARRAY.  
 2 MASK (optional) shall be of type logical and shall be **conformable** with ARRAY.  
 3 KIND (optional) shall be a scalar integer **constant expression**.  
 4 BACK (optional) shall be a logical scalar.

5 4 **Result Characteristics.** Integer. If KIND is present, the kind type parameter is that specified by the value of  
 6 KIND; otherwise the kind type parameter is that of default integer type. If DIM does not appear, the result is  
 7 an array of **rank** one and of size equal to the **rank** of ARRAY; otherwise, the result is of **rank**  $n - 1$  and shape  
 8  $[d_1, d_2, \dots, d_{\text{DIM}-1}, d_{\text{DIM}+1}, \dots, d_n]$ , where  $[d_1, d_2, \dots, d_n]$  is the shape of ARRAY.

9 5 **Result Value.**

10 *Case (i):* The result of FINDLOC (ARRAY, VALUE) is a rank-one array whose element values are the values  
 11 of the subscripts of an element of ARRAY whose value matches VALUE. If there is such a value,  
 12 the  $i^{\text{th}}$  element value is in the range 1 to  $e_i$ , where  $e_i$  is the extent of the  $i^{\text{th}}$  dimension of ARRAY.  
 13 If no elements match VALUE or ARRAY has size zero, all elements of the result are zero.

14 *Case (ii):* The result of FINDLOC (ARRAY, VALUE, MASK = MASK) is a rank-one array whose element  
 15 values are the values of the subscripts of an element of ARRAY, corresponding to a true element  
 16 of MASK, whose value matches VALUE. If there is such a value, the  $i^{\text{th}}$  element value is in the  
 17 range 1 to  $e_i$ , where  $e_i$  is the extent of the  $i^{\text{th}}$  dimension of ARRAY. If no elements match VALUE,  
 18 ARRAY has size zero, or every element of MASK has the value false, all elements of the result are  
 19 zero.

20 *Case (iii):* If ARRAY has **rank** one, the result of  
 21 FINDLOC (ARRAY, VALUE, DIM=DIM [, MASK = MASK]) is a scalar whose value is equal to  
 22 that of the first element of FINDLOC (ARRAY, VALUE [, MASK = MASK]). Otherwise, the value  
 23 of element  $(s_1, s_2, \dots, s_{\text{DIM}-1}, s_{\text{DIM}+1}, \dots, s_n)$  of the result is equal to FINDLOC (ARRAY  $(s_1,$   
 24  $s_2, \dots, s_{\text{DIM}-1}, :, s_{\text{DIM}+1}, \dots, s_n)$ , VALUE, DIM=1 [, MASK = MASK  $(s_1, s_2, \dots, s_{\text{DIM}-1}, :,$   
 25  $s_{\text{DIM}+1}, \dots, s_n)$ )).

26 6 If both ARRAY and VALUE are of type logical, the comparison is performed with the **.EQV.** operator; otherwise,  
 27 the comparison is performed with the **==** operator. If the value of the comparison is true, that element of ARRAY  
 28 matches VALUE.

29 7 If DIM is not present, more than one element matches VALUE, and BACK is absent or present with the value  
 30 false, the value returned indicates the first such element, taken in array element order. If DIM is not present and  
 31 BACK is present with the value true, the value returned indicates the last such element, taken in array element  
 32 order.

33 8 **Examples.**

34 *Case (i):* The value of FINDLOC ([2, 6, 4, 6], VALUE=6) is [2], and the value of FINDLOC ([2, 6, 4, 6],  
 35 VALUE=6, BACK = **.TRUE.**) is [4].

36 *Case (ii):* If A has the value  $\begin{bmatrix} 0 & -5 & 7 & 7 \\ 3 & 4 & -1 & 2 \\ 1 & 5 & 6 & 7 \end{bmatrix}$ , and M has the value  $\begin{bmatrix} T & T & F & T \\ T & T & F & T \\ T & T & F & T \end{bmatrix}$ , FINDLOC (A, 7,  
 37 MASK = M) has the value [1, 4] and FINDLOC (A, 7, MASK = M, BACK = **.TRUE.**) has the  
 38 value [3, 4]. This is independent of the declared lower bounds for A.

39 *Case (iii):* The value of FINDLOC ([2, 6, 4], VALUE=6, DIM=1) is 2. If B has the value  
 40  $\begin{bmatrix} 1 & 2 & -9 \\ 2 & 2 & 6 \end{bmatrix}$ , FINDLOC (B, VALUE=2, DIM=1) has the value [2, 1, 0] and FINDLOC (B,  
 41 VALUE=2, DIM=2) has the value [2, 1]. This is independent of the declared lower bounds for B.

## 42 13.9.79 FLOOR (A [, KIND])

43 1 **Description.** Greatest integer less than or equal to A.

44 2 **Class.** **Elemental** function.

1 3 **Arguments.**

2 A shall be of type real.

3 KIND (optional) shall be a scalar integer [constant expression](#).

4 4 **Result Characteristics.** Integer. If KIND is present, the kind type parameter is that specified by the value of  
5 KIND; otherwise, the kind type parameter is that of default integer type.

6 5 **Result Value.** The result has a value equal to the greatest integer less than or equal to A.

7 6 **Examples.** FLOOR (3.7) has the value 3. FLOOR (−3.7) has the value −4.

8 **13.9.80 FRACTION (X)**

9 1 **Description.** Fractional part of number.

10 2 **Class.** [Elemental](#) function.

11 3 **Argument.** X shall be of type real.

12 4 **Result Characteristics.** Same as X.

13 5 **Result Value.** The result has the value  $X \times b^{-e}$ , where  $b$  and  $e$  are as defined in [13.4](#) for the representation of  
14 X in the extended real model for the kind of X. If X has the value zero, the result is zero. If X is an [IEEE NaN](#),  
15 the result is that [NaN](#). If X is an IEEE infinity, the result is an [IEEE NaN](#).

16 6 **Example.** FRACTION (3.0) has the value 0.75 for reals whose model is as in Note [13.4](#).

17 **13.9.81 GAMMA (X)**

18 1 **Description.** Gamma function.

19 2 **Class.** [Elemental](#) function.

20 3 **Argument.** X shall be of type real. Its value shall not be a negative integer or zero.

21 4 **Result Characteristics.** Same as X.

5 **Result Value.** The result has a value equal to a processor-dependent approximation to the gamma function of X,

$$\Gamma(X) = \begin{cases} \int_0^\infty t^{X-1} \exp(-t) dt & X > 0 \\ \int_0^\infty t^{X-1} \left( \exp(-t) - \sum_{k=0}^n \frac{(-t)^k}{k!} \right) dt & -n-1 < X < -n, n \text{ an integer } \geq 0 \end{cases}$$

22 6 **Example.** GAMMA (1.0) has the value 1.000 (approximately).

23 **13.9.82 GET\_COMMAND ([COMMAND, LENGTH, STATUS, ERRMSG])**

24 1 **Description.** Get program invocation command.

25 2 **Class.** Subroutine.

26 3 **Arguments.**

27 COMMAND (optional) shall be a default character scalar. It is an [INTENT \(OUT\)](#) argument. It is assigned  
28 the entire command by which the program was invoked. If the command cannot be determined,  
29 COMMAND is assigned all blanks.

30 LENGTH (optional) shall be a scalar of type integer with a decimal exponent range of at least four. It is an  
31 [INTENT \(OUT\)](#) argument. It is assigned the significant length of the command by which the

program was invoked. The significant length may include trailing blanks if the processor allows commands with significant trailing blanks. This length does not consider any possible truncation or padding in assigning the command to the COMMAND argument; in fact the COMMAND argument need not even be present. If the command length cannot be determined, a length of 0 is assigned.

STATUS (optional) shall be a scalar of type integer with a decimal exponent range of at least four. It is an **INTENT (OUT)** argument. It is assigned the value  $-1$  if the COMMAND argument is present and has a length less than the significant length of the command. It is assigned a processor-dependent positive value if the command retrieval fails. Otherwise it is assigned the value 0.

ERRMSG (optional) shall be a default character scalar. It is an **INTENT (INOUT)** argument. It is assigned a processor-dependent explanatory message if the command retrieval fails. Otherwise, it is unchanged.

- 4 **Example.** If the program below is invoked with the command “example” on a processor that supports command retrieval, it will display “Hello example”.

```

PROGRAM hello
  CHARACTER(:), ALLOCATABLE :: cmd
  INTEGER :: cmdlen
  CALL GET_COMMAND(LENGTH=cmdlen)
  IF (cmdlen>0) THEN
    ALLOCATE(CHARACTER(cmdlen) :: cmd)
    CALL GET_COMMAND(cmd)
    PRINT *, 'Hello ', cmd
  END IF
END PROGRAM

```

### 13.9.83 GET\_COMMAND\_ARGUMENT (NUMBER [, VALUE, LENGTH, STATUS, ERRMSG])

- 1 **Description.** Get program invocation argument.

- 2 **Class.** Subroutine.

- 3 **Arguments.**

NUMBER shall be an integer scalar. It is an **INTENT (IN)** argument that specifies the number of the command argument that the other arguments give information about.

Command argument 0 always exists, and is the command name by which the program was invoked if the processor has such a concept; otherwise, the value of command argument 0 is processor dependent. The remaining command arguments are numbered consecutively from 1 to the argument count in an order determined by the processor.

VALUE (optional) shall be a default character scalar. It is an **INTENT (OUT)** argument. If the command argument specified by NUMBER exists, its value is assigned to VALUE; otherwise, VALUE is assigned all blanks.

LENGTH (optional) shall be a scalar of type integer with a decimal exponent range of at least four. It is an **INTENT (OUT)** argument. If the command argument specified by NUMBER exists, its significant length is assigned to LENGTH; otherwise, LENGTH is assigned the value zero. It is processor dependent whether the significant length includes trailing blanks. This length does not consider any possible truncation or padding in assigning the command argument value to the VALUE argument; in fact the VALUE argument need not even be present.

STATUS (optional) shall be a scalar of type integer with a decimal exponent range of at least four. It is an **INTENT (OUT)** argument. If NUMBER is less than zero or greater than the argument count that would be returned by the intrinsic function **COMMAND\_ARGUMENT\_COUNT**, or command retrieval fails, STATUS is assigned a processor-dependent positive value. Otherwise, if VALUE is

present and has a length less than the significant length of the specified command argument, it is assigned the value  $-1$ . Otherwise it is assigned the value 0.

ERRMSG (optional) shall be a default character scalar. It is an **INTENT (INOUT)** argument. It is assigned a processor-dependent explanatory message if the optional argument STATUS is, or would be if present, assigned a positive value. Otherwise, it is unchanged.

4 **Example.** On a processor that supports command arguments, the following program displays the arguments of the command by which it was invoked.

```

PROGRAM show_arguments
  INTEGER :: i
  CHARACTER :: command*32, arg*128
  CALL get_command_argument(0, command)
  WRITE (*,*) "Command name is: ", command
  DO i = 1, command_argument_count()
    CALL get_command_argument(i, arg)
    WRITE (*,*) "Argument ", i, " is ", arg
  END DO
END PROGRAM show_arguments

```

### 13.9.84 GET\_ENVIRONMENT\_VARIABLE (NAME [, VALUE, LENGTH, STATUS, TRIM\_NAME, ERRMSG])

1 **Description.** Get environment variable.

2 **Class.** Subroutine.

3 **Arguments.**

NAME shall be a default character scalar. It is an **INTENT (IN)** argument. The interpretation of case is processor dependent.

VALUE (optional) shall be a default character scalar. It is an **INTENT (OUT)** argument. It is assigned the value of the environment variable specified by NAME. VALUE is assigned all blanks if the environment variable does not exist or does not have a value or if the processor does not support environment variables.

LENGTH (optional) shall be a scalar of type integer with a decimal exponent range of at least four. It is an **INTENT (OUT)** argument. If the specified environment variable exists and has a value, LENGTH is assigned the value of its length. Otherwise LENGTH is assigned the value zero.

STATUS (optional) shall be a scalar of type integer with a decimal exponent range of at least four. It is an **INTENT (OUT)** argument. If the environment variable exists and either has no value, its value is successfully assigned to VALUE, or the VALUE argument is not present, STATUS is assigned the value zero. STATUS is assigned the value  $-1$  if the VALUE argument is present and has a length less than the significant length of the environment variable. It is assigned the value 1 if the specified environment variable does not exist, or 2 if the processor does not support environment variables. Processor-dependent values greater than 2 may be assigned for other error conditions.

TRIM\_NAME (optional) shall be a logical scalar. It is an **INTENT (IN)** argument. If TRIM\_NAME is present with the value false then trailing blanks in NAME are considered significant if the processor supports trailing blanks in environment variable names. Otherwise trailing blanks in NAME are not considered part of the environment variable's name.

ERRMSG (optional) shall be a default character scalar. It is an **INTENT (INOUT)** argument. It is assigned a processor-dependent explanatory message if the optional argument STATUS is, or would be if present, assigned a positive value. Otherwise, it is unchanged.

4 It is processor dependent whether an environment variable that exists on an [image](#) also exists on another [image](#), and if it does exist on both [images](#), whether the values are the same or different.

5 **Example.** If the value of the environment variable DATAFILE is datafile.dat, executing the statement sequence below will assign the value 'datafile.dat' to FILENAME.

```

5      CHARACTER(:),ALLOCATABLE :: FILENAME
6      INTEGER :: NAMELEN
7      CALL GET_ENVIRONMENT_VARIABLE ("DATAFILE", LENGTH=NAMELEN)
8      IF (LENGTH>0) THEN
9          ALLOCATE(CHARACTER(LENGTH) :: FILENAME)
10         CALL GET_ENVIRONMENT_VARIABLE("DATAFILE", FILENAME)
11     END IF

```

## 12 13.9.85 GET\_TEAM([LEVEL])

13 1 **Description.** Team.

14 2 **Class.** Transformational function.

15 3 **Argument.** LEVEL (optional) shall be a scalar integer whose value shall be equal to one of the named constants [INITIAL\\_TEAM](#), [PARENT\\_TEAM](#), or [CURRENT\\_TEAM](#) from the intrinsic module [ISO\\_FORTRAN\\_ENV](#).

17 4 **Result Characteristics.** Scalar of type [TEAM\\_TYPE](#) from the intrinsic module [ISO\\_FORTRAN\\_ENV](#).

18 5 **Result Value.** The result is a [TEAM\\_TYPE](#) value specifying the [current team](#) if LEVEL is not present, present with the value [CURRENT\\_TEAM](#), or the [current team](#) is the [initial team](#). Otherwise, the result specifies the [parent team](#) if LEVEL is present with the value [PARENT\\_TEAM](#), and specifies the [initial team](#) if LEVEL is present with the value [INITIAL\\_TEAM](#).

22 6 **Examples.**

```

23      PROGRAM EXAMPLE1
24          USE,INTRINSIC :: ISO_FORTRAN_ENV, ONLY: TEAM_TYPE
25          TYPE(TEAM_TYPE) :: WORLD_TEAM, TEAM2
26
27          ! Define a team variable representing the initial team
28          WORLD_TEAM = GET_TEAM()
29      END PROGRAM
30
31      SUBROUTINE EXAMPLE2 (A)
32          USE,INTRINSIC :: ISO_FORTRAN_ENV, ONLY: TEAM_TYPE
33          REAL A[*]
34          TYPE(TEAM_TYPE) :: NEW_TEAM, PARENT_TEAM
35
36          ... ! Form NEW_TEAM
37
38          PARENT_TEAM = GET_TEAM ()
39
40          CHANGE TEAM (NEW_TEAM)
41
42          ! Reference image 1 in parent's team

```

```

1           A [1,TEAM=PARENT_TEAM] = 4.2
2
3           ! Reference image 1 in current team
4           A [1] = 9.0
5       END TEAM
6   END SUBROUTINE TT

```

### 13.9.86 HUGE (X)

1 **Description.** Largest model number.

2 **Class.** [Inquiry function](#).

3 **Argument.** X shall be of type integer or real. It may be a scalar or an array.

4 **Result Characteristics.** Scalar of the same type and kind type parameter as X.

5 **Result Value.** The result has the value  $r^q - 1$  if X is of type integer and  $(1 - b^{-p})b^{e_{\max}}$  if X is of type real, where  $r$ ,  $q$ ,  $b$ ,  $p$ , and  $e_{\max}$  are as defined in [13.4](#) for the model representing numbers of the same type and kind type parameter as X.

6 **Example.** HUGE (X) has the value  $(1 - 2^{-24}) \times 2^{127}$  for real X whose model is as in Note [13.4](#).

### 13.9.87 HYPOT (X, Y)

1 **Description.** Euclidean distance function.

2 **Class.** [Elemental function](#).

3 **Arguments.**

X shall be of type real.

Y shall be of type real with the same kind type parameter as X.

4 **Result Characteristics.** Same as X.

5 **Result Value.** The result has a value equal to a processor-dependent approximation to the Euclidean distance,  $\sqrt{X^2 + Y^2}$ , without undue overflow or underflow.

6 **Example.** HYPOT (3.0, 4.0) has the value 5.0 (approximately).

### 13.9.88 IACHAR (C [, KIND])

1 **Description.** ASCII code value for character.

2 **Class.** [Elemental function](#).

3 **Arguments.**

C shall be of type character and of length one.

KIND (optional) shall be a scalar integer [constant expression](#).

4 **Result Characteristics.** Integer. If KIND is present, the kind type parameter is that specified by the value of KIND; otherwise, the kind type parameter is that of default integer type.

5 **Result Value.** If C is in the [collating sequence](#) defined by the codes specified in ISO/IEC 646:1991 (International Reference Version), the result is the position of C in that sequence; it is nonnegative and less than or equal to 127. The value of the result is processor dependent if C is not in the [ASCII collating sequence](#). The results

are consistent with the LGE, LGT, LLE, and LLT comparison functions. For example, if LLE (C, D) is true, IACHAR (C) <= IACHAR (D) is true where C and D are any two characters representable by the processor.

**Example.** IACHAR ('X') has the value 88.

### 13.9.89 IALL (ARRAY, DIM [, MASK]) or IALL (ARRAY [, MASK])

**Description.** Array reduced by IAND function.

**Class.** Transformational function.

**Arguments.**

ARRAY shall be an array of type integer.

DIM shall be an integer scalar with a value in the range  $1 \leq \text{DIM} \leq n$ , where  $n$  is the **rank** of ARRAY.

MASK (optional) shall be of type logical and shall be **conformable** with ARRAY.

**Result Characteristics.** The result is of the same type and kind type parameter as ARRAY. It is scalar if DIM does not appear or if ARRAY has **rank** one; otherwise, the result is an array of **rank**  $n - 1$  and shape  $[d_1, d_2, \dots, d_{\text{DIM}-1}, d_{\text{DIM}+1}, \dots, d_n]$  where  $[d_1, d_2, \dots, d_n]$  is the shape of ARRAY.

**Result Value.**

*Case (i):* If ARRAY has size zero the result value is equal to NOT (INT (0, KIND (ARRAY))). Otherwise, the result of IALL (ARRAY) has a value equal to the bitwise AND of all the elements of ARRAY.

*Case (ii):* The result of IALL (ARRAY, MASK=MASK) has a value equal to IALL (PACK (ARRAY, MASK)).

*Case (iii):* The result of IALL (ARRAY, DIM=DIM [, MASK=MASK]) has a value equal to that of IALL (ARRAY [, MASK=MASK]) if ARRAY has **rank** one. Otherwise, the value of element  $(s_1, s_2, \dots, s_{\text{DIM}-1}, s_{\text{DIM}+1}, \dots, s_n)$  of the result is equal to IALL (ARRAY  $(s_1, s_2, \dots, s_{\text{DIM}-1}, :, s_{\text{DIM}+1}, \dots, s_n)$  [, MASK = MASK  $(s_1, s_2, \dots, s_{\text{DIM}-1}, :, s_{\text{DIM}+1}, \dots, s_n)$ ]).

**Examples.** IALL ([14, 13, 11]) has the value 8. IALL ([14, 13, 11], MASK=[.true., .false., .true]) has the value 10.

### 13.9.90 IAND (I, J)

**Description.** Bitwise AND.

**Class.** Elemental function.

**Arguments.**

I shall be of type integer or a **boz-literal-constant**.

J shall be of type integer or a **boz-literal-constant**. If both I and J are of type integer, they shall have the same kind type parameter. I and J shall not both be **boz-literal-constants**.

**Result Characteristics.** Same as I if I is of type integer; otherwise, same as J.

**Result Value.** If either I or J is a **boz-literal-constant**, it is first converted as if by the intrinsic function INT to type integer with the kind type parameter of the other. The result has the value obtained by combining I and J bit-by-bit according to the following table:

I	J	IAND (I, J)
1	1	1
1	0	0
0	1	0
0	0	0

6 The model for the interpretation of an integer value as a sequence of bits is in 13.3.

7 **Example.** IAND (1, 3) has the value 1.

### 13.9.91 IANY (ARRAY, DIM [, MASK]) or IANY (ARRAY [, MASK])

1 **Description.** Reduce array with bitwise OR operation.

2 **Class.** Transformational function.

3 **Arguments.**

4 ARRAY shall be of type integer. It shall be an array.

5 DIM shall be an integer scalar with a value in the range  $1 \leq DIM \leq n$ , where  $n$  is the rank of ARRAY.

6 MASK (optional) shall be of type logical and shall be conformable with ARRAY.

7 **Result Characteristics.** The result is of the same type and kind type parameter as ARRAY. It is scalar if DIM does not appear or if ARRAY has rank one; otherwise, the result is an array of rank  $n - 1$  and shape  $[d_1, d_2, \dots, d_{DIM-1}, d_{DIM+1}, \dots, d_n]$  where  $[d_1, d_2, \dots, d_n]$  is the shape of ARRAY.

8 **Result Value.**

9 *Case (i):* The result of IANY (ARRAY) is the bitwise OR of all the elements of ARRAY. If ARRAY has size zero the result value is equal to zero.

10 *Case (ii):* The result of IANY (ARRAY, MASK=MASK) has a value equal to IANY (PACK (ARRAY, MASK)).

11 *Case (iii):* The result of IANY (ARRAY, DIM=DIM [, MASK=MASK]) has a value equal to that of IANY (ARRAY [, MASK=MASK]) if ARRAY has rank one. Otherwise, the value of element  $(s_1, s_2, \dots, s_{DIM-1}, s_{DIM+1}, \dots, s_n)$  of the result is equal to IANY (ARRAY  $(s_1, s_2, \dots, s_{DIM-1}, :, s_{DIM+1}, \dots, s_n)$  [, MASK = MASK  $(s_1, s_2, \dots, s_{DIM-1}, :, s_{DIM+1}, \dots, s_n)$ ]).

12 **Examples.** IANY ([14, 13, 8]) has the value 15. IANY ([14, 13, 8], MASK=[.true., .false., .true]) has the value 14.

### 13.9.92 IBCLR (I, POS)

1 **Description.** I with bit POS replaced by zero.

2 **Class.** Elemental function.

3 **Arguments.**

4 I shall be of type integer.

5 POS shall be of type integer. It shall be nonnegative and less than BIT\_SIZE (I).

6 **Result Characteristics.** Same as I.

7 **Result Value.** The result has the value of the sequence of bits of I, except that bit POS is zero. The model for the interpretation of an integer value as a sequence of bits is in 13.3.

8 **Examples.** IBCLR (14, 1) has the value 12. If V has the value [1, 2, 3, 4], the value of IBCLR (POS = V, I = 31) is [29, 27, 23, 15].

### 13.9.93 IBITS (I, POS, LEN)

1 **Description.** Specified sequence of bits.

2 **Class.** Elemental function.

3 **Arguments.**

- 1 I shall be of type integer.
- 2 POS shall be of type integer. It shall be nonnegative and POS + LEN shall be less than or equal to
- 3 BIT\_SIZE (I).
- 4 LEN shall be of type integer and nonnegative.

5 4 **Result Characteristics.** Same as I.

6 5 **Result Value.** The result has the value of the sequence of LEN bits in I beginning at bit POS, right-adjusted

7 and with all other bits zero. The model for the interpretation of an integer value as a sequence of bits is in [13.3](#).

8 6 **Example.** IBITS (14, 1, 3) has the value 7.

### 9 13.9.94 IBSET (I, POS)

10 1 **Description.** I with bit POS replaced by one.

11 2 **Class.** [Elemental](#) function.

12 3 **Arguments.**

13 I shall be of type integer.

14 POS shall be of type integer. It shall be nonnegative and less than BIT\_SIZE (I).

15 4 **Result Characteristics.** Same as I.

16 5 **Result Value.** The result has the value of the sequence of bits of I, except that bit POS is one. The model for

17 the interpretation of an integer value as a sequence of bits is in [13.3](#).

18 6 **Examples.** IBSET (12, 1) has the value 14. If V has the value [1, 2, 3, 4], the value of IBSET (POS = V, I = 0)

19 is [2, 4, 8, 16].

### 20 13.9.95 ICHAR (C [, KIND])

21 1 **Description.** Code value for character.

22 2 **Class.** [Elemental](#) function.

23 3 **Arguments.**

24 C shall be of type character and of length one. Its value shall be that of a character capable of

25 representation in the processor.

26 KIND (optional) shall be a scalar integer [constant expression](#).

27 4 **Result Characteristics.** Integer. If KIND is present, the kind type parameter is that specified by the value of

28 KIND; otherwise, the kind type parameter is that of default integer type.

29 5 **Result Value.** The result is the position of C in the processor [collating sequence](#) associated with the kind type

30 parameter of C; it is nonnegative and less than  $n$ , where  $n$  is the number of characters in the [collating sequence](#).

31 The [kind type parameter](#) of the result shall specify an integer kind that is capable of representing  $n$ . For any char-

32 acters C and D capable of representation in the processor,  $C \leq D$  is true if and only if ICHAR (C)  $\leq$  ICHAR (D)

33 is true and  $C == D$  is true if and only if ICHAR (C)  $==$  ICHAR (D) is true.

34 6 **Example.** ICHAR ('X') has the value 88 on a processor using the [ASCII collating sequence](#) for default characters.

### 35 13.9.96 IEOR (I, J)

36 1 **Description.** Bitwise exclusive OR.

37 2 **Class.** [Elemental](#) function.

1 3 **Arguments.**

2 I shall be of type integer or a *boz-literal-constant*.  
 3 J shall be of type integer or a *boz-literal-constant*. If both I and J are of type integer, they shall have  
 4 the same kind type parameter. I and J shall not both be *boz-literal-constants*.

5 4 **Result Characteristics.** Same as I if I is of type integer; otherwise, same as J.

6 5 **Result Value.** If either I or J is a *boz-literal-constant*, it is first converted as if by the intrinsic function **INT** to  
 7 type integer with the kind type parameter of the other. The result has the value obtained by combining I and J  
 8 bit-by-bit according to the following table:

I	J	IEOR (I, J)
1	1	0
1	0	1
0	1	1
0	0	0

9 6 The model for the interpretation of an integer value as a sequence of bits is in [13.3](#).

10 7 **Example.** IEOR (1, 3) has the value 2.

### 11 13.9.97 IMAGE\_INDEX (COARRAY, SUB) or (COARRAY, SUB, TEAM) or (COARRAY, SUB, TEAM\_NUMBER)

12 1 **Description.** *Image index* from *cosubscripts*.

13 2 **Class.** *Inquiry function*.

#### 14 3 **Arguments.**

15 COARRAY shall be a *coarray* of any type. If TEAM\_NUMBER appears and the *current team* is not the *initial*  
 16 *team*, it shall be established in an *ancestor* of the *current team*. Otherwise, if TEAM appears, it  
 17 shall be established in the that *team*. Otherwise, it shall be established in the *current team*.

18 SUB shall be a rank-one integer array of size equal to the *corank* of COARRAY.

19 TEAM shall be a scalar of type **TEAM\_TYPE** from the intrinsic module **ISO\_FORTRAN\_ENV**, with a  
 20 value that identifies the *current* or an *ancestor team*.

21 TEAM\_NUMBER shall be a positive integer scalar. If the *current team* is not the *initial team*, it shall identify  
 22 a *team* whose *parent* is the same as that of the *current team*.

23 4 **Result Characteristics.** Default integer scalar.

24 5 **Result Value.** If the value of SUB is a valid sequence of *cosubscripts* for COARRAY in the *team* specified by  
 25 TEAM or the *current team* if TEAM does not appear, the result is the *index* of the corresponding *image* in that  
 26 *team*. Otherwise, the result is zero.

27 6 **Examples.** If A and B are declared as A [0:\*) and B (10, 20) [10, 0:9, 0:\*) respectively, IMAGE\_INDEX (A, [0])  
 28 has the value 1 and IMAGE\_INDEX (B, [3, 1, 2]) has the value 213 (on any *image*).

### 29 13.9.98 IMAGE\_STATUS(IMAGE, [TEAM])

30 1 **Description.** Image execution state.

31 2 **Class.** *Elemental function*.

#### 32 3 **Arguments.**

1 IMAGE shall be of type integer. Its value shall be positive and less than or equal to the number of `images`  
 2 in the specified `team`.  
 3 TEAM (optional) shall be a scalar of type `TEAM_TYPE` from the intrinsic module `ISO_FORTRAN_ENV`. Its  
 4 value shall represent the `current` or an `ancestor` team. If TEAM is absent, the team specified is the  
 5 `current team`.

6 4 **Result Characteristics.** Default integer.

7 5 **Result Value.** The result value is `STAT_FAILED_IMAGE` from the intrinsic module `ISO_FORTRAN_ENV`  
 8 if the specified image has failed, `STAT_STOPPED_IMAGE` from the intrinsic module `ISO_FORTRAN_ENV`  
 9 if that image has initiated `normal termination`, and zero otherwise.

#### Unresolved Technical Issue 030

There is, and should be, no concept of an “error image”.

I deleted “, a positive processor-dependent value different from `STAT_FAILED_IMAGE` or `STAT_STOPPED_IMAGE` if some other error has occurred for that image”, because an image is either failed, stopped, or active. There is, and should not be, any other image execution state.

If it felt desirable for an image execution state to be “some other error”, lots of extra edits will be needed to establish that concept and make it consistent with the rest of the fault tolerance design.

10 6 **Example.** If image 3 of the current team has failed, `IMAGE_STATUS` (3) has the value `STAT_FAILED_`  
 11 `IMAGE`.

## 12 13.9.99 INDEX (STRING, SUBSTRING [, BACK, KIND])

13 1 **Description.** Character string search.

14 2 **Class.** Elemental function.

15 3 **Arguments.**

16 STRING shall be of type character.

17 SUBSTRING shall be of type character with the same `kind type parameter` as STRING.

18 BACK (optional) shall be of type logical.

19 KIND (optional) shall be a scalar integer `constant expression`.

20 4 **Result Characteristics.** Integer. If KIND is present, the `kind type parameter` is that specified by the value of  
 21 KIND; otherwise the `kind type parameter` is that of default integer type.

22 5 **Result Value.**

23 *Case (i):* If `STRING % LEN < SUBSTRING % LEN`, the result has the value zero.

24 *Case (ii):* Otherwise, if there is an integer I in the range  $1 \leq I \leq \text{STRING \% LEN} - \text{SUBSTRING \% LEN} + 1$ , such that `STRING(I : I + SUBSTRING % LEN - 1)` is equal to SUBSTRING, the result has  
 25 the value of the smallest such I if BACK is absent or present with the value false, and the greatest  
 26 such I if BACK is present with the value true.

27 *Case (iii):* Otherwise, the result has the value zero.

29 6 **Examples.** `INDEX ('FORTRAN', 'R')` has the value 3.  
 30 `INDEX ('FORTRAN', 'R', BACK = .TRUE.)` has the value 5.

### 13.9.100 INT (A [, KIND])

**Description.** Conversion to integer type.

**Class.** [Elemental](#) function.

**Arguments.**

A shall be of type integer, real, or complex, or a *boz-literal-constant*.

KIND (optional) shall be a scalar integer [constant expression](#).

**Result Characteristics.** Integer. If KIND is present, the kind type parameter is that specified by the value of KIND; otherwise, the kind type parameter is that of default integer type.

**Result Value.**

*Case (i):* If A is of type integer,  $\text{INT}(A) = A$ .

*Case (ii):* If A is of type real, there are two cases: if  $|A| < 1$ ,  $\text{INT}(A)$  has the value 0; if  $|A| \geq 1$ ,  $\text{INT}(A)$  is the integer whose magnitude is the largest integer that does not exceed the magnitude of A and whose sign is the same as the sign of A.

*Case (iii):* If A is of type complex,  $\text{INT}(A) = \text{INT}(\text{REAL}(A, \text{KIND}(A)))$ .

*Case (iv):* If A is a *boz-literal-constant*, the value of the result is the value whose bit sequence according to the model in [13.3](#) is the same as that of A as modified by padding or truncation according to [13.3.3](#). The interpretation of a bit sequence whose most significant bit is 1 is processor dependent.

**Example.**  $\text{INT}(-3.7)$  has the value  $-3$ .

### 13.9.101 IOR (I, J)

**Description.** Bitwise inclusive OR.

**Class.** [Elemental](#) function.

**Arguments.**

I shall be of type integer or a *boz-literal-constant*.

J shall be of type integer or a *boz-literal-constant*. If both I and J are of type integer, they shall have the same kind type parameter. I and J shall not both be *boz-literal-constants*.

**Result Characteristics.** Same as I if I is of type integer; otherwise, same as J.

**Result Characteristics.** Same as I.

**Result Value.** If either I or J is a *boz-literal-constant*, it is first converted as if by the intrinsic function [INT](#) to type integer with the kind type parameter of the other. The result has the value obtained by combining I and J bit-by-bit according to the following table:

I	J	IOR (I, J)
1	1	1
1	0	1
0	1	1
0	0	0

The model for the interpretation of an integer value as a sequence of bits is in [13.3](#).

**Example.**  $\text{IOR}(5, 3)$  has the value 7.

### 13.9.102 IPARITY (ARRAY, DIM [, MASK]) or IPARITY (ARRAY [, MASK])

**1 Description.** Array reduced by IEOR function.

**2 Class.** Transformational function.

**3 Arguments.**

ARRAY shall be of type integer. It shall be an array.

DIM shall be an integer scalar with a value in the range  $1 \leq \text{DIM} \leq n$ , where  $n$  is the [rank](#) of ARRAY.

MASK (optional) shall be of type logical and shall be [conformable](#) with ARRAY.

**4 Result Characteristics.** The result is of the same type and kind type parameter as ARRAY. It is scalar if DIM does not appear; otherwise, the result has [rank](#)  $n - 1$  and shape  $[d_1, d_2, \dots, d_{\text{DIM}-1}, d_{\text{DIM}+1}, \dots, d_n]$  where  $[d_1, d_2, \dots, d_n]$  is the shape of ARRAY.

**5 Result Value.**

*Case (i):* The result of IPARITY (ARRAY) has a value equal to the bitwise exclusive OR of all the elements of ARRAY. If ARRAY has size zero the result has the value zero.

*Case (ii):* The result of IPARITY (ARRAY, MASK=MASK) has a value equal to that of IPARITY (PACK (ARRAY, MASK)).

*Case (iii):* The result of IPARITY (ARRAY, DIM=DIM [, MASK=MASK]) has a value equal to that of IPARITY (ARRAY [, MASK=MASK]) if ARRAY has [rank](#) one. Otherwise, the value of element  $(s_1, s_2, \dots, s_{\text{DIM}-1}, s_{\text{DIM}+1}, \dots, s_n)$  of the result is equal to IPARITY (ARRAY  $(s_1, s_2, \dots, s_{\text{DIM}-1}, :, s_{\text{DIM}+1}, \dots, s_n)$  [, MASK = MASK  $(s_1, s_2, \dots, s_{\text{DIM}-1}, :, s_{\text{DIM}+1}, \dots, s_n)$ ])).

**6 Examples.** IPARITY ([14, 13, 8]) has the value 11. IPARITY ([14, 13, 8], MASK=[.true., .false., .true]) has the value 6.

### 13.9.103 ISHFT (I, SHIFT)

**1 Description.** Logical shift.

**2 Class.** Elemental function.

**3 Arguments.**

I shall be of type integer.

SHIFT shall be of type integer. The absolute value of SHIFT shall be less than or equal to BIT\_SIZE (I).

**4 Result Characteristics.** Same as I.

**5 Result Value.** The result has the value obtained by shifting the bits of I by SHIFT positions. If SHIFT is positive, the shift is to the left; if SHIFT is negative, the shift is to the right; if SHIFT is zero, no shift is performed. Bits shifted out from the left or from the right, as appropriate, are lost. Zeros are shifted in from the opposite end. The model for the interpretation of an integer value as a sequence of bits is in [13.3](#).

**6 Example.** ISHFT (3, 1) has the value 6.

### 13.9.104 ISHFTC (I, SHIFT [, SIZE])

**1 Description.** Circular shift of the rightmost bits.

**2 Class.** Elemental function.

**3 Arguments.**

I shall be of type integer.

SHIFT shall be of type integer. The absolute value of SHIFT shall be less than or equal to SIZE.

1 SIZE (optional) shall be of type integer. The value of SIZE shall be positive and shall not exceed BIT\_SIZE (I).  
 2 If SIZE is absent, it is as if it were present with the value of BIT\_SIZE (I).

3 4 **Result Characteristics.** Same as I.

4 5 **Result Value.** The result has the value obtained by shifting the SIZE rightmost bits of I circularly by SHIFT  
 5 positions. If SHIFT is positive, the shift is to the left; if SHIFT is negative, the shift is to the right; and if SHIFT  
 6 is zero, no shift is performed. No bits are lost. The unshifted bits are unaltered. The model for the interpretation  
 7 of an integer value as a sequence of bits is in 13.3.

8 6 **Example.** ISHFTC (3, 2, 3) has the value 5.

## 9 13.9.105 IS\_CONTIGUOUS (ARRAY)

10 1 **Description.** Array contiguity test (5.5.7).

11 2 **Class.** Inquiry function.

12 3 **Argument.** ARRAY may be of any type. It shall be assumed-rank or an array. If it is a pointer it shall be  
 13 associated.

14 4 **Result Characteristics.** Default logical scalar.

15 5 **Result Value.** The result has the value true if ARRAY has rank zero or is contiguous, and false otherwise.

16 6 **Example.** After the pointer assignment AP => TARGET (1:10:2), IS\_CONTIGUOUS (AP) has the value  
 17 false.

## 18 13.9.106 IS\_IOSTAT\_END (I)

19 1 **Description.** IOSTAT value test for end of file.

20 2 **Class.** Elemental function.

21 3 **Argument.** I shall be of type integer.

22 4 **Result Characteristics.** Default logical.

23 5 **Result Value.** The result has the value true if and only if I is a value for the *scalar-int-variable* in an IOSTAT=  
 24 specifier (9.11.5) that would indicate an end-of-file condition.

## 25 13.9.107 IS\_IOSTAT\_EOR (I)

26 1 **Description.** IOSTAT value test for end of record.

27 2 **Class.** Elemental function.

28 3 **Argument.** I shall be of type integer.

29 4 **Result Characteristics.** Default logical.

30 5 **Result Value.** The result has the value true if and only if I is a value for the *scalar-int-variable* in an IOSTAT=  
 31 specifier (9.11.5) that would indicate an end-of-record condition.

## 32 13.9.108 KIND (X)

33 1 **Description.** Value of the kind type parameter of X.

34 2 **Class.** Inquiry function.

1 3 **Argument.** X may be of any intrinsic type. It may be a scalar or an array.

2 4 **Result Characteristics.** Default integer scalar.

3 5 **Result Value.** The result has a value equal to the kind type parameter value of X.

4 6 **Example.** KIND (0.0) has the kind type parameter value of default real.

### 5 13.9.109 LBOUND (ARRAY [, DIM, KIND])

6 1 **Description.** Lower bound(s).

7 2 **Class.** Inquiry function.

8 3 **Arguments.**

9 ARRAY shall be [assumed-rank](#) or an array. It shall not be an unallocated [allocatable](#) variable or a pointer  
10 that is not associated.

11 DIM (optional) shall be an integer scalar with a value in the range  $1 \leq \text{DIM} \leq n$ , where  $n$  is the [rank](#) of ARRAY.  
12 The corresponding [actual argument](#) shall not be an optional dummy argument, a disassociated  
13 pointer, or an unallocated allocatable.

14 KIND (optional) shall be a scalar integer [constant expression](#).

15 4 **Result Characteristics.** Integer. If KIND is present, the kind type parameter is that specified by the value of  
16 KIND; otherwise the kind type parameter is that of default integer type. The result is scalar if DIM is present;  
17 otherwise, the result is an array of [rank](#) one and size  $n$ , where  $n$  is the [rank](#) of ARRAY.

18 5 **Result Value.**

19 *Case (i):* If DIM is present, ARRAY is a [whole array](#), and either ARRAY is an [assumed-size array](#) of [rank](#)  
20 DIM or dimension DIM of ARRAY has nonzero extent, the result has a value equal to the lower  
21 bound for subscript DIM of ARRAY. Otherwise, if DIM is present, the result value is 1.

22 *Case (ii):* LBOUND (ARRAY) has a value whose  $i^{th}$  element is equal to LBOUND (ARRAY,  $i$ ), for  $i = 1, 2,$   
23  $\dots, n$ , where  $n$  is the [rank](#) of ARRAY. LBOUND (ARRAY, KIND=KIND) has a value whose  $i^{th}$   
24 element is equal to LBOUND (ARRAY,  $i$ , KIND), for  $i = 1, 2, \dots, n$ , where  $n$  is the [rank](#) of  
25 ARRAY.

#### NOTE 13.15

If ARRAY is [assumed-rank](#) and has rank zero, DIM cannot be present since it cannot satisfy the requirement  $1 \leq \text{DIM} \leq 0$ .

26 6 **Examples.** If A is declared by the statement

27 7 REAL A (2:3, 7:10)

28 8 then LBOUND (A) is [2, 7] and LBOUND (A, DIM=2) is 7.

### 29 13.9.110 LCOBOUND (COARRAY [, DIM, KIND])

30 1 **Description.** Lower [cobound](#)(s) of a [coarray](#).

31 2 **Class.** Inquiry function.

32 3 **Arguments.**

33 COARRAY shall be a [coarray](#) and may be of any type. It may be a scalar or an array. If it is [allocatable](#) it  
34 shall be allocated.

35 DIM (optional) shall be an integer scalar with a value in the range  $1 \leq \text{DIM} \leq n$ , where  $n$  is the [corank](#)  
36 of COARRAY. The corresponding [actual argument](#) shall not be an optional dummy argument, a  
37 disassociated pointer, or an unallocated allocatable.

1 KIND (optional) shall be a scalar integer [constant expression](#).

2 4 **Result Characteristics.** Integer. If KIND is present, the kind type parameter is that specified by the value of  
 3 KIND; otherwise, the kind type parameter is that of default integer type. The result is scalar if DIM is present;  
 4 otherwise, the result is an array of [rank](#) one and size  $n$ , where  $n$  is the [corank](#) of COARRAY.

5 5 **Result Value.**

6 *Case (i):* If DIM is present, the result has a value equal to the lower [cobound](#) for [codimension](#) DIM of  
 7 COARRAY.

8 *Case (ii):* If DIM is absent, the result has a value whose  $i^{th}$  element is equal to the lower [cobound](#) for [codi-](#)  
 9 [mension](#)  $i$  of COARRAY, for  $i = 1, 2, \dots, n$ , where  $n$  is the [corank](#) of COARRAY.

10 6 **Examples.** If A is allocated by the statement ALLOCATE (A [2:3, 7:\*) then LCOBOUND (A) is [2, 7] and  
 11 LCOBOUND (A, DIM=2) is 7.

## 12 13.9.111 LEADZ (I)

13 1 **Description.** Number of leading zero bits.

14 2 **Class.** [Elemental](#) function.

15 3 **Argument.** I shall be of type integer.

16 4 **Result Characteristics.** Default integer.

17 5 **Result Value.** If all of the bits of I are zero, the result has the value BIT\_SIZE (I). Otherwise, the result has  
 18 the value BIT\_SIZE (I) – 1 –  $k$ , where  $k$  is the position of the leftmost 1 bit in I. The model for the interpretation  
 19 of an integer value as a sequence of bits is in [13.3](#).

20 6 **Examples.** LEADZ (1) has the value 31 if BIT\_SIZE (1) has the value 32.

## 21 13.9.112 LEN (STRING [, KIND])

22 1 **Description.** Length of a character entity.

23 2 **Class.** [Inquiry function](#).

24 3 **Arguments.**

25 STRING shall be of type character. If it is an unallocated [allocatable](#) variable or a pointer that is not  
 26 associated, its [length type parameter](#) shall not be [deferred](#).

27 KIND (optional) shall be a scalar integer [constant expression](#).

28 4 **Result Characteristics.** Integer scalar. If KIND is present, the kind type parameter is that specified by the  
 29 value of KIND; otherwise the kind type parameter is that of default integer type.

30 5 **Result Value.** The result has a value equal to the number of characters in STRING if it is scalar or in an  
 31 element of STRING if it is an array.

32 6 **Example.** If C is declared by the statement

33 7 CHARACTER (11) C (100)

34 8 LEN (C) has the value 11.

## 35 13.9.113 LEN\_TRIM (STRING [, KIND])

36 1 **Description.** Length without trailing blanks.

37 2 **Class.** [Elemental](#) function.

1 3 **Arguments.**

2 STRING shall be of type character.

3 KIND (optional) shall be a scalar integer [constant expression](#).4 4 **Result Characteristics.** Integer. If KIND is present, the kind type parameter is that specified by the value of  
5 KIND; otherwise the kind type parameter is that of default integer type.6 5 **Result Value.** The result has a value equal to the number of characters remaining after any trailing blanks in  
7 STRING are removed. If the argument contains no nonblank characters, the result is zero.8 6 **Examples.** LEN\_TRIM (' A B ') has the value 4 and LEN\_TRIM (' ') has the value 0.9 **13.9.114 LGE (STRING\_A, STRING\_B)**10 1 **Description.** ASCII greater than or equal.11 2 **Class.** [Elemental](#) function.12 3 **Arguments.**13 STRING\_A shall be default character or [ASCII character](#).

14 STRING\_B shall be of type character with the same kind type parameter as STRING\_A.

15 4 **Result Characteristics.** Default logical.16 5 **Result Value.** If the strings are of unequal length, the comparison is made as if the shorter string were extended  
17 on the right with blanks to the length of the longer string. If either string contains a character not in the [ASCII](#)  
18 [character](#) set, the result is processor dependent. The result is true if the strings are equal or if STRING\_A follows  
19 STRING\_B in the [ASCII collating sequence](#); otherwise, the result is false.

## NOTE 13.16

The result is true if both STRING\_A and STRING\_B are of zero length.

20 6 **Example.** LGE ('ONE', 'TWO') has the value false.21 **13.9.115 LGT (STRING\_A, STRING\_B)**22 1 **Description.** ASCII greater than.23 2 **Class.** [Elemental](#) function.24 3 **Arguments.**25 STRING\_A shall be default character or [ASCII character](#).

26 STRING\_B shall be of type character with the same kind type parameter as STRING\_A.

27 4 **Result Characteristics.** Default logical.28 5 **Result Value.** If the strings are of unequal length, the comparison is made as if the shorter string were extended  
29 on the right with blanks to the length of the longer string. If either string contains a character not in the [ASCII](#)  
30 [character](#) set, the result is processor dependent. The result is true if STRING\_A follows STRING\_B in the  
31 [ASCII collating sequence](#); otherwise, the result is false.

## NOTE 13.17

The result is false if both STRING\_A and STRING\_B are of zero length.

32 6 **Example.** LGT ('ONE', 'TWO') has the value false.

**13.9.116 LLE (STRING\_A, STRING\_B)**

**Description.** ASCII less than or equal.

**Class.** [Elemental](#) function.

**Arguments.**

STRING\_A shall be default character or [ASCII character](#).

STRING\_B shall be of type character with the same kind type parameter as STRING\_A.

**Result Characteristics.** Default logical.

**Result Value.** If the strings are of unequal length, the comparison is made as if the shorter string were extended on the right with blanks to the length of the longer string. If either string contains a character not in the [ASCII character](#) set, the result is processor dependent. The result is true if the strings are equal or if STRING\_A precedes STRING\_B in the [ASCII collating sequence](#); otherwise, the result is false.

**NOTE 13.18**

The result is true if both STRING\_A and STRING\_B are of zero length.

**Example.** LLE ('ONE', 'TWO') has the value true.

**13.9.117 LLT (STRING\_A, STRING\_B)**

**Description.** ASCII less than.

**Class.** [Elemental](#) function.

**Arguments.**

STRING\_A shall be default character or [ASCII character](#).

STRING\_B shall be of type character with the same kind type parameter as STRING\_A.

**Result Characteristics.** Default logical.

**Result Value.** If the strings are of unequal length, the comparison is made as if the shorter string were extended on the right with blanks to the length of the longer string. If either string contains a character not in the [ASCII character](#) set, the result is processor dependent. The result is true if STRING\_A precedes STRING\_B in the [ASCII collating sequence](#); otherwise, the result is false.

**NOTE 13.19**

The result is false if both STRING\_A and STRING\_B are of zero length.

**Example.** LLT ('ONE', 'TWO') has the value true.

**13.9.118 LOG (X)**

**Description.** Natural logarithm.

**Class.** [Elemental](#) function.

**Argument.** X shall be of type real or complex. If X is real, its value shall be greater than zero. If X is complex, its value shall not be zero.

**Result Characteristics.** Same as X.

**Result Value.** The result has a value equal to a processor-dependent approximation to  $\log_e X$ . A result of type complex is the principal value with imaginary part  $\omega$  in the range  $-\pi \leq \omega \leq \pi$ . If the real part of X is less

than zero and the imaginary part of X is zero, then the imaginary part of the result is approximately  $\pi$  if the imaginary part of X is positive real zero or the processor does not distinguish between positive and negative real zero, and approximately  $-\pi$  if the imaginary part of X is negative real zero.

**Example.** LOG (10.0) has the value 2.3025851 (approximately).

### 13.9.119 LOG\_GAMMA (X)

**Description.** Logarithm of the absolute value of the gamma function.

**Class.** [Elemental](#) function.

**Argument.** X shall be of type real. Its value shall not be a negative integer or zero.

**Result Characteristics.** Same as X.

**Result Value.** The result has a value equal to a processor-dependent approximation to the natural logarithm of the absolute value of the gamma function of X.

**Example.** LOG\_GAMMA (3.0) has the value 0.693 (approximately).

### 13.9.120 LOG10 (X)

**Description.** Common logarithm.

**Class.** [Elemental](#) function.

**Argument.** X shall be of type real. The value of X shall be greater than zero.

**Result Characteristics.** Same as X.

**Result Value.** The result has a value equal to a processor-dependent approximation to  $\log_{10}X$ .

**Example.** LOG10 (10.0) has the value 1.0 (approximately).

### 13.9.121 LOGICAL (L [, KIND])

**Description.** Conversion between kinds of logical.

**Class.** [Elemental](#) function.

**Arguments.**

L shall be of type logical.

KIND (optional) shall be a scalar integer [constant expression](#).

**Result Characteristics.** Logical. If KIND is present, the kind type parameter is that specified by the value of KIND; otherwise, the kind type parameter is that of default logical.

**Result Value.** The value is that of L.

**Example.** LOGICAL (L [.OR.](#) [.NOT.](#) L) has the value true and is default logical, regardless of the kind type parameter of the logical variable L.

### 13.9.122 MASKL (I [, KIND])

**Description.** Left justified mask.

**Class.** [Elemental](#) function.

**Arguments.**

1 I shall be of type integer. It shall be nonnegative and less than or equal to the number of bits  $z$  of  
 2 the model integer defined for bit manipulation contexts in 13.3 for the kind of the result.  
 3 KIND (optional) shall be a scalar integer [constant expression](#).

4 4 **Result Characteristics.** Integer. If KIND is present, the kind type parameter is that specified by the value of  
 5 KIND; otherwise, the kind type parameter is that of default integer type.

6 5 **Result Value.** The result value has its leftmost  $I$  bits set to 1 and the remaining bits set to 0. The model for  
 7 the interpretation of an integer value as a sequence of bits is in 13.3.

8 6 **Example.** MASKL (3) has the value SHIFTL (7, BIT\_SIZE (0) – 3).

### 9 13.9.123 MASKR (I [, KIND])

10 1 **Description.** Right justified mask.

11 2 **Class.** [Elemental function](#).

12 3 **Arguments.**

13 I shall be of type integer. It shall be nonnegative and less than or equal to the number of bits  $z$  of  
 14 the model integer defined for bit manipulation contexts in 13.3 for the kind of the result.

15 KIND (optional) shall be a scalar integer [constant expression](#).

16 4 **Result Characteristics.** Integer. If KIND is present, the kind type parameter is that specified by the value of  
 17 KIND; otherwise, the kind type parameter is that of default integer type.

18 5 **Result Value.** The result value has its rightmost  $I$  bits set to 1 and the remaining bits set to 0. The model for  
 19 the interpretation of an integer value as a sequence of bits is in 13.3.

20 6 **Example.** MASKR (3) has the value 7.

### 21 13.9.124 MATMUL (MATRIX\_A, MATRIX\_B)

22 1 **Description.** Matrix multiplication.

23 2 **Class.** [Transformational function](#).

24 3 **Arguments.**

25 MATRIX\_A shall be a rank-one or rank-two array of [numeric type](#) or logical type.

26 MATRIX\_B shall be of [numeric type](#) if MATRIX\_A is of [numeric type](#) and of logical type if MATRIX\_A is of  
 27 logical type. It shall be an array of [rank](#) one or two. MATRIX\_A and MATRIX\_B shall not both  
 28 have [rank](#) one. The size of the first (or only) dimension of MATRIX\_B shall equal the size of the  
 29 last (or only) dimension of MATRIX\_A.

30 4 **Result Characteristics.** If the arguments are of [numeric type](#), the type and kind type parameter of the result  
 31 are determined by the types of the arguments as specified in 7.1.9.3 for the \* operator. If the arguments are of  
 32 type logical, the result is of type logical with the kind type parameter of the arguments as specified in 7.1.9.3 for  
 33 the .AND. operator. The shape of the result depends on the shapes of the arguments as follows:

34 *Case (i):* If MATRIX\_A has shape  $[n, m]$  and MATRIX\_B has shape  $[m, k]$ , the result has shape  $[n, k]$ .

35 *Case (ii):* If MATRIX\_A has shape  $[m]$  and MATRIX\_B has shape  $[m, k]$ , the result has shape  $[k]$ .

36 *Case (iii):* If MATRIX\_A has shape  $[n, m]$  and MATRIX\_B has shape  $[m]$ , the result has shape  $[n]$ .

37 5 **Result Value.**

38 *Case (i):* Element  $(i, j)$  of the result has the value SUM (MATRIX\_A  $(i, :)$  \* MATRIX\_B  $(:, j)$ ) if the  
 39 arguments are of [numeric type](#) and has the value ANY (MATRIX\_A  $(i, :)$  .AND. MATRIX\_B  $(:, j)$ ) if the arguments are of logical type.  
 40

1 *Case (ii):* Element ( $j$ ) of the result has the value SUM (MATRIX\_A (:) \* MATRIX\_B (:,  $j$ )) if the arguments  
 2 are of [numeric type](#) and has the value ANY (MATRIX\_A (:) [.AND.](#) MATRIX\_B (:,  $j$ )) if the  
 3 arguments are of logical type.

4 *Case (iii):* Element ( $i$ ) of the result has the value SUM (MATRIX\_A ( $i$ , :) \* MATRIX\_B (:)) if the arguments  
 5 are of [numeric type](#) and has the value ANY (MATRIX\_A ( $i$ , :) [.AND.](#) MATRIX\_B (:)) if the  
 6 arguments are of logical type.

7 6 **Examples.** Let A and B be the matrices  $\begin{bmatrix} 1 & 2 & 3 \\ 2 & 3 & 4 \end{bmatrix}$  and  $\begin{bmatrix} 1 & 2 \\ 2 & 3 \\ 3 & 4 \end{bmatrix}$ ; let X and Y be the vectors [1, 2] and  
 8 [1, 2, 3].

9 *Case (i):* The result of MATMUL (A, B) is the matrix-matrix product AB with the value  $\begin{bmatrix} 14 & 20 \\ 20 & 29 \end{bmatrix}$ .

10 *Case (ii):* The result of MATMUL (X, A) is the vector-matrix product XA with the value [5, 8, 11].

11 *Case (iii):* The result of MATMUL (A, Y) is the matrix-vector product AY with the value [14, 20].

## 12 13.9.125 MAX (A1, A2 [, A3, ...])

13 1 **Description.** Maximum value.

14 2 **Class.** [Elemental function](#).

15 3 **Arguments.** The arguments shall all have the same type which shall be integer, real, or character and they shall  
 16 all have the same kind type parameter.

17 4 **Result Characteristics.** The type and kind type parameter of the result are the same as those of the arguments.  
 18 For arguments of character type, the length of the result is the length of the longest argument.

19 5 **Result Value.** The value of the result is that of the largest argument. For arguments of character type, the  
 20 result is the value that would be selected by application of intrinsic relational operators; that is, the [collating](#)  
 21 [sequence](#) for characters with the kind type parameter of the arguments is applied. If the selected argument is  
 22 shorter than the longest argument, the result is extended with blanks on the right to the length of the longest  
 23 argument.

24 6 **Examples.** MAX (−9.0, 7.0, 2.0) has the value 7.0, MAX ('Z', 'BB') has the value 'Z ', and MAX ([ 'A', 'Z'],  
 25 [ 'BB', 'Y ']) has the value [ 'BB', 'Z '].

## 26 13.9.126 MAXEXPONENT (X)

27 1 **Description.** Maximum exponent of a real model.

28 2 **Class.** [Inquiry function](#).

29 3 **Argument.** X shall be of type real. It may be a scalar or an array.

30 4 **Result Characteristics.** Default integer scalar.

31 5 **Result Value.** The result has the value  $e_{\max}$ , as defined in [13.4](#) for the model representing numbers of the same  
 32 type and kind type parameter as X.

33 6 **Example.** MAXEXPONENT (X) has the value 127 for real X whose model is as in Note [13.4](#).

## 34 13.9.127 MAXLOC (ARRAY, DIM [, MASK, KIND, BACK]) or MAXLOC (ARRAY [, MASK, KIND, BACK])

35 1 **Description.** Location(s) of maximum value.

36 2 **Class.** [Transformational function](#).

### 3 Arguments.

ARRAY shall be an array of type integer, real, or character.

DIM shall be an integer scalar with a value in the range  $1 \leq \text{DIM} \leq n$ , where  $n$  is the [rank](#) of ARRAY.

MASK (optional) shall be of type logical and shall be [conformable](#) with ARRAY.

KIND (optional) shall be a scalar integer [constant expression](#).

BACK (optional) shall be a logical scalar.

**4 Result Characteristics.** Integer. If KIND is present, the kind type parameter is that specified by the value of KIND; otherwise the kind type parameter is that of default integer type. If DIM does not appear, the result is an array of [rank](#) one and of size equal to the [rank](#) of ARRAY; otherwise, the result is of [rank](#)  $n - 1$  and shape  $[d_1, d_2, \dots, d_{\text{DIM}-1}, d_{\text{DIM}+1}, \dots, d_n]$ , where  $[d_1, d_2, \dots, d_n]$  is the shape of ARRAY.

### 5 Result Value.

*Case (i):* If DIM does not appear and MASK is absent, the result is a rank-one array whose element values are the values of the subscripts of an element of ARRAY whose value equals the maximum value of all of the elements of ARRAY. The  $i^{\text{th}}$  subscript returned lies in the range 1 to  $e_i$ , where  $e_i$  is the extent of the  $i^{\text{th}}$  dimension of ARRAY. If ARRAY has size zero, all elements of the result are zero.

*Case (ii):* If DIM does not appear and MASK is present, the result is a rank-one array whose element values are the values of the subscripts of an element of ARRAY, corresponding to a true element of MASK, whose value equals the maximum value of all such elements of ARRAY. The  $i^{\text{th}}$  subscript returned lies in the range 1 to  $e_i$ , where  $e_i$  is the extent of the  $i^{\text{th}}$  dimension of ARRAY. If ARRAY has size zero or every element of MASK has the value false, all elements of the result are zero.

*Case (iii):* If ARRAY has [rank](#) one and DIM is specified, the result has a value equal to that of the first element of MAXLOC (ARRAY [, MASK = MASK, KIND = KIND, BACK = BACK]). Otherwise, if DIM is specified, the value of element  $(s_1, s_2, \dots, s_{\text{DIM}-1}, s_{\text{DIM}+1}, \dots, s_n)$  of the result is equal to

MAXLOC (ARRAY ( $s_1, s_2, \dots, s_{\text{DIM}-1}, :, s_{\text{DIM}+1}, \dots, s_n$ ),

DIM = 1

[, MASK = MASK ( $s_1, s_2, \dots, s_{\text{DIM}-1}, :, s_{\text{DIM}+1}, \dots, s_n$ ),

KIND = KIND,

BACK = BACK]).

6 If only one element has the maximum value, that element's subscripts are returned. Otherwise, if more than one element has the maximum value and BACK is absent or present with the value false, the element whose subscripts are returned is the first such element, taken in array element order. If BACK is present with the value true, the element whose subscripts are returned is the last such element, taken in array element order.

7 If ARRAY has type character, the result is the value that would be selected by application of intrinsic relational operators; that is, the [collating sequence](#) for characters with the kind type parameter of the arguments is applied.

### 8 Examples.

*Case (i):* The value of MAXLOC ([2, 6, 4, 6]) is [2] and the value of MAXLOC ([2, 6, 4, 6], BACK=.TRUE.) is [4].

*Case (ii):* If A has the value  $\begin{bmatrix} 0 & -5 & 8 & -3 \\ 3 & 4 & -1 & 2 \\ 1 & 5 & 6 & -4 \end{bmatrix}$ , MAXLOC (A, MASK = A < 6) has the value [3, 2]. This is independent of the declared lower bounds for A.

*Case (iii):* The value of MAXLOC ([5, -9, 3], DIM = 1) is 1. If B has the value  $\begin{bmatrix} 1 & 3 & -9 \\ 2 & 2 & 6 \end{bmatrix}$ , MAXLOC (B, DIM = 1) is [2, 1, 2] and MAXLOC (B, DIM = 2) is [2, 3]. This is independent of the declared lower bounds for B.

## 13.9.128 MAXVAL (ARRAY, DIM [, MASK]) or MAXVAL (ARRAY [, MASK])

**1 Description.** Maximum value(s) of array.

**2 Class.** [Transformational function](#).

1 3 **Arguments.**

2 ARRAY shall be an array of type integer, real, or character.

3 DIM shall be an integer scalar with a value in the range  $1 \leq \text{DIM} \leq n$ , where  $n$  is the **rank** of ARRAY.4 MASK (optional) shall be of type logical and shall be **conformable** with ARRAY.5 4 **Result Characteristics.** The result is of the same type and type parameters as ARRAY. It is scalar if DIM  
6 does not appear; otherwise, the result has **rank**  $n - 1$  and shape  $[d_1, d_2, \dots, d_{\text{DIM}-1}, d_{\text{DIM}+1}, \dots, d_n]$  where  
7  $[d_1, d_2, \dots, d_n]$  is the shape of ARRAY.8 5 **Result Value.**9 *Case (i):* The result of MAXVAL (ARRAY) has a value equal to the maximum value of all the elements of  
10 ARRAY if the size of ARRAY is not zero. If ARRAY has size zero and type integer or real, the  
11 result has the value of the negative number of the largest magnitude supported by the processor  
12 for numbers of the type and kind type parameter of ARRAY. If ARRAY has size zero and type  
13 character, the result has the value of a string of characters of length LEN (ARRAY), with each  
14 character equal to CHAR (0, KIND (ARRAY)).15 *Case (ii):* The result of MAXVAL (ARRAY, MASK = MASK) has a value equal to that of MAXVAL (PACK  
16 (ARRAY, MASK)).17 *Case (iii):* The result of MAXVAL (ARRAY, DIM = DIM [,MASK = MASK]) has a value equal to that of  
18 MAXVAL (ARRAY [,MASK = MASK]) if ARRAY has **rank** one. Otherwise, the value of element  
19  $(s_1, s_2, \dots, s_{\text{DIM}-1}, s_{\text{DIM}+1}, \dots, s_n)$  of the result is equal to20  $\text{MAXVAL}(\text{ARRAY}(s_1, s_2, \dots, s_{\text{DIM}-1}, :, s_{\text{DIM}+1}, \dots, s_n)$   
21  $[:, \text{MASK} = \text{MASK}(s_1, s_2, \dots, s_{\text{DIM}-1}, :, s_{\text{DIM}+1}, \dots, s_n)])$ .22 6 If ARRAY is of type character, the result is the value that would be selected by application of intrinsic relational  
23 operators; that is, the **collating sequence** for characters with the kind type parameter of the arguments is applied.24 7 **Examples.**25 *Case (i):* The value of MAXVAL ([1, 2, 3]) is 3.26 *Case (ii):* MAXVAL (C, MASK = C < 0.0) is the maximum of the negative elements of C.27 *Case (iii):* If B is the array  $\begin{bmatrix} 1 & 3 & 5 \\ 2 & 7 & 6 \end{bmatrix}$ , MAXVAL (B, DIM = 1) is [2, 7, 6] and MAXVAL (B, DIM = 2) is  
28 [5, 7].29 **13.9.129 MERGE (TSOURCE, FSOURCE, MASK)**30 1 **Description.** Expression value selection.31 2 **Class.** **Elemental** function.32 3 **Arguments.**

33 TSOURCE may be of any type.

34 FSOURCE shall be of the same type and type parameters as TSOURCE.

35 MASK shall be of type logical.

36 4 **Result Characteristics.** Same as TSOURCE.37 5 **Result Value.** The result is TSOURCE if MASK is true and FSOURCE otherwise.38 6 **Examples.** If TSOURCE is the array  $\begin{bmatrix} 1 & 6 & 5 \\ 2 & 4 & 6 \end{bmatrix}$ , FSOURCE is the array  $\begin{bmatrix} 0 & 3 & 2 \\ 7 & 4 & 8 \end{bmatrix}$  and MASK is the  
39 array  $\begin{bmatrix} \text{T} & \cdot & \text{T} \\ \cdot & \cdot & \text{T} \end{bmatrix}$ , where “T” represents true and “.” represents false, then MERGE (TSOURCE, FSOURCE,

1 MASK) is  $\begin{bmatrix} 1 & 3 & 5 \\ 7 & 4 & 6 \end{bmatrix}$ . The value of MERGE (1.0, 0.0,  $K > 0$ ) is 1.0 for  $K = 5$  and 0.0 for  $K = -2$ .

## 2 13.9.130 MERGE\_BITS (I, J, MASK)

3 1 **Description.** Merge of bits under mask.

4 2 **Class.** [Elemental](#) function.

5 3 **Arguments.**

6 I shall be of type integer or a [boz-literal-constant](#).

7 J shall be of type integer or a [boz-literal-constant](#). If both I and J are of type integer they shall have  
8 the same kind type parameter. I and J shall not both be [boz-literal-constants](#).

9 MASK shall be of type integer or a [boz-literal-constant](#). If MASK is of type integer, it shall have the same  
10 kind type parameter as each other argument of type integer.

11 4 **Result Characteristics.** Same as I if I is of type integer; otherwise, same as J.

12 5 **Result Value.** If any argument is a [boz-literal-constant](#), it is first converted as if by the intrinsic function  
13 [INT](#) to the type and kind type parameter of the result. The result has the value of IOR (IAND (I, MASK),  
14 IAND (J, NOT (MASK))).

15 6 **Example.** MERGE\_BITS (13, 18, 22) has the value 4.

## 16 13.9.131 MIN (A1, A2 [, A3, ...])

17 1 **Description.** Minimum value.

18 2 **Class.** [Elemental](#) function.

19 3 **Arguments.** The arguments shall all be of the same type which shall be integer, real, or character and they  
20 shall all have the same kind type parameter.

21 4 **Result Characteristics.** The type and kind type parameter of the result are the same as those of the arguments.  
22 For arguments of character type, the length of the result is the length of the longest argument.

23 5 **Result Value.** The value of the result is that of the smallest argument. For arguments of character type, the  
24 result is the value that would be selected by application of intrinsic relational operators; that is, the [collating](#)  
25 [sequence](#) for characters with the kind type parameter of the arguments is applied. If the selected argument is  
26 shorter than the longest argument, the result is extended with blanks on the right to the length of the longest  
27 argument.

28 6 **Examples.** MIN (−9.0, 7.0, 2.0) has the value −9.0, MIN ('A', 'YY') has the value 'A ', and  
29 MIN (['Z', 'A'], ['YY', 'B ']) has the value ['YY', 'A '].

## 30 13.9.132 MINEXPONENT (X)

31 1 **Description.** Minimum exponent of a real model.

32 2 **Class.** [Inquiry function](#).

33 3 **Argument.** X shall be of type real. It may be a scalar or an array.

34 4 **Result Characteristics.** Default integer scalar.

35 5 **Result Value.** The result has the value  $e_{\min}$ , as defined in [13.4](#) for the model representing numbers of the same  
36 type and kind type parameter as X.

1 6 **Example.** MINEXPONENT (X) has the value  $-126$  for real X whose model is as in Note 13.4.

## 2 13.9.133 MINLOC (ARRAY, DIM [, MASK, KIND, BACK]) or MINLOC (ARRAY [, MASK, KIND, BACK])

3 1 **Description.** Location(s) of minimum value.

4 2 **Class.** Transformational function.

5 3 **Arguments.**

6 ARRAY shall be an array of type integer, real, or character.

7 DIM shall be an integer scalar with a value in the range  $1 \leq \text{DIM} \leq n$ , where  $n$  is the [rank](#) of ARRAY.

8 MASK (optional) shall be of type logical and shall be [conformable](#) with ARRAY.

9 KIND (optional) shall be a scalar integer [constant expression](#).

10 BACK (optional) shall be a logical scalar.

11 4 **Result Characteristics.** Integer. If KIND is present, the kind type parameter is that specified by the value of  
12 KIND; otherwise the kind type parameter is that of default integer type. If DIM does not appear, the result is  
13 an array of [rank](#) one and of size equal to the [rank](#) of ARRAY; otherwise, the result is of [rank](#)  $n - 1$  and shape  
14  $[d_1, d_2, \dots, d_{\text{DIM}-1}, d_{\text{DIM}+1}, \dots, d_n]$ , where  $[d_1, d_2, \dots, d_n]$  is the shape of ARRAY.

15 5 **Result Value.**

16 *Case (i):* If DIM does not appear and MASK is absent the result is a rank-one array whose element values  
17 are the values of the subscripts of an element of ARRAY whose value equals the minimum value  
18 of all the elements of ARRAY. The  $i^{\text{th}}$  subscript returned lies in the range 1 to  $e_i$ , where  $e_i$  is the  
19 extent of the  $i^{\text{th}}$  dimension of ARRAY. If ARRAY has size zero, all elements of the result are zero.

20 *Case (ii):* If DIM does not appear and MASK is present, the result is a rank-one array whose element values  
21 are the values of the subscripts of an element of ARRAY, corresponding to a true element of MASK,  
22 whose value equals the minimum value of all such elements of ARRAY. The  $i^{\text{th}}$  subscript returned  
23 lies in the range 1 to  $e_i$ , where  $e_i$  is the extent of the  $i^{\text{th}}$  dimension of ARRAY. If ARRAY has size  
24 zero or every element of MASK has the value false, all elements of the result are zero.

25 *Case (iii):* If ARRAY has [rank](#) one and DIM is specified, the result has a value equal to that of the first element  
26 of MINLOC (ARRAY [, MASK = MASK, KIND = KIND, BACK = BACK]). Otherwise, if DIM  
27 is specified, the value of element  $(s_1, s_2, \dots, s_{\text{DIM}-1}, s_{\text{DIM}+1}, \dots, s_n)$  of the result is equal to  
MINLOC (ARRAY  $(s_1, s_2, \dots, s_{\text{DIM}-1}, :, s_{\text{DIM}+1}, \dots, s_n)$ ,

28 DIM = 1  
[, MASK = MASK  $(s_1, s_2, \dots, s_{\text{DIM}-1}, :, s_{\text{DIM}+1}, \dots, s_n)$ ,  
KIND = KIND,  
BACK = BACK] ),

29 6 If only one element has the minimum value, that element's subscripts are returned. Otherwise, if more than one  
30 element has the minimum value and BACK is absent or present with the value false, the element whose subscripts  
31 are returned is the first such element, taken in array element order. If BACK is present with the value true, the  
32 element whose subscripts are returned is the last such element, taken in array element order.

33 7 If ARRAY is of type character, the result is the value that would be selected by application of intrinsic relational  
34 operators; that is, the [collating sequence](#) for characters with the kind type parameter of the arguments is applied.

35 8 **Examples.**

36 *Case (i):* The value of MINLOC ([4, 3, 6, 3]) is [2] and the value of MINLOC ([4, 3, 6, 3], BACK = .TRUE.)  
37 is [4].

38 *Case (ii):* If A has the value  $\begin{bmatrix} 0 & -5 & 8 & -3 \\ 3 & 4 & -1 & 2 \\ 1 & 5 & 6 & -4 \end{bmatrix}$ , MINLOC (A, MASK = A > -4) has the value [1, 4].  
39 This is independent of the declared lower bounds for A.

1 *Case (iii):* The value of MINLOC ([5, -9, 3], DIM = 1) is 2. If B has the value  $\begin{bmatrix} 1 & 3 & -9 \\ 2 & 2 & 6 \end{bmatrix}$ , MIN-  
 2 LOC (B, DIM = 1) is [1, 2, 1] and MINLOC (B, DIM = 2) is [3, 1]. This is independent of  
 3 the declared lower bounds for B.

#### 4 **13.9.134 MINVAL (ARRAY, DIM [, MASK]) or MINVAL (ARRAY [, MASK])**

5 1 **Description.** Minimum value(s) of array.

6 2 **Class.** Transformational function.

7 3 **Arguments.**

8 ARRAY shall be an array of type integer, real, or character.

9 DIM shall be an integer scalar with a value in the range  $1 \leq \text{DIM} \leq n$ , where  $n$  is the [rank](#) of ARRAY.

10 MASK (optional) shall be of type logical and shall be [conformable](#) with ARRAY.

11 4 **Result Characteristics.** The result is of the same type and type parameters as ARRAY. It is scalar if DIM  
 12 does not appear; otherwise, the result has [rank](#)  $n - 1$  and shape  $[d_1, d_2, \dots, d_{\text{DIM}-1}, d_{\text{DIM}+1}, \dots, d_n]$  where  
 13  $[d_1, d_2, \dots, d_n]$  is the shape of ARRAY.

14 5 **Result Value.**

15 *Case (i):* The result of MINVAL (ARRAY) has a value equal to the minimum value of all the elements of  
 16 ARRAY if the size of ARRAY is not zero. If ARRAY has size zero and type integer or real, the  
 17 result has the value of the positive number of the largest magnitude supported by the processor  
 18 for numbers of the type and kind type parameter of ARRAY. If ARRAY has size zero and type  
 19 character, the result has the value of a string of characters of length LEN (ARRAY), with each  
 20 character equal to CHAR ( $n - 1$ , KIND (ARRAY)), where  $n$  is the number of characters in the  
 21 [collating sequence](#) for characters with the kind type parameter of ARRAY.

22 *Case (ii):* The result of MINVAL (ARRAY, MASK = MASK) has a value equal to that of MINVAL (PACK  
 23 (ARRAY, MASK)).

24 *Case (iii):* The result of MINVAL (ARRAY, DIM = DIM [, MASK = MASK]) has a value equal to that of  
 25 MINVAL (ARRAY [, MASK = MASK]) if ARRAY has [rank](#) one. Otherwise, the value of element  
 26  $(s_1, s_2, \dots, s_{\text{DIM}-1}, s_{\text{DIM}+1}, \dots, s_n)$  of the result is equal to

27  $\text{MINVAL} (\text{ARRAY} (s_1, s_2, \dots, s_{\text{DIM}-1}, :, s_{\text{DIM}+1}, \dots, s_n)$   
 28  $[\text{, MASK} = \text{MASK} (s_1, s_2, \dots, s_{\text{DIM}-1}, :, s_{\text{DIM}+1}, \dots, s_n) ] )$ .

29 6 If ARRAY is of type character, the result is the value that would be selected by application of intrinsic relational  
 30 operators; that is, the [collating sequence](#) for characters with the kind type parameter of the arguments is applied.

31 7 **Examples.**

32 *Case (i):* The value of MINVAL ([1, 2, 3]) is 1.

33 *Case (ii):* MINVAL (C, MASK = C > 0.0) is the minimum of the positive elements of C.

34 *Case (iii):* If B is the array  $\begin{bmatrix} 1 & 3 & 5 \\ 2 & 4 & 6 \end{bmatrix}$ , MINVAL (B, DIM = 1) is [1, 3, 5] and MINVAL (B, DIM = 2) is  
 35 [1, 2].

#### 36 **13.9.135 MOD (A, P)**

37 1 **Description.** Remainder function.

38 2 **Class.** Elemental function.

39 3 **Arguments.**

40 A shall be of type integer or real.

1 P shall be of the same type and kind type parameter as A. P shall not be zero.

2 4 **Result Characteristics.** Same as A.

3 5 **Result Value.** The value of the result is  $A - \text{INT}(A/P) * P$ .

4 6 **Examples.** MOD (3.0, 2.0) has the value 1.0 (approximately). MOD (8, 5) has the value 3. MOD (−8, 5) has the value −3. MOD (8, −5) has the value 3. MOD (−8, −5) has the value −3.

## 6 13.9.136 MODULO (A, P)

7 1 **Description.** Modulo function.

8 2 **Class.** [Elemental](#) function.

9 3 **Arguments.**

10 A shall be of type integer or real.

11 P shall be of the same type and kind type parameter as A. P shall not be zero.

12 4 **Result Characteristics.** Same as A.

13 5 **Result Value.**

14 *Case (i):* A is of type integer. MODULO (A, P) has the value R such that  $A = Q \times P + R$ , where Q is an integer, the inequalities  $0 \leq R < P$  hold if  $P > 0$ , and  $P < R \leq 0$  hold if  $P < 0$ .

16 *Case (ii):* A is of type real. The value of the result is  $A - \text{FLOOR}(A / P) * P$ .

17 6 **Examples.** MODULO (8, 5) has the value 3. MODULO (−8, 5) has the value 2. MODULO (8, −5) has the value −2. MODULO (−8, −5) has the value −3.

## 19 13.9.137 MOVE\_ALLOC (FROM, TO [, STAT, ERRMSG])

20 1 **Description.** Move an allocation.

21 2 **Class.** [Pure subroutine](#).

22 3 **Arguments.**

23 FROM may be of any type, [rank](#), and [corank](#). It shall be [allocatable](#). It is an [INTENT \(INOUT\)](#) argument.

24 TO shall be [type compatible](#) (4.3.2.3) with FROM and have the same [rank](#) and [corank](#). It shall be [allocatable](#). It shall be [polymorphic](#) if FROM is [polymorphic](#). It is an [INTENT \(OUT\)](#) argument. Each nondeferred parameter of the [declared type](#) of TO shall have the same value as the corresponding parameter of the [declared type](#) of FROM.

28 STAT (optional) shall be a noncoindexed default integer scalar. It is an [INTENT \(OUT\)](#) argument.

29 ERRMSG (optional) shall be a noncoindexed default character scalar. It is an [INTENT \(INOUT\)](#) argument.

30 4 If execution of MOVE\_ALLOC is successful, or if [STAT\\_FAILED\\_IMAGE](#) is assigned to STAT,

- 31 • On invocation of MOVE\_ALLOC, if the allocation status of TO is allocated, it is deallocated. Then, if FROM has an allocation status of allocated on entry to MOVE\_ALLOC, TO becomes allocated with [dynamic type](#), [type parameters](#), [bounds](#), [cobounds](#), and value identical to those that FROM had on entry to MOVE\_ALLOC.
- 35 • If TO has the [TARGET attribute](#), any pointer associated with FROM on entry to MOVE\_ALLOC becomes correspondingly associated with TO. If TO does not have the [TARGET attribute](#), the pointer association status of any pointer associated with FROM on entry becomes undefined.
- 38 • The allocation status of FROM becomes unallocated.

39 5 When a reference to MOVE\_ALLOC is executed for which the FROM argument is a [coarray](#), there is an implicit

synchronization of all **active images** of the **current team**. On those **images**, execution of the segment (8.6.2) following the **CALL statement** is delayed until all other **active images** of the **current team** have executed the same statement the same number of times.

6 If STAT is present and execution is successful, it is assigned the value zero.

7 If an error condition occurs and STAT is absent, **error termination** is initiated. Otherwise, if the **current team** contains a **stopped image**, STAT is assigned the value **STAT\_STOPPED\_IMAGE** from the intrinsic module **ISO\_FORTRAN\_ENV**. Otherwise, if the **current team** contains a **failed image**, no **stopped image**, and no other error condition occurs, STAT is assigned the value **STAT\_FAILED\_IMAGE** from the intrinsic module **ISO\_FORTRAN\_ENV**. Otherwise, STAT is assigned a processor-dependent positive value that differs from that of **STAT\_STOPPED\_IMAGE** or **STAT\_FAILED\_IMAGE**.

8 If the ERRMSG argument is present and an error condition occurs, it is assigned an explanatory message. If no error condition occurs, the definition status and value of ERRMSG are unchanged.

9 **Example.** The example below demonstrates reallocation of GRID to twice its previous size, with its previous contents evenly distributed over the new elements so that intermediate points can be inserted.

```
REAL,ALLOCATABLE :: GRID(:),TEMPGRID(:)
...
ALLOCATE(GRID(-N:N))      ! initial allocation of GRID
...
ALLOCATE(TEMPGRID(-2*N:2*N)) ! allocate bigger grid
TEMPGRID(::2)=GRID ! distribute values to new locations
CALL MOVE_ALLOC(TO=GRID,FROM=TEMPGRID)
```

The old grid is deallocated because TO is **INTENT (OUT)**, and GRID then takes over the new grid allocation.

#### NOTE 13.20

It is expected that the implementation of **allocatable** objects will typically involve descriptors to locate the allocated storage; **MOVE\_ALLOC** could then be implemented by transferring the contents of the descriptor for FROM to the descriptor for TO and clearing the descriptor for FROM.

### 13.9.138 MVBITS (FROM, FROMPOS, LEN, TO, TOPOS)

1 **Description.** Copy a sequence of bits.

2 **Class.** **Elemental** subroutine.

3 **Arguments.**

FROM shall be of type integer. It is an **INTENT (IN)** argument.

FROMPOS shall be of type integer and nonnegative. It is an **INTENT (IN)** argument. FROMPOS + LEN shall be less than or equal to **BIT\_SIZE** (FROM). The model for the interpretation of an integer value as a sequence of bits is in 13.3.

LEN shall be of type integer and nonnegative. It is an **INTENT (IN)** argument.

TO shall be a variable of the same type and kind type parameter value as FROM and may be associated with FROM (12.8.3). It is an **INTENT (INOUT)** argument. TO is defined by copying the sequence of bits of length LEN, starting at position FROMPOS of FROM to position TOPOS of TO. No other bits of TO are altered. On return, the LEN bits of TO starting at TOPOS are equal to the value that the LEN bits of FROM starting at FROMPOS had on entry. The model for the interpretation of an integer value as a sequence of bits is in 13.3.

TOPOS shall be of type integer and nonnegative. It is an **INTENT (IN)** argument. TOPOS + LEN shall be less than or equal to **BIT\_SIZE** (TO).

1 4 **Example.** If TO has the initial value 6, its value after the statement CALL MVBITS (7, 2, 2, TO, 0) is 5.

## 2 13.9.139 NEAREST (X, S)

3 1 **Description.** Adjacent machine number.

4 2 **Class.** [Elemental](#) function.

5 3 **Arguments.**

6 X shall be of type real.

7 S shall be of type real and not equal to zero.

8 4 **Result Characteristics.** Same as X.

9 5 **Result Value.** The result has a value equal to the machine-representable number distinct from X and nearest  
10 to it in the direction of the  $\infty$  with the same sign as S.

11 6 **Example.** NEAREST (3.0, 2.0) has the value  $3 + 2^{-22}$  on a machine whose representation is that of the model  
12 in Note 13.4.

### NOTE 13.21

Unlike other floating-point manipulation functions, NEAREST operates on machine-representable numbers rather than model numbers. On many systems there are machine-representable numbers that lie between adjacent model numbers.

## 13 13.9.140 NEW\_LINE (A)

14 1 **Description.** Newline character.

15 2 **Class.** [Inquiry function](#).

16 3 **Argument.** A shall be of type character. It may be a scalar or an array.

17 4 **Result Characteristics.** Character scalar of length one with the same kind type parameter as A.

18 5 **Result Value.**

19 *Case (i):* If A is default character and the character in position 10 of the [ASCII collating sequence](#) is repres-  
20 entable in the default character set, then the result is [ACHAR](#) (10).

21 *Case (ii):* If A is [ASCII character](#) or [ISO 10646 character](#), then the result is [CHAR](#) (10, [KIND](#) (A)).

22 *Case (iii):* Otherwise, the result is a processor-dependent character that represents a newline in output to files  
23 connected for formatted stream output if there is such a character.

24 *Case (iv):* Otherwise, the result is the blank character.

25 6 **Example.** If there is a suitable newline character, and unit 10 is connected for formatted stream output, the  
26 statement

27 `WRITE (10, '(A)') 'New'//NEW_LINE('a')//'Line'`

28 will write a record containing “New” and then a record containing “Line”.

## 29 13.9.141 NINT (A [, KIND])

30 1 **Description.** Nearest integer.

31 2 **Class.** [Elemental](#) function.

1 3 **Arguments.**

2 A shall be of type real.

3 KIND (optional) shall be a scalar integer [constant expression](#).4 4 **Result Characteristics.** Integer. If KIND is present, the kind type parameter is that specified by the value of  
5 KIND; otherwise, the kind type parameter is that of default integer type.6 5 **Result Value.** The result is the integer nearest A, or if there are two integers equally near A, the result is  
7 whichever such integer has the greater magnitude.8 6 **Example.** NINT (2.783) has the value 3.9 **13.9.142 NORM2 (X) or NORM2 (X, DIM)**10 1 **Description.**  $L_2$  norm of an array.11 2 **Class.** [Transformational function](#).12 3 **Arguments.**

13 X shall be a real array.

14 DIM shall be an integer scalar with a value in the range  $1 \leq \text{DIM} \leq n$ , where  $n$  is the [rank](#) of X.15 4 **Result Characteristics.** The result is of the same type and type parameters as X. It is scalar if DIM does not  
16 appear; otherwise the result has [rank](#)  $n - 1$  and shape  $[d_1, d_2, \dots, d_{\text{DIM}-1}, d_{\text{DIM}+1}, \dots, d_n]$ , where  $n$  is the [rank](#)  
17 of X and  $[d_1, d_2, \dots, d_n]$  is the shape of X.18 5 **Result Value.**19 *Case (i):* The result of NORM2 (X) has a value equal to a processor-dependent approximation to the gener-  
20 alized  $L_2$  norm of X, which is the square root of the sum of the squares of the elements of X. If X  
21 has size zero, the result has the value zero.22 *Case (ii):* The result of NORM2 (X, DIM=DIM) has a value equal to that of NORM2 (X) if X has [rank](#)  
23 one. Otherwise, the value of element  $(s_1, s_2, \dots, s_{\text{DIM}-1}, s_{\text{DIM}+1}, \dots, s_n)$  of the result is equal to  
24 NORM2 (X( $s_1, s_2, \dots, s_{\text{DIM}-1}, :, s_{\text{DIM}+1}, \dots, s_n$ )).

25 6 It is recommended that the processor compute the result without undue overflow or underflow.

26 7 **Example.** The value of NORM2 ([3.0, 4.0]) is 5.0 (approximately). If X has the value  $\begin{bmatrix} 1.0 & 2.0 \\ 3.0 & 4.0 \end{bmatrix}$  then the  
27 value of NORM2 (X, DIM=1) is [3.162, 4.472] (approximately) and the value of NORM2 (X, DIM=2) is [2.236,  
28 5.0] (approximately).29 **13.9.143 NOT (I)**30 1 **Description.** Bitwise complement.31 2 **Class.** [Elemental function](#).32 3 **Argument.** I shall be of type integer.33 4 **Result Characteristics.** Same as I.34 5 **Result Value.** The result has the value obtained by complementing I bit-by-bit according to the following table:

I	NOT (I)
1	0
0	1

- 6 The model for the interpretation of an integer value as a sequence of bits is in 13.3.
- 7 **Example.** If I is represented by the string of bits 01010101, NOT (I) has the binary value 10101010.

### 13.9.144 NULL ([MOLD])

- 1 **Description.** Disassociated pointer or unallocated allocatable entity.
- 2 **Class.** Transformational function.
- 3 **Argument.** MOLD shall be a pointer or allocatable. It may be of any type or may be a procedure pointer. If MOLD is a pointer its pointer association status may be undefined, disassociated, or associated. If MOLD is allocatable its allocation status may be allocated or unallocated. It need not be defined with a value.
- 4 **Result Characteristics.** If MOLD is present, the characteristics are the same as MOLD. If MOLD has deferred type parameters, those type parameters of the result are deferred.
- 5 If MOLD is absent, the characteristics of the result are determined by the entity with which the reference is associated. See Table 13.5. MOLD shall not be absent in any other context. If any type parameters of the contextual entity are deferred, those type parameters of the result are deferred. If any type parameters of the contextual entity are assumed, MOLD shall be present.
- 6 If the context of the reference to NULL is an actual argument in a generic procedure reference, MOLD shall be present if the type, type parameters, or rank are required to resolve the generic reference.

Table 13.5: Characteristics of the result of NULL ( )

Appearance of NULL ( )	Type, type parameters, and rank of result:
right side of a pointer assignment	pointer on the left side
initialization for an object in a declaration	the object
default initialization for a component	the component
in a structure constructor	the corresponding component
as an actual argument	the corresponding dummy argument
in a DATA statement	the corresponding pointer object

- 7 **Result.** The result is a disassociated pointer or an unallocated allocatable entity.
- 8 **Examples.**
- Case (i): REAL, POINTER, DIMENSION (:) :: VEC => NULL ( ) defines the initial association status of VEC to be disassociated.
- Case (ii): The MOLD argument is required in the following:
- ```

INTERFACE GEN
  SUBROUTINE S1 (J, PI)
    INTEGER J
    INTEGER, POINTER :: PI
  END SUBROUTINE S1
  SUBROUTINE S2 (K, PR)
    INTEGER K
    REAL, POINTER :: PR
  END SUBROUTINE S2
END INTERFACE
REAL, POINTER :: REAL_PTR
CALL GEN (7, NULL (REAL_PTR) )      ! Invokes S2

```

### 13.9.145 NUM\_IMAGES ( ) or NUM\_IMAGES (TEAM) or NUM\_IMAGES (TEAM\_NUMBER)

**Description.** Number of *images*.

**Class.** Transformational function.

**Arguments.**

TEAM shall be a scalar of type *TEAM\_TYPE* from the intrinsic module *ISO\_FORTRAN\_ENV*, with a value that identifies the *current* or an *ancestor team*.

TEAM\_NUMBER shall be a positive integer scalar. If the *current team* is not the *initial team*, it shall identify a *team* whose *parent* is the same as that of the *current team*.

**Result Characteristics.** Default integer scalar.

**Result Value.** The number of *images* in the specified *team*, or in the *initial team* if no team is specified.

**Example.** The following code uses *image* 1 to read data and broadcast it to other *images*.

```
REAL :: P[*]
IF (THIS_IMAGE()==1) THEN
  READ (6,*) P
  DO I = 2, NUM_IMAGES()
    P[I] = P
  END DO
END IF
SYNC ALL
```

### 13.9.146 OUT\_OF\_RANGE (X, MOLD, [, ROUND])

**Description.** Whether a value cannot be converted safely.

**Class.** Elemental function.

**Arguments.**

X shall be of type integer or real.

MOLD shall be an integer or real scalar. If it is a variable, it need not be defined.

ROUND (optional) shall be a logical scalar. ROUND shall be present only if X is of type real and MOLD is of type integer.

**Result Characteristics.** Default logical scalar.

**Result Value.**

*Case (i):* If MOLD is of type integer, and ROUND is absent or present with the value false, the result is true if and only if the value of X is an IEEE infinity or NaN, or if the integer with largest magnitude that lies between zero and X inclusive is not representable by objects with the type and kind of MOLD.

*Case (ii):* If MOLD is of type integer, and ROUND is present with the value true, the result is true if and only if the value of X is an IEEE infinity or NaN, or if the integer nearest X, or the integer of greater magnitude if two integers are equally near to X, is not representable by objects with the type and kind of MOLD.

*Case (iii):* Otherwise, the result is true if and only if the value of X is an IEEE infinity or NaN that is not supported by objects of the type and kind of MOLD, or if X is a finite number and the result of

rounding the value of  $X$  (according to the IEEE rounding mode if appropriate) to the extended model for the kind of MOLD has magnitude larger than that of the largest finite number with the same sign as  $X$  that is representable by objects with the type and kind of MOLD.

**Examples.** If INT8 is the kind value for an 8-bit binary integer type, OUT\_OF\_RANGE (−128.5, 0\_INT8) will have the value false and OUT\_OF\_RANGE (−128.5, 0\_INT8, .TRUE.) will have the value true.

### 13.9.147 PACK (ARRAY, MASK [, VECTOR])

**Description.** Array packed into a vector.

**Class.** Transformational function.

**Arguments.**

ARRAY shall be an array of any type.

MASK shall be of type logical and shall be conformable with ARRAY.

VECTOR (optional) shall be of the same type and type parameters as ARRAY and shall have rank one. VECTOR shall have at least as many elements as there are true elements in MASK. If MASK is scalar with the value true, VECTOR shall have at least as many elements as there are in ARRAY.

**Result Characteristics.** The result is an array of rank one with the same type and type parameters as ARRAY. If VECTOR is present, the result size is that of VECTOR; otherwise, the result size is the number  $t$  of true elements in MASK unless MASK is scalar with the value true, in which case the result size is the size of ARRAY.

**Result Value.** Element  $i$  of the result is the element of ARRAY that corresponds to the  $i^{th}$  true element of MASK, taking elements in array element order, for  $i = 1, 2, \dots, t$ . If VECTOR is present and has size  $n > t$ , element  $i$  of the result has the value VECTOR ( $i$ ), for  $i = t + 1, \dots, n$ .

**Examples.** The nonzero elements of an array M with the value  $\begin{bmatrix} 0 & 0 & 0 \\ 9 & 0 & 0 \\ 0 & 0 & 7 \end{bmatrix}$  can be “gathered” by the function PACK. The result of PACK (M, MASK = M /= 0) is [9, 7] and the result of PACK (M, M /= 0, VECTOR = [2, 4, 6, 8, 10, 12]) is [9, 7, 6, 8, 10, 12].

### 13.9.148 PARITY (MASK) or PARITY (MASK, DIM)

**Description.** Array reduced by .NEQV. operation.

**Class.** Transformational function.

**Arguments.**

MASK shall be a logical array.

DIM shall be an integer scalar with a value in the range  $1 \leq \text{DIM} \leq n$ , where  $n$  is the rank of MASK.

**Result Characteristics.** The result is of type logical with the same kind type parameter as MASK. It is scalar if DIM does not appear; otherwise, the result has rank  $n - 1$  and shape  $[d_1, d_2, \dots, d_{\text{DIM}-1}, d_{\text{DIM}+1}, \dots, d_n]$  where  $[d_1, d_2, \dots, d_n]$  is the shape of MASK.

**Result Value.**

*Case (i):* The result of PARITY (MASK) has the value true if an odd number of the elements of MASK are true, and false otherwise.

*Case (ii):* If MASK has rank one, PARITY (MASK, DIM) is equal to PARITY (MASK). Otherwise, the value of element  $(s_1, s_2, \dots, s_{\text{DIM}-1}, s_{\text{DIM}+1}, \dots, s_n)$  of PARITY (MASK, DIM) is equal to PARITY (MASK ( $s_1, s_2, \dots, s_{\text{DIM}-1}, :, s_{\text{DIM}+1}, \dots, s_n$ )).

**Examples.**

1 *Case (i):* The value of PARITY ([T, T, T, F]) is true if T has the value true and F has the value false.

2 *Case (ii):* If B is the array  $\begin{bmatrix} T & T & F \\ T & T & T \end{bmatrix}$ , where T has the value true and F has the value false, then  
 3 PARITY (B, DIM=1) has the value [F, F, T] and PARITY (B, DIM=2) has the value [F, T].

#### 4 **13.9.149 POPCNT (I)**

5 1 **Description.** Number of one bits.

6 2 **Class.** [Elemental](#) function.

7 3 **Argument.** I shall be of type integer.

8 4 **Result Characteristics.** Default integer.

9 5 **Result Value.** The result value is equal to the number of one bits in the sequence of bits of I. The model for  
 10 the interpretation of an integer value as a sequence of bits is in [13.3](#).

11 6 **Examples.** POPCNT ([1, 2, 3, 4, 5, 6]) has the value [1, 1, 2, 1, 2, 2].

#### 12 **13.9.150 POPPAR (I)**

13 1 **Description.** Parity expressed as 0 or 1.

14 2 **Class.** [Elemental](#) function.

15 3 **Argument.** I shall be of type integer.

16 4 **Result Characteristics.** Default integer.

17 5 **Result Value.** POPPAR (I) has the value 1 if POPCNT (I) is odd, and 0 if POPCNT (I) is even.

18 6 **Examples.** POPPAR ([1, 2, 3, 4, 5, 6]) has the value [1, 1, 0, 1, 0, 0].

#### 19 **13.9.151 PRECISION (X)**

20 1 **Description.** Decimal precision of a real model.

21 2 **Class.** [Inquiry function](#).

22 3 **Argument.** X shall be of type real or complex. It may be a scalar or an array.

23 4 **Result Characteristics.** Default integer scalar.

24 5 **Result Value.** The result has the value  $\text{INT}((p-1) * \text{LOG}_{10}(b)) + k$ , where  $b$  and  $p$  are as defined in [13.4](#)  
 25 for the model representing real numbers with the same value for the kind type parameter as X, and where  $k$  is 1  
 26 if  $b$  is an integral power of 10 and 0 otherwise.

27 6 **Example.** PRECISION (X) has the value  $\text{INT}(23 * \text{LOG}_{10}(2.)) = \text{INT}(6.92\dots) = 6$  for real X whose model  
 28 is as in Note [13.4](#).

#### 29 **13.9.152 PRESENT (A)**

30 1 **Description.** Presence of optional argument.

31 2 **Class.** [Inquiry function](#).

32 3 **Argument.** A shall be the name of an optional dummy argument that is accessible in the subprogram in which  
 33 the PRESENT function reference appears. There are no other requirements on A.

34 4 **Result Characteristics.** Default logical scalar.

1 5 **Result Value.** The result has the value true if A is present (12.5.2.12) and otherwise has the value false.

## 2 13.9.153 PRODUCT (ARRAY, DIM [, MASK]) or PRODUCT (ARRAY [, MASK])

3 1 **Description.** Array reduced by multiplication.

4 2 **Class.** Transformational function.

5 3 **Arguments.**

6 ARRAY shall be an array of numeric type.

7 DIM shall be an integer scalar with a value in the range  $1 \leq \text{DIM} \leq n$ , where  $n$  is the rank of ARRAY.

8 MASK (optional) shall be of type logical and shall be conformable with ARRAY.

9 4 **Result Characteristics.** The result is of the same type and kind type parameter as ARRAY. It is scalar if  
10 DIM does not appear; otherwise, the result has rank  $n - 1$  and shape  $[d_1, d_2, \dots, d_{\text{DIM}-1}, d_{\text{DIM}+1}, \dots, d_n]$  where  
11  $[d_1, d_2, \dots, d_n]$  is the shape of ARRAY.

12 5 **Result Value.**

13 *Case (i):* The result of PRODUCT (ARRAY) has a value equal to a processor-dependent approximation to  
14 the product of all the elements of ARRAY or has the value one if ARRAY has size zero.

15 *Case (ii):* The result of PRODUCT (ARRAY, MASK = MASK) has a value equal to a processor-dependent  
16 approximation to the product of the elements of ARRAY corresponding to the true elements of  
17 MASK or has the value one if there are no true elements.

18 *Case (iii):* If ARRAY has rank one, PRODUCT (ARRAY, DIM = DIM [, MASK = MASK]) has a value equal  
19 to that of PRODUCT (ARRAY [, MASK = MASK]). Otherwise, the value of element  $(s_1, s_2, \dots,$   
20  $s_{\text{DIM}-1}, s_{\text{DIM}+1}, \dots, s_n)$  of PRODUCT (ARRAY, DIM = DIM [, MASK = MASK]) is equal to  
21  $\text{PRODUCT}(\text{ARRAY}(s_1, s_2, \dots, s_{\text{DIM}-1}, :, s_{\text{DIM}+1}, \dots, s_n) [, \text{MASK} = \text{MASK}(s_1, s_2, \dots,$   
22  $s_{\text{DIM}-1}, :, s_{\text{DIM}+1}, \dots, s_n) ])$ .

23 6 **Examples.**

24 *Case (i):* The value of PRODUCT ([1, 2, 3]) is 6.

25 *Case (ii):* PRODUCT (C, MASK = C > 0.0) forms the product of the positive elements of C.

26 *Case (iii):* If B is the array  $\begin{bmatrix} 1 & 3 & 5 \\ 2 & 4 & 6 \end{bmatrix}$ , PRODUCT (B, DIM = 1) is [2, 12, 30] and PRODUCT (B, DIM = 2)  
27 is [15, 48].

## 28 13.9.154 RADIX (X)

29 1 **Description.** Base of a numeric model.

30 2 **Class.** Inquiry function.

31 3 **Argument.** X shall be of type integer or real. It may be a scalar or an array.

32 4 **Result Characteristics.** Default integer scalar.

33 5 **Result Value.** The result has the value  $r$  if X is of type integer and the value  $b$  if X is of type real, where  $r$  and  
34  $b$  are as defined in 13.4 for the model representing numbers of the same type and kind type parameter as X.

35 6 **Example.** RADIX (X) has the value 2 for real X whose model is as in Note 13.4.

### 13.9.155 RANDOM\_INIT (REPEATABLE, IMAGE\_DISTINCT)

**Description.** Initialize the pseudorandom number generator.

**Class.** Subroutine.

**Arguments.**

REPEATABLE shall be a logical scalar. It is an **INTENT (IN)** argument. If it has the value true, the seed accessed by the pseudorandom number generator is set to a processor-dependent value that is the same each time RANDOM\_INIT is called from the same **image**. If it has the value false, the seed is set to a processor-dependent, unpredictably different value on each call.

IMAGE\_DISTINCT shall be a logical scalar. It is an **INTENT (IN)** argument. If it has the value true, the seed accessed by the pseudorandom number generator is set to a processor-dependent value that is distinct from the value that would be set by a call to RANDOM\_INIT by another **image**. If it has the value false, the value to which the seed is set does not depend on which **image** calls RANDOM\_INIT.

**Example.** The following statement initializes the pseudorandom number generator so that the seed is different on each call and that the sequence generated will differ from that of another image:

```
CALL RANDOM_INIT (REPEATABLE=.FALSE., IMAGE_DISTINCT=.TRUE.)
```

### 13.9.156 RANDOM\_NUMBER (HARVEST)

**Description.** Generate pseudorandom number(s).

**Class.** Subroutine.

**Argument.** HARVEST shall be of type real. It is an **INTENT (OUT)** argument. It may be a scalar or an array. It is assigned pseudorandom numbers from the uniform distribution in the interval  $0 \leq x < 1$ . If **images** use a common generator, the interleaving of values assigned in unordered segments is processor dependent.

**Example.**

```
REAL X, Y (10, 10)
! Initialize X with a pseudorandom number
CALL RANDOM_NUMBER (HARVEST = X)
CALL RANDOM_NUMBER (Y)
! X and Y contain uniformly distributed random numbers
```

### 13.9.157 RANDOM\_SEED ([SIZE, PUT, GET])

**Description.** Restart or query the pseudorandom number generator.

**Class.** Subroutine.

**Arguments.** There shall either be exactly one or no arguments present.

SIZE (optional) shall be a default integer scalar. It is an **INTENT (OUT)** argument. It is assigned the number  $N$  of integers that the processor uses to hold the value of the seed.

PUT (optional) shall be a default integer array of **rank** one and size  $\geq N$ . It is an **INTENT (IN)** argument. It is used in a processor-dependent manner to compute the seed value accessed by the pseudorandom number generator.

GET (optional) shall be a default integer array of **rank** one and size  $\geq N$ . It is an **INTENT (OUT)** argument. It is assigned the value of the seed.

1 4 If no argument is present, the processor assigns a processor-dependent value to the seed.

2 5 The pseudorandom number generator used by `RANDOM_NUMBER` maintains a seed that is updated during the  
3 execution of `RANDOM_NUMBER` and that can be retrieved or changed by `RANDOM_SEED`. Computation  
4 of the seed from the argument `PUT` is performed in a processor-dependent manner. The value assigned to `GET`  
5 need not be the same as the value of `PUT` in an immediately preceding reference to `RANDOM_SEED`. For  
6 example, following execution of the statements

```
7      CALL RANDOM_SEED (PUT=SEED1)
8      CALL RANDOM_SEED (GET=SEED2)
```

9 SEED2 need not equal SEED1. When the values differ, the use of either value as the `PUT` argument in a  
10 subsequent call to `RANDOM_SEED` shall result in the same sequence of pseudorandom numbers being generated.  
11 For example, after execution of the statements

```
12      CALL RANDOM_SEED (PUT=SEED1)
13      CALL RANDOM_SEED (GET=SEED2)
14      CALL RANDOM_NUMBER (X1)
15      CALL RANDOM_SEED (PUT=SEED2)
16      CALL RANDOM_NUMBER (X2)
```

17 X2 equals X1.

18 6 **Examples.**

```
19      CALL RANDOM_SEED                ! Processor initialization
20      CALL RANDOM_SEED (SIZE = K)      ! Puts size of seed in K
21      CALL RANDOM_SEED (PUT = SEED (1 : K)) ! Define seed
22      CALL RANDOM_SEED (GET = OLD (1 : K)) ! Read current seed
```

### 23 13.9.158 RANGE (X)

24 1 **Description.** Decimal exponent range of a numeric model (13.4).

25 2 **Class.** Inquiry function.

26 3 **Argument.** X shall be of type integer, real, or complex. It may be a scalar or an array.

27 4 **Result Characteristics.** Default integer scalar.

28 5 **Result Value.**

29 *Case (i):* If X is of type integer, the result has the value `INT (LOG10 (HUGE (X)))`.

30 *Case (ii):* If X is of type real, the result has the value `INT (MIN (LOG10 (HUGE (X)), -LOG10 (TINY (X))))`.

31 *Case (iii):* If X is of type complex, the result has the value `RANGE (REAL (X))`.

32 6 **Examples.** `RANGE (X)` has the value 38 for real X whose model is as in Note 13.4, because in this case  
33 `HUGE (X) = (1 - 2-24) × 2127` and `TINY (X) = 2-127`.

### 34 13.9.159 RANK (A)

35 1 **Description.** Rank of a data object.

36 2 **Class.** Inquiry function.

37 3 **Argument.** A shall be a data object of any type.

1 4 **Result Characteristics.** Default integer scalar.

2 5 **Result Value.** The value of the result is the rank of A.

3 6 **Example.** If X is an [assumed-rank](#) dummy argument and its associated [effective argument](#) is an array of rank  
4 3, RANK(X) has the value 3.

### 5 13.9.160 REAL (A [, KIND])

6 1 **Description.** Conversion to real type.

7 2 **Class.** [Elemental](#) function.

8 3 **Arguments.**

9 A shall be of type integer, real, or complex, or a [boz-literal-constant](#).

10 KIND (optional) shall be a scalar integer [constant expression](#).

11 4 **Result Characteristics.** Real.

12 *Case (i):* If A is of type integer or real and KIND is present, the kind type parameter is that specified by the  
13 value of KIND. If A is of type integer or real and KIND is not present, the kind type parameter is  
14 that of default real kind.

15 *Case (ii):* If A is of type complex and KIND is present, the kind type parameter is that specified by the value  
16 of KIND. If A is of type complex and KIND is not present, the kind type parameter is the kind  
17 type parameter of A.

18 *Case (iii):* If A is a [boz-literal-constant](#) and KIND is present, the kind type parameter is that specified by the  
19 value of KIND. If A is a [boz-literal-constant](#) and KIND is not present, the kind type parameter is  
20 that of default real kind.

21 5 **Result Value.**

22 *Case (i):* If A is of type integer or real, the result is equal to a processor-dependent approximation to A.

23 *Case (ii):* If A is of type complex, the result is equal to a processor-dependent approximation to the real part  
24 of A.

25 *Case (iii):* If A is a [boz-literal-constant](#), the value of the result is the value whose internal representation as a  
26 bit sequence is the same as that of A as modified by padding or truncation according to [13.3.3](#). The  
27 interpretation of the bit sequence is processor dependent.

28 6 **Examples.** REAL (−3) has the value −3.0. REAL (Z) has the same kind type parameter and the same value  
29 as the real part of the complex variable Z.

### 30 13.9.161 REDUCE (ARRAY, OPERATION [, MASK, IDENTITY, ORDERED]) or REDUCE (ARRAY, OPERATION, DIM [, MASK, IDENTITY, ORDERED])

31 1 **Description.** General reduction of array.

32 2 **Class.** [Transformational](#) function.

33 3 **Arguments.**

34 ARRAY shall be an array of any type.

35 OPERATION shall be a [pure function](#) with two arguments of the same type and [type parameters](#) as ARRAY.  
36 Its result shall have the same type and [type parameters](#) as ARRAY. The arguments and result shall  
37 not be [polymorphic](#). OPERATION should implement a mathematically associative operation. It  
38 need not be commutative.

39 DIM shall be an integer scalar with a value in the range  $1 \leq \text{DIM} \leq n$ , where  $n$  is the rank of ARRAY.

1 MASK (optional) shall be of type logical and shall be [conformable](#) with ARRAY.  
 2 IDENTITY (optional) shall be scalar with the same type and [type parameters](#) as ARRAY.  
 3 ORDERED (optional) shall be a logical scalar.

4 **4 Result Characteristics.** The result is of the same type and [type parameters](#) as ARRAY. It is scalar if DIM  
 5 does not appear; otherwise, the result has [rank](#)  $n - 1$  and shape  $[d_1, d_2, \dots, d_{\text{DIM}-1}, d_{\text{DIM}+1}, \dots, d_n]$  where  
 6  $[d_1, d_2, \dots, d_n]$  is the shape of ARRAY.

7 **5 Result Value.**

8 *Case (i):* The result of REDUCE (ARRAY, OPERATION [, IDENTITY = IDENTITY, ORDERED =  
 9 ORDERED]) over the sequence of values in ARRAY is the result of an iterative process. The  
 10 initial order of the sequence is array element order. While the sequence has more than one element,  
 11 each iteration involves the execution of  $r = \text{OPERATION}(x, y)$  for adjacent  $x$  and  $y$  in the sequence,  
 12 with  $x$  immediately preceding  $y$ , and the subsequent replacement of  $x$  and  $y$  with  $r$ ; if ORDERED  
 13 is present with the value true,  $x$  and  $y$  shall be the first two elements of the sequence. The process  
 14 continues until the sequence has only one element which is the value of the reduction. If the initial  
 15 sequence is empty, the result has the value IDENTITY if IDENTITY is present, and otherwise,  
 16 [error termination](#) is initiated.

17 *Case (ii):* The result of REDUCE (ARRAY, OPERATION, MASK = MASK [, IDENTITY = IDENTITY,  
 18 ORDERED = ORDERED]) is as for Case (i) except that the initial sequence is only those elements  
 19 of ARRAY for which the corresponding elements of MASK are true.

20 *Case (iii):* If ARRAY has [rank](#) one, REDUCE (ARRAY, OPERATION, DIM = DIM [, MASK = MASK,  
 21 IDENTITY = IDENTITY, ORDERED = ORDERED]) has a value equal to that of REDUCE (AR-  
 22 RAY, OPERATION [, MASK = MASK, IDENTITY = IDENTITY, ORDERED = ORDERED]).  
 23 Otherwise, the value of element  $(s_1, s_2, \dots, s_{\text{DIM}-1}, s_{\text{DIM}+1}, \dots, s_n)$  of REDUCE (ARRAY,  
 24 DIM = DIM [, MASK = MASK, IDENTITY = IDENTITY]) is equal to  
     REDUCE (ARRAY  $(s_1, s_2, \dots, s_{\text{DIM}-1}, :, s_{\text{DIM}+1}, \dots, s_n)$ ,  
             OPERATION = OPERATION,  
             DIM=1  
             [, MASK = MASK  $(s_1, s_2, \dots, s_{\text{DIM}-1}, :, s_{\text{DIM}+1}, \dots, s_n)$ ,  
             IDENTITY = IDENTITY,  
             ORDERED = ORDERED] ).

26 **6 Examples.** The following examples all use the function MY\_MULT, which returns the product of its two integer  
 27 arguments.

28 *Case (i):* The value of REDUCE ([1, 2, 3], MY\_MULT) is 6.

29 *Case (ii):* REDUCE (C, MY\_MULT, MASK= C > 0, IDENTITY=1) forms the product of the positive  
 30 elements of C.

31 *Case (iii):* If B is the array  $\begin{bmatrix} 1 & 3 & 5 \\ 2 & 4 & 6 \end{bmatrix}$ , REDUCE (B, MY\_MULT, DIM = 1) is [2, 12, 30] and REDUCE (B,  
 32 MY\_MULT, DIM = 2) is [15, 48].

#### NOTE 13.22

If OPERATION is not computationally associative, REDUCE without ORDERED=[.TRUE.](#) with the same argument values might not always produce the same result, as the processor can apply the associative law to the evaluation.

### 33 13.9.162 REPEAT (STRING, NCOPIES)

34 **1 Description.** Repetitive string concatenation.

35 **2 Class.** [Transformational function](#).

36 **3 Arguments.**

1     STRING     shall be a character scalar.

2     NCOPIES    shall be an integer scalar. Its value shall not be negative.

3     4 **Result Characteristics.** Character scalar of length NCOPIES times that of STRING, with the same kind type  
4     parameter as STRING.

5     5 **Result Value.** The value of the result is the concatenation of NCOPIES copies of STRING.

6     6 **Examples.** REPEAT ('H', 2) has the value HH. REPEAT ('XYZ', 0) has the value of a zero-length string.

### 7     **13.9.163   RESHAPE (SOURCE, SHAPE [, PAD, ORDER])**

8     1 **Description.** Arbitrary shape array construction.

9     2 **Class.** [Transformational function](#).

10    3 **Arguments.**

11    SOURCE     shall be an array of any type. If PAD is absent or of size zero, the size of SOURCE shall be greater  
12    than or equal to PRODUCT (SHAPE). The size of the result is the product of the values of the  
13    elements of SHAPE.

14    SHAPE       shall be a rank-one integer array. [SIZE](#) ( $x$ ), where  $x$  is the [actual argument](#) corresponding to  
15    [SHAPE](#), shall be a [constant expression](#) whose value is positive and less than 16. It shall not have  
16    an element whose value is negative.

17    PAD (optional) shall be an array of the same type and type parameters as SOURCE.

18    ORDER (optional) shall be of type integer, shall have the same shape as SHAPE, and its value shall be a  
19    permutation of  $(1, 2, \dots, n)$ , where  $n$  is the size of SHAPE. If absent, it is as if it were present with  
20    value  $(1, 2, \dots, n)$ .

21    4 **Result Characteristics.** The result is an array of shape SHAPE (that is, SHAPE (RESHAPE (SOURCE,  
22    SHAPE, PAD, ORDER)) is equal to SHAPE) with the same type and type parameters as SOURCE.

23    5 **Result Value.** The elements of the result, taken in permuted subscript order ORDER  $(1), \dots, \text{ORDER}(n)$ , are  
24    those of SOURCE in normal array element order followed if necessary by those of PAD in array element order,  
25    followed if necessary by additional copies of PAD in array element order.

26    6 **Examples.** RESHAPE ([1, 2, 3, 4, 5, 6], [2, 3]) has the value  $\begin{bmatrix} 1 & 3 & 5 \\ 2 & 4 & 6 \end{bmatrix}$ .  
27    RESHAPE ([1, 2, 3, 4, 5, 6], [2, 4], [0, 0], [2, 1]) has the value  $\begin{bmatrix} 1 & 2 & 3 & 4 \\ 5 & 6 & 0 & 0 \end{bmatrix}$ .

### 28    **13.9.164   RRSPACING (X)**

29    1 **Description.** Reciprocal of relative spacing of model numbers.

30    2 **Class.** [Elemental function](#).

31    3 **Argument.** X shall be of type real.

32    4 **Result Characteristics.** Same as X.

33    5 **Result Value.** The result has the value  $|Y \times b^{-e}| \times b^p = \text{ABS}(\text{FRACTION}(Y)) * \text{RADIX}(X) / \text{EPSILON}(X)$ ,  
34    where  $b$ ,  $e$ , and  $p$  are as defined in 13.4 for Y, the value nearest to X in the model for real values whose kind type  
35    parameter is that of X; if there are two such values, the value of greater absolute value is taken. If X is an IEEE  
36    infinity, the result is an [IEEE NaN](#). If X is an [IEEE NaN](#), the result is that [NaN](#).

37    6 **Example.** RRSPACING  $(-3.0)$  has the value  $0.75 \times 2^{24}$  for reals whose model is as in Note 13.4.

### 13.9.165 SAME\_TYPE\_AS (A, B)

**1 Description.** [Dynamic type](#) equality test.

**2 Class.** [Inquiry function](#).

**3 Arguments.**

A shall be an object of [extensible declared](#) type or [unlimited polymorphic](#). If it is a [polymorphic](#) pointer, it shall not have an undefined association status.

B shall be an object of [extensible declared](#) type or [unlimited polymorphic](#). If it is a [polymorphic](#) pointer, it shall not have an undefined association status.

**4 Result Characteristics.** Default logical scalar.

**5 Result Value.** If the [dynamic type](#) of A or B is [extensible](#), the result is true if and only if the [dynamic type](#) of A is the same as the [dynamic type](#) of B. If neither A nor B has [extensible dynamic](#) type, the result is processor dependent.

#### NOTE 13.23

The [dynamic type](#) of a [disassociated](#) pointer or unallocated [allocatable](#) variable is its [declared type](#). An [unlimited polymorphic](#) entity has no [declared type](#).

**6 Example.** Given the declarations and assignments

```
TYPE T1
  REAL C
END TYPE
TYPE, EXTENDS(T1) :: T2
END TYPE
CLASS(T1), POINTER :: P, Q, R
ALLOCATE(P, Q)
ALLOCATE(T2 :: R)
```

the value of SAME\_TYPE\_AS (P, Q) will be true, and the value of SAME\_TYPE\_AS (P, R) will be false.

### 13.9.166 SCALE (X, I)

**1 Description.** Real number scaled by radix power.

**2 Class.** [Elemental](#) function.

**3 Arguments.**

X shall be of type real.

I shall be of type integer.

**4 Result Characteristics.** Same as X.

**5 Result Value.** The result has the value  $X \times b^I$ , where  $b$  is defined in [13.4](#) for model numbers representing values of X, provided this result is representable; if not, the result is processor dependent.

**6 Example.** SCALE (3.0, 2) has the value 12.0 for reals whose model is as in Note [13.4](#).

### 13.9.167 SCAN (STRING, SET [, BACK, KIND])

**1 Description.** Character set membership search.

**2 Class.** [Elemental](#) function.

**3 Arguments.**

STRING shall be of type character.

SET shall be of type character with the same kind type parameter as STRING.

BACK (optional) shall be of type logical.

KIND (optional) shall be a scalar integer [constant expression](#).

**4 Result Characteristics.** Integer. If KIND is present, the kind type parameter is that specified by the value of KIND; otherwise the kind type parameter is that of default integer type.

**5 Result Value.**

*Case (i):* If BACK is absent or is present with the value false and if STRING contains at least one character that is in SET, the value of the result is the position of the leftmost character of STRING that is in SET.

*Case (ii):* If BACK is present with the value true and if STRING contains at least one character that is in SET, the value of the result is the position of the rightmost character of STRING that is in SET.

*Case (iii):* The value of the result is zero if no character of STRING is in SET or if the length of STRING or SET is zero.

**6 Examples.**

*Case (i):* SCAN ('FORTRAN', 'TR') has the value 3.

*Case (ii):* SCAN ('FORTRAN', 'TR', BACK = [.TRUE.](#)) has the value 5.

*Case (iii):* SCAN ('FORTRAN', 'BCD') has the value 0.

### 13.9.168 SELECTED\_CHAR\_KIND (NAME)

**1 Description.** Character kind selection.

**2 Class.** [Transformational](#) function.

**3 Argument.** NAME shall be default character scalar.

**4 Result Characteristics.** Default integer scalar.

**5 Result Value.** If NAME has the value DEFAULT, then the result has a value equal to that of the kind type parameter of default character. If NAME has the value ASCII, then the result has a value equal to that of the kind type parameter of [ASCII character](#) if the processor supports such a kind; otherwise the result has the value -1. If NAME has the value ISO\_10646, then the result has a value equal to that of the kind type parameter of the [ISO 10646 character](#) kind (corresponding to UCS-4 as specified in ISO/IEC 10646) if the processor supports such a kind; otherwise the result has the value -1. If NAME is a processor-defined name of some other character kind supported by the processor, then the result has a value equal to that kind type parameter value. If NAME is not the name of a supported character type, then the result has the value -1. The NAME is interpreted without respect to case or trailing blanks.

**6 Examples.** SELECTED\_CHAR\_KIND ('ASCII') has the value 1 on a processor that uses 1 as the kind type parameter for the [ASCII character](#) set. The following subroutine produces a Japanese date stamp.

```
SUBROUTINE create_date_string(string)
  INTRINSIC date_and_time,selected_char_kind
  INTEGER,PARAMETER :: ucs4 = selected_char_kind("ISO_10646")
```

```

1      CHARACTER(1,UCS4),PARAMETER :: nen=CHAR(INT(Z'5e74'),UCS4), & !year
2      gatsu=CHAR(INT(Z'6708'),UCS4), & !month
3      nichi=CHAR(INT(Z'65e5'),UCS4) !day
4      CHARACTER(len= *, kind= ucs4) string
5      INTEGER values(8)
6      CALL date_and_time(values=values)
7      WRITE(string,1) values(1),nen,values(2),gatsu,values(3),nichi
8 1  FORMAT(IO,A,IO,A,IO,A)
9      END SUBROUTINE

```

### 13.9.169 SELECTED\_INT\_KIND (R)

1 **Description.** Integer kind selection.

2 **Class.** [Transformational function](#).

3 **Argument.** R shall be an integer scalar.

4 **Result Characteristics.** Default integer scalar.

5 **Result Value.** The result has a value equal to the value of the kind type parameter of an integer type that represents all values  $n$  in the range  $-10^R < n < 10^R$ , or if no such kind type parameter is available on the processor, the result is  $-1$ . If more than one kind type parameter meets the criterion, the value returned is the one with the smallest decimal exponent range, unless there are several such values, in which case the smallest of these kind values is returned.

6 **Example.** Assume a processor supports two integer kinds, 32 with representation method  $r = 2$  and  $q = 31$ , and 64 with representation method  $r = 2$  and  $q = 63$ . On this processor SELECTED\_INT\_KIND (9) has the value 32 and SELECTED\_INT\_KIND (10) has the value 64.

### 13.9.170 SELECTED\_REAL\_KIND ([P, R, RADIX])

1 **Description.** Real kind selection.

2 **Class.** [Transformational function](#).

3 **Arguments.** At least one argument shall be present.

P (optional) shall be an integer scalar.

R (optional) shall be an integer scalar.

RADIX (optional) shall be an integer scalar.

4 **Result Characteristics.** Default integer scalar.

5 **Result Value.** If P or R is absent, the result value is the same as if it were present with the value zero. If RADIX is absent, there is no requirement on the radix of the selected kind.

6 The result has a value equal to a value of the kind type parameter of a real type with decimal precision, as returned by the function PRECISION, of at least P digits, a decimal exponent range, as returned by the function RANGE, of at least R, and a radix, as returned by the function RADIX, of RADIX, if such a kind type parameter is available on the processor.

7 Otherwise, the result is  $-1$  if the processor supports a real type with radix RADIX and exponent range of at least R but not with precision of at least P,  $-2$  if the processor supports a real type with radix RADIX and precision of at least P but not with exponent range of at least R,  $-3$  if the processor supports a real type with radix RADIX but with neither precision of at least P nor exponent range of at least R,  $-4$  if the processor supports a real type

with radix RADIX and either precision of at least P or exponent range of at least R but not both together, and  $-5$  if the processor supports no real type with radix RADIX.

If more than one kind type parameter value meets the criteria, the value returned is the one with the smallest decimal precision, unless there are several such values, in which case the smallest of these kind values is returned.

**Example.** `SELECTED_REAL_KIND (6, 70)` has the value `KIND (0.0)` on a machine that supports a default real approximation method with  $b = 16$ ,  $p = 6$ ,  $e_{\min} = -64$ , and  $e_{\max} = 63$  and does not have a less precise approximation method.

### 13.9.171 SET\_EXPONENT (X, I)

**Description.** Real value with specified exponent.

**Class.** [Elemental](#) function.

**Arguments.**

X shall be of type real.

I shall be of type integer.

**Result Characteristics.** Same as X.

**Result Value.** If X has the value zero, the result has the same value as X. If X is an IEEE infinity, the result is an [IEEE NaN](#). If X is an [IEEE NaN](#), the result is the same [NaN](#). Otherwise, the result has the value  $X \times b^{I-e}$ , where  $b$  and  $e$  are as defined in [13.4](#) for the representation for the value of X in the extended real model for the kind of X.

**Example.** `SET_EXPONENT (3.0, 1)` has the value 1.5 for reals whose model is as in Note [13.4](#).

### 13.9.172 SHAPE (SOURCE [, KIND])

**Description.** Shape of an array or a scalar.

**Class.** [Inquiry](#) function.

**Arguments.**

SOURCE shall be a scalar or array of any type. It shall not be an unallocated [allocatable](#) variable or a pointer that is not associated. It shall not be an [assumed-size array](#).

KIND (optional) shall be a scalar integer [constant expression](#).

**Result Characteristics.** Integer. If KIND is present, the kind type parameter is that specified by the value of KIND; otherwise the kind type parameter is that of default integer type. The result is an array of [rank](#) one whose size is equal to the [rank](#) of SOURCE.

**Result Value.** The result has a value whose  $i^{th}$  element is equal to the extent of dimension  $i$  of SOURCE, except that if SOURCE is [assumed-rank](#), and associated with an [assumed-size array](#), the last element is equal to  $-1$ .

**Examples.** The value of `SHAPE (A (2:5, -1:1))` is `[4, 3]`. The value of `SHAPE (3)` is the rank-one array of size zero.

### 13.9.173 SHIFTA (I, SHIFT)

**Description.** Right shift with fill.

**Class.** [Elemental](#) function.

**Arguments.**

1 I shall be of type integer.  
 2 SHIFT shall be of type integer. It shall be nonnegative and less than or equal to BIT\_SIZE (I).

3 4 **Result Characteristics.** Same as I.

4 5 **Result Value.** The result has the value obtained by shifting the bits of I to the right SHIFT bits and replicating  
 5 the leftmost bit of I in the left SHIFT bits.

6 6 If SHIFT is zero the result is I. Bits shifted out from the right are lost. The model for the interpretation of an  
 7 integer value as a sequence of bits is in [13.3](#).

8 7 **Example.** SHIFTA (IBSET (0, BIT\_SIZE (0) - 1), 2) is equal to SHIFTL (7, BIT\_SIZE (0) - 3).

### 9 **13.9.174 SHIFTL (I, SHIFT)**

10 1 **Description.** Left shift.

11 2 **Class.** [Elemental](#) function.

12 3 **Arguments.**

13 I shall be of type integer.

14 SHIFT shall be of type integer. It shall be nonnegative and less than or equal to BIT\_SIZE (I).

15 4 **Result Characteristics.** Same as I.

16 5 **Result Value.** The value of the result is ISHFT (I, SHIFT).

17 6 **Examples.** SHIFTL (3, 1) has the value 6.

### 18 **13.9.175 SHIFTR (I, SHIFT)**

19 1 **Description.** Right shift.

20 2 **Class.** [Elemental](#) function.

21 3 **Arguments.**

22 I shall be of type integer.

23 SHIFT shall be of type integer. It shall be nonnegative and less than or equal to BIT\_SIZE (I).

24 4 **Result Characteristics.** Same as I.

25 5 **Result Value.** The value of the result is ISHFT (I, -SHIFT).

26 6 **Examples.** SHIFTR (3, 1) has the value 1.

### 27 **13.9.176 SIGN (A, B)**

28 1 **Description.** Magnitude of A with the sign of B.

29 2 **Class.** [Elemental](#) function.

30 3 **Arguments.**

31 A shall be of type integer or real.

32 B shall be of the same type as A.

33 4 **Result Characteristics.** Same as A.

34 5 **Result Value.**

- 1 *Case (i):* If  $B > 0$ , the value of the result is  $|A|$ .
- 2 *Case (ii):* If  $B < 0$ , the value of the result is  $-|A|$ .
- 3 *Case (iii):* If  $B$  is of type integer and  $B=0$ , the value of the result is  $|A|$ .
- 4 *Case (iv):* If  $B$  is of type real and is zero, then:
- 5 • if the processor does not distinguish between positive and negative real zero, or if  $B$  is positive
  - 6 real zero, the value of the result is  $|A|$ ;
  - 7 • if the processor distinguishes between positive and negative real zero, and  $B$  is negative real
  - 8 zero, the value of the result is  $-|A|$ .

9 6 **Example.** SIGN ( $-3.0, 2.0$ ) has the value 3.0.

## 10 13.9.177 SIN (X)

11 1 **Description.** Sine function.

12 2 **Class.** [Elemental](#) function.

13 3 **Argument.** X shall be of type real or complex.

14 4 **Result Characteristics.** Same as X.

15 5 **Result Value.** The result has a value equal to a processor-dependent approximation to  $\sin(X)$ . If X is of type

16 real, it is regarded as a value in radians. If X is of type complex, its real part is regarded as a value in radians.

17 6 **Example.** SIN (1.0) has the value 0.84147098 (approximately).

## 18 13.9.178 SINH (X)

19 1 **Description.** Hyperbolic sine function.

20 2 **Class.** [Elemental](#) function.

21 3 **Argument.** X shall be of type real or complex.

22 4 **Result Characteristics.** Same as X.

23 5 **Result Value.** The result has a value equal to a processor-dependent approximation to  $\sinh(X)$ . If X is of type

24 complex its imaginary part is regarded as a value in radians.

25 6 **Example.** SINH (1.0) has the value 1.1752012 (approximately).

## 26 13.9.179 SIZE (ARRAY [, DIM, KIND])

27 1 **Description.** Size of an array or one extent.

28 2 **Class.** [Inquiry function](#).

29 3 **Arguments.**

30 ARRAY shall be [assumed-rank](#) or an array. It shall not be an unallocated [allocatable](#) variable or a pointer

31 that is not associated. If ARRAY is an [assumed-size array](#), DIM shall be present with a value less

32 than the [rank](#) of ARRAY.

33 DIM (optional) shall be an integer scalar with a value in the range  $1 \leq \text{DIM} \leq n$ , where  $n$  is the [rank](#) of ARRAY.

34 KIND (optional) shall be a scalar integer [constant expression](#).

35 4 **Result Characteristics.** Integer scalar. If KIND is present, the kind type parameter is that specified by the

36 value of KIND; otherwise the kind type parameter is that of default integer type.

5 **Result Value.** If DIM is present, the result has a value equal to the extent of dimension DIM of ARRAY, except  
 2 that if ARRAY is [assumed-rank](#) and associated with an [assumed-size array](#) and DIM is present with a value equal  
 3 to the rank of ARRAY, the value is  $-1$ .

6 If DIM is absent and ARRAY is [assumed-rank](#), the result has a value equal to [PRODUCT](#)([SHAPE](#)(ARRAY,  
 5 KIND)). Otherwise, the result has a value equal to the total number of elements of ARRAY.

7 **Examples.** The value of [SIZE](#) (A (2:5, -1:1), DIM=2) is 3. The value of [SIZE](#) (A (2:5, -1:1)) is 12.

#### NOTE 13.24

If ARRAY is [assumed-rank](#) and has rank zero, DIM cannot be present since it cannot satisfy the requirement  $1 \leq \text{DIM} \leq 0$ .

### 13.9.180 SPACING (X)

1 **Description.** Spacing of model numbers.

2 **Class.** [Elemental](#) function.

3 **Argument.** X shall be of type real.

4 **Result Characteristics.** Same as X.

5 **Result Value.** If X does not have the value zero and is not an IEEE infinity or [NaN](#), the result has the value  
 $b^{\max(e-p, e_{\min}-1)}$ , where  $b$ ,  $e$ , and  $p$  are as defined in [13.4](#) for the value nearest to X in the model for real values  
 whose kind type parameter is that of X; if there are two such values the value of greater absolute value is taken.  
 If X has the value zero, the result is the same as that of [TINY](#) (X). If X is an IEEE infinity, the result is an [IEEE](#)  
[NaN](#). If X is an [IEEE NaN](#), the result is that [NaN](#).

6 **Example.** [SPACING](#) (3.0) has the value  $2^{-22}$  for reals whose model is as in [Note 13.4](#).

### 13.9.181 SPREAD (SOURCE, DIM, NCOPIES)

1 **Description.** Value replicated in a new dimension.

2 **Class.** [Transformational](#) function.

3 **Arguments.**

SOURCE shall be a scalar or array of any type. The [rank](#) of SOURCE shall be less than 15.

DIM shall be an integer scalar with value in the range  $1 \leq \text{DIM} \leq n+1$ , where  $n$  is the [rank](#) of SOURCE.

NCOPIES shall be an integer scalar.

4 **Result Characteristics.** The result is an array of the same type and type parameters as SOURCE and of [rank](#)  
 $n+1$ , where  $n$  is the [rank](#) of SOURCE.

Case (i): If SOURCE is scalar, the shape of the result is (MAX (NCOPIES, 0)).

Case (ii): If SOURCE is an array with shape  $[d_1, d_2, \dots, d_n]$ , the shape of the result is  $[d_1, d_2, \dots, d_{\text{DIM}-1},$   
 MAX (NCOPIES, 0),  $d_{\text{DIM}}, \dots, d_n]$ .

5 **Result Value.**

Case (i): If SOURCE is scalar, each element of the result has a value equal to SOURCE.

Case (ii): If SOURCE is an array, the element of the result with subscripts  $(r_1, r_2, \dots, r_{n+1})$  has the value  
 SOURCE  $(r_1, r_2, \dots, r_{\text{DIM}-1}, r_{\text{DIM}+1}, \dots, r_{n+1})$ .

6 **Examples.** If A is the array [2, 3, 4], [SPREAD](#) (A, DIM=1, NCOPIES=NC) is the array  $\begin{bmatrix} 2 & 3 & 4 \\ 2 & 3 & 4 \\ 2 & 3 & 4 \end{bmatrix}$  if NC  
 has the value 3 and is a zero-sized array if NC has the value 0.

### 13.9.182 SQRT (X)

**Description.** Square root.

**Class.** [Elemental](#) function.

**Argument.** X shall be of type real or complex. If X is real, its value shall be greater than or equal to zero.

**Result Characteristics.** Same as X.

**Result Value.** The result has a value equal to a processor-dependent approximation to the square root of X. A result of type complex is the principal value with the real part greater than or equal to zero. When the real part of the result is zero, the imaginary part has the same sign as the imaginary part of X.

**Example.** SQRT (4.0) has the value 2.0 (approximately).

### 13.9.183 STOPPED\_IMAGES([TEAM, KIND])

**Description.** Indices of stopped images.

**Class.** [Transformational](#) function.

**Arguments.**

TEAM (optional) shall be a scalar of type [TEAM\\_TYPE](#) from the intrinsic module [ISO\\_FORTRAN\\_ENV](#), whose value identifies the [current](#) or an [ancestor](#) team. If TEAM is absent the team specified is the [current team](#).

KIND (optional) shall be a scalar integer [constant expression](#). Its value shall specify a representation method whose decimal range is at least as large as that of default integer.

**Result Characteristics.** Integer. If KIND is present, the kind type parameter is that specified by the value of KIND; otherwise, the kind type parameter is that of default integer type. The result is an array of rank one whose size is equal to the number of images in the specified [team](#) that have initiated [normal termination](#).

**Result Value.** The elements of the result are the values of the indices of the images that have are known to have initiated [normal termination](#) in the specified team, in numerically increasing order. If the executing image has previously executed an [image control statement](#) whose STAT= specifier assigned the value [STAT\\_STOPPED\\_IMAGE](#) from the intrinsic module [ISO\\_FORTRAN\\_ENV](#) or invoked a [collective subroutine](#) whose STAT argument was assigned [STAT\\_STOPPED\\_IMAGE](#), at least one of the images participating in that [image control statement](#) or collective invocation shall be known to have initiated [normal termination](#).

**Examples.** If image 3 is the only image in the [current team](#) that is known to have initiated [normal termination](#), STOPPED\_IMAGES() will have the value [3]. If there are no images in the [current team](#) that have initiated [normal termination](#), the value of STOPPED\_IMAGES() will be a zero-sized array.

### 13.9.184 STORAGE\_SIZE (A [, KIND])

**Description.** Storage size in bits.

**Class.** [Inquiry](#) function.

**Arguments.**

A shall be a data object of any type. If it is [polymorphic](#) it shall not be an undefined [pointer](#). If it is [unlimited polymorphic](#) or has any [deferred type parameters](#), it shall not be an unallocated [allocatable](#) variable or a [disassociated](#) or undefined pointer.

KIND (optional) shall be a scalar integer [constant expression](#).

**Result Characteristics.** Integer scalar. If KIND is present, the kind type parameter is that specified by the value of KIND; otherwise, the kind type parameter is that of default integer type.

1 5 **Result Value.** The result value is the size expressed in bits for an element of an array that has the [dynamic](#)  
 2 [type](#) and type parameters of A. If the type and type parameters are such that [storage association](#) (16.5.3) applies,  
 3 the result is consistent with the [named constants](#) defined in the intrinsic module [ISO\\_FORTRAN\\_ENV](#).

**NOTE 13.25**

An array element might take more bits to store than an isolated scalar, since any hardware-imposed alignment requirements for array elements might not apply to a simple scalar variable.

**NOTE 13.26**

This is intended to be the size in memory that an object takes when it is stored; this might differ from the size it takes during expression handling (which might be the native register size) or when stored in a file. If an object is never stored in memory but only in a register, this function nonetheless returns the size it would take if it were stored in memory.

4 6 **Example.** `STORAGE_SIZE` (1.0) has the same value as the [named constant](#) `NUMERIC_STORAGE_SIZE` in  
 5 the intrinsic module [ISO\\_FORTRAN\\_ENV](#).

## 6 13.9.185 SUM (ARRAY, DIM [, MASK]) or SUM (ARRAY [, MASK])

7 1 **Description.** Array reduced by addition.

8 2 **Class.** [Transformational function](#).

9 3 **Arguments.**

10 ARRAY shall be an array of [numeric type](#).

11 DIM shall be an integer scalar with a value in the range  $1 \leq \text{DIM} \leq n$ , where  $n$  is the [rank](#) of ARRAY.

12 MASK (optional) shall be of type logical and shall be [conformable](#) with ARRAY.

13 4 **Result Characteristics.** The result is of the same type and kind type parameter as ARRAY. It is scalar if  
 14 DIM does not appear; otherwise, the result has [rank](#)  $n - 1$  and shape  $[d_1, d_2, \dots, d_{\text{DIM}-1}, d_{\text{DIM}+1}, \dots, d_n]$  where  
 15  $[d_1, d_2, \dots, d_n]$  is the shape of ARRAY.

16 5 **Result Value.**

17 *Case (i):* The result of `SUM (ARRAY)` has a value equal to a processor-dependent approximation to the sum  
 18 of all the elements of ARRAY or has the value zero if ARRAY has size zero.

19 *Case (ii):* The result of `SUM (ARRAY, MASK = MASK)` has a value equal to a processor-dependent approx-  
 20 imation to the sum of the elements of ARRAY corresponding to the true elements of MASK or has  
 21 the value zero if there are no true elements.

22 *Case (iii):* If ARRAY has [rank](#) one, `SUM (ARRAY, DIM = DIM [, MASK = MASK])` has a value equal to that  
 23 of `SUM (ARRAY [, MASK = MASK])`. Otherwise, the value of element  $(s_1, s_2, \dots, s_{\text{DIM}-1}, s_{\text{DIM}+1},$   
 24  $\dots, s_n)$  of `SUM (ARRAY, DIM = DIM [, MASK = MASK])` is equal to

25  $\text{SUM} (\text{ARRAY} (s_1, s_2, \dots, s_{\text{DIM}-1}, :, s_{\text{DIM}+1}, \dots, s_n) [, \text{MASK} = \text{MASK} (s_1, s_2, \dots, s_{\text{DIM}-1},$   
 26  $:, s_{\text{DIM}+1}, \dots, s_n) ] )$ .

27 6 **Examples.**

28 *Case (i):* The value of `SUM ([1, 2, 3])` is 6.

29 *Case (ii):* `SUM (C, MASK = C > 0.0)` forms the sum of the positive elements of C.

30 *Case (iii):* If B is the array  $\begin{bmatrix} 1 & 3 & 5 \\ 2 & 4 & 6 \end{bmatrix}$ , `SUM (B, DIM = 1)` is [3, 7, 11] and `SUM (B, DIM = 2)` is [9, 12].

### 13.9.186 SYSTEM\_CLOCK ([COUNT, COUNT\_RATE, COUNT\_MAX])

**Description.** Query system clock.

**Class.** Subroutine.

**Arguments.**

COUNT (optional) shall be an integer scalar. It is an **INTENT (OUT)** argument. It is assigned a processor-dependent value based on the value of a processor clock, or **—HUGE (COUNT)** if there is no clock for the invoking **image**. The processor-dependent value is incremented by one for each clock count until the value COUNT\_MAX is reached and is reset to zero at the next count. It lies in the range 0 to COUNT\_MAX if there is a clock.

COUNT\_RATE (optional) shall be an integer or real scalar. It is an **INTENT (OUT)** argument. It is assigned a processor-dependent approximation to the number of processor clock counts per second, or zero if there is no clock for the invoking **image**.

COUNT\_MAX (optional) shall be an integer scalar. It is an **INTENT (OUT)** argument. It is assigned the maximum value that COUNT can have, or zero if there is no clock for the invoking **image**.

Whether an **image** has no clock, has a single clock of its own, or shares a clock with another **image**, is processor dependent.

**Example.** If the processor clock is a 24-hour clock that registers time at approximately 18.20648193 ticks per second, at 11:30 A.M. the reference

CALL SYSTEM\_CLOCK (COUNT = C, COUNT\_RATE = R, COUNT\_MAX = M)

defines  $C = (11 \times 3600 + 30 \times 60) \times 18.20648193 = 753748$ ,  $R = 18.20648193$ , and  $M = 24 \times 3600 \times 18.20648193 - 1 = 1573039$ .

### 13.9.187 TAN (X)

**Description.** Tangent function.

**Class.** **Elemental** function.

**Argument.** X shall be of type real or complex.

**Result Characteristics.** Same as X.

**Result Value.** The result has a value equal to a processor-dependent approximation to  $\tan(X)$ . If X is of type real, it is regarded as a value in radians. If X is of type complex, its real part is regarded as a value in radians.

**Example.** TAN (1.0) has the value 1.5574077 (approximately).

### 13.9.188 TANH (X)

**Description.** Hyperbolic tangent function.

**Class.** **Elemental** function.

**Argument.** X shall be of type real or complex.

**Result Characteristics.** Same as X.

**Result Value.** The result has a value equal to a processor-dependent approximation to  $\tanh(X)$ . If X is of type complex its imaginary part is regarded as a value in radians.

**Example.** TANH (1.0) has the value 0.76159416 (approximately).

### 13.9.189 TEAM\_NUMBER([TEAM])

**1 Description.** Team number.

**2 Class.** Transformational function.

**3 Argument.** TEAM (optional) shall be a scalar of type `TEAM_TYPE` from the intrinsic module `ISO_FORTRAN_ENV`, whose value identifies the `current` or an `ancestor` team. If TEAM is absent, the team specified is the `current` team.

**4 Result Characteristics.** Default integer scalar.

**5 Result Value.** If the specified team is the `initial` team, the result value is  $-1$ ; otherwise, the result value is the team number identifying the team of the invoking image within the specified team.

**6 Example.** The team number can be used to control which statements get executed, for example:

```
TYPE(Team_Type) :: ODD_EVEN
...
FORM TEAM (2-MOD(ME,2), ODD_EVEN)
...
CHANGE TEAM (ODD_EVEN)
  SELECT CASE (TEAM_NUMBER())
  CASE (1)
    ! Case for images with odd image indices in the parent team.
  CASE (2)
    ! Case for images with even image indices in the parent team.
  END SELECT
END TEAM
```

### 13.9.190 THIS\_IMAGE ([TEAM]) or THIS\_IMAGE (COARRAY [, TEAM]) or THIS\_IMAGE (COARRAY, DIM [, TEAM])

**1 Description.** `Cosubscript(s)` for this `image`.

**2 Class.** Transformational function.

**3 Arguments.**

COARRAY shall be a `coarray` of any type except `TEAM_TYPE` from the intrinsic module `ISO_FORTRAN_ENV`. If it is `allocatable` it shall be allocated.

#### Unresolved Technical Issue 029

##### Resolving ambiguity in THIS\_IMAGE.

The TS introduces an ambiguity between `THIS_IMAGE(X)` when X is a coarray of type `TEAM_TYPE`: is this the `THIS_IMAGE(TEAM)` form or the `THIS_IMAGE(COARRAY)` form? I resolved this by prohibiting the COARRAY argument from having type `TEAM_TYPE`. If that is the correct resolution, this UTI can be deleted. Otherwise, additional edits will be required.

DIM shall be an integer scalar. Its value shall be in the range  $1 \leq \text{DIM} \leq n$ , where  $n$  is the `corank` of COARRAY.

TEAM (optional) shall be a scalar of type `TEAM_TYPE` from the intrinsic module `ISO_FORTRAN_ENV`, whose value identifies the `current` or an `ancestor` team. If COARRAY appears, it shall be established in that team.

1 4 **Result Characteristics.** Default integer. It is scalar if COARRAY does not appear or DIM appears; otherwise,  
2 the result has [rank](#) one and its size is equal to the [corank](#) of COARRAY.

3 5 **Result Value.**

4 *Case (i):* The result of THIS\_IMAGE ([TEAM]) is a scalar with a value equal to the [index](#) of the invoking  
5 [image](#) in the [team](#) specified by TEAM, if present, or in the [current team](#) if absent.

6 *Case (ii):* The result of THIS\_IMAGE (COARRAY [, TEAM = TEAM]) is the sequence of [cosubscript](#) values  
7 for COARRAY that would specify the invoking [image](#) in the [team](#) specified by TEAM, if present,  
8 or in the [current team](#) if absent.

9 *Case (iii):* The result of THIS\_IMAGE (COARRAY, DIM [, TEAM = TEAM]) is the value of [cosubscript](#)  
10 DIM in the sequence of [cosubscript](#) values for COARRAY that would specify the invoking [image](#) in  
11 the [team](#) specified by TEAM, if present, or in the [current team](#) if absent.

12 6 **Examples.** If A is declared by the statement

13 REAL A (10, 20) [10, 0:9, 0:~]

14 then on [image](#) 5, THIS\_IMAGE ( ) has the value 5 and THIS\_IMAGE (A) has the value [5, 0, 0]. For the same  
15 [coarray](#) on [image](#) 213, THIS\_IMAGE (A) has the value [3, 1, 2].

16 7 The following code uses [image](#) 1 to read data. The other [images](#) then copy the data.

17 IF (THIS\_IMAGE()==1) READ (\*,\*) P

18 SYNC ALL

19 P = P[1]

## 20 13.9.191 TINY (X)

21 1 **Description.** Smallest positive model number.

22 2 **Class.** [Inquiry function](#).

23 3 **Argument.** X shall be a real scalar or array.

24 4 **Result Characteristics.** Scalar with the same type and kind type parameter as X.

25 5 **Result Value.** The result has the value  $b^{e_{\min}-1}$  where  $b$  and  $e_{\min}$  are as defined in [13.4](#) for the model representing  
26 numbers of the same type and kind type parameter as X.

27 6 **Example.** TINY (X) has the value  $2^{-127}$  for real X whose model is as in Note [13.4](#).

## 28 13.9.192 TRAILZ (I)

29 1 **Description.** Number of trailing zero bits.

30 2 **Class.** [Elemental function](#).

31 3 **Argument.** I shall be of type integer.

32 4 **Result Characteristics.** Default integer.

33 5 **Result Value.** If all of the bits of I are zero, the result value is BIT\_SIZE (I). Otherwise, the result value is the  
34 position of the rightmost 1 bit in I. The model for the interpretation of an integer value as a sequence of bits is  
35 in [13.3](#).

36 6 **Examples.** TRAILZ (8) has the value 3.

### 13.9.193 TRANSFER (SOURCE, MOLD [, SIZE])

**Description.** Transfer physical representation.

**Class.** Transformational function.

**Arguments.**

SOURCE shall be a scalar or array of any type.

MOLD shall be a scalar or array of any type. If it is a variable, it need not be defined.

SIZE (optional) shall be an integer scalar. The corresponding **actual argument** shall not be an optional dummy argument.

**Result Characteristics.** The result is of the same type and type parameters as MOLD.

*Case (i):* If MOLD is a scalar and SIZE is absent, the result is a scalar.

*Case (ii):* If MOLD is an array and SIZE is absent, the result is an array and of **rank** one. Its size is as small as possible such that its physical representation is not shorter than that of SOURCE.

*Case (iii):* If SIZE is present, the result is an array of **rank** one and size SIZE.

**Result Value.** If the physical representation of the result has the same length as that of SOURCE, the physical representation of the result is that of SOURCE. If the physical representation of the result is longer than that of SOURCE, the physical representation of the leading part is that of SOURCE and the remainder is processor dependent. If the physical representation of the result is shorter than that of SOURCE, the physical representation of the result is the leading part of SOURCE. If D and E are scalar variables such that the physical representation of D is as long as or longer than that of E, the value of TRANSFER (TRANSFER (E, D), E) shall be the value of E. IF D is an array and E is an array of **rank** one, the value of TRANSFER (TRANSFER (E, D), E, SIZE (E)) shall be the value of E.

**Examples.**

*Case (i):* TRANSFER (1082130432, 0.0) has the value 4.0 on a processor that represents the values 4.0 and 1082130432 as the string of binary digits 0100 0000 1000 0000 0000 0000 0000 0000.

*Case (ii):* TRANSFER ([1.1, 2.2, 3.3], [(0.0, 0.0)]) is a complex rank-one array of length two whose first element has the value (1.1, 2.2) and whose second element has a real part with the value 3.3. The imaginary part of the second element is processor dependent.

*Case (iii):* TRANSFER ([1.1, 2.2, 3.3], [(0.0, 0.0)], 1) is a complex rank-one array of length one whose only element has the value (1.1, 2.2).

### 13.9.194 TRANSPOSE (MATRIX)

**Description.** Transpose of an array of **rank** two.

**Class.** Transformational function.

**Argument.** MATRIX shall be a rank-two array of any type.

**Result Characteristics.** The result is an array of the same type and type parameters as MATRIX and with **rank** two and shape  $[n, m]$  where  $[m, n]$  is the shape of MATRIX.

**Result Value.** Element  $(i, j)$  of the result has the value MATRIX  $(j + \text{LBOUND}(\text{MATRIX}, 1) - 1, i + \text{LBOUND}(\text{MATRIX}, 2) - 1)$ .

**Example.** If A is the array  $\begin{bmatrix} 1 & 2 & 3 \\ 4 & 5 & 6 \\ 7 & 8 & 9 \end{bmatrix}$ , then TRANSPOSE (A) has the value  $\begin{bmatrix} 1 & 4 & 7 \\ 2 & 5 & 8 \\ 3 & 6 & 9 \end{bmatrix}$ .

### 13.9.195 TRIM (STRING)

**1 Description.** String without trailing blanks.

**2 Class.** Transformational function.

**3 Argument.** STRING shall be a character scalar.

**4 Result Characteristics.** Character with the same kind type parameter value as STRING and with a length that is the length of STRING less the number of trailing blanks in STRING. If STRING contains no nonblank characters, the result has zero length.

**5 Result Value.** The value of the result is the same as STRING except any trailing blanks are removed.

**6 Example.** TRIM ( ' A B ' ) has the value ' A B '.

### 13.9.196 UBOUND (ARRAY [, DIM, KIND])

**1 Description.** Upper bound(s).

**2 Class.** Inquiry function.

**3 Arguments.**

ARRAY shall be assumed-rank or an array. It shall not be an unallocated allocatable array or a pointer that is not associated. If ARRAY is an assumed-size array, DIM shall be present with a value less than the rank of ARRAY.

DIM (optional) shall be an integer scalar with a value in the range  $1 \leq \text{DIM} \leq n$ , where  $n$  is the rank of ARRAY. The corresponding actual argument shall not be an optional dummy argument, a disassociated pointer, or an unallocated allocatable.

KIND (optional) shall be a scalar integer constant expression.

**4 Result Characteristics.** Integer. If KIND is present, the kind type parameter is that specified by the value of KIND; otherwise the kind type parameter is that of default integer type. The result is scalar if DIM is present; otherwise, the result is an array of rank one and size  $n$ , where  $n$  is the rank of ARRAY.

**5 Result Value.**

*Case (i):* If DIM is present, ARRAY is a whole array, and dimension DIM of ARRAY has nonzero extent, the result has a value equal to the upper bound for subscript DIM of ARRAY. Otherwise, if DIM is present and ARRAY is assumed-rank, the value of the result is as if ARRAY were a whole array, with the extent of the final dimension of ARRAY when ARRAY is associated with an assumed-size array being considered to be  $-1$ . Otherwise, if DIM is present, the result has a value equal to the number of elements in dimension DIM of ARRAY.

*Case (ii):* If ARRAY has rank zero, UBOUND (ARRAY) has a value that is a zero-sized array. Otherwise, UBOUND (ARRAY) has a value whose  $i^{\text{th}}$  element is equal to UBOUND (ARRAY,  $i$ ), for  $i = 1, 2, \dots, n$ , where  $n$  is the rank of ARRAY. UBOUND (ARRAY, KIND=KIND) has a value whose  $i^{\text{th}}$  element is equal to UBOUND (ARRAY,  $i$ , KIND=KIND), for  $i = 1, 2, \dots, n$ , where  $n$  is the rank of ARRAY.

**6 Examples.** If A is declared by the statement

REAL A (2:3, 7:10)

then UBOUND (A) is [3, 10] and UBOUND (A, DIM = 2) is 10.

#### NOTE 13.27

If ARRAY is assumed-rank and has rank zero, DIM cannot be present since it cannot satisfy the requirement  $1 \leq \text{DIM} \leq 0$ .

### 13.9.197 UCBOUND (COARRAY [, DIM, KIND])

**1 Description.** Upper `cobound`(s) of a `coarray`.

**2 Class.** Inquiry function.

**3 Arguments.**

COARRAY shall be a `coarray` of any type. It may be a scalar or an array. If it is `allocatable` it shall be allocated.

DIM (optional) shall be an integer scalar with a value in the range  $1 \leq \text{DIM} \leq n$ , where  $n$  is the `corank` of COARRAY. The corresponding `actual argument` shall not be an optional dummy argument, a disassociated pointer, or an unallocated allocatable.

KIND (optional) shall be a scalar integer `constant expression`.

**4 Result Characteristics.** Integer. If KIND is present, the kind type parameter is that specified by the value of KIND; otherwise, the kind type parameter is that of default integer type. The result is scalar if DIM is present; otherwise, the result is an array of `rank` one and size  $n$ , where  $n$  is the `corank` of COARRAY.

**5 Result Value.** The final upper `cobound` is the final `cosubscript` in the `cosubscript` list for the `coarray` that selects the `image` whose `index` is equal to the number of in the `team current` when COARRAY was established.

*Case (i):* If DIM is present, the result has a value equal to the upper `cobound` for `codimension` DIM of COARRAY.

*Case (ii):* If DIM is absent, the result has a value whose  $i^{\text{th}}$  element is equal to the upper `cobound` for `codimension`  $i$  of COARRAY, for  $i = 1, 2, \dots, n$ , where  $n$  is the `corank` of COARRAY.

**6 Examples.** If `NUM_IMAGES( )` has the value 30 and A is allocated by the statement

```
ALLOCATE (A [2:3, 0:7, *])
```

then `UCBOUND (A)` is [3, 7, 2] and `UCBOUND (A, DIM=2)` is 7. Note that the `cosubscripts` [3, 7, 2] do not correspond to an actual `image`.

### 13.9.198 UNPACK (VECTOR, MASK, FIELD)

**1 Description.** Vector unpacked into an array.

**2 Class.** Transformational function.

**3 Arguments.**

VECTOR shall be a rank-one array of any type. Its size shall be at least  $t$  where  $t$  is the number of true elements in MASK.

MASK shall be a logical array.

FIELD shall be of the same type and type parameters as VECTOR and shall be `conformable` with MASK.

**4 Result Characteristics.** The result is an array of the same type and type parameters as VECTOR and the same shape as MASK.

**5 Result Value.** The element of the result that corresponds to the  $i^{\text{th}}$  true element of MASK, in array element order, has the value VECTOR ( $i$ ) for  $i = 1, 2, \dots, t$ , where  $t$  is the number of true values in MASK. Each other element has a value equal to FIELD if FIELD is scalar or to the corresponding element of FIELD if it is an array.

**6 Examples.** Particular values can be “scattered” to particular positions in an array by using UNPACK. If M is the

array  $\begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}$ , V is the array [1, 2, 3], and Q is the logical mask  $\begin{bmatrix} . & \text{T} & . \\ \text{T} & . & . \\ . & . & \text{T} \end{bmatrix}$ , where “T” represents true

and “.” represents false, then the result of `UNPACK (V, MASK = Q, FIELD = M)` has the value  $\begin{bmatrix} 1 & 2 & 0 \\ 1 & 1 & 0 \\ 0 & 0 & 3 \end{bmatrix}$

and the result of UNPACK (V, MASK = Q, FIELD = 0) has the value  $\begin{bmatrix} 0 & 2 & 0 \\ 1 & 0 & 0 \\ 0 & 0 & 3 \end{bmatrix}$ .

### 13.9.199 VERIFY (STRING, SET [, BACK, KIND])

**1 Description.** Character set non-membership search.

**2 Class.** [Elemental](#) function.

**3 Arguments.**

STRING shall be of type character.

SET shall be of type character with the same kind type parameter as STRING.

BACK (optional) shall be of type logical.

KIND (optional) shall be a scalar integer [constant expression](#).

**4 Result Characteristics.** Integer. If KIND is present, the kind type parameter is that specified by the value of KIND; otherwise the kind type parameter is that of default integer type.

**5 Result Value.**

*Case (i):* If BACK is absent or has the value false and if STRING contains at least one character that is not in SET, the value of the result is the position of the leftmost character of STRING that is not in SET.

*Case (ii):* If BACK is present with the value true and if STRING contains at least one character that is not in SET, the value of the result is the position of the rightmost character of STRING that is not in SET.

*Case (iii):* The value of the result is zero if each character in STRING is in SET or if STRING has zero length.

**6 Examples.**

*Case (i):* VERIFY ('ABBA', 'A') has the value 2.

*Case (ii):* VERIFY ('ABBA', 'A', BACK = [.TRUE.](#)) has the value 3.

*Case (iii):* VERIFY ('ABBA', 'AB') has the value 0.

## 13.10 Standard intrinsic modules

### 13.10.1 General

1 This part of ISO/IEC 1539 defines five standard intrinsic modules: a Fortran environment module, a set of three modules to support floating-point exceptions and IEEE arithmetic, and a module to support interoperability with the C programming language.

2 The intrinsic modules IEEE\_EXCEPTIONS, IEEE\_ARITHMETIC, and IEEE\_FEATURES are described in Clause 14. The intrinsic module [ISO\\_C\\_BINDING](#) is described in Clause 15. The [module procedures](#) described in 13.10.2 are [pure](#).

#### NOTE 13.28

The types and procedures defined in [standard intrinsic](#) modules are not themselves [intrinsic](#).

3 A processor may extend the [standard intrinsic](#) modules to provide public entities in them in addition to those specified in this part of ISO/IEC 1539.

## 13.10.2 The ISO\_FORTRAN\_ENV intrinsic module

### 13.10.2.1 General

- 1 The intrinsic module ISO\_FORTRAN\_ENV provides public entities relating to the Fortran environment.
- 2 The processor shall provide the [named constants](#), derived types, and procedures described in subclause 13.10.2. In the detailed descriptions below, procedure names are generic and not specific.

### 13.10.2.2 ATOMIC\_INT\_KIND

- 1 The value of the default integer scalar constant ATOMIC\_INT\_KIND is the kind type parameter value of type integer variables for which the processor supports atomic operations specified by [atomic subroutines](#).

### 13.10.2.3 ATOMIC\_LOGICAL\_KIND

- 1 The value of the default integer scalar constant ATOMIC\_LOGICAL\_KIND is the kind type parameter value of type logical variables for which the processor supports atomic operations specified by [atomic subroutines](#).

### 13.10.2.4 CHARACTER\_KINDS

- 1 The values of the elements of the default integer array constant CHARACTER\_KINDS are the kind values supported by the processor for variables of type character. The order of the values is processor dependent. The rank of the array is one, its lower bound is one, and its size is the number of character kinds supported.

### 13.10.2.5 CHARACTER\_STORAGE\_SIZE

- 1 The value of the default integer scalar constant CHARACTER\_STORAGE\_SIZE is the size expressed in bits of the [character storage unit](#) (16.5.3.2).

### 13.10.2.6 COMPILER\_OPTIONS ( )

- 1 **Description.** Processor-dependent string describing the options that controlled the program translation phase.
- 2 **Class.** [Transformational function](#).
- 3 **Argument.** None.
- 4 **Result Characteristics.** Default character scalar with processor-dependent length.
- 5 **Result Value.** A processor-dependent value which describes the options that controlled the translation phase of program execution. This value should include relevant information that could be useful for diagnosing problems at a later date.
- 6 **Example.** COMPILER\_OPTIONS ( ) might have the value '/OPTIMIZE /FLOAT=IEEE'.

### 13.10.2.7 COMPILER\_VERSION ( )

- 1 **Description.** Processor-dependent string identifying the program translation phase.
- 2 **Class.** [Transformational function](#).
- 3 **Argument.** None.
- 4 **Result Characteristics.** Default character scalar with processor-dependent length.

1 5 **Result Value.** A processor-dependent value that identifies the name and version of the program translation  
 2 phase of the processor. This value should include relevant information that could be useful for diagnosing problems  
 3 at a later date.

4 6 **Example.** COMPILER\_VERSION ( ) might have the value 'Fast KL-10 Compiler Version 7'.

#### NOTE 13.29

Relevant information that could be useful for diagnosing problems at a later date might include compiler release and patch level, default compiler arguments, environment variable values, and run time library requirements. A processor might include this information in an object file automatically, without the user needing to save the result of this function in a variable.

### 5 13.10.2.8 CURRENT\_TEAM

6 1 The value of the default integer scalar constant CURRENT\_TEAM identifies the [current team](#) when it is used  
 7 as the LEVEL argument to [GET\\_TEAM](#).

### 8 13.10.2.9 ERROR\_UNIT

9 1 The value of the default integer scalar constant ERROR\_UNIT identifies the processor-dependent [preconnected](#)  
 10 [external unit](#) used for the purpose of error reporting (9.5). This [unit](#) may be the same as OUTPUT\_UNIT. The  
 11 value shall not be -1.

### 12 13.10.2.10 EVENT\_TYPE

13 1 EVENT\_TYPE is a derived type with private components. It is an extensible type with no type parameters.  
 14 Each nonallocatable component is fully default-initialized.

15 2 A scalar variable of type EVENT\_TYPE is an [event variable](#). The value of an event variable includes its event  
 16 count, which is updated by execution of a sequence of [EVENT POST](#) or [EVENT WAIT](#) statements. The effect  
 17 of each change is as if the intrinsic subroutine [ATOMIC\\_ADD](#) were executed with a variable that stores the  
 18 event count as its ATOM argument. A coarray that is of type EVENT\_TYPE may be referenced or defined  
 19 during execution of a segment that is unordered relative to the execution of another segment in which that  
 20 coarray is defined. The event count is of type integer with kind [ATOMIC\\_INT\\_KIND](#) from the intrinsic module  
 21 [ISO\\_FORTRAN\\_ENV](#). The initial value of the event count of an event variable is zero.

22 C1303 A named variable of type EVENT\_TYPE shall be a coarray. A named variable with a noncoarray  
 23 subcomponent of type EVENT\_TYPE shall be a coarray.

24 C1304 An [event variable](#) shall not appear in a variable definition context except as the *event-variable* in an  
 25 [EVENT POST](#) or [EVENT WAIT](#) statement, as an *allocate-object*, or as an [actual argument](#) in a reference  
 26 to a procedure with an [explicit interface](#) if the corresponding [dummy argument](#) has [INTENT \(INOUT\)](#).

27 C1305 A variable with a nonpointer subobject of type EVENT\_TYPE shall not appear in a variable definition  
 28 context except as an *allocate-object* in an [ALLOCATE statement](#) without a [SOURCE= specifier](#), as an  
 29 *allocate-object* in a [DEALLOCATE statement](#), or as an [actual argument](#) in a reference to a procedure  
 30 with an [explicit interface](#) if the corresponding [dummy argument](#) has [INTENT \(INOUT\)](#).

#### NOTE 13.30

The restrictions against changing an [event variable](#) except via [EVENT POST](#) and [EVENT WAIT](#) statements ensure the integrity of its value and facilitate efficient implementation, particularly when special synchronization is needed for correct event handling.

**NOTE 13.31**

Updates to variables via atomic subroutines are coherent but not necessarily consistent, so a processor might have to use extra synchronization to obtain the consistency required for the segments ordered by [EVENT POST](#) and [EVENT WAIT](#) statements.

**13.10.2.11 FILE\_STORAGE\_SIZE**

- 1 The value of the default integer scalar constant `FILE_STORAGE_SIZE` is the size expressed in bits of the [file storage unit](#) (9.3.5).

**13.10.2.12 INITIAL\_TEAM**

- 1 The value of the default integer scalar constant `INITIAL_TEAM` identifies the [initial team](#) when it is used as the `LEVEL` argument to [GET\\_TEAM](#).

**13.10.2.13 INPUT\_UNIT**

- 1 The value of the default integer scalar constant `INPUT_UNIT` identifies the same processor-dependent [external unit preconnected](#) for sequential formatted input as the one identified by an asterisk in a [READ statement](#); this [unit](#) is the one used for a [READ statement](#) that does not contain an input/output control list (9.6.4.3). The value shall not be `-1`.

**13.10.2.14 INT8, INT16, INT32, and INT64**

- 1 The values of these default integer scalar constants shall be those of the [kind type parameters](#) that specify an `INTEGER` type whose storage size expressed in bits is 8, 16, 32, and 64 respectively. If, for any of these constants, the processor supports more than one kind of that size, it is processor dependent which kind value is provided. If the processor supports no kind of a particular size, that constant shall be equal to `-2` if the processor supports a kind with larger size and `-1` otherwise.

**13.10.2.15 INTEGER\_KINDS**

- 1 The values of the elements of the default integer array constant `INTEGER_KINDS` are the kind values supported by the processor for variables of type integer. The order of the values is processor dependent. The rank of the array is one, its lower bound is one, and its size is the number of integer kinds supported.

**13.10.2.16 IOSTAT\_END**

- 1 The value of the default integer scalar constant `IOSTAT_END` is assigned to the variable specified in an [IOSTAT=specifier](#) (9.11.5) if an end-of-file condition occurs during execution of an [input statement](#) and no error condition occurs. This value shall be negative.

**13.10.2.17 IOSTAT\_EOR**

- 1 The value of the default integer scalar constant `IOSTAT_EOR` is assigned to the variable specified in an [IOSTAT=specifier](#) (9.11.5) if an end-of-record condition occurs during execution of an [input statement](#) and no end-of-file or error condition occurs. This value shall be negative and different from the value of `IOSTAT_END`.

**13.10.2.18 IOSTAT\_INQUIRE\_INTERNAL\_UNIT**

- 1 The value of the default integer scalar constant `IOSTAT_INQUIRE_INTERNAL_UNIT` is assigned to the variable specified in an [IOSTAT=specifier](#) in an [INQUIRE statement](#) (9.10) if a [file-unit-number](#) identifies an [internal unit](#) in that statement.

**NOTE 13.32**

This can only occur when a [defined input/output](#) procedure is called by the processor as the result of executing a parent [data transfer statement](#) (9.6.4.8.2) for an [internal unit](#).

**13.10.2.19 LOCK\_TYPE**

LOCK\_TYPE is a [derived type](#) with private components; no component is allocatable or a pointer. It is an [extensible type](#) with no [type parameters](#). Therefore it does not have the [BIND attribute](#) and is not a [sequence type](#). All components have default initialization.

A scalar variable of type LOCK\_TYPE is a [lock variable](#). A lock variable can have one of two states: locked and unlocked. The unlocked state is represented by the one value that is the default value of a LOCK\_TYPE variable; this is the value specified by the [structure constructor](#) LOCK\_TYPE ( ). The locked state is represented by all other values. The value of a lock variable can be changed with the [LOCK](#) and [UNLOCK](#) statements (8.6.10).

C1306 A named variable of type LOCK\_TYPE shall be a [coarray](#). A named variable with a noncoarray [subcomponent](#) of type LOCK\_TYPE shall be a [coarray](#).

C1307 A lock variable shall not appear in a variable definition context except as the [lock-variable](#) in a [LOCK](#) or [UNLOCK](#) statement, as an [allocate-object](#), or as an [actual argument](#) in a reference to a procedure with an [explicit interface](#) where the corresponding [dummy argument](#) has [INTENT \(INOUT\)](#).

C1308 A variable with a subobject of type LOCK\_TYPE shall not appear in a variable definition context except as an [allocate-object](#) or as an [actual argument](#) in a reference to a procedure with an [explicit interface](#) where the corresponding [dummy argument](#) has [INTENT \(INOUT\)](#).

**NOTE 13.33**

The restrictions against changing a lock variable except via the [LOCK](#) and [UNLOCK](#) statements ensure the integrity of its value and facilitate efficient implementation, particularly when special synchronization is needed for correct lock operation.

**13.10.2.20 LOGICAL\_KINDS**

The values of the elements of the default integer array constant LOGICAL\_KINDS are the kind values supported by the processor for variables of type logical. The order of the values is processor dependent. The rank of the array is one, its lower bound is one, and its size is the number of logical kinds supported.

**13.10.2.21 NUMERIC\_STORAGE\_SIZE**

The value of the default integer scalar constant NUMERIC\_STORAGE\_SIZE is the size expressed in bits of the [numeric storage unit](#) (16.5.3.2).

**13.10.2.22 OUTPUT\_UNIT**

The value of the default integer scalar constant OUTPUT\_UNIT identifies the same processor-dependent [external unit preconnected](#) for sequential formatted output as the one identified by an asterisk in a [WRITE statement](#) (9.6.4.3). The value shall not be  $-1$ .

**13.10.2.23 PARENT\_TEAM**

The value of the default integer scalar constant PARENT\_TEAM identifies the [parent team](#) when it is used as the LEVEL argument to [GET\\_TEAM](#).

#### 13.10.2.24 REAL\_KINDS

- 1 The values of the elements of the default integer array constant `REAL_KINDS` are the kind values supported by the processor for variables of type real. The order of the values is processor dependent. The rank of the array is one, its lower bound is one, and its size is the number of real kinds supported.

#### 13.10.2.25 REAL32, REAL64, and REAL128

- 1 The values of these default integer scalar [named constants](#) shall be those of the kind type parameters that specify a REAL type whose storage size expressed in bits is 32, 64, and 128 respectively. If, for any of these constants, the processor supports more than one kind of that size, it is processor dependent which kind value is provided. If the processor supports no kind of a particular size, that constant shall be equal to `-2` if the processor supports kinds of a larger size and `-1` otherwise.

#### 13.10.2.26 STAT\_FAILED\_IMAGE

- 1 If the processor has the ability to detect that an image has failed, the value of the default integer scalar constant `STAT_FAILED_IMAGE` is positive; otherwise, the value of `STAT_FAILED_IMAGE` is negative. If an image involved in execution of an [image control statement](#), a reference to a coindexed object, or execution of a [collective](#) or [atomic](#) subroutine has failed, and no other error condition occurs, the value of `STAT_FAILED_IMAGE` is assigned to the variable specified in a `STAT= specifier` in the execution of an [image control statement](#) or reference to a coindexed object, or to the `STAT` argument in an invocation of a [collective](#) or [atomic](#).
- 2 If the `STAT= specifier` of an execution of a [CHANGE TEAM](#), [END TEAM](#), [FORM TEAM](#), [SYNC ALL](#), [SYNC IMAGES](#), or [SYNC TEAM](#) statement is assigned the value `STAT_FAILED_IMAGE`, the effect on all the active images involved in that statement shall be the same as for successful execution, except for the value assigned to the `STAT= specifier`.

#### 13.10.2.27 STAT\_LOCKED

- 1 The value of the default integer scalar constant `STAT_LOCKED` is assigned to the variable specified in a `STAT= specifier` (8.6.11) of a [LOCK statement](#) if the [lock variable](#) is locked by the executing [image](#).

#### 13.10.2.28 STAT\_LOCKED\_OTHER\_IMAGE

- 1 The value of the default integer scalar constant `STAT_LOCKED_OTHER_IMAGE` is assigned to the variable specified in a `STAT= specifier` (8.6.11) of an [UNLOCK statement](#) if the [lock variable](#) is locked by another [image](#).

#### 13.10.2.29 STAT\_STOPPED\_IMAGE

- 1 The value of the default integer scalar constant `STAT_STOPPED_IMAGE` is assigned to the variable specified in a `STAT= specifier` (6.7.4, 8.6.11), or a `STAT` argument in a reference to a [collective subroutine](#), if execution of the statement with that specifier or argument requires synchronization with an image that has initiated [normal termination](#). This value shall be positive.

#### 13.10.2.30 STAT\_UNLOCKED

- 1 The value of the default integer scalar constant `STAT_UNLOCKED` is assigned to the variable specified in a `STAT= specifier` (8.6.11) of an [UNLOCK statement](#) if the [lock variable](#) is unlocked.

#### 13.10.2.31 TEAM\_TYPE

- 1 `TEAM_TYPE` is a [derived type](#) with private components; no component is a pointer. It is an [extensible type](#) with no [type parameters](#). Therefore it does not have the [BIND attribute](#) and is not a [sequence type](#). Each component is fully default-initialized.

1 2 A scalar variable of type TEAM\_\_TYPE is a [team variable](#), and can identify a team. The default initial value of  
2 a team variable does not identify any team.

3 **13.10.2.32 Uniqueness of named constant values**

4 1 The values of the named constants [IOSTAT\\_\\_INQUIRE\\_\\_INTERNAL\\_\\_UNIT](#), [STAT\\_\\_FAILED\\_\\_IMAGE](#), [STAT\\_\\_](#)  
5 [LOCKED](#), [STAT\\_\\_LOCKED\\_\\_OTHER\\_\\_IMAGE](#), [STAT\\_\\_STOPPED\\_\\_IMAGE](#), and [STAT\\_\\_UNLOCKED](#) shall be  
6 distinct.

## 14 Exceptions and IEEE arithmetic

### 14.1 Overview of IEEE arithmetic support

- 1 The intrinsic modules IEEE\_EXCEPTIONS, IEEE\_ARITHMETIC, and IEEE\_FEATURES provide support for the facilities defined by ISO/IEC/IEEE 60559:2011\*. Whether the modules are provided is processor dependent. If the module IEEE\_FEATURES is provided, which of the named constants defined in this part of ISO/IEC 1539 are included is processor dependent. The module IEEE\_ARITHMETIC behaves as if it contained a [USE statement](#) for IEEE\_EXCEPTIONS; everything that is public in IEEE\_EXCEPTIONS is public in IEEE\_ARITHMETIC.

#### NOTE 14.1

The types and procedures defined in these modules are not themselves intrinsic.

- 2 If IEEE\_EXCEPTIONS or IEEE\_ARITHMETIC is accessible in a [scoping unit](#), the exceptions IEEE\_OVERFLOW and IEEE\_DIVIDE\_BY\_ZERO are supported in the [scoping unit](#) for all kinds of real and complex IEEE floating-point data. Which other exceptions are supported can be determined by the [inquiry function](#) IEEE\_SUPPORT\_FLAG (14.11.45); whether control of [halting](#) is supported can be determined by the [inquiry function](#) IEEE\_SUPPORT\_HALTING. The extent of support of the other exceptions may be influenced by the accessibility of the [named constants](#) IEEE\_INEXACT\_FLAG, IEEE\_INVALID\_FLAG, and IEEE\_UNDERFLOW\_FLAG of the module IEEE\_FEATURES. If a [scoping unit](#) has access to IEEE\_UNDERFLOW\_FLAG of IEEE\_FEATURES, within the [scoping unit](#) the processor shall support underflow and return true from IEEE\_SUPPORT\_FLAG (IEEE\_UNDERFLOW, X) for at least one kind of real. Similarly, if IEEE\_INEXACT\_FLAG or IEEE\_INVALID\_FLAG is accessible, within the [scoping unit](#) the processor shall support the exception and return true from the corresponding [inquiry function](#) for at least one kind of real. If IEEE\_HALTING is accessible, within the [scoping unit](#) the processor shall support control of [halting](#) and return true from IEEE\_SUPPORT\_HALTING (FLAG) for the flag.

#### NOTE 14.2

IEEE\_INVALID is not required to be supported whenever IEEE\_EXCEPTIONS is accessed. This is to allow a processor whose arithmetic does not conform to ISO/IEC/IEEE 60559:2011 to provide support for overflow and divide\_by\_zero. On a processor which does support ISO/IEC/IEEE 60559:2011, invalid is an equally serious condition.

- 3 If a [scoping unit](#) does not access IEEE\_FEATURES, IEEE\_EXCEPTIONS, or IEEE\_ARITHMETIC, the level of support is processor dependent, and need not include support for any exceptions. If a flag is signaling on entry to such a [scoping unit](#), the processor ensures that it is signaling on exit. If a flag is quiet on entry to such a [scoping unit](#), whether it is signaling on exit is processor dependent.
- 4 Additional ISO/IEC/IEEE 60559:2011 facilities are available from the module IEEE\_ARITHMETIC. The extent of support may be influenced by the accessibility of the [named constants](#) of the module IEEE\_FEATURES. If a [scoping unit](#) has access to IEEE\_DATATYPE of IEEE\_FEATURES, within the [scoping unit](#) the processor shall support IEEE arithmetic and return true from IEEE\_SUPPORT\_DATATYPE (X) (14.11.42) for at least one kind of real. Similarly, if IEEE\_DENORMAL, IEEE\_DIVIDE, IEEE\_INF, IEEE\_NAN, IEEE\_ROUNDING, IEEE\_SQRT, or IEEE\_SUBNORMAL is accessible, within the [scoping unit](#) the processor shall support the feature and return true from the corresponding [inquiry function](#) for at least one kind of real. In the case of IEEE\_ROUNDING, it shall return true for the rounding modes IEEE\_NEAREST, IEEE\_TO\_ZERO, IEEE\_UP, and IEEE\_DOWN; support for [IEEE\\_AWAY](#) is also required if there is at least one kind of real X for

\* Because ISO/IEC/IEEE 60559:2011 was originally an IEEE standard, its facilities are widely known as “IEEE arithmetic”, and this terminology is used by this part of ISO/IEC 1539.

which `IEEE_SUPPORT_DATATYPE` (X) is true and `RADIX` (X) is equal to ten. Note that the effect of `IEEE_DENORMAL` is the same as that of `IEEE_SUBNORMAL`.

Execution might be slowed on some processors by the support of some features. If `IEEE_EXCEPTIONS` or `IEEE_ARITHMETIC` is accessed but `IEEE_FEATURES` is not accessed, the supported subset of features is processor dependent. The processor's fullest support is provided when all of `IEEE_FEATURES` is accessed as in

```
USE, INTRINSIC :: IEEE_ARITHMETIC; USE, INTRINSIC :: IEEE_FEATURES
```

but execution might then be slowed by the presence of a feature that is not needed. In all cases, the extent of support can be determined by the [inquiry functions](#).

## 14.2 Derived types and constants defined in the modules

The modules `IEEE_EXCEPTIONS`, `IEEE_ARITHMETIC`, and `IEEE_FEATURES` define five derived types, whose components are all private. No [direct component](#) of any of these types is [allocatable](#) or a pointer.

The module `IEEE_EXCEPTIONS` defines the following types.

- `IEEE_FLAG_TYPE` is for identifying a particular exception flag. Its only possible values are those of [named constants](#) defined in the module: `IEEE_INVALID`, `IEEE_OVERFLOW`, `IEEE_DIVIDE_BY_ZERO`, `IEEE_UNDERFLOW`, and `IEEE_INEXACT`. The module also defines the array [named constants](#) `IEEE_USUAL` = [ `IEEE_OVERFLOW`, `IEEE_DIVIDE_BY_ZERO`, `IEEE_INVALID` ] and `IEEE_ALL` = [ `IEEE_USUAL`, `IEEE_UNDERFLOW`, `IEEE_INEXACT` ].
- `IEEE_MODES_TYPE` is for representing the floating-point modes.
- `IEEE_STATUS_TYPE` is for representing the floating-point status.

The module `IEEE_ARITHMETIC` defines the following.

- The type `IEEE_CLASS_TYPE`, for identifying a class of floating-point values. Its only possible values are those of [named constants](#) defined in the module: `IEEE_SIGNALING_NAN`, `IEEE_QUIET_NAN`, `IEEE_NEGATIVE_INF`, `IEEE_NEGATIVE_NORMAL`, `IEEE_NEGATIVE_DENORMAL`, `IEEE_NEGATIVE_ZERO`, `IEEE_POSITIVE_ZERO`, `IEEE_POSITIVE_SUBNORMAL`, `IEEE_POSITIVE_NORMAL`, `IEEE_POSITIVE_INF`, and `IEEE_OTHER_VALUE`. The named constants `IEEE_NEGATIVE_DENORMAL` and `IEEE_POSITIVE_DENORMAL` are defined with the same value as [IEEE\\_NEGATIVE\\_SUBNORMAL](#) and [IEEE\\_POSITIVE\\_SUBNORMAL](#) respectively.
- The type `IEEE_ROUND_TYPE`, for identifying a particular rounding mode. Its only possible values are those of [named constants](#) defined in the module: `IEEE_NEAREST`, `IEEE_TO_ZERO`, `IEEE_UP`, `IEEE_DOWN`, and [IEEE\\_AWAY](#) for the rounding modes specified in ISO/IEC/IEEE 60559:2011, and `IEEE_OTHER` for any other mode.
- The [elemental operator](#) `==` for two values of one of these types to return true if the values are the same and false otherwise.
- The [elemental operator](#) `/=` for two values of one of these types to return true if the values differ and false otherwise.

The module `IEEE_FEATURES` defines the type `IEEE_FEATURES_TYPE`, for expressing the need for particular ISO/IEC/IEEE 60559:2011 features. Its only possible values are those of [named constants](#) defined in the module: `IEEE_DATATYPE`, `IEEE_DENORMAL`, `IEEE_DIVIDE`, `IEEE_HALTING`, `IEEE_INEXACT_FLAG`, `IEEE_INF`, `IEEE_INVALID_FLAG`, `IEEE_NAN`, `IEEE_ROUNDING`, `IEEE_SQRT`, `IEEE_SUBNORMAL`, and `IEEE_UNDERFLOW_FLAG`.

## 14.3 The exceptions

The exceptions are the following.

- IEEE\_OVERFLOW occurs in an intrinsic real addition, subtraction, multiplication, division, or conversion by the intrinsic function `REAL`, as specified by ISO/IEC/IEEE 60559:2011 if `IEEE_SUPPORT_DATA_TYPE` is true for the operands of the operation or conversion, and as determined by the processor otherwise. It occurs in an intrinsic real exponentiation as determined by the processor. It occurs in a complex operation, or conversion by the intrinsic function `CMPLX`, if it is caused by the calculation of the real or imaginary part of the result.
- IEEE\_DIVIDE\_BY\_ZERO occurs in a real division as specified by ISO/IEC/IEEE 60559:2011 if `IEEE_SUPPORT_DATATYPE` is true for the operands of the division, and as determined by the processor otherwise. It is processor-dependent whether it occurs in a real exponentiation with a negative exponent. It occurs in a complex division if it is caused by the calculation of the real or imaginary part of the result.
- IEEE\_INVALID occurs when a real or complex operation or assignment is invalid; possible examples are `SQRT(X)` when `X` is real and has a nonzero negative value, and conversion to an integer (by assignment, an intrinsic procedure, or a procedure defined in an intrinsic module) when the result is too large to be representable. In a numeric relational operation  $x_1 \text{ rel-op } x_2$ , if  $x_1 + x_2$  is of type real and `IEEE_SUPPORT_NAN` ( $x_1 + x_2$ ) is true, IEEE\_INVALID shall signal as specified by ISO/IEC/IEEE 60559:2011 for the compareSignaling{relation} operations; that is, if  $x_1$  and  $x_2$  are unordered.
- IEEE\_UNDERFLOW occurs when the result for an intrinsic real operation or assignment has an absolute value less than a processor-dependent limit and loss of accuracy is detected, or the real or imaginary part of the result for an intrinsic complex operation or assignment has an absolute value less than a processor-dependent limit and loss of accuracy is detected.
- IEEE\_INEXACT occurs when the result of a real or complex operation or assignment is not exact.

2 Each exception has a flag whose value is either quiet or signaling. The value can be determined by the subroutine `IEEE_GET_FLAG`. Its initial value is quiet and it signals when the associated exception occurs. Its status can also be changed by the subroutine `IEEE_SET_FLAG` or the subroutine `IEEE_SET_STATUS`. Once signaling within a procedure, it remains signaling unless set quiet by an invocation of the subroutine `IEEE_SET_FLAG` or the subroutine `IEEE_SET_STATUS`.

3 If a flag is signaling on entry to a procedure other than `IEEE_GET_FLAG` or `IEEE_GET_STATUS`, the processor will set it to quiet on entry and restore it to signaling on return. If a flag signals during execution of a procedure, the processor shall not set it to quiet on return.

4 Evaluation of a `specification expression` might cause an exception to signal.

5 In a `scoping unit` that has access to `IEEE_EXCEPTIONS` or `IEEE_ARITHMETIC`, if an intrinsic procedure or a procedure defined in an intrinsic module executes normally, the values of the flags `IEEE_OVERFLOW`, `IEEE_DIVIDE_BY_ZERO`, and `IEEE_INVALID` shall be as on entry to the procedure, even if one or more of them signals during the calculation. If a real or complex result is too large for the procedure to handle, `IEEE_OVERFLOW` may signal. If a real or complex result is a NaN because of an invalid operation (for example, `LOG(-1.0)`), `IEEE_INVALID` may signal. Similar rules apply to format processing and to intrinsic operations: no signaling flag shall be set quiet and no quiet flag shall be set signaling because of an intermediate calculation that does not affect the result.

6 In a sequence of statements that has no invocations of `IEEE_GET_FLAG`, `IEEE_SET_FLAG`, `IEEE_GET_STATUS`, `IEEE_SET_HALTING_MODE`, or `IEEE_SET_STATUS`, if the execution of an operation would cause an exception to signal but after execution of the sequence no value of a variable depends on the operation, whether the exception is signaling is processor dependent. For example, when `Y` has the value zero, whether the code

```
X = 1.0/Y
X = 3.0
```

signals `IEEE_DIVIDE_BY_ZERO` is processor dependent. Another example is the following:

```
REAL, PARAMETER :: X=0.0, Y=6.0
IF (1.0/X == Y) PRINT *, 'Hello world'
```

where the processor is permitted to discard the `IF` statement because the logical expression can never be true and no value of a variable depends on it.

- 7 An exception shall not signal if this could arise only during execution of an operation beyond those required or permitted by the standard. For example, the statement

`IF (F (X) > 0.0) Y = 1.0/Z`

shall not signal `IEEE_DIVIDE_BY_ZERO` when both `F (X)` and `Z` are zero and the statement

`WHERE (A > 0.0) A = 1.0/A`

shall not signal `IEEE_DIVIDE_BY_ZERO`. On the other hand, when `X` has the value 1.0 and `Y` has the value 0.0, the expression

`X>0.00001 .OR. X/Y>0.00001`

is permitted to cause the signaling of `IEEE_DIVIDE_BY_ZERO`.

- 8 The processor need not support `IEEE_INVALID`, `IEEE_UNDERFLOW`, and `IEEE_INEXACT`. If an exception is not supported, its flag is always quiet. The inquiry function `IEEE_SUPPORT_FLAG` can be used to inquire whether a particular flag is supported.

## 14.4 The rounding modes

- 1 ISO/IEC/IEEE 60559:2011 specifies a binary rounding mode that affects floating-point arithmetic with radix two, and a decimal rounding mode that affects floating-point arithmetic with radix ten. Unqualified references to the rounding mode with respect to a particular arithmetic operation or operands refers to the mode for the radix of the operation or operands, and other unqualified references to the rounding mode refers to both binary and decimal rounding modes.

- 2 ISO/IEC/IEEE 60559:2011 specifies five possible modes for rounding:

- `IEEE_NEAREST` rounds the exact result to the nearest representable value;
- `IEEE_TO_ZERO` rounds the exact result towards zero to the next representable value;
- `IEEE_UP` rounds the exact result towards  $+\infty$  to the next representable value;
- `IEEE_DOWN` rounds the exact result towards  $-\infty$  to the next representable value;
- `IEEE_AWAY` rounds the exact result away from zero to the next representable value; ISO/IEC/IEEE 60559:2011 requires this mode for decimal floating-point, but it is optional for binary floating-point.

- 3 The subroutine `IEEE_GET_ROUNDING_MODE` can be used to get the rounding modes. The initial rounding modes are processor dependent.

- 4 If the processor supports the alteration of the rounding modes during execution, the subroutine `IEEE_SET_ROUNDING_MODE` can be used to alter them. The inquiry function `IEEE_SUPPORT_ROUNDING` can be used to inquire whether this facility is available for a particular mode. The inquiry function `IEEE_SUPPORT_IO` can be used to inquire whether rounding for base conversion in formatted input/output (9.5.6.16, 9.6.2.13, 10.7.2.3.8) is as specified in ISO/IEC/IEEE 60559:2011.

- 5 In a procedure other than `IEEE_SET_ROUNDING_MODE` or `IEEE_SET_STATUS`, the processor shall not change the rounding modes on entry, and on return shall ensure that the rounding modes are the same as they were on entry.

### NOTE 14.3

Within a program, all literal constants that have the same form have the same value (4.1.4). Therefore, the value of a literal constant is not affected by the rounding modes.

## 14.5 Underflow mode

- 1 Some processors allow control during program execution of whether underflow produces a subnormal number in conformance with ISO/IEC/IEEE 60559:2011 (gradual underflow) or produces zero instead (abrupt underflow). On some processors, floating-point performance is typically better in abrupt underflow mode than in gradual underflow mode.
- 2 Control over the underflow mode is exercised by invocation of `IEEE_SET_UNDERFLOW_MODE`. The subroutine `IEEE_GET_UNDERFLOW_MODE` can be used to get the underflow mode. The [inquiry function](#) `IEEE_SUPPORT_UNDERFLOW_CONTROL` can be used to inquire whether this facility is available. The initial underflow mode is processor dependent. In a procedure other than `IEEE_SET_UNDERFLOW_MODE` or `IEEE_SET_STATUS`, the processor shall not change the underflow mode on entry, and on return shall ensure that the underflow mode is the same as it was on entry.
- 3 The underflow mode affects only floating-point calculations whose type is that of an X for which `IEEE_SUPPORT_UNDERFLOW_CONTROL` returns true.

## 14.6 Halting

- 1 Some processors allow control during program execution of whether to abort or continue execution after an exception. Such control is exercised by invocation of the subroutine `IEEE_SET_HALTING_MODE`. Halting is not precise and may occur any time after the exception has occurred. The [inquiry function](#) `IEEE_SUPPORT_HALTING` can be used to inquire whether this facility is available. The initial halting mode is processor dependent. In a procedure other than `IEEE_SET_HALTING_MODE` or `IEEE_SET_STATUS`, the processor shall not change the halting mode on entry, and on return shall ensure that the halting mode is the same as it was on entry.

## 14.7 The floating-point modes and status

- 1 The values of the rounding modes, [underflow mode](#), and [halting mode](#) are collectively called the floating-point modes. The values of all the supported flags for exceptions and the floating-point modes are collectively called the floating-point status. The floating-point modes can be stored in a scalar variable of type `IEEE_MODES_TYPE` with the subroutine `IEEE_GET_MODES` and restored with the subroutine `IEEE_SET_MODES`. The floating-point status can be stored in a scalar variable of type `IEEE_STATUS_TYPE` with the subroutine `IEEE_GET_STATUS` and restored with the subroutine `IEEE_SET_STATUS`. There are no facilities for finding the values of particular flags represented by such a variable.

### NOTE 14.4

Each [image](#) has its own floating-point status (2.3.4).

### NOTE 14.5

Some processors hold all these flags and modes in one or two status registers that can be obtained and set as a whole faster than all individual flags and modes can be obtained and set. These procedures are provided to exploit this feature.

### NOTE 14.6

The processor is required to ensure that a call to a Fortran procedure does not change the floating-point status other than by setting exception flags to signaling.

## 14.8 Exceptional values

- 1 ISO/IEC/IEEE 60559:2011 specifies the following exceptional floating-point values.

- Subnormal values have very small absolute values and reduced precision.
- Infinite values (+infinity and -infinity) are created by overflow or division by zero.
- Not-a-Number ( NaN) values are undefined values or values created by an invalid operation.

2 A value that does not fall into the above classes is called a normal number.

3 The functions IEEE\_IS\_FINITE, IEEE\_IS\_NAN, IEEE\_IS\_NEGATIVE, and IEEE\_IS\_NORMAL are provided to test whether a value is finite, NaN, negative, or normal. The function IEEE\_VALUE is provided to generate an IEEE number of any class, including an infinity or a NaN. The inquiry functions IEEE\_SUPPORT\_SUBNORMAL, IEEE\_SUPPORT\_INF, and IEEE\_SUPPORT\_NAN are provided to determine whether these facilities are available for a particular kind of real.

## 14.9 IEEE arithmetic

1 The inquiry function IEEE\_SUPPORT\_DATATYPE can be used to inquire whether IEEE arithmetic is supported for a particular kind of real. Complete conformance with ISO/IEC/IEEE 60559:2011 is not required, but

- the normal numbers shall be exactly those of an ISO/IEC/IEEE 60559:2011 floating-point format,
- for at least one rounding mode, the intrinsic operations of addition, subtraction and multiplication shall conform whenever the operands and result specified by ISO/IEC/IEEE 60559:2011 are normal numbers,
- the IEEE function abs shall be provided by the intrinsic function ABS,
- the IEEE operation rem shall be provided by the function IEEE\_REM, and
- the IEEE functions copysign, scalb, logb, nextafter, and unordered shall be provided by the functions IEEE\_COPY\_SIGN, IEEE\_SCALB, IEEE\_LOGB, IEEE\_NEXT\_AFTER, and IEEE\_UNORDERED, respectively,

for that kind of real.

2 The inquiry function IEEE\_SUPPORT\_NAN is provided to inquire whether the processor supports IEEE NaNs. Where these are supported, the result of the intrinsic operations +, −, and \*, and the functions IEEE\_REM and IEEE\_RINT from the intrinsic module IEEE\_ARITHMETIC, shall conform to ISO/IEC/IEEE 60559:2011 when the result is an IEEE NaN.

3 The inquiry function IEEE\_SUPPORT\_INF is provided to inquire whether the processor supports IEEE infinities. Where these are supported, the result of the intrinsic operations +, −, and \*, and the functions IEEE\_REM and IEEE\_RINT from the intrinsic module IEEE\_ARITHMETIC, shall conform to ISO/IEC/IEEE 60559:2011 when exactly one operand or the result specified by ISO/IEC/IEEE 60559:2011 is an IEEE infinity.

4 The inquiry function IEEE\_SUPPORT\_SUBNORMAL is provided to inquire whether the processor supports subnormal numbers. Where these are supported, the result of the intrinsic operations +, −, and \*, and the functions IEEE\_REM and IEEE\_RINT from the intrinsic module IEEE\_ARITHMETIC, shall conform to ISO/IEC/IEEE 60559:2011 when the result specified by ISO/IEC/IEEE 60559:2011 is subnormal, or any operand is subnormal and either the result is not an IEEE infinity or IEEE\_SUPPORT\_INF is true.

5 The inquiry function IEEE\_SUPPORT\_DIVIDE is provided to inquire whether, on kinds of real for which IEEE\_SUPPORT\_DATATYPE returns true, the intrinsic division operation conforms to ISO/IEC/IEEE 60559:2011 when both operands and the result specified by ISO/IEC/IEEE 60559:2011 are normal numbers. If IEEE\_SUPPORT\_NAN is also true for a particular kind of real, the intrinsic division operation on that kind conforms to ISO/IEC/IEEE 60559:2011 when the result specified by ISO/IEC/IEEE 60559:2011 is a NaN. If IEEE\_SUPPORT\_INF is also true for a particular kind of real, the intrinsic division operation on that kind conforms to ISO/IEC/IEEE 60559:2011 when one operand or the result specified by ISO/IEC/IEEE 60559:2011 is an IEEE infinity. If IEEE\_SUPPORT\_SUBNORMAL is also true for a particular kind of real, the intrinsic division operation on that kind conforms to ISO/IEC/IEEE 60559:2011 when the result specified by ISO/IEC/IEEE 60559:2011 is subnormal, or when any operand is subnormal and either the result specified by ISO/IEC/IEEE 60559:2011 is not an infinity or IEEE\_SUPPORT\_INF is true.

- 6 ISO/IEC/IEEE 60559:2011 specifies a square root function that returns negative real zero for the square root of negative real zero and has certain accuracy requirements. The [inquiry function](#) IEEE\_SUPPORT\_SQRT can be used to inquire whether the intrinsic function [SQRT](#) conforms to ISO/IEC/IEEE 60559:2011 for a particular kind of real. If IEEE\_SUPPORT\_NAN is also true for a particular kind of real, the intrinsic function [SQRT](#) on that kind conforms to ISO/IEC/IEEE 60559:2011 when the result specified by ISO/IEC/IEEE 60559:2011 is a NaN. If IEEE\_SUPPORT\_INF is also true for a particular kind of real, the intrinsic function [SQRT](#) on that kind conforms to ISO/IEC/IEEE 60559:2011 when the result specified by ISO/IEC/IEEE 60559:2011 is an IEEE infinity. If [IEEE\\_SUPPORT\\_SUBNORMAL](#) is also true for a particular kind of real, the intrinsic function [SQRT](#) on that kind conforms to ISO/IEC/IEEE 60559:2011 when the argument is subnormal.
- 7 The [inquiry function](#) IEEE\_SUPPORT\_STANDARD is provided to inquire whether the processor supports all the ISO/IEC/IEEE 60559:2011 facilities defined in this part of ISO/IEC 1539 for a particular kind of real.

## 14.10 Summary of the procedures

- 1 For all of the procedures defined in the modules, the arguments shown are the names that shall be used for [argument keywords](#) if the keyword form is used for the [actual arguments](#).
- 2 A procedure classified in 14.10 as an [inquiry function](#) depends on the properties of one or more of its arguments instead of their values; in fact, these argument values may be undefined. Unless the description of one of these [inquiry functions](#) states otherwise, these arguments are permitted to be unallocated [allocatable](#) variables or pointers that are undefined or [disassociated](#). A procedure that is classified as a [transformational function](#) is neither an [inquiry function](#) nor [elemental](#).
- 3 In the Class column of Tables 14.1 and 14.2,
- E indicates that the procedure is an [elemental](#) function,
  - ES indicates that the procedure is an [elemental](#) subroutine,
  - I indicates that the procedure is an [inquiry function](#),
  - PS indicates that the procedure is a [pure](#) subroutine,
  - S indicates that the procedure is an impure subroutine, and
  - T indicates that the procedure is a [transformational function](#).

Table 14.1: **IEEE\_ARITHMETIC** module procedure summary

| Procedure               | Arguments                  | Class | Description                                      |
|-------------------------|----------------------------|-------|--------------------------------------------------|
| IEEE_CLASS              | (X)                        | E     | Classify number.                                 |
| IEEE_COPY_SIGN          | (X, Y)                     | E     | Copy sign.                                       |
| IEEE_FMA                | (A, B, C)                  | E     | Fused multiply-add operation.                    |
| IEEE_GET_ROUNDING_MODE  | (ROUND_VALUE<br>[, RADIX]) | S     | Get rounding mode.                               |
| IEEE_GET_UNDERFLOW_MODE | (GRADUAL)                  | S     | Get <a href="#">underflow mode</a> .             |
| IEEE_INT                | (A, ROUND [, KIND])        | E     | Conversion to integer type.                      |
| IEEE_IS_FINITE          | (X)                        | E     | Whether a value is finite.                       |
| IEEE_IS_NAN             | (X)                        | E     | Whether a value is an <a href="#">IEEE NaN</a> . |
| IEEE_IS_NEGATIVE        | (X)                        | E     | Whether a value is negative.                     |
| IEEE_IS_NORMAL          | (X)                        | E     | Whether a value is a normal number.              |
| IEEE_LOGB               | (X)                        | E     | Exponent.                                        |
| IEEE_MAX_NUM            | (X, Y)                     | E     | Maximum numeric value.                           |
| IEEE_MAX_NUM_MAG        | (X, Y)                     | E     | Maximum magnitude numeric value.                 |
| IEEE_MIN_NUM            | (X, Y)                     | E     | Minimum numeric value.                           |
| IEEE_MIN_NUM_MAG        | (X, Y)                     | E     | Minimum magnitude numeric value.                 |
| IEEE_NEXT_AFTER         | (X, Y)                     | E     | Adjacent machine number.                         |
| IEEE_NEXT_DOWN          | (X)                        | E     | Adjacent lower machine number.                   |
| IEEE_NEXT_UP            | (X)                        | E     | Adjacent higher machine number.                  |

Table 14.1: IEEE\_ARITHMETIC module procedure summary (cont.)

| Procedure                           | Arguments                  | Class | Description                                      |
|-------------------------------------|----------------------------|-------|--------------------------------------------------|
| IEEE_QUIET_EQ                       | (A, B)                     | E     | Quiet compares equal.                            |
| IEEE_QUIET_GE                       | (A, B)                     | E     | Quiet compares greater than or equal.            |
| IEEE_QUIET_GT                       | (A, B)                     | E     | Quiet compares greater than.                     |
| IEEE_QUIET_LE                       | (A, B)                     | E     | Quiet compares less than or equal.               |
| IEEE_QUIET_LT                       | (A, B)                     | E     | Quiet compares less than.                        |
| IEEE_QUIET_NE                       | (A, B)                     | E     | Quiet compares not equal.                        |
| IEEE_REAL                           | (A [, KIND])               | E     | Conversion to real type.                         |
| IEEE_REM                            | (X, Y)                     | E     | Exact remainder.                                 |
| IEEE_RINT                           | (X)                        | E     | Round to integer.                                |
| IEEE_SCALB                          | (X, I)                     | E     | $X \times 2^I$ .                                 |
| IEEE_SELECTED_REAL_KIND             | ([P, R, RADIX])            | T     | IEEE kind type parameter value.                  |
| IEEE_SET_ROUNDING_MODE              | (ROUND_VALUE<br>[, RADIX]) | S     | Set rounding mode.                               |
| IEEE_SET_UNDERFLOW_<br>MODE         | (GRADUAL)                  | S     | Set <a href="#">underflow mode</a> .             |
| IEEE_SIGNBIT                        | (X)                        | E     | Test sign bit.                                   |
| IEEE_SUPPORT_DATATYPE               | ([X])                      | I     | Query IEEE arithmetic support.                   |
| IEEE_SUPPORT_DENORMAL               | ([X])                      | I     | Query subnormal number support.                  |
| IEEE_SUPPORT_DIVIDE                 | ([X])                      | I     | Query IEEE division support.                     |
| IEEE_SUPPORT_INF                    | ([X])                      | I     | Query IEEE infinity support.                     |
| IEEE_SUPPORT_IO                     | ([X])                      | I     | Query IEEE formatting support.                   |
| IEEE_SUPPORT_NAN                    | ([X])                      | I     | Query <a href="#">IEEE NaN</a> support.          |
| IEEE_SUPPORT_ROUNDING               | (ROUND_VALUE<br>[, X])     | I     | Query IEEE rounding support.                     |
| IEEE_SUPPORT_SQRT                   | ([X])                      | I     | Query IEEE square root support.                  |
| IEEE_SUPPORT_SUBNORMAL              | ([X])                      | I     | Query subnormal number support.                  |
| IEEE_SUPPORT_STANDARD               | ([X])                      | I     | Query IEEE standard support.                     |
| IEEE_SUPPORT_UNDER-<br>FLOW_CONTROL | ([X])                      | I     | Query <a href="#">underflow</a> control support. |
| IEEE_UNORDERED                      | (X, Y)                     | E     | Whether two values are unordered.                |
| IEEE_VALUE                          | (X, CLASS)                 | E     | Return number in a class.                        |

Table 14.2: IEEE\_EXCEPTIONS module procedure summary

| Procedure             | Arguments          | Class | Description                                 |
|-----------------------|--------------------|-------|---------------------------------------------|
| IEEE_GET_FLAG         | (FLAG, FLAG_VALUE) | ES    | Get an exception flag.                      |
| IEEE_GET_HALTING_MODE | (FLAG, HALTING)    | ES    | Get a <a href="#">halting mode</a> .        |
| IEEE_GET_MODES        | (MODES)            | S     | Get floating-point modes.                   |
| IEEE_GET_STATUS       | (STATUS_VALUE)     | S     | Get floating-point status.                  |
| IEEE_SET_FLAG         | (FLAG, FLAG_VALUE) | PS    | Set an exception flag.                      |
| IEEE_SET_HALTING_MODE | (FLAG, HALTING)    | PS    | Set a <a href="#">halting mode</a> .        |
| IEEE_SET_MODES        | (MODES)            | S     | Set floating-point modes.                   |
| IEEE_SET_STATUS       | (STATUS_VALUE)     | S     | Restore floating-point status.              |
| IEEE_SUPPORT_FLAG     | (FLAG [, X])       | I     | Query exception support.                    |
| IEEE_SUPPORT_HALTING  | (FLAG)             | I     | Query <a href="#">halting mode</a> support. |

1 4 In the intrinsic module IEEE\_ARITHMETIC, the [elemental](#) functions listed are provided for all reals X and Y.

## 14.11 Specifications of the procedures

### 14.11.1 General

- 1 In the detailed descriptions in 14.11, procedure names are generic and are not specific. All the functions are pure and all the subroutines are impure unless otherwise stated. All dummy arguments have **INTENT (IN)** if the intent is not stated explicitly. In the examples, it is assumed that the processor supports IEEE arithmetic for default real.
- 2 For the **elemental** functions of IEEE\_ARITHMETIC that return a floating-point result, if X or Y has a value that is an infinity or a NaN, the result shall be consistent with the general rules in 6.1 and 6.2 of ISO/IEC/IEEE 60559:2011. For example, the result for an infinity shall be constructed as the limiting case of the result with a value of arbitrarily large magnitude, if such a limit exists.
- 3 A program may contain statements that, if executed, would violate the requirements listed in a **Restriction** paragraph.

#### NOTE 14.7

A program can avoid violating those requirements by using **IF constructs** to check whether particular features are supported. For example,

```
IF (IEEE_SUPPORT_DATATYPE (X)) THEN
  C = IEEE_CLASS (X)
ELSE
  ...
END IF
```

avoids invoking **IEEE\_CLASS** except on a processor which supports that facility.

### 14.11.2 IEEE\_CLASS (X)

- 1 **Description.** Classify number.
- 2 **Class.** **Elemental** function.
- 3 **Argument.** X shall be of type real.
- 4 **Restriction.** IEEE\_CLASS (X) shall not be invoked if IEEE\_SUPPORT\_DATATYPE (X) has the value false.
- 5 **Result Characteristics.** **IEEE\_CLASS\_TYPE**.
- 6 **Result Value.** The result value shall be IEEE\_SIGNALING\_NAN or IEEE\_QUIET\_NAN if IEEE\_SUPPORT\_NAN (X) has the value true and the value of X is a signaling or quiet NaN, respectively. The result value shall be IEEE\_NEGATIVE\_INF or IEEE\_POSITIVE\_INF if IEEE\_SUPPORT\_INF (X) has the value true and the value of X is negative or positive infinity, respectively. The result value shall be **IEEE\_NEGATIVE\_SUBNORMAL** or **IEEE\_POSITIVE\_SUBNORMAL** if **IEEE\_SUPPORT\_SUBNORMAL** (X) has the value true and the value of X is a negative or positive subnormal value, respectively. The result value shall be IEEE\_NEGATIVE\_NORMAL, IEEE\_NEGATIVE\_ZERO, IEEE\_POSITIVE\_ZERO, or IEEE\_POSITIVE\_NORMAL if the value of X is negative normal, negative zero, positive zero, or positive normal, respectively. Otherwise, the result value shall be IEEE\_OTHER\_VALUE.
- 7 **Example.** IEEE\_CLASS (−1.0) has the value IEEE\_NEGATIVE\_NORMAL.

#### NOTE 14.8

The result value IEEE\_OTHER\_VALUE is useful on systems that are almost IEEE-compatible, but do not implement all of it. For example, if a subnormal value is encountered on a system that does not support

## NOTE 14.8 (cont.)

|       |
|-------|
| them. |
|-------|

**14.11.3 IEEE\_COPY\_SIGN (X, Y)****1 Description.** Copy sign.**2 Class.** [Elemental](#) function.**3 Arguments.** The arguments shall be of type real.**4 Restriction.** IEEE\_COPY\_SIGN (X, Y) shall not be invoked if IEEE\_SUPPORT\_DATATYPE (X) or IEEE\_SUPPORT\_DATATYPE (Y) has the value false.**5 Result Characteristics.** Same as X.**6 Result Value.** The result has the absolute value of X with the sign of Y. This is true even for IEEE special values, such as a [NaN](#) or an infinity (on processors supporting such values).**7 Example.** The value of IEEE\_COPY\_SIGN (X, 1.0) is [ABS](#) (X) even when X is a [NaN](#).**14.11.4 IEEE\_FMA (A, B, C)****1 Description.** Fused multiply-add operation.**2 Class.** [Elemental](#) function.**3 Arguments.**

A shall be of type real.

B shall be of the same type and kind type parameter as A.

C shall be of the same type and kind type parameter as A.

**4 Restriction.** IEEE\_FMA (A, B, C) shall not be invoked if [IEEE\\_SUPPORT\\_DATATYPE](#) (A) has the value false.**5 Result Characteristics.** Same as A.**6 Result Value.** The result has the value specified by ISO/IEC/IEEE 60559:2011 for the fusedMultiplyAdd operation; that is, when the result is in range, its value is equal to the mathematical value of  $(A \times B) + C$  rounded to the representation method of A according to the rounding mode. IEEE\_OVERFLOW, IEEE\_UNDERFLOW, and IEEE\_INEXACT shall be signaled according to the final step in the calculation and not by any intermediate calculation.**7 Example.** The value of IEEE\_FMA ([TINY](#) (0.0), [TINY](#) (0.0), 1.0), when the rounding mode is IEEE\_NEAREST, is equal to 1.0; only the IEEE\_INEXACT exception is signaled.**14.11.5 IEEE\_GET\_FLAG (FLAG, FLAG\_VALUE)****1 Description.** Get an exception flag.**2 Class.** [Elemental](#) subroutine.**3 Arguments.**FLAG shall be of type [IEEE\\_FLAG\\_TYPE](#). It specifies the exception flag to be obtained.FLAG\_VALUE shall be of type logical. It is an [INTENT \(OUT\)](#) argument. If the value of FLAG is IEEE\_INVALID, IEEE\_OVERFLOW, IEEE\_DIVIDE\_BY\_ZERO, IEEE\_UNDERFLOW, or IEEE\_INEXACT, FLAG\_VALUE is assigned the value true if the corresponding exception flag is signaling

and is assigned the value false otherwise.

- Example.** Following CALL IEEE\_GET\_FLAG (IEEE\_OVERFLOW, FLAG\_VALUE), FLAG\_VALUE is true if the IEEE\_OVERFLOW flag is signaling and is false if it is quiet.

#### 14.11.6 IEEE\_GET\_HALTING\_MODE (FLAG, HALTING)

- Description.** Get a [halting mode](#).

- Class.** [Elemental](#) subroutine.

- Arguments.**

FLAG shall be of type [IEEE\\_FLAG\\_TYPE](#). It specifies the exception flag. It shall have one of the values IEEE\_INVALID, IEEE\_OVERFLOW, IEEE\_DIVIDE\_BY\_ZERO, IEEE\_UNDERFLOW, or IEEE\_INEXACT.

HALTING shall be of type logical. It is an [INTENT \(OUT\)](#) argument. It is assigned the value true if the exception specified by FLAG will cause halting. Otherwise, it is assigned the value false.

- Example.** To store the [halting mode](#) for IEEE\_OVERFLOW, do a calculation without halting, and restore the [halting mode](#) later:

```
USE, INTRINSIC :: IEEE_ARITHMETIC
LOGICAL HALTING
...
CALL IEEE_GET_HALTING_MODE (IEEE_OVERFLOW, HALTING) ! Store halting mode
CALL IEEE_SET_HALTING_MODE (IEEE_OVERFLOW, .FALSE.) ! No halting
... ! calculation without halting
CALL IEEE_SET_HALTING_MODE (IEEE_OVERFLOW, HALTING) ! Restore halting mode
```

#### 14.11.7 IEEE\_GET\_MODES (MODES)

- Description.** Get floating-point modes.

- Class.** Subroutine.

- Argument.** MODES shall be a scalar of type [IEEE\\_MODES\\_TYPE](#). It is an [INTENT \(OUT\)](#) argument that is assigned the value of the floating-point modes.

- Example.** To save the floating-point modes, do a calculation with specific rounding and [underflow](#) modes, and restore them later:

```
USE, INTRINSIC :: IEEE_ARITHMETIC
TYPE (IEEE_MODES_TYPE) SAVE_MODES
...
CALL IEEE_GET_MODES (SAVE_MODES) ! Save all modes.
CALL IEEE_SET_ROUNDING_MODE (IEEE_TO_ZERO)
CALL IEEE_SET_UNDERFLOW_MODE (GRADUAL=.FALSE.)
... ! calculation with abrupt round-to-zero.
CALL IEEE_SET_MODES (SAVE_MODES) ! Restore all modes.
```

### 14.11.8 IEEE\_GET\_ROUNDING\_MODE (ROUND\_VALUE [, RADIX])

1 **Description.** Get rounding mode.

2 **Class.** Subroutine.

3 **Arguments.**

ROUND\_VALUE shall be a scalar of type `IEEE_ROUND_TYPE`. It is an `INTENT (OUT)` argument. It is assigned the value `IEEE_NEAREST`, `IEEE_TO_ZERO`, `IEEE_UP`, `IEEE_DOWN`, or `IEEE_AWAY` if the corresponding rounding mode is in operation and `IEEE_OTHER` otherwise.

RADIX (optional) shall be an integer scalar with the value two or ten. If RADIX is present with the value ten, the rounding mode queried is the decimal rounding mode, otherwise it is the binary rounding mode.

4 **Example.** To save the binary rounding mode, do a calculation with round to nearest, and restore the rounding mode later:

```
USE, INTRINSIC :: IEEE_ARITHMETIC
TYPE (IEEE_ROUND_TYPE) ROUND_VALUE
...
CALL IEEE_GET_ROUNDING_MODE (ROUND_VALUE) ! Store the rounding mode
CALL IEEE_SET_ROUNDING_MODE (IEEE_NEAREST)
... ! calculation with round to nearest
CALL IEEE_SET_ROUNDING_MODE (ROUND_VALUE) ! Restore the rounding mode
```

### 14.11.9 IEEE\_GET\_STATUS (STATUS\_VALUE)

1 **Description.** Get floating-point status.

2 **Class.** Subroutine.

3 **Argument.** STATUS\_VALUE shall be a scalar of type `IEEE_STATUS_TYPE`. It is an `INTENT (OUT)` argument. It is assigned the value of the floating-point status.

4 **Example.** To store all the exception flags, do a calculation involving exception handling, and restore them later:

```
USE, INTRINSIC :: IEEE_ARITHMETIC
TYPE (IEEE_STATUS_TYPE) STATUS_VALUE
...
CALL IEEE_GET_STATUS (STATUS_VALUE) ! Get the flags
CALL IEEE_SET_FLAG (IEEE_ALL, .FALSE.) ! Set the flags quiet.
... ! calculation involving exception handling
CALL IEEE_SET_STATUS (STATUS_VALUE) ! Restore the flags
```

### 14.11.10 IEEE\_GET\_UNDERFLOW\_MODE (GRADUAL)

1 **Description.** Get `underflow mode`.

2 **Class.** Subroutine.

3 **Argument.** GRADUAL shall be a logical scalar. It is an `INTENT (OUT)` argument. It is assigned the value true if the `underflow mode` is gradual underflow, and false if the `underflow mode` is abrupt underflow.

4 **Restriction.** `IEEE_GET_UNDERFLOW_MODE` shall not be invoked unless `IEEE_SUPPORT_UNDERFLOW_CONTROL (X)` is true for some X.

1 5 **Example.** After CALL [IEEE\\_SET\\_UNDERFLOW\\_MODE](#) ([.FALSE.](#)), a subsequent CALL [IEEE\\_GET\\_UNDERFLOW\\_MODE](#) ([GRADUAL](#)) will set [GRADUAL](#) to false.

### 3 14.11.11 [IEEE\\_INT](#) (A, ROUND [, KIND])

4 1 **Description.** Conversion to integer type.

5 2 **Class.** [Elemental](#) function.

6 3 **Arguments.**

7 A shall be of type real.

8 ROUND shall be of type [IEEE\\_ROUND\\_TYPE](#).

9 KIND (optional) shall be a scalar integer constant expression.

10 4 **Restriction.** [IEEE\\_INT](#) (A, ROUND, KIND) shall not be invoked if [IEEE\\_SUPPORT\\_DATATYPE](#) (A) has the value false.

12 5 **Result Characteristics.** Integer. If KIND is present, the kind type parameter is that specified by the value of KIND; otherwise, the kind type parameter is that of default integer.

14 6 **Result Value.** The result has the value specified by ISO/IEC/IEEE 60559:2011 for the `convertToInteger{round}` or the `convertToIntegerExact{round}` operation; the processor shall consistently choose which operation it provides. That is, the value of A is converted to an integer according to the rounding mode specified by ROUND; if this value is representable in the representation method of the result, the result has this value, otherwise [IEEE\\_INVALID](#) is signaled and the result is processor dependent. If the processor provides the `convertToIntegerExact` operation, [IEEE\\_INVALID](#) did not signal, and the value of the result differs from that of A, [IEEE\\_INEXACT](#) will be signaled.

21 7 **Example.** The value of [IEEE\\_INT](#) (12.5, [IEEE\\_UP](#)) is 13; [IEEE\\_INEXACT](#) will be signaled if the processor provides the `convertToIntegerExact` operation.

### 23 14.11.12 [IEEE\\_IS\\_FINITE](#) (X)

24 1 **Description.** Whether a value is finite.

25 2 **Class.** [Elemental](#) function.

26 3 **Argument.** X shall be of type real.

27 4 **Restriction.** [IEEE\\_IS\\_FINITE](#) (X) shall not be invoked if [IEEE\\_SUPPORT\\_DATATYPE](#) (X) has the value false.

29 5 **Result Characteristics.** Default logical.

30 6 **Result Value.** The result has the value true if the value of X is finite, that is, [IEEE\\_CLASS](#) (X) has one of the values [IEEE\\_NEGATIVE\\_NORMAL](#), [IEEE\\_NEGATIVE\\_SUBNORMAL](#), [IEEE\\_NEGATIVE\\_ZERO](#), [IEEE\\_POSITIVE\\_ZERO](#), [IEEE\\_POSITIVE\\_SUBNORMAL](#), or [IEEE\\_POSITIVE\\_NORMAL](#); otherwise, the result has the value false.

34 7 **Example.** [IEEE\\_IS\\_FINITE](#) (1.0) has the value true.

### 35 14.11.13 [IEEE\\_IS\\_NAN](#) (X)

36 1 **Description.** Whether a value is an [IEEE NaN](#).

37 2 **Class.** [Elemental](#) function.

38 3 **Argument.** X shall be of type real.

1 4 **Restriction.** IEEE\_IS\_NAN (X) shall not be invoked if IEEE\_SUPPORT\_NAN (X) has the value false.

2 5 **Result Characteristics.** Default logical.

3 6 **Result Value.** The result has the value true if the value of X is an IEEE NaN; otherwise, it has the value false.

4 7 **Example.** IEEE\_IS\_NAN (SQRT (−1.0)) has the value true if IEEE\_SUPPORT\_SQRT (1.0) has the value  
5 true.

#### 6 14.11.14 IEEE\_IS\_NEGATIVE (X)

7 1 **Description.** Whether a value is negative.

8 2 **Class.** Elemental function.

9 3 **Argument.** X shall be of type real.

10 4 **Restriction.** IEEE\_IS\_NEGATIVE (X) shall not be invoked if IEEE\_SUPPORT\_DATATYPE (X) has the  
11 value false.

12 5 **Result Characteristics.** Default logical.

13 6 **Result Value.** The result has the value true if IEEE\_CLASS (X) has one of the values IEEE\_NEGATIVE\_-  
14 NORMAL, IEEE\_NEGATIVE\_SUBNORMAL, IEEE\_NEGATIVE\_ZERO or IEEE\_NEGATIVE\_INF; oth-  
15 erwise, the result has the value false.

16 7 **Example.** IEEE\_IS\_NEGATIVE (0.0) has the value false.

#### 17 14.11.15 IEEE\_IS\_NORMAL (X)

18 1 **Description.** Whether a value is a normal number.

19 2 **Class.** Elemental function.

20 3 **Argument.** X shall be of type real.

21 4 **Restriction.** IEEE\_IS\_NORMAL (X) shall not be invoked if IEEE\_SUPPORT\_DATATYPE (X) has the  
22 value false.

23 5 **Result Characteristics.** Default logical.

24 6 **Result Value.** The result has the value true if IEEE\_CLASS (X) has one of the values IEEE\_NEGATIVE\_-  
25 NORMAL, IEEE\_NEGATIVE\_ZERO, IEEE\_POSITIVE\_ZERO or IEEE\_POSITIVE\_NORMAL; otherwise,  
26 the result has the value false.

27 7 **Example.** IEEE\_IS\_NORMAL (SQRT (−1.0)) has the value false if IEEE\_SUPPORT\_SQRT (1.0) has the  
28 value true.

#### 29 14.11.16 IEEE\_LOGB (X)

30 1 **Description.** Exponent.

31 2 **Class.** Elemental function.

32 3 **Argument.** X shall be of type real.

33 4 **Restriction.** IEEE\_LOGB (X) shall not be invoked if IEEE\_SUPPORT\_DATATYPE (X) has the value false.

34 5 **Result Characteristics.** Same as X.

35 6 **Result Value.**

- 1 *Case (i):* If the value of X is neither zero, infinity, nor NaN, the result has the value of the unbiased exponent  
2 of X. Note: this value is equal to **EXPONENT** (X) – 1.
- 3 *Case (ii):* If X==0, the result is –infinity if **IEEE\_SUPPORT\_INF** (X) is true and –**HUGE** (X) otherwise;  
4 **IEEE\_DIVIDE\_BY\_ZERO** signals.
- 5 *Case (iii):* If **IEEE\_SUPPORT\_INF** (X) is true and X is infinite, the result is +infinity.
- 6 *Case (iv):* If **IEEE\_SUPPORT\_NAN** (X) is true and X is a NaN, the result is a NaN.

7 7 **Example.** **IEEE\_LOGB** (–1.1) has the value 0.0.

#### 8 14.11.17 **IEEE\_MAX\_NUM** (X, Y)

9 1 **Description.** Maximum numeric value.

10 2 **Class.** **Elemental** function.

11 3 **Arguments.**

12 X shall be of type real.

13 Y shall be of the same type and kind type parameter as X.

14 4 **Restriction.** **IEEE\_MAX\_NUM** shall not be invoked if **IEEE\_SUPPORT\_DATATYPE** (X) has the value  
15 false.

16 5 **Result Characteristics.** Same as X.

17 6 **Result Value.** The result has the value specified for the maxNum operation in ISO/IEC/IEEE 60559:2011;  
18 that is,

- 19 • if  $X < Y$  the result has the value of Y;
- 20 • if  $Y < X$  the result has the value of X;
- 21 • if exactly one of X and Y is a quiet NaN the result has the value of the other argument;
- 22 • if both X and Y are quiet NaNs the result is either X or Y (processor dependent);
- 23 • if one or both of X and Y are signaling NaNs, **IEEE\_INVALID** signals and the result is a NaN.

24 7 Except when X or Y is a signaling NaN, no exception is signaled.

25 8 **Example.** The value of **IEEE\_MAX\_NUM** (1.5, **IEEE\_VALUE** (**IEEE\_QUIET\_NAN**)) is 1.5.

#### 26 14.11.18 **IEEE\_MAX\_NUM\_MAG** (X, Y)

27 1 **Description.** Maximum magnitude numeric value.

28 2 **Class.** **Elemental** function.

29 3 **Arguments.**

30 X shall be of type real.

31 Y shall be of the same type and kind type parameter as X.

32 4 **Restriction.** **IEEE\_MAX\_NUM\_MAG** shall not be invoked if **IEEE\_SUPPORT\_DATATYPE** (X) has the  
33 value false.

34 5 **Result Characteristics.** Same as X.

35 6 **Result Value.** The result has the value specified for the maxNumMag operation in ISO/IEC/IEEE 60559:2011;  
36 that is,

- 37 • if **ABS** (X) < **ABS** (Y) the result has the value of Y;
- 38 • if **ABS** (Y) < **ABS** (X) the result has the value of X;

- if exactly one of X and Y is a quiet NaN the result has the value of the other argument;
- if both X and Y are quiet NaNs the result is either X or Y (processor dependent);
- if one or both of X and Y are signaling NaNs, IEEE\_INVALID signals and the result is a NaN.

Except when X or Y is a signaling NaN, no exception is signaled.

**Example.** The value of IEEE\_MAX\_NUM\_MAG (1.5, -2.5) is -2.5.

### 14.11.19 IEEE\_MIN\_NUM (X, Y)

**Description.** Minimum numeric value.

**Class.** [Elemental](#) function.

**Arguments.**

X shall be of type real.

Y shall be of the same type and kind type parameter as X.

**Restriction.** IEEE\_MIN\_NUM shall not be invoked if [IEEE\\_SUPPORT\\_DATATYPE](#) (X) has the value false.

**Result Characteristics.** Same as X.

**Result Value.**

The result has the value specified for the minNum operation in ISO/IEC/IEEE 60559:2011; that is,

- if  $X < Y$  the result has the value of X;
- if  $Y < X$  the result has the value of Y;
- if exactly one of X and Y is a quiet NaN the result has the value of the other argument;
- if both X and Y are quiet NaNs the result is either X or Y (processor dependent);
- if one or both of X and Y are signaling NaNs, IEEE\_INVALID signals and the result is a NaN.

Except when X or Y is a signaling NaN, no exception is signaled.

**Example.** The value of IEEE\_MIN\_NUM (1.5, [IEEE\\_VALUE](#) (IEEE\_QUIET\_NAN)) is 1.5.

### 14.11.20 IEEE\_MIN\_NUM\_MAG (X, Y)

**Description.** Minimum magnitude numeric value.

**Class.** [Elemental](#) function.

**Arguments.**

X shall be of type real.

Y shall be of the same type and kind type parameter as X.

**Restriction.** IEEE\_MIN\_NUM\_MAG shall not be invoked if [IEEE\\_SUPPORT\\_DATATYPE](#) (X) has the value false.

**Result Characteristics.** Same as X.

**Result Value.** The result has the value specified for the minNumMag operation in ISO/IEC/IEEE 60559:2011; that is,

- if  $\text{ABS}(X) < \text{ABS}(Y)$  the result has the value of X;
- if  $\text{ABS}(Y) < \text{ABS}(X)$  the result has the value of Y;
- if exactly one of X and Y is a quiet NaN the result has the value of the other argument;

- if both X and Y are quiet NaNs the result is either X or Y (processor dependent);
- if one or both of X and Y are signaling NaNs, IEEE\_INVALID signals and the result is a NaN.

Except when X or Y is a signaling NaN, no exception is signaled.

**Example.** The value of IEEE\_MIN\_NUM\_MAG (1.5, -2.5) is 1.5.

### 14.11.21 IEEE\_NEXT\_AFTER (X, Y)

**Description.** Adjacent machine number.

**Class.** Elemental function.

**Arguments.** The arguments shall be of type real.

**Restriction.** IEEE\_NEXT\_AFTER (X, Y) shall not be invoked if IEEE\_SUPPORT\_DATATYPE (X) or IEEE\_SUPPORT\_DATATYPE (Y) has the value false.

**Result Characteristics.** Same as X.

**Result Value.**

*Case (i):* If  $X == Y$ , the result is X and no exception is signaled.

*Case (ii):* If  $X \neq Y$ , the result has the value of the next representable neighbor of X in the direction of Y. The neighbors of zero (of either sign) are both nonzero. IEEE\_OVERFLOW is signaled when X is finite but IEEE\_NEXT\_AFTER (X, Y) is infinite; IEEE\_UNDERFLOW is signaled when IEEE\_NEXT\_AFTER (X, Y) is subnormal; in both cases, IEEE\_INEXACT signals.

**Example.** The value of IEEE\_NEXT\_AFTER (1.0, 2.0) is  $1.0 + \text{EPSILON}(X)$ .

### 14.11.22 IEEE\_NEXT\_DOWN (X)

**Description.** Adjacent lower machine number.

**Class.** Elemental function.

**Argument.** X shall be of type real.

**Restriction.** IEEE\_NEXT\_DOWN (X) shall not be invoked if IEEE\_SUPPORT\_DATATYPE (X) has the value false. IEEE\_NEXT\_DOWN ( $-\text{HUGE}(X)$ ) shall not be invoked if IEEE\_SUPPORT\_INF (X) has the value false.

**Result Characteristics.** Same as X.

**Result Value.** The result has the value specified for the nextDown operation in ISO/IEC/IEEE 60559:2011; that is, it is the greatest value in the representation method of X that compares less than X, except when X is equal to  $-\infty$  the result has the value  $-\infty$ , and when X is a NaN the result is a NaN. If X is a signaling NaN, IEEE\_INVALID signals; otherwise, no exception is signaled.

**Example.** If IEEE\_SUPPORT\_SUBNORMAL (0.0) is true, the value of IEEE\_NEXT\_DOWN (+0.0) is the negative subnormal number with least magnitude.

### 14.11.23 IEEE\_NEXT\_UP (X)

**Description.** Adjacent higher machine number.

**Class.** Elemental function.

**Argument.** X shall be of type real.

1 4 **Restriction.** IEEE\_NEXT\_UP (X) shall not be invoked if IEEE\_SUPPORT\_DATATYPE (X) has the value  
2 false. IEEE\_NEXT\_UP (HUGE (X)) shall not be invoked if IEEE\_SUPPORT\_INF (X) has the value false.

3 5 **Result Characteristics.** Same as X.

4 6 **Result Value.** The result has the value specified for the nextUp operation in ISO/IEC/IEEE 60559:2011; that  
5 is, it is the least value in the representation method of X that compares greater than X, except when X is  
6 equal to  $+\infty$  the result has the value  $+\infty$ , and when X is a NaN the result is a NaN. If X is a signaling NaN,  
7 IEEE\_INVALID\_signals; otherwise, no exception is signaled.

8 7 **Example.** If IEEE\_SUPPORT\_INF (X) is true, the value of IEEE\_NEXT\_UP (HUGE (X)) is  $+\infty$ .

#### 9 14.11.24 IEEE\_QUIET\_EQ (A, B)

10 1 **Description.** Quiet compares equal.

11 2 **Class.** Elemental function.

12 3 **Arguments.**

13 A shall be of type real.

14 B shall have the same type and type parameters as A.

15 4 **Restriction.** IEEE\_QUIET\_EQ (A) shall not be invoked if IEEE\_SUPPORT\_DATATYPE (A) has the value  
16 false.

17 5 **Result Characteristics.** Default logical.

18 6 **Result Value.** The result has the value specified for the compareQuietEqual operation in ISO/IEC/IEEE  
19 60559:2011; that is, it is true if and only if A compares equal to B; if A or B is a NaN, the result will be false  
20 and no exception will be signaled.

21 7 **Example.** IEEE\_QUIET\_EQ (1.0, IEEE\_VALUE (IEEE\_QUIET\_NAN)) has the value false and no exception  
22 is signaled.

#### 23 14.11.25 IEEE\_QUIET\_GE (A, B)

24 1 **Description.** Quiet compares greater than or equal.

25 2 **Class.** Elemental function.

26 3 **Arguments.**

27 A shall be of type real.

28 B shall have the same type and type parameters as A.

29 4 **Restriction.** IEEE\_QUIET\_GE (A) shall not be invoked if IEEE\_SUPPORT\_DATATYPE (A) has the value  
30 false.

31 5 **Result Characteristics.** Default logical.

32 6 **Result Value.** The result has the value specified for the compareQuietGreaterEqual operation in ISO/IEC/IEEE  
33 60559:2011; that is, it is true if and only if A compares greater than or equal to B; if A or B is a NaN, the result  
34 will be false and no exception will be signaled.

35 7 **Example.** IEEE\_QUIET\_GE (1.0, IEEE\_VALUE (IEEE\_QUIET\_NAN)) has the value false and no exception  
36 is signaled.

### 14.11.26 IEEE\_QUIET\_GT (A, B)

**Description.** Quiet compares greater than.

**Class.** [Elemental](#) function.

**Arguments.**

A shall be of type real.

B shall have the same type and type parameters as A.

**Restriction.** IEEE\_QUIET\_GT (A) shall not be invoked if [IEEE\\_SUPPORT\\_DATATYPE](#) (A) has the value false.

**Result Characteristics.** Default logical.

**Result Value.** The result has the value specified for the compareQuietGreater operation in ISO/IEC/IEEE 60559:2011; that is, it is true if and only if A compares greater than B; if A or B is a NaN, the result will be false and no exception will be signaled.

**Example.** IEEE\_QUIET\_GT (1.0, [IEEE\\_VALUE](#) (IEEE\_QUIET\_NAN)) has the value false and no exception is signaled.

### 14.11.27 IEEE\_QUIET\_LE (A, B)

**Description.** Quiet compares less than or equal.

**Class.** [Elemental](#) function.

**Arguments.**

A shall be of type real.

B shall have the same type and type parameters as A.

**Restriction.** IEEE\_QUIET\_LE (A) shall not be invoked if [IEEE\\_SUPPORT\\_DATATYPE](#) (A) has the value false.

**Result Characteristics.** Default logical.

**Result Value.** The result has the value specified for the compareQuietLessEqual operation in ISO/IEC/IEEE 60559:2011; that is, it is true if and only if A compares less than or equal to B; if A or B is a NaN, the result will be false and no exception will be signaled.

**Example.** IEEE\_QUIET\_LE (1.0, [IEEE\\_VALUE](#) (IEEE\_QUIET\_NAN)) has the value false and no exception is signaled.

### 14.11.28 IEEE\_QUIET\_LT (A, B)

**Description.** Quiet compares less than.

**Class.** [Elemental](#) function.

**Arguments.**

A shall be of type real.

B shall have the same type and type parameters as A.

**Restriction.** IEEE\_QUIET\_LT (A) shall not be invoked if [IEEE\\_SUPPORT\\_DATATYPE](#) (A) has the value false.

**Result Characteristics.** Default logical.

1 6 **Result Value.** The result has the value specified for the compareQuietLess operation in ISO/IEC/IEEE  
 2 60559:2011; that is, it is true if and only if A compares less than B; if A or B is a NaN, the result will be  
 3 false and no exception will be signaled.

4 7 **Example.** IEEE\_QUIET\_LT (1.0, IEEE\_VALUE (IEEE\_QUIET\_NAN)) has the value false and no exception  
 5 is signaled.

## 6 14.11.29 IEEE\_QUIET\_NE (A, B)

7 1 **Description.** Quiet compares not equal.

8 2 **Class.** Elemental function.

9 3 **Arguments.**

10 A shall be of type real.

11 B shall have the same type and type parameters as A.

12 4 **Restriction.** IEEE\_QUIET\_NE (A) shall not be invoked if IEEE\_SUPPORT\_DATATYPE (A) has the value  
 13 false.

14 5 **Result Characteristics.** Default logical.

15 6 **Result Value.** The result has the value specified for the compareQuietNotEqual operation in ISO/IEC/IEEE  
 16 60559:2011; that is, it is true if and only if A compares not equal to B; if A or B is a NaN, the result will be true  
 17 and no exception will be signaled.

18 7 **Example.** IEEE\_QUIET\_NE (1.0, IEEE\_VALUE (IEEE\_QUIET\_NAN)) has the value true and no exception  
 19 is signaled.

## 20 14.11.30 IEEE\_REAL (A, [, KIND])

21 1 **Description.** Conversion to real type.

22 2 **Class.** Elemental function.

23 3 **Arguments.**

24 A shall be of type integer or real.

25 KIND (optional) shall be a scalar integer constant expression.

26 4 **Restriction.** IEEE\_REAL shall not be invoked if A is of type real and IEEE\_SUPPORT\_DATATYPE (A)  
 27 has the value false, or if IEEE\_SUPPORT\_DATATYPE (IEEE\_REAL (A, KIND)) has the value false.

28 5 **Result Characteristics.** Real. If KIND is present, the kind type parameter is that specified by the value of  
 29 KIND; otherwise, the kind type parameter is that of default real.

30 6 **Result Value.** The result has the same value as A if that value is representable in the representation method  
 31 of the result, and is rounded according to the rounding mode otherwise. This shall be consistent with the  
 32 specification of ISO/IEC/IEEE 60559:2011 for the convertFromInt operation when A is of type integer, and with  
 33 the convertFormat operation otherwise.

34 7 **Example.** The value of IEEE\_REAL (123) is 123.0.

## 35 14.11.31 IEEE\_REM (X, Y)

36 1 **Description.** Exact remainder.

37 2 **Class.** Elemental function.

- 1 3 **Arguments.** The arguments shall be of type real and have the same radix.
- 2 4 **Restriction.** IEEE\_REM (X, Y) shall not be invoked if IEEE\_SUPPORT\_DATATYPE (X) or IEEE\_SUPPORT\_DATATYPE (Y) has the value false.
- 3
- 4 5 **Result Characteristics.** Real with the kind type parameter of whichever argument has the greater precision.
- 5 6 **Result Value.** The result value, regardless of the rounding mode, shall be exactly  $X - Y \cdot N$ , where N is the integer nearest to the exact value  $X/Y$ ; whenever  $|N - X/Y| = \frac{1}{2}$ , N shall be even. If the result value is zero, the sign shall be that of X. This function computes the remainder operation specified in ISO/IEC/IEEE 60559:2011.
- 6
- 7
- 8 7 **Examples.** The value of IEEE\_REM (4.0, 3.0) is 1.0, the value of IEEE\_REM (3.0, 2.0) is -1.0, and the value of IEEE\_REM (5.0, 2.0) is 1.0.
- 9

### 10 14.11.32 IEEE\_RINT (X [, ROUND])

- 11 1 **Description.** Round to integer.
- 12 2 **Class.** Elemental function.
- 13 3 **Arguments.**
- 14 X shall be of type real.
- 15 ROUND (optional) shall be of type IEEE\_ROUND\_TYPE.
- 16 4 **Restriction.** IEEE\_RINT (X) shall not be invoked if IEEE\_SUPPORT\_DATATYPE (X) has the value false.
- 17 5 **Result Characteristics.** Same as X.
- 18 6 **Result Value.** If ROUND is present, the value of the result is the value of X rounded to an integer according to the mode specified by ROUND; this is the ISO/IEC/IEEE 60559:2011 operation roundToInteger{rounding}. Otherwise, the value of the result is that specified for the operation roundIntegralToExact in ISO/IEC/IEEE 60559:2011; this is the value of X rounded to an integer according to the rounding mode. If the result has the value zero, the sign is that of X.
- 19
- 20
- 21
- 22
- 23 7 **Examples.** If the rounding mode is round to nearest, the value of IEEE\_RINT (1.1) is 1.0. The value of IEEE\_RINT (1.1, IEEE\_UP) is 2.0.
- 24

### 25 14.11.33 IEEE\_SCALB (X, I)

- 26 1 **Description.**  $X \times 2^I$ .
- 27 2 **Class.** Elemental function.
- 28 3 **Arguments.**
- 29 X shall be of type real.
- 30 I shall be of type integer.
- 31 4 **Restriction.** IEEE\_SCALB (X) shall not be invoked if IEEE\_SUPPORT\_DATATYPE (X) has the value false.
- 32 5 **Result Characteristics.** Same as X.
- 33 6 **Result Value.**
- 34 *Case (i):* If  $X \times 2^I$  is representable as a normal number, the result has this value.
- 35 *Case (ii):* If X is finite and  $X \times 2^I$  is too large, the IEEE\_OVERFLOW exception shall occur. If IEEE\_SUPPORT\_INF (X) is true, the result value is infinity with the sign of X; otherwise, the result value is SIGN (HUGE (X), X).
- 36
- 37

- 1 *Case (iii):* If  $X \times 2^I$  is too small and there is loss of accuracy, the IEEE\_UNDERFLOW exception shall occur.  
 2 The result is the representable number having a magnitude nearest to  $|2^I|$  and the same sign as X.  
 3 *Case (iv):* If X is infinite, the result is the same as X; no exception signals.

4 7 **Example.** The value of IEEE\_SCALB (1.0, 2) is 4.0.

### 5 14.11.34 IEEE\_SELECTED\_REAL\_KIND ([P, R, RADIX])

6 1 **Description.** IEEE kind type parameter value.

7 2 **Class.** Transformational function.

8 3 **Arguments.** At least one argument shall be present.

9 P (optional) shall be an integer scalar.

10 R (optional) shall be an integer scalar.

11 RADIX (optional) shall be an integer scalar.

12 4 **Result Characteristics.** Default integer scalar.

13 5 **Result Value.** If P or R is absent, the result value is the same as if it were present with the value zero. If  
 14 RADIX is absent, there is no requirement on the radix of the selected kind. The result has a value equal to a  
 15 value of the kind type parameter of an ISO/IEC/IEEE 60559:2011 floating-point format with decimal precision,  
 16 as returned by the intrinsic function PRECISION, of at least P digits, a decimal exponent range, as returned  
 17 by the intrinsic function RANGE, of at least R, and a radix, as returned by the intrinsic function RADIX, of  
 18 RADIX, if such a kind type parameter is available on the processor.

19 6 Otherwise, the result is  $-1$  if the processor supports an IEEE real type with radix RADIX and exponent range  
 20 of at least R but not with precision of at least P,  $-2$  if the processor supports an IEEE real type with radix  
 21 RADIX and precision of at least P but not with exponent range of at least R,  $-3$  if the processor supports an  
 22 IEEE real type with radix RADIX but with neither precision of at least P nor exponent range of at least R,  $-4$  if  
 23 the processor supports an IEEE real type with radix RADIX and either precision of at least P or exponent range  
 24 of at least R but not both together, and  $-5$  if the processor supports no IEEE real type with radix RADIX.

25 7 If more than one kind type parameter value meets the criteria, the value returned is the one with the smallest  
 26 decimal precision, unless there are several such values, in which case the smallest of these kind values is returned.

27 8 **Example.** IEEE\_SELECTED\_REAL\_KIND (6, 30) has the value KIND (0.0) on a machine that supports  
 28 ISO/IEC/IEEE 60559:2011 single precision arithmetic for its default real approximation method.

### 29 14.11.35 IEEE\_SET\_FLAG (FLAG, FLAG\_VALUE)

30 1 **Description.** Set an exception flag.

31 2 **Class.** Pure subroutine.

32 3 **Arguments.**

33 FLAG shall be a scalar or array of type IEEE\_FLAG\_TYPE. If a value of FLAG is IEEE\_INVALID,  
 34 IEEE\_OVERFLOW, IEEE\_DIVIDE\_BY\_ZERO, IEEE\_UNDERFLOW, or IEEE\_INEXACT,  
 35 the corresponding exception flag is assigned a value. No two elements of FLAG shall have the same  
 36 value.

37 FLAG\_VALUE shall be a logical scalar or array. It shall be conformable with FLAG. If an element has the value  
 38 true, the corresponding flag is set to be signaling; otherwise, the flag is set to be quiet.

39 4 **Example.** CALL IEEE\_SET\_FLAG (IEEE\_OVERFLOW, .TRUE.) sets the IEEE\_OVERFLOW flag to be  
 40 signaling.

### 14.11.36 IEEE\_SET\_HALTING\_MODE (FLAG, HALTING)

**1 Description.** Set a [halting mode](#).

**2 Class.** Pure subroutine.

**3 Arguments.**

FLAG shall be a scalar or array of type [IEEE\\_FLAG\\_TYPE](#). It shall have only the values IEEE\_INVALID, IEEE\_OVERFLOW, IEEE\_DIVIDE\_BY\_ZERO, IEEE\_UNDERFLOW, or IEEE\_INEXACT. No two elements of FLAG shall have the same value.

HALTING shall be a logical scalar or array. It shall be [conformable](#) with FLAG. If an element has the value true, the corresponding exception specified by FLAG will cause halting. Otherwise, execution will continue after this exception.

**4 Restriction.** IEEE\_SET\_HALTING\_MODE (FLAG, HALTING) shall not be invoked if IEEE\_SUPPORT\_HALTING (FLAG) has the value false.

**5 Example.** CALL IEEE\_SET\_HALTING\_MODE (IEEE\_DIVIDE\_BY\_ZERO, [.TRUE.](#)) causes halting after a divide\_by\_zero exception.

### 14.11.37 IEEE\_SET\_MODES (MODES)

**1 Description.** Set floating-point modes.

**2 Class.** Subroutine.

**3 Argument.** MODES shall be a scalar of type [IEEE\\_MODES\\_TYPE](#). Its value shall be one that was assigned by a previous invocation of [IEEE\\_GET\\_MODES](#) to its MODES argument. The floating-point modes ([14.7](#)) are restored to the state at that invocation.

**4 Example.**

To save the floating-point modes, do a calculation with specific rounding and [underflow](#) modes, and restore them later:

```
USE, INTRINSIC :: IEEE_ARITHMETIC
TYPE (IEEE_MODES_TYPE) SAVE_MODES
...
CALL IEEE_GET_MODES (SAVE_MODES) ! Save all modes.
CALL IEEE_SET_ROUNDING_MODE (IEEE_TO_ZERO))
CALL IEEE_SET_UNDERFLOW_MODE (GRADUAL=.FALSE.)
... ! calculation with abrupt round-to-zero.
CALL IEEE_SET_MODES (SAVE_MODES) ! Restore all modes.
```

### 14.11.38 IEEE\_SET\_ROUNDING\_MODE (ROUND\_VALUE [, RADIX])

**1 Description.** Set rounding mode.

**2 Class.** Subroutine.

**3 Arguments.**

ROUND\_VALUE shall be a scalar of type [IEEE\\_ROUND\\_TYPE](#). It specifies the rounding mode to be set.

RADIX (optional) shall be an integer scalar with the value two or ten. If RADIX is present with the value ten, the rounding mode set is the decimal rounding mode; otherwise it is the binary rounding mode.

**4 Restriction.** IEEE\_SET\_ROUNDING\_MODE (ROUND\_VALUE) shall not be invoked unless [IEEE\\_SUPPORT\\_ROUNDING](#) (ROUND\_VALUE, X) is true for some X such that [IEEE\\_SUPPORT\\_DATATYPE](#) (X)

1 is true. IEEE\_SET\_ROUNDING\_MODE (ROUND\_VALUE, RADIX) shall not be invoked unless IEEE\_SUPPORT\_ROUNDING (ROUND\_VALUE, X) is true for some X with radix RADIX such that IEEE\_SUPPORT\_DATATYPE (X) is true.

2  
3  
4 **Example.** To save the binary rounding mode, do a calculation with round to nearest, and restore the rounding mode later:

```
6      USE, INTRINSIC :: IEEE_ARITHMETIC
7      TYPE (IEEE_ROUND_TYPE) ROUND_VALUE
8      ...
9      CALL IEEE_GET_ROUNDING_MODE (ROUND_VALUE) ! Store the rounding mode
10     CALL IEEE_SET_ROUNDING_MODE (IEEE_NEAREST)
11     ... ! calculation with round to nearest
12     CALL IEEE_SET_ROUNDING_MODE (ROUND_VALUE) ! Restore the rounding mode
```

### 13 14.11.39 IEEE\_SET\_STATUS (STATUS\_VALUE)

14 1 **Description.** Restore floating-point status.

15 2 **Class.** Subroutine.

16 3 **Argument.** STATUS\_VALUE shall be a scalar of type IEEE\_STATUS\_TYPE. Its value shall be one that was assigned by a previous invocation of IEEE\_GET\_STATUS to its STATUS\_VALUE argument. The floating-point status (14.7 is restored to the state at that invocation).

19 4 **Example.** To store all the exceptions flags, do a calculation involving exception handling, and restore them later:

```
21     USE, INTRINSIC :: IEEE_EXCEPTIONS
22     TYPE (IEEE_STATUS_TYPE) STATUS_VALUE
23     ...
24     CALL IEEE_GET_STATUS (STATUS_VALUE) ! Store the flags
25     CALL IEEE_SET_FLAG (IEEE_ALL, .FALSE.) ! Set them quiet
26     ... ! calculation involving exception handling
27     CALL IEEE_SET_STATUS (STATUS_VALUE) ! Restore the flags
```

### 28 14.11.40 IEEE\_SET\_UNDERFLOW\_MODE (GRADUAL)

29 1 **Description.** Set underflow mode.

30 2 **Class.** Subroutine.

31 3 **Argument.** GRADUAL shall be a logical scalar. If it is true, the underflow mode is set to gradual underflow. If it is false, the underflow mode is set to abrupt underflow.

33 4 **Restriction.** IEEE\_SET\_UNDERFLOW\_MODE shall not be invoked unless IEEE\_SUPPORT\_UNDERFLOW\_CONTROL (X) is true for some X.

35 5 **Example.** To perform some calculations with abrupt underflow and then restore the previous mode:

```
36     USE, INTRINSIC :: IEEE_ARITHMETIC
37     LOGICAL SAVE_UNDERFLOW_MODE
38     ...
```

```

1      CALL IEEE_GET_UNDERFLOW_MODE (SAVE_UNDERFLOW_MODE)
2      CALL IEEE_SET_UNDERFLOW_MODE (GRADUAL=.FALSE.)
3      ... ! Perform some calculations with abrupt underflow
4      CALL IEEE_SET_UNDERFLOW_MODE (SAVE_UNDERFLOW_MODE)

```

#### 14.11.41 IEEE\_SIGNBIT (X)

1 **Description.** Test sign bit.

2 **Class.** [Elemental](#) function.

3 **Argument.** X shall be of type real.

4 **Restriction.** IEEE\_SIGNBIT (X) shall not be invoked if [IEEE\\_SUPPORT\\_DATATYPE](#) (X) has the value false.

5 **Result Characteristics.** Default logical.

6 **Result Value.** The result has the value specified for the isSignMinus operation in ISO/IEC/IEEE 60559:2011; that is, it is true if and only if the sign bit of X is nonzero. No exception is signaled even if X is a signaling NaN.

7 **Example.** IEEE\_SIGNBIT (−1.0) has the value true.

#### 14.11.42 IEEE\_SUPPORT\_DATATYPE () or IEEE\_SUPPORT\_DATATYPE (X)

1 **Description.** Query IEEE arithmetic support.

2 **Class.** [Inquiry](#) function.

3 **Argument.** X shall be of type real. It may be a scalar or an array.

4 **Result Characteristics.** Default logical scalar.

5 **Result Value.** The result has the value true if the processor supports IEEE arithmetic for all reals (X does not appear) or for real variables of the same kind type parameter as X; otherwise, it has the value false. Here, support is as defined in the first paragraph of [14.9](#).

6 **Example.** If default real kind conforms to ISO/IEC/IEEE 60559:2011 except that underflow values flush to zero instead of being subnormal, IEEE\_SUPPORT\_DATATYPE (1.0) has the value true.

#### 14.11.43 IEEE\_SUPPORT\_DENORMAL () or IEEE\_SUPPORT\_DENORMAL (X)

1 **Description.** Query subnormal number support.

2 **Class.** [Inquiry](#) function.

3 **Argument.** X shall be of type real. It may be a scalar or an array.

4 **Result Characteristics.** Default logical scalar.

5 **Result Value.**

*Case (i):* IEEE\_SUPPORT\_DENORMAL (X) has the value true if [IEEE\\_SUPPORT\\_DATATYPE](#) (X) has the value true and the processor supports arithmetic operations and assignments with subnormal numbers (biased exponent  $e = 0$  and fraction  $f \neq 0$ , see subclause 3.2 of ISO/IEC/IEEE 60559:2011) for real variables of the same kind type parameter as X; otherwise, it has the value false.

1 *Case (ii):* IEEE\_SUPPORT\_DENORMAL () has the value true if IEEE\_SUPPORT\_DENORMAL (X) has  
2 the value true for all real X; otherwise, it has the value false.

3 6 **Example.** IEEE\_SUPPORT\_DENORMAL (X) has the value true if the processor supports subnormal values  
4 for X.

#### NOTE 14.9

A reference to IEEE\_SUPPORT\_DENORMAL will have the same result value as a reference to IEEE\_SUPPORT\_SUBNORMAL with the same argument list.

### 5 14.11.44 IEEE\_SUPPORT\_DIVIDE () or IEEE\_SUPPORT\_DIVIDE (X)

6 1 **Description.** Query IEEE division support.

7 2 **Class.** Inquiry function.

8 3 **Argument.** X shall be of type real. It may be a scalar or an array.

9 4 **Result Characteristics.** Default logical scalar.

10 5 **Result Value.**

11 *Case (i):* IEEE\_SUPPORT\_DIVIDE (X) has the value true if the processor supports division with the  
12 accuracy specified by ISO/IEC/IEEE 60559:2011 for real variables of the same kind type parameter  
13 as X; otherwise, it has the value false.

14 *Case (ii):* IEEE\_SUPPORT\_DIVIDE () has the value true if IEEE\_SUPPORT\_DIVIDE (X) has the value  
15 true for all real X; otherwise, it has the value false.

16 6 **Example.** IEEE\_SUPPORT\_DIVIDE (X) has the value true if division of operands with the same kind as X  
17 conforms to ISO/IEC/IEEE 60559:2011.

### 18 14.11.45 IEEE\_SUPPORT\_FLAG (FLAG) or IEEE\_SUPPORT\_FLAG (FLAG, X)

19 1 **Description.** Query exception support.

20 2 **Class.** Inquiry function.

21 3 **Arguments.**

22 FLAG shall be a scalar of type IEEE\_FLAG\_TYPE. Its value shall be one of IEEE\_INVALID, IEEE\_  
23 OVERFLOW, IEEE\_DIVIDE\_BY\_ZERO, IEEE\_UNDERFLOW, or IEEE\_INEXACT.

24 X shall be of type real. It may be a scalar or an array.

25 4 **Result Characteristics.** Default logical scalar.

26 5 **Result Value.**

27 *Case (i):* IEEE\_SUPPORT\_FLAG (FLAG, X) has the value true if the processor supports detection of the  
28 specified exception for real variables of the same kind type parameter as X; otherwise, it has the  
29 value false.

30 *Case (ii):* IEEE\_SUPPORT\_FLAG (FLAG) has the value true if IEEE\_SUPPORT\_FLAG (FLAG, X) has  
31 the value true for all real X; otherwise, it has the value false.

32 6 **Example.** IEEE\_SUPPORT\_FLAG (IEEE\_INEXACT) has the value true if the processor supports the inexact  
33 exception.

#### 14.11.46 IEEE\_SUPPORT\_HALTING (FLAG)

**1 Description.** Query [halting mode](#) support.

**2 Class.** [Inquiry function](#).

**3 Argument.** FLAG shall be a scalar of type [IEEE\\_FLAG\\_TYPE](#). Its value shall be one of IEEE\_INVALID, IEEE\_OVERFLOW, IEEE\_DIVIDE\_BY\_ZERO, IEEE\_UNDERFLOW, or IEEE\_INEXACT.

**4 Result Characteristics.** Default logical scalar.

**5 Result Value.** The result has the value true if the processor supports the ability to control during program execution whether to abort or continue execution after the exception specified by FLAG; otherwise, it has the value false. Support includes the ability to change the mode by CALL [IEEE\\_SET\\_HALTING\\_MODE](#) (FLAG).

**6 Example.** IEEE\_SUPPORT\_HALTING (IEEE\_OVERFLOW) has the value true if the processor supports control of [halting](#) after an overflow.

#### 14.11.47 IEEE\_SUPPORT\_INF () or IEEE\_SUPPORT\_INF (X)

**1 Description.** Query IEEE infinity support.

**2 Class.** [Inquiry function](#).

**3 Argument.** X shall be of type real. It may be a scalar or an array.

**4 Result Characteristics.** Default logical scalar.

**5 Result Value.**

*Case (i):* IEEE\_SUPPORT\_INF (X) has the value true if the processor supports IEEE infinities (positive and negative) for real variables of the same kind type parameter as X; otherwise, it has the value false.

*Case (ii):* IEEE\_SUPPORT\_INF () has the value true if IEEE\_SUPPORT\_INF (X) has the value true for all real X; otherwise, it has the value false.

**6 Example.** IEEE\_SUPPORT\_INF (X) has the value true if the processor supports IEEE infinities for X.

#### 14.11.48 IEEE\_SUPPORT\_IO () or IEEE\_SUPPORT\_IO (X)

**1 Description.** Query IEEE formatting support.

**2 Class.** [Inquiry function](#).

**3 Argument.** X shall be of type real. It may be a scalar or an array.

**4 Result Characteristics.** Default logical scalar.

**5 Result Value.**

*Case (i):* IEEE\_SUPPORT\_IO (X) has the value true if base conversion during formatted input/output ([9.5.6.16](#), [9.6.2.13](#), [10.7.2.3.8](#)) conforms to ISO/IEC/IEEE 60559:2011 for the modes UP, DOWN, ZERO, and NEAREST for real variables of the same kind type parameter as X; otherwise, it has the value false.

*Case (ii):* IEEE\_SUPPORT\_IO () has the value true if IEEE\_SUPPORT\_IO (X) has the value true for all real X; otherwise, it has the value false.

**6 Example.** IEEE\_SUPPORT\_IO (X) has the value true if formatted input/output base conversions conform to ISO/IEC/IEEE 60559:2011.

#### 14.11.49 IEEE\_SUPPORT\_NAN () or IEEE\_SUPPORT\_NAN (X)

**1 Description.** Query IEEE NaN support.

**2 Class.** Inquiry function.

**3 Argument.** X shall be of type real. It may be a scalar or an array.

**4 Result Characteristics.** Default logical scalar.

**5 Result Value.**

*Case (i):* IEEE\_SUPPORT\_NAN (X) has the value true if the processor supports IEEE NaNs for real variables of the same kind type parameter as X; otherwise, it has the value false.

*Case (ii):* IEEE\_SUPPORT\_NAN () has the value true if IEEE\_SUPPORT\_NAN (X) has the value true for all real X; otherwise, it has the value false.

**6 Example.** IEEE\_SUPPORT\_NAN (X) has the value true if the processor supports IEEE NaNs for X.

#### 14.11.50 IEEE\_SUPPORT\_ROUNDING (ROUND\_VALUE) or IEEE\_SUPPORT\_ROUNDING (ROUND\_VALUE, X)

**1 Description.** Query IEEE rounding support.

**2 Class.** Inquiry function.

**3 Arguments.**

ROUND\_VALUE shall be of type IEEE\_ROUND\_TYPE.

X shall be of type real. It may be a scalar or an array.

**4 Result Characteristics.** Default logical scalar.

**5 Result Value.**

*Case (i):* IEEE\_SUPPORT\_ROUNDING (ROUND\_VALUE, X) has the value true if the processor supports the rounding mode defined by ROUND\_VALUE for real variables of the same kind type parameter as X; otherwise, it has the value false. Support includes the ability to change the mode by CALL IEEE\_SET\_ROUNDING\_MODE (ROUND\_VALUE).

*Case (ii):* IEEE\_SUPPORT\_ROUNDING (ROUND\_VALUE) has the value true if IEEE\_SUPPORT\_ROUNDING (ROUND\_VALUE, X) has the value true for all real X; otherwise, it has the value false.

**6 Example.** IEEE\_SUPPORT\_ROUNDING (IEEE\_TO\_ZERO) has the value true if the processor supports rounding to zero for all reals.

#### 14.11.51 IEEE\_SUPPORT\_SQRT () or IEEE\_SUPPORT\_SQRT (X)

**1 Description.** Query IEEE square root support.

**2 Class.** Inquiry function.

**3 Argument.** X shall be of type real. It may be a scalar or an array.

**4 Result Characteristics.** Default logical scalar.

**5 Result Value.**

*Case (i):* IEEE\_SUPPORT\_SQRT (X) has the value true if the intrinsic function SQRT conforms to ISO/IEC/IEEE 60559:2011 for real variables of the same kind type parameter as X; otherwise, it has the value false.

1 *Case (ii):* IEEE\_SUPPORT\_SQRT () has the value true if IEEE\_SUPPORT\_SQRT (X) has the value true  
2 for all real X; otherwise, it has the value false.

3 6 **Example.** If IEEE\_SUPPORT\_SQRT (1.0) has the value true, SQRT (−0.0) will have the value −0.0.

#### 4 14.11.52 IEEE\_SUPPORT\_STANDARD () or IEEE\_SUPPORT\_STANDARD (X)

5 1 **Description.** Query IEEE standard support.

6 2 **Class.** Inquiry function.

7 3 **Argument.** X shall be of type real. It may be a scalar or an array.

8 4 **Result Characteristics.** Default logical scalar.

9 5 **Result Value.**

10 *Case (i):* IEEE\_SUPPORT\_STANDARD (X) has the value true if the results of all the functions  
11 IEEE\_SUPPORT\_DATATYPE (X), IEEE\_SUPPORT\_DIVIDE (X), IEEE\_SUPPORT-  
12 FLAG (FLAG, X) for valid FLAG, IEEE\_SUPPORT\_HALTING (FLAG) for valid  
13 FLAG, IEEE\_SUPPORT\_INF (X), IEEE\_SUPPORT\_NAN (X), IEEE\_SUPPORT\_-  
14 ROUNDING (ROUND\_VALUE, X) for valid ROUND\_VALUE, IEEE\_SUPPORT\_SQRT  
15 (X), and IEEE\_SUPPORT\_SUBNORMAL (X) are all true; otherwise, it has the value false.

16 *Case (ii):* IEEE\_SUPPORT\_STANDARD () has the value true if IEEE\_SUPPORT\_STANDARD (X) has  
17 the value true for all real X; otherwise, it has the value false.

18 6 **Example.** IEEE\_SUPPORT\_STANDARD () has the value false if some but not all kinds of reals conform to  
19 ISO/IEC/IEEE 60559:2011.

#### 20 14.11.53 IEEE\_SUPPORT\_SUBNORMAL () or IEEE\_SUPPORT\_SUBNORMAL (X)

21 1 **Description.** Query subnormal number support.

22 2 **Class.** Inquiry function.

23 3 **Argument.** X shall be of type real. It may be a scalar or an array.

24 4 **Result Characteristics.** Default logical scalar.

25 5 **Result Value.**

26 *Case (i):* IEEE\_SUPPORT\_SUBNORMAL (X) has the value true if IEEE\_SUPPORT\_DATATYPE (X)  
27 has the value true and the processor supports arithmetic operations and assignments with subnormal  
28 numbers (biased exponent  $e = 0$  and fraction  $f \neq 0$ , see subclause 3.2 of ISO/IEC/IEEE 60559:2011)  
29 for real variables of the same kind type parameter as X; otherwise, it has the value false.

30 *Case (ii):* IEEE\_SUPPORT\_SUBNORMAL () has the value true if IEEE\_SUPPORT\_SUBNORMAL (X)  
31 has the value true for all real X; otherwise, it has the value false.

32 6 **Example.** IEEE\_SUPPORT\_SUBNORMAL (X) has the value true if the processor supports subnormal values  
33 for X.

#### NOTE 14.10

The subnormal numbers are not included in the 13.4 model for real numbers; they satisfy the inequality  $ABS(X) < TINY(X)$ . They usually occur as a result of an arithmetic operation whose exact result is less than  $TINY(X)$ . Such an operation causes IEEE\_UNDERFLOW to signal unless the result is exact. IEEE\_SUPPORT\_SUBNORMAL (X) is false if the processor never returns a subnormal number as the result of an arithmetic operation.

#### 14.11.54 IEEE\_SUPPORT\_UNDERFLOW\_CONTROL () or IEEE\_SUPPORT\_UNDERFLOW\_CONTROL (X)

**Description.** Query [underflow](#) control support.

**Class.** [Inquiry function](#).

**Argument.** X shall be of type real. It may be a scalar or an array.

**Result Characteristics.** Default logical scalar.

**Result Value.**

*Case (i):* IEEE\_SUPPORT\_UNDERFLOW\_CONTROL (X) has the value true if the processor supports control of the [underflow mode](#) for floating-point calculations with the same type as X, and false otherwise.

*Case (ii):* IEEE\_SUPPORT\_UNDERFLOW\_CONTROL () has the value true if the processor supports control of the [underflow mode](#) for all floating-point calculations, and false otherwise.

**Example.** IEEE\_SUPPORT\_UNDERFLOW\_CONTROL (2.5) has the value true if the processor supports [underflow mode](#) control for default real calculations.

#### 14.11.55 IEEE\_UNORDERED (X, Y)

**Description.** Whether two values are unordered.

**Class.** [Elemental function](#).

**Arguments.** The arguments shall be of type real.

**Restriction.** IEEE\_UNORDERED (X, Y) shall not be invoked if IEEE\_SUPPORT\_DATATYPE (X) or IEEE\_SUPPORT\_DATATYPE (Y) has the value false.

**Result Characteristics.** Default logical.

**Result Value.** The result has the value true if X or Y is a NaN or both are NaNs; otherwise, it has the value false.

**Example.** IEEE\_UNORDERED (0.0, [SQRT](#) (−1.0)) has the value true if IEEE\_SUPPORT\_SQRT (1.0) has the value true.

#### 14.11.56 IEEE\_VALUE (X, CLASS)

**Description.** Return number in a class.

**Class.** [Elemental function](#).

**Arguments.**

X shall be of type real.

CLASS shall be of type [IEEE\\_CLASS\\_TYPE](#). The value is permitted to be: IEEE\_SIGNALING\_NAN or IEEE\_QUIET\_NAN if [IEEE\\_SUPPORT\\_NAN](#) (X) has the value true, IEEE\_NEGATIVE\_INF or IEEE\_POSITIVE\_INF if [IEEE\\_SUPPORT\\_INF](#) (X) has the value true, IEEE\_NEGATIVE\_SUBNORMAL or IEEE\_POSITIVE\_SUBNORMAL if [IEEE\\_SUPPORT\\_SUBNORMAL](#) (X) has the value true, IEEE\_NEGATIVE\_NORMAL, IEEE\_NEGATIVE\_ZERO, IEEE\_POSITIVE\_ZERO or IEEE\_POSITIVE\_NORMAL.

**Restriction.** IEEE\_VALUE (X, CLASS) shall not be invoked if [IEEE\\_SUPPORT\\_DATATYPE](#) (X) has the value false.

**Result Characteristics.** Same as X.

- 1 6 **Result Value.** The result value is an IEEE value as specified by CLASS. Although in most cases the value is  
 2 processor dependent, the value shall not vary between invocations for any particular X kind type parameter and  
 3 CLASS value.
- 4 7 **Example.** IEEE\_VALUE (1.0, IEEE\_NEGATIVE\_INF) has the value  $-\infty$ .
- 5 8 Whenever IEEE\_VALUE returns a signaling NaN, it is processor dependent whether or not invalid is raised and  
 6 processor dependent whether or not the signaling NaN is converted into a quiet NaN.

**NOTE 14.11**

If the *expr* in an *assignment statement* is a reference to the IEEE\_VALUE function that returns a signaling NaN and the *variable* is of the same type and kind as the function result, it is recommended that the signaling NaN be preserved.

## 7 14.12 Examples

**NOTE 14.12**

```

MODULE DOT
  ! Module for dot product of two real arrays of rank 1.
  ! The caller needs to ensure that exceptions do not cause halting.
  USE, INTRINSIC :: IEEE_EXCEPTIONS
  LOGICAL :: MATRIX_ERROR = .FALSE.
  INTERFACE OPERATOR(.dot.)
    MODULE PROCEDURE MULT
  END INTERFACE
CONTAINS
  REAL FUNCTION MULT (A, B)
    REAL, INTENT (IN) :: A(:), B(:)
    INTEGER I
    LOGICAL OVERFLOW
    IF (SIZE(A) /= SIZE(B)) THEN
      MATRIX_ERROR = .TRUE.
      RETURN
    END IF
    ! The processor ensures that IEEE_OVERFLOW is quiet.
    MULT = 0.0
    DO I = 1, SIZE (A)
      MULT = MULT + A(I)*B(I)
    END DO
    CALL IEEE_GET_FLAG (IEEE_OVERFLOW, OVERFLOW)
    IF (OVERFLOW) MATRIX_ERROR = .TRUE.
  END FUNCTION MULT
END MODULE DOT

```

This module provides a function that computes the dot product of two real arrays of *rank* 1. If the sizes of the arrays are different, an immediate return occurs with MATRIX\_ERROR true. If overflow occurs during the actual calculation, the IEEE\_OVERFLOW flag will signal and MATRIX\_ERROR will be true.

## NOTE 14.13

```

USE, INTRINSIC :: IEEE_EXCEPTIONS
USE, INTRINSIC :: IEEE_FEATURES, ONLY: IEEE_INVALID_FLAG
! The other exceptions of IEEE_USUAL (IEEE_OVERFLOW and
! IEEE_DIVIDE_BY_ZERO) are always available with IEEE_EXCEPTIONS
TYPE (IEEE_STATUS_TYPE) STATUS_VALUE
LOGICAL, DIMENSION(3) :: FLAG_VALUE
...
CALL IEEE_GET_STATUS (STATUS_VALUE)
CALL IEEE_SET_HALTING_MODE (IEEE_USUAL, .FALSE.) ! Needed in case the
!               default on the processor is to halt on exceptions
CALL IEEE_SET_FLAG (IEEE_USUAL, .FALSE.)
! First try the "fast" algorithm for inverting a matrix:
MATRIX1 = FAST_INV (MATRIX) ! This shall not alter MATRIX.
CALL IEEE_GET_FLAG (IEEE_USUAL, FLAG_VALUE)
IF (ANY(FLAG_VALUE)) THEN
  ! "Fast" algorithm failed; try "slow" one:
  CALL IEEE_SET_FLAG (IEEE_USUAL, .FALSE.)
  MATRIX1 = SLOW_INV (MATRIX)
  CALL IEEE_GET_FLAG (IEEE_USUAL, FLAG_VALUE)
  IF (ANY (FLAG_VALUE)) THEN
    WRITE (*, *) 'Cannot invert matrix'
    STOP
  END IF
END IF
CALL IEEE_SET_STATUS (STATUS_VALUE)

```

In this example, the function FAST\_INV might cause a condition to signal. If it does, another try is made with SLOW\_INV. If this still fails, a message is printed and the program stops. Note, also, that it is important to set the flags quiet before the second try. The state of all the flags is stored and restored.

## NOTE 14.14

```

USE, INTRINSIC :: IEEE_EXCEPTIONS
LOGICAL FLAG_VALUE
...
CALL IEEE_SET_HALTING_MODE (IEEE_OVERFLOW, .FALSE.)
! First try a fast algorithm for inverting a matrix.
CALL IEEE_SET_FLAG (IEEE_OVERFLOW, .FALSE.)
DO K = 1, N
  ...
  CALL IEEE_GET_FLAG (IEEE_OVERFLOW, FLAG_VALUE)
  IF (FLAG_VALUE) EXIT
END DO
IF (FLAG_VALUE) THEN
  ! Alternative code which knows that K-1 steps have executed normally.

```

**NOTE 14.14 (cont.)**

...  
END IF

Here the code for matrix inversion is in line and the transfer is made more precise by adding extra tests of the flag.



# 15 Interoperability with C

## 15.1 General

- 1 Fortran provides a means of referencing procedures that are defined by means of the C programming language or procedures that can be described by C prototypes as defined in 6.7.6.3 of ISO/IEC 9899:2011, even if they are not actually defined by means of C. Conversely, there is a means of specifying that a procedure defined by a Fortran subprogram can be referenced from a function defined by means of C. In addition, there is a means of declaring global variables that are associated with C variables whose names have external linkage as defined in 6.2.2 of ISO/IEC 9899:2011.
- 2 The ISO\_C\_BINDING module provides access to [named constants](#) that represent [kind type parameters](#) of data representations compatible with C types. Fortran also provides facilities for defining derived types (4.5) and enumerations (4.6) that correspond to C types.
- 3 The source file ISO\_Fortran\_binding.h provides definitions and prototypes to enable a C function to interoperate with a Fortran procedure that has a dummy data object that is [allocatable](#), [assumed-shape](#), [assumed-rank](#), [pointer](#), or is of type character with an assumed length.

## 15.2 The ISO\_C\_BINDING intrinsic module

### 15.2.1 Summary of contents

- 1 The processor shall provide the intrinsic module ISO\_C\_BINDING. This module shall make accessible the following entities: the [named constants](#) C\_NULL\_PTR, C\_NULL\_FUNPTR, and those with names listed in the first column of Table 15.1 and the second column of Table 15.2, the types C\_PTR and C\_FUNPTR, and the procedures in 15.2.3. A processor may provide other public entities in the ISO\_C\_BINDING intrinsic module in addition to those listed here.

### 15.2.2 Named constants and derived types in the module

- 1 The entities listed in the second column of Table 15.2 shall be default integer [named constants](#).
- 2 A Fortran [intrinsic type](#) whose [kind type parameter](#) is one of the values in the module shall have the same representation as the C type with which it interoperates, for each value that a variable of that type can have. For C\_BOOL, the internal representation of .TRUE.\_C\_BOOL and .FALSE.\_C\_BOOL shall be the same as those of the C values (\_Bool)1 and (\_Bool)0 respectively.
- 3 The value of C\_INT shall be a valid value for an integer kind parameter on the processor. The values of C\_SHORT, C\_LONG, C\_LONG\_LONG, C\_SIGNED\_CHAR, C\_SIZE\_T, C\_INT8\_T, C\_INT16\_T, C\_INT32\_T, C\_INT64\_T, C\_INT\_LEAST8\_T, C\_INT\_LEAST16\_T, C\_INT\_LEAST32\_T, C\_INT\_LEAST64\_T, C\_INT\_FAST8\_T, C\_INT\_FAST16\_T, C\_INT\_FAST32\_T, C\_INT\_FAST64\_T, C\_INTMAX\_T, and C\_INTPTR\_T shall each be a valid value for an integer [kind type parameter](#) on the processor or shall be -1 if the [companion processor](#) (2.5.7) defines the corresponding C type and there is no interoperating Fortran processor kind, or -2 if the [companion processor](#) does not define the corresponding C type.
- 4 The values of C\_FLOAT, C\_DOUBLE, and C\_LONG\_DOUBLE shall each be a valid value for a real [kind type parameter](#) on the processor or shall be -1 if the [companion processor](#)'s type does not have a precision equal to the precision of any of the Fortran processor's real kinds, -2 if the [companion processor](#)'s type does not have a range equal to the range of any of the Fortran processor's real kinds, -3 if the [companion processor](#)'s type

has neither the precision nor range of any of the Fortran processor's real kinds, and equal to  $-4$  if there is no interoperating Fortran processor kind for other reasons. The values of `C_FLOAT_COMPLEX`, `C_DOUBLE_COMPLEX`, and `C_LONG_DOUBLE_COMPLEX` shall be the same as those of `C_FLOAT`, `C_DOUBLE`, and `C_LONG_DOUBLE`, respectively.

The value of `C_BOOL` shall be a valid value for a logical kind parameter on the processor or shall be  $-1$ .

The value of `C_CHAR` shall be a valid value for a character [kind type parameter](#) on the processor or shall be  $-1$ . If the value of `C_CHAR` is nonnegative, the character kind specified is the C character kind; otherwise, there is no C character kind.

The following entities shall be [named constants](#) of type character with a length parameter of one. The kind parameter value shall be equal to the value of `C_CHAR` unless `C_CHAR` =  $-1$ , in which case the kind parameter value shall be the same as for default kind. The values of these constants are specified in Table 15.1. In the case that `C_CHAR`  $\neq -1$  the value is specified using C syntax. The semantics of these values are explained in 5.2.1 and 5.2.2 of ISO/IEC 9899:2011.

Table 15.1: Names of C characters with special semantics

| Name                           | C definition    | Value                      |                               |
|--------------------------------|-----------------|----------------------------|-------------------------------|
|                                |                 | <code>C_CHAR</code> = $-1$ | <code>C_CHAR</code> $\neq -1$ |
| <code>C_NULL_CHAR</code>       | null character  | <code>CHAR(0)</code>       | <code>'\0'</code>             |
| <code>C_ALERT</code>           | alert           | <code>ACHAR(7)</code>      | <code>'\a'</code>             |
| <code>C_BACKSPACE</code>       | backspace       | <code>ACHAR(8)</code>      | <code>'\b'</code>             |
| <code>C_FORM_FEED</code>       | form feed       | <code>ACHAR(12)</code>     | <code>'\f'</code>             |
| <code>C_NEW_LINE</code>        | new line        | <code>ACHAR(10)</code>     | <code>'\n'</code>             |
| <code>C_CARRIAGE_RETURN</code> | carriage return | <code>ACHAR(13)</code>     | <code>'\r'</code>             |
| <code>C_HORIZONTAL_TAB</code>  | horizontal tab  | <code>ACHAR(9)</code>      | <code>'\t'</code>             |
| <code>C_VERTICAL_TAB</code>    | vertical tab    | <code>ACHAR(11)</code>     | <code>'\v'</code>             |

The entities `C_PTR` and `C_FUNPTR` are described in 15.3.3.

The entity `C_NULL_PTR` shall be a [named constant](#) of type `C_PTR`. The value of `C_NULL_PTR` shall be the same as the value `NULL` in C. The entity `C_NULL_FUNPTR` shall be a [named constant](#) of type `C_FUNPTR`. The value of `C_NULL_FUNPTR` shall be that of a null pointer to a function in C.

#### NOTE 15.1

The value of `NEW_LINE(C_NEW_LINE)` is `C_NEW_LINE` (13.9.140).

## 15.2.3 Procedures in the module

### 15.2.3.1 General

In the detailed descriptions below, procedure names are generic and not specific. The `C_F_POINTER` subroutine is impure; all other procedures in the module are pure.

#### 15.2.3.2 `C_ASSOCIATED(C_PTR_1 [, C_PTR_2])`

**Description.** Query C pointer status.

**Class.** [Inquiry function](#).

**Arguments.**

`C_PTR_1` shall be a scalar of type `C_PTR` or `C_FUNPTR`.

`C_PTR_2` (optional) shall be a scalar of the same type as `C_PTR_1`.

1 4 **Result Characteristics.** Default logical scalar.

2 5 **Result Value.**

- 3 *Case (i):* If C\_PTR\_2 is absent, the result is false if C\_PTR\_1 is a C null pointer and true otherwise.
- 4 *Case (ii):* If C\_PTR\_2 is present, the result is false if C\_PTR\_1 is a C null pointer. If C\_PTR\_1 is not a
- 5 C null pointer, the result is true if C\_PTR\_1 compares equal to C\_PTR\_2 in the sense of 6.3.2.3
- 6 and 6.5.9 of ISO/IEC 9899:2011, and false otherwise.

#### NOTE 15.2

The following example illustrates the use of C\_LOC and C\_ASSOCIATED.

```
USE, INTRINSIC :: ISO_C_BINDING, ONLY: C_PTR, C_FLOAT, C_ASSOCIATED, C_LOC
INTERFACE
  SUBROUTINE FOO(GAMMA) BIND(C)
    IMPORT C_PTR
    TYPE(C_PTR), VALUE :: GAMMA
  END SUBROUTINE FOO
END INTERFACE
REAL(C_FLOAT), TARGET, DIMENSION(100) :: ALPHA
TYPE(C_PTR) :: BETA
...
IF (.NOT. C_ASSOCIATED(BETA)) THEN
  BETA = C_LOC(ALPHA)
ENDIF
CALL FOO(BETA)
```

#### 7 15.2.3.3 C\_F\_POINTER (CPTR, FPTR [, SHAPE])

8 1 **Description.** Associate a data pointer with the [target](#) of a C pointer and specify its shape.

9 2 **Class.** Subroutine.

10 3 **Arguments.**

- 11 CPTR shall be a scalar of type [C\\_PTR](#). It is an [INTENT \(IN\)](#) argument. Its value shall be
- 12 • the [C address](#) of an [interoperable](#) data entity,
  - 13 • the result of a reference to C\_LOC with a noninteroperable argument, or
  - 14 • the [C address](#) of a storage sequence that is not in use by any other Fortran entity.

15 The value of CPTR shall not be the [C address](#) of a Fortran variable that does not have the [TARGET](#)

16 [attribute](#).

17 FPTR shall be a pointer, shall not have a [deferred type parameter](#), and shall not be a [coindexed object](#).

18 It is an [INTENT \(OUT\)](#) argument. If FPTR is an array, its shape is specified by SHAPE and each

19 lower bound is equal to 1.

- 20 *Case (i):* If the value of CPTR is the [C address](#) of an [interoperable](#) data entity, FPTR
- 21 shall be a [data pointer](#) with type and [type parameter](#) values [interoperable](#) with
- 22 the type of the entity. If the target T of CPTR is scalar, FPTR becomes pointer
- 23 associated with T; if FPTR is an array, SHAPE shall specify a size of 1. If T is
- 24 an array, and FPTR is scalar, FPTR becomes associated with the first element of
- 25 T. If both T and FPTR are arrays, SHAPE shall specify a size that is less than or
- 26 equal to the size of T, and FPTR becomes associated with the first [PRODUCT](#)
- 27 (SHAPE) elements of T (this could be the entirety of T).

*Case (ii):* If the value of CPTR is the result of a reference to C\_LOC with a noninteroperable [effective argument](#) X, FPTR shall be a nonpolymorphic pointer with the same type and type parameters as X. In this case, X shall not have been deallocated or have become undefined due to execution of a [RETURN](#) or [END](#) statement since the reference. If X is scalar, FPTR becomes [pointer associated](#) with X; if FPTR is an array, SHAPE shall specify a size of 1. If X is an array and FPTR is scalar, FPTR becomes associated with the first element of X. If both X and FPTR are arrays, SHAPE shall specify a size that is less than or equal to the size of X, and FPTR becomes associated with the first [PRODUCT](#) (SHAPE) elements of X (this could be the entirety of X).

*Case (iii):* If the value of CPTR is the [C address](#) of a storage sequence that is not in use by any other Fortran entity, FPTR becomes associated with that storage sequence. If FPTR is an array, its shape is specified by SHAPE and each lower bound is 1. The storage sequence shall be large enough to contain the target object described by FPTR and shall satisfy any other processor-dependent requirement for association.

SHAPE (optional) shall be a rank-one integer array. It is an [INTENT \(IN\)](#) argument. SHAPE shall be present if and only if FPTR is an array; its size shall be equal to the [rank](#) of FPTR.

#### 4 Examples.

*Case (i):*

```

extern double c_x;
void *address_of_x (void)
{
    return &c_x;
}

! Assume interface to "address_of_x" is available.
Real (C_double), Pointer :: xp
Call C_F_Pointer (address_of_x (), xp)

```

*Case (ii):*

```

Type t
    Real, Allocatable :: v(:, :)
End Type
Type(t), Target :: x
Type(C_ptr) :: xloc
xloc = C_Loc (x)
...
Type(t), Pointer :: y
Call C_F_Pointer (xloc, y)

```

*Case (iii):*

```

void *getmem (int nbits)
{
    return malloc ((nbits+CHAR_BIT-1)/CHAR_BIT);
}

! Assume interface to "getmem" is available,
! and there is a derived type "mytype" accessible.
Type(mytype), Pointer :: x
Call C_F_Pointer (getmem (Storage\_Size (x)), x)

```

The following statements illustrate the use of C\_F\_POINTER when the pointer to be set has a [deferred type](#)

parameter:

```
Character(42), Pointer :: C1
Character(:), Pointer :: C2
Call C_F_Pointer (CPTR, C1)
C2 => C1
```

This will associate C2 with the entity at the C address specified by CPTR, and specify its length to be the same as that of C1.

#### NOTE 15.3

In the case of associating FPTR with a storage sequence, there might be processor-dependent requirements such as alignment of the memory address or placement in memory.

### 15.2.3.4 C\_F\_PROCPOINTER (CPTR, FPTR)

1 **Description.** Associate a procedure pointer with the target of a C function pointer.

2 **Class.** Pure subroutine.

3 **Arguments.**

CPTR shall be a scalar of type C\_FUNPTR. It is an INTENT (IN) argument. Its value shall be the C address of a procedure that is interoperable, or the result of a reference to the function C\_FUNLOC from the intrinsic module ISO\_C\_BINDING.

FPTR shall be a procedure pointer, and shall not be a component of a coindexed object. It is an INTENT (OUT) argument. If the target of CPTR is interoperable, the interface for FPTR shall be interoperable with the prototype that describes the target of CPTR; otherwise, the interface for FPTR shall have the same characteristics as that target. FPTR becomes pointer associated with the target of CPTR.

#### NOTE 15.4

The term “target” in the descriptions of C\_F\_POINTER and C\_F\_PROCPOINTER denotes the entity referenced by a C pointer, as described in 6.2.5 of ISO/IEC 9899:2011.

### 15.2.3.5 C\_FUNLOC (X)

1 **Description.** C address of the argument.

2 **Class.** Inquiry function.

3 **Argument.** X shall be a procedure; if it is a procedure pointer it shall be associated. It shall not be a coindexed object.

4 **Result Characteristics.** Scalar of type C\_FUNPTR.

5 **Result Value.** The result value is described using the result name FUNPTR. The result is determined as if C\_FUNPTR were a derived type containing a procedure pointer component PX with an implicit interface and the pointer assignment FUNPTR%PX => X were executed.

6 The result is a value that can be used as an actual CPTR argument in a call to C\_F\_PROCPOINTER where the FPTR argument has attributes that would allow the pointer assignment FPTR => X. Such a call to C\_F\_PROCPOINTER shall have the effect of the pointer assignment FPTR => X.

### 15.2.3.6 C\_LOC (X)

**Description.** C address of the argument.

**Class.** Inquiry function.

**Argument.** X shall have either the **POINTER** or **TARGET** attribute. It shall not be a **coindexed object**. It shall either be a variable with **interoperable** type and **kind type parameters**, or be a nonpolymorphic variable with no **length type parameters**. If it is **allocatable**, it shall be allocated. If it is a pointer, it shall be associated. If it is an array, it shall be **contiguous** and have nonzero size. It shall not be a zero-length string.

**Result Characteristics.** Scalar of type **C\_PTR**.

**Result Value.** The result value is described using the result name CPTR.

If X is a scalar data entity, the result is determined as if **C\_PTR** were a derived type containing a scalar pointer component PX of the type and type parameters of X and the **pointer assignment** CPTR%PX => X were executed.

If X is an array data entity, the result is determined as if **C\_PTR** were a derived type containing a scalar pointer component PX of the type and type parameters of X and the **pointer assignment** of CPTR%PX to the first element of X were executed.

If X is a data entity that is **interoperable** or has **interoperable** type and type parameters, the result is the value that the C processor returns as the result of applying the unary "&" operator (as defined in ISO/IEC 9899:2011, 6.5.3.2) to the **target** of CPTR%PX.

The result is a value that can be used as an **actual** CPTR argument in a call to **C\_F\_POINTER** where FPTR has attributes that would allow the **pointer assignment** FPTR => X. Such a call to **C\_F\_POINTER** shall have the effect of the **pointer assignment** FPTR => X.

#### NOTE 15.5

Where the **actual argument** is of noninteroperable type or type parameters, the result of **C\_LOC** provides an opaque "handle" for it. In an actual implementation, this handle might be the **C address** of the argument; however, only a C function that treats it as a void (generic) C pointer that cannot be dereferenced (6.5.3.2 in ISO/IEC 9899:2011) is likely to be portable.

### 15.2.3.7 C\_SIZEOF (X)

**Description.** Size of X in bytes.

**Class.** Inquiry function.

**Argument.** X shall be an **interoperable** data entity that is not an **assumed-size array**.

**Result Characteristics.** Scalar integer of kind **C\_SIZE\_T** (15.3.2).

**Result Value.** If X is scalar, the result value is the value that the **companion processor** returns as the result of applying the sizeof operator (ISO/IEC 9899:2011, subclause 6.5.3.4) to an object of a type that interoperates with the type and type parameters of X.

If X is an array, the result value is the value that the **companion processor** returns as the result of applying the sizeof operator to an object of a type that interoperates with the type and type parameters of X, multiplied by the number of elements in X.

## 15.3 Interoperability between Fortran and C entities

### 15.3.1 General

- 1 Subclause 15.3 defines the conditions under which a Fortran entity is [interoperable](#). If a Fortran entity is [interoperable](#), an equivalent entity could be defined by means of C and the Fortran entity would interoperate with the C entity. There does not have to be such an interoperating C entity.

#### NOTE 15.6

A Fortran entity can be [interoperable](#) with more than one C entity.

### 15.3.2 Interoperability of intrinsic types

- 1 Table 15.2 shows the interoperability between Fortran intrinsic types and C types. A Fortran intrinsic type with particular type parameter values is [interoperable](#) with a C type if the type and [kind type parameter](#) value are listed in the table on the same row as that C type. If the type is character, the [length type parameter](#) is [interoperable](#) if and only if its value is one. A combination of Fortran type and type parameters that is [interoperable](#) with a C type listed in the table is also [interoperable](#) with any unqualified C type that is compatible with the listed C type.
- 2 The second column of the table refers to the [named constants](#) made accessible by the ISO\_C\_BINDING intrinsic module. If the value of any of these [named constants](#) is negative, there is no combination of Fortran type and type parameters [interoperable](#) with the C type shown in that row.
- 3 A combination of intrinsic type and type parameters is [interoperable](#) if it is [interoperable](#) with a C type. The C types mentioned in table 15.2 are defined in subclauses 6.2.5, 7.19, and 7.20.1 of ISO/IEC 9899:2011.

Table 15.2: Interoperability between Fortran and C types

| Fortran type | Named constant from the ISO_C_BINDING module<br>(kind type parameter if value is positive) | C type                       |
|--------------|--------------------------------------------------------------------------------------------|------------------------------|
| INTEGER      | C_INT                                                                                      | int                          |
|              | C_SHORT                                                                                    | short int                    |
|              | C_LONG                                                                                     | long int                     |
|              | C_LONG_LONG                                                                                | long long int                |
|              | C_SIGNED_CHAR                                                                              | signed char<br>unsigned char |
|              | C_SIZE_T                                                                                   | size_t                       |
|              | C_INT8_T                                                                                   | int8_t                       |
|              | C_INT16_T                                                                                  | int16_t                      |
|              | C_INT32_T                                                                                  | int32_t                      |
|              | C_INT64_T                                                                                  | int64_t                      |
|              | C_INT_LEAST8_T                                                                             | int_least8_t                 |
|              | C_INT_LEAST16_T                                                                            | int_least16_t                |
|              | C_INT_LEAST32_T                                                                            | int_least32_t                |
|              | C_INT_LEAST64_T                                                                            | int_least64_t                |
|              | C_INT_FAST8_T                                                                              | int_fast8_t                  |
|              | C_INT_FAST16_T                                                                             | int_fast16_t                 |
|              | C_INT_FAST32_T                                                                             | int_fast32_t                 |
|              | C_INT_FAST64_T                                                                             | int_fast64_t                 |
|              | C_INTMAX_T                                                                                 | intmax_t                     |

## Interoperability between Fortran and C types

(cont.)

| Fortran type | Named constant from the ISO_C_BINDING module<br>(kind type parameter if value is positive) | C type               |
|--------------|--------------------------------------------------------------------------------------------|----------------------|
|              | C_INTPTR_T                                                                                 | intptr_t             |
|              | C_PTRDIFF_T                                                                                | ptrdiff_t            |
|              | C_FLOAT                                                                                    | float                |
| REAL         | C_DOUBLE                                                                                   | double               |
|              | C_LONG_DOUBLE                                                                              | long double          |
|              | C_FLOAT_COMPLEX                                                                            | float _Complex       |
| COMPLEX      | C_DOUBLE_COMPLEX                                                                           | double _Complex      |
|              | C_LONG_DOUBLE_COMPLEX                                                                      | long double _Complex |
| LOGICAL      | C_BOOL                                                                                     | _Bool                |
| CHARACTER    | C_CHAR                                                                                     | char                 |

## NOTE 15.7

ISO/IEC 9899:2011 specifies that the representations for nonnegative signed integers are the same as the corresponding values of unsigned integers. Because Fortran does not provide direct support for unsigned kinds of integers, the ISO\_C\_BINDING module does not make accessible [named constants](#) for their [kind type parameter](#) values. A user can use the signed kinds of integers to interoperate with the unsigned types and all their qualified versions as well. This has the potentially surprising side effect that the C type unsigned char is [interoperable](#) with the type integer with a [kind type parameter](#) of C\_SIGNED\_CHAR.

## 15.3.3 Interoperability with C pointer types

- 1 C\_PTR and C\_FUNPTR shall be derived types with only private components. No [direct component](#) of either of these types is [allocatable](#) or a pointer. C\_PTR is [interoperable](#) with any C object pointer type. C\_FUNPTR is [interoperable](#) with any C function pointer type.

## NOTE 15.8

This means that only a C processor with the same representation method for all C object pointer types, and the same representation method for all C function pointer types, can be the target of interoperability of a Fortran processor. ISO/IEC 9899:2011 does not require this to be the case.

## NOTE 15.9

The function C\_LOC can be used to return a value of type C\_PTR that is the [C address](#) of an allocated [allocatable](#) variable. The function C\_FUNLOC can be used to return a value of type C\_FUNPTR that is the [C address](#) of a procedure. For C\_LOC and C\_FUNLOC the returned value is of an [interoperable](#) type and thus can be used in contexts where the procedure or [allocatable](#) variable is not directly allowed. For example, it could be passed as an [actual argument](#) to a C function.

Similarly, type C\_FUNPTR or C\_PTR can be used in a dummy argument or [structure component](#) and can have a value that is the [C address](#) of a procedure or [allocatable](#) variable, even in contexts where a procedure or [allocatable](#) variable is not directly allowed.

## 15.3.4 Interoperability of derived types and C struct types

- 1 Interoperability between a [derived type](#) in Fortran and a struct type in C is provided by the [BIND attribute](#) on the Fortran type.
- C1501 (R426) A [derived type](#) with the [BIND attribute](#) shall not have the [SEQUENCE attribute](#).

- 1 C1502 (R426) A **derived type** with the **BIND attribute** shall not have **type parameters**.
- 2 C1503 (R426) A **derived type** with the **BIND attribute** shall not have the **EXTENDS attribute**.
- 3 C1504 (R426) A *derived-type-def* that defines a **derived type** with the **BIND attribute** shall not have a *type-bound-procedure-part*.
- 4
- 5 C1505 (R426) A **derived type** with the **BIND attribute** shall have at least one component.
- 6 C1506 (R426) Each component of a **derived type** with the **BIND attribute** shall be a nonpointer, nonallocatable
- 7 data component with **interoperable** type and **type parameters**.

**NOTE 15.10**

The syntax rules and their constraints require that a **derived type** that is **interoperable** with a C struct type have components that are all data entities that are **interoperable**. No component is permitted to be **allocatable** or a pointer, but the value of a component of type **C\_FUNPTR** or **C\_PTR** can be the **C address** of such an entity.

- 8 2 A **derived type** is **interoperable** with a C struct type if and only if the **derived type** has the **BIND attribute** (4.5.2),
- 9 the **derived type** and the C struct type have the same number of components, and the components of the **derived**
- 10 **type** would interoperate with corresponding components of the C struct type as described in 15.3.5 and 15.3.6 if
- 11 the components were variables. A component of a **derived type** and a component of a C struct type correspond
- 12 if they are declared in the same relative position in their respective type definitions.

**NOTE 15.11**

The names of the corresponding components of the **derived type** and the C struct type need not be the same.

- 13 3 There is no Fortran type that is **interoperable** with a C struct type that contains a bit field or that contains a
- 14 flexible array member. There is no Fortran type that is **interoperable** with a C union type.

**NOTE 15.12**

For example, the C type myctype, declared below, is **interoperable** with the Fortran type myftype, declared below.

```
typedef struct {
    int m, n;
    float r;
} myctype;

USE, INTRINSIC :: ISO_C_BINDING
TYPE, BIND(C) :: MYFTYPE
    INTEGER(C_INT) :: I, J
    REAL(C_FLOAT) :: S
END TYPE MYFTYPE
```

The names of the types and the names of the components are not significant for the purposes of determining whether a Fortran derived type is **interoperable** with a C struct type.

**NOTE 15.13**

ISO/IEC 9899:2011 requires the names and component names to be the same in order for the types to be compatible (ISO/IEC 9899:2011, subclause 6.2.7). This is similar to Fortran's rule describing when different derived type definitions describe the same **sequence type**. This rule was not extended to determine whether

**NOTE 15.13 (cont.)**

a Fortran derived type is [interoperable](#) with a C struct type because the case of identifiers is significant in C but not in Fortran.

**15.3.5 Interoperability of scalar variables**

- 1 A named scalar Fortran variable is [interoperable](#) if and only if its type and type parameters are [interoperable](#), it is not a [coarray](#), it has neither the [ALLOCATABLE](#) nor the [POINTER](#) attribute, and if it is of type character its length is not assumed or declared by an expression that is not a [constant expression](#).
- 2 An [interoperable](#) scalar Fortran variable is [interoperable](#) with a scalar C entity if their types and type parameters are [interoperable](#).

**15.3.6 Interoperability of array variables**

- 1 A Fortran variable that is a named array is [interoperable](#) if and only if its type and type parameters are [interoperable](#), it is not a [coarray](#), it is of explicit shape or assumed size, and if it is of type character its length is not assumed or declared by an expression that is not a [constant expression](#).
- 2 An [explicit-shape](#) or [assumed-size](#) array of [rank](#)  $r$ , with a shape of  $[ e_1 \dots e_r ]$  is [interoperable](#) with a C array if its size is nonzero and
  - (1) either
    - (a) the array is [assumed-size](#), and the C array does not specify a size, or
    - (b) the array is an [explicit-shape array](#), and the extent of the last dimension ( $e_r$ ) is the same as the size of the C array, and
  - (2) either
    - (a)  $r$  is equal to one, and an element of the array is [interoperable](#) with an element of the C array, or
    - (b)  $r$  is greater than one, and an [explicit-shape array](#) with shape of  $[ e_1 \dots e_{r-1} ]$ , with the same type and type parameters as the original array, is [interoperable](#) with a C array of a type equal to the element type of the original C array.

**NOTE 15.14**

An element of a multi-dimensional C array is an array type, so a Fortran array of [rank](#) one is not [interoperable](#) with a multidimensional C array.

**NOTE 15.15**

An [allocatable](#) array or [array pointer](#) is never [interoperable](#). Such an array does not meet the requirement of being an [explicit-shape](#) or [assumed-size](#) array.

**NOTE 15.16**

For example, a Fortran array declared as

```
INTEGER(C_INT) :: A(18, 3:7, *)
```

is [interoperable](#) with a C array declared as

```
int b[][5][18];
```

**NOTE 15.17**

The C programming language defines null-terminated strings, which are actually arrays of the C type char that have a C null character in them to indicate the last valid element. A Fortran array of type character with a **kind type parameter** equal to C\_CHAR is **interoperable** with a C string.

Fortran's rules of sequence association (12.5.2.11) permit a character scalar **actual argument** to correspond to a dummy argument array. This makes it possible to argument associate a Fortran character string with a C string.

Note 15.21 has an example of interoperation between Fortran and C strings.

**15.3.7 Interoperability of procedures and procedure interfaces**

- 1 A Fortran procedure is **interoperable** if it has the **BIND attribute**, that is, if its **interface** is specified with a *proc-language-binding-spec*.
- 2 A Fortran procedure **interface** is **interoperable** with a C function prototype if
  - (1) the **interface** has the **BIND attribute**,
  - (2) either
    - (a) the **interface** describes a function whose **result** is a scalar variable that is **interoperable** with the result of the prototype or
    - (b) the **interface** describes a subroutine and the prototype has a result type of void,
  - (3) the number of **dummy arguments** of the **interface** is equal to the number of formal parameters of the prototype,
  - (4) any scalar dummy argument with the **VALUE attribute** is **interoperable** with the corresponding formal parameter of the prototype,
  - (5) any dummy argument without the **VALUE attribute** corresponds to a formal parameter of the prototype that is of a pointer type, and either
    - the dummy argument is **interoperable** with an entity of the referenced type (ISO/IEC 9899:2011, 6.2.5, 7.19, and 7.20.1) of the formal parameter,
    - the dummy argument is a nonallocatable nonpointer variable of type CHARACTER with assumed character length and the formal parameter is a pointer to CFI\_cdesc\_t,
    - the dummy argument is **allocatable**, **assumed-shape**, **assumed-rank**, or a **pointer** without the **CONTIGUOUS attribute**, and the formal parameter is a pointer to CFI\_cdesc\_t, or
    - the dummy argument is **assumed-type** and not **allocatable**, **assumed-shape**, **assumed-rank**, or a **pointer**, and the formal parameter is a pointer to void,
  - (6) each **allocatable** or **pointer** dummy argument of type CHARACTER has **deferred character length**, and
  - (7) the prototype does not have variable arguments as denoted by the ellipsis (...).

**NOTE 15.18**

The **referenced type** of a C pointer type is the C type of the object that the C pointer type points to. For example, the referenced type of the pointer type `int *` is `int`.

**NOTE 15.19**

The C language allows specification of a C function that can take a variable number of arguments (ISO/IEC 9899:2011, 7.16). This part of ISO/IEC 1539 does not provide a mechanism for Fortran procedures to interoperate with such C functions.

- 3 A formal parameter of a C function prototype corresponds to a **dummy argument** of a Fortran **interface** if they are in the same relative positions in the C parameter list and the **dummy argument** list, respectively.

- 1 4 In a reference from C to a Fortran procedure with an interoperable [interface](#), a C actual argument shall be the  
 2 address of a [C descriptor](#) for the intended [effective argument](#) if the corresponding [dummy argument](#) interoperates  
 3 with a C formal parameter that is a pointer to `CFI_cdesc_t`. In this [C descriptor](#), the members other than  
 4 `attribute` and `type` shall describe an object with the same characteristics as the intended [effective argument](#).  
 5 The value of the `attribute` member of the [C descriptor](#) shall be compatible with the characteristics of the [dummy](#)  
 6 [argument](#). The `type` member shall have a value that depends on the intended [effective argument](#) as follows:
- 7 • if the [dynamic type](#) of the intended [effective argument](#) is an interoperable type listed in Table 15.4, the  
 8 corresponding value for that type;
  - 9 • if the [dynamic type](#) of the intended [effective argument](#) is an [intrinsic type](#) for which the processor defines  
 10 a nonnegative type specifier value not listed in Table 15.4, that type specifier value;
  - 11 • otherwise, `CFI_type_other`.
- 12 5 When an interoperable Fortran procedure with a [simply contiguous dummy argument](#) is invoked from C and the  
 13 actual argument is the address of a [C descriptor](#) for a discontinuous object, the Fortran processor shall handle  
 14 the difference in contiguity.
- 15 6 When an interoperable C procedure whose Fortran interface has a [simply contiguous dummy argument](#) is invoked  
 16 from Fortran and the [effective argument](#) is discontinuous, the Fortran processor shall ensure that the C procedure  
 17 receives a descriptor for a contiguous object.
- 18 7 When an interoperable C procedure whose Fortran interface has a [simply contiguous dummy argument](#) is invoked  
 19 from C, and the actual argument is the address of a [C descriptor](#) for a discontinuous object, the C code within  
 20 the procedure shall be prepared to handle the discontinuous argument.
- 21 8 If an interoperable procedure defined by means other than Fortran has an optional [dummy argument](#), and the  
 22 corresponding [actual argument](#) in a reference from Fortran is absent, the procedure is invoked with a null pointer  
 23 for that argument. If an interoperable procedure defined by means of Fortran is invoked by a C function, an  
 24 optional [dummy argument](#) is absent if and only if the corresponding argument in the invocation is a null pointer.

**NOTE 15.20**

For example, a Fortran procedure [interface](#) described by

```

INTERFACE
  FUNCTION FUNC(I, J, K, L, M) BIND(C)
    USE, INTRINSIC :: ISO_C_BINDING
    INTEGER(C_SHORT) :: FUNC
    INTEGER(C_INT), VALUE :: I
    REAL(C_DOUBLE) :: J
    INTEGER(C_INT) :: K, L(10)
    TYPE(C_PTR), VALUE :: M
  END FUNCTION FUNC
END INTERFACE

```

is [interoperable](#) with the C function prototype

```
short func(int i, double *j, int *k, int l[10], void *m);
```

A C pointer can correspond to a Fortran dummy argument of type `C_PTR` with the [VALUE attribute](#) or to a Fortran scalar that does not have the [VALUE attribute](#). In the above example, the C pointers `j` and `k` correspond to the Fortran scalars `J` and `K`, respectively, and the C pointer `m` corresponds to the Fortran dummy argument `M` of type `C_PTR`.

**NOTE 15.21**

The interoperability of Fortran procedure [interfaces](#) with C function prototypes is only one part of invocation of a C function from Fortran. There are four pieces to consider in such an invocation: the procedure reference, the Fortran procedure [interface](#), the C function prototype, and the C function. Conversely, the invocation of a Fortran procedure from C involves the function reference, the C function prototype, the Fortran procedure [interface](#), and the Fortran procedure. In order to determine whether a reference is allowed, it is necessary to consider all four pieces.

For example, consider a C function that can be described by the C function prototype

```
void copy(char in[], char out[]);
```

Such a function can be invoked from Fortran as follows:

```
USE, INTRINSIC :: ISO_C_BINDING, ONLY: C_CHAR, C_NULL_CHAR
INTERFACE
  SUBROUTINE COPY(IN, OUT) BIND(C)
    IMPORT C_CHAR
    CHARACTER(KIND=C_CHAR), DIMENSION(*) :: IN, OUT
  END SUBROUTINE COPY
END INTERFACE

CHARACTER(LEN=10, KIND=C_CHAR) :: &
&    DIGIT_STRING = C_CHAR_'123456789' // C_NULL_CHAR
CHARACTER(KIND=C_CHAR) :: DIGIT_ARR(10)

CALL COPY(DIGIT_STRING, DIGIT_ARR)
PRINT '(1X, A1)', DIGIT_ARR(1:9)
END
```

The procedure reference has character string [actual arguments](#). These correspond to character array [dummy arguments](#) in the procedure [interface body](#) as allowed by Fortran's rules of sequence association ([12.5.2.11](#)). Those array [dummy arguments](#) in the procedure [interface](#) are [interoperable](#) with the formal parameters of the C function prototype. The C function is not shown here, but is assumed to be compatible with the C function prototype.

**NOTE 15.22**

The requirements on the Fortran processor when a discontinuous [actual argument](#) is being passed to a [simply contiguous dummy argument](#) are such that the C programmer has the same functionality (and safety) as the Fortran programmer, when either the calling or called procedure are Fortran procedures. The requirements on the C procedure, in the case when both the called and calling procedure are C procedures, effectively require that the C procedure not access memory through the descriptor that is not described by the descriptor; otherwise the program will not conform to this part of ISO/IEC 1539. A [dummy argument](#) is [simply contiguous](#) if it has the [CONTIGUOUS attribute](#) or is an array that is not [assumed-shape](#), [assumed-rank](#), or a [pointer](#).

## 1 15.4 C descriptors

- 2 1 A [C descriptor](#) is a C structure of type [CFI\\_cdesc\\_t](#). Together with library functions that have standard
- 3 prototypes, it provides a means for describing and manipulating Fortran data objects from within a C function.
- 4 This C structure is defined in the source file [ISO\\_Fortran\\_binding.h](#).

## 15.5 The source file ISO\_Fortran\_binding.h

### 15.5.1 Summary of contents

- 1 The source file ISO\_Fortran\_binding.h shall contain the C structure definitions, typedef declarations, macro definitions, and function prototypes specified in subclauses 15.5.2 to 15.5.5. The definitions and declarations in ISO\_Fortran\_binding.h can be used by a C function to interpret and manipulate a C descriptor. These provide a means to specify a C prototype that interoperates with a Fortran interface that has a non-interoperable dummy variable (15.3.7).
- 2 The source file ISO\_Fortran\_binding.h may be included in any order relative to the standard C headers, and may be included more than once in a given scope, with no effect different from being included only once, other than the effect on line numbers.
- 3 A C source file that includes the ISO\_Fortran\_binding.h header file shall not use any names starting with CFI\_ that are not defined in the header, and shall not define any of the structure names defined in the header as macro names. All names other than structure member names defined in the header begin with CFI\_ or an underscore character, or are defined by a standard C header that it includes.

### 15.5.2 The CFI\_dim\_t structure type

- 1 CFI\_dim\_t is a typedef name for a C structure. It is used to represent lower bound, extent, and memory stride information for one dimension of an array. The type CFI\_index\_t is described in 15.5.4. CFI\_dim\_t contains at least the following members in any order.

**CFI\_index\_t lower\_bound;** The value is equal to the value of the lower bound for the dimension being described.

**CFI\_index\_t extent;** The value is equal to the number of elements in the dimension being described, or -1 for the final dimension of an assumed-size array.

**CFI\_index\_t sm;** The value is equal to the memory stride for a dimension; this is the difference in bytes between the addresses of successive elements in the dimension being described.

### 15.5.3 The CFI\_cdesc\_t structure type

- 1 CFI\_cdesc\_t is a typedef name for a C structure, which contains a flexible array member. It shall contain at least the members described in this subclause. The values of these members of a structure of type CFI\_cdesc\_t that is produced by the functions and macros specified in this part of ISO/IEC 1539, or received by a C function when invoked by a Fortran procedure, shall have the properties described in this subclause.

- 2 The first three members of the structure shall be **base\_addr**, **elem\_len**, and **version** in that order. The final member shall be **dim**. All other members shall be between **version** and **dim**, in any order. The types CFI\_attribute\_t, CFI\_rank\_t, and CFI\_type\_t are described in 15.5.4. The type CFI\_dim\_t is described in 15.5.2.

**void \* base\_addr;** If the object is an unallocated allocatable variable or a pointer that is disassociated, the value is a null pointer. If the object has zero size, the value is not a null pointer but is otherwise processor-dependent. Otherwise, the value is the base address of the object being described. The base address of a scalar is its C address. The base address of an array is the C address of the first element in Fortran array element order.

**size\_t elem\_len;** If the object is scalar, the value is the storage size in bytes of the object; otherwise, the value is the storage size in bytes of an element of the object.

**int version;** The value is equal to the value of CFI\_VERSION in the source file ISO\_Fortran\_binding.h that defined the format and meaning of this C descriptor when the descriptor was established.

- 1 **CFI\_rank\_t rank**; The value is equal to the number of dimensions of the Fortran object being described; if  
2 the object is scalar, the value is zero.
- 3 **CFI\_type\_t type**; The value is equal to the specifier for the type of the object. Each interoperable intrinsic C  
4 type has a specifier. Specifiers are also provided to indicate that the type of the object is an interoperable  
5 structure, or is unknown. The macros listed in Table 15.4 provide values that correspond to each specifier.
- 6 **CFI\_attribute\_t attribute**; The value is equal to the value of an attribute code that indicates whether the  
7 object described is [allocatable](#), a [data pointer](#), or a nonallocatable nonpointer data object. The macros  
8 listed in Table 15.3 provide values that correspond to each code.
- 9 **CFI\_dim\_t dim**; The number of elements in the **dim** array is equal to the rank of the object. Each element of  
10 the array contains the lower bound, extent, and memory stride information for the corresponding dimension  
11 of the Fortran object.
- 12 3 For a [C descriptor](#) of an array pointer or [allocatable](#) array, the value of the **lower\_bound** member of each element  
13 of the **dim** member of the descriptor is determined by argument association, allocation, or pointer association.  
14 For a [C descriptor](#) of a nonallocatable nonpointer object, the value of the **lower\_bound** member of each element  
15 of the **dim** member of the descriptor is zero.
- 16 4 There shall be an ordering of the dimensions such that the absolute value of the **sm** member of the first dimension  
17 is not less than the **elem\_len** member of the [C descriptor](#) and the absolute value of the **sm** member of each  
18 subsequent dimension is not less than the absolute value of the **sm** member of the previous dimension multiplied  
19 by the extent of the previous dimension.
- 20 5 In a [C descriptor](#) of an [assumed-size array](#), the **extent** member of the last element of the **dim** member has the  
21 value  $-1$ .

**NOTE 15.23**

The reason for the restriction on the absolute values of the **sm** members is to ensure that there is no overlap between the elements of the array that is being described, while allowing for the reordering of subscripts. Within Fortran, such a reordering can be achieved with the intrinsic function TRANSPOSE or the intrinsic function RESHAPE with the optional argument ORDER, and an optimizing compiler can accommodate it without making a copy by constructing the appropriate descriptor whenever it can determine that a copy is not needed.

**NOTE 15.24**

The value of **elem\_len** for a Fortran CHARACTER object is equal to the character length times the number of bytes of a single character of that kind. If the kind is C\_CHAR, this value will be equal to the character length.

## 22 15.5.4 Macros and typedefs in ISO\_Fortran\_binding.h

- 23 1 Except for CFI\_CDESC\_T, each macro defined in ISO\_Fortran\_binding.h expands to an integer constant  
24 expression that is either a single token or a parenthesized expression that is suitable for use in #if preprocessing  
25 directives.
- 26 2 CFI\_CDESC\_T is a function-like macro that takes one argument, which is the rank of the [C descriptor](#) to create,  
27 and evaluates to an unqualified type of suitable size and alignment for defining a variable to use as a [C descriptor](#)  
28 of that rank. The argument shall be an integer constant expression with a value that is greater than or equal to  
29 zero and less than or equal to CFI\_MAX\_RANK. A pointer to a variable declared using CFI\_CDESC\_T can  
30 be cast to [CFI\\_cdesc\\_t](#) \*. A variable declared using CFI\_CDESC\_T shall not have an initializer.

**NOTE 15.25**

The CFI\_CDESC\_T macro provides the memory for a C descriptor. The address of an entity declared using the macro is not usable as an actual argument corresponding to a formal parameter of type CFI\_cdesc\_t \* without an explicit cast. For example, the following code uses CFI\_CDESC\_T to declare a C descriptor of rank 5 and pass it to CFI\_deallocate (15.5.5.4).

```
CFI_CDESC_T(5) object;
int ind;
... code to define and use C descriptor ...
ind = CFI_deallocate((CFI_cdesc_t *)&object);
```

- 1 3 CFI\_index\_t is a typedef name for a standard signed integer type capable of representing the result of subtracting  
2 two pointers.
- 3 4 The CFI\_MAX\_RANK macro has a processor-dependent value equal to the largest rank supported. The value  
4 shall be greater than or equal to 15. CFI\_rank\_t is a typedef name for a standard integer type capable of  
5 representing the largest supported rank.
- 6 5 The CFI\_VERSION macro has a processor-dependent value that encodes the version of the ISO\_Fortran\_  
7 binding.h source file containing this macro. This value should be increased if a new version of the source file is  
8 incompatible with the previous version.
- 9 6 The macros in Table 15.3 are for use as attribute codes. The values shall be nonnegative and distinct. CFI\_  
10 attribute\_t is a typedef name for a standard integer type capable of representing the values of the attribute  
11 codes.

Table 15.3: ISO\_Fortran\_binding.h macros for attribute codes

| Macro name                | Attribute                 |
|---------------------------|---------------------------|
| CFI_attribute_pointer     | data pointer              |
| CFI_attribute_allocatable | allocatable               |
| CFI_attribute_other       | nonallocatable nonpointer |

- 12 7 CFI\_attribute\_pointer specifies a data object with the Fortran POINTER attribute. CFI\_attribute\_allocatable  
13 specifies an object with the Fortran ALLOCATABLE attribute. CFI\_attribute\_other specifies a nonallocatable  
14 nonpointer object.
- 15 8 The macros in Table 15.4 are for use as type specifiers. The value for CFI\_type\_other shall be negative and  
16 distinct from all other type specifiers. CFI\_type\_struct specifies a C structure that is interoperable with a  
17 Fortran derived type; its value shall be positive and distinct from all other type specifiers. If a C type is not  
18 interoperable with a Fortran type and kind supported by the Fortran processor, its macro shall evaluate to a  
19 negative value. Otherwise, the value for an intrinsic type shall be positive.
- 20 9 If the processor supports interoperability of a Fortran intrinsic type with a C type not listed in Table 15.4,  
21 the processor shall define a type specifier value for that type which is positive and distinct from all other type  
22 specifiers.
- 23 10 CFI\_type\_t is a typedef name for a standard integer type capable of representing the values for the supported  
24 type specifiers.

Table 15.4: ISO\_Fortran\_binding.h macros for type codes

| Macro name           | C Type      |
|----------------------|-------------|
| CFI_type_signed_char | signed char |
| CFI_type_short       | short int   |

## ISO\_Fortran\_binding.h macros for type codes (cont.)

| Macro name                   | C Type                    |
|------------------------------|---------------------------|
| CFI_type_int                 | int                       |
| CFI_type_long                | long int                  |
| CFI_type_long_long           | long long int             |
| CFI_type_size_t              | size_t                    |
| CFI_type_int8_t              | int8_t                    |
| CFI_type_int16_t             | int16_t                   |
| CFI_type_int32_t             | int32_t                   |
| CFI_type_int64_t             | int64_t                   |
| CFI_type_int_least8_t        | int_least8_t              |
| CFI_type_int_least16_t       | int_least16_t             |
| CFI_type_int_least32_t       | int_least32_t             |
| CFI_type_int_least64_t       | int_least64_t             |
| CFI_type_int_fast8_t         | int_fast8_t               |
| CFI_type_int_fast16_t        | int_fast16_t              |
| CFI_type_int_fast32_t        | int_fast32_t              |
| CFI_type_int_fast64_t        | int_fast64_t              |
| CFI_type_intmax_t            | intmax_t                  |
| CFI_type_intptr_t            | intptr_t                  |
| CFI_type_ptrdiff_t           | ptrdiff_t                 |
| CFI_type_float               | float                     |
| CFI_type_double              | double                    |
| CFI_type_long_double         | long double               |
| CFI_type_float_Complex       | float _Complex            |
| CFI_type_double_Complex      | double _Complex           |
| CFI_type_long_double_Complex | long double _Complex      |
| CFI_type_Bool                | _Bool                     |
| CFI_type_char                | char                      |
| CFI_type_cptr                | void *                    |
| CFI_type_struct              | interoperable C structure |
| CFI_type_other               | Not otherwise specified   |

**NOTE 15.26**

The values for different C types can be the same; for example, CFI\_type\_int and CFI\_type\_int32\_t might have the same value.

- 11 The macros in Table 15.5 are for use as error codes. The macro CFI\_SUCCESS shall be defined to be the integer constant 0. The value of each macro other than CFI\_SUCCESS shall be nonzero and shall be different from the values of the other macros specified in this subclause. Error conditions other than those listed in this subclause should be indicated by error codes different from the values of the macros named in this subclause.
- 12 The values of the macros in Table 15.5 indicate the error condition described.

Table 15.5: ISO\_Fortran\_binding.h macros for error codes

| Macro name                   | Error condition                                                                                                                  |
|------------------------------|----------------------------------------------------------------------------------------------------------------------------------|
| CFI_SUCCESS                  | No error detected.                                                                                                               |
| CFI_ERROR_BASE_ADDR_NULL     | The base address member of a <a href="#">C descriptor</a> is a null pointer in a context that requires a non-null pointer value. |
| CFI_ERROR_BASE_ADDR_NOT_NULL | The base address member of a <a href="#">C descriptor</a> is not a null pointer in a context that requires a null pointer value. |
| CFI_INVALID_ELEM_LEN         | The value supplied for the element length member of a <a href="#">C descriptor</a> is not valid.                                 |

## ISO\_Fortran\_binding.h macros for error codes

(cont.)

| Macro name               | Error condition                                                                             |
|--------------------------|---------------------------------------------------------------------------------------------|
| CFI_INVALID_RANK         | The value supplied for the rank member of a <a href="#">C descriptor</a> is not valid.      |
| CFI_INVALID_TYPE         | The value supplied for the type member of a <a href="#">C descriptor</a> is not valid.      |
| CFI_INVALID_ATTRIBUTE    | The value supplied for the attribute member of a <a href="#">C descriptor</a> is not valid. |
| CFI_INVALID_EXTENT       | The value supplied for the extent member of a CFI_dim_t structure is not valid.             |
| CFI_INVALID_DESCRIPTOR   | A <a href="#">C descriptor</a> is invalid in some way.                                      |
| CFI_ERROR_MEM_ALLOCATION | Memory allocation failed.                                                                   |
| CFI_ERROR_OUT_OF_BOUNDS  | A reference is out of bounds.                                                               |

## 15.5.5 Functions declared in ISO\_Fortran\_binding.h

## 15.5.5.1 Arguments and results of the functions

- 1 Some of the functions described in [15.5.5](#) return an error indicator; this is an integer value that indicates whether an error condition was detected. The value zero indicates that no error condition was detected, and a nonzero value indicates which error condition was detected. Table [15.5](#) lists standard error conditions and macro names for their corresponding error codes. A processor is permitted to detect other error conditions. If an invocation of a function defined in [15.5.5](#) could detect more than one error condition and an error condition is detected, which error condition is detected is processor dependent.
- 2 In function arguments representing subscripts, bounds, extents, or strides, the ordering of the elements is the same as the ordering of the elements of the `dim` member of a [C descriptor](#).
- 3 Prototypes for these functions, or equivalent macros, are provided in the `ISO_Fortran_binding.h` file as described in [15.5.5](#). It is unspecified whether the functions defined by this header are macros or identifiers declared with external linkage. If a macro definition is suppressed in order to access an actual function, the behavior is undefined.

**NOTE 15.27**

These functions are allowed to be macros to provide extra implementation flexibility. For example, `CFI_`-`establish` could include the value of `CFI_VERSION` in the header used to compile the call to `CFI_``establish` as an extra argument of the actual function used to establish the [C descriptor](#).

## 15.5.5.2 The CFI\_address function

- 1 **Synopsis.** C address of an object described by a [C descriptor](#).

```
void *CFI_address(const CFI\_cdesc\_t *dv, const CFI_index_t subscripts[]);
```

- 2 **Formal Parameters.**

`dv` shall be the address of a [C descriptor](#) describing the object. The object shall not be an unallocated allocatable variable or a pointer that is not associated.

`subscripts` shall be a null pointer or the address of an array of type `CFI_index_t`. If the object is an array, `subscripts` shall be the address of an array of `CFI_index_t` with at least  $n$  elements, where  $n$  is the rank of the object. The value of `subscripts[i]` shall be within the bounds of dimension  $i$  specified by the `dim` member of the [C descriptor](#).

- 3 **Result Value.** If the object is an array of rank  $n$ , the result is the [C address](#) of the element of the object that

1 the first  $n$  elements of the **subscripts** argument would specify if used as subscripts. If the object is scalar, the  
 2 result is its **C address**.

3 4 **Example.** If **dv** is the address of a **C descriptor** for the Fortran array **A** declared as

```
4 REAL(C_FLOAT) :: A(100, 100)
```

5 the following code calculates the **C address** of **A(5, 10)**:

```
6 CFI_index_t subscripts[2];
7 float *address;
8 subscripts[0] = 4;
9 subscripts[1] = 9;
10 address = (float *) CFI_address(dv, subscripts );
```

### 11 15.5.5.3 The CFI\_allocate function

12 1 **Synopsis.** Allocate memory for an object described by a **C descriptor**.

```
13 int CFI_allocate(CFI_cdesc_t *dv, const CFI_index_t lower_bounds[],
14                 const CFI_index_t upper_bounds[], size_t elem_len);
```

15 2 **Formal Parameters.**

16 **dv** shall be the address of a **C descriptor** specifying the rank and type of the object. The **base\_**  
 17 **addr** member of the **C descriptor** shall be a null pointer. If the type is not a character type, the  
 18 **elem\_len** member shall specify the element length. The **attribute** member shall have a value of  
 19 **CFI\_attribute\_allocatable** or **CFI\_attribute\_pointer**.

20 **lower\_bounds** shall be the address of an array with at least **dv->rank** elements, if **dv->rank**>0.

21 **upper\_bounds** shall be the address of an array with at least **dv->rank** elements, if **dv->rank**>0.

22 **elem\_len** If the type specified in the **C descriptor** type is a Fortran character type, the value of **elem\_len**  
 23 shall be the storage size in bytes of an element of the object; otherwise, **elem\_len** is ignored.

24 3 **Description.** Successful execution of **CFI\_allocate** allocates memory for the object described by the **C**  
 25 **descriptor** with the address **dv** using the same mechanism as the Fortran **ALLOCATE statement**, and assigns the  
 26 address of that memory to **dv->base\_addr**. The first **dv->rank** elements of the **lower\_bounds** and **upper\_bounds**  
 27 arguments provide the lower and upper Fortran bounds, respectively, for each corresponding dimension of the  
 28 object. The supplied lower and upper bounds override any current dimension information in the **C descriptor**.  
 29 If the rank is zero, the **lower\_bounds** and **upper\_bounds** arguments are ignored. If the type specified in the **C**  
 30 **descriptor** is a character type, the supplied element length overrides the current element-length information in  
 31 the descriptor.

32 If an error is detected, the **C descriptor** is not modified.

33 4 **Result Value.** The result is an error indicator.

34 5 **Example.** If **dv** is the address of a **C descriptor** for the Fortran array **A** declared as

```
35 REAL, ALLOCATABLE :: A(:, :)
```

36 and the array is not allocated, the following code allocates it to be of shape [100, 500]:

```
37 CFI_index_t lower[2], upper[2];
38 int ind;
39 lower[0] = 1; lower[1] = 1;
40 upper[0] = 100; upper[1] = 500;
41 ind = CFI_allocate(dv, lower, upper, 0);
```

#### 15.5.5.4 The CFI\_deallocate function

**Synopsis.** Deallocate memory for an object described by a [C descriptor](#).

```
int CFI_deallocate(CFI_cdesc_t *dv);
```

**Formal Parameter.** `dv` shall be the address of a [C descriptor](#) describing the object. It shall have been allocated using the same mechanism as the Fortran [ALLOCATE statement](#). If the object is a pointer, it shall be associated with a target satisfying the conditions for successful deallocation by the Fortran [DEALLOCATE statement \(6.7.3\)](#).

**Description.** Successful execution of `CFI_deallocate` deallocates memory for the object using the same mechanism as the Fortran [DEALLOCATE statement](#), and the `base_addr` member of the [C descriptor](#) becomes a null pointer.

If an error is detected, the [C descriptor](#) is not modified.

**Result Value.** The result is an error indicator.

**Example.** If `dv` is the address of a [C descriptor](#) for the Fortran array `A` declared as

```
REAL, ALLOCATABLE :: A(:, :)
```

and the array is allocated, the following code deallocates it:

```
int ind;
ind = CFI_deallocate(dv);
```

#### 15.5.5.5 The CFI\_establish function

**Synopsis.** Establish a [C descriptor](#).

```
int CFI_establish(CFI_cdesc_t *dv, void *base_addr, CFI_attribute_t attribute,
                  CFI_type_t type, size_t elem_len, CFI_rank_t rank,
                  const CFI_index_t extents[]);
```

**Formal Parameters.**

`dv` shall be the address of a data object large enough to hold a [C descriptor](#) of the rank specified by `rank`. It shall not have the same value as either a C formal parameter that corresponds to a Fortran actual argument or a C actual argument that corresponds to a Fortran dummy argument. It shall not be the address of a [C descriptor](#) that describes an allocated allocatable object.

`base_addr` shall be a null pointer or the base address of the object to be described. If it is not a null pointer, it shall be the address of a contiguous storage sequence that is appropriately aligned (ISO/IEC 9899:2011 3.2) for an object of the type specified by `type`.

`attribute` shall be one of the attribute codes in Table 15.3. If it is `CFI_attribute_allocatable`, `base_addr` shall be a null pointer.

`type` shall be one of the type codes in Table 15.4.

`elem_len` If the type is equal to `CFI_type_struct`, `CFI_type_other`, or a Fortran character type code, `elem_len` shall be greater than zero and equal to the storage size in bytes of an element of the object. Otherwise, `type` will be ignored.

`rank` shall have a value in the range  $0 \leq \text{rank} \leq \text{CFI\_MAX\_RANK}$ . It specifies the rank of the object.

`extents` is ignored if `rank` is equal to zero or if `base_addr` is a null pointer. Otherwise, it shall be the address of an array with `rank` elements; the value of each element shall be nonnegative, and `extents[i]` specifies the extent of dimension *i* of the object.

3 **Description.** Successful execution of CFI\_establish updates the object with the address `dv` to be an established C descriptor for a nonallocatable nonpointer data object of known shape, an unallocated allocatable object, or a data pointer. If `base_addr` is not a null pointer, it is for a nonallocatable entity that is a scalar or a contiguous array; if the `attribute` argument has the value CFI\_attribute\_pointer, the lower bounds of the object described by `dv` are set to zero. If `base_addr` is a null pointer, the established C descriptor is for an unallocated allocatable, a disassociated pointer, or is a C descriptor that has the `attribute` CFI\_attribute\_other but does not describe a data object. If `base_addr` is the C address of a Fortran data object, the `type` and `elem_len` arguments shall be consistent with the type and type parameters of the Fortran data object. The remaining properties of the object are given by the other arguments.

10 If an error is detected, the object with the address `dv` is not modified.

11 4 **Result Value.** The result is an error indicator.

#### NOTE 15.28

CFI\_establish is used to initialize a C descriptor declared in C with CFI\_CDESC\_T before passing it to any other functions as an actual argument, in order to set the rank, attribute, type and element length.

#### NOTE 15.29

A C descriptor with `attribute` CFI\_attribute\_other and `base_addr` a null pointer can be used as the argument `result` in calls to CFI\_section or CFI\_select\_part, which will produce a C descriptor for a nonallocatable nonpointer data object.

#### 12 5 Examples.

13 *Case (i):* The following code fragment establishes a C descriptor for an unallocated rank-one allocatable array that can be passed to Fortran for allocation there.

```
15     CFI_rank_t rank;
16     CFI_CDESC_T(1) field;
17     int ind;
18     rank = 1;
19     ind = CFI_establish((CFI_cdesc_t *)&field, NULL, CFI_attribute_allocatable,
20                       CFI_type_double, 0, rank, NULL);
```

21 *Case (ii):* Given the Fortran type definition

```
22     TYPE, BIND(C) :: T
23         REAL(C_DOUBLE) :: X
24         COMPLEX(C_DOUBLE_COMPLEX) :: Y
25     END TYPE
```

26 and a Fortran subprogram that has an assumed-shape dummy argument of type T, the following code fragment creates a descriptor `a_fortran` for an array of size 100 that can be used as the actual argument in an invocation of the subprogram from C:

```
29     typedef struct {double x; double _Complex y;} t;
30     t a_c[100];
31     CFI_CDESC_T(1) a_fortran;
32     int ind;
33     CFI_index_t extent[1];

34
35     extent[0] = 100;
36     ind = CFI_establish((CFI_cdesc_t *)&a_fortran, a_c, CFI_attribute_other,
37                       CFI_type_struct, sizeof(t), 1, extent);
```

### 15.5.5.6 The CFI\_is\_contiguous function

1 **Synopsis.** Test contiguity of an array.

```
2 int CFI_is_contiguous(const CFI_cdesc_t * dv);
```

2 **Formal Parameter.** `dv` shall be the address of a **C descriptor** describing an array. The `base_addr` member of the **C descriptor** shall not be a null pointer.

3 **Result Value.** The value of the result is 1 if the array described by `dv` is contiguous, and 0 otherwise.

#### NOTE 15.30

**Assumed-size** and **allocatable** arrays are always contiguous, and therefore the result of `CFI_is_contiguous` on a **C descriptor** for such an array will be equal to 1.

### 15.5.5.7 The CFI\_section function

1 **Synopsis.** Update a **C descriptor** for an array section for which each element is an element of a given array.

```
2 int CFI_section(CFI_cdesc_t *result, const CFI_cdesc_t *source,
3               const CFI_index_t lower_bounds[], const CFI_index_t upper_bounds[],
4               const CFI_index_t strides[]);
```

2 **Formal Parameters.**

**result** shall be the address of a **C descriptor** with rank equal to the rank of **source** minus the number of zero strides. The **attribute** member shall have the value `CFI_attribute_other` or `CFI_attribute_pointer`. If the value of **result** is the same as either a C formal parameter that corresponds to a Fortran actual argument or a C actual argument that corresponds to a Fortran dummy argument, the **attribute** member shall have the value `CFI_attribute_pointer`.

Successful execution of `CFI_section` updates the **base\_addr** and **dim** members of the **C descriptor** with the address **result** to describe the array section determined by **source**, **lower\_bounds**, **upper\_bounds**, and **strides**, as follows.

The array section is equivalent to the Fortran array section `SOURCE(sectsub1, sectsub2, ... sectsubn)`, where `SOURCE` is the array described by **source**,  $n$  is the rank of that array, and `sectsubi` is the subscript  $lower_i$  if  $stride_i$  is zero, and the section subscript  $lower_i : upper_i : stride_i$  otherwise. The value of  $lower_i$  is the lower bound of dimension  $i$  of `SOURCE` if **lower\_bounds** is a null pointer and **lower\_bounds**[ $i$ ] otherwise. The value of  $upper_i$  is the upper bound of dimension  $i$  of `SOURCE` if **upper\_bounds** is a null pointer and **upper\_bounds**[ $i$ ] otherwise. The value of  $stride_i$  is 1 if **strides** is a null pointer and **strides**[ $i$ ] otherwise. If  $stride_i$  has the value zero,  $lower_i$  shall have the same value as  $upper_i$ .

If an error is detected, the **C descriptor** with the address **result** is not modified.

**source** shall be the address of a **C descriptor** that describes a nonallocatable nonpointer array, an allocated allocatable array, or an associated array pointer. The **elem\_len** and **type** members of **source** shall have the same values as the corresponding members of **result**.

**lower\_bounds** shall be a null pointer or the address of an array with at least **source**->**rank** elements. If it is not a null pointer, and  $stride_i$  is zero or  $(upper_i - lower\_bounds[i] + stride_i)/stride_i > 0$ , the value of **lower\_bounds**[ $i$ ] shall be within the bounds of dimension  $i$  of `SOURCE`.

**upper\_bounds** shall be a null pointer or the address of an array with at least **source**->**rank** elements. If **source** describes an **assumed-size array**, **upper\_bounds** shall not be a null pointer. If it is not a null pointer and  $stride_i$  is zero or  $(upper\_bounds[i] - lower_i + stride_i)/stride_i > 0$ , the value of **upper\_bounds**[ $i$ ] shall be within the bounds of dimension  $i$  of `SOURCE`.

**strides** shall be a null pointer or the address of an array with at least **source**->**rank** elements.

1 3 **Result Value.** The result is an error indicator.

2 4 **Examples.**

3 *Case (i):* If **source** is already the address of a **C descriptor** for the rank-one Fortran array **A**, the lower  
 4 bounds of **A** are equal to 1, and the lower bounds in the **C descriptor** are equal to 0, the following  
 5 code fragment establishes a new **C descriptor** section and updates it to describe the array section  
 6 **A(3::5)**:

```

7     CFI_index_t lower[1], strides[1];
8     CFI_CDESC_T(1) section;
9     int ind;
10    lower[0] = 2;
11    strides[0] = 5;
12    ind = CFI_establish((CFI_cdesc_t *)&section, NULL, CFI_attribute_other,
13                      CFI_type_float, 0, 1, NULL);
14    ind = CFI_section((CFI_cdesc_t *)&section, source, lower, NULL, strides);
```

15 *Case (ii):* If **source** is already the address of a **C descriptor** for a rank-two Fortran assumed-shape array **A**  
 16 with lower bounds equal to 1, the following code fragment establishes a **C descriptor** and updates  
 17 it to describe the rank-one array section **A(:, 42)**.

```

18    CFI_index_t lower[2], upper[2], strides[2];
19    CFI_CDESC_T(1) section;
20    int ind;
21    lower[0] = source->dim[0].lower_bound;
22    upper[0] = source->dim[0].lower_bound + source->dim[0].extent - 1;
23    strides[0] = 1;
24    lower[1] = upper[1] = source->dim[1].lower_bound + 41;
25    strides[1] = 0;
26    ind = CFI_establish((CFI_cdesc_t *)&section, NULL, CFI_attribute_other,
27                      CFI_type_float, 0, 1, NULL);
28    ind = CFI_section((CFI_cdesc_t *)&section, source, lower, upper, strides);
```

#### 29 15.5.5.8 The CFI\_select\_part function

30 1 **Synopsis.** Update a **C descriptor** for an array section for which each element is a part of the corresponding  
 31 element of an array.

```

32    int CFI_select_part(CFI_cdesc_t *result, const CFI_cdesc_t *source, size_t displacement,
33                      size_t elem_len);
```

#### 34 2 **Formal Parameters.**

35 **result** shall be the address of a **C descriptor**; **result->rank** shall have the same value as **source->rank**  
 36 and **result->attribute** shall have the value **CFI\_attribute\_other** or **CFI\_attribute\_pointer**. If  
 37 the address specified by **result** is the value of a C formal parameter that corresponds to a For-  
 38 tran actual argument or of a C actual argument that corresponds to a Fortran dummy argument,  
 39 **result->attribute** shall have the value **CFI\_attribute\_pointer**. The value of **result->type** spe-  
 40 cifies the type of the array section.

41 **source** shall be the address of a **C descriptor** for a nonallocatable nonpointer array, an allocated allocatable  
 42 array, or an associated array pointer.

43 **displacement** shall have a value  $0 \leq \text{displacement} \leq \text{source->elem\_len} - 1$ , and the sum of the displacement  
 44 and the size in bytes of an element of the array section shall be less than or equal to **source->elem\_-**  
 45 **len**. The address **displacement** bytes greater than the value of **source->base\_addr** is the base of

the array section and shall be appropriately aligned (ISO/IEC 9899:2011, 3.2) for an object of the type of the array section.

**elem\_len** shall have a value equal to the storage size in bytes of an element of the array section if **result->type** specifies a Fortran character type; otherwise, **elem\_len** is ignored.

**3 Description.** Successful execution of **CFI\_select\_part** updates the **base\_addr**, **dim**, and **elem\_len** members of the **C descriptor** with the address **result** for an array section for which each element is a part of the corresponding element of the array described by the **C descriptor** with the address **source**. The part shall be a component of a structure, a substring, or the real or imaginary part of a complex value.

If an error is detected, the **C descriptor** with the address **result** is not modified.

**4 Result Value.** The result is an error indicator.

**5 Example.** If **source** is already the address of a **C descriptor** for the Fortran array **A** declared with

```
TYPE, BIND(C) :: T
  REAL(C_DOUBLE) :: X
  COMPLEX(C_DOUBLE_COMPLEX) :: Y
END TYPE
TYPE(T) A(100)
```

the following code fragment establishes a **C descriptor** for the array **A%Y**:

```
typedef struct {
  double x; double _Complex y;
} t;
CFI_CDESC_T(1) component;
CFI_cdesc_t * comp_cdesc = (CFI_cdesc_t *)&component;
CFI_index_t extent[] = { 100 };
(void)CFI_establish(comp_cdesc, NULL, CFI_attribute_other, CFI_type_double_Complex,
  sizeof(double _Complex), 1, extent);
(void)CFI_select_part(comp_cdesc, source, offsetof(t,y), 0);
```

#### 15.5.5.9 The **CFI\_setpointer** function

**1 Synopsis.** Update a **C descriptor** for a Fortran pointer to be associated with the whole of a given object or to be disassociated.

```
int CFI_setpointer(CFI_cdesc_t *result, CFI_cdesc_t *source,
  const CFI_index_t lower_bounds[]);
```

#### 2 Formal Parameters.

**result** shall be the address of a **C descriptor** for a Fortran pointer. It is updated using information from the **source** and **lower\_bounds** arguments.

**source** shall be a null pointer or the address of a **C descriptor** for a nonallocatable nonpointer data object, an allocated allocatable object, or a data pointer object. If **source** is not a null pointer, the corresponding values of the **elem\_len**, **rank**, and **type** members shall be the same in the **C descriptors** with the addresses **source** and **result**.

**lower\_bounds** If **source** is not a null pointer and **source->rank** is nonzero, **lower\_bounds** shall be a null pointer or the address of an array with at least **source->rank** elements.

1 3 **Result Value.** The result is an error indicator.

2 4 **Description.** Successful execution of `CFI_setpointer` updates the `base_addr` and `dim` members of the `C descriptor` with the address `result` as follows:

- 4 • if `source` is a null pointer or the address of a `C descriptor` for a disassociated pointer, the updated `C descriptor` describes a disassociated pointer;
- 5 • otherwise, the `C descriptor` with the address `result` becomes a `C descriptor` for the object described by the `C descriptor` with the address `source`, except that if `source->rank` is nonzero and `lower_bounds` is not a null pointer, the lower bounds are replaced by the values of the first `source->rank` elements of the `lower_bounds` array.

10 If an error is detected, the `C descriptor` with the address `result` is not modified.

11 5 **Example.** If `ptr` is already the address of a `C descriptor` for an array pointer of rank 1, the following code updates it to be a `C descriptor` for a pointer to the same array with lower bound 0.

```
13     CFI_index_t lower_bounds[1];
14     int ind;
15     lower_bounds[0] = 0;
16     ind = CFI_setpointer(ptr, ptr, lower_bounds);
```

## 17 15.6 Restrictions on C descriptors

18 1 A `C descriptor` shall not be initialized, updated, or copied other than by calling the functions specified in 15.5.5.

19 2 If the address of a `C descriptor` is a formal parameter that corresponds to a Fortran actual argument or a C actual argument that corresponds to a Fortran dummy argument,

- 21 • the `C descriptor` shall not be modified if either the corresponding dummy argument in the Fortran interface has the `INTENT (IN) attribute` or the `C descriptor` is for a nonallocatable nonpointer object, and
- 23 • the `base_addr` member of the `C descriptor` shall not be accessed before it is given a value if the corresponding dummy argument in the Fortran interface has the `POINTER` and `INTENT (OUT)` attributes.

### NOTE 15.31

In this context, modification refers to any change to the location or contents of the `C descriptor`, including establishment and update. The intent of these restrictions is that `C descriptors` remain intact at all times they are accessible to an active Fortran procedure, so that the Fortran code is not required to copy them.

## 25 15.7 Restrictions on formal parameters

26 1 Within a C function, an `allocatable` object shall be allocated or deallocated only by execution of the `CFI_allocate` and `CFI_deallocate` functions. A Fortran pointer can become associated with a target by execution of the `CFI_allocate` function.

29 2 Calling `CFI_allocate` or `CFI_deallocate` for a `C descriptor` changes the allocation status of the Fortran variable it describes.

31 3 If the address of an object is the value of a formal parameter that corresponds to a nonpointer dummy argument in an interface with the `BIND attribute`, then

- 33 • if the dummy argument has the `INTENT (IN) attribute`, the object shall not be defined or become undefined, and

- if the dummy argument has the **INTENT (OUT) attribute**, the object shall not be referenced before it is defined.

- 4 If a formal parameter that is a pointer to **CFI\_cdesc\_t** corresponds to a dummy argument in an interoperable procedure interface, a pointer based on the **base\_addr** in that **C descriptor** shall not be used to access memory that is not part of the object described by the **C descriptor**.

## 15.8 Restrictions on lifetimes

- 1 A **C descriptor** of, or C pointer to, any part of a Fortran object becomes undefined under the same conditions that the association status of a Fortran pointer associated with that object would become undefined, and any further use of it is undefined behavior (ISO/IEC 9899:2011, 3.4.3).
- 2 A **C descriptor** whose address is a formal parameter that corresponds to a Fortran dummy argument becomes undefined on return from a call to the function from Fortran. If the dummy argument does not have either the **TARGET** or **ASYNCHRONOUS** attribute, all C pointers to any part of the object described by the **C descriptor** become undefined on return from the call, and any further use of them is undefined behavior.
- 3 If the address of a **C descriptor** is passed as an actual argument to a Fortran procedure, the lifetime (ISO/IEC 9899:2011, 6.2.4) of the **C descriptor** shall not end before the return from the procedure call. If an object is passed to a Fortran procedure as a nonallocatable, nonpointer dummy argument, its lifetime shall not end before the return from the procedure call.
- 4 If the lifetime of a **C descriptor** for an allocatable object that was established by C ends before the program exits, the object shall be unallocated at that time.
- 5 If a Fortran pointer becomes associated with a data object defined by the companion processor, the association status of the Fortran pointer becomes undefined when the lifetime of that data object ends.

### NOTE 15.32

The following example illustrates how a **C descriptor** becomes undefined upon returning from a call to a C function.

```
REAL, TARGET :: X(1000), B
INTERFACE
  REAL FUNCTION CFUN(ARRAY) BIND(C, NAME="Cfun")
    REAL ARRAY(:)
  END FUNCTION
END INTERFACE
B = CFUN(X)
```

Cfun is a C function. Before or during the invocation of Cfun, the processor will create a **C descriptor** for the array x. On return from Cfun, that **C descriptor** will become undefined. In addition, because the dummy argument ARRAY does not have the **TARGET** or **ASYNCHRONOUS** attribute, a C pointer whose value was set during execution of Cfun to be the address of any part of X will become undefined.

## 15.9 Interoperation with C global variables

### 15.9.1 General

- 1 A C variable whose name has external linkage may interoperate with a **common block** or with a variable declared in the scope of a module. The **common block** or variable shall be specified to have the **BIND attribute**.

- 1 2 At most one variable that is associated with a particular C variable whose name has external linkage is permitted  
 2 to be declared within all the Fortran **program units** of a program. A variable shall not be initially defined by  
 3 more than one processor.
- 4 3 If a **common block** is specified in a **BIND statement**, it shall be specified in a **BIND statement** with the same **binding label** in each  
 5 **scoping unit** in which it is declared. A C variable whose name has external linkage interoperates with a **common block** that has been  
 6 specified in a **BIND statement** if
- 7 • the C variable is of a struct type and the variables that are members of the **common block** are **interoperable** with corresponding  
 8 components of the struct type, or
  - 9 • the **common block** contains a single variable, and the variable is **interoperable** with the C variable.
- 10 4 There does not have to be an associated C entity for a Fortran entity with the **BIND attribute**.

**NOTE 15.33**

The following are examples of the usage of the **BIND attribute** for variables and for a **common block**. The Fortran variables, C\_EXTERN and C2, interoperate with the C variables, c\_extern and myVariable, respectively. The Fortran **common blocks**, COM and SINGLE, interoperate with the C variables, com and single, respectively.

```

MODULE LINK_TO_C_VARS
  USE, INTRINSIC :: ISO_C_BINDING
  INTEGER(C_INT), BIND(C) :: C_EXTERN
  INTEGER(C_LONG) :: C2
  BIND(C, NAME='myVariable') :: C2

  COMMON /COM/ R, S
  REAL(C_FLOAT) :: R, S, T
  BIND(C) :: /COM/, /SINGLE/
  COMMON /SINGLE/ T
END MODULE LINK_TO_C_VARS

/* Global variables. */
int c_extern;
long myVariable;
struct { float r, s; } com;
float single;

```

## 11 15.9.2 Binding labels for common blocks and variables

- 12 1 The **binding label** of a variable or **common block** is a default character value that specifies the name by which the  
 13 variable or **common block** is known to the **companion processor**.
- 14 2 If a variable or **common block** has the **BIND attribute** with the NAME= specifier and the value of its expression,  
 15 after discarding leading and trailing blanks, has nonzero length, the variable or **common block** has this as its **binding**  
 16 **label**. The case of letters in the **binding label** is significant. If a variable or **common block** has the **BIND attribute**  
 17 specified without a NAME= specifier, the **binding label** is the same as the name of the entity using lower case  
 18 letters. Otherwise, the variable or **common block** has no **binding label**.
- 19 3 The **binding label** of a C variable whose name has external linkage is the same as the name of the C variable. A  
 20 Fortran variable or **common block** with the **BIND attribute** that has the same **binding label** as a C variable whose  
 21 name has external linkage is **linkage associated** (16.5.1.5) with that variable.

## 15.10 Interoperation with C functions

### 15.10.1 Definition and reference of interoperable procedures

- 1 A procedure that is [interoperable](#) may be defined either by means other than Fortran or by means of a Fortran subprogram, but not both.
- 2 If the procedure is defined by means other than Fortran, it shall
  - be describable by a C prototype that is [interoperable](#) with the [interface](#),
  - have a name that has external linkage as defined by 6.2.2 of ISO/IEC 9899:2011, and
  - have the same [binding label](#) as the [interface](#).
- 3 A reference to such a procedure causes the function described by the C prototype to be called as specified in ISO/IEC 9899:2011.
- 4 A reference in C to a procedure that has the [BIND attribute](#), has the same [binding label](#), and is defined by means of Fortran, causes the Fortran procedure to be invoked. A C function shall not invoke a function pointer whose value is the result of a reference to [C\\_FUNLOC](#) with a noninteroperable argument.
- 5 A procedure defined by means of Fortran shall not invoke setjmp or longjmp (ISO/IEC 9899:2011, 7.13). If a procedure defined by means other than Fortran invokes setjmp or longjmp, that procedure shall not cause any procedure defined by means of Fortran to be invoked. A procedure defined by means of Fortran shall not be invoked as a signal handler (ISO/IEC 9899:2011, 7.14.1).
- 6 If a procedure defined by means of Fortran and a procedure defined by means other than Fortran perform input/output operations on the same [external file](#), the results are processor dependent (9.5.4).
- 7 If the value of a C function pointer will be the result of a reference to [C\\_FUNLOC](#) with a noninteroperable argument, it is recommended that the C function pointer be declared to have the type `void (*)()`.

### 15.10.2 Binding labels for procedures

- 1 The [binding label](#) of a procedure is a default character value that specifies the name by which a procedure with the [BIND attribute](#) is known to the [companion processor](#).
  - 2 If a procedure has the [BIND attribute](#) with the NAME= specifier and the value of its expression, after discarding leading and trailing blanks, has nonzero length, the procedure has this as its [binding label](#). The case of letters in the [binding label](#) is significant. If a procedure has the [BIND attribute](#) with no NAME= specifier, and the procedure is not a [dummy procedure](#), [internal procedure](#), or procedure pointer, then the [binding label](#) of the procedure is the same as the name of the procedure using lower case letters. Otherwise, the procedure has no [binding label](#).
- C1507 A procedure defined in a submodule shall not have a [binding label](#) unless its [interface](#) is declared in the ancestor module.
- 3 The [binding label](#) for a C function whose name has external linkage is the same as the C function name.

#### NOTE 15.34

In the following sample, the [binding label](#) of C\_SUB is "c\_sub", and the [binding label](#) of C\_FUNC is "C\_funC".

```
SUBROUTINE C_SUB() BIND(C)
...
END SUBROUTINE C_SUB
```

**NOTE 15.34 (cont.)**

```

INTEGER(C_INT) FUNCTION C_FUNC() BIND(C, NAME="C_func")
  USE, INTRINSIC :: ISO_C_BINDING
  ...
END FUNCTION C_FUNC

```

ISO/IEC 9899:2011 permits functions to have names that are not permitted as Fortran names; it also distinguishes between names that would be considered as the same name in Fortran. For example, a C name can begin with an underscore, and C names that differ in case are distinct names.

The specification of a [binding label](#) allows a program to use a Fortran name to refer to a procedure defined by a [companion processor](#).

**15.10.3 Exceptions and IEEE arithmetic procedures**

- 1 A procedure defined by means other than Fortran shall not use signal (ISO/IEC 9899:2011, 7.14.1) to change the [handling of any exception](#) that is being handled by the Fortran processor.
- 2 A procedure defined by means other than Fortran shall not alter the floating-point status ([14.7](#)) other than by setting an exception flag to signaling.
- 3 The values of the floating-point exception flags on entry to a procedure defined by means other than Fortran are processor dependent.

**15.10.4 Asynchronous communication**

- 1 Asynchronous communication for a Fortran variable with the [ASYNCHRONOUS attribute](#) occurs through the action of procedures defined by means other than Fortran. It is initiated by execution of an asynchronous communication initiation procedure and completed by execution of an asynchronous communication completion procedure. Between the execution of the initiation and completion procedures, any variable of which any part is associated with any part of the asynchronous communication variable is a pending communication affector. Whether a procedure is an asynchronous communication initiation or completion procedure is processor dependent.
- 2 Asynchronous communication is either input communication or output communication. For input communication, a pending communication affector shall not be referenced, become defined, become undefined, become associated with a dummy argument that has the [VALUE attribute](#), or have its pointer association status changed. For output communication, a pending communication affector shall not be redefined, become undefined, or have its pointer association status changed. The restrictions for asynchronous input communication are the same as for asynchronous input data transfer. The restrictions for asynchronous output communication are the same as for asynchronous output data transfer.

**NOTE 15.35**

Asynchronous communication can be used for nonblocking MPI calls such as MPI\_Irecv and MPI\_ISEND. For example,

```

REAL :: BUF(100,100)
... ! Code that involves BUF
BLOCK
  ASYNCHRONOUS :: BUF
  CALL MPI_Irecv(BUF,...REQ,...)
  ... ! Code that does not involve BUF
  CALL MPI_WAIT(REQ,...)
END BLOCK

```

**NOTE 15.35 (cont.)**

... ! Code that involves BUF

In this example, there is asynchronous input communication and BUF is a pending communication affecter between the two calls. MPI\_Irecv can return while the communication (reading values into BUF) is still underway. The intent is that the code between MPI\_Irecv and MPI\_WAIT can execute without waiting for this communication to complete.

Similar code with the call of MPI\_Irecv replaced by a call of MPI\_Isend is asynchronous output communication.

## 16 Scope, association, and definition

### 16.1 Scopes, identifiers, and entities

1 An entity is identified by an identifier.

2 The scope of

- a global identifier is a program (2.2.2),
- a local identifier is an [inclusive scope](#),
- an identifier of a [construct entity](#) is that construct (7.2.4, 8.1), and
- an identifier of a [statement entity](#) is that statement or part of that statement (3.3),

excluding any nested scope where the identifier is treated as the identifier of a different entity (16.3, 16.4).

3 An entity may be identified by

- an [image index](#) (1.3.87),
- a [name](#) (1.3.106),
- a [statement label](#) (1.3.138),
- an [external input/output unit](#) number (9.5),
- an identifier of a pending data transfer operation (9.6.2.9, 9.7),
- a submodule identifier (11.2.3),
- a [generic identifier](#) (1.3.81), or
- a [binding label](#) (1.3.16).

4 By means of association, an entity may be referred to by the same identifier or a different identifier in a different scope, or by a different identifier in the same scope.

### 16.2 Global identifiers

1 [Program units](#), [common blocks](#), [external procedures](#), entities with [binding labels](#), [external input/output units](#), pending data transfer operations, and [images](#) are global entities of a program. The name of a [common block](#) with no [binding label](#), [external procedure](#) with no [binding label](#), or [program unit](#) that is not a submodule is a global identifier. The submodule identifier of a submodule is a global identifier. A [binding label](#) of an entity of the program is a global identifier. An entity of the program shall not be identified by more than one [binding label](#).

2 The global identifier of an entity shall not be the same as the global identifier of any other entity. Furthermore, a [binding label](#) shall not be the same as the global identifier of any other global entity, ignoring differences in case. A processor may assign a global identifier to an entity that is not specified by this part of ISO/IEC 1539 to have a global identifier (such as an intrinsic procedure); in such a case, the processor shall ensure that this assigned global identifier differs from all other global identifiers in the program.

#### NOTE 16.1

An intrinsic module is not a [program unit](#), so a global identifier can be the same as the name of an intrinsic module.

**NOTE 16.2**

Submodule identifiers are global identifiers, but because they consist of a module name and a descendant submodule name, the name of a submodule can be the same as the name of another submodule so long as they do not have the same ancestor module.

## 16.3 Local identifiers

### 16.3.1 Classes of local identifiers

#### 1 Identifiers of entities in the classes

- (1) except for statement or **construct** entities (16.4), named variables, **named constants**, named constructs, statement functions, **internal procedures**, **module procedures**, **dummy procedures**, **intrinsic procedures**, **external procedures** that have **binding labels**, intrinsic modules, **abstract interfaces**, **generic interfaces**, derived types, namelist groups, **external procedures** accessed via **USE**, and statement labels,
- (2) type parameters, components, and **type-bound procedure bindings**, in a separate class for each type,
- (3) **argument keywords**, in a separate class for each procedure with an **explicit interface**, and
- (4) common blocks that have **binding labels**

are local identifiers.

#### 2 Within its scope, a local identifier of an entity of class (1) or class (4) shall not be the same as a global identifier used in that scope unless the global identifier

- is used only as the *use-name* of a **rename** in a **USE statement**,
- is a **common block** name (16.3.2),
- is an **external procedure** name that is also a generic name, or
- is an external function name and the **inclusive scope** is its defining subprogram (16.3.3).

#### 3 Within its scope, a local identifier of one class shall not be the same as another local identifier of the same class, except that a generic name may be the same as the name of a procedure as explained in 12.4.3.5 or the same as the name of a derived type (4.5.10). A local identifier of one class may be the same as a local identifier of another class.

**NOTE 16.3**

An intrinsic procedure is inaccessible by its own name in a **scoping unit** that uses the same name as a local identifier of class (1) for a different entity. For example, in the program fragment

```
SUBROUTINE SUB
  ...
  A = SIN (K)
  ...
CONTAINS
  FUNCTION SIN (X)
    ...
  END FUNCTION SIN
END SUBROUTINE SUB
```

any reference to function SIN in subroutine SUB refers to the internal function SIN, not to the intrinsic function of the same name.

4 A local identifier identifies an entity in a scope and may be used to identify an entity in another scope except in the following cases.

- The name that appears as a *subroutine-name* in a *subroutine-stmt* has limited use within the scope established by the *subroutine-stmt*. It can be used to identify recursive references of the subroutine or to identify a **common block** (the latter is possible only for internal and module subroutines).
- The name that appears as a *function-name* in a *function-stmt* has limited use within the scope established by that *function-stmt*. It can be used to identify the **function result**, to identify recursive references of the function, or to identify a **common block** (the latter is possible only for internal and module functions).
- The name that appears as an *entry-name* in an *entry-stmt* has limited use within the scope of the subprogram in which the *entry-stmt* appears. It can be used to identify the **function result** if the subprogram is a function, to identify recursive references, or to identify a **common block** (the latter is possible only if the *entry-stmt* is in a module subprogram).

### 16.3.2 Local identifiers that are the same as common block names

- 1 A name that identifies a **common block** in a **scoping unit** shall not be used to identify a constant or an intrinsic procedure in that **scoping unit**. If a local identifier of class (1) is also the name of a **common block**, the appearance of that name in any context other than as a **common block** name in a **BIND**, **COMMON**, or **SAVE** statement is an appearance of the local identifier.

### 16.3.3 Function results

- 1 For each **FUNCTION statement** or **ENTRY statement** in a function subprogram, there is a **function result**. If there is no **RESULT** clause, the **function result** has the same name as the function being defined; otherwise, the **function result** has the name specified in the **RESULT** clause.

### 16.3.4 Components, type parameters, and bindings

- 1 A component name has the scope of its derived-type definition. Outside the type definition, it may also appear within a designator of a component of a structure of that type or as a **component keyword** in a **structure constructor** for that type.
- 2 A type parameter name has the scope of its derived-type definition. Outside the derived-type definition, it may also appear as a type parameter keyword in a *derived-type-spec* for the type or as the *type-param-name* of a *type-param-inquiry*.
- 3 The **binding name** (4.5.5) of a **type-bound procedure** has the scope of its derived-type definition. Outside of the derived-type definition, it may also appear as the *binding-name* in a procedure reference.
- 4 A generic **binding** for which the *generic-spec* is not a *generic-name* has a scope that consists of all **scoping units** in which an entity of the type is accessible.
- 5 A component name or **binding name** may appear only in a scope in which it is accessible.
- 6 The accessibility of components and **bindings** is specified in 4.5.4.8 and 4.5.5.

### 16.3.5 Argument keywords

- 1 As an **argument keyword**, a **dummy argument** name in an **internal procedure**, **module procedure**, or an **interface body** has a scope of the **scoping unit** of the **host** of the procedure or **interface body**. As an **argument keyword**, the name of a **dummy argument** of a procedure declared by a **procedure declaration statement** that specifies an **explicit interface** has a scope of the **scoping unit** containing the **procedure declaration statement**. It may appear only in a procedure reference for the procedure of which it is a **dummy argument**. If the procedure is accessible

1 in another [scoping unit](#) by [use](#) or [host](#) association (16.5.1.3, 16.5.1.4), the [argument keyword](#) is accessible for  
 2 procedure references for that procedure in that [scoping unit](#).

3 2 A [dummy argument](#) name in an [intrinsic](#) procedure has a scope as an [argument keyword](#) of the [scoping unit](#)  
 4 in which the reference to the procedure occurs. As an [argument keyword](#), it may appear only in a procedure  
 5 reference for the procedure of which it is a [dummy argument](#).

## 6 16.4 Statement and construct entities

7 1 A variable that appears as a [data-i-do-variable](#) in a [DATA statement](#) or an [ac-do-variable](#) in an array constructor,  
 8 as a dummy argument in a [statement function statement](#), or as an [index-name](#) in a [FORALL statement](#) is a [statement entity](#). A  
 9 variable that appears as an [index-name](#) in a [FORALL](#) or [DO CONCURRENT](#), as an [associate-name](#) in a [SELECT](#)  
 10 [TYPE](#) or [ASSOCIATE](#) construct, or as a [coarray-name](#) in a [codimension-decl](#) in a [CHANGE TEAM](#) construct  
 11 is a [construct entity](#). A variable that has [LOCAL](#) or [LOCAL\\_INIT](#) locality in a [DO CONCURRENT](#) construct  
 12 is a [construct entity](#). An entity that is explicitly declared in the specification part of a [BLOCK construct](#), other  
 13 than only in [ASYNCHRONOUS](#) and [VOLATILE](#) statements, is a [construct entity](#). Two [construct entities](#) of the  
 14 same construct shall not have the same identifier.

15 2 Even if the name of a statement entity is the same as another identifier and the statement is in the scope of that  
 16 identifier, within the scope of the statement entity the name is interpreted as that of the statement entity.

17 3 The name of a statement entity shall not be the same as an accessible global identifier or local identifier of class  
 18 (1) (16.3.1), except for a [common block](#) name or a scalar variable name. Within the scope of a [statement entity](#),  
 19 another [statement entity](#) shall not have the same name.

20 4 The name of a [data-i-do-variable](#) in a [DATA statement](#) or an [ac-do-variable](#) in an array constructor has a scope  
 21 of its [data-implied-do](#) or [ac-implied-do](#). It is a scalar variable. If [integer-type-spec](#) appears in [data-implied-do](#) or  
 22 [ac-implied-do-control](#) it has the specified type and type parameters; otherwise it has the type and type parameters  
 23 that it would have if it were the name of a variable in the innermost executable construct or [scoping unit](#) that  
 24 includes the [DATA statement](#) or array constructor, and this type shall be integer type. It has no other attributes.  
 25 The appearance of a name as a [data-i-do-variable](#) of an implied DO in a [DATA statement](#) or an [ac-do-variable](#)  
 26 in an array constructor is not an implicit declaration of a variable whose scope is the [scoping unit](#) that contains  
 27 the statement.

28 5 The name of a variable that appears as an [index-name](#) in a [DO CONCURRENT construct](#), [FORALL statement](#), or  
 29 [FORALL construct](#) has a scope of the statement or construct. It is a scalar variable. If [integer-type-spec](#) appears in  
 30 [concurrent-header](#) it has the specified type and type parameters; otherwise it has the type and type parameters  
 31 that it would have if it were the name of a variable in the innermost executable construct or [scoping unit](#) that  
 32 includes the [DO CONCURRENT](#) or [FORALL](#), and this type shall be integer type. It has no other attributes.  
 33 The appearance of a name as an [index-name](#) in a [DO CONCURRENT construct](#), [FORALL statement](#), or [FORALL](#)  
 34 [construct](#) is not an implicit declaration of a variable whose scope is the [scoping unit](#) that contains the statement or  
 35 construct.

36 6 A variable that has [LOCAL](#) or [LOCAL\\_INIT](#) locality in a [DO CONCURRENT construct](#) has the scope of that  
 37 construct. Its attributes are specified in 8.1.7.5.

38 7 If [integer-type-spec](#) does not appear in a [concurrent-header](#), an [index-name](#) shall not be the same as an accessible  
 39 global identifier, local identifier, or identifier of an outer construct entity, except for a [common block](#) name or

a scalar variable name. An *index-name* of a contained **DO CONCURRENT** construct, **FORALL** statement, or **FORALL** construct shall not be the same as an *index-name* of any of its containing **DO CONCURRENT** or **FORALL** constructs.

The **associate name** of a **SELECT TYPE** construct has a separate scope for each block of the construct. Within each block, it has the **declared type**, **dynamic type**, type parameters, **rank**, and bounds specified in 8.1.11.2.

The **associate names** of an **ASSOCIATE** construct have the scope of the block. They have the **declared type**, **dynamic type**, type parameters, **rank**, and **bounds** specified in 8.1.3.2.

The **associate names** of a **CHANGE TEAM** construct have the scope of the block. They have the **declared type**, **dynamic type**, type parameters, **rank**, **corank**, **bounds**, and **cobounds** specified in 8.1.5.

The name of a variable that appears as a dummy argument in a **statement function statement** has a scope of the statement in which it appears. It is a scalar that has the type and type parameters that it would have if it were the name of a variable in the **scoping unit** that includes the **statement function**; it has no other attributes.

## 16.5 Association

### 16.5.1 Name association

#### 16.5.1.1 Forms of name association

There are five forms of **name association**: **argument association**, **use association**, **host association**, **linkage association**, and **construct association**. **Argument**, **use**, and **host** association provide mechanisms by which entities known in one scope may be accessed in another scope.

#### 16.5.1.2 Argument association

The rules governing **argument association** are given in Clause 12. As explained in 12.5, execution of a procedure reference establishes a correspondence between each dummy argument and an **actual argument** and thus an association between each dummy argument and its **effective argument**. Argument association may be sequence association (12.5.2.11).

The name of the dummy argument may be different from the name, if any, of its **effective argument**. The dummy argument name is the name by which the **effective argument** is known, and by which it may be accessed, in the referenced procedure.

#### NOTE 16.4

An **effective argument** can be a nameless data entity, such as the result of evaluating an expression that is not simply a variable or constant.

Upon termination of execution of a procedure reference, all **argument associations** established by that reference are terminated. A dummy argument of that procedure may be associated with an entirely different **effective argument** in a subsequent invocation of the procedure.

### 16.5.1.3 Use association

**Use association** is the association of names in different scopes specified by a **USE statement**. The rules governing **use association** are given in 11.2.2. They allow for renaming of entities being accessed. **Use association** allows access in one scope to entities defined or declared in another scope; it remains in effect throughout the execution of the program.

### 16.5.1.4 Host association

A nested **scoping unit** has access to named entities from its host as specified in 12.4.3.4. In the case of an **internal subprogram**, the access is to the entities in its **host instance**. The accessed entities are identified by the same identifier and have the same attributes as in the **host**, except that a local entity may have the **ASYNCHRONOUS attribute** even if the host entity does not, and a noncoarray local entity may have the **VOLATILE attribute** even if the host entity does not. The accessed entities are named data objects, derived types, **abstract interfaces**, procedures, **generic identifiers**, and namelist groups.

If an entity that is accessed by **use association** has the same nongeneric name as a host entity, the host entity is inaccessible by that name. The name of an **external procedure** that is given the **EXTERNAL attribute** (5.5.9) within the **scoping unit**, or a name that appears within the **scoping unit** as a *module-name* in a *use-stmt* is a global identifier; any entity of the **host** that has this as its nongeneric name is inaccessible by that name. A name that appears in the **scoping unit** as

- (1) a *function-name* in a *stmt-function-stmt* or in an *entity-decl* in a *type-declaration-stmt*, unless it is a global identifier,
- (2) an *object-name* in an *entity-decl* in a *type-declaration-stmt*, in a *pointer-stmt*, in a *save-stmt*, in an *allocatable-stmt*, or in a *target-stmt*,
- (3) a *type-param-name* in a *derived-type-stmt*,
- (4) a *named-constant* in a *named-constant-def* in a *parameter-stmt*,
- (5) a *coarray-name* in a *codimension-stmt*,
- (6) an *array-name* in a *dimension-stmt*,
- (7) a *variable-name* in a *common-block-object* in a *common-stmt*,
- (8) a **procedure pointer** given the **EXTERNAL attribute** in the **scoping unit**,
- (9) the name of a variable that is wholly or partially initialized in a *data-stmt*,
- (10) the name of an object that is wholly or partially equivalenced in an *equivalence-stmt*,
- (11) a *dummy-arg-name* in a *function-stmt*, in a *subroutine-stmt*, in an *entry-stmt*, or in a *stmt-function-stmt*,
- (12) a *result-name* in a *function-stmt* or in an *entry-stmt*,
- (13) the name of an entity declared by an **interface body**, unless it is a global identifier,
- (14) an *intrinsic-procedure-name* in an *intrinsic-stmt*,
- (15) a *namelist-group-name* in a *namelist-stmt*,
- (16) a *generic-name* in a *generic-spec* in an *interface-stmt*, or
- (17) the name of a named construct

is a local identifier in the **scoping unit** and any entity of the **host** that has this as its nongeneric name is inaccessible by that name by **host association**. If a **scoping unit** is the **host** of a derived-type definition or a subprogram that does not define a separate module procedure, the name of the derived type or of any procedure defined by the subprogram is a local identifier in the **scoping unit**; any entity of the **host** that has this as its nongeneric name is

1 inaccessible by that name. Local identifiers of a subprogram are not accessible to its [host](#).

#### NOTE 16.5

A name that appears in an [ASYNCHRONOUS](#) or [VOLATILE](#) statement is not necessarily the name of a local variable. In an [internal](#) or module procedure, if a variable that is accessible via [host association](#) is specified in an [ASYNCHRONOUS](#) or [VOLATILE](#) statement, that host variable is given the [ASYNCHRONOUS](#) or [VOLATILE](#) attribute in the local scope.

2 3 If a host entity is inaccessible only because a local variable with the same name is wholly or partially initialized  
3 in a [DATA statement](#), the local variable shall not be referenced or defined prior to the [DATA statement](#).

4 4 If a derived-type name of a [host](#) is inaccessible, data entities of that type or subobjects of such data entities still  
5 can be accessible.

#### NOTE 16.6

An [interface body](#) that is not a module procedure interface body accesses by [host association](#) only those entities made accessible by [IMPORT statements](#).

6 5 If an [external](#) or [dummy](#) procedure with an [implicit interface](#) is accessed via [host association](#), then it shall have  
7 the [EXTERNAL attribute](#) in the [host scoping unit](#); if it is invoked as a function in the inner [scoping unit](#), its type  
8 and type parameters shall be established in the [host scoping unit](#). The type and type parameters of a function  
9 with the [EXTERNAL attribute](#) are established in a [scoping unit](#) if that [scoping unit](#) explicitly declares them,  
10 invokes the function, accesses the function from a module, or accesses the function from its [host](#) where its type  
11 and type parameters are established.

12 6 If an intrinsic procedure is accessed via [host association](#), then it shall be established to be intrinsic in the [host](#)  
13 [scoping unit](#). An intrinsic procedure is established to be intrinsic in a [scoping unit](#) if that [scoping unit](#) explicitly  
14 gives it the [INTRINSIC attribute](#), invokes it as an intrinsic procedure, accesses it from a module, or accesses it  
15 from its [host](#) where it is established to be intrinsic.

#### NOTE 16.7

A host subprogram and an internal subprogram can contain the same and differing use-associated entities, as illustrated in the following example.

```
MODULE B; REAL BX, Q; INTEGER IX, JX; END MODULE B
MODULE C; REAL CX; END MODULE C
MODULE D; REAL DX, DY, DZ; END MODULE D
MODULE E; REAL EX, EY, EZ; END MODULE E
MODULE F; REAL FX; END MODULE F
MODULE G; USE F; REAL GX; END MODULE G
PROGRAM A
USE B; USE C; USE D
...
CONTAINS
  SUBROUTINE INNER_PROC (Q)
    USE C          ! Not needed
    USE B, ONLY: BX ! Entities accessible are BX, IX, and JX
                  ! if no other IX or JX
```

## NOTE 16.7 (cont.)

```

! is accessible to INNER_PROC
! Q is local to INNER_PROC,
! because Q is a dummy argument
USE D, X => DX ! Entities accessible are DX, DY, and DZ
! X is local name for DX in INNER_PROC
! X and DX denote same entity if no other
! entity DX is local to INNER_PROC
USE E, ONLY: EX ! EX is accessible in INNER_PROC, not in program A
! EY and EZ are not accessible in INNER_PROC
! or in program A
USE G          ! FX and GX are accessible in INNER_PROC
...
END SUBROUTINE INNER_PROC
END PROGRAM A

```

Because program A contains the statement

```
USE B
```

all of the entities in module B, except for Q, are accessible in INNER\_PROC, even though INNER\_PROC contains the statement

```
USE B, ONLY: BX
```

The [USE statement](#) with the [ONLY](#) option means that this particular statement brings in only the entity named, not that this is the only variable from the module accessible in this [scoping unit](#).

## NOTE 16.8

For more examples of [host association](#), see subclause [C.13](#).

## 1 16.5.1.5 Linkage association

2 1 [Linkage association](#) occurs between a module variable that has the [BIND attribute](#) and the C variable with which  
3 it interoperates, or between a Fortran [common block](#) and the C variable with which it interoperates ([15.9](#)). Such association  
4 remains in effect throughout the execution of the program.

## 5 16.5.1.6 Construct association

6 1 Execution of a [SELECT RANK](#) or [SELECT TYPE](#) statement establishes an association between the selector and  
7 the [associate name](#) of the construct. Execution of an [ASSOCIATE statement](#) establishes an association between  
8 each selector and the corresponding [associate name](#) of the construct. Execution of a [CHANGE TEAM statement](#)  
9 establishes an association between each coselector and the corresponding [associate name](#) of the construct.

10 2 In an [ASSOCIATE](#) or [SELECT TYPE](#) construct, the following rules apply.

- 11 • If a selector is [allocatable](#), it shall be allocated; the [associate name](#) is associated with the data object and  
12 does not have the [ALLOCATABLE attribute](#).

- If a selector has the **POINTER attribute**, it shall be associated; the **associate name** is associated with the **target** of the pointer and does not have the **POINTER attribute**.

3 If the selector is a variable other than an **array section** having a **vector subscript**, the association is with the data object specified by the selector; otherwise, the association is with the value of the selector expression, which is evaluated prior to execution of the block.

4 Each **associate name** remains associated with the corresponding selector throughout the execution of the executed block. Within the block, each selector is known by and may be accessed by the corresponding **associate name**. On completion of execution of the construct, the association is terminated.

#### NOTE 16.9

The association between the **associate name** and a data object is established prior to execution of the block and is not affected by subsequent changes to variables that were used in subscripts or substring ranges in the **selector**.

## 16.5.2 Pointer association

### 16.5.2.1 General

1 **Pointer association** between a **pointer** and a **target** allows the **target** to be referenced by a reference to the pointer. At different times during the execution of a program, a pointer may be **undefined**, associated with different **targets** on its own **image**, or be **disassociated**. If a pointer is associated with a **target**, the definition status of the pointer is either defined or undefined, depending on the definition status of the **target**. If the pointer has deferred **type parameters** or shape, their values are assumed from the **target**. If the pointer is **polymorphic**, its **dynamic type** is assumed from the **dynamic type** of the **target**.

### 16.5.2.2 Pointer association status

1 A pointer may have a pointer association status of associated, **disassociated**, or **undefined**. Its association status may change during execution of a program. Unless a pointer is initialized (**explicitly** or by **default**), it has an initial association status of **undefined**. A pointer may be initialized to have an association status of **disassociated** or associated.

#### NOTE 16.10

A pointer from a module **program unit** might be accessible in a subprogram via **use association**. Such pointers have a lifetime that is greater than **targets** that are declared in the subprogram, unless such **targets** are **saved**. Therefore, if such a pointer is associated with a local **target**, there is the possibility that when a procedure defined by the subprogram completes execution, the **target** will cease to exist, leaving the pointer “dangling”. This part of ISO/IEC 1539 considers such pointers to have an undefined association status. They are neither associated nor **disassociated**. They cannot be used again in the program until their status has been reestablished. A processor is not required to detect when a pointer **target** ceases to exist.

### 16.5.2.3 Events that cause pointers to become associated

1 A pointer becomes associated when any of the following events occur.

- (1) The pointer is allocated (6.7.1) as the result of the successful execution of an **ALLOCATE statement** referencing the pointer.

- (2) The pointer is pointer-assigned to a [target \(7.2.2\)](#) that is associated or is specified with the [TARGET attribute](#) and, if [allocatable](#), is allocated.
- (3) The pointer is a [subobject](#) of an object that is allocated by an [ALLOCATE statement](#) in which [SOURCE=](#) appears and the corresponding [subobject](#) of *source-expr* is associated.
- (4) The pointer is a dummy argument and its corresponding [actual argument](#) is not a pointer.
- (5) The pointer is a [default-initialized subcomponent](#) of an object, the corresponding initializer is not a reference to the intrinsic function [NULL](#), and
  - (a) a procedure is invoked with this object as an [actual argument](#) corresponding to a nonpointer nonallocatable dummy argument with [INTENT \(OUT\)](#),
  - (b) a procedure with this object as an [unsaved](#) nonpointer nonallocatable local variable is invoked,
  - (c) a [BLOCK construct](#) is entered and this object is an [unsaved](#) local nonpointer nonallocatable local variable of the [BLOCK construct](#),  
or
  - (d) this object is allocated other than by an [ALLOCATE statement](#) in which [SOURCE=](#) appears.

#### 16.5.2.4 Events that cause pointers to become disassociated

1 A pointer becomes [disassociated](#) when

- (1) the pointer is nullified ([6.7.2](#)),
- (2) the pointer is deallocated ([6.7.3](#)),
- (3) the pointer is [pointer-assigned \(7.2.2\)](#) to a [disassociated](#) pointer,
- (4) the pointer is a [subobject](#) of an object that is allocated by an [ALLOCATE statement](#) in which [SOURCE=](#) appears and the corresponding [subobject](#) of *source-expr* is [disassociated](#),  
or
- (5) the pointer is a [default-initialized subcomponent](#) of an object, the corresponding initializer is a reference to the intrinsic function [NULL](#), and
  - (a) a procedure is invoked with this object as an [actual argument](#) corresponding to a nonpointer nonallocatable dummy argument with [INTENT \(OUT\)](#),
  - (b) a procedure with this object as an [unsaved](#) nonpointer nonallocatable local variable is invoked,
  - (c) a [BLOCK construct](#) is entered and this object is an [unsaved](#) local nonpointer nonallocatable local variable of the [BLOCK construct](#),  
or
  - (d) this object is allocated other than by an [ALLOCATE statement](#) in which [SOURCE=](#) appears.

#### 16.5.2.5 Events that cause the association status of pointers to become undefined

1 The association status of a pointer becomes undefined when

- (1) the pointer is pointer-assigned to a [target](#) that has an undefined association status,
- (2) the pointer is pointer-assigned to a [target](#) on a different [image](#),
- (3) the [target](#) of the pointer is deallocated other than through the pointer,
- (4) the [target](#) of the pointer is a data object defined by the companion processor and the lifetime of that data object ends,

- (5) the allocation transfer procedure (13.9.137) is executed, the pointer is associated with the argument FROM, and the argument TO does not have the TARGET attribute,
- (6) completion of execution of an instance of a subprogram causes the pointer's target to become undefined (item (3) of 16.6.6),
- (7) completion of execution of a BLOCK construct causes the pointer's target to become undefined (item (23) of 16.6.6),
- (8) execution of the host instance of a procedure pointer is completed,
- (9) execution of an instance of a subprogram completes and the pointer is declared or accessed in the subprogram that defines the procedure unless the pointer
  - (a) has the SAVE attribute,
  - (b) is in blank common,
  - (c) is in a named common block that is declared in at least one other scoping unit that is in execution,
  - (d) is accessed by host association, or
  - (e) is the return value of a function declared to have the POINTER attribute,
- (10) execution of an instance of a subprogram completes, the pointer is associated with a dummy argument of the procedure, and
  - (a) the effective argument does not have the TARGET attribute or is an array section with a vector subscript, or
  - (b) the dummy argument has the VALUE attribute,
- (11) a BLOCK construct completes execution and the pointer is an unsaved construct entity of that BLOCK construct,
- (12) a DO CONCURRENT construct is terminated and the pointer's association status was changed in more than one iteration of the construct,
- (13) an iteration of a DO CONCURRENT construct completes and the pointer is associated with a variable of that construct that has LOCAL or LOCAL\_INIT locality,
- (14) the pointer is a subcomponent of an object that is allocated and either
  - (a) the pointer is not default-initialized and SOURCE= does not appear, or
  - (b) SOURCE= appears and the association status of the corresponding subcomponent of source-expr is undefined,
- (15) the pointer is a subcomponent of an object, the pointer is not default-initialized, and a procedure is invoked with this object as an actual argument corresponding to a dummy argument with INTENT (OUT), or
- (16) a procedure is invoked with the pointer as an actual argument corresponding to a pointer dummy argument with INTENT (OUT).

#### 16.5.2.6 Other events that change the association status of pointers

- 1 When a pointer becomes associated with another pointer by argument association, construct association, or host association, the effects on its association status are specified in 16.5.5.
- 2 While two pointers are name associated, storage associated, or inheritance associated, if the association status of one pointer changes, the association status of the other changes accordingly.

3 The association status of a pointer object with the [VOLATILE attribute](#) might change by means not specified  
 2 by the program.

### 3 16.5.2.7 Pointer definition status

4 1 The definition status of an associated pointer is that of its [target](#). If a pointer is associated with a [definable](#)  
 5 [target](#), it may be defined or become undefined according to the rules for a variable (16.6). The definition status  
 6 of a pointer that is not associated is undefined.

### 7 16.5.2.8 Relationship between association status and definition status

8 1 If the association status of a pointer is [disassociated](#) or undefined, the pointer shall not be referenced or dealloc-  
 9 ated. Whatever its association status, a pointer always may be nullified, allocated, or pointer-assigned. A nullified  
 10 pointer is [disassociated](#). When a pointer is allocated, it becomes associated but undefined. When a pointer is  
 11 pointer-assigned, its association and definition status become those of the specified [data-target](#) or [proc-target](#).

## 12 16.5.3 Storage association

### 13 16.5.3.1 General

14 1 [Storage sequences](#) are used to describe relationships that exist among variables and [common blocks](#). [Storage asso-](#)  
 15 [ciation](#) is the association of two or more data objects that occurs when two or more [storage sequences](#) share or  
 16 are aligned with one or more [storage units](#).

### 17 16.5.3.2 Storage sequence

18 1 A [storage sequence](#) is a sequence of [storage units](#). The size of a [storage sequence](#) is the number of [storage units](#)  
 19 in the [storage sequence](#). A [storage unit](#) is a [character storage unit](#), a [numeric storage unit](#), a [file storage unit](#)  
 20 (9.3.5), or an [unspecified storage unit](#). The sizes of the [numeric storage unit](#), the [character storage unit](#) and the  
 21 [file storage unit](#) are the values of constants in the ISO\_FORTRAN\_ENV intrinsic module (13.10.2).

22 2 In a storage association context

- 23 (1) a nonpointer scalar object that is default integer, default real, or default logical occupies a single  
 24 [numeric storage unit](#),
- 25 (2) a nonpointer scalar object that is double precision real or default complex occupies two [contiguous](#)  
 26 [numeric storage units](#),
- 27 (3) a default character nonpointer scalar object of character length *len* occupies *len* [contiguous character](#)  
 28 [storage units](#),
- 29 (4) if C character kind is not the same as default character kind a nonpointer scalar object of type char-  
 30 acter with the C character kind (15.2.2) and character length *len* occupies *len* [contiguous unspecified](#)  
 31 [storage units](#),
- 32 (5) a nonpointer scalar object of [sequence type](#) occupies a sequence of [storage sequences](#) corresponding  
 33 to the sequence of its [ultimate components](#),
- 34 (6) a nonpointer scalar object of any type not specified in items (1)-(5) occupies a single [unspecified](#)  
 35 [storage unit](#) that is different for each case and each set of type parameter values, and that is different  
 36 from the [unspecified storage units](#) of item (4),
- 37 (7) a nonpointer array occupies a sequence of [contiguous storage sequences](#), one for each array element,  
 38 in array element order (6.5.3.2), and

- (8) a **data pointer** occupies a single **unspecified storage unit** that is different from that of any nonpointer object and is different for each combination of type, type parameters, and **rank**. A **data pointer** that has the **CONTIGUOUS attribute** occupies a **storage unit** that is different from that of a **data pointer** that does not have the **CONTIGUOUS attribute**.

3 A sequence of **storage sequences** forms a **storage sequence**. The order of the **storage units** in such a composite **storage sequence** is that of the individual **storage units** in each of the constituent **storage sequences** taken in succession, ignoring any zero-sized constituent sequences.

4 Each **common block** has a **storage sequence** (5.9.2.2).

### 16.5.3.3 Association of storage sequences

1 Two nonzero-sized **storage sequences**  $s_1$  and  $s_2$  are storage associated if the  $i$ th **storage unit** of  $s_1$  is the same as the  $j$ th **storage unit** of  $s_2$ . This causes the  $(i + k)$ th **storage unit** of  $s_1$  to be the same as the  $(j + k)$ th **storage unit** of  $s_2$ , for each integer  $k$  such that  $1 \leq i + k \leq \text{size of } s_1$  and  $1 \leq j + k \leq \text{size of } s_2$  where *size of* measures the number of **storage units**.

2 Storage association also is defined between two zero-sized **storage sequences**, and between a zero-sized **storage sequence** and a **storage unit**. A zero-sized **storage sequence** in a sequence of **storage sequences** is storage associated with its successor, if any. If the successor is another zero-sized **storage sequence**, the two **sequences** are storage associated. If the successor is a nonzero-sized **storage sequence**, the zero-sized **sequence** is storage associated with the first **storage unit** of the successor. Two **storage units** that are each storage associated with the same zero-sized **storage sequence** are the same **storage unit**.

### 16.5.3.4 Association of scalar data objects

1 Two scalar data objects are storage associated if their **storage sequences** are storage associated. Two scalar entities are totally associated if they have the same **storage sequence**. Two scalar entities are partially associated if they are associated without being totally associated.

2 The definition status and value of a data object affects the definition status and value of any storage associated entity. An **EQUIVALENCE statement**, a **COMMON statement**, or an **ENTRY statement** can cause storage association of **storage sequences**.

3 An **EQUIVALENCE statement** causes storage association of data objects only within one **scoping unit**, unless one of the equivalenced entities is also in a **common block** (5.9.1.2, 5.9.2.2).

4 **COMMON statements** cause data objects in one **scoping unit** to become storage associated with data objects in another **scoping unit**.

5 A **common block** is permitted to contain a sequence of differing **storage units**. All **scoping units** that access named **common blocks** with the same name shall specify an identical sequence of **storage units**. **Blank common** blocks may be declared with differing sizes in different **scoping units**. For any two **blank common** blocks, the initial sequence of **storage units** of the longer **blank common** block shall be identical to the sequence of **storage units** of the shorter **common block**. If two **blank common** blocks are the same length, they shall have the same sequence of **storage units**.

6 An **ENTRY statement** in a function subprogram causes storage association of the **function results** that are variables.

7 Partial association shall exist only between

- an object that is default character or of **character sequence type** and an object that is default character or

of [character sequence type](#), or

- an object that is default complex, double precision real, or of [numeric sequence type](#) and an object that is default integer, default real, default logical, double precision real, default complex, or of [numeric sequence type](#).

For noncharacter entities, partial association may occur only through the use of [COMMON](#), [EQUIVALENCE](#), or [ENTRY](#) statements. For character entities, partial association may occur only through [argument association](#) or the use of [COMMON](#) or [EQUIVALENCE](#) statements.

Partial association of character entities occurs when some, but not all, of the [storage units](#) of the entities are the same.

A [storage unit](#) shall not be [explicitly initialized](#) more than once in a program. [Explicit initialization](#) overrides [default initialization](#), and [default initialization](#) for an object of derived type overrides [default initialization](#) for a component of the object (4.5.4.6). [Default initialization](#) may be specified for a [storage unit](#) that is storage associated provided the objects supplying the [default initialization](#) are of the same type and type parameters, and supply the same value for the [storage unit](#).

#### 16.5.4 Inheritance association

[Inheritance association](#) occurs between components of the [parent component](#) and components [inherited](#) by type extension into an [extended type](#) (4.5.7.2). This association is persistent; it is not affected by the accessibility of the [inherited](#) components.

#### 16.5.5 Establishing associations

When an association is established between two entities by [argument association](#), [host association](#), or [construct association](#), certain properties of the [associating entity](#) become those of the pre-existing entity.

For [argument association](#), the pre-existing entity is the [effective argument](#) and the [associating entity](#) is the dummy argument.

For [host association](#), the [associating entity](#) is the entity in the contained [scoping unit](#). When a procedure is invoked, the pre-existing entity that participates in the association is the one from its [host instance](#) (12.6.2.4). Otherwise the pre-existing entity that participates in the association is the entity in the [host scoping unit](#).

For [construct association](#), the [associating entity](#) is identified by the [associate name](#) and the pre-existing entity is the selector.

When an association is established by [argument association](#), [host association](#), or [construct association](#), the following applies.

- If the entities have the [POINTER attribute](#), the [pointer association](#) status of the [associating entity](#) becomes the same as that of the pre-existing entity. If the pre-existing entity has a [pointer association](#) status of associated, the [associating entity](#) becomes [pointer associated](#) with the same [target](#) and, if they are arrays, the bounds of the [associating entity](#) become the same as those of the pre-existing entity.
- If the [associating entity](#) has the [ALLOCATABLE attribute](#), its allocation status becomes the same as that of the pre-existing entity. If the pre-existing entity is allocated, the bounds (if it is an array), values of [deferred type parameters](#), definition status, and value (if it is defined) become the same as those of the

pre-existing entity. If the **associating entity** is **polymorphic** and the pre-existing entity is allocated, the **dynamic type** of the **associating entity** becomes the same as that of the pre-existing entity.

- If the **associating entity** is neither a pointer nor allocatable, its definition status, value (if it is defined), and **dynamic type** (if it is **polymorphic**) become the same as those of the pre-existing entity. If the entities are arrays and the association is not **argument association**, the bounds of the **associating entity** become the same as those of the pre-existing entity.
- If the **associating entity** is a pointer dummy argument and the pre-existing entity is a nonpointer **actual argument** the **associating entity** becomes **pointer associated** with the pre-existing entity and, if the entities are arrays, the bounds of the **associating entity** become the same as those of the pre-existing entity.

## 16.6 Definition and undefinition of variables

### 16.6.1 Definition of objects and subobjects

- 1 A variable may be defined or may be undefined and its definition status may change during execution of a program. An action that causes a variable to become undefined does not imply that the variable was previously defined. An action that causes a variable to become defined does not imply that the variable was previously undefined.
- 2 Arrays, including sections, and variables of derived, character, or complex type are objects that consist of zero or more subobjects. Associations may be established between variables and subobjects and between subobjects of different variables. These subobjects may become defined or undefined.
- 3 An array is defined if and only if all of its elements are defined.
- 4 A derived-type scalar object is defined if and only if all of its nonpointer components are defined.
- 5 A complex or character scalar object is defined if and only if all of its subobjects are defined.
- 6 If an object is undefined, at least one (but not necessarily all) of its subobjects are undefined.

### 16.6.2 Variables that are always defined

- 1 Zero-sized arrays and zero-length strings are always defined.

### 16.6.3 Variables that are initially defined

- 1 The following variables are initially defined:
  - (1) variables specified to have initial values by **DATA statements**;
  - (2) variables specified to have initial values by **type declaration statements**;
  - (3) nonpointer **default-initialized subcomponents** of **saved** variables that do not have the **ALLOCATABLE** or **POINTER** attribute;
  - (4) pointers specified to be initially associated with a variable that is initially defined;
  - (5) variables that are always defined;
  - (6) variables with the **BIND attribute** that are initialized by means other than Fortran.

**NOTE 16.11**

Fortran code:

```
module mod
  integer, bind(c,name="blivet") :: foo
end module mod
```

C code:

```
int blivet = 123;
```

In the above example, the Fortran variable foo is initially defined to have the value 123 by means other than Fortran.

**16.6.4 Variables that are initially undefined**

- 1 Variables that are not initially defined are initially undefined.

**16.6.5 Events that cause variables to become defined**

- 1 Variables become defined by the following events.

- (1) Execution of an [intrinsic assignment statement](#) other than a [masked array assignment](#) or [FORALL assignment statement](#) causes the variable that precedes the equals to become defined.
- (2) Execution of a [masked array assignment](#) or [FORALL assignment statement](#) might cause some or all of the array elements in the [assignment statement](#) to become defined ([7.2.3](#)).
- (3) As execution of an [input statement](#) proceeds, each variable that is assigned a value from the input file becomes defined at the time that data are transferred to it. (See (4) in [16.6.6](#).) Execution of a [WRITE statement](#) whose [unit](#) specifier identifies an [internal file](#) causes each record that is written to become defined.
- (4) Execution of a [DO statement](#) causes the DO variable, if any, to become defined.
- (5) Beginning of execution of the action specified by an [io-implied-do](#) in a synchronous data transfer statement causes the [do-variable](#) to become defined.
- (6) A reference to a procedure causes an entire dummy data object to become defined if the dummy data object does not have [INTENT \(OUT\)](#) and the entire [effective argument](#) is defined.  
A reference to a procedure causes a subobject of a dummy argument to become defined if the dummy argument does not have [INTENT \(OUT\)](#) and the corresponding subobject of the [effective argument](#) is defined.
- (7) Execution of an input/output statement containing an [IOSTAT= specifier](#) causes the specified integer variable to become defined.
- (8) Execution of a synchronous [input statement](#) containing a [SIZE= specifier](#) causes the specified integer variable to become defined.
- (9) Execution of a wait operation ([9.7.1](#)) corresponding to an asynchronous [input statement](#) containing a [SIZE= specifier](#) causes the specified integer variable to become defined.
- (10) Execution of an [INQUIRE statement](#) causes any variable that is assigned a value during the execution of the statement to become defined if no error condition exists.

- 1           (11) If an error, end-of-file, or end-of-record condition occurs during execution of an input/output state-  
2           ment that has an **IOMSG= specifier**, the *iomsg-variable* becomes defined.
- 3           (12) When a **character storage unit** becomes defined, all associated **character storage units** become defined.  
4           When a **numeric storage unit** becomes defined, all associated **numeric storage units** of the same type  
5           become defined. When an entity of double precision real type becomes defined, all totally associated  
6           entities of double precision real type become defined.  
7           When an **unspecified storage unit** becomes defined, all associated **unspecified storage units** become  
8           defined.
- 9           (13) When a default complex entity becomes defined, all partially associated default real entities become  
10          defined.
- 11          (14) When both parts of a default complex entity become defined as a result of partially associated default  
12          real or default complex entities becoming defined, the default complex entity becomes defined.
- 13          (15) When all components of a structure of a **numeric sequence type** or **character sequence type** become  
14          defined as a result of partially associated objects becoming defined, the structure becomes defined.
- 15          (16) Execution of a statement with a **STAT= specifier** causes the variable specified by the **STAT= specifier**  
16          to become defined.
- 17          (17) If an error condition occurs during execution of a statement that has an **ERRMSG= specifier**, the  
18          variable specified by the **ERRMSG= specifier** becomes defined.
- 19          (18) Allocation of a zero-sized array or zero-length character variable causes the array or variable to  
20          become defined.
- 21          (19) Allocation of an object that has a nonpointer **default-initialized subcomponent**, except by an **AL-**  
22          **LOCATE statement** with a **SOURCE= specifier**, causes that **subcomponent** to become defined.
- 23          (20) Successful execution of an **ALLOCATE statement** with a **SOURCE= specifier** causes a subobject of  
24          the allocated object to become defined if the corresponding subobject of the **SOURCE=** expression  
25          is defined.
- 26          (21) Invocation of a procedure causes any **automatic data object** of zero size or zero character length in  
27          that procedure to become defined.
- 28          (22) When a pointer becomes associated with a **target** that is defined, the pointer becomes defined.
- 29          (23) Invocation of a procedure that contains an **unsaved** nonpointer nonallocatable local variable causes  
30          all nonpointer **default-initialized subcomponents** of the object to become defined.
- 31          (24) Invocation of a procedure that has a nonpointer nonallocatable **INTENT (OUT)** dummy argument  
32          causes all nonpointer **default-initialized subcomponents** of the dummy argument to become defined.
- 33          (25) In a **DO CONCURRENT** or **FORALL** construct, the *index-name* becomes defined when the *index-*  
34          *name* value set is evaluated.
- 35          (26) In a **DO CONCURRENT construct**, a variable with **LOCAL\_INIT** locality becomes defined at the  
36          beginning of each iteration.
- 37          (27) An object with the **VOLATILE attribute** that is changed by a means not specified by the program  
38          might become defined (see 5.5.19).
- 39          (28) Execution of the **BLOCK statement** of a **BLOCK construct** that has an **unsaved** nonpointer non-  
40          allocatable local variable causes all nonpointer **default-initialized subcomponents** of the variable to  
41          become defined.
- 42          (29) Execution of an **OPEN statement** containing a **NEWUNIT= specifier** causes the specified integer  
43          variable to become defined.

- (30) Execution of a **LOCK statement** containing an **ACQUIRED\_LOCK= specifier** causes the specified logical variable to become defined. If the logical variable becomes defined with the value true, the **lock variable** in the **LOCK statement** also becomes defined.
- (31) Successful execution of a **LOCK statement** that does not contain an **ACQUIRED\_LOCK= specifier** causes the **lock variable** to become defined.
- (32) Successful execution of an **UNLOCK statement** causes the **lock variable** to become defined.
- (33) Failure of an **image** that locked a **lock variable** without unlocking it causes the **lock variable** to become defined.
- (34) Successful execution of an **EVENT POST** or **EVENT WAIT** statement causes the **event variable** to become defined.

### 16.6.6 Events that cause variables to become undefined

1 Variables become undefined by the following events.

- (1) When a scalar variable of **intrinsic type** becomes defined, all totally associated variables of different type become undefined. When a double precision scalar variable becomes defined, all partially associated scalar variables become undefined. When a scalar variable becomes undefined, all partially associated double precision scalar variables become undefined.
  - (2) If the evaluation of a function would cause a variable to become defined and if a reference to the function appears in an expression in which the value of the function is not needed to determine the value of the expression, the variable becomes undefined when the expression is evaluated.
  - (3) When execution of an instance of a subprogram completes,
    - (a) its **unsaved** local variables become undefined,
    - (b) **unsaved** variables in a named **common block** that appears in the subprogram become undefined if they have been defined or redefined, unless another active **scoping unit** is referencing the **common block**, and
    - (c) a variable of type **C\_PTR** from the intrinsic module **ISO\_C\_BINDING** whose value is the **C address** of an **unsaved** local variable of the subprogram becomes undefined.
  - (4) When an error condition or end-of-file condition occurs during execution of an **input statement**, all of the variables specified by the input list or namelist group of the statement become undefined.
  - (5) When an error condition occurs during execution of an **output statement** in which the **unit** is an **internal file**, the **internal file** becomes undefined.
  - (6) When an error condition, end-of-file condition, or end-of-record condition occurs during execution of an input/output statement and the statement contains any *io-implied-dos*, all of the *do-variables* in the statement become undefined (9.11).
  - (7) Execution of a direct access **input statement** that specifies a record that has not been written previously causes all of the variables specified by the input list of the statement to become undefined.
  - (8) Execution of an **INQUIRE statement** might cause the **NAME=**, **RECL=**, and **NEXTREC=** variables to become undefined (9.10).
  - (9) When a **character storage unit** becomes undefined, all associated **character storage units** become undefined.
- When a **numeric storage unit** becomes undefined, all associated **numeric storage units** become undefined unless the undefinition is a result of defining an associated **numeric storage unit** of different type (see (1) above).

When an entity of double precision real type becomes undefined, all totally associated entities of double precision real type become undefined.

When an [unspecified storage unit](#) becomes undefined, all associated [unspecified storage units](#) become undefined.

- (10) When an [allocatable](#) entity is deallocated, it becomes undefined.
- (11) When the allocation transfer procedure ([13.9.137](#)) causes the allocation status of an [allocatable](#) entity to become unallocated, the entity becomes undefined.
- (12) Successful execution of an [ALLOCATE statement](#) with no [SOURCE= specifier](#) causes a [subcomponent](#) of an allocated object to become undefined if [default initialization](#) has not been specified for that [subcomponent](#).
- (13) Successful execution of an [ALLOCATE statement](#) with a [SOURCE= specifier](#) causes a subobject of the allocated object to become undefined if the corresponding subobject of the [SOURCE=](#) expression is undefined.
- (14) Execution of an [INQUIRE statement](#) causes all inquiry specifier variables to become undefined if an error condition exists, except for any variable in an [IOSTAT=](#) or [IOMSG=](#) specifier.
- (15) When a procedure is invoked
  - (a) an optional dummy argument that has no corresponding [actual argument](#) becomes undefined,
  - (b) a dummy argument with [INTENT \(OUT\)](#) becomes undefined except for any nonpointer [default-initialized subcomponents](#) of the argument,
  - (c) an [actual argument](#) corresponding to a dummy argument with [INTENT \(OUT\)](#) becomes undefined except for any nonpointer [default-initialized subcomponents](#) of the argument,
  - (d) a subobject of a dummy argument that does not have [INTENT \(OUT\)](#) becomes undefined if the corresponding subobject of the [effective argument](#) is undefined, and
  - (e) a variable that is the [function result](#) of that procedure becomes undefined except for any of its nonpointer [default-initialized subcomponents](#).
- (16) When the association status of a pointer becomes undefined or [disassociated](#) ([16.5.2.4](#), [16.5.2.5](#)), the pointer becomes undefined.
- (17) When a [DO CONCURRENT construct](#) terminates, a variable that is defined or becomes undefined during more than one iteration of the construct becomes undefined.
- (18) When execution of an iteration of a [DO CONCURRENT construct](#) completes, a construct entity of that construct which has [LOCAL](#) or [LOCAL\\_INIT](#) locality becomes undefined.
- (19) Execution of an asynchronous [READ statement](#) causes all of the variables specified by the input list or [SIZE= specifier](#) to become undefined. Execution of an asynchronous namelist [READ statement](#) causes any variable in the namelist group to become undefined if that variable will subsequently be defined during the execution of the [READ statement](#) or the corresponding wait operation ([9.7.1](#)).
- (20) When a variable with the [TARGET attribute](#) is deallocated, a variable of type [C\\_PTR](#) from the intrinsic module [ISO\\_C\\_BINDING](#) becomes undefined if its value is the [C address](#) of any part of the variable that is deallocated.
- (21) When a pointer is deallocated, a variable of type [C\\_PTR](#) from the intrinsic module [ISO\\_C\\_BINDING](#) becomes undefined if its value is the [C address](#) of any part of the [target](#) that is deallocated.
- (22) Execution of the allocation transfer procedure ([13.9.137](#)) where the argument TO does not have the [TARGET attribute](#) causes a variable of type [C\\_PTR](#) from the intrinsic module [ISO\\_C\\_BINDING](#) to become undefined if its value is the [C address](#) of any part of the argument FROM.

- (23) When a **BLOCK construct** completes execution,
  - its **unsaved** local variables become undefined, and
  - a variable of type **C\_PTR** from the intrinsic module **ISO\_C\_BINDING**, whose value is the **C address** of an **unsaved** local variable of the **BLOCK construct**, becomes undefined.
- (24) When execution of the **host instance** of the **target** of a variable of type **C\_FUNPTR** from the intrinsic module **ISO\_C\_BINDING** is completed by execution of a **RETURN** or **END** statement, the variable becomes undefined.
- (25) Execution of an intrinsic assignment of the type **C\_PTR** or **C\_FUNPTR** from the intrinsic module **ISO\_C\_BINDING** in which the variable and *expr* are not on the same **image** causes the variable to become undefined.
- (26) An object with the **VOLATILE attribute** (5.5.19) might become undefined by means not specified by the program.
- (27) When a pointer becomes associated with a target that is undefined, the pointer becomes undefined.
- (28) When an image fails during execution of a segment, a data object on a nonfailed image becomes undefined if it is not a **lock variable** and it might become undefined by execution of a statement of the segment other than an invocation of an **atomic subroutine** with the object as an **actual argument** corresponding to the **ATOM dummy argument**.

**NOTE 16.12**

Execution of a **defined assignment statement** could leave all or part of the variable undefined.

**16.6.7 Variable definition context**

- 1 Some variables are prohibited from appearing in a syntactic context that would imply definition or undefinition of the variable (5.5.10, 5.5.15, 12.7). The following are the contexts in which the appearance of a variable implies such definition or undefinition of the variable:

- (1) the *variable* of an *assignment-stmt*;
- (2) a *do-variable* in a *do-stmt* or *io-implied-do*;
- (3) an *input-item* in a *read-stmt*;
- (4) a *variable-name* in a *namelist-stmt* if the *namelist-group-name* appears in a **NML= specifier** in a *read-stmt*;
- (5) an *internal-file-variable* in a *write-stmt*;
- (6) an **IOSTAT=**, **SIZE=**, or **IOMSG=** specifier in an input/output statement;
- (7) a specifier in an **INQUIRE statement** other than **FILE=**, **ID=**, and **UNIT=**;
- (8) a **NEWUNIT=** specifier in an **OPEN statement**;
- (9) a *stat-variable*, *allocate-object*, or *errmsg-variable*;
- (10) an **actual argument** in a reference to a procedure with an **explicit interface** if the corresponding **dummy argument** is not a pointer and has **INTENT (OUT)** or **INTENT (INOUT)**;
- (11) a *variable* that is a *selector* in an **ASSOCIATE**, **SELECT RANK**, or **SELECT TYPE** construct if the corresponding **associate name** or any subobject thereof appears in a variable definition context;
- (12) a *coselector-name* in a **CHANGE TEAM statement** if the corresponding **associate name** appears in a variable definition context;
- (13) an *event-variable* in a **EVENT POST** or **EVENT WAIT** statement;

- (14) a *lock-variable* in a LOCK or UNLOCK statement;
- (15) a *scalar-logical-variable* in an ACQUIRED\_LOCK= specifier;
- (16) a *team-variable* in a FORM TEAM statement.

2 If a reference to a function appears in a variable definition context the result of the function reference shall be a pointer that is associated with a definable target. That target is the variable that becomes defined or undefined.

### 16.6.8 Pointer association context

1 Some pointers are prohibited from appearing in a syntactic context that would imply alteration of the pointer association status (16.5.2.2, 5.5.10, 5.5.15, 12.7). The following are the contexts in which the appearance of a pointer implies such alteration of its pointer association status:

- a *pointer-object* in a *nullify-stmt*;
- a *data-pointer-object* or *proc-pointer-object* in a *pointer-assignment-stmt*;
- an *allocate-object* in an *allocate-stmt* or *deallocate-stmt*;
- an *actual argument* in a reference to a procedure if the corresponding *dummy argument* is a pointer with the INTENT (OUT) or INTENT (INOUT) attribute.



# Annex A

(Informative)

## Processor Dependencies

### A.1 Unspecified Items

1 This part of ISO/IEC 1539 does not specify the following:

- the properties excluded in 1.1;
- a processor's error detection capabilities beyond those listed in 1.5;
- which additional intrinsic procedures or modules a processor provides (1.5);
- the number and kind of companion processors (2.5.7);
- the number of representation methods and associated kind type parameter values of the intrinsic types (4.4), except that there shall be at least two representation methods for type real, and a representation method of type complex that corresponds to each representation method for type real.

### A.2 Processor Dependencies

1 According to this part of ISO/IEC 1539, the following are processor dependent:

- the order of evaluation of the specification expressions within the specification part of an invoked Fortran procedure (2.3.5);
- how soon an image terminates if another image initiates error termination (2.3.5);
- the value of a reference to a coindexed object on a failed image (2.3.6);
- the conditions that cause an image to fail (2.3.6);
- whether the processor supports a concept of process exit status, and if so, the process exit status on program termination (2.3.7);
- the mechanism of a companion processor, and the means of selecting between multiple companion processors (2.5.7);
- the processor character set (3.1);
- the means for specifying the source form of a program unit (3.3);
- the maximum number of characters allowed on a source line containing characters not of default kind (3.3.2, 3.3.3);
- the maximum depth of nesting of include lines (3.4);
- the interpretation of the *char-literal-constant* in the include line (3.4);
- the set of values supported by an intrinsic type, other than logical (4.1.3);
- the kind of a character length type parameter (4.4.4.1);
- the blank padding character for nondefault character kind (4.4.4.2)
- whether particular control characters can appear within a character literal constant in fixed source form (4.4.4.3);
- the collating sequence for each character set (4.4.4.4);

- 1 • the order of [finalization](#) of components of objects of derived type (4.5.6.2);
- 2 • the order of [finalization](#) when several objects are finalized as the consequence of a single event (4.5.6.2);
- 3 • whether and when an object is finalized if it is allocated by pointer allocation and it later becomes un-
- 4 reachable due to all pointers associated with the object having their pointer association status changed
- 5 ([4.5.6.3](#));
- 6 • whether an object is finalized by a deallocation in which an error condition occurs ([4.5.6.3](#));
- 7 • the kind type parameter of each enumeration and its enumerators ([4.6](#));
- 8 • whether an array is contiguous, except as specified in [5.5.7](#);
- 9 • the set of error conditions that can occur in [ALLOCATE](#) and [DEALLOCATE](#) statements ([6.7.1](#), [6.7.3](#));
- 10 • the allocation status of a variable after evaluation of an expression if the evaluation of a function would
- 11 change the allocation status of the variable and if a reference to the function appears in the expression in
- 12 which the value of the function is not needed to determine the value of the expression ([6.7.1.3](#));
- 13 • the order of deallocation when several objects are deallocated by a [DEALLOCATE statement](#) ([6.7.3](#));
- 14 • the order of deallocation when several objects are deallocated due to the occurrence of an event described
- 15 in [6.7.3.2](#);
- 16 • whether an allocated allocatable subobject is deallocated when an error condition occurs in the deallocation
- 17 of an object ([6.7.3.2](#));
- 18 • the positive integer values assigned to the *stat-variable* in a [STAT= specifier](#) as the result of an error
- 19 condition ([6.7.4](#), [8.6.11](#));
- 20 • the allocation status or pointer association status of an *allocate-object* if an error occurs during execution
- 21 of an [ALLOCATE](#) or [DEALLOCATE](#) statement ([6.7.4](#));
- 22 • the value assigned to the *errmsg-variable* in an [ERRMSG= specifier](#) as the result of an error condition
- 23 ([6.7.5](#), [8.6.11](#));
- 24 • the kind type parameter value of the result of a numeric intrinsic binary operation where
  - 25 – both operands are of type integer but with different kind type parameters, and the decimal exponent
  - 26 ranges are the same,
  - 27 – one operand is of type real or complex and the other is of type real or complex with a different kind
  - 28 type parameter, and the decimal precisions are the same,
  - 29 and for a logical intrinsic binary operation where the operands have different kind type parameters ([7.1.9.3](#));
- 30 • the character assigned to the variable in an [intrinsic assignment statement](#) if the kind of the expression is
- 31 different and the character is not representable in the kind of the variable ([7.2.1.3](#));
- 32 • the order of evaluation of the [specification expressions](#) within the specification part of a [BLOCK construct](#)
- 33 when the construct is executed ([8.1.4](#));
- 34 • the pointer association status of a pointer that has its pointer association changed in more than one iteration
- 35 of a [DO CONCURRENT construct](#), on termination of the construct ([8.1.7](#));
- 36 • the ordering between records written by different iterations of a [DO CONCURRENT construct](#) if the records
- 37 are written to a file connected for sequential access by more than one iteration ([8.1.7](#));
- 38 • the manner in which the stop code of a [STOP](#) or [ERROR STOP](#) statement is made available ([8.4](#));
- 39 • the mechanisms available for creating dependencies for cooperative synchronization ([8.6.5](#));
- 40 • the value of the count of the event variable in an [EVENT POST](#) or [EVENT WAIT](#) statement if an error
- 41 condition occurs ([8.6.7](#), [8.6.8](#));
- 42 • the [image index](#) value established for each image in a [team](#) by a [FORM TEAM statement](#) without a
- 43 [NEW\\_INDEX= specifier](#) ([8.6.9](#));

- the set of error conditions that can occur in [image control statements](#) (8.6.11);
- the relationship between the [file storage units](#) when viewing a file as a [stream file](#), and the records when viewing that file as a [record file](#) (9);
- whether particular control characters can appear in a formatted [record](#) or a formatted [stream file](#) (9.2.2);
- the form of values in an unformatted record (9.2.3);
- at any time, the set of allowed access methods, set of allowed forms, set of allowed actions, and set of allowed record lengths for a [file](#) (9.3);
- the set of allowable names for a [file](#) (9.3);
- whether a named file on one [image](#) is the same as a file with the same name on another [image](#) (9.3.1);
- the set of [external files](#) that exist for a program (9.3.2);
- the relationship between positions of successive [file storage units](#) in an [external file](#) that is connected for formatted stream access (9.3.3.4);
- the [external unit](#) preconnected for sequential formatted input and identified by an asterisk or the [named constant](#) INPUT\_UNIT of the ISO\_FORTRAN\_ENV intrinsic module (9.5);
- the [external unit](#) preconnected for sequential formatted output and identified by an asterisk or the [named constant](#) OUTPUT\_UNIT of the ISO\_FORTRAN\_ENV intrinsic module (9.5);
- the [external unit](#) preconnected for sequential formatted output and identified by the [named constant](#) ERROR\_UNIT of the ISO\_FORTRAN\_ENV intrinsic module, and whether this unit is the same as OUTPUT\_UNIT (9.5);
- at any time, the set of [external units](#) that exist for an [image](#) (9.5.3);
- whether a [unit](#) can be connected to a file that is also connected to a C stream (9.5.4);
- whether a file can be connected to more than one [unit](#) at the same time (9.5.4);
- the effect of performing input/output operations on multiple units while they are connected to the same [external file](#) (9.5.4);
- the result of performing input/output operations on a [unit](#) connected to a file that is also connected to a C stream (9.5.4);
- whether the files connected to the [units](#) INPUT\_UNIT, OUTPUT\_UNIT, and ERROR\_UNIT correspond to the predefined C text streams standard input, standard output, and standard error, respectively (9.5.4);
- the results of performing input/output operations on an [external file](#) both from Fortran and from a procedure defined by means other than Fortran (9.5.4);
- the default value for the [ACTION= specifier](#) in an [OPEN statement](#) (9.5.6.4);
- the encoding of a file opened with [ENCODING='DEFAULT'](#) (9.5.6.9);
- the file connected by an [OPEN statement](#) with [STATUS='SCRATCH'](#) (9.5.6.10);
- the interpretation of case in a file name (9.5.6.10, 9.10.2.2);
- the position of a file after executing an [OPEN statement](#) with a [POSITION= specifier](#) of ASIS, when the file previously existed but was not connected (9.5.6.14);
- the default value for the [RECL= specifier](#) in an [OPEN statement](#) (9.5.6.15);
- the effect of [RECL=](#) on a record containing any nondefault characters (9.5.6.15);
- the default input/output rounding mode (9.5.6.16);
- the default sign mode (9.5.6.17);
- the file status when [STATUS='UNKNOWN'](#) is specified in an [OPEN statement](#) (9.5.6.18);
- the value assigned to the variable in the [ID= specifier](#) in an asynchronous [data transfer statement](#) when

- execution of the statement is successfully completed (9.6.2.9);
- whether `POS=` is permitted with particular files, and whether `POS=` can position a particular file to a position prior to its current position (9.6.2.11);
  - the form in which a single value of derived type is treated in an unformatted input/output statement if the *effective item* is not processed by a *defined input/output* procedure (9.6.3);
  - the result of unformatted input when the value stored in the file has a different type or *type parameters* from that of the input list item (9.6.4.5.2);
  - the negative value of the *unit* argument to a *defined input/output* procedure if the parent *data transfer statement* accesses an *internal file* (9.6.4.8.3);
  - the manner in which the processor makes the value of the *iomsg* argument of a *defined input/output* procedure available if the procedure assigns a nonzero value to the *iostat* argument and the processor therefore *terminates execution* of the program (9.6.4.8.3);
  - the action caused by the flush operation, whether the processor supports the flush operation for the specified *unit*, and the negative value assigned to the `IOSTAT=` variable if the processor does not support the flush operation for the specified *unit* (9.9);
  - the case of characters assigned to the variable in a `NAME= specifier` in an *INQUIRE statement* (9.10.2.15);
  - which of the connected external unit numbers is assigned to the *scalar-int-variable* in the `NUMBER= specifier` in an *INQUIRE by file statement*, if more than one unit on an *image* is connected to the file (9.10.2.18);
  - the value of the variable in a `POSITION= specifier` in an *INQUIRE statement* if the file has been repositioned since connection (9.10.2.23);
  - the relationship between file size and the data stored in records in a sequential or direct access file (9.10.2.30);
  - the number of *file storage units* needed to store data in an unformatted file (9.10.3);
  - the set of error conditions that can occur in input/output statements (9.11.1);
  - when an input/output error condition occurs or is detected (9.11.1);
  - the positive integer value assigned to the variable in an `IOSTAT= specifier` as the result of an error condition (9.11.5);
  - the value assigned to the variable in an `IOMSG= specifier` as the result of an error condition (9.11.6);
  - the result of output of non-representable characters to a Unicode file (10.7.1);
  - the interpretation of the optional non-blank characters within the parentheses of a real NaN input field (10.7.2.3.2);
  - the interpretation of a sign in a NaN input field (10.7.2.3.2);
  - for output of an IEEE NaN, whether after the letters 'NaN', the processor produces additional alphanumeric characters enclosed in parentheses (10.7.2.3.2);
  - the effect of the input/output rounding mode `PROCESSOR_DEFINED` (10.7.2.3.8);
  - which value is chosen if the input/output rounding mode is `NEAREST` and the value to be converted is exactly halfway between the two nearest representable values in the result format (10.7.2.3.8);
  - the field width, decimal part width, and exponent width used for the `G0` edit descriptor (10.7.5);
  - the file position when position editing skips a character of nondefault kind in an *internal file* of default character kind or an *external unit* that is not connected to a Unicode file (10.8.1.1);
  - when the sign mode is `PROCESSOR_DEFINED`, whether a plus sign appears in a numeric output field for a nonnegative value (10.8.4);
  - the results of list-directed output (10.10.4);

- 1 • the results of namelist output (10.11.4);
- 2 • the interaction between [argument association](#) and pointer association, (12.5.2.4);
- 3 • the values returned by some intrinsic functions (13);
- 4 • how the sequences of atomic actions in unordered segments interleave (13.5);
- 5 • the value assigned to a STAT argument in a reference to an atomic subroutine when an error condition
- 6 occurs (13.5);
- 7 • the effects of calling [COMMAND\\_ARGUMENT\\_COUNT](#), [EXECUTE\\_COMMAND\\_LINE](#), [GET\\_COM-](#)
- 8 [MAND](#), and [GET\\_COMMAND\\_ARGUMENT](#) on any [image](#) other than [image 1](#) in the [initial team](#) (13.7);
- 9 • whether each [image](#) uses a separate random number generator, or if some or all [images](#) use common random
- 10 number generators (13.7);
- 11 • whether the results returned from [CPU\\_TIME](#), [DATE\\_AND\\_TIME](#) and [SYSTEM\\_CLOCK](#) are depend-
- 12 ent on which [image](#) calls them (13.7);
- 13 • the set of error conditions that can occur in some intrinsic subroutines (13.9);
- 14 • the value assigned to a CMDSTAT, ERRMSG, EXITSTAT, STAT, or STATUS argument to indicate a
- 15 processor-dependent error condition (13.9);
- 16 • the computed value of the the intrinsic subroutine [CO\\_REDUCE](#) (13.9.49) and the intrinsic subroutine
- 17 [CO\\_SUM](#) (13.9.50);
- 18 • the value assigned to the TIME argument by the intrinsic subroutine [CPU\\_TIME](#) (13.9.57);
- 19 • whether date, clock, and time zone information is available (13.9.59);
- 20 • whether date, clock, and time zone information on one [image](#) is the same as that on another [image](#) (13.9.59);
- 21 • the value of command argument zero, if the processor does not support the concept of a command name
- 22 (13.9.83);
- 23 • the order of command arguments (13.9.83);
- 24 • whether the significant length of a command argument includes trailing blanks (13.9.83);
- 25 • the interpretation of case for the NAME argument of the intrinsic subroutine [GET\\_ENVIRONMENT\\_-](#)
- 26 [VARIABLE](#) (13.9.84);
- 27 • whether an environment variable that exists on an [image](#) also exists on another [image](#), and if it does exist
- 28 on both [images](#), whether the values are the same or different (13.9.84);
- 29 • the value assigned to the pseudorandom number generator by the intrinsic subroutine [RANDOM\\_INIT](#)
- 30 (13.9.155);
- 31 • the computation of the seed value used by the pseudorandom number generator (13.9.157);
- 32 • on [images](#) that use a common random number generator, the interleaving of values assigned by [RANDOM\\_-](#)
- 33 [NUMBER](#) in unordered segments(13.7);
- 34 • the value assigned to the seed by the intrinsic subroutine [RANDOM\\_SEED](#) when no argument is present
- 35 (13.9.157);
- 36 • the values assigned to its arguments by the intrinsic subroutine [SYSTEM\\_CLOCK](#) (13.9.186);
- 37 • the values of the named constants in the intrinsic module [ISO\\_FORTRAN\\_ENV](#)(13.10.2);
- 38 • the values returned by the functions [COMPILER\\_OPTIONS](#) and [COMPILER\\_VERSION](#) in the intrinsic
- 39 module [ISO\\_FORTRAN\\_ENV](#)(13.10.2);
- 40 • the extent to which a processor supports IEEE arithmetic (14);
- 41 • whether a flag that is quiet on entry to a scoping unit that does not access [IEEE\\_FEATURES](#), [IEEE\\_-](#)
- 42 [EXCEPTIONS](#), or [IEEE\\_ARITHMETIC](#) is signaling on exit (14.1);

- 1 • the conditions under which IEEE\_OVERFLOW is raised in a calculation involving non-ISO/IEC/IEEE
- 2 60559:2011 floating-point data (14.3);
- 3 • the conditions under which IEEE\_OVERFLOW and IEEE\_DIVIDE\_BY\_ZERO are raised in a floating-
- 4 point exponentiation operation (14.3);
- 5 • the conditions under which IEEE\_DIVIDE\_BY\_ZERO is raised in a calculation involving non-ISO/IEC/IEEE
- 6 60559:2011 floating-point data (14.3);
- 7 • whether an exception signals at the end of a sequence of statements that has no invocations of IEEE\_GET\_-
- 8 FLAG, IEEE\_SET\_FLAG, IEEE\_GET\_STATUS, IEEE\_SET\_STATUS, or IEEE\_SET\_HALTING\_-
- 9 MODE, in which execution of an operation would cause it to signal, if no value of a variable depends upon
- 10 the result of the operation (14.3);
- 11 • the initial rounding modes (14.4);
- 12 • the initial underflow mode (14.5);
- 13 • the initial halting mode (14.6);
- 14 • whether IEEE\_INT implements the convertToInteger{round} or convertToIntegerExact{round} operation
- 15 specified by ISO/IEC/IEEE 60559:2011 (14.11.11);
- 16 • which argument is the result value of IEEE\_MAX\_NUM, IEEE\_MAX\_NUM\_MAG, IEEE\_MIN\_NUM,
- 17 or IEEE\_MIN\_NUM\_MAG when both arguments are quiet NaNs (14.11.17, 14.11.18, 14.11.19, 14.11.20);
- 18 • the requirements on the storage sequence to be associated with the pointer FPTR by the C\_F\_POINTER
- 19 subroutine (15.2.3.4);
- 20 • the order of the members of the CFI\_dim\_t structure defined in the source file CFI\_Fortran\_binding.h
- 21 (15.5.2);
- 22 • members of the CFI\_cdesc\_t structure defined in the source file CFI\_Fortran\_binding.h beyond the re-
- 23 quirements of 15.5.3;
- 24 • the value of CFI\_MAX\_RANK in the source file CFI\_Fortran\_binding.h (15.5.4);
- 25 • the value of CFI\_VERSION in the source file CFI\_Fortran\_binding.h (15.5.4);
- 26 • which error condition is detected if more than one error condition could be detected for an invocation of
- 27 one of the functions declared in the source file CFI\_Fortran\_binding.h (15.5.5.1);
- 28 • the values of the attribute specifier macros defined in the source file CFI\_Fortran\_binding.h (15.5.4);
- 29 • the values of the type specifier macros defined in the source file CFI\_Fortran\_binding.h;
- 30 • which additional type specifier values are defined in the source file CFI\_Fortran\_binding.h (15.5.4);
- 31 • the values of the error code macros other than CFI\_SUCCESS that are defined in the source file CFI\_-
- 32 Fortran\_binding.h (15.5.4);
- 33 • the base address of a zero-sized array (15.5.3);
- 34 • the values of the floating-point exception flags on entry to a procedure defined by means other than Fortran
- 35 (15.10.3);
- 36 • whether a procedure defined by means other than Fortran is an asynchronous communication initiation or
- 37 completion procedure (15.10.4).

## Annex B

(Informative)

### Deleted and obsolescent features

#### B.1 Deleted features from Fortran 90

1 These deleted features are those features of Fortran 90 that were redundant and considered largely unused.

2 The following Fortran 90 features are not required.

- (1) Real and double precision DO variables.

In FORTRAN 77 and Fortran 90, a DO variable was allowed to be of type real or double precision in addition to type integer; this has been deleted. A similar result can be achieved by using a **DO construct** with no loop control and the appropriate exit test.

- (2) Branching to an **END IF statement** from outside its block.

In FORTRAN 77 and Fortran 90, it was possible to branch to an **END IF statement** from outside the **IF construct**; this has been deleted. A similar result can be achieved by branching to a **CONTINUE statement** that is immediately after the **END IF statement**.

- (3) PAUSE statement.

The PAUSE statement, provided in FORTRAN 66, FORTRAN 77, and Fortran 90, has been deleted. A similar result can be achieved by writing a message to the appropriate **unit**, followed by reading from the appropriate **unit**.

- (4) ASSIGN and assigned GO TO statements, and assigned format specifiers.

The ASSIGN statement and the related assigned GO TO statement, provided in FORTRAN 66, FORTRAN 77, and Fortran 90, have been deleted. Further, the ability to use an assigned integer as a format, provided in FORTRAN 77 and Fortran 90, has been deleted. A similar result can be achieved by using other control constructs instead of the assigned GO TO statement and by using a default character variable to hold a format specification instead of using an assigned integer.

- (5) H edit descriptor.

In FORTRAN 77 and Fortran 90, there was an alternative form of character string edit descriptor, which had been the only such form in FORTRAN 66; this has been deleted. A similar result can be achieved by using a character string edit descriptor.

- (6) Vertical format control.

In FORTRAN 66, FORTRAN 77, Fortran 90, and Fortran 95 formatted output to certain **units** resulted in the first character of each record being interpreted as controlling vertical spacing. There was no standard way to detect whether output to a **unit** resulted in this vertical format control, and no way to specify that it should be applied; this has been deleted. The effect can be achieved by post-processing a formatted file.

3 See ISO/IEC 1539:1991 for detailed rules of how these deleted features worked.

## B.2 Deleted features from Fortran 2008

1 These deleted features are those features of Fortran 2008 that were redundant and considered largely unused.

2 The following Fortran 2008 features are not required.

(1) Arithmetic IF statement.

The arithmetic IF statement is incompatible with ISO/IEC/IEEE 60559:2011 and necessarily involves the use of [statement labels](#); [statement labels](#) can hinder optimization, and make code hard to read and maintain. Similar logic can be more clearly encoded using other conditional statements.

(2) Nonblock DO construct

The nonblock forms of the DO loop were confusing and hard to maintain. Shared termination and dual use of labeled action statements as do termination and branch targets were especially error-prone.

## B.3 Obsolescent features

### B.3.1 General

1 The obsolescent features are those features of Fortran 90 that were redundant and for which better methods were available in Fortran 90. Subclause [1.7.3](#) describes the nature of the obsolescent features. The obsolescent features in this part of ISO/IEC 1539 are the following.

(1) Alternate return — see [B.3.2](#).

(2) [Computed GO TO](#) — see [B.3.3](#).

(3) Statement functions — see [B.3.4](#).

(4) [DATA statements](#) amongst executable statements — see [B.3.5](#).

(5) Assumed length character functions — see [B.3.6](#).

(6) Fixed form source — see [B.3.7](#).

(7) CHARACTER\* form of CHARACTER declaration — see [B.3.8](#).

(8) [ENTRY statements](#) — see [B.3.9](#).

(9) Label form of [DO statement](#) — see [B.3.10](#).

(10) [COMMON](#) and [EQUIVALENCE](#) statements, and the [block data program unit](#) — see [B.3.11](#).

(11) Specific names for intrinsic functions — see [B.3.12](#).

(12) [FORALL](#) construct and statement — see [B.3.13](#)

### B.3.2 Alternate return

1 An alternate return introduces labels into an argument list to allow the called procedure to direct the execution of the caller upon return. The same effect can be achieved with a return code that is used in a [SELECT CASE construct](#) on return. This avoids an irregularity in the syntax and semantics of [argument association](#). For example,

```
CALL SUBR_NAME (X, Y, Z, *100, *200, *300)
```

can be replaced by

```
CALL SUBR_NAME (X, Y, Z, RETURN_CODE)
```

```
SELECT CASE (RETURN_CODE)
```

```

1      CASE (1)
2          ...
3      CASE (2)
4          ...
5      CASE (3)
6          ...
7      CASE DEFAULT
8          ...
9  END SELECT

```

### B.3.3 Computed GO TO statement

- 1 The [computed GO TO statement](#) has been superseded by the [SELECT CASE construct](#), which is a generalized, easier to use, and clearer means of expressing the same computation.

### B.3.4 Statement functions

- 1 [Statement functions](#) are subject to a number of nonintuitive restrictions and are a potential source of error because their syntax is easily confused with that of an assignment statement.
- 2 The internal function is a more generalized form of the [statement function](#) and completely supersedes it.

### B.3.5 DATA statements among executables

- 1 The statement ordering rules allow [DATA statements](#) to appear anywhere in a [program unit](#) after the specification statements. The ability to position [DATA statements](#) amongst executable statements is very rarely used, unnecessary, and a potential source of error.

### B.3.6 Assumed character length functions

- 1 Assumed character length for functions is an irregularity in the language in that elsewhere in Fortran the philosophy is that the attributes of a function result depend only on the [actual arguments](#) of the invocation and on any data accessible by the function through host or use association. Some uses of this facility can be replaced with an [automatic](#) character length function, where the length of the function result is declared in a [specification expression](#). Other uses can be replaced by the use of a subroutine whose arguments correspond to the function result and the function arguments.
- 2 Note that dummy arguments of a function can have assumed character length.

### B.3.7 Fixed form source

- 1 Fixed form source was designed when the principal machine-readable input medium for new programs was punched cards. Now that new and amended programs are generally entered via keyboards with screen displays, it is an unnecessary overhead, and is potentially error-prone, to have to locate positions 6, 7, or 72 on a line. Free form source was designed expressly for this more modern technology.
- 2 It is a simple matter for a software tool to convert from fixed to free form source.

### B.3.8 CHARACTER\* form of CHARACTER declaration

- 1 In addition to the CHARACTER\**char-length* form introduced in FORTRAN 77, Fortran 90 provided the CHARACTER([ LEN = ] *type-param-value*) form. The older form (CHARACTER\**char-length*) is redundant.

### B.3.9 ENTRY statements

- 1 *ENTRY statements* allow more than one entry point to a subprogram, facilitating sharing of data items and executable statements local to that subprogram.
- 2 This can be replaced by a module containing the (private) data items, with a module procedure for each entry point and the shared code in a private module procedure.

### B.3.10 Label DO statement

- 1 The label in the *DO statement* is redundant with the construct name. Furthermore, the label allows unrestricted branches and, for its main purpose (the target of a conditional branch to skip the rest of the current iteration), is redundant with the *CYCLE statement*, which is clearer.

### B.3.11 COMMON and EQUIVALENCE statements and the block data program unit

- 1 *Common blocks* are error-prone and have largely been superseded by *modules*. *EQUIVALENCE* similarly is error-prone. Whilst use of these statements was invaluable prior to Fortran 90 they are now redundant and can inhibit performance. The *block data program unit* exists only to serve *common blocks* and hence is also redundant.

### B.3.12 Specific names for intrinsic functions

- 1 The *specific names of the intrinsic functions* are often obscure and hinder portability. They have been redundant since Fortran 90. Use generic names for references to intrinsic procedures.

### B.3.13 FORALL construct and statement

- 1 The *FORALL* construct and statement were added to the language in the expectation that they would enable highly efficient execution, especially on parallel processors. However, experience indicates that they are too complex and have too many restrictions for compilers to take advantage of them. They are redundant with the *DO CONCURRENT construct*, and many of the manipulations for which they might be used can be done more effectively using pointers, especially using pointer rank remapping.

# Annex C

(Informative)

## Extended notes

### C.1 Clause 4 notes

#### C.1.1 Selection of the approximation methods (4.4.3.2)

- 1 One can select the real approximation method for an entire program through the use of a module and the parameterized real type. This is accomplished by defining a named integer constant to have a particular kind type parameter value and using that [named constant](#) in all real, complex, and derived-type declarations. For example, the specification statements

```
INTEGER, PARAMETER :: LONG_FLOAT = 8
REAL (LONG_FLOAT) X, Y
COMPLEX (LONG_FLOAT) Z
```

specify that the approximation method corresponding to a kind type parameter value of 8 is supplied for the data objects X, Y, and Z in the [program unit](#). The kind type parameter value LONG\_FLOAT can be made available to an entire program by placing the INTEGER specification statement in a module and accessing the [named constant](#) LONG\_FLOAT with a [USE statement](#). Note that by changing 8 to 4 once in the module, a different approximation method is selected.

- 2 To avoid the use of the processor-dependent values 4 or 8, replace 8 by [KIND \(0.0\)](#) or [KIND \(0.0D0\)](#). Another way to avoid these processor-dependent values is to select the kind value using the intrinsic function [SELECTED\\_REAL\\_KIND \(13.9.170\)](#). In the above specification statement, the 8 might be replaced by, for instance, [SELECTED\\_REAL\\_KIND \(10, 50\)](#), which requires an approximation method to be selected with at least 10 decimal digits of precision and a range from  $10^{-50}$  to  $10^{50}$ . There are no magnitude or ordering constraints placed on kind values, in order that implementers have flexibility in assigning such values and can add new kinds without changing previously assigned kind values.

- 3 As kind values have no portable meaning, a good practice is to use them in programs only through [named constants](#) as described above (for example, SINGLE, IEEE\_SINGLE, DOUBLE, and QUAD), rather than using the kind values directly.

#### C.1.2 Type extension and component accessibility (4.5.2.2, 4.5.4)

- 1 The default accessibility of the components of an [extended type](#) can be specified in the type definition. The accessibility of its components can be specified individually. For example:

```
module types
  type base_type
    private                !-- Sets default accessibility
    integer :: i            !-- a private component
    integer, private :: j  !-- another private component
```

```

1      integer, public :: k    !-- a public component
2  end type base_type
3
4      type, extends(base_type) :: my_type
5      private                !-- Sets default for components declared in my_type
6      integer :: l            !-- A private component.
7      integer, public :: m    !-- A public component.
8  end type my_type
9
10 end module types
11
12 subroutine sub
13 use types
14 type (my_type) :: x
15
16 ...
17
18 call another_sub( &
19     x%base_type, &    !-- ok because base_type is a public subobject of x
20     x%base_type%k, &  !-- ok because x%base_type is ok and has k as a
21                       !-- public component.
22     x%k, &            !-- ok because it is shorthand for x%base_type%k
23     x%base_type%i, &  !-- Invalid because i is private.
24     x%i)              !-- Invalid because it is shorthand for x%base_type%i
25 end subroutine sub

```

### 26 C.1.3 Generic type-bound procedures (4.5.5)

27 Example of a derived type with generic type-bound procedures:

- 28 1 The only difference between this example and the same thing rewritten to use generic [interface blocks](#) is that  
29 with [type-bound procedures](#),

```
30 USE rational_numbers, ONLY: rational
```

31 does not block the [type-bound procedures](#); the user still gets access to the defined assignment and extended  
32 operations.

```

33 MODULE rational_numbers
34 IMPLICIT NONE
35 PRIVATE
36 TYPE, PUBLIC :: rational
37     PRIVATE
38     INTEGER n,d
39 CONTAINS
40     ! ordinary type-bound procedure
41     PROCEDURE :: real => rat_to_real
42     ! specific type-bound procedures for generic support

```

```

1      PROCEDURE,PRIVATE :: rat_asgn_i, rat_plus_i, rat_plus_rat => rat_plus
2      PROCEDURE,PRIVATE,PASS(b) :: i_plus_rat
3      ! generic type-bound procedures
4      GENERIC :: ASSIGNMENT(=) => rat_asgn_i
5      GENERIC :: OPERATOR(+) => rat_plus_rat, rat_plus_i, i_plus_rat
6  END TYPE
7  CONTAINS
8      ELEMENTAL REAL FUNCTION rat_to_real(this) RESULT(r)
9          CLASS(rational),INTENT(IN) :: this
10         r = REAL(this%n)/this%d
11     END FUNCTION
12     ELEMENTAL SUBROUTINE rat_asgn_i(a,b)
13         CLASS(rational),INTENT(OUT) :: a
14         INTEGER,INTENT(IN) :: b
15         a%n = b
16         a%d = 1
17     END SUBROUTINE
18     ELEMENTAL TYPE(rational) FUNCTION rat_plus_i(a,b) RESULT(r)
19         CLASS(rational),INTENT(IN) :: a
20         INTEGER,INTENT(IN) :: b
21         r%n = a%n + b*a%d
22         r%d = a%d
23     END FUNCTION
24     ELEMENTAL TYPE(rational) FUNCTION i_plus_rat(a,b) RESULT(r)
25         INTEGER,INTENT(IN) :: a
26         CLASS(rational),INTENT(IN) :: b
27         r%n = b%n + a*b%d
28         r%d = b%d
29     END FUNCTION
30     ELEMENTAL TYPE(rational) FUNCTION rat_plus(a,b) RESULT(r)
31         CLASS(rational),INTENT(IN) :: a,b
32         r%n = a%n*b%d + b%n*a%d
33         r%d = a%d*b%d
34     END FUNCTION
35 END

```

### 36 C.1.4 Abstract types (4.5.7.1)

- 37 1 The following illustrates how an abstract type can be used as the basis for a collection of related types, and how  
38 a non-abstract member of that collection can be created by type extension.

```

39     TYPE, ABSTRACT :: DRAWABLE_OBJECT
40         REAL, DIMENSION(3) :: RGB_COLOR = (/1.0,1.0,1.0/) ! White
41         REAL, DIMENSION(2) :: POSITION = (/0.0,0.0/) ! Centroid
42     CONTAINS
43         PROCEDURE(RENDER_X), PASS(OBJECT), DEFERRED :: RENDER
44     END TYPE DRAWABLE_OBJECT

```

```

1
2      ABSTRACT INTERFACE
3          SUBROUTINE RENDER_X(OBJECT, WINDOW)
4              IMPORT DRAWABLE_OBJECT, X_WINDOW
5              CLASS(DRAWABLE_OBJECT), INTENT(IN) :: OBJECT
6              CLASS(X_WINDOW), INTENT(INOUT) :: WINDOW
7          END SUBROUTINE RENDER_X
8      END INTERFACE
9
10     ...
11
12     TYPE, EXTENDS(DRAWABLE_OBJECT) :: DRAWABLE_TRIANGLE ! Not ABSTRACT
13         REAL, DIMENSION(2,3) :: VERTICES ! In relation to centroid
14     CONTAINS
15         PROCEDURE, PASS(OBJECT) :: RENDER=>RENDER_TRIANGLE_X
16     END TYPE DRAWABLE_TRIANGLE

```

17 2 The actual drawing procedure will draw a triangle in WINDOW with vertices at x and y coordinates at  
18 OBJECT%POSITION(1)+OBJECT%VERTICES(1,1:3) and OBJECT%POSITION(2)+OBJECT%VERTICES(2,1:3):

```

19      SUBROUTINE RENDER_TRIANGLE_X(OBJECT, WINDOW)
20          CLASS(DRAWABLE_TRIANGLE), INTENT(IN) :: OBJECT
21          CLASS(X_WINDOW), INTENT(INOUT) :: WINDOW
22          ...
23      END SUBROUTINE RENDER_TRIANGLE_X

```

## 24 C.1.5 Structure constructors and generic names (4.5.10)

25 1 A generic name can be the same as a type name. This can be used to emulate user-defined [structure constructors](#)  
26 for that type, even if the type has private components. For example:

```

27      MODULE mytype_module
28          TYPE mytype
29              PRIVATE
30              COMPLEX value
31              LOGICAL exact
32          END TYPE
33          INTERFACE mytype
34              MODULE PROCEDURE int_to_mytype
35          END INTERFACE
36          ! Operator definitions etc.
37          ...
38      CONTAINS
39          TYPE(mytype) FUNCTION int_to_mytype(i)
40              INTEGER, INTENT(IN) :: i
41              int_to_mytype%value = i
42              int_to_mytype%exact = .TRUE.

```

```

1      END FUNCTION
2      ! Procedures to support operators etc.
3      ...
4  END
5
6  PROGRAM example
7      USE mytype_module
8      TYPE(mytype) x
9      x = mytype(17)
10     END

```

11    2 The type name can still be used as a generic name if the type has type parameters. For example:

```

12     MODULE m
13         TYPE t(kind)
14             INTEGER, KIND :: kind
15             COMPLEX(kind) value
16         END TYPE
17         INTEGER, PARAMETER :: single = KIND(0.0), double = KIND(0d0)
18         INTERFACE t
19             MODULE PROCEDURE real_to_t1, dble_to_t2, int_to_t1, int_to_t2
20         END INTERFACE
21         ...
22     CONTAINS
23         TYPE(t(single)) FUNCTION real_to_t1(x)
24             REAL(single) x
25             real_to_t1%value = x
26         END FUNCTION
27         TYPE(t(double)) FUNCTION dble_to_t2(x)
28             REAL(double) x
29             dble_to_t2%value = x
30         END FUNCTION
31         TYPE(t(single)) FUNCTION int_to_t1(x,mold)
32             INTEGER x
33             TYPE(t(single)) mold
34             int_to_t1%value = x
35         END FUNCTION
36         TYPE(t(double)) FUNCTION int_to_t2(x,mold)
37             INTEGER x
38             TYPE(t(double)) mold
39             int_to_t2%value = x
40         END FUNCTION
41         ...
42     END
43
44     PROGRAM example
45         USE m

```

```

1      TYPE(t(single)) x
2      TYPE(t(double)) y
3      x = t(1.5)                ! References real_to_t1
4      x = t(17,mold=x)          ! References int_to_t1
5      y = t(1.5d0)              ! References dble_to_t2
6      y = t(42,mold=y)          ! References int_to_t2
7      y = t(kind(0d0)) ((0,1)) ! Uses the structure constructor for type t
8  END

```

### 9 C.1.6 Final subroutines (4.5.6, 4.5.6.2, 4.5.6.3, 4.5.6.4)

10 Example of a parameterized derived type with final subroutines:

```

11  MODULE m
12      TYPE t(k)
13          INTEGER, KIND :: k
14          REAL(k),POINTER :: vector(:) => NULL()
15      CONTAINS
16          FINAL :: finalize_t1s, finalize_t1v, finalize_t2e
17      END TYPE
18  CONTAINS
19      SUBROUTINE finalize_t1s(x)
20          TYPE(t(KIND(0.0))) x
21          IF (ASSOCIATED(x%vector)) DEALLOCATE(x%vector)
22      END SUBROUTINE
23      SUBROUTINE finalize_t1v(x)
24          TYPE(t(KIND(0.0))) x(:)
25          DO i=LBOUND(x,1),UBOUND(x,1)
26              IF (ASSOCIATED(x(i)%vector)) DEALLOCATE(x(i)%vector)
27          END DO
28      END SUBROUTINE
29      ELEMENTAL SUBROUTINE finalize_t2e(x)
30          TYPE(t(KIND(0.0d0))),INTENT(INOUT) :: x
31          IF (ASSOCIATED(x%vector)) DEALLOCATE(x%vector)
32      END SUBROUTINE
33  END MODULE
34
35  SUBROUTINE example(n)
36      USE m
37      TYPE(t(KIND(0.0))) a,b(10),c(n,2)
38      TYPE(t(KIND(0.0d0))) d(n,n)
39      ...
40      ! Returning from this subroutine will effectively do
41      !   CALL finalize_t1s(a)
42      !   CALL finalize_t1v(b)
43      !   CALL finalize_t2e(d)
44      ! No final subroutine will be called for variable C because the user

```

```

1          ! omitted to define a suitable specific procedure for it.
2      END SUBROUTINE

```

3 **Example of [extended types](#) with [final subroutines](#):**

```

4      MODULE m
5          TYPE t1
6              REAL a,b
7          END TYPE
8          TYPE,EXTENDS(t1) :: t2
9              REAL,POINTER :: c(:),d(:)
10         CONTAINS
11             FINAL :: t2f
12         END TYPE
13         TYPE,EXTENDS(t2) :: t3
14             REAL,POINTER :: e
15         CONTAINS
16             FINAL :: t3f
17         END TYPE
18         ...
19     CONTAINS
20         SUBROUTINE t2f(x) ! Finalizer for TYPE(t2)'s extra components
21             TYPE(t2) :: x
22             IF (ASSOCIATED(x%c)) DEALLOCATE(x%c)
23             IF (ASSOCIATED(x%d)) DEALLOCATE(x%d)
24         END SUBROUTINE
25         SUBROUTINE t3f(y) ! Finalizer for TYPE(t3)'s extra components
26             TYPE(t3) :: y
27             IF (ASSOCIATED(y%e)) DEALLOCATE(y%e)
28         END SUBROUTINE
29     END MODULE
30
31     SUBROUTINE example
32         USE m
33         TYPE(t1) x1
34         TYPE(t2) x2
35         TYPE(t3) x3
36         ...
37         ! Returning from this subroutine will effectively do
38         !     ! Nothing to x1; it is not finalizable
39         !     CALL t2f(x2)
40         !     CALL t3f(x3)
41         !     CALL t2f(x3%t2)
42     END SUBROUTINE

```

## C.2 Clause 5 notes: The VOLATILE attribute (5.5.19)

- 1 The following example shows the use of a variable with the [VOLATILE attribute](#) to communicate with an asynchronous process, in this case the operating system. The program detects a user keystroke on the terminal and reacts at a convenient point in its processing.
- 2 The [VOLATILE attribute](#) is necessary to prevent an optimizing compiler from storing the communication variable in a register or from doing flow analysis and deciding that the [EXIT statement](#) can never be executed.

```

SUBROUTINE TERMINATE_ITERATIONS
  LOGICAL, VOLATILE :: USER_HIT_ANY_KEY

  ! Have the OS start to look for a user keystroke and set the variable
  ! "USER_HIT_ANY_KEY" to TRUE as soon as it detects a keystroke.
  ! This call is operating system dependent.

  CALL OS_BEGIN_DETECT_USER_KEYSTROKE( USER_HIT_ANY_KEY )
  USER_HIT_ANY_KEY = .FALSE.      ! This will ignore any recent keystrokes.
  PRINT *, " Hit any key to terminate iterations!"

  DO I = 1,100
    ...                          ! Compute a value for R.
    PRINT *, I, R
    IF (USER_HIT_ANY_KEY)        EXIT
  ENDDO

  ! Have the OS stop looking for user keystrokes.
  CALL OS_STOP_DETECT_USER_KEYSTROKE

END SUBROUTINE TERMINATE_ITERATIONS

```

## C.3 Clause 6 notes

### C.3.1 Structure components (6.4.2)

- 1 Components of a structure are referenced by writing the components of successive levels of the structure hierarchy until the desired component is described. For example,

```

TYPE ID_NUMBERS
  INTEGER SSN
  INTEGER EMPLOYEE_NUMBER
END TYPE ID_NUMBERS

TYPE PERSON_ID
  CHARACTER (LEN=30) LAST_NAME
  CHARACTER (LEN=1) MIDDLE_INITIAL
  CHARACTER (LEN=30) FIRST_NAME

```

```

1      TYPE (ID_NUMBERS) NUMBER
2  END TYPE PERSON_ID
3
4  TYPE PERSON
5      INTEGER AGE
6      TYPE (PERSON_ID) ID
7  END TYPE PERSON
8
9  TYPE (PERSON) GEORGE, MARY
10
11  PRINT *, GEORGE % AGE           ! Print the AGE component
12  PRINT *, MARY % ID % LAST_NAME ! Print LAST_NAME of MARY
13  PRINT *, MARY % ID % NUMBER % SSN ! Print SSN of MARY
14  PRINT *, GEORGE % ID % NUMBER ! Print SSN and EMPLOYEE_NUMBER of GEORGE

```

- 15 2 A [structure component](#) can be a data object of intrinsic type as in the case of GEORGE % AGE or it can be  
 16 of derived type as in the case of GEORGE % ID % NUMBER. The resultant component can be a scalar or an  
 17 array of intrinsic or derived type.

```

18  TYPE LARGE
19      INTEGER ELT (10)
20      INTEGER VAL
21  END TYPE LARGE
22
23  TYPE (LARGE) A (5)           ! 5 element array, each of whose elements
24                               ! includes a 10 element array ELT and
25                               ! a scalar VAL.
26  PRINT *, A (1)               ! Prints 10 element array ELT and scalar VAL.
27  PRINT *, A (1) % ELT (3) ! Prints scalar element 3
28                               ! of array element 1 of A.
29  PRINT *, A (2:4) % VAL      ! Prints scalar VAL for array elements
30                               ! 2 to 4 of A.

```

- 31 3 Components of an object of [extensible type](#) that are [inherited](#) from the [parent type](#) can be accessed as a whole  
 32 by using the [parent component](#) name, or individually, either with or without qualifying them by the [parent](#)  
 33 [component](#) name. For example:

```

34  TYPE POINT           ! A base type
35      REAL :: X, Y
36  END TYPE POINT
37  TYPE, EXTENDS(POINT) :: COLOR_POINT ! An extension of TYPE(POINT)
38      ! Components X and Y, and component name POINT, inherited from parent
39      INTEGER :: COLOR
40  END TYPE COLOR_POINT
41
42  TYPE(POINT) :: PV = POINT(1.0, 2.0)
43  TYPE(COLOR_POINT) :: CPV = COLOR_POINT(POINT=PV, COLOR=3)

```

```

1
2      PRINT *, CPV%POINT           ! Prints 1.0 and 2.0
3      PRINT *, CPV%POINT%X, CPV%POINT%Y ! And this does, too
4      PRINT *, CPV%X, CPV%Y         ! And this does, too

```

### 5 C.3.2 Allocation with dynamic type (6.7.1)

6 1 The following example illustrates the use of allocation with the value and [dynamic type](#) of the allocated object  
7 given by another object. The example copies a list of objects of any type. It copies the list starting at IN\_LIST.  
8 After copying, each element of the list starting at LIST\_COPY has a polymorphic component, ITEM, for which  
9 both the value and type are taken from the ITEM component of the corresponding element of the list starting at  
10 IN\_LIST.

```

11      TYPE :: LIST ! A list of anything
12      TYPE(LIST), POINTER :: NEXT => NULL()
13      CLASS(*), ALLOCATABLE :: ITEM
14      END TYPE LIST
15      ...
16      TYPE(LIST), POINTER :: IN_LIST, LIST_COPY => NULL()
17      TYPE(LIST), POINTER :: IN_WALK, NEW_TAIL
18      ! Copy IN_LIST to LIST_COPY
19      IF (ASSOCIATED(IN_LIST)) THEN
20          IN_WALK => IN_LIST
21          ALLOCATE(LIST_COPY)
22          NEW_TAIL => LIST_COPY
23          DO
24              ALLOCATE(NEW_TAIL%ITEM, SOURCE=IN_WALK%ITEM)
25              IN_WALK => IN_WALK%NEXT
26              IF (.NOT. ASSOCIATED(IN_WALK)) EXIT
27              ALLOCATE(NEW_TAIL%NEXT)
28              NEW_TAIL => NEW_TAIL%NEXT
29          END DO
30      END IF

```

## 31 C.4 Clause 7 notes

### 32 C.4.1 Evaluation of function references (7.1.7)

33 1 If more than one function reference appears in a statement, they can be executed in any order (subject to a  
34 function result being evaluated after the evaluation of its arguments) and their values cannot depend on the order  
35 of execution. This lack of dependence on order of evaluation enables parallel execution of the function references.

### 36 C.4.2 Pointers in expressions (7.1.9.2)

37 1 A pointer is considered to be like any other variable when it is used as a primary in an expression. If a pointer  
38 is used as an operand to an operator that expects a value, the pointer will automatically deliver the value stored  
39 in the space described by the pointer, that is, the value of the target object associated with the pointer.

### 1 C.4.3 Pointers in variable definition contexts (7.2.1.3, 16.6.7)

- 1 The appearance of a pointer in a context that requires its value is a reference to its target. Similarly, where a  
 3 pointer appears in a variable definition context the variable that is defined is the target of the pointer.

- 2 Executing the program fragment

```

5      REAL, POINTER :: A
6      REAL, TARGET :: B = 10.0
7      A => B
8      A = 42.0
9      PRINT '(F4.1)', B

```

10 produces “42.0” as output.

## 11 C.5 Clause 8 notes

### 12 C.5.1 The SELECT CASE construct (8.1.9)

- 13 1 At most one case block is selected for execution within a [SELECT CASE construct](#), and there is no fall-through  
 14 from one block into another block within a [SELECT CASE construct](#). Thus there is no requirement for the user  
 15 to exit explicitly from a block.

### 16 C.5.2 Loop control (8.1.7)

- 17 1 Fortran provides several forms of loop control:

- 18 (1) With an iteration count and a DO variable. This is the classic Fortran DO loop.
- 19 (2) Test a logical condition before each execution of the loop (DO WHILE).
- 20 (3) DO “forever”.

### 21 C.5.3 Examples of DO constructs (8.1.7)

- 22 1 The following are all valid examples of [DO constructs](#).

#### 23 Example 1:

```

24      SUM = 0.0
25      READ (IUN) N
26      OUTER: DO L = 1, N                ! A DO with a construct name
27          READ (IUN) IQUAL, M, ARRAY (1:M)
28          IF (IQUAL < IQUAL_MIN) CYCLE OUTER ! Skip inner loop
29          INNER: DO 40 I = 1, M          ! A DO with a label and a name
30              CALL CALCULATE (ARRAY (I), RESULT)
31              IF (RESULT < 0.0) CYCLE
32              SUM = SUM + RESULT
33              IF (SUM > SUM_MAX) EXIT OUTER
34      40      END DO INNER
35      END DO OUTER

```

2 The outer loop has an iteration count of **MAX** (N, 0), and will execute that number of times or until SUM exceeds SUM\_MAX, in which case the **EXIT** OUTER statement terminates both loops. The inner loop is skipped by the first **CYCLE statement** if the quality flag, IQUAL, is too low. If CALCULATE returns a negative RESULT, the second **CYCLE statement** prevents it from being summed. Both loops have construct names and the inner loop also has a label. A construct name is required in the **EXIT statement** in order to terminate both loops, but is optional in the **CYCLE statements** because each belongs to its innermost loop.

#### 7 Example 2:

```

8         N = 0
9         DO 50, I = 1, 10
10            J = I
11            DO K = 1, 5
12                L = K
13                N = N + 1 ! This statement executes 50 times
14            END DO        ! Nonlabeled DO inside a labeled DO
15    50 CONTINUE

```

3 After execution of the above program fragment, I = 11, J = 10, K = 6, L = 5, and N = 50.

#### 17 Example 3:

```

18         N = 0
19         DO I = 1, 10
20             J = I
21             DO 60, K = 5, 1 ! This inner loop is never executed
22                 L = K
23                 N = N + 1
24         60 CONTINUE        ! Labeled DO inside a nonlabeled DO
25     END DO

```

4 After execution of the above program fragment, I = 11, J = 10, K = 5, N = 0, and L is not defined by these statements.

### 28 C.5.4 Examples of invalid DO constructs (8.1.7)

1 The following are all examples of invalid skeleton **DO constructs**:

#### 30 Example 1:

```

31         DO I = 1, 10
32             ...
33     END DO LOOP ! No matching construct name

```

#### 34 Example 2:

```

35         LOOP: DO 1000 I = 1, 10 ! No matching construct name
36             ...
37     1000 CONTINUE

```

1     **Example 3:**

```

2           LOOP1: DO
3           ...
4           END DO LOOP2    ! Construct names don't match

```

5     **Example 4:**

```

6           DO I = 1, 10    ! Label required or ...
7           ...
8           1010 CONTINUE  ! ... END DO required

```

9     **Example 5:**

```

10          DO 1020 I = 1, 10
11          ...
12          1021 END DO      ! Labels don't match

```

13    **Example 6:**

```

14          FIRST: DO I = 1, 10
15              SECOND: DO J = 1, 5
16              ...
17              END DO FIRST    ! Improperly nested DOs
18          END DO SECOND

```

19    **C.5.5 Simple example using events**

- 20    1 A tree is a graph in which every node except one has a single “parent” node to which it is connected by an edge.  
 21    The node without a parent is the “root” of the tree. The nodes that have a particular node as their parent are  
 22    the “children” of that node. The root is at level 1, its children are at level 2, and so on.
- 23    2 A multifrontal code to solve a sparse set of linear equations involves a tree. Work at a node can start after all of  
 24    its children’s work is complete and their data have been passed to it.
- 25    3 Here we assume that each node has been assigned to an image. Each image has a list of its nodes and these  
 26    are ordered in decreasing tree level (all those at level  $L$  preceding those at level  $L - 1$ ). For each node, array  
 27    elements hold the number of children, details about the parent, and an event variable. This allows the processing  
 28    to proceed asynchronously subject to the rule that a parent must wait for all its children.

29    **Outline of example code:**

```

30          PROGRAM TREE
31              USE, INTRINSIC :: ISO_FORTRAN_ENV
32              INTEGER, ALLOCATABLE :: NODE (:) ! Tree nodes that this image handles.
33              INTEGER, ALLOCATABLE :: NC (:)    ! NODE(I) has NC(I) children.
34              INTEGER, ALLOCATABLE :: PARENT (:), SUB (:)
35              ! The parent of NODE (I) is NODE (SUB (I)) [PARENT (I)].
36              TYPE (EVENT_TYPE), ALLOCATABLE :: DONE (:) [*]

```

```

1      INTEGER :: I, J, STATUS
2      ! Set up the tree, including allocation of all arrays.
3      DO I = 1, SIZE (NODE)
4          ! Wait for children to complete
5          IF (NC (I) > 0) THEN
6              EVENT WAIT (DONE (I), UNTIL_COUNT=NC (I), STAT=STATUS)
7              IF (STATUS/=0) EXIT
8          END IF
9
10         ! Process node, using data from children.
11         IF (PARENT (I)>0) THEN
12             ! Node is not the root.
13             ! Place result on image PARENT (I) for node NODE (SUB) [PARENT (I)]
14             ! Tell PARENT (I) that this has been done.
15             EVENT POST (DONE (SUB (I)) [PARENT (I)], STAT=STATUS)
16             IF (STATUS/=0) EXIT
17         END IF
18     END DO
19 END PROGRAM TREE

```

## C.5.6 Example using three teams

- 1 The following example illustrates the structure of a routine that will compute fluxes based on surface properties over land, sea, and ice, each in a different team. Each image will deal with areas containing exactly one of the three surface types.

```

24      SUBROUTINE COMPUTE_FLUXES (FLUX_MOM, FLUX_SENS, FLUX_LAT)
25      USE, INTRINSIC :: ISO_FORTRAN_ENV, ONLY: TEAM_TYPE
26      REAL, INTENT (OUT) :: FLUX_MOM (:,:), FLUX_SENS (:,:), FLUX_LAT (:,:)
27      INTEGER, PARAMETER :: LAND = 1, SEA = 2, ICE = 3
28      CHARACTER (LEN=10) :: SURFACE_TYPE
29      INTEGER :: MY_SURFACE_TYPE, N_IMAGE
30      TYPE (TEAM_TYPE) :: TEAM_SURFACE_TYPE
31
32      CALL GET_SURFACE_TYPE(THIS_IMAGE (), SURFACE_TYPE)
33      SELECT CASE (SURFACE_TYPE)
34      CASE ("LAND")
35          MY_SURFACE_TYPE = LAND
36      CASE ("SEA")
37          MY_SURFACE_TYPE = SEA
38      CASE ("ICE")
39          MY_SURFACE_TYPE = ICE
40      CASE DEFAULT
41          ERROR STOP
42      END SELECT
43      FORM TEAM (MY_SURFACE_TYPE, TEAM_SURFACE_TYPE)
44

```

```

1      CHANGE TEAM (TEAM_SURFACE_TYPE)
2      SELECT CASE (TEAM_NUMBER ( ))
3      CASE (LAND)      ! Compute fluxes over land surface
4          CALL COMPUTE_FLUXES_LAND (FLUX_MOM, FLUX_SENS, FLUX_LAT)
5      CASE (SEA)      ! Compute fluxes over sea surface
6          CALL COMPUTE_FLUXES_SEA (FLUX_MOM, FLUX_SENS, FLUX_LAT)
7      CASE (ICE)      ! Compute fluxes over ice surface
8          CALL COMPUTE_FLUXES_ICE (FLUX_MOM, FLUX_SENS, FLUX_LAT)
9      CASE DEFAULT
10         ERROR STOP
11     END SELECT
12 END TEAM
13 END SUBROUTINE COMPUTE_FLUXES

```

### C.5.7 Accessing coarrays in sibling teams

- 1 The following program illustrates subdividing a  $4 \times 4$  grid into  $2 \times 2$  teams, and the denotation of sibling teams.

```

16 PROGRAM DEMO
17     ! Initial team : 16 images. Algorithm design is a 4 by 4 grid.
18     ! Desire 4 teams, for the upper left (UL), upper right (UR),
19     !                               lower left (LL), lower right (LR)
20     USE, INTRINSIC :: ISO_FORTRAN_ENV, ONLY: TEAM_TYPE
21     TYPE (TEAM_TYPE) :: t
22     INTEGER, PARAMETER :: UL=11, UR=22, LL=33, LR=44
23     REAL :: A(10,10)[4,*]
24     INTEGER :: MYPE, TEAMNUM, NEWPE
25     INTEGER :: UL_IMAGE_LIST(4) = [1, 2, 5, 6], &
26             LL_IMAGE_LIST(4) = UL_IMAGE_LIST + 2, &
27             UR_IMAGE_LIST(4) = UL_IMAGE_LIST + 8, &
28             LR_IMAGE_LIST(4) = UL_IMAGE_LIST + 10
29
30     MYPE = THIS_IMAGE()
31     IF (ANY (MYPE==UL_IMAGE_LIST)) TEAMNUM = UL
32     IF (ANY (MYPE==LL_IMAGE_LIST)) TEAMNUM = LL
33     IF (ANY (MYPE==UR_IMAGE_LIST)) TEAMNUM = UR
34     IF (ANY (MYPE==LR_IMAGE_LIST)) TEAMNUM = LR
35     FORM TEAM (TEAMNUM, T)
36
37     A = 3.14
38
39     CHANGE TEAM (T, B[2,*] => A)
40     ! Inside change team, image pattern for B is a 2 by 2 grid.
41     B (5, 5) = B (1, 1)[2, 1]
42
43     ! Outside the team addressing:
44

```

```

1      NEWPE = THIS_IMAGE()
2      SELECT CASE (TEAM_NUMBER ())
3      CASE (UL)
4          IF (NEWPE==3) THEN
5              ! Right column of UL gets left column of UR.
6              B (:, 10) = B (:, 1)[1, 1, TEAM_NUMBER=UR]
7          ELSE IF (NEWPE==4) THEN
8              B (:, 10) = B (:, 1)[2, 1, TEAM_NUMBER=UR]
9          END IF
10     CASE (LL)
11         ! Similar to complete column exchange across middle of the original grid.
12         ...
13     END SELECT
14 END TEAM
15 END PROGRAM DEMO

```

### C.5.8 Example involving failed images

- 1 Parallel algorithms often use work sharing schemes based on a specific mapping between image indices and global data addressing. To allow such programs to continue when one or more images fail, spare images can be used to re-establish execution of the algorithm with the failed images replaced by spare images, while retaining the previous image mapping for nonfailed images.
- 2 The following example illustrates how this might be done. In this example, failure cannot be tolerated for image 1 in the initial team.

```

23 PROGRAM possibly_recoverable_simulation
24     USE, INTRINSIC :: ISO_FORTRAN_ENV, ONLY: TEAM_TYPE, STAT_FAILED_IMAGE
25     IMPLICIT NONE
26     INTEGER, ALLOCATABLE :: failures (:)
27     INTEGER :: images_used, i, images_spare, status
28     INTEGER :: id [*], me [*]
29     TYPE (TEAM_TYPE) :: simulation_team
30     LOGICAL :: read_checkpoint, done [*]
31
32     ! Keep 1% spare images if we have a lot, just 1 if 10-199 images, 0 if <10.
33     images_spare = MAX (INT (0.01*NUM_IMAGES ()), 0, MIN (NUM_IMAGES () - 10, 1)
34     images_used = NUM_IMAGES () - images_spare
35     read_checkpoint = THIS_IMAGE () > images_used
36
37     setup : DO
38         me = THIS_IMAGE ()
39         id = MERGE (1, 2, me<=images_used)
40         !
41         ! Set up spare images as replacement for failed ones.
42         !
43         IF (IMAGE_STATUS (1) == STAT_FAILED_IMAGE) &

```

```

1      ERROR STOP "cannot recover"
2  IF (me == 1) THEN
3      failures = FAILED_IMAGES ()
4      k = images_used
5      DO i = 1, SIZE (failures)
6          DO k = k+1, NUM_IMAGES ()
7              IF (IMAGE_STATUS (k) == 0) EXIT
8          END DO
9          IF (k > NUM_IMAGES ()) ERROR STOP "cannot recover"
10         me [k] = failures (i)
11         id [k] = 1
12     END DO
13     images_used = k
14 END IF
15 !
16 ! Set up a simulation team of constant size.
17 ! Team 2 is the set of spares, so does not participate in the simulation.
18 !
19 FORM TEAM (id, simulation_team, NEW_INDEX=me, STAT=status)
20 simulation : CHANGE TEAM (simulation_team, STAT=status)
21 IF (status == STAT_FAILED_IMAGE) EXIT simulation
22 IF (TEAM_NUMBER () == 1) THEN
23     iter : DO
24         CALL simulation_procedure (read_checkpoint, status, done)
25         ! The simulation_procedure:
26         ! - sets up and performs some part of the simulation;
27         ! - resets to the last checkpoint if requested;
28         ! - sets status from its internal synchronizations;
29         ! - sets done to .TRUE. when the simulation has completed.
30     IF (status == STAT_FAILED_IMAGE) THEN
31         read_checkpoint = .TRUE.
32         EXIT simulation
33     ELSE IF (done)
34         EXIT iter
35     END IF
36     read_checkpoint = .FALSE.
37 END DO iter
38 END IF
39 END TEAM simulation (STAT=status)
40 SYNC ALL (STAT=status)
41 IF (THIS_IMAGE () > images_used) done = done[1]
42 IF (done) EXIT setup
43 END DO setup
44 END PROGRAM possibly_recoverable_simulation

```

3 Supporting fault-tolerant execution imposes obligations on library writers who use the parallel language facilities.

Every synchronization statement, allocation or deallocation of coarrays, or invocation of a collective procedure must be prepared to handle error conditions, and implicit deallocation of coarrays must be avoided. Also, coarray module variables that are allocated inside the team execution context are not persistent.

### C.5.9 EVENT\_QUERY example that tolerates image failure

- 1 This example is an adaptation of the earlier EVENT\_QUERY example of A.2.1 to make it able to execute in the presence of the failure of one or more of the worker images. The function create\_work\_item now accepts an integer argument to indicate which work item is required. It is assumed that the work items are indexed 1, 2, ... . It is also assumed that if an image fails while processing a work item, that work item can subsequently be processed by another image.

```

PROGRAM work_share
  USE, INTRINSIC :: ISO_FORTRAN_ENV, ONLY: EVENT_TYPE
  USE :: mod_work, ONLY: & ! Module that creates work items
    work, & ! Type for holding a work item
    create_work_item, & ! Function that creates work item
    process_item, & ! Function that processes an item
    work_done ! Logical function that returns true
               ! if all work done

  TYPE :: worker_type
    TYPE (/linktypeEVENT_TYPE), ALLOCATABLE :: free (:)
  END TYPE

  TYPE (EVENT_TYPE) :: submit [*] ! Whether work ready for a worker
  TYPE (worker_type) :: worker [*] ! Whether worker is free
  TYPE (work) :: work_item [*] ! Holds the data for a work item
  INTEGER :: count, i, k, kk, nbusy [*, np, status]
  INTEGER, ALLOCATABLE :: working (:) ! Items being worked on
  INTEGER, ALLOCATABLE :: pending (:) ! Items pending after image failure

  IF (THIS_IMAGE () == 1) THEN
    ! Get started
    ALLOCATE (worker%free (2:NUM_IMAGES ()))
    ALLOCATE (working (2: NUM_IMAGES ()), pending(NUM_IMAGES ()-1))
    nbusy = 0 ! This holds the number of workers working
    k = 1 ! Index of next work item
    np = 0 ! Number of work items in array pending
    DO i = 2, NUM_IMAGES () ! Start the workers working
      IF (work_done ()) EXIT
      working (i) = 0
      IF (IMAGE_STATUS (i) == STAT_FAILED_IMAGE) CYCLE
      work_item [i] = create_work_item (k)
      working (i) = k
      k = k + 1
      nbusy = nbusy + 1
      EVENT POST (submit [i], STAT=status)
    
```

```

1      END DO
2      ! Main work distribution loop
3      master : DO
4          image : DO i = 2, NUM_IMAGES ()
5              IF (IMAGE_STATUS (i) == STAT_FAILED_IMAGE) THEN
6                  IF (working (i)>0) THEN                ! It failed while working
7                      np = np + 1
8                      pending (np) = working (i)
9                      working (i) = 0
10                     END IF
11                     CYCLE image
12                 END IF
13                 CALL EVENT_QUERY (worker%free (i), count)
14                 IF (count == 0) CYCLE image                ! Worker is not free
15                 EVENT WAIT (worker%free (i))
16                 nbusy = nbusy - 1
17                 IF (np>0) THEN
18                     kk = pending (np)
19                     np = np - 1
20                 ELSE
21                     IF (work_done ()) CYCLE image
22                     kk = k
23                     k = k + 1
24                 END IF
25                 nbusy = nbusy + 1
26                 working (i) = kk
27                 work_item [i] = create_work_item (kk)
28                 EVENT POST (submit [i], STAT=status)
29                 ! If image i has failed, the failure will be handled on
30                 ! the next iteration of the master loop.
31             END DO image
32             IF ( nbusy==0 ) THEN ! All done. Exit on all images.
33                 DO i = 2, NUM_IMAGES ()
34                     EVENT POST (submit [i], STAT=status)
35                     IF (status == STAT_FAILED_IMAGE) CYCLE
36                 END DO
37                 EXIT master
38             END IF
39         END DO master
40     ELSE
41         ! Work processing loop
42         worker : DO
43             EVENT WAIT (submit)
44             IF (nbusy [1] == 0) EXIT worker
45             CALL process_item(work_item)
46             EVENT POST (worker[1]%free (THIS_IMAGE ()))

```

```

1          END DO worker
2          END IF
3      END PROGRAM work_share

```

## C.6 Clause 9 notes

### C.6.1 External files (9.3)

- 1 This part of ISO/IEC 1539 accommodates, but does not require, file cataloging. To do this, several concepts are introduced.

#### C.6.1.1 File existence (9.3.2)

- 1 Totally independent of the connection state is the property of existence, this being a file property. The processor “knows” of a set of files that exist at a given time for a given program. This set would include tapes ready to read, files in a catalog, a keyboard, a printer, etc. The set might exclude files inaccessible to the program because of security, because they are already in use by another program, etc. This part of ISO/IEC 1539 does not specify which files exist, hence wide latitude is available to a processor to implement security, locks, privilege techniques, etc. Existence is a convenient concept to designate all of the files that a program can potentially process.

- 2 All four combinations of connection and existence can occur:

| Connect | Exist | Examples                                             |
|---------|-------|------------------------------------------------------|
| Yes     | Yes   | A card reader loaded and ready to be read            |
| Yes     | No    | A printer before the first line is written           |
| No      | Yes   | A file named 'JOAN' in the catalog                   |
| No      | No    | A file on a reel of tape, not known to the processor |

- 3 Means are provided to create, delete, connect, and disconnect files.

#### C.6.1.2 File access (9.3.3)

- 1 This part of ISO/IEC 1539 does not address problems of security, protection, locking, and many other concepts that might be part of the concept of “right of access”. Such concepts are considered to be in the province of an operating system.

- 2 The [OPEN](#) and [INQUIRE](#) statements can be extended naturally to consider these things.

- 3 Possible access methods for a file are: sequential, stream and direct. The processor might implement three different types of files, each with its own access method. It might instead implement one type of file with three different access methods.

- 4 Direct access to files is of a simple and commonly available type, that is, fixed-length records. The key is a positive integer.

### 1 C.6.1.3 File connection (9.5)

2 1 Before any input/output can be performed on a file, it needs to be connected to a [unit](#). The [unit](#) then serves as a  
3 designator for that file as long as it is connected. To be connected does not imply that “buffers” have or have not  
4 been allocated, that “file-control tables” have or have not been filled, or that any other method of implementation  
5 has been used. Connection means that (barring some other fault) a [READ](#) or [WRITE](#) statement can be executed  
6 on the [unit](#), hence on the file. Without a connection, a [READ](#) or [WRITE](#) statement cannot be executed.

### 7 C.6.1.4 File names (9.5.6.10)

8 1 A file can have a name. The form of a file name is not specified. If a system does not have some form of cataloging  
9 or tape labeling for at least some of its files, all file names disappear at the [termination of execution](#). This is  
10 a valid implementation. Nowhere does this part of ISO/IEC 1539 require names to survive for any period of  
11 time longer than the execution time span of a program. Therefore, this part of ISO/IEC 1539 does not impose  
12 cataloging as a prerequisite. The naming feature is intended to enable use of a cataloging system where one  
13 exists.

## 14 C.6.2 Nonadvancing input/output (9.3.4.2)

15 1 [Data transfer statements](#) affect the positioning of an [external file](#). In FORTRAN 77, if no error or end-of-file  
16 condition exists, the file is positioned after the record just read or written and that record becomes the preceding  
17 record. This part of ISO/IEC 1539 contains the [ADVANCE= specifier](#) in a [data transfer statement](#) that provides  
18 the capability of maintaining a position within the current record from one formatted [data transfer statement](#) to  
19 the next [data transfer statement](#). The value NO provides this capability. The value YES positions the file after  
20 the record just read or written. The default is YES.

21 2 The tab edit descriptor and the slash are still appropriate for use with this type of record access but the tab  
22 cannot reposition before the left tab limit.

23 3 A [BACKSPACE](#) of a file that is positioned within a record causes the specified unit to be positioned before the  
24 current record.

25 4 If the next input/output operation on a file after a nonadvancing write is a [rewind](#), [backspace](#), [end file](#) or [close](#)  
26 operation, the file is positioned implicitly after the current record before an ENDFILE record is written to the  
27 file, that is, a [REWIND](#), [BACKSPACE](#), or [ENDFILE](#) statement following a nonadvancing [WRITE statement](#)  
28 causes the file to be positioned at the end of the current output record before the endfile record is written to the  
29 file.

30 5 This part of ISO/IEC 1539 provides a [SIZE= specifier](#) to be used with formatted [data transfer statements](#). The  
31 variable in the [SIZE= specifier](#) is assigned the count of the number of characters that make up the sequence of  
32 values read by the data edit descriptors in the [input statement](#).

33 6 The count is especially helpful if there is only one list item in the input list because it is the number of characters  
34 that appeared for the item.

35 7 The [EOR= specifier](#) is provided to indicate when an EOR condition is encountered during a nonadvancing [data](#)  
36 [transfer statement](#). The EOR condition is not an error condition. If this specifier appears, an input list item that  
37 requires more characters than the record contained is padded with blanks if [PAD= 'YES'](#) is in effect. This means  
38 that the input list item completed successfully. The file is positioned after the current record. If the [IOSTAT=](#)  
39 [specifier](#) appears, the specified variable is defined with the value of the [named constant IOSTAT\\_EOR](#) from

the intrinsic module `ISO_FORTRAN_ENV` and the `data transfer statement` is terminated. Program execution continues with the statement specified in the `EOR= specifier`. The `EOR= specifier` gives the capability of taking control of execution when the EOR condition is encountered. The *do-variables* in *io-implied-dos* retain their last defined value and any remaining items in the *input-item-list* retain their definition status when an EOR condition occurs. If the `SIZE= specifier` appears, the specified variable is assigned the number of characters read with the data edit descriptors during the `READ statement`.

For nonadvancing input, the processor is not required to read partial records. The processor could read the entire record into an internal buffer and make successive portions of the record available to successive `input statements`.

In an implementation of nonadvancing input/output in which a nonadvancing write to a terminal device causes immediate display of the output, such a write can be used as a mechanism to output a prompt. In this case, the statement

```
WRITE (*, FMT='(A)', ADVANCE='NO') 'CONTINUE?(Y/N): '
```

would result in the prompt

```
CONTINUE?(Y/N):
```

being displayed with no subsequent line feed.

The response, which might be read by a statement of the form

```
READ (*, FMT='(A)') ANSWER
```

can then be entered on the same line as the prompt as in

```
CONTINUE?(Y/N): Y
```

This part of ISO/IEC 1539 does not require that an implementation of nonadvancing input/output operate in this manner. For example, an implementation of nonadvancing output in which the display of the output is deferred until the current record is complete is also standard-conforming. Such an implementation will not, however, allow a prompting mechanism of this kind to operate.

### C.6.3 OPEN statement (9.5.6)

1 A file can become connected to a `unit` either by preconnection or by execution of an `OPEN statement`. Preconnection is performed prior to the beginning of execution of a program by means external to Fortran. For example, it could be done by job control action or by processor-established defaults. Execution of an `OPEN statement` is not required in order to access preconnected files (9.5.5).

2 The `OPEN statement` provides a means to access existing files that are not preconnected. An `OPEN statement` can be used in either of two ways: with a file name (open-by-name) and without a file name (open-by-unit). A `unit` is given in either case. Open-by-name connects the specified file to the specified `unit`. Open-by-unit connects a processor-dependent default file to the specified `unit`. (The default file might or might not have a name.)

3 Therefore, there are three ways a file can become connected and hence processed: preconnection, open-by-name, and open-by-unit. Once a file is connected, there is no means in standard Fortran to determine how it became connected.

4 An `OPEN statement` can also be used to create a new file. In fact, any of the foregoing three connection methods can be performed on a file that does not exist. When a `unit` is preconnected, writing the first record creates the file. With the other two methods, execution of the `OPEN statement` creates the file.

- 5 When an **OPEN statement** is executed, the **unit** specified in the **OPEN statement** might or might not already be connected to a file. If it is already connected to a file (either through preconnection or by prior execution of an **OPEN statement**), then omitting the **FILE= specifier** in the **OPEN statement** implies that the file is to remain connected to the **unit**. Such an **OPEN statement** can be used to change the values of the blank interpretation mode, decimal edit mode, pad mode, input/output rounding mode, delimiter mode, and sign mode.
- 6 If the value of the **ACTION= specifier** is WRITE, then a **READ statement** cannot refer to the connection. **ACTION = 'WRITE'** does not restrict positioning by a **BACKSPACE statement** or positioning specified by the **POSITION= specifier** with the value APPEND. However, a **BACKSPACE statement** or an **OPEN statement** containing **POSITION = 'APPEND'** might fail if the processor needs to read the file to achieve the positioning.
- 7 The following examples illustrate these rules. In the first example, **unit 10** is preconnected to a SCRATCH file; the **OPEN statement** changes the value of **PAD=** to YES.

```

CHARACTER (LEN = 20) CH1
WRITE (10, '(A)') 'THIS IS RECORD 1'
OPEN (UNIT = 10, STATUS = 'OLD', PAD = 'YES')
REWIND 10
READ (10, '(A20)') CH1    ! CH1 now has the value
                           ! 'THIS IS RECORD 1    '

```

- 8 In the next example, **unit 12** is first connected to a file named FRED, with a status of OLD. The second **OPEN statement** then opens **unit 12** again, retaining the connection to the file FRED, but changing the value of the **DELIM= specifier** to QUOTE.

```

CHARACTER (LEN = 25) CH2, CH3
OPEN (12, FILE = 'FRED', STATUS = 'OLD', DELIM = 'NONE')
CH2 = '"THIS STRING HAS QUOTES."'
      ! Quotes in string CH2
WRITE (12, *) CH2          ! Written with no delimiters
OPEN (12, DELIM = 'QUOTE') ! Now quote is the delimiter
REWIND 12
READ (12, *) CH3    ! CH3 now has the value
                   ! 'THIS STRING HAS QUOTES.    '

```

- 9 The next example is invalid because it attempts to change the value of the **STATUS= specifier**.

```

OPEN (10, FILE = 'FRED', STATUS = 'OLD')
WRITE (10, *) A, B, C
OPEN (10, STATUS = 'SCRATCH')    ! Attempts to make FRED a SCRATCH file

```

- 10 The previous example could be made valid by closing the **unit** first, as in the next example.

```

OPEN (10, FILE = 'FRED', STATUS = 'OLD')
WRITE (10, *) A, B, C
CLOSE (10)
OPEN (10, STATUS = 'SCRATCH')    ! Opens a different SCRATCH file

```

#### C.6.4 Connection properties (9.5.4)

1 When a [unit](#) becomes connected to a file, either by execution of an [OPEN statement](#) or by preconnection, the following connection properties, among others, are established.

- (1) An access method, which is sequential, direct, or stream, is established for the connection ([9.5.6.3](#)).
- (2) A form, which is formatted or unformatted, is established for a connection to a file that exists or is created by the connection. For a connection that results from execution of an [OPEN statement](#), a default form (which depends on the access method, as described in [9.3.3](#)) is established if no form is specified. For a preconnected file that exists, a form is established by preconnection. For a preconnected file that does not exist, a form might be established, or the establishment of a form might be delayed until the file is created (for example, by execution of a formatted or unformatted [WRITE statement](#)) ([9.5.6.11](#)).
- (3) A record length might be established. If the access method is direct, the connection establishes a record length that specifies the length of each record of the file. A direct access file can only contain records that are all of equal length.
- (4) A sequential file can contain records of varying lengths. In this case, the record length established specifies the maximum length of a record in the file ([9.5.6.15](#)).

2 A processor has wide latitude in adapting these concepts and actions to its own cataloging and job control conventions. Some processors might need job control action to specify the set of files that exist or that will be created by a program. Some processors might not need any job control action prior to execution. This part of ISO/IEC 1539 enables processors to perform dynamic open, close, or file creation operations, but it does not require such capabilities of the processor.

3 The meaning of “open” in contexts other than Fortran might include such things as mounting a tape, console messages, spooling, label checking, security checking, etc. These actions might occur upon job control action external to Fortran, upon execution of an [OPEN statement](#), or upon execution of the first read or write of the file. The [OPEN statement](#) describes properties of the connection to the file and might or might not cause physical activities to take place.

#### C.6.5 Asynchronous input/output (9.6.2.5)

1 Rather than limit support for asynchronous input/output to what has been traditionally provided by facilities such as BUFFERIN/BUFFEROUT, this part of ISO/IEC 1539 builds upon existing Fortran syntax. This permits alternative approaches for implementing asynchronous input/output, and simplifies the task of adapting existing standard-conforming programs to use asynchronous input/output.

2 Not all processors actually perform input/output asynchronously, nor will every processor that does be able to handle [data transfer statements](#) with complicated input/output item lists in an asynchronous manner. Such processors can still be standard-conforming.

3 This part of ISO/IEC 1539 allows for at least two different conceptual models for asynchronous input/output.

4 Model 1: the processor performs asynchronous input/output when the item list is simple (perhaps one contiguous named array) and the input/output is unformatted. The implementation cost is reduced, and this is the scenario most likely to be beneficial on traditional “big-iron” machines.

5 Model 2: The processor is free to do any of the following:

- (1) on output, create a buffer inside the input/output library, completely formatted, and then start an asynchronous write of the buffer, and immediately return to the next statement in the program. The processor is free to wait for previously issued WRITES, or not, or
- (2) pass the input/output list addresses to another processor/process, which processes the list items independently of the processor that executes the user's code. The addresses of the list items must be computed before the asynchronous READ/WRITE statement completes. There is still an ordering requirement on list item processing to handle things like READ (...) N,(a(i),i=1,N).

6 A program can issue a large number of asynchronous input/output requests, without waiting for any of them to complete, and then wait for any or all of them. That does not constitute a requirement for the processor to keep track of each individual request separately.

7 It is not necessary for all requests to be tracked by the runtime library. If an ID= specifier does not appear in on a READ or WRITE statement, the runtime library can forget about this particular request once it has successfully completed. If an error or end-of-file condition occurs for a request, the processor can report this during any input/output operation to that unit. If an ID= specifier appears, the processor's runtime input/output library will need to keep track of any end-of-file or error conditions for that particular input/output request. However, if the input/output request succeeds without any exceptional conditions occurring, then the runtime can forget that ID= value. A runtime library might only keep track of the last request made, or perhaps a very few. Then, when a user WAITs for a particular request, either the library will know about it (and does the right thing with respect to error handling, etc.), or can assume it is a request that successfully completed and was forgotten about (and will just return without signaling any end-of-file or error condition). A standard-conforming program can only pass valid ID= values, but there is no requirement on the processor to detect invalid ID= values. There might be a processor dependent limit on how many outstanding input/output requests that generate an end-of-file or error condition can be handled before the processor runs out of memory to keep track of such conditions. The restrictions on the SIZE= variables are designed to enable the processor to update such variables at any time (after the request has been processed, but before the wait operation), and then forget about them. Only error and end-of-file conditions are expected to be tracked by individual request by the runtime, and then only if an ID= specifier appears. The END= and EOR= specifiers have not been added to all statements that can perform wait operations. Instead, the IOSTAT variable can be queried after a wait operation to handle this situation. This choice was made because the WAIT statement is expected to be the usual method of waiting for input/output to complete (and WAIT does support the END= and EOR= specifiers). This particular choice is philosophical, and was not based on significant technical difficulties.

8 The requirement to set the IOSTAT variable correctly means that a processor will need to remember which input/output requests encountered an end-of-record condition, so that a subsequent wait operation can return the correct IOSTAT value. Therefore there might be a processor defined limit on the number of outstanding nonadvancing input/output requests that have encountered an end-of-record condition (constrained by available memory to keep track of this information, similar to end-of-file and error conditions).

## C.7 Clause 10 notes

### C.7.1 Number of records (10.4, 10.5, 10.8.2)

1 The number of records read by an explicitly formatted advancing input statement can be determined from the following rule: a record is read at the beginning of the format scan (even if the input list is empty unless the most recently previous operation on the unit was not a nonadvancing read operation), at each slash edit descriptor

encountered in the format, and when a format rescan occurs at the end of the format.

The number of records written by an explicitly formatted advancing output statement can be determined from the following rule: a record is written when a slash edit descriptor is encountered in the format, when a format rescan occurs at the end of the format, and at completion of execution of an advancing [output statement](#) (even if the output list is empty). Thus, the occurrence of  $n$  successive slashes between two other edit descriptors causes  $n - 1$  blank lines if the records are printed. The occurrence of  $n$  slashes at the beginning or end of a complete format specification causes  $n$  blank lines if the records are printed. However, a complete format specification containing  $n$  slashes ( $n > 0$ ) and no other edit descriptors causes  $n + 1$  blank lines if the records are printed. For example, the statements

```
PRINT 3
3 FORMAT (/)
```

will write two records that cause two blank lines if the records are printed.

### C.7.2 List-directed input ([10.10.3](#))

The following examples illustrate list-directed input. A blank character is represented by b.

#### Example 1:

##### Program:

```
J = 3
READ *, I
READ *, J
```

##### Sequential input file:

```
record 1:  b1b,4bbbbbb
record 2:  ,2bbbbbbbbb
```

Result:  $I = 1$ ,  $J = 3$ .

Explanation: The second [READ statement](#) reads the second record. The initial comma in the record designates a null value; therefore,  $J$  is not redefined.

#### Example 2:

##### Program:

```
CHARACTER A *8, B *1
READ *, A, B
```

##### Sequential input file:

```
record 1:  'bbbbbbbbb'
record 2:  'QXY'b'Z'
```

1 6 Result: A = 'bbbbbbbb', B = 'Q'

2 7 Explanation: In the first record, the rightmost apostrophe is interpreted as delimiting the constant (it cannot  
3 be the first of a pair of embedded apostrophes representing a single apostrophe because this would involve  
4 the prohibited “splitting” of the pair by the end of a record); therefore, A is assigned the character constant  
5 'bbbbbbbb'. The end of a record acts as a blank, which in this case is a value separator because it occurs between  
6 two constants.

## 7 C.8 Clause 11 notes

### 8 C.8.1 Main program and block data program unit (11.1, 11.3)

9 1 The name of the main program or of a block data program unit has no explicit use within the Fortran language.  
10 It is available for documentation and for possible use by a processor.

11 2 A processor might implement an unnamed program unit by assigning it a global identifier that is not used  
12 elsewhere in the program. This could be done by using a default name that does not satisfy the rules for Fortran  
13 names.

### 14 C.8.2 Dependent compilation (11.2)

15 1 This part of ISO/IEC 1539, like its predecessors, is intended to enable the implementation of conforming processors  
16 in which a program can be broken into multiple units, each of which can be separately translated in preparation  
17 for execution. Such processors are commonly described as supporting separate compilation. There is an important  
18 difference between the way separate compilation can be implemented under this part of ISO/IEC 1539 and the  
19 way it could be implemented under the FORTRAN 77 International Standard. Under the FORTRAN 77 standard,  
20 any information required to translate a [program unit](#) was specified in that [program unit](#). Each translation was  
21 thus totally independent of all others. Under this part of ISO/IEC 1539, a [program unit](#) can use information  
22 that was specified in a separate module and thus can be dependent on that module. The implementation of this  
23 dependency in a processor might be that the translation of a [program unit](#) depends on the results of translating  
24 one or more modules. Processors implementing the dependency this way are commonly described as supporting  
25 dependent compilation.

26 2 The dependencies involved here are new only in the sense that the Fortran processor is now aware of them. The  
27 same information dependencies existed under the FORTRAN 77 International Standard, but it was the program-  
28 mer's responsibility to transport the information necessary to resolve them by making redundant specifications of  
29 the information in multiple [program units](#). The availability of separate but dependent compilation offers several  
30 potential advantages over the redundant textual specification of information.

31 (1) Specifying information at a single place in the program ensures that different [program units](#) using that  
32 information are translated consistently. Redundant specification leaves the possibility that different  
33 information can be erroneously be specified. Even if an INCLUDE line is used to ensure that the  
34 text of the specifications is identical in all involved [program units](#), the presence of other specifications  
35 (for example, an [IMPLICIT statement](#)) could change the interpretation of that text.

36 (2) During the revision of a program, it is possible for a processor to assist in determining whether differ-  
37 ent [program units](#) have been translated using different (incompatible) versions of a module, although  
38 there is no requirement that a processor provide such assistance. Inconsistencies in redundant textual  
39 specification of information, on the other hand, tend to be much more difficult to detect.

(3) Putting information in a module provides a way of packaging it. Without modules, redundant specifications frequently are interleaved with other specifications in a [program unit](#), making convenient packaging of such information difficult.

(4) Because a processor can be implemented such that the specifications in a module are translated once and then repeatedly referenced, there is the potential for greater efficiency than when the processor translates redundant specifications of information in multiple [program units](#).

3 The exact meaning of the requirement that the public portions of a module be available at the time of reference is processor dependent. For example, a processor could consider a module to be available only after it has been compiled and require that if the module has been compiled separately, the result of that compilation be identified to the compiler when compiling [program units](#) that use it.

#### C.8.2.1 USE statement and dependent compilation (11.2.2)

1 Another benefit of the [USE statement](#) is its enhanced facilities for name management. If one needs to use only selected entities in a module, one can do so without having to worry about the names of all the other entities in that module. If one needs to use two different modules that happen to contain entities with the same name, there are several ways to deal with the conflict. If none of the entities with the same name are to be used, they can simply be ignored. If the name happens to refer to the same entity in both modules (for example, if both modules obtained it from a third module), then there is no confusion about what the name denotes and the name can be freely used. If the entities are different and one or both is to be used, the local renaming facility in the [USE statement](#) makes it possible to give those entities different names in the [program unit](#) containing the [USE statements](#).

2 A benefit of using the [ONLY](#) option consistently, as compared to [USE](#) without it, is that the module from which each accessed entity is accessed is explicitly specified in each [program unit](#). This means that one need not search other [program units](#) to find where each one is defined. This reduces maintenance costs.

3 A typical implementation of dependent but separate compilation might involve storing the result of translating a module in a file whose name is derived from the name of the module. Note, however, that the name of a module is limited only by the Fortran rules and not by the names allowed in the file system. Thus the processor might have to provide a mapping between Fortran names and file system names.

4 The result of translating a module could reasonably either contain only the information textually specified in the module (with “pointers” to information originally textually specified in other modules) or contain all information specified in the module (including copies of information originally specified in other modules). Although the former approach would appear to save on storage space, the latter approach can greatly simplify the logic necessary to process a [USE statement](#) and can avoid the necessity of imposing a limit on the logical “nesting” of modules via the [USE statement](#).

5 There is an increased potential for undetected errors in a [scoping unit](#) that uses both implicit typing and the [USE statement](#). For example, in the program fragment

```
SUBROUTINE SUB
  USE MY_MODULE
  IMPLICIT INTEGER (I-N), REAL (A-H, O-Z)
  X = F (B)
  A = G (X) + H (X + 1)
END SUBROUTINE SUB
```

X could be either an implicitly typed real variable or a variable obtained from the module MY\_MODULE and might change from one to the other because of changes in MY\_MODULE unrelated to the action performed by SUB. Logic errors resulting from this kind of situation can be extremely difficult to locate. Thus, the use of these features together is discouraged.

### C.8.2.2 Accessibility attributes (5.5.2)

The **PUBLIC** and **PRIVATE** attributes, which can be declared only in modules, divide the entities in a module into those that are actually relevant to a **scoping unit** referencing the module and those that are not. This information might be used to improve the performance of a Fortran processor. For example, it might be possible to discard much of the information about the private entities once a module has been translated, thus saving on both storage and the time to search it. Similarly, it might be possible to recognize that two versions of a module differ only in the private entities they contain and avoid retranslating **program units** that use that module when switching from one version of the module to the other.

## C.8.3 Examples of the use of modules (11.2.1)

### C.8.3.1 Global data (11.2.1)

A module could contain only data objects, for example:

```
MODULE DATA_MODULE
  SAVE
  REAL A (10), B, C (20,20)
  INTEGER :: I=0
  INTEGER, PARAMETER :: J=10
  COMPLEX D (J,J)
END MODULE DATA_MODULE
```

Data objects made global in this manner can have any combination of data types.

Access to some of these can be made by a **USE statement** with the **ONLY** option, such as:

```
USE DATA_MODULE, ONLY: A, B, D
```

and access to all of them can be made by the following **USE statement**:

```
USE DATA_MODULE
```

Access to all of them with some renaming to avoid name conflicts can be made by, for example:

```
USE DATA_MODULE, AMODULE => A, DMODULE => D
```

### C.8.3.2 Derived types (11.2.1)

A derived type can be defined in a module and accessed in a number of **program units**. For example,

```
MODULE SPARSE
  TYPE NONZERO
    REAL A
    INTEGER I, J
```

```

1          END TYPE NONZERO
2      END MODULE SPARSE

```

3 defines a type consisting of a real component and two integer components for holding the numerical value of a  
 4 nonzero matrix element and its row and column indices.

### 5 **C.8.3.3 Global allocatable arrays (11.2.1)**

6 1 Many programs need large global allocatable arrays whose sizes are not known before program execution. A  
 7 simple form for such a program is:

```

8      PROGRAM GLOBAL_WORK
9          CALL CONFIGURE_ARRAYS      ! Perform the appropriate allocations
10         CALL COMPUTE               ! Use the arrays in computations
11     END PROGRAM GLOBAL_WORK
12     MODULE WORK_ARRAYS             ! An example set of work arrays
13         INTEGER N
14         REAL, ALLOCATABLE :: A (:), B (:, :), C (:, :, :)
15     END MODULE WORK_ARRAYS
16     SUBROUTINE CONFIGURE_ARRAYS    ! Process to set up work arrays
17         USE WORK_ARRAYS
18         READ (*, *) N
19         ALLOCATE (A (N), B (N, N), C (N, N, 2 * N))
20     END SUBROUTINE CONFIGURE_ARRAYS
21     SUBROUTINE COMPUTE
22         USE WORK_ARRAYS
23         ... ! Computations involving arrays A, B, and C
24     END SUBROUTINE COMPUTE

```

25 2 Typically, many subprograms need access to the work arrays, and all such subprograms would contain the  
 26 statement

```

27         USE WORK_ARRAYS

```

### 28 **C.8.3.4 Procedure libraries (11.2.2)**

29 1 [Interface bodies](#) for [external procedures](#) in a library can be gathered into a module. An [interface body](#) specifies  
 30 an [explicit interface](#) (12.4.2.2).

31 2 An example is the following library module:

```

32     MODULE LIBRARY_LLS
33         INTERFACE
34             SUBROUTINE LLS (X, A, F, FLAG)
35                 REAL X (:, :)
36                 ! The SIZE in the next statement is an intrinsic function
37                 REAL, DIMENSION (SIZE (X, 2)) :: A, F
38                 INTEGER FLAG
39             END SUBROUTINE LLS

```

```

1          ...
2      END INTERFACE
3          ...
4  END MODULE LIBRARY_LLS

```

3 This module provides an [explicit interface](#) that is necessary for the subroutine LLS to be invoked. for example:

```

6      USE LIBRARY_LLS
7          ...
8      CALL LLS (X = ABC, A = D, F = XX, FLAG = IFLAG)
9          ...

```

4 Because [dummy argument](#) names in an [interface body](#) for an [external procedure](#) are not required to be the same as in the procedure definition, different versions can be constructed for different applications using [argument keywords](#) appropriate to each application.

### C.8.3.5 Operator extensions (11.2.2)

1 In order to extend an intrinsic operator symbol to have an additional meaning, an [interface block](#) specifying that operator symbol in the [OPERATOR](#) option of the [INTERFACE statement](#) could be placed in a module.

2 For example, // can be extended to perform concatenation of two derived-type objects serving as varying length character strings and + can be extended to specify matrix addition for type MATRIX or interval arithmetic addition for type INTERVAL.

3 A module might contain several such [interface blocks](#). An operator can be defined by an external function (either in Fortran or some other language) and its procedure [interface](#) placed in the module.

### C.8.3.6 Data abstraction (11.2.2)

1 In addition to providing a portable means of avoiding the redundant specification of information in multiple [program units](#), a module provides a convenient means of “packaging” related entities, such as the definitions of the representation and operations of an abstract data type. The following example of a module defines a data abstraction for a SET type where the elements of each set are of type integer. The usual set operations of UNION, INTERSECTION, and DIFFERENCE are provided. The CARDINALITY function returns the cardinality of (number of elements in) its set argument. Two functions returning logical values are included, ELEMENT and SUBSET. ELEMENT defines the operator .IN. and SUBSET extends the operator <=. ELEMENT determines if a given scalar integer value is an element of a given set, and SUBSET determines if a given set is a subset of another given set. (Two sets can be checked for equality by comparing cardinality and checking that one is a subset of the other, or checking to see if each is a subset of the other.)

2 The transfer function SETF converts a vector of integer values to the corresponding set, with duplicate values removed. Thus, a vector of constant values can be used as set constants. An inverse transfer function VECTOR returns the elements of a set as a vector of values in ascending order. In this SET implementation, set data objects have a maximum cardinality of 200.

3 Here is the example module:

```

37      MODULE INTEGER_SETS

```

```

1      ! This module is intended to illustrate use of the module facility
2      ! to define a new type, along with suitable operators.
3
4      INTEGER, PARAMETER :: MAX_SET_CARD = 200
5
6      TYPE SET                                ! Define SET type
7          PRIVATE
8          INTEGER CARD
9          INTEGER ELEMENT (MAX_SET_CARD)
10     END TYPE SET
11
12     INTERFACE OPERATOR (.IN.)
13         MODULE PROCEDURE ELEMENT
14     END INTERFACE OPERATOR (.IN.)
15
16     INTERFACE OPERATOR (<=)
17         MODULE PROCEDURE SUBSET
18     END INTERFACE OPERATOR (<=)
19
20     INTERFACE OPERATOR (+)
21         MODULE PROCEDURE UNION
22     END INTERFACE OPERATOR (+)
23
24     INTERFACE OPERATOR (-)
25         MODULE PROCEDURE DIFFERENCE
26     END INTERFACE OPERATOR (-)
27
28     INTERFACE OPERATOR (*)
29         MODULE PROCEDURE INTERSECTION
30     END INTERFACE OPERATOR (*)
31
32     CONTAINS
33
34     INTEGER FUNCTION CARDINALITY (A)          ! Returns cardinality of set A
35         TYPE (SET), INTENT (IN) :: A
36         CARDINALITY = A % CARD
37     END FUNCTION CARDINALITY
38
39     LOGICAL FUNCTION ELEMENT (X, A)           ! Determines if
40         INTEGER, INTENT(IN) :: X             ! element X is in set A
41         TYPE (SET), INTENT(IN) :: A
42         ELEMENT = ANY (A % ELEMENT (1 : A % CARD) == X)
43     END FUNCTION ELEMENT
44
45     FUNCTION UNION (A, B)                    ! Union of sets A and B
46         TYPE (SET) UNION

```

```

1      TYPE (SET), INTENT(IN) :: A, B
2      INTEGER J
3      UNION = A
4      DO J = 1, B % CARD
5          IF (.NOT. (B % ELEMENT (J) .IN. A)) THEN
6              IF (UNION % CARD < MAX_SET_CARD) THEN
7                  UNION % CARD = UNION % CARD + 1
8                  UNION % ELEMENT (UNION % CARD) = B % ELEMENT (J)
9              ELSE
10                 ! Maximum set size exceeded . . .
11             END IF
12         END IF
13     END DO
14 END FUNCTION UNION

15
16 FUNCTION DIFFERENCE (A, B)           ! Difference of sets A and B
17     TYPE (SET) DIFFERENCE
18     TYPE (SET), INTENT(IN) :: A, B
19     INTEGER J, X
20     DIFFERENCE % CARD = 0             ! The empty set
21     DO J = 1, A % CARD
22         X = A % ELEMENT (J)
23         IF (.NOT. (X .IN. B)) DIFFERENCE = DIFFERENCE + SET (1, X)
24     END DO
25 END FUNCTION DIFFERENCE

26
27 FUNCTION INTERSECTION (A, B)         ! Intersection of sets A and B
28     TYPE (SET) INTERSECTION
29     TYPE (SET), INTENT(IN) :: A, B
30     INTERSECTION = A - (A - B)
31 END FUNCTION INTERSECTION

32
33 LOGICAL FUNCTION SUBSET (A, B)        ! Determines if set A is
34     TYPE (SET), INTENT(IN) :: A, B    ! a subset of set B
35     INTEGER I
36     SUBSET = A % CARD <= B % CARD
37     IF (.NOT. SUBSET) RETURN           ! For efficiency
38     DO I = 1, A % CARD
39         SUBSET = SUBSET .AND. (A % ELEMENT (I) .IN. B)
40     END DO
41 END FUNCTION SUBSET

42
43 TYPE (SET) FUNCTION SETF (V)          ! Transfer function between a vector
44     INTEGER V (:)                     ! of elements and a set of elements
45     INTEGER J                           ! removing duplicate elements
46     SETF % CARD = 0

```

```

1      DO J = 1, SIZE (V)
2          IF (.NOT. (V (J) .IN. SETF)) THEN
3              IF (SETF % CARD < MAX_SET_CARD) THEN
4                  SETF % CARD = SETF % CARD + 1
5                  SETF % ELEMENT (SETF % CARD) = V (J)
6              ELSE
7                  ! Maximum set size exceeded . . .
8              END IF
9          END IF
10     END DO
11 END FUNCTION SETF
12
13 FUNCTION VECTOR (A)                ! Transfer the values of set A
14     TYPE (SET), INTENT (IN) :: A    ! into a vector in ascending order
15     INTEGER, POINTER :: VECTOR (:)
16     INTEGER I, J, K
17     ALLOCATE (VECTOR (A % CARD))
18     VECTOR = A % ELEMENT (1 : A % CARD)
19     DO I = 1, A % CARD - 1          ! Use a better sort if
20         DO J = I + 1, A % CARD      ! A % CARD is large
21             IF (VECTOR (I) > VECTOR (J)) THEN
22                 K = VECTOR (J); VECTOR (J) = VECTOR (I); VECTOR (I) = K
23             END IF
24         END DO
25     END DO
26 END FUNCTION VECTOR
27
28 END MODULE INTEGER_SETS

```

29 4 Examples of using INTEGER\_SETS (A, B, and C are variables of type SET; X is an integer variable):

```

30     ! Check to see if A has more than 10 elements
31     IF (CARDINALITY (A) > 10) ...
32
33     ! Check for X an element of A but not of B
34     IF (X .IN. (A - B)) ...
35
36     ! C is the union of A and the result of B intersected
37     ! with the integers 1 to 100
38     C = A + B * SETF ([I, I = 1, 100])
39
40     ! Does A have any even numbers in the range 1:100?
41     IF (CARDINALITY (A * SETF ([I, I = 2, 100, 2])) > 0) ...
42
43     PRINT *, VECTOR (B) ! Print out the elements of set B, in ascending order

```

### 1 C.8.3.7 Public entities renamed (11.2.2)

2 1 At times it might be necessary to rename entities that are accessed with [USE statements](#).

3 2 The following example illustrates renaming features of the [USE statement](#).

```

4      MODULE J; REAL JX, JY, JZ; END MODULE J
5      MODULE K
6          USE J, ONLY : KX => JX, KY => JY
7          ! KX and KY are local names to module K
8          REAL KZ      ! KZ is local name to module K
9          REAL JZ      ! JZ is local name to module K
10     END MODULE K
11     PROGRAM RENAME
12         USE J; USE K
13         ! Module J's entity JX is accessible under names JX and KX
14         ! Module J's entity JY is accessible under names JY and KY
15         ! Module K's entity KZ is accessible under name KZ
16         ! Module J's entity JZ and K's entity JZ are different entities
17         ! and cannot be referenced
18         ...
19     END PROGRAM RENAME

```

### 20 C.8.4 Modules with submodules (11.2.3)

21 1 Each submodule specifies that it is the child of exactly one parent module or submodule. Therefore, a module  
22 and all of its descendant submodules stand in a tree-like relationship one to another.

23 2 A separate module procedure that is declared in a module to have public accessibility can be accessed by use  
24 association even if it is defined in a submodule. No other entity in a submodule can be accessed by use association.  
25 Each [program unit](#) that references a module by use association depends on it, and each submodule depends on  
26 its ancestor module. Therefore, if one changes a separate module procedure body in a submodule but does not  
27 change its corresponding module procedure interface, a tool for automatic program translation would not need  
28 to reprocess [program units](#) that reference the module by use association. This is so even if the tool exploits the  
29 relative modification times of files as opposed to comparing the result of translating the module to the result of  
30 a previous translation.

31 3 By constructing taller trees, one can put entities at intermediate levels that are shared by submodules at lower  
32 levels; changing these entities cannot change the interpretation of anything that is accessible from the module  
33 by use association. Developers of modules that embody large complicated concepts can exploit this possibility  
34 to organize components of the concept into submodules, while preserving the privacy of entities that are shared  
35 by the submodules and that ought not to be exposed to users of the module. Putting these shared entities at an  
36 intermediate level also prevents cascades of reprocessing and testing if some of them are changed.

37 4 The following example illustrates a module, `color_points`, with a submodule, `color_points_a`, that in turn has  
38 a submodule, `color_points_b`. Public entities declared within `color_points` can be accessed by use association.  
39 The submodules `color_points_a` and `color_points_b` can be changed without causing retranslation of [program](#)  
40 [units](#) that reference the module `color_points`.

5 The module `color_points` does not have a *module-subprogram-part*, but a *module-subprogram-part* is not prohibited. The module could be published as definitive specification of the interface, without revealing trade secrets contained within `color_points_a` or `color_points_b`. Of course, a similar module without the `module` prefix in the interface bodies would serve equally well as documentation – but the procedures would be *external procedures*. It would make little difference to the consumer, but the developer would forfeit all of the advantages of modules.

```

6      module color_points
7
8          type color_point
9              private
10                 real :: x, y
11                 integer :: color
12             end type color_point
13
14         interface                ! Interfaces for procedures with separate
15                                 ! bodies in the submodule color_points_a
16             module subroutine color_point_del ( p ) ! Destroy a color_point object
17                 type(color_point), allocatable :: p
18             end subroutine color_point_del
19             ! Distance between two color_point objects
20             real module function color_point_dist ( a, b )
21                 type(color_point), intent(in) :: a, b
22             end function color_point_dist
23             module subroutine color_point_draw ( p ) ! Draw a color_point object
24                 type(color_point), intent(in) :: p
25             end subroutine color_point_draw
26             module subroutine color_point_new ( p ) ! Create a color_point object
27                 type(color_point), allocatable :: p
28             end subroutine color_point_new
29         end interface
30
31     end module color_points

```

6 The only entities within `color_points_a` that can be accessed by use association are the separate module procedures that were declared in `color_points`. If the procedures are changed but their interfaces are not, the interface from *program units* that access them by use association is unchanged. If the module and submodule are in separate files, utilities that examine the time of modification of a file would notice that changes in the module could affect the translation of its submodules or of *program units* that reference the module by use association, but that changes in submodules could not affect the translation of the parent module or *program units* that reference it by use association.

7 The variable `instance_count` in the following example is not accessible by use association of `color_points`, but is accessible within `color_points_a`, and its submodules.

```

41     submodule ( color_points ) color_points_a ! Submodule of color_points
42
43         integer :: instance_count = 0

```

```

1
2      interface                ! Interface for a procedure with a separate
3                               ! body in submodule color_points_b
4      module subroutine inquire_palette ( pt, pal )
5          use palette_stuff      ! palette_stuff, especially submodules
6                               ! thereof, can reference color_points by use
7                               ! association without causing a circular
8                               ! dependence during translation because this
9                               ! use is not in the module. Furthermore,
10                              ! changes in the module palette_stuff do not
11                              ! affect the translation of color_points.
12          type(color_point), intent(in) :: pt
13          type(palette), intent(out) :: pal
14      end subroutine inquire_palette
15
16  end interface
17
18  contains ! Invisible bodies for public separate module procedures
19          ! declared in the module
20
21  module subroutine color_point_del ( p )
22      type(color_point), allocatable :: p
23      instance_count = instance_count - 1
24      deallocate ( p )
25  end subroutine color_point_del
26  real module function color_point_dist ( a, b ) result ( dist )
27      type(color_point), intent(in) :: a, b
28      dist = sqrt( (b%x - a%x)**2 + (b%y - a%y)**2 )
29  end function color_point_dist
30  module subroutine color_point_new ( p )
31      type(color_point), allocatable :: p
32      instance_count = instance_count + 1
33      allocate ( p )
34  end subroutine color_point_new
35
36  end submodule color_points_a

```

- 8 The subroutine `inquire_palette` is accessible within `color_points_a` because its interface is declared therein. It is not, however, accessible by use association, because its interface is not declared in the module, `color_points`. Since the interface is not declared in the module, changes in the interface cannot affect the translation of [program units](#) that reference the module by use association.

```

41  module palette_stuff
42      type :: palette ; ... ; end type palette
43  contains
44      subroutine test_palette ( p )
45          ! Draw a color wheel using procedures from the color_points module

```

```

1      use color_points ! This does not cause a circular dependency because
2                      ! the "use palette_stuff" that is logically within
3                      ! color_points is in the color_points_a submodule.
4      type(palette), intent(in) :: p
5      ...
6      end subroutine test_palette
7  end module palette_stuff
8
9  submodule ( color_points:color_points_a ) color_points_b ! Subsidiary**2 submodule
10
11  contains
12      ! Invisible body for interface declared in the ancestor module
13      module subroutine color_point_draw ( p )
14          use palette_stuff, only: palette
15          type(color_point), intent(in) :: p
16          type(palette) :: MyPalette
17          ...; call inquire_palette ( p, MyPalette ); ...
18      end subroutine color_point_draw
19
20      ! Invisible body for interface declared in the parent submodule
21      module procedure inquire_palette
22          ... implementation of inquire_palette
23      end procedure inquire_palette
24
25      subroutine private_stuff ! not accessible from color_points_a
26          ...
27      end subroutine private_stuff
28
29  end submodule color_points_b

```

30 9 There is a use palette\_stuff in color\_points\_a, and a use color\_points in palette\_stuff. The use  
31 palette\_stuff would cause a circular reference if it appeared in color\_points. In this case, it does not cause  
32 a circular dependence because it is in a submodule. Submodules cannot be referenced by use association, and  
33 therefore what would be a circular appearance of use palette\_stuff is not accessed.

```

34  program main
35      use color_points
36      ! "instance_count" and "inquire_palette" are not accessible here
37      ! because they are not declared in the "color_points" module.
38      ! "color_points_a" and "color_points_b" cannot be referenced by
39      ! use association.
40      interface draw ! just to demonstrate it's possible
41          module procedure color_point_draw
42      end interface
43      type(color_point) :: C_1, C_2
44      real :: RC
45      ...

```

```

1      call color_point_new (c_1)      ! body in color_points_a, interface in color_points
2      ...
3      call draw (c_1)                 ! body in color_points_b, specific interface
4                                      ! in color_points, generic interface here.
5      ...
6      rc = color_point_dist (c_1, c_2) ! body in color_points_a, interface in color_points
7      ...
8      call color_point_del (c_1)      ! body in color_points_a, interface in color_points
9      ...
10     end program main

```

10 A multilevel submodule system can be used to package and organize a large and interconnected concept without exposing entities of one subsystem to other subsystems.

11 Consider a **Plasma** module from a Tokamak simulator. A plasma simulation requires attention at least to fluid flow, thermodynamics, and electromagnetism. Fluid flow simulation requires simulation of subsonic, supersonic, and hypersonic flow. This problem decomposition can be reflected in the submodule structure of the **Plasma** module:

| Plasma module      |                      |                      |                   |                            |
|--------------------|----------------------|----------------------|-------------------|----------------------------|
| Flow submodule     |                      |                      | Thermal submodule | Electromagnetics submodule |
| Subsonic submodule | Supersonic submodule | Hypersonic submodule |                   |                            |

12 Entities can be shared among the **Subsonic**, **Supersonic**, and **Hypersonic** submodules by putting them within the **Flow** submodule. One then need not worry about accidental use of these entities by use association or by the **Thermal** or **Electromagnetics** submodules, or the development of a dependency of correct operation of those subsystems upon the representation of entities of the **Flow** subsystem as a consequence of maintenance. Since these entities are not accessible by use association, if any of them are changed, the new values cannot be accessed in **program units** that reference the **Plasma** module by use association; the answer to the question “where are these entities used” is therefore confined to the set of descendant submodules of the **Flow** submodule.

## C.9 Clause 12 notes

### C.9.1 Portability problems with external procedures (12.4.3.6)

1 There is a potential portability problem in a **scoping unit** that references an **external procedure** without explicitly declaring it to have the **EXTERNAL attribute** (5.5.9). On a different processor, the name of that procedure might be the name of a **nonstandard intrinsic** procedure and in such a case the processor would interpret those procedure references as references to that intrinsic procedure. (On that processor, the program would also be viewed as not conforming to this part of ISO/IEC 1539 because of the references to the **nonstandard intrinsic** procedure.) Declaration of the **EXTERNAL attribute** causes the references to be to the **external procedure** regardless of the availability of an intrinsic procedure with the same name. Note that declaration of the type of a procedure is not enough to make it **external**, even if the type is inconsistent with the type of the result of an intrinsic procedure of the same name.

## C.9.2 Procedures defined by means other than Fortran (12.6.3)

- 1 A processor is not required to provide any means other than Fortran for defining [external procedures](#). Among the means that might be supported are the machine assembly language, other high level languages, the Fortran language extended with nonstandard features, and the Fortran language as supported by another Fortran processor (for example, a previously existing FORTRAN 77 processor). The means other than Fortran for defining [external procedures](#), including any restrictions on the structure for organization of those procedures, are not specified by this part of ISO/IEC 1539.
- 2 A Fortran processor might limit its support of procedures defined by means other than Fortran such that these procedures can affect entities in the Fortran environment only on the same basis as procedures written in Fortran. For example, it might not support the value of a local variable from being changed by a procedure reference unless that variable were one of the arguments to the procedure.

## C.9.3 Abstract interfaces and procedure pointer components (12.4, 4.5)

- 1 This is an example of a library module providing lists of callbacks that the user can register and invoke.

```

MODULE callback_list_module
    !
    ! Type for users to extend with their own data, if they so desire
    !
    TYPE callback_data
    END TYPE
    !
    ! Abstract interface for the callback procedures
    !
    ABSTRACT INTERFACE
        SUBROUTINE callback_procedure(data)
            IMPORT callback_data
            CLASS(callback_data),OPTIONAL :: data
        END SUBROUTINE
    END INTERFACE
    !
    ! The callback list type.
    !
    TYPE callback_list
        PRIVATE
        CLASS(callback_record),POINTER :: first => NULL()
    END TYPE
    !
    ! Internal: each callback registration creates one of these
    !
    TYPE,PRIVATE :: callback_record
        PROCEDURE(callback_procedure),POINTER,NOPASS :: proc
        CLASS(callback_record),POINTER :: next
        CLASS(callback_data),POINTER :: data => NULL();
    END TYPE

```

```
1      PRIVATE invoke,forward_invoke
2  CONTAINS
3      !
4      ! Register a callback procedure with optional data
5      !
6      SUBROUTINE register_callback(list, entry, data)
7          TYPE(callback_list),INTENT(INOUT) :: list
8          PROCEDURE(callback_procedure) :: entry
9          CLASS(callback_data),OPTIONAL :: data
10         TYPE(callback_record),POINTER :: new,last
11         ALLOCATE(new)
12         new%proc => entry
13         IF (PRESENT(data)) ALLOCATE(new%data,SOURCE=data)
14         new%next => list%first
15         list%first => new
16     END SUBROUTINE
17     !
18     ! Internal: Invoke a single callback and destroy its record
19     !
20     SUBROUTINE invoke(callback)
21         TYPE(callback_record),POINTER :: callback
22         IF (ASSOCIATED(callback%data) THEN
23             CALL callback%proc(list%first%data)
24             DEALLOCATE(callback%data)
25         ELSE
26             CALL callback%proc
27         END IF
28         DEALLOCATE(callback)
29     END SUBROUTINE
30     !
31     ! Call the procedures in reverse order of registration
32     !
33     SUBROUTINE invoke_callback_reverse(list)
34         TYPE(callback_list),INTENT(INOUT) :: list
35         TYPE(callback_record),POINTER :: next,current
36         current => list%first
37         NULLIFY(list%first)
38         DO WHILE (ASSOCIATED(current))
39             next => current%next
40             CALL invoke(current)
41             current => next
42         END DO
43     END SUBROUTINE
44     !
45     ! Internal: Forward mode invocation
46     !
```

```

1      SUBROUTINE forward_invoke(callback)
2          IF (ASSOCIATED(callback%next)) CALL forward_invoke(callback%next)
3          CALL invoke(callback)
4      END SUBROUTINE
5      !
6      ! Call the procedures in forward order of registration
7      !
8      SUBROUTINE invoke_callback_forward(list)
9          TYPE(callback_list), INTENT(INOUT) :: list
10         IF (ASSOCIATED(list%first)) CALL forward_invoke(list%first)
11     END SUBROUTINE
12 END

```

### C.9.4 Pointers and targets as arguments (12.5.2.4, 12.5.2.6, 12.5.2.7)

- 1 If a dummy argument is declared to be a pointer the corresponding **actual argument** could be a pointer, or could be a nonpointer variable. Consider the two cases separately.

*Case (i):* The **actual argument** is a pointer. When procedure execution commences the pointer association status of the dummy argument becomes the same as that of the **actual argument**. If the pointer association status of the dummy argument is changed, the pointer association status of the **actual argument** changes in the same way.

*Case (ii):* The **actual argument** is not a pointer. This only occurs when the **actual argument** has the **TARGET attribute** and the dummy argument has the **INTENT (IN) attribute**. The dummy argument becomes pointer associated with the **actual argument**.

- 2 When execution of a procedure completes, any pointer that remains defined and that is associated with a dummy argument that has the **TARGET attribute** and is either a scalar or an **assumed-shape array**, remains associated with the corresponding **actual argument** if the **actual argument** has the **TARGET attribute** and is not an array section with a **vector subscript**.

- 3 For example, consider:

```

28      REAL, POINTER      :: PBEST
29      REAL, TARGET       :: B (10000)
30      CALL BEST (PBEST, B)      ! Upon return PBEST is associated
31      ...                     ! with the ‘‘best’’ element of B
32  CONTAINS
33      SUBROUTINE BEST (P, A)
34          REAL, POINTER, INTENT (OUT) :: P
35          REAL, TARGET, INTENT (IN)   :: A (:)
36          ...                         ! Find the ‘‘best’’ element A(I)
37          P => A (I)
38      RETURN
39      END SUBROUTINE BEST
40  END

```

When procedure BEST completes, the pointer PBEST is associated with an element of B.

- 4 An [actual argument](#) without the [TARGET attribute](#) can become associated with a dummy argument with the [TARGET attribute](#). This enables a pointer to become associated with the dummy argument during execution of the procedure that contains the dummy argument. For example:

```

4      INTEGER LARGE(100,100)
5      CALL SUB (LARGE)
6      ...
7      CALL SUB ( )
8      CONTAINS
9      SUBROUTINE SUB(ARG)
10         INTEGER, TARGET, OPTIONAL :: ARG(100,100)
11         INTEGER, POINTER, DIMENSION(:,:) :: PARG
12         IF (PRESENT(ARG)) THEN
13             PARG => ARG
14         ELSE
15             ALLOCATE (PARG(100,100))
16             PARG = 0
17         ENDIF
18         ... ! Code with lots of references to PARG
19         IF (.NOT. PRESENT(ARG)) DEALLOCATE(PARG)
20     END SUBROUTINE SUB
21 END

```

Within subroutine SUB the pointer PARG is either associated with the dummy argument ARG or it is associated with an allocated target. The bulk of the code can reference PARG without further calls to the intrinsic function [PRESENT](#).

- 5 If a nonpointer dummy argument has the [TARGET attribute](#) and the corresponding [actual argument](#) does not, any pointers that become associated with the dummy argument, and therefore with the [actual argument](#), during execution of the procedure, become undefined when execution of the procedure completes.

## C.9.5 Polymorphic Argument Association (12.5.2.9)

- 1 The following example illustrates the polymorphic [argument association](#) rules using the derived types defined in Note 4.55.

```

31      TYPE(POINT) :: T2
32      TYPE(COLOR_POINT) :: T3
33      CLASS(POINT) :: P2
34      CLASS(COLOR_POINT) :: P3
35      ! Dummy argument is polymorphic and actual argument is of fixed type
36      SUBROUTINE SUB2 ( X2 ); CLASS(POINT) :: X2; ...
37      SUBROUTINE SUB3 ( X3 ); CLASS(COLOR_POINT) :: X3; ...
38
39      CALL SUB2 ( T2 ) ! Valid -- The declared type of T2 is the same as the
40                        !           declared type of X2.
41      CALL SUB2 ( T3 ) ! Valid -- The declared type of T3 is extended from
42                        !           the declared type of X2.

```

```

1      CALL SUB3 ( T2 ) ! Invalid -- The declared type of T2 is neither the
2                                ! same as nor extended from the declared type
3                                ! type of X3.
4      CALL SUB3 ( T3 ) ! Valid -- The declared type of T3 is the same as the
5                                ! declared type of X3.
6      ! Actual argument is polymorphic and dummy argument is of fixed type
7      SUBROUTINE TUB2 ( D2 ); TYPE(POINT) :: D2; ...
8      SUBROUTINE TUB3 ( D3 ); TYPE(COLOR_POINT) :: D3; ...
9
10     CALL TUB2 ( P2 ) ! Valid -- The declared type of P2 is the same as the
11                               ! declared type of D2.
12     CALL TUB2 ( P3 ) ! Invalid -- The declared type of P3 differs from the
13                               ! declared type of D2.
14     CALL TUB2 ( P3%POINT ) ! Valid alternative to the above
15     CALL TUB3 ( P2 ) ! Invalid -- The declared type of P2 differs from the
16                               ! declared type of D3.
17     SELECT TYPE ( P2 ) ! Valid conditional alternative to the above
18     CLASS IS ( COLOR_POINT ) ! Works if the dynamic type of P2 is the same
19         CALL TUB3 ( P2 )      ! as the declared type of D3, or a type
20                               ! extended therefrom.
21     CLASS DEFAULT
22                               ! Cannot work if not.
23     END SELECT
24     CALL TUB3 ( P3 ) ! Valid -- The declared type of P3 is the same as the
25                               ! declared type of D3.
26     ! Both the actual and dummy arguments are of polymorphic type.
27     CALL SUB2 ( P2 ) ! Valid -- The declared type of P2 is the same as the
28                               ! declared type of X2.
29     CALL SUB2 ( P3 ) ! Valid -- The declared type of P3 is extended from
30                               ! the declared type of X2.
31     CALL SUB3 ( P2 ) ! Invalid -- The declared type of P2 is neither the
32                               ! same as nor extended from the declared
33                               ! type of X3.
34     SELECT TYPE ( P2 ) ! Valid conditional alternative to the above
35     CLASS IS ( COLOR_POINT ) ! Works if the dynamic type of P2 is the
36         CALL SUB3 ( P2 )      ! same as the declared type of X3, or a
37                               ! type extended therefrom.
38     CLASS DEFAULT
39                               ! Cannot work if not.
40     END SELECT
41     CALL SUB3 ( P3 ) ! Valid -- The declared type of P3 is the same as the
42                               ! declared type of X3.

```

### C.9.6 Rules ensuring unambiguous generics (12.4.3.5.5)

1 The rules in 12.4.3.5.5 are intended to ensure

- that it is possible to reference each specific procedure or binding in the generic collection,
- that for any valid generic procedure reference, the determination of the specific procedure referenced is unambiguous, and
- that the determination of the specific procedure or binding referenced can be made before execution of the program begins (during compilation).

Interfaces of specific procedures or bindings are distinguished by fixed properties of their arguments, specifically type, kind type parameters, [rank](#), and whether the dummy argument has the [POINTER](#) or [ALLOCATABLE](#) attribute. A valid reference to one procedure in a generic collection will differ from another because it has an argument that the other cannot accept, because it is missing an argument that the other requires, or because one of these fixed properties is different.

Although the declared type of a data entity is a fixed property, polymorphic variables allow for a limited degree of type mismatch between dummy arguments and [actual arguments](#), so the requirement for distinguishing two dummy arguments is type incompatibility, not merely different types. (This is illustrated in the BAD6 example later in this note.)

That same limited type mismatch means that two dummy arguments that are not type incompatible can be distinguished on the basis of the values of the kind type parameters they have in common; if one of them has a kind type parameter that the other does not, that is irrelevant in distinguishing them.

[Rank](#) is a fixed property, but some forms of array dummy arguments allow [rank](#) mismatches when a procedure is referenced by its specific name. In order to allow [rank](#) to always be usable in distinguishing generics, such [rank](#) mismatches are disallowed for those arguments when the procedure is referenced as part of a generic. Additionally, the fact that elemental procedures can accept array arguments is not taken into account when applying these rules, so apparent ambiguity between elemental and nonelemental procedures is possible; in such cases, the reference is interpreted as being to the nonelemental procedure.

For procedures referenced as operators or defined-assignment, syntactically distinguished arguments are mapped to specific positions in the argument list, so the rule for distinguishing such procedures is that it be possible to distinguish the arguments at one of the argument positions.

For [defined input/output](#) procedures, only the `dtv` argument corresponds to something explicitly written in the program, so it is the `dtv` that is required to be distinguished. Because `dtv` arguments are required to be scalar, they cannot differ in [rank](#). Thus this rule effectively involves only type and kind type parameters.

For generic procedure names, the rules are more complicated because optional arguments can be omitted and because arguments can be specified either positionally or by name.

In the special case of type-bound procedures with [passed-object dummy arguments](#), the passed-object argument is syntactically distinguished in the reference, so rule (3) in 12.4.3.5.5 can be applied. The type of passed-object arguments is constrained in ways that prevent passed-object arguments in the same [scoping unit](#) from being type incompatible. Thus this rule effectively involves only kind type parameters and [rank](#).

The primary means of distinguishing named generics is rule (4). The most common application of that rule is a single argument satisfying both (4a) and (4b):

```
INTERFACE GOOD1
  FUNCTION F1A(X)
    REAL :: F1A,X
```

```

1      END FUNCTION F1A
2      FUNCTION F1B(X)
3          INTEGER :: F1B,X
4      END FUNCTION F1B
5      END INTERFACE GOOD1

```

11 Whether one writes `GOOD1(1.0)` or `GOOD1(X=1.0)`, the reference is to `F1A` because `F1B` would require an integer argument whereas these references provide the real constant 1.0.

12 This example and those that follow are expressed using interface bodies, with type as the distinguishing property. This was done to make it easier to write and describe the examples. The principles being illustrated are equally applicable when the procedures get their explicit interfaces in some other way or when kind type parameters or [rank](#) are the distinguishing property.

13 Another common variant is the argument that satisfies (4a) and (4b) by being required in one specific and completely missing in the other:

```

14      INTERFACE GOOD2
15          FUNCTION F2A(X)
16              REAL :: F2A,X
17          END FUNCTION F2A
18          FUNCTION F2B(X,Y)
19              COMPLEX :: F2B
20              REAL :: X,Y
21          END FUNCTION F2B
22      END INTERFACE GOOD2

```

14 Whether one writes `GOOD2(0.0,1.0)`, `GOOD2(0.0,Y=1.0)`, or `GOOD2(Y=1.0,X=0.0)`, the reference is to `F2B`, because `F2A` has no argument in the second position or with the name `Y`. This approach is used as an alternative to optional arguments when one wants a function to have different result type, kind type parameters, or [rank](#), depending on whether the argument is present. In many of the intrinsic functions, the `DIM` argument works this way.

15 It is possible to construct cases where different arguments are used to distinguish positionally and by name:

```

29      INTERFACE GOOD3
30          SUBROUTINE S3A(W,X,Y,Z)
31              REAL :: W,Y
32              INTEGER :: X,Z
33          END SUBROUTINE S3A
34          SUBROUTINE S3B(X,W,Z,Y)
35              REAL :: W,Z
36              INTEGER :: X,Y
37          END SUBROUTINE S3B
38      END INTERFACE GOOD3

```

16 If one writes `GOOD3(1.0,2,3.0,4)` to reference `S3A`, then the third and fourth arguments are consistent with a reference to `S3B`, but the first and second are not. If one switches to writing the first two arguments as keyword

arguments in order for them to be consistent with a reference to S3B, the latter two arguments must also be written as keyword arguments, GOOD3(X=2,W=1.0,Z=4,Y=3.0), and the named arguments Y and Z are distinguished.

17 The ordering requirement in rule (4) is critical:

```

4      INTERFACE BAD4  ! this interface is invalid !
5          SUBROUTINE S4A(W,X,Y,Z)
6              REAL :: W,Y
7              INTEGER :: X,Z
8          END SUBROUTINE S4A
9          SUBROUTINE S4B(X,W,Z,Y)
10             REAL :: X,Y
11             INTEGER :: W,Z
12         END SUBROUTINE S4B
13     END INTERFACE BAD4

```

18 In this example, the positionally distinguished arguments are Y and Z, and it is W and X that are distinguished by name. In this order it is possible to write BAD4(1.0,2,Y=3.0,Z=4), which is a valid reference for both S4A and S4B.

19 Rule (1) can be used to distinguish some cases that are not covered by rule (4):

```

18     INTERFACE GOOD5
19         SUBROUTINE S5A(X)
20             REAL :: X
21         END SUBROUTINE S5A
22         SUBROUTINE S5B(Y,X)
23             REAL :: Y,X
24         END SUBROUTINE S5B
25     END INTERFACE GOOD5

```

20 In attempting to apply rule (4), position 2 and name Y are distinguished, but they are in the wrong order, just like the BAD4 example. However, when we try to construct a similarly ambiguous reference, we get GOOD5(1.0,X=2.0), which can't be a reference to S5A because it would be attempting to associate two different [actual arguments](#) with the dummy argument X. Rule (4) catches this case by recognizing that S5B requires two real arguments, and S5A cannot possibly accept more than one.

21 The application of rule (1) becomes more complicated when [extensible types](#) are involved. If FRUIT is an [extensible type](#), PEAR and APPLE are [extensions](#) of FRUIT, and BOSC is an [extension](#) of PEAR, then

```

33     INTERFACE BAD6  ! this interface is invalid !
34         SUBROUTINE S6A(X,Y)
35             CLASS(PEAR) :: X,Y
36         END SUBROUTINE S6A
37         SUBROUTINE S6B(X,Y)
38             CLASS(FRUIT) :: X
39             CLASS(BOSC) :: Y
40         END SUBROUTINE S6B
41     END INTERFACE BAD6

```

1 might, at first glance, seem distinguishable this way, but because of the limited type mismatching allowed,  
 2 BAD6(A\_PEAR,A\_BOSC) is a valid reference to both S6A and S6B.

3 22 It is important to try rule (1) for each type that appears:

```

4      INTERFACE GOOD7
5          SUBROUTINE S7A(X,Y,Z)
6              CLASS(PEAR) :: X,Y,Z
7          END SUBROUTINE S7A
8          SUBROUTINE S7B(X,Z,W)
9              CLASS(FRUIT) :: X
10             CLASS(BOSC) :: Z
11             CLASS(APPLE),OPTIONAL :: W
12          END SUBROUTINE S7B
13      END INTERFACE GOOD7

```

14 23 Looking at the most general type, S7A has a minimum and maximum of 3 FRUIT arguments, while S7B has a  
 15 minimum of 2 and a maximum of three. Looking at the most specific, S7A has a minimum of 0 and a maximum  
 16 of 3 BOSC arguments, while S7B has a minimum of 1 and a maximum of 2. However, when we look at the  
 17 intermediate, S7A has a minimum and maximum of 3 PEAR arguments, while S7B has a minimum of 1 and a  
 18 maximum of 2. Because S7A's minimum exceeds S7B's maximum, they can be distinguished.

19 24 In identifying the minimum number of arguments with a particular set of properties, we exclude optional argu-  
 20 ments and test TKR compatibility, so the corresponding [actual arguments](#) are required to have those properties.  
 21 In identifying the maximum number of arguments with those properties, we include the optional arguments and  
 22 test not distinguishable, so we include [actual arguments](#) which could have those properties but are not required  
 23 to have them.

24 25 These rules are sufficient to ensure that references to procedures that meet them are unambiguous, but there  
 25 remain examples that fail to meet these rules but which can be shown to be unambiguous:

```

26 26      INTERFACE BAD8 ! this interface is invalid !
27          ! despite the fact that it is unambiguous !
28          SUBROUTINE S8A(X,Y,Z)
29              REAL,OPTIONAL :: X
30              INTEGER :: Y
31              REAL :: Z
32          END SUBROUTINE S8A
33          SUBROUTINE S8B(X,Z,Y)
34              INTEGER,OPTIONAL :: X
35              INTEGER :: Z
36              REAL :: Y
37          END SUBROUTINE S8B
38      END INTERFACE BAD8

```

39 27 This interface fails rule (4) because there are no required arguments that can be distinguished from the positionally  
 40 corresponding argument, but in order for the mismatch of the optional arguments not to be relevant, the later  
 41 arguments must be specified as keyword arguments, so distinguishing by name does the trick. This interface is

nevertheless invalid so a standard-conforming Fortran processor is not required to do such reasoning. The rules to cover all cases are too complicated to be useful.

The real data objects that would be valid arguments for **S9A** are entirely disjoint from procedures that are valid arguments to **S9B** and **S9C**, and the procedures that valid arguments for **S9B** are disjoint from the procedures that are valid arguments to **S9C** because the former are required to accept real arguments and the latter integer arguments. Again, this interface is invalid, so a standard-conforming Fortran processor need not examine such properties when deciding whether a generic collection is valid. Again, the rules to cover all cases are too complicated to be useful.

If one dummy argument has the **POINTER** attribute and a corresponding argument in the other interface body has the **ALLOCATABLE** attribute the generic interface is not ambiguous. If one dummy argument has either the **POINTER** or **ALLOCATABLE** attribute and a corresponding argument in the other interface body has neither attribute, the generic interface might be ambiguous.

## C.10 Clause 13 notes

### C.10.1 Atomic memory consistency

#### C.10.1.1 Relaxed memory model

Parallel programs sometimes have apparently impossible behavior because data transfers and other messages can be delayed, reordered and even repeated, by hardware, communication software, and caching and other forms of optimization. Requiring processors to deliver globally consistent behavior is incompatible with performance on many systems. This part of ISO/IEC 1539 specifies that all ordered actions will be consistent (2.3.5 and 8.6), but all consistency between unordered segments is deliberately left processor dependent. Depending on the hardware, this can be observed even when only two images and one mechanism are involved.

#### C.10.1.2 Examples with atomic operations

1 When variables are being referenced (atomically) from segments that are unordered with respect to the segment that is atomically defining or redefining the variables, the results are processor dependent. This supports use of so-called “relaxed memory model” architectures, which can enable more efficient execution on some hardware implementations.

2 The following examples assume these declarations:

```
MODULE EXAMPLE
  USE, INTRINSIC :: ISO_FORTRAN_ENV
  INTEGER(ATOMIC_INT_KIND) :: X [*] = 0, Y [*] = 0
```

#### 3 Example 1

4 With  $X[j]$  and  $Y[j]$  still in their initial state (both zero), image  $j$  executes the following sequence of statements:

```
CALL ATOMIC_DEFINE (X, 1)
CALL ATOMIC_DEFINE (Y, 1)
```

and a different image,  $k$ , executes the following sequence of statements:

```

1      DO
2          CALL ATOMIC_REF (TMP, Y [j])
3          IF (TMP==1) EXIT
4      END DO
5      CALL ATOMIC_REF (TMP, X [j])
6      PRINT *, TMP

```

7 5 The final value of TMP on image  $k$  can be either 0 or 1. That is, even though image  $j$  thinks that it defined X  
8  $[j]$  before it defined Y  $[j]$ , this ordering is not guaranteed to be observed on image  $k$ . There are many aspects of  
9 hardware and software implementation that can cause this effect, but conceptually this example can be thought  
10 of as the change in the value of Y propagating faster through the inter-image connections than the change in the  
11 value of X.

12 6 Even if image  $j$  executed the sequence

```

13      CALL ATOMIC_DEFINE (X, 1)
14      SYNC MEMORY
15      CALL ATOMIC_DEFINE (Y, 1)

```

16 the same effect can be seen. That is because even though X and Y are defined in ordered segments, the references  
17 from image  $k$  are both from a segment that is unordered with respect to image  $j$ .

18 7 Only if the reference on image  $k$  to Y  $[j]$  is in a segment that is ordered after the segment on image  $j$  that defined  
19 Y, will TMP be guaranteed to have the value 1.

#### 20 8 **Example 2:**

21 With the initial state of X and Y on image  $j$  (i.e. X  $[j]$  and Y  $[j]$ ) still being zero, execution of

```

22      CALL ATOMIC_REF (TMP, X [j])
23      CALL ATOMIC_DEFINE (Y [j], 1)
24      PRINT *, TMP

```

25 on image  $k_1$ , and execution of

```

26      CALL ATOMIC_REF (TMP, Y [j])
27      CALL ATOMIC_DEFINE (X [j], 1)
28      PRINT *, TMP

```

29 on image  $k_2$ , in unordered segments, might print the value 1 both times.

30 9 This can happen by such mechanisms as “load buffering”; one might imagine that what is happening is that  
31 the definitions (ATOMIC\_DEFINE) are overtaking the references (ATOMIC\_REF). On some processors it is  
32 possible that insertion of SYNC MEMORY statements between the calls to ATOMIC\_REF and ATOMIC\_  
33 DEFINE might be sufficient to make the output print the value 1 at most one time (or even exactly one time),  
34 but this is still processor dependent unless the SYNC MEMORY statement executions cause the relevant segments  
35 on images  $k_1$  and  $k_2$  to be ordered.

### 36 **C.10.2 EVENT\_QUERY example**

37 1 The following example illustrates the use of events via a program in which image 1 acts as the master image,  
38 distributing work items to the other images. Only one work item at a time can be active on a worker image, and  
39 each deals with the result (e.g. via input/output) without directly feeding data back to the master image.

1 2 Because the work items are not expected to be balanced, the master keeps cycling through all images to find one  
2 that is waiting for work.

3 3 An event is posted by each worker to indicate that it has completed its work item. Since the corresponding  
4 variables are needed only on the master, we place them in an allocatable array component of a coarray. An event  
5 on each worker is needed for the master to post the fact that it has made a work item available for it.

6 **Example code:**

```

7      PROGRAM work_share
8          USE, INTRINSIC :: ISO_FORTRAN_ENV, ONLY: EVENT_TYPE
9          USE :: mod_work, ONLY:                & ! Module that creates work items
10             work,                               & ! Type for holding a work item
11             create_work_item, & ! Function that creates work item
12             process_item,    & ! Function that processes an item
13             work_done        ! Logical function that returns true
14                               ! if all work has been done.
15
16      TYPE :: worker_type
17          TYPE (EVENT_TYPE), ALLOCATABLE :: free (:)
18      END TYPE
19      TYPE (EVENT_TYPE) :: submit [*]    ! Post when work ready for a worker
20      TYPE (worker_type) :: worker [*]   ! Post when worker is free
21      TYPE (work) :: work_item [*] ! Holds the data for a work item
22      INTEGER :: count, i, nbusy [*]
23
24      IF (THIS_IMAGE ()==1) THEN
25          ! Get started
26          ALLOCATE (worker%free (2:NUM_IMAGES ()))
27          nbusy = 0 ! This holds the number of workers working
28          DO i = 2, NUM_IMAGES () ! Start the workers working
29              IF (work_done ()) EXIT
30              nbusy = nbusy + 1
31              work_item [i] = create_work_item ()
32              EVENT POST (submit [i])
33          END DO
34          ! Main work distribution loop
35      master: DO
36          image: DO i = 2, NUM_IMAGES ()
37              CALL EVENT_QUERY (worker%free (i), count)
38              IF (count==0) CYCLE image ! Worker is not free
39              EVENT WAIT (worker%free (i))
40              nbusy = nbusy - 1
41              IF (work_done ()) CYCLE
42              nbusy = nbusy + 1
43              work_item [i] = create_work_item ()
44              EVENT POST (submit [i])
45          END DO image

```

```

1          IF (nbusy==0) THEN
2              ! All done. Exit on all images.
3              DO i = 2, NUM_IMAGES ()
4                  EVENT POST (submit [i])
5              END DO
6              EXIT master
7          END IF
8      END DO master
9  ELSE
10         ! Work processing loop
11  worker: DO
12         EVENT WAIT (submit)
13         IF (nbusy[1] == 0) EXIT
14         CALL process_item (work_item)
15         EVENT POST (worker [1]%free (THIS_IMAGE ()))
16     END DO worker
17 END IF
18 END PROGRAM work_share

```

## C.11 Clause 13 notes

### C.11.1 Collective subroutine examples

- 1 The following example computes a dot product of two scalar coarrays using `CO_SUM` to store the result in a noncoarray scalar variable.

```

23      SUBROUTINE codot (x, y, x_dot_y)
24          REAL :: x [*], y [*], x_dot_y
25          x_dot_y = x*y
26          CALL CO_SUM (x_dot_y)
27      END SUBROUTINE codot

```

- 2 The function below demonstrates passing a noncoarray dummy argument to `CO_MAX`. The function uses `CO_MAX` to find the maximum value of the dummy argument across all images. Then the function flags all images that hold values matching the maximum. The function then returns the maximum image index for an image that holds the maximum value.

```

32      FUNCTION find_max (j) RESULT (j_max_location)
33          INTEGER, INTENT (IN) :: j
34          INTEGER j_max, j_max_location
35          j_max = j
36          CALL CO_MAX (j_max)
37          ! Flag images that hold the maximum j.
38          IF (j==j_max) THEN
39              j_max_location = THIS_IMAGE ()
40          ELSE
41              j_max_location = 0

```

```

1      END IF
2      ! Return highest image index associated with a maximal j.
3      CALL CO_MAX(j_max_location)
4      END FUNCTION find_max

```

## 5 C.11.2 Atomic memory consistency

### 6 C.11.2.1 Relaxed memory model

- 7 1 Parallel programs sometimes have apparently impossible behavior because data transfers and other messages  
8 can be delayed, reordered and even repeated, by hardware, communication software, caching, and other forms  
9 of optimization. Requiring processors to deliver globally consistent behavior is incompatible with performance  
10 on many systems. Therefore all ordered actions are required to be consistent (8.6), but consistency between  
11 unordered segments is processor dependent. Depending on the hardware, this could be observed even when only  
12 two images and one mechanism are involved.

### 13 C.11.2.2 Examples with atomic operations

- 14 1 When variables are being referenced (atomically) from segments that are unordered with respect to the segment  
15 that is is atomically defining or redefining the variables, the results are processor dependent. This supports use  
16 of so-called “relaxed memory model” architectures, which can enable more efficient execution on some hardware.
- 17 2 The following examples assume the following declarations:

```

18      MODULE example
19      USE, INTRINSIC :: ISO_FORTRAN_ENV
20      INTEGER (ATOMIC_INT_KIND) :: X [*] = 0, Y [*] = 0

```

#### 21 Example 1:

- 22 3 With  $X[j]$  and  $Y[j]$  still in their initial state (both zero), image  $j$  executes the following sequence of statements:

```

23      CALL ATOMIC_DEFINE (X, 1)
24      CALL ATOMIC_DEFINE (Y, 1)

```

25 and image  $k$  executes the following sequence of statements:

```

26      DO
27      CALL ATOMIC_REF (TMP, Y [j])
28      IF (tmp==1) EXIT
29      END DO
30      CALL ATOMIC_REF (TMP, X [j])
31      PRINT *, tmp

```

- 32 4 The final value of TMP on image  $k$  could be either 0 or 1. That is, even though image  $j$  thinks it wrote  $X[j]$   
33 before writing  $Y[j]$ , this ordering is not guaranteed to be that observed on image  $k$ .
- 34 5 There are many aspects of hardware and software implementation that can cause this effect, but conceptually this  
35 example can be thought of as the change in the value of  $Y$  propagating faster across the inter-image connections  
36 than the change in the value of  $X$ .

1 6 Changing the execution on image  $j$  by inserting

2       **SYNC MEMORY**

3       in between the definitions of  $X$  and  $Y$  is not sufficient to prevent unexpected results; even though  $X$  and  $Y$  are  
4       being updated in ordered segments, the references from image  $k$  are both from a segment that is unordered with  
5       respect to those on image  $j$ .

6 7 To guarantee the expected value for  $TMP$  of 1 at the end of the code sequence on image  $k$ , it is necessary to  
7       ensure that the atomic reference on image  $k$  is in a segment that is ordered relative to the segment on image  $j$   
8       that defined  $X[j]$ ; **SYNC MEMORY** is certainly necessary, but not sufficient unless it is somehow synchronized.

9       **Example 2:**

10 8 With the initial state of  $X$  and  $Y$  on image  $j$  (i.e.  $X[j]$  and  $Y[j]$ ) still being zero, execution of

```
11       CALL ATOMIC_REF (TMP, X [j])
12       CALL ATOMIC_DEFINE (Y [j], 1)
13       PRINT *, TMP
```

14       on image  $k_1$ , and execution of

```
15       CALL ATOMIC_REF (TMP, Y [j])
16       CALL ATOMIC_DEFINE (X [j], 1)
17       PRINT *, TMP
```

18       on image  $k_2$ , in unordered segments, might print the value 1 both times.

19 9 This can happen by such mechanisms as “load buffering”; one might imagine that what is happening is that the  
20       writes (**ATOMIC\_DEFINE**) are overtaking the reads (**ATOMIC\_REF**).

21 10 It is likely that insertion of **SYNC MEMORY** between the calls to **ATOMIC\_REF** and **ATOMIC\_DEFINE**  
22       would be sufficient to prevent this anomalous behavior, but that is only guaranteed by the standard if the **SYNC**  
23       **MEMORY** executions cause an ordering between the relevant segments on images  $k_1$  and  $k_2$ .

24       **Example 3:**

25 11 Because there are no segment boundaries implied by collective subroutines, with the initial state as before,  
26       execution of

```
27       IF (THIS_IMAGE ()==1) THEN
28       CALL ATOMIC_DEFINE (X [3], 23)
29       Y = 42
30       END IF
31       CALL CO_BROADCAST (Y, 1)
32       IF (THIS_IMAGE ()==2) THEN
33       CALL ATOMIC_REF (TMP, X [3])
34       PRINT *, Y, TMP
35       END IF
```

36       could print the values 42 and 0.

37       **Example 4:**

1 12 Assuming the declarations

2           INTEGER (ATOMIC\_INT\_KIND) :: X [\*] = 0, Z = 0

3 the statements

4           CALL ATOMIC\_ADD (X [1], 1)           ! (A)

5           IF (THIS\_IMAGE() == 2) THEN

6           wait: DO

7           CALL ATOMIC\_REF (Z, X [1])       ! (B)

8           IF (Z == NUM\_IMAGES ()) EXIT wait

9           END DO wait                       ! (C)

10          END IF

11 will execute the “wait” loop on image 2 until all images have completed statement (A). The updates of X [1] are  
 12 performed by each image in the same manner, but in an arbitrary order. Because the result from the complete  
 13 set of updates will eventually become visible by execution of statement (B) for some loop iteration on image 2,  
 14 the termination condition is guaranteed to be eventually fulfilled, provided that no image failure occurs, every  
 15 image executes the above code, and no other code is executed in an unordered segment that performs an update  
 16 to X [1]. Furthermore, if two SYNC MEMORY statements are inserted in the above code before statement (A)  
 17 and after statement (C), respectively, the segment started by the second SYNC MEMORY on image 2 is ordered  
 18 after the segments on all images that end with the first SYNC MEMORY.

## 19 C.12 Clause 15 notes

### 20 C.12.1 Runtime environments (15.1)

21 1 This part of ISO/IEC 1539 allows programs to contain procedures defined by means other than Fortran. That  
 22 raises the issues of initialization of and interaction between the runtime environments involved.

23 2 Implementations are free to solve these issues as they see fit, provided that

- 24 • heap allocation/deallocation (e.g., (DE)ALLOCATE in a Fortran subprogram and malloc/free in a C func-  
 25 tion) can be performed without interference,
- 26 • input/output to and from external files can be performed without interference, as long as procedures defined  
 27 by different means do not do input/output with the same external file,
- 28 • input/output preconnections exist as required by the respective standards, and
- 29 • initialized data are initialized according to the respective standards.

### 30 C.12.2 Example of Fortran calling C (15.3)

31 C Function Prototype:

32           int C\_Library\_Function(void\* sendbuf, int sendcount, int \*recvcounts);

33 Fortran Module:

34           MODULE CLIBFUN\_INTERFACE

35           INTERFACE

```

1      INTEGER (C_INT) FUNCTION C_LIBRARY_FUNCTION (SENDBUF, SENDCOUNT, RECVCOUNTS) &
2          BIND(C, NAME='C_Library_Function')
3          USE, INTRINSIC :: ISO_C_BINDING
4          IMPLICIT NONE
5          TYPE (C_PTR), VALUE :: SENDBUF
6          INTEGER (C_INT), VALUE :: SENDCOUNT
7          INTEGER (C_INT) :: RECVCOUNTS(*)
8      END FUNCTION C_LIBRARY_FUNCTION
9  END INTERFACE
10 END MODULE CLIBFUN_INTERFACE

```

- 11 1 The module CLIBFUN\_INTERFACE contains the declaration of the Fortran dummy arguments, which correspond to the C formal parameters. The **NAME=** is used in the **BIND attribute** in order to handle the case-sensitive name change between Fortran and C from “c\_library\_function” to “C\_Library\_Function”.
- 14 2 The first C formal parameter is the pointer to void **sendbuf**, which corresponds to the Fortran dummy argument **SENDBUF**, which has the type **C\_PTR** and the **VALUE attribute**.
- 16 3 The second C formal parameter is the int **sendcount**, which corresponds to the Fortran dummy argument **SENDCOUNT**, which has the type **INTEGER (C\_INT)** and the **VALUE attribute**.
- 18 4 The third C formal parameter is the pointer to int **recvcounts**, which corresponds to the Fortran dummy argument **RECVCOUNTS**, which is an **assumed-size array** of type **INTEGER (C\_INT)**.
- 20 5 This example show how **C\_Library\_Function** might be referenced in a Fortran **program unit**:

```

21      USE, INTRINSIC :: ISO_C_BINDING, ONLY: C_INT, C_FLOAT, C_LOC
22      USE CLIBFUN_INTERFACE
23      ...
24      REAL (C_FLOAT), TARGET :: SEND(100)
25      INTEGER (C_INT) :: SENDCOUNT, RET
26      INTEGER (C_INT), ALLOCATABLE :: RECVCOUNTS(:)
27      ...
28      ALLOCATE( RECVCOUNTS(100) )
29      ...
30      RET = C_LIBRARY_FUNCTION(C_LOC(SEND), SENDCOUNT, RECVCOUNTS)
31      ...

```

- 32 6 The first Fortran actual argument is a reference to the function **C\_LOC** which returns the value of the C address of its argument, **SEND**. This value becomes the value of the first formal parameter, the pointer **sendbuf**, in **C\_Library\_Function**.
- 35 7 The second Fortran actual argument is **SENDCOUNT** of type **INTEGER (C\_INT)**. Its value becomes the initial value of the second formal parameter, the int **sendcount**, in **C\_Library\_Function**.
- 37 8 The third Fortran actual argument is the allocatable array **RECVCOUNTS** of type **INTEGER (C\_INT)**. The base C address of this array becomes the value of the third formal parameter, the pointer **recvcounts**, in **C\_Library\_Function**. Note that interoperability is based on the characteristics of the dummy arguments in the specified interface and not on those of the actual arguments. Thus, the fact that the actual argument is

allocatable is not relevant here.

### C.12.3 Example of C calling Fortran (15.3)

#### Fortran Code:

```

SUBROUTINE SIMULATION(ALPHA, BETA, GAMMA, DELTA, ARRAYS) BIND(C)
  USE, INTRINSIC :: ISO_C_BINDING
  IMPLICIT NONE
  INTEGER (C_LONG), VALUE                :: ALPHA
  REAL (C_DOUBLE), INTENT(INOUT)         :: BETA
  INTEGER (C_LONG), INTENT(OUT)          :: GAMMA
  REAL (C_DOUBLE), DIMENSION(*), INTENT(IN) :: DELTA
  TYPE, BIND(C) :: PASS
    INTEGER (C_INT) :: LENC, LENF
    TYPE (C_PTR)    :: C, F
  END TYPE PASS
  TYPE (PASS), INTENT(INOUT) :: ARRAYS
  REAL (C_FLOAT), ALLOCATABLE, TARGET, SAVE :: ETA(:)
  REAL (C_FLOAT), POINTER :: C_ARRAY(:)
  ...
  ! Associate C_ARRAY with an array allocated in C
  CALL C_F_POINTER (ARRAYS%C, C_ARRAY, [ ARRAYS%LENC ])
  ...
  ! Allocate an array and make it available in C
  ARRAYS%LENF = 100
  ALLOCATE (ETA(ARRAYS%LENF))
  ARRAYS%F = C_LOC(ETA)
  ...
END SUBROUTINE SIMULATION

```

#### C Struct Declaration:

```

struct pass {
  int lenc, lenf;
  float *c, *f;
};

```

#### C Function Prototype:

```

void simulation(long alpha, double *beta, long *gamma, double delta[],
               struct pass *arrays);

```

#### C Calling Sequence:

```

simulation(alpha, beta, gamma, delta, arrays);

```

- 1 1 The above-listed Fortran code specifies a subroutine SIMULATION. This subroutine corresponds to the C void  
2 function `simulation`.
- 3 2 The Fortran subroutine references the intrinsic module `ISO_C_BINDING`.
- 4 3 The first Fortran dummy argument of the subroutine is ALPHA, which has the type `INTEGER(C_LONG)` and  
5 the [VALUE attribute](#). This dummy argument corresponds to the C formal parameter `alpha`, which is a long.  
6 The C actual argument is also a long.
- 7 4 The second Fortran dummy argument of the subroutine is BETA, which has the type `REAL(C_DOUBLE)` and  
8 the [INTENT \(INOUT\) attribute](#). This dummy argument corresponds to the C formal parameter `beta`, which is  
9 a pointer to double. An address is passed as the C actual argument.
- 10 5 The third Fortran dummy argument of the subroutine is GAMMA, which has the type `INTEGER(C_LONG)`  
11 and the [INTENT \(OUT\) attribute](#). This dummy argument corresponds to the C formal parameter `gamma`, which  
12 is a pointer to long. An address is passed as the C actual argument.
- 13 6 The fourth Fortran dummy argument is the [assumed-size array](#) DELTA, which has the type `REAL (C_DOUBLE)`  
14 and the [INTENT \(IN\) attribute](#). This dummy argument corresponds to the C formal parameter `delta`, which is  
15 a double array. The C actual argument is also a double array.
- 16 7 The fifth Fortran dummy argument is ARRAYS, which is a structure for accessing an array allocated in C and  
17 an array allocated in Fortran. The lengths of these arrays are held in the components `LENC` and `LENF`; their [C](#)  
18 [addresses](#) are held in components `C` and `F`.

#### 19 C.12.4 Example of calling C functions with noninteroperable data (15.10)

- 20 1 Many Fortran processors support 16-byte real numbers, which might not be supported by the C processor.  
21 Assume a Fortran programmer wants to use a C procedure from a message passing library for an array of these  
22 reals. The C prototype of this procedure is

```
23 void ProcessBuffer(void *buffer, int n_bytes);
```

24 with the corresponding Fortran interface

```
25 USE, INTRINSIC :: ISO_C_BINDING
26 INTERFACE
27   SUBROUTINE PROCESS_BUFFER(BUFFER,N_BYTES) BIND(C,NAME="ProcessBuffer")
28     IMPORT :: C_PTR, C_INT
29     TYPE(C_PTR), VALUE :: BUFFER ! The 'C address' of the array buffer
30     INTEGER (C_INT), VALUE :: N_BYTES ! Number of bytes in buffer
31   END SUBROUTINE PROCESS_BUFFER
32 END INTERFACE
```

- 33 2 This can be done using [C\\_LOC](#) if the particular Fortran processor specifies that [C\\_LOC](#) returns an appropriate  
34 address:

```
35 REAL(R_QUAD), DIMENSION(:), ALLOCATABLE, TARGET :: QUAD_ARRAY
36 ...
37 CALL PROCESS_BUFFER(C_LOC(QUAD_ARRAY), INT(16*SIZE(QUAD_ARRAY),C_INT))
38 ! One quad real takes 16 bytes on this processor
```

## C.12.5 Example of opaque communication between C and Fortran (15.3)

- 1 The following example demonstrates how a Fortran processor can make a modern object-oriented random number generator written in Fortran available to a C program.

```

4      USE, INTRINSIC :: ISO_C_BINDING
5      ! Assume this code is inside a module
6
7      TYPE RANDOM_STREAM
8      ! A (uniform) random number generator (URNG)
9      CONTAINS
10     PROCEDURE(RANDOM_UNIFORM), DEFERRED, PASS(STREAM) :: NEXT
11     ! Generates the next number from the stream
12 END TYPE RANDOM_STREAM
13
14 ABSTRACT INTERFACE
15     ! Abstract interface of Fortran URNG
16     SUBROUTINE RANDOM_UNIFORM(STREAM, NUMBER)
17         IMPORT :: RANDOM_STREAM, C_DOUBLE
18         CLASS(RANDOM_STREAM), INTENT(INOUT) :: STREAM
19         REAL(C_DOUBLE), INTENT(OUT) :: NUMBER
20     END SUBROUTINE RANDOM_UNIFORM
21 END INTERFACE

```

- 2 A polymorphic object with declared type RANDOM\_STREAM is not [interoperable](#) with C. However, we can make such a random number generator available to C by packaging it inside another nonpolymorphic, nonparameterized derived type:

```

25     TYPE :: URNG_STATE ! No BIND(C), as this type is not interoperable
26     CLASS(RANDOM_STREAM), ALLOCATABLE :: STREAM
27 END TYPE URNG_STATE

```

- 3 The following two procedures will enable a C program to use our Fortran uniform random number generator:

```

29     ! Initialize a uniform random number generator:
30     SUBROUTINE INITIALIZE_URNG(STATE_HANDLE, METHOD) &
31         BIND(C, NAME="InitializeURNG")
32     TYPE(C_PTR), INTENT(OUT) :: STATE_HANDLE
33     ! An opaque handle for the URNG
34     CHARACTER(C_CHAR), DIMENSION(*), INTENT(IN) :: METHOD
35     ! The algorithm to be used
36
37     TYPE(URNG_STATE), POINTER :: STATE
38     ! An actual URNG object
39
40     ALLOCATE(STATE)
41     ! There needs to be a corresponding finalization
42     ! procedure to avoid memory leaks, not shown in this example

```

```

1      ! Allocate STATE%STREAM with a dynamic type depending on METHOD
2      ...
3      STATE_HANDLE=C_LOC(STATE)
4      ! Obtain an opaque handle to return to C
5      END SUBROUTINE INITIALIZE_URNG
6
7      ! Generate a random number:
8      SUBROUTINE GENERATE_UNIFORM(STATE_HANDLE, NUMBER) &
9          BIND(C, NAME="GenerateUniform")
10         TYPE(C_PTR), INTENT(IN), VALUE :: STATE_HANDLE
11         ! An opaque handle: Obtained via a call to INITIALIZE_URNG
12         REAL(C_DOUBLE), INTENT(OUT) :: NUMBER
13
14         TYPE(URNG_STATE), POINTER :: STATE
15         ! A pointer to the actual URNG
16
17         CALL C_F_POINTER(CPTR=STATE_HANDLE, FPTR=STATE)
18         ! Convert the opaque handle into a usable pointer
19         CALL STATE%STREAM%NEXT(NUMBER)
20         ! Use the type-bound procedure NEXT to generate NUMBER
21     END SUBROUTINE GENERATE_UNIFORM

```

## C.12.6 Using assumed type to interoperate with C

### C.12.6.1 Overview

1 The mechanism for handling [unlimited polymorphic](#) entities whose [dynamic type](#) is interoperable with C is designed to handle the following two situations:

- (1) A formal parameter that is a C pointer to void. This is an address, and no further information about the entity is provided. The formal parameter corresponds to a dummy argument that is a nonallocatable nonpointer scalar or is an [assumed-size array](#).
- (2) A formal parameter that is the address of a [C descriptor](#). Additional information on the status, type, size, and shape is implicitly provided. The formal parameter corresponds to a dummy argument that is [assumed-shape](#) or [assumed-rank](#).

2 In the first situation, it is the programmer's responsibility to explicitly provide any information needed on the status, type, size, and shape of the entity.

### C.12.6.2 Mapping of interfaces with void \* C parameters to Fortran

1 A C interface for message passing or input/output functionality could be provided in the form

```
int EXAMPLE_send(const void *buffer, size_t buffer_size, const HANDLE_t *handle);
```

where the `buffer_size` argument is given in units of bytes, and the `handle` argument (which is of a type aliased to `int`) provides information about the target the buffer is to be transferred to. In this example, type resolution is not required.

2 The first method provides a thin binding; a call to `EXAMPLE_send` from Fortran directly invokes the C function.

```

1      INTERFACE
2          INTEGER (C_INT) FUNCTION example_send(buffer, buffer_size, handle) &
3          BIND(C, NAME='EXAMPLE_send')
4          USE, INTRINSIC :: ISO_C_BINDING
5          TYPE(*), INTENT (IN) :: buffer(*)
6          INTEGER (C_SIZE_T), VALUE :: buffer_size
7          INTEGER (C_INT), INTENT (IN) :: handle
8      END FUNCTION
9  END INTERFACE

```

10 3 It is assumed that this interface is declared in the specification part of the module MOD\_EXAMPLE\_OLD. An  
 11 example of its use follows:

```

12      USE, INTRINSIC :: ISO_C_BINDING
13      USE MOD_EXAMPLE_OLD
14
15      REAL(C_FLOAT) :: x(100)
16      INTEGER(C_INT) :: y(10,10)
17      REAL(C_DOUBLE) :: z
18      INTEGER(C_INT) :: status, handle
19      ...
20      ! Assign values to x, y, z and initialize handle.
21      ...
22      ! Send values in x, y, and z using EXAMPLE_send.
23      status = example_send(x, C_SIZEOF(x), handle)
24      status = example_send(y, C_SIZEOF(y), handle)
25      status = example_send([ z ], C_SIZEOF(z), handle)

```

26 4 In those invocations, x and y are passed directly with sequence association, but it is necessary to make an array  
 27 expression containing the value of z to pass it.

28 5 The second method provides a Fortran interface which is easier to use, but requires writing a separate C wrapper  
 29 routine. With this method, a [C descriptor](#) is created because the buffer is [assumed-rank](#) in the Fortran interface;  
 30 the use of an optional argument is also demonstrated.

```

31      INTERFACE
32          SUBROUTINE example_send(buffer, handle, status) BIND(C, NAME="EG_send_fortran")
33          USE, INTRINSIC :: ISO_C_BINDING
34          TYPE(*), CONTIGUOUS, INTENT (IN) :: buffer(..)
35          INTEGER (C_INT), INTENT (IN) :: handle
36          INTEGER (C_INT), INTENT(OUT), OPTIONAL :: status
37      END SUBROUTINE
38  END INTERFACE

```

39 6 It is assumed that this interface is declared in the specification part of a module MOD\_EXAMPLE\_NEW.  
 40 Example invocations from Fortran are then

```

1      USE, INTRINSIC :: iso_c_binding
2      USE mod_example_new
3
4      TYPE, BIND(C) :: my_derived
5          INTEGER(C_INT) :: len_used
6          REAL(C_FLOAT) :: stuff(100)
7      END TYPE
8      TYPE(my_derived) :: w(3)
9      REAL(C_FLOAT) :: x(100)
10     INTEGER(C_INT) :: y(10,10)
11     REAL(C_DOUBLE) :: z
12     INTEGER(C_INT) :: status, handle
13     ...
14     ! Assign values to w, x, y, z and initialize handle.
15     ...
16     ! Send values in w, x, y, and z using example_send.
17     CALL example_send(w, handle, status)
18     CALL example_send(x, handle)
19     CALL example_send(y, handle)
20     CALL example_send(z, handle)
21     CALL example_send(y(:,5), handle) ! Fifth column of y.
22     CALL example_send(y(1,5), handle) ! Scalar y(1,5) passed by descriptor.

```

23    7 The wrapper routine can be written in C as follows.

```

24     #include "ISO_Fortran_binding.h"
25
26     void EXAMPLE_send_fortran(const CFI_cdesc_t *buffer, const HANDLE_t *handle,
27                             int *status)
28     {
29         int status_local;
30         size_t buffer_size;
31         int i;
32
33         buffer_size = buffer->elem_len;
34         for (i=0; i<buffer->rank; i++) {
35             buffer_size *= buffer->dim[i].extent;
36         }
37         status_local = EXAMPLE_send(buffer->base_addr,buffer_size, handle);
38         if (status != NULL) *status = status_local;
39     }

```

## 40    C.12.7 Using assumed-type variables in Fortran

41    1 An [assumed-type](#) dummy argument in a Fortran procedure can be used as an actual argument corresponding to an  
42    [assumed-type](#) dummy in a call to another procedure. In the following example, the Fortran subroutine SIMPLE\_  
43    SEND serves as a wrapper to hide the complications associated with calls to a C function named ACTUAL\_Send.

Module COMM\_INFO contains node and address information for the current data transfer operations.

```

SUBROUTINE SIMPLE_SEND(buffer, nbytes)
  USE comm_info, ONLY: my_node, r_node, r_addr
  USE, INTRINSIC :: ISO_C_BINDING
  IMPLICIT NONE

  TYPE(*), INTENT (IN) :: buffer(*)
  INTEGER                :: nbytes, ierr

  INTERFACE
    SUBROUTINE actual_Send(buffer, nbytes, node, addr, ierr) &
      BIND(C, NAME="ACTUAL_Send")
      IMPORT :: C_SIZE_T, C_INT, C_INTPTR_T
      TYPE(*), INTENT (IN)      :: buffer(*)
      INTEGER(C_SIZE_T), VALUE  :: nbytes
      INTEGER(C_INT), VALUE     :: node
      INTEGER(C_INTPTR_T), VALUE :: addr
      INTEGER(C_INT), INTENT(OUT) :: ierr
    END SUBROUTINE actual_Send
  END INTERFACE

  CALL actual_Send(buffer, INT(nbytes, C_SIZE_T), r_node, r_addr, ierr)

  IF (ierr /= 0) THEN
    PRINT *, "Error sending from node", my_node, "to node", r_node
    PRINT *, "Program Aborting"      ! Or call a recovery procedure
    ERROR STOP                       ! Omit in the recovery case
  END IF
END SUBROUTINE simple_Send

```

## C.12.8 Simplifying interfaces for arbitrary rank procedures

- 1 There are situations where an [assumed-rank](#) dummy argument can be useful in Fortran, although a Fortran procedure cannot itself access its value. For example, the IEEE inquiry functions in Clause 14 could be written using an [assumed-rank](#) dummy argument instead of writing 16 separate specific routines, one for each possible rank.
- 2 In particular, the specific procedures for the IEEE\_SUPPORT\_DIVIDE function could possibly be implemented in Fortran as follows:

```

INTERFACE ieee_support_divide
  MODULE PROCEDURE ieee_support_divide_noarg, ieee_support_divide_onearg_r, &
    ieee_support_divide_onearg_d
END INTERFACE ieee_support_divide

...

```

```

1      LOGICAL FUNCTION ieee_support_divide_noarg ()
2          ieee_support_divide_noarg = .TRUE.
3      END FUNCTION ieee_support_divide_noarg
4
5      LOGICAL FUNCTION ieee_support_divide_onearg_r (x)
6          REAL, INTENT (IN) :: x(..)
7          ieee_support_divide_onearg_r4 = .TRUE.
8      END FUNCTION ieee_support_divide_onearg_r
9
10     LOGICAL FUNCTION ieee_support_divide_onearg_d (x)
11         DOUBLE PRECISION, INTENT (IN) :: x(..)
12         ieee_support_divide_onearg_r8 = .TRUE.
13     END FUNCTION ieee_support_divide_onearg_d

```

### 14 C.12.9 Processing assumed-shape arrays in C

- 15 1 The example shown below calculates the product of individual elements of arrays A and B and returns the result  
16 in array C. The Fortran interface of `elemental_mult` will accept arguments of any type and rank. However, the  
17 C function will return an error code if any argument is not a two-dimensional `int` array. Note that the arguments  
18 are permitted to be array sections, so the C function does not assume that any argument is contiguous.

- 19 2 The Fortran interface is:

```

20     INTERFACE
21         FUNCTION elemental_mult(a, b, c) BIND(C, NAME="elemental_mult_c") RESULT(err)
22             USE, INTRINSIC :: ISO_C_BINDING
23             INTEGER(C_INT) :: err
24             TYPE(*), DIMENSION(..) :: a, b, c
25         END FUNCTION elemental_mult
26     END INTERFACE

```

- 27 3 The definition of the C function is:

```

28     #include "ISO_Fortran_binding.h"
29
30     int elemental_mult_c(CFI_cdesc_t * a_desc, CFI_cdesc_t * b_desc, CFI_cdesc_t * c_desc)
31     {
32         size_t i, j, ni, nj;
33
34         int err = 1; /* this error code represents all errors */
35
36         char * a_col = (char*) a_desc->base_addr;
37         char * b_col = (char*) b_desc->base_addr;
38         char * c_col = (char*) c_desc->base_addr;
39         char *a_elt, *b_elt, *c_elt;
40
41         /* Only support int. */
42         if (a_desc->type != CFI_type_int || b_desc->type != CFI_type_int ||

```

```

1         c_desc->type != CFI_type_int) {
2     return err;
3 }
4
5     /* Only support two dimensions. */
6     if (a_desc->rank != 2 || b_desc->rank != 2 || c_desc->rank != 2) {
7         return err;
8     }
9
10    ni = a_desc->dim[0].extent;
11    nj = a_desc->dim[1].extent;
12
13    /* Ensure the shapes conform. */
14    if (ni != b_desc->dim[0].extent || ni != c_desc->dim[0].extent) return err;
15    if (nj != b_desc->dim[1].extent || nj != c_desc->dim[1].extent) return err;
16
17    /* Multiply the elements of the two arrays. */
18    for (j = 0; j < nj; j++) {
19        a_elt = a_col;
20        b_elt = b_col;
21        c_elt = c_col;
22        for (i = 0; i < ni; i++) {
23            *(int*)a_elt = *(int*)b_elt * *(int*)c_elt;
24            a_elt += a_desc->dim[0].sm;
25            b_elt += b_desc->dim[0].sm;
26            c_elt += c_desc->dim[0].sm;
27        }
28        a_col += a_desc->dim[1].sm;
29        b_col += b_desc->dim[1].sm;
30        c_col += c_desc->dim[1].sm;
31    }
32    return 0;
33 }

```

### 34 C.12.10 Creating a contiguous copy of an array

35 1 A C function might need to create a contiguous copy of an array section, for example, to pass the array section  
36 as an actual argument corresponding to a dummy argument with the CONTIGUOUS attribute. The following  
37 example provides functions that can be used to copy an array described by a [CFI\\_cdesc\\_t](#) descriptor to a  
38 contiguous buffer. The input array need not be contiguous.

39 2 The C functions are:

```

40     #include "ISO_Fortran_binding.h"
41     /* Other necessary includes omitted. */
42
43     /*
44     * Returns the number of elements in the object described by desc.

```

```

1      * If it is an array, it need not be contiguous.
2      * (The number of elements could be zero).
3      */
4      size_t numElements(const CFI_cdesc_t * desc)
5      {
6          CFI_rank_t r;
7          size_t num = 1;
8
9          for (r = 0; r < desc->rank; r++) {
10             num *= desc->dim[r].extent;
11         }
12         return num;
13     }
14
15     /*
16     * Auxiliary recursive function to copy an array of a given rank.
17     * Recursion is useful because an array of rank n is composed of an
18     * ordered set of arrays of rank n-1.
19     */
20     static void *_copyToContiguous (const CFI_cdesc_t *vald, void *output,
21                                     const void *input, CFI_rank_t rank)
22     {
23         CFI_index_t e;
24
25         if (rank == 0) {
26             /* Copy scalar element. */
27             memcpy (output, input, vald->elem_len);
28             output = (void *)((char *)output + vald->elem_len);
29         }
30         else {
31             for (e = 0; e < vald->dim[rank-1].extent; e++) {
32                 /* Recurse on subarrays of lesser rank. */
33                 output = _copyToContiguous (vald, output, input, rank-1);
34                 input = (void *) ((char *)input + vald->dim[rank].sm);
35             }
36         }
37         return output;
38     }
39
40     /*
41     * General routine to copy the elements in the array described by vald
42     * to buffer, as done by sequence association. The array itself can
43     * be non-contiguous. This is not the most efficient approach.
44     */
45     void copyToContiguous (void * buffer, const CFI_cdesc_t * vald) {
46         _copyToContiguous (vald, buffer, vald->base_addr, vald->rank);

```

1            }

## 2       **C.12.11   Changing the attributes of an array**

3   1 A C programmer might want to call more than one Fortran procedure and the attributes of an array involved  
 4   might differ between the procedures. In this case, it is necessary to set up more than one **C descriptor** for the  
 5   array. For example, this code fragment initializes the first **C descriptor** for an allocatable entity of rank 2, calls  
 6   a procedure that allocates the array described by the first **C descriptor**, constructs the second **C descriptor** by  
 7   invoking CFI\_section with the value CFI\_attribute\_other for the **attribute** parameter, then calls a procedure  
 8   that expects an **assumed-shape array**.

```

9         CFI_CDESC_T(2) loc_alloc, loc_assum;
10        CFI_cdesc_t * desc_alloc = (CFI_cdesc_t *)&loc_alloc,
11                                     * desc_assum = (CFI_cdesc_t *)&loc_assum;
12        CFI_index_t extents[2];
13        CFI_rank_t rank = 2;
14        int flag;

15
16        flag = CFI_establish(desc_alloc,
17                             NULL,
18                             CFI_attribute_allocatable,
19                             CFI_type_double,
20                             sizeof(double),
21                             rank,
22                             NULL);

23
24        Fortran_factor (desc_alloc, ...); /* Allocates array described by desc_alloc. */

25
26        /* Extract extents from descriptor. */
27        extents[0] = desc_alloc->dim[0].extent;
28        extents[1] = desc_alloc->dim[1].extent;

29
30        flag = CFI_establish(desc_assum,
31                             desc_alloc->base_addr,
32                             CFI_attribute_other,
33                             CFI_type_double,
34                             sizeof(double),
35                             rank,
36                             extents);

37
38        Fortran_solve (desc_assum, ...); /* Uses array allocated in Fortran_factor. */

```

39   2 After invocation of the second CFI\_establish, the lower bounds stored in the dim member of desc\_assum will  
 40   have the value 0 even if the corresponding entries in desc\_alloc have different values.

## C.12.12 Creating an array section in C using CFI\_section

- 1 The C function `set_odd` sets every second element of an array to a specific value, beginning with the first element. It does this by making an array section descriptor for the elements to be set, and calling a Fortran subroutine `SET_ALL` that sets every element of an assumed-shape array to a specific value. An interface block for `set_odd` permits it to be also called from Fortran.

```

SUBROUTINE set_all(int_array, val) BIND(C)
  INTEGER(C_INT) :: int_array(:)
  INTEGER(C_INT), VALUE :: val
  int_array = val
END SUBROUTINE

INTERFACE
  SUBROUTINE set_odd(int_array, val) BIND(C)
    USE, INTRINSIC :: ISO_C_BINDING, ONLY : C_INT
    INTEGER(C_INT) :: int_array(:)
    INTEGER(C_INT), VALUE :: val
  END SUBROUTINE
END INTERFACE

#include "ISO_Fortran_binding.h"

void set_odd(CFI_cdesc_t *int_array, int val)
{
  CFI_index_t lower_bound[1], upper_bound[1], stride[1];
  CFI_CDESC_T(1) array;
  int status;
  /* Create a new descriptor which will contain the section. */
  status = CFI_establish((CFI_cdesc_t *)&array,
                        NULL,
                        CFI_attribute_other,
                        int_array->type,
                        int_array->elem_len,
                        /* rank */ 1,
                        /* extents is ignored */ NULL);

  lower_bound[0] = int_array->dim[0].lower_bound;
  upper_bound[0] = lower_bound[0] + (int_array->dim[0].extent - 1);
  stride[0] = 2;

  status = CFI_section((CFI_cdesc_t *)&array,
                      int_array,
                      lower_bound,
                      upper_bound,
                      stride);

```

```

1      set_all( (CFI_cdesc_t *) &array, val);
2
3      /* Here one could make use of int_array and access all its data. */
4  }
```

5 2 The `set_odd` procedure can be called from Fortran as follows:

```

6      INTEGER(C_INT) :: d(5)
7      d = (/ 1, 2, 3, 4, 5 /)
8      CALL set_odd(d, -1)
9      PRINT *, d
```

10 3 This program will print something like:

```

11      -1    2    -1    4    -1
```

12 4 During execution of the subroutine `SET_ALL`, its dummy argument `INT_ARRAY` would have size (and upper bound) 3.

14 5 It is also possible to invoke `set_odd()` from C. However, it would be the C programmer's responsibility to make sure that all members of the [C descriptor](#) have the correct value on entry to the function. Inserting additional checking into the function could alleviate this problem.

17 6 Following is an example C function that dynamically generates a [C descriptor](#) for an [assumed-shape array](#) and calls `set_odd`.

```

19      #include <stdio.h>
20      #include <stdlib.h>
21      #include "ISO_Fortran_binding.h"
22
23      #define ARRAY_SIZE 5
24
25      void example_of_calling_set_odd(void)
26      {
27          CFI_CDESC_T(1) d;
28          CFI_index_t extent[1];
29          CFI_index_t subscripts[1];
30          void *base;
31          int i, status;
32
33          base = malloc(ARRAY_SIZE*sizeof(int));
34          extent[0] = ARRAY_SIZE;
35          status = CFI_establish((CFI_cdesc_t *)&d,
36                                base,
37                                CFI_attribute_other,
38                                CFI_type_int,
39                                /* element length is ignored */ 0,
40                                /* rank */ 1,
```

```

1          extent);
2
3      set_odd((CFI_cdesc_t *)&d, -1);
4
5      for (i=0; i<ARRAY_SIZE; i++) {
6          subscripts[0] = i;
7          printf(" %d",*((int *)CFI_address((CFI_cdesc_t *)&d, subscripts)));
8      }
9      putc(10, stdout);
10     free(base);
11 }

```

12 The above C function will print similar output to that of the preceding Fortran program.

### 13 C.12.13 Use of CFI\_setpointer

14 1 The C function `change_target` modifies a pointer to an integer variable to become associated with a global  
15 variable defined inside C:

```

16     #include "ISO_Fortran_binding.h"
17
18     int y = 2;
19
20     void change_target(CFI_cdesc_t *ip) {
21         CFI_CDESC_T(0) yp;
22         int status;
23         /* Make local yp point at y. */
24         status = CFI_establish((CFI_cdesc_t *)&yp,
25                               &y,
26                               CFI_attribute_pointer,
27                               CFI_type_int,
28                               /* elem_len is ignored */ sizeof(int),
29                               /* rank */ 0,
30                               /* extents are ignored */ NULL);
31         /* Pointer-associate ip with (the target of) yp. */
32         status = CFI_setpointer(ip, (CFI_cdesc_t *)&yp, NULL);
33         if (status != CFI_SUCCESS) {
34             ... report run time error...
35         }
36     }

```

37 2 The restrictions on the use of `CFI_establish` prohibit direct modification of the incoming pointer entity `ip` by  
38 invoking that function on it.

39 3 The following program illustrates the usage of `change_target` from Fortran.

```

40     PROGRAM change_target_example
41     USE, INTRINSIC :: ISO_C_BINDING

```

```

1      INTERFACE
2          SUBROUTINE change_target(ip) BIND(C)
3              IMPORT :: C_INT
4              INTEGER(C_INT), POINTER :: ip
5          END SUBROUTINE
6      END INTERFACE
7      INTEGER(C_INT), TARGET :: it = 1
8      INTEGER(C_INT), POINTER :: it_ptr
9      it_ptr => it
10     WRITE (*,*) it_ptr
11     CALL change_target(it_ptr)
12     WRITE (*,*) it_ptr

```

13 4 This will print something similar to

```

14     1
15     2

```

## 16 C.12.14 Mapping of MPI interfaces to Fortran

17 1 The Message Passing Interface (MPI) specifies procedures for exchanging data between MPI processes. This  
18 example shows the usage of `MPI_Send` and is similar to the second variant of `EXAMPLE_Send` in C.12.6.2. It also  
19 shows the usage of assumed-length character dummy arguments and optional dummy arguments.

20 2 `MPI_Send` has the C prototype:

```

21     int MPI_Send(void *buf, int count, MPI_Datatype datatype, int dest, int tag,
22                  MPI_Comm comm);

```

23 where `MPI_Datatype` and `MPI_Comm` are opaque handles. Most MPI C functions return an error code, which in  
24 Fortran is the last dummy argument to the corresponding subroutine and can be made optional. Thus, the use  
25 of a Fortran subroutine requires a wrapper function, declared as

```

26     void MPI_Send_f(CFI_cdesc_t *buf, int count, MPI_Datatype_f datatype, int dest,
27                     int tag, MPI_Datatype_f comm, int *ierror);

```

28 3 This wrapper function will convert `MPI_Datatype_f` and `MPI_Comm_f` to `MPI_Datatype` and `MPI_Comm`, and pro-  
29 duce a contiguous `void *` buffer from `CFI_cdesc_t *buf` (if necessary).

30 4 Similarly, the wrapper function for `MPI_Comm_set_name` could have the C prototype:

```

31     void MPI_Comm_set_name_f(MPI_Comm comm, CFI_cdesc_t *comm_name, int *ierror);

```

32 5 The Fortran handle types and interfaces are defined in the module `MPI_F08`. For example,

```

33     MODULE mpi_f08
34     ...
35     TYPE, BIND(C) :: mpi_comm
36     PRIVATE
37     INTEGER(C_INT) :: mpi_val

```

```

1      END TYPE mpi_comm
2
3      INTERFACE
4          SUBROUTINE MPI_SEND(buf,count,datatype,dest,tag,comm,ierror) &
5              BIND(C, NAME='MPI_Send_f')
6              USE, INTRINSIC :: ISO_C_BINDING
7              IMPORT :: MPI_Datatype, MPI_Comm
8              TYPE(*), DIMENSION(..), INTENT (IN) :: buf
9              INTEGER(C_INT), VALUE, INTENT (IN) :: count, dest, tag
10             TYPE(mpi_datatype), INTENT (IN) :: datatype
11             TYPE(mpi_comm), INTENT (IN) :: comm
12             INTEGER(C_INT), OPTIONAL, INTENT (OUT) :: ierror
13         END SUBROUTINE mpi_send
14
15         SUBROUTINE mpi_comm_set_name(comm,comm_name,ierror) &
16             BIND(C, NAME='MPI_Comm_set_name_f')
17             USE, INTRINSIC :: ISO_C_BINDING
18             IMPORT :: mpi_comm
19             TYPE(mpi_comm), INTENT (IN) :: comm
20             CHARACTER(KIND=C_CHAR, LEN=*), INTENT (IN) :: comm_name
21             INTEGER(C_INT), OPTIONAL, INTENT (OUT) :: ierror
22         END SUBROUTINE mpi_comm_set_name
23     END INTERFACE
24     ...
25 END MODULE mpi_f08

```

26 6 Some examples of invocation from Fortran are:

```

27     USE, INTRINSIC :: ISO_C_BINDING
28     USE :: MPI_f08
29
30     TYPE(mpi_comm) :: comm
31     REAL :: x(100)
32     INTEGER :: y(10,10)
33     REAL(KIND(1.0d0)) :: z
34     INTEGER :: dest, tag, ierror
35     ...
36     ! Assign values to x, y, z and initialize MPI variables.
37     ...
38
39     ! Set the name of the communicator.
40     CALL mpi_comm_set_name(comm, "Communicator Name", ierror)
41
42     ! Send values in x, y, and z.
43     CALL mpi_send(x, 100, MPI_REAL, dest, tag, comm, ierror)
44     IF (ierror/=0) PRINT *, 'WARNING: X send error', ierror
45     CALL mpi_send(y(3,:), 10, MPI_INTEGER, dest, tag, comm)

```

```
CALL mpi_send(z, 1, MPI_DOUBLE_PRECISION, dest, tag, comm)
```

- 7 The first example sends the entire array X and includes the optional error argument return value. The second example sends a noncontiguous subarray (the third row of Y) and the third example sends a scalar Z. Note the differences between the calls in this example and those in [C.12.6.2](#).

### C.13 Clause 16 notes : Examples of host association (16.5.1.4)

- 1 The first two examples are examples of valid [host association](#). The third example is an example of invalid [host association](#).

#### Example 1:

```
PROGRAM A
  INTEGER I, J
  ...
CONTAINS
  SUBROUTINE B
    INTEGER I ! Declaration of I hides
              ! program A's declaration of I
    ...
    I = J      ! Use of variable J from program A
              ! through host association
  END SUBROUTINE B
END PROGRAM A
```

#### Example 2:

```
PROGRAM A
  TYPE T
  ...
END TYPE T
...
CONTAINS
  SUBROUTINE B
    IMPLICIT TYPE (T) (C) ! Refers to type T declared below
                        ! in subroutine B, not type T
                        ! declared above in program A
    ...
    TYPE T
    ...
  END TYPE T
  ...
  END SUBROUTINE B
END PROGRAM A
```

#### Example 3:

```
1      PROGRAM Q
2          REAL (KIND = 1) :: C
3          ...
4      CONTAINS
5          SUBROUTINE R
6              REAL (KIND = KIND (C)) :: D ! Invalid declaration
7  ! See below
8              REAL (KIND = 2) :: C
9              ...
10         END SUBROUTINE R
11     END PROGRAM Q
```

12 2 In the declaration of D in subroutine R, the use of C would refer to the declaration of C in subroutine R, not  
13 program Q. However, it is invalid because the declaration of C is required to occur before it is used in the  
14 declaration of D (7.1.12).

## Index

In the index, entries in *italics* denote BNF terms, and page numbers in **bold face** denote primary text or definitions.

## Symbols

- , 147
- <, 151
- <=, 151
- >, 151
- >=, 151
- \*, 49, 52, 54, 59, 94, 98, 108, 130, 147, 234, 235, 265, 277, 282, 307, 326
- \*\*, 147
- +, 147
- stmt*, 18
- .AND., 142, 143, 146, **150**, 150, 348
- .EQ., 141, 143, 146, 150, **151**, 151, 153, 296
- .EQV., 142, 143, 146, **150**, 150
- .FALSE., **62**, 473
- .GE., 141, 143, 146, 150, **151**, 153, 296
- .GT., 141, 143, 146, 150, **151**, 153, 296
- .LE., 141, 143, 146, 150, **151**, 153, 296
- .LT., 141, 143, 146, 150, **151**, 153, 296
- .NE., 141, 143, 146, 150, **151**, 151, 153, 296
- .NEQV., 142, 143, 146, **150**, 150, 409
- .NOT., 142, 143, 146, **150**, 150
- .OR., 142, 143, 146, **150**, 150, 349
- .TRUE., **62**, 473
- /, 147
- / edit descriptor, 275
- //, 149
- /=, 151
- : edit descriptor, 275
- ;;, 48
- ==, 151
- &, 48, 282
- A**
- A edit descriptor, 272
- ABS, **346**, 444
- ABSTRACT, **63**, 63, 79, **296**, 297
- ABSTRACT attribute, 20, 63, **79**
- abstract interface, **2**, 12, 288, **295**, 297, 305, 324, 504, 508
- abstract interface block, **12**, 12, 297
- abstract type, **20**, 53, 76, **79**, 79, 82, 123, 131
- ac-do-variable* (R476), **87**, 87, 156, 158, 506
- ac-implied-do* (R474), **87**, 87, 145, 506
- ac-implied-do-control* (R475), **87**, 87, 145, 156–158, 506
- ac-spec* (R470), 86, **87**
- ac-value* (R473), **87**, 87
- access-id* (R528), **105**, 105
- access-name*, 105
- access-spec* (R507), 63, 68, 69, 74–77, 89, **92**, 92, 105, 106, 298, 304
- access-stmt* (R527), 30, 92, **105**, 105, 106
- ACCESS= specifier, 220, **222**, 248, **249**
- accessibility attribute, **92**, 105, 288
- accessibility statement, 105
- ACHAR, 62, 162, **346**
- ACOS, **346**
- ACOSH, **347**
- ACQUIRED\_LOCK= specifier, **207**, 520, 523
- action, 212
- action-stmt* (R215), 4, **31**, 31, 145, 188, 196
- ACTION= specifier, 221, **222**, 248, **249**, 557
- active image, **11**, 132, 135, 136, 177–179, 206, 337, 361, 362, 404
- actual argument, **2**, 13, 26, 36, 39–41, 54, 56, 66, 78, 84, 97–102, 123, 125, 133, 135, 145, 154, 155, 240, 301, 302, 306–318, 320, 328, 329, 331–334, 338, 344, 345, 366, 391, 407, 416, 429–431, 434, 445, 477, 478, 480, 483–485, 507, 512, 513, 517, 521–523, 533, 576, 577, 579, 581, 582
- actual-arg* (R1225), **307**, 307

- actual-arg-spec* (R1224), 82, **307**, 307
- add-op* (R709), 45, **140**, 140
- add-operand* (R705), **140**, 140, 143, 144
- ADJUSTL, **347**
- ADJUSTR, **347**
- ADVANCE= specifier, 226, 227, **228**, 238, 555
- advancing input/output statement, 215
- AIMAG, **347**
- AINIT, **348**
- ALL, **298**, **348**
- alloc-opt* (R628), **130**, 130, 131
- allocatable, **2**, **2**, 17, 39, 40, 52, 54, 65, 70, 71, 73, 78, 81, 83, 84, 90, 94, 97–100, 104, 108, 116, 117, 123, 124, 130, 131, 133–136, 155, 157, 159–162, 164, 165, 185, 199, 231, 232, 237, 294, 295, 307, 311, 313, 314, 317, 325, 331, 333, 348, 349, 365, 375, 391, 392, 403, 404, 407, 417, 420, 422, 424, 427, 430, 431, 440, 445, 473, 478, 480–482, 487, 494, 497, 510, 512, 521
- ALLOCATABLE attribute, **2**, 52–54, 62, 68, 69, **92**, 92–94, 97, 102, 104, 106, 123, 126, 175, 181, 191, 192, 292, 295, 302, 303, 313, 317, 324, 332, 482, 510, 516, 517, 579, 583
- ALLOCATABLE statement, **106**
- allocatable-decl* (R530), **106**, 106
- allocatable-stmt* (R529), 30, **106**, 508
- ALLOCATE statement, 52, 53, 60, 92, 95, 97, **130**, 136, 137, 164, 198, 434, 491, 492, 511, 512, 519, 521, 526
- allocate-coarray-spec* (R637), **130**, 130, 131
- allocate-coshape-spec* (R638), **130**, 130, 131
- allocate-object* (R633), 60, **130**, 130–134, 136, 137, 198, 434, 436, 522, 523, 526
- allocate-shape-spec* (R634), **130**, 130–133
- allocate-stmt* (R627), 31, **130**, 523
- ALLOCATED, 68, 69, 133, 137, **348**
- allocation* (R632), **130**, 130–133
- alphanumeric-character* (R301), **43**, 43, 44
- alt-return-spec* (R1226), 4, 196, **307**, 307
- ancestor component, **80**
- ancestor-module-name*, 291
- and-op* (R719), 45, **142**, 142
- and-operand* (R714), **142**, 142
- ANINT, **349**
- ANY, **349**
- arg-name*, 69, 71, 76
- argument
  - dummy, 312
- argument association, **3**, **3**, 21, 52, 60, 70, 71, 94, 97, 104, 105, 135, 136, 293, 309, 310, 320, 326, 507, 513, 516, 517, 532, 577
- argument keyword, **9**, **13**, 41, 295, 298, 309, 333, 338, 445, 504, **505**, **506**, 565
- arithmetic IF statement, **532**
- array, **2**, 4, 10, 17, 40, 96–98, 125–128
  - assumed-shape, **2**, 54, 95, 97, 129, 295, 299, 300, 311–313, 315, 318, 324, 473, 483, 485, 576, 594, 601, 603
  - assumed-size, **2**, 98, 99, 104, 115, 125, 126, 139, 156, 159, 181, 192, 231, 311, 312, 317, 391, 420, 422, 423, 430, 478, 482, 487, 494, 590, 592, 594
  - deferred-shape, **2**, 97
  - explicit-shape, **2**, 54, 70, 94, 97, 159, 312, 317, 482
- array bound, **4**, 69, 72, 91
- array constructor, **86**, **86**
- array element, **2**, 39, 126
- array element order, 126–127
- array pointer, **2**, 95, 97, 98, 155, 351, 482
- array section, **2**, 95, 107, 108, 124, 126–129, 175, 217, 311, 312, 318, 511, 513
- array-constructor* (R469), **86**, 87, 139
- array-element* (R617), 107, 108, 115, 121, 122, **125**
- array-name*, 109, 508
- array-section* (R618), **2**, 121, **125**, 126, 127
- array-spec* (R515), 24, 89–91, **96**, 96–99, 106, 109, 111, 117, 118
- ASCII character, **3**, **59**, **62**, 159, 217, 218, 232, 264, 278, 346, 361, 382, 385, 393, 394, 405, 418
- ASCII collating sequence, **62**, 346, 361, 382, 385, 393, 394, 405
- ASIN, **350**
- ASINH, **350**
- ASSIGN statement, **531**
- assigned format, **531**
- assigned GO TO statement, **531**
- ASSIGNMENT, 75, 163, **296**, 296, 302, 303
- assignment, 159–171
  - defined, 75, 163, 302
  - elemental, **9**, 163
  - elemental array (FORALL), 169
  - masked array (WHERE), 167

- pointer, **163**
- assignment statement, **14, 15, 38, 52, 78, 159, 171, 469, 518**
- assignment-stmt* (R732), **31, 159, 159, 167, 170, 522**
- ASSOCIATE construct, **174, 317, 506, 507, 510, 522**
- associate name, **3, 3, 21, 52, 54, 82, 135, 174, 175, 194, 507, 510, 511, 516, 522**
- ASSOCIATE statement, **174, 510**
- associate-construct* (R802), **31, 174, 174**
- associate-construct-name*, **174**
- associate-name*, **174, 191–195, 506**
- associate-stmt* (R803), **4, 174, 174, 196**
- ASSOCIATED, **68, 69, 134, 137, 334, 350**
- associating entity, **3, 60, 174, 175, 194, 195, 326, 516, 516, 517**
- association, **3**
  - argument, **3, 3, 21, 52, 60, 70, 71, 94, 97, 104, 105, 135, 136, 293, 309, 310, 320, 326, 507, 513, 516, 517, 532, 577**
  - common, **118**
  - construct, **3, 3, 135, 136, 507, 510, 513, 516**
  - equivalence, **116**
  - host, **3, 3, 33, 40, 53, 60, 92, 105, 107, 112, 118, 156, 157, 165, 291, 293, 317, 329–331, 506–510, 513, 516, 607**
  - inheritance, **3, 3, 6, 80, 82, 513, 516**
  - linkage, **3, 3, 499, 507, 510, 510**
  - name, **3, 3, 507, 513**
  - pointer, **3, 3, 9, 19, 21, 38, 79, 81, 83, 95, 100, 102, 104, 105, 123, 135, 137, 163, 165, 166, 185, 199, 203, 234, 294, 310, 312, 314, 316, 317, 325, 326, 339, 350, 407, 476, 477, 511–514, 516, 517, 523**
  - sequence, **316**
  - storage, **3, 3, 40, 41, 115–117, 327, 330, 425, 513–516**
  - use, **3, 3, 33, 41, 60, 80, 92, 103, 105, 112, 115–117, 157, 165, 288, 287–291, 297, 327, 330, 506–508, 511**
- association* (R804), **174, 174**
- assumed type parameter, **21, 21, 52, 311, 313, 314**
- assumed-implied-spec* (R521), **98, 98, 99**
- assumed-rank dummy data object, **4, 54, 95, 96, 129, 191, 192, 294, 295, 303, 311–313, 318, 324, 390, 391, 414, 420, 422, 423, 430, 473, 483, 485, 594, 595, 597**
- assumed-rank-spec* (R525), **96, 99**
- assumed-shape array, **2, 54, 95, 97, 129, 295, 299, 300, 311–313, 315, 318, 324, 473, 483, 485, 576, 594, 601, 603**
- assumed-shape-spec* (R519), **96, 97, 97**
- assumed-size array, **2, 98, 99, 104, 115, 125, 126, 139, 156, 159, 181, 192, 231, 311, 312, 317, 391, 420, 422, 423, 430, 478, 482, 487, 494, 590, 592, 594**
- assumed-size-spec* (R522), **96, 98, 98**
- assumed-type, **4, 53, 54, 311, 324, 483, 596**
- ASYNCHRONOUS attribute, **92, 92, 93, 106, 175, 185, 229, 288, 290, 294, 295, 312, 313, 363, 498, 501, 508, 509**
- asynchronous communication, **92, 501**
- asynchronous input/output, **92, 220, 222, 224, 228–230, 233, 234, 241, 244, 245, 247, 249, 251, 252**
- ASYNCHRONOUS statement, **106, 176, 292, 506, 509**
- asynchronous-stmt* (R531), **31, 106**
- ASYNCHRONOUS= specifier, **221, 222, 226, 227, 228, 248, 249**
- ATAN, **351**
- ATAN2, **27, 351**
- ATANH, **352**
- atomic subroutine, **19, 36, 199, 200, 333, 336, 338, 352–356, 373, 433, 437, 522**
- ATOMIC\_ADD, **352, 434**
- ATOMIC\_AND, **352**
- ATOMIC\_CAS, **353**
- ATOMIC\_DEFINE, **353, 583, 584**
- ATOMIC\_FETCH\_ADD, **354**
- ATOMIC\_FETCH\_AND, **354**
- ATOMIC\_FETCH\_OR, **354**
- ATOMIC\_FETCH\_XOR, **355**
- ATOMIC\_INT\_KIND, **352–356, 433**
- ATOMIC\_LOGICAL\_KIND, **353, 356, 433**
- ATOMIC\_OR, **355**
- ATOMIC\_REF, **356, 584**
- ATOMIC\_XOR, **356**
- attr-spec* (R502), **89, 89, 91, 111**
- attribute, **4, 53, 62, 66, 89, 91–105, 290**
  - ABSTRACT, **20, 63, 79**
  - accessibility, **92, 105, 288**
  - ALLOCATABLE, **2, 52–54, 62, 68, 69, 92, 92–94, 97, 102, 104, 106, 123, 126, 175, 181, 191, 192, 292, 295, 302, 303, 313, 317, 324, 332, 482, 510, 516, 517, 579, 583**

- ASYNCHRONOUS, [92](#), [92](#), [93](#), [106](#), [175](#), [185](#), [229](#), [288](#), [290](#), [294](#), [295](#), [312](#), [313](#), [363](#), [498](#), [501](#), [508](#), [509](#)
- BIND, [3](#), [4](#), [38](#), [63–65](#), [79](#), [84](#), [93](#), [93](#), [104](#), [106](#), [116](#), [117](#), [164](#), [165](#), [185](#), [194](#), [292](#), [294](#), [295](#), [323](#), [325](#), [436](#), [437](#), [480](#), [481](#), [483](#), [497–500](#), [510](#), [517](#), [590](#)
- CODIMENSION, [54](#), [69](#), [90](#), [93](#), [93](#), [99](#), [107](#)
- CONTIGUOUS, [69](#), [71](#), [95](#), [95](#), [96](#), [107](#), [128](#), [129](#), [165](#), [185](#), [294](#), [311–315](#), [318](#), [483](#), [485](#), [515](#)
- DEFERRED, [75](#), [76](#), [79](#)
- DIMENSION, [69](#), [90](#), [96](#), [96](#), [109](#), [118](#)
- EXTENDS, [20](#), [79](#), [79](#), [481](#)
- EXTERNAL, [14](#), [25](#), [26](#), [99](#), [99](#), [102](#), [111](#), [112](#), [114](#), [165](#), [288](#), [292](#), [294](#), [297](#), [304](#), [316](#), [321](#), [322](#), [508](#), [509](#), [573](#)
- INTENT, [99](#), [99–101](#), [110](#), [185](#)
- INTENT (IN), [99](#), [100](#), [101](#), [105](#), [181](#), [301–303](#), [311](#), [314](#), [316](#), [318](#), [330](#), [331](#), [334](#), [352–356](#), [362–364](#), [373](#), [374](#), [379](#), [380](#), [404](#), [412](#), [447](#), [475–477](#), [497](#), [576](#), [592](#)
- INTENT (INOUT), [99](#), [100](#), [101](#), [104](#), [302](#), [312](#), [314](#), [320](#), [331](#), [332](#), [352–356](#), [362–364](#), [374](#), [379](#), [380](#), [403](#), [404](#), [434](#), [436](#), [522](#), [523](#), [592](#)
- INTENT (OUT), [26](#), [54](#), [77](#), [78](#), [98](#), [99](#), [99–101](#), [104](#), [135](#), [156](#), [302](#), [312](#), [314](#), [320](#), [330–332](#), [352–356](#), [362–364](#), [366](#), [368](#), [373](#), [374](#), [378–380](#), [403](#), [404](#), [412](#), [426](#), [448–450](#), [475](#), [477](#), [497](#), [498](#), [512](#), [513](#), [518](#), [519](#), [521–523](#), [592](#)
- INTRINSIC, [99](#), [101](#), [101](#), [102](#), [288](#), [306](#), [321](#), [509](#)
- NON\_OVERRIDABLE, [75](#), [76](#)
- OPTIONAL, [102](#), [102](#), [104](#), [110](#), [156](#), [175](#), [181](#), [295](#)
- PARAMETER, [7](#), [38](#), [85](#), [91](#), [102](#), [102](#), [110](#), [122](#)
- PASS, [69](#), [71](#), [76](#), [307](#)
- POINTER, [2](#), [14](#), [52–54](#), [62](#), [68–70](#), [90](#), [92](#), [97](#), [99](#), [102](#), [102](#), [104](#), [109](#), [111](#), [123](#), [126](#), [134](#), [164](#), [175](#), [185](#), [191](#), [192](#), [293–295](#), [297](#), [302](#), [303](#), [305](#), [313](#), [316–319](#), [324](#), [330–332](#), [478](#), [482](#), [497](#), [511](#), [513](#), [516](#), [517](#), [579](#), [583](#)
- PRIVATE, [65](#), [80](#), [92](#), [92](#), [105](#), [115](#), [331](#), [563](#)
- PROTECTED, [103](#), [103](#), [111](#), [116](#), [185](#), [289](#)
- PUBLIC, [80](#), [92](#), [92](#), [105](#), [115](#), [563](#)
- SAVE, [16](#), [22](#), [28](#), [72](#), [79](#), [91](#), [93](#), [94](#), [103](#), [103](#), [104](#), [107](#), [111](#), [116](#), [118](#), [136](#), [185](#), [305](#), [330](#), [513](#)
- SEQUENCE, [17](#), [63](#), [64](#), [65](#), [66](#), [79](#), [117](#), [164](#), [165](#), [194](#), [480](#)
- TARGET, [3](#), [19](#), [72](#), [102](#), [104](#), [104](#), [111](#), [116](#), [118](#), [134](#), [135](#), [164](#), [175](#), [185](#), [192](#), [295](#), [302](#), [311](#), [312](#), [314](#), [318](#), [319](#), [363](#), [403](#), [475](#), [478](#), [498](#), [512](#), [513](#), [521](#), [576](#), [577](#)
- VALUE, [54](#), [71](#), [77](#), [99](#), [104](#), [104](#), [112](#), [185](#), [234](#), [294](#), [295](#), [297](#), [301](#), [302](#), [310–312](#), [314](#), [330](#), [332](#), [363](#), [483](#), [484](#), [501](#), [513](#), [590](#), [592](#)
- VOLATILE, [104](#), [104](#), [105](#), [112](#), [164](#), [165](#), [175](#), [185](#), [288](#), [290](#), [294](#), [295](#), [312](#), [313](#), [315](#), [508](#), [509](#), [514](#), [519](#), [522](#), [542](#)
- attribute specification statements, [105–119](#)
- automatic data object, [4](#), [27](#), [90](#), [91](#), [94](#), [104](#), [108](#), [116](#), [117](#), [519](#), [533](#)
- ## B
- B edit descriptor, [271](#)
- BACKSPACE statement, [212](#), [215](#), [241](#), [244](#), [245](#), [246](#), [555](#), [557](#)
- backspace-stmt* (R924), [31](#), [245](#), [331](#)
- base object, [4](#), [92](#), [95](#), [115](#), [123](#), [129](#), [156](#), [229](#), [317](#), [330](#), [332](#)
- BESSEL\_J0, [356](#)
- BESSEL\_J1, [357](#)
- BESSEL\_JN, [357](#)
- BESSEL\_Y0, [357](#)
- BESSEL\_Y1, [358](#)
- BESSEL\_YN, [358](#)
- BGE, [358](#)
- BGT, [359](#)
- binary-constant* (R465), [86](#), [86](#)
- BIND (C), *see* BIND attribute
- BIND attribute, [3](#), [4](#), [38](#), [63–65](#), [79](#), [84](#), [93](#), [93](#), [104](#), [106](#), [116](#), [117](#), [164](#), [165](#), [185](#), [194](#), [292](#), [294](#), [295](#), [323](#), [325](#), [436](#), [437](#), [480](#), [481](#), [483](#), [497–500](#), [510](#), [517](#), [590](#)
- BIND statement, [106](#), [292](#), [499](#), [505](#)
- bind-entity* (R533), [106](#), [106](#)
- bind-stmt* (R532), [31](#), [106](#)
- binding, [4](#), [75](#), [76](#), [76](#), [79–81](#), [152](#), [163](#), [238](#), [243](#), [303](#), [323](#), [504](#), [505](#)
- binding label, [4](#), [93](#), [295](#), [305](#), [324](#), [326](#), [499–501](#), [503](#), [504](#)
- binding name, [4](#), [75](#), [76](#), [80](#), [307](#), [505](#)
- binding-attr* (R452), [75](#), [76](#), [76](#)
- binding-name*, [75](#), [76](#), [307](#), [323](#), [505](#)
- binding-private-stmt* (R447), [75](#), [75](#), [77](#)
- bit model, [335](#)
- BIT\_SIZE, [335](#), [359](#), [404](#)

- blank common, [6](#), [90](#), [107](#), [117](#), [118](#), [513](#), [515](#)
- blank interpretation mode, [222](#)
- blank-interp-edit-desc* (R1018), [261](#), [262](#)
- BLANK= specifier, [221](#), [222](#), [226](#), [227](#), [229](#), [241](#), [248](#), [249](#), [276](#)
- BLE, [359](#)
- block, [4](#)
  - interface, [289](#)
- block* (R801), [4](#), [173](#), [174](#), [175](#), [177](#), [179](#), [181–183](#), [186](#), [187](#), [189](#), [191](#), [193](#)
- BLOCK construct, [16](#), [35](#), [38](#), [40](#), [78](#), [91](#), [92](#), [95](#), [97](#), [103](#), [105](#), [112](#), [114](#), [135](#), [156](#), [175](#), [330](#), [506](#), [512](#), [513](#), [519](#), [522](#), [526](#)
- block data program unit, [291](#)
- BLOCK DATA statement, [47](#), [287](#), [291](#)
- block scoping unit, [12](#), [17](#)
- BLOCK statement, [40](#), [91](#), [95](#), [97](#), [175](#), [519](#)
- block-construct* (R807), [31](#), [175](#), [176](#)
- block-construct-name*, [175](#), [176](#)
- block-data* (R1120), [29](#), [291](#), [292](#), [298](#)
- block-data-name*, [291](#)
- block-data-stmt* (R1121), [29](#), [291](#), [291](#)
- block-specification-part* (R809), [18](#), [175](#), [176](#), [176](#)
- block-stmt* (R808), [4](#), [175](#), [175](#), [176](#), [196](#)
- BLT, [360](#)
- BN edit descriptor, [276](#)
- bound, [2](#), [4](#), [4](#), [39](#), [40](#), [68](#), [69](#), [81](#), [84](#), [97](#), [130](#), [132](#), [136](#), [137](#), [165](#), [403](#), [507](#)
- bounds, [96–99](#), [125–128](#)
- bounds-remapping* (R736), [163](#), [164](#), [164–166](#)
- bounds-spec* (R735), [163](#), [164](#), [164](#), [166](#)
- boz-literal-constant* (R464), [45](#), [86](#), [86](#), [109](#), [271](#), [335](#), [358–361](#), [369–371](#), [383](#), [386](#), [388](#), [400](#), [414](#)
- branch, [196](#), [328](#), [531](#)
- branch target statement, [4](#), [34](#), [46](#), [168](#), [184](#), [196](#), [196](#), [221](#), [225](#), [226](#), [244](#), [245](#), [247](#), [249](#), [308](#), [328](#)
- BTEST, [360](#)
- BZ edit descriptor, [276](#)
- C**
- C address, [5](#), [475–478](#), [480](#), [481](#), [486](#), [490](#), [491](#), [493](#), [520–522](#), [592](#)
- C descriptor, [5](#), [135](#), [484–498](#)
- C\_ALERT, [474](#)
- C\_ASSOCIATED, [474](#)
- C\_BACKSPACE, [474](#)
- C\_BOOL, [473](#), [474](#)
- C\_CARRIAGE\_RETURN, [474](#)
- C\_CHAR, [474](#)
- C\_DOUBLE, [473](#)
- C\_DOUBLE\_COMPLEX, [474](#)
- C\_F\_POINTER, [474](#), [475](#)
- C\_F\_PROCPTR, [477](#)
- C\_FLOAT, [473](#)
- C\_FLOAT\_COMPLEX, [474](#)
- C\_FORM\_FEED, [474](#)
- C\_FUNLOC, [477](#), [477](#), [500](#)
- C\_FUNPTR, [68](#), [79](#), [93](#), [123](#), [131](#), [132](#), [162](#), [473](#), [474](#), [477](#), [480](#), [481](#), [522](#)
- C\_HORIZONTAL\_TAB, [474](#)
- C\_INT, [473](#)
- C\_INT16\_T, [473](#)
- C\_INT32\_T, [473](#)
- C\_INT64\_T, [473](#)
- C\_INT8\_T, [473](#)
- C\_INT\_FAST16\_T, [473](#)
- C\_INT\_FAST32\_T, [473](#)
- C\_INT\_FAST64\_T, [473](#)
- C\_INT\_FAST8\_T, [473](#)
- C\_INT\_LEAST16\_T, [473](#)
- C\_INT\_LEAST32\_T, [473](#)
- C\_INT\_LEAST64\_T, [473](#)
- C\_INT\_LEAST8\_T, [473](#)
- C\_INTMAX\_T, [473](#)
- C\_INTPTR\_T, [473](#)
- C\_LOC, [54](#), [99](#), [331](#), [478](#)
- C\_LONG, [473](#)
- C\_LONG\_DOUBLE, [473](#)
- C\_LONG\_DOUBLE\_COMPLEX, [474](#)
- C\_LONG\_LONG, [473](#)
- C\_NEW\_LINE, [474](#)
- C\_NULL\_CHAR, [474](#)
- C\_NULL\_FUNPTR, [473](#), [474](#)
- C\_NULL\_PTR, [473](#), [474](#)
- C\_PTR, [68](#), [79](#), [93](#), [123](#), [131](#), [132](#), [162](#), [473–475](#), [478](#), [480](#), [481](#), [484](#), [520–522](#), [590](#)
- C\_SHORT, [473](#)
- C\_SIGNED\_CHAR, [473](#)
- C\_SIZE\_T, [473](#)
- C\_SIZEOF, [156](#), [478](#)
- C\_VERTICAL\_TAB, [474](#)
- CALL statement, [19](#), [196](#), [198](#), [293](#), [307](#), [320](#), [328](#), [404](#)
- call-stmt* (R1222), [31](#), [307](#), [308](#), [310](#)

- CASE statement, **189**
- case-construct* (R840), **31**, **188**, **189**
- case-construct-name*, **189**
- case-expr* (R844), **189**, **189**
- case-selector* (R845), **189**, **189**
- case-stmt* (R842), **188**, **189**, **189**
- case-value* (R847), **189**, **189**
- case-value-range* (R846), **189**, **189**
- CEILING, **360**
- CFI\_address, **490**
- CFI\_allocate, **491**, **497**
- CFI\_cdesc\_t, **5**, **483–485**, **486**, **486–488**, **490–496**, **498**
- CFI\_deallocate, **488**, **492**, **497**
- CFI\_establish, **492**
- CFI\_is\_contiguous, **494**
- CFI\_section, **494**
- CFI\_select\_part, **495**
- CFI\_setpointer, **496**
- CHANGE TEAM construct, **19**, **40**, **177**, **184**, **196**, **506**, **507**
- CHANGE TEAM statement, **19**, **35**, **40**, **177**, **198**, **337**, **437**, **510**, **522**
- change-team-construct* (R811), **31**, **177**, **177**
- change-team-stmt* (R812), **177**, **177**
- changeable mode, **218**
- CHAR, **61**, **361**
- char-length* (R423), **59**, **59**, **60**, **68**, **69**, **89–91**, **534**
- char-literal-constant* (R424), **45**, **49**, **50**, **60**, **240**, **261**, **262**, **525**
- char-selector* (R421), **55**, **59**, **60**
- char-string-edit-desc* (R1021), **260**, **262**
- char-variable* (R605), **121**, **121**, **217**, **218**
- character context, **5**, **43**, **47–49**, **61**
- character literal constant, **60**
- character sequence type, **17**, **65**, **116**, **118**, **515**, **516**, **519**
- character set, **43**
- character storage unit, **18**, **18**, **98**, **116–118**, **433**, **514**, **519**, **520**
- character string edit descriptor, **260**, **277**
- character type, **59–62**
- CHARACTER\_KINDS, **433**
- CHARACTER\_STORAGE\_SIZE, **433**
- characteristics, **5**, **80**, **166**, **238**, **239**, **294**, **295**, **297**, **305**, **306**, **316**, **320**, **324**, **325**, **327**, **346**, **407**
- dummy argument, **294**
- procedure, **294**
- child data transfer statement, **216**, **217**, **228–230**, **233**, **237**, **237–241**, **256**, **280**
- CLASS, **53**, **53**, **54**, **238**
- CLASS DEFAULT statement, **193**
- CLASS IS statement, **193**
- CLOSE statement, **212**, **213**, **217**, **219**, **220**, **224**, **224**, **241**, **244**, **555**
- close-spec* (R909), **225**, **225**
- close-stmt* (R908), **31**, **225**, **331**
- CMPLX, **161**, **335**, **361**, **441**
- CO\_BROADCAST, **361**
- CO\_MAX, **362**
- CO\_MIN, **362**
- CO\_REDUCE, **363**, **529**
- CO\_SUM, **364**, **529**
- coarray, **5**, **5**, **7**, **35**, **39**, **40**, **64**, **68**, **70**, **78**, **93–95**, **100**, **102**, **104**, **105**, **116**, **117**, **130–133**, **135–137**, **159**, **162**, **164**, **165**, **175**, **181**, **198–200**, **204**, **288**, **295**, **307**, **315**, **337**, **341**, **344**, **352–356**, **365**, **386**, **391**, **403**, **427**, **428**, **431**, **436**, **482**
- established, **5**
- coarray-association* (R813), **40**, **177**, **177**, **178**
- coarray-name*, **107**, **177**, **506**, **508**
- coarray-spec* (R509), **68–70**, **89**, **90**, **93**, **93**, **94**, **106**, **107**, **111**
- cobound, **5**, **39**, **40**, **93–95**, **129**, **132**, **175**, **315**, **341**, **344**, **391**, **392**, **403**, **431**, **507**
- codimension, **5**, **5**, **7**, **39**, **95**, **129**, **175**, **294**, **365**, **366**, **392**, **431**
- CODIMENSION attribute, **54**, **69**, **90**, **93**, **93**, **99**, **107**
- codimension-decl* (R535), **107**, **107**, **177**, **178**, **506**
- codimension-stmt* (R534), **31**, **107**, **508**
- coindexed object, **5**, **39**, **40**, **107**, **123**, **129**, **159**, **162**, **164**, **165**, **174**, **205**, **307**, **308**, **311–314**, **330**, **352–356**, **373**, **475**, **477**, **478**
- coindexed-named-object* (R614), **121**, **122**, **124**, **124**
- collating sequence, **6**, **61**, **62**, **151**, **264**, **346**, **361**, **382**, **385**, **393**, **394**, **397–402**, **405**, **525**
- collective subroutine, **19**, **333**, **337**, **338**, **361–364**, **376**, **424**, **437**
- COMMAND\_ARGUMENT\_COUNT, **158**, **364**, **379**
- comment, **48**, **49**, **284**
- common association, **118**
- common block, **6**, **28**, **32**, **38**, **90**, **91**, **93**, **103**, **104**, **106**, **107**, **116–119**, **156**, **291**, **292**, **498**, **499**, **503–506**, **510**, **513–515**, **520**, **534**

- common block storage sequence, **118**
- COMMON statement, **6**, **117**, 117–119, 176, 290, 292, 505, 515, 516, 532
- common-block-name*, 106, 111, 117, 176, 290
- common-block-object* (R573), **117**, 117, 290, 508
- common-stmt* (R572), 31, **117**, 508
- companion processor, 4, 5, **6**, 13, 36, 41, 42, 63, 85, 93, 473, 478, 499–501, 525
- compatibility
  - FORTRAN 77, 27
  - Fortran 2003, 26
  - Fortran 2008, 26
  - Fortran 90, 27
  - Fortran 95, 27
- COMPILER\_OPTIONS, 156, **433**
- COMPILER\_VERSION, 156, **433**
- completion step, 36, 225
- complex part designator, **8**, 37, 124
- complex type, 58
- complex-literal-constant* (R418), 45, **58**
- complex-part-designator* (R615), 121, **124**, 124, 125, 129
- component, **6**, 8, 12, 13, 17, 19, 63–65, **68**, 112, 505
  - direct, **6**, 6, 62, 63, 72, 312, 440, 480
  - parent, **3**, **6**, 74, 78, 80, 83, 516, 543
  - potential subobject, **6**, 62–64, 331
  - ultimate, **6**, 62–64, 68, 93, 95, 98, 99, 104, 116, 117, 131, 133, 158, 237, 311, 514
- component definition statement, **68**
- component keyword, **13**, 41, 74, 82, 83, 505
- component order, **6**, 74, 83, 232
- component-array-spec* (R440), **68**, 68, 69
- component-attr-spec* (R438), **68**, 68–70, 72
- component-data-source* (R458), **82**, 82, 83
- component-decl* (R439), 60, **68**, 68–70, 72
- component-def-stmt* (R436), 6, **68**, 68, 69
- component-initialization* (R443), 68, **71**, 71, 72
- component-name*, 68, 72
- component-part* (R435), 63, **68**, 74, 77
- component-spec* (R457), **82**, 82, 157
- computed GO TO statement, 4, **196**, 196, 532, 533
- computed-goto-stmt* (R858), 32, **196**, 196
- concat-op* (R711), 45, **141**, 141
- CONCURRENT, **181**
- concurrent-control* (R826), 169, 171, **181**, 181, 183
- concurrent-header* (R825), 170, 171, **181**, 181, 182, 506
- concurrent-limit* (R827), 145, 171, **181**, 181–183
- concurrent-locality* (R829), **181**, 181
- concurrent-step* (R828), 145, 171, **181**, 181–183
- conformable, **6**, 39, 133, 145, 152, 159, 163, 320, 332, 377, 383, 384, 389, 398, 399, 401, 402, 409, 411, 425, 431, 460, 461
- CONJG, **365**
- connect-spec* (R905), **220**, 220, 221
- connected, **6**, 10, 13, 15, 212–216, 219–221, 223–225, 230, 234, 236, 237
- connection mode, 218
- constant, **7**, 38, 45, 51
  - integer, 55
  - named, 110
- constant* (R304), **45**, 45, 108, 122, 139
- constant expression, 4, **7**, 21, 27, 51, 52, 60, 67, 69, 70, 72, 87, 91, 95, 97, 98, 107, 108, 110, 116, 156, **157**, **158**, 158, 228, 294, 295, 317, 334, 346, 348, 349, 360, 361, 366, 376–378, 382, 385, 387, 388, 391–393, 395, 396, 398, 401, 406, 414, 416, 418, 420, 422, 424, 430–432, 482
- constant-expr* (R729), 21, 52, 71, 72, 90, 91, 98, 99, 102, 110, **158**, 158, 189
- constant-subobject* (R547), **108**, 108
- construct
  - ASSOCIATE, 173, **174**, 191, 317, 506, 507, 510, 522
  - BLOCK, xix, 12, 16–18, 35, 38, 40, 78, 91, 92, 95, 97, 103, 105, 112, 114, 135, 156, 173, **175**, 330, 506, 512, 513, 519, 522, 526
  - CHANGE TEAM, xix, 19, 40, 173, **177**, 184, 196, 506, 507
  - CRITICAL, 173, **179**, 179, 184, 196, 337
  - DO, 35, 46, 87, 173, **180**, 196, 232, 531, 545, 546
  - DO CONCURRENT, xviii, **180**, 184, 187, 196, 331, 506, 507, 513, 519, 521, 526, 534
  - FORALL, **169**, 331, 506, 507, 519, 532, 534
  - IF, 35, 173, **187**, 447, 531
  - nonblock DO, xviii, **532**
  - SELECT CASE, 35, 173, **188**, 532, 533, 545
  - SELECT RANK, xviii, 35, 98, 99, 173, **191**, 317, 522
  - SELECT TYPE, 35, 52, 53, 173, **193**, 317, 506, 507, 510, 522
  - WHERE, 14, **167**
- construct association, **3**, 3, 135, 136, 507, **510**, 513, 516
- construct entity, 3, **7**, 105, 174, 176, 185, 193, 503, 504,

506, 513  
*construct-name*, 196  
 constructor  
   array, 86  
   derived-type, 82  
   structure, 82  
 CONTAINS statement, 34, **75**, **328**  
*contains-stmt* (R1244), 30, **75**, 288, **328**  
 contiguous, **7**, 17, 65, 71, **95**, 122, 129, 165, 167, 175, 229, 236, 390, 478, 514  
 CONTIGUOUS attribute, 69, 71, **95**, 95, 96, 107, 128, 129, 165, 185, 294, 311–315, 318, 483, 485, 515  
 CONTIGUOUS statement, **107**  
*contiguous-stmt* (R536), 31, **107**  
 continuation, 48, 49  
 CONTINUE statement, **197**, 531  
*continue-stmt* (R859), 31, 182, **197**  
 control character, 43, 60, 211, 214  
 control edit descriptor, **260**, 274–277  
 control information list, **226**  
 control mask, 168  
*control-edit-desc* (R1013), 260, **261**  
 conversion  
   numeric, 161  
 corank, **7**, 39, 40, 70, 93, 94, 96, 123, 129, 131, 139, 175, 294, 314, 366, 386, 391, 392, 403, 427, 428, 431, 507  
 COS, **365**  
*coselector-name*, 3, 177, 522  
 COSH, **365**  
 COSHAPE, **365**  
 cosubscript, **7**, 39, 40, 95, 129, 341, 343, 386, 427, 428, 431  
*cosubscript* (R625), 123, **129**, 129  
 COUNT, 334, **366**  
 CPU\_TIME, **366**  
 CRITICAL construct, **179**, 184, 196  
 CRITICAL statement, 153, **179**, 198  
*critical-construct* (R816), 31, **179**, 179  
*critical-construct-name*, 179  
*critical-stmt* (R817), 4, **179**, 179, 196  
 CSHIFT, **367**  
 current record, **215**  
 current team, **19**, 78, 129, 132, 135, 136, 177, 178, 200, 201, 204, 206, 337, 338, 361, 362, 376, 381, 386, 387, 404, 408, 424, 427, 428, 431, 434

CURRENT\_TEAM, 381, **434**  
 CYCLE statement, 174, 180, **184**, 184, 534  
*cycle-stmt* (R833), 31, 177, **184**, 184  
**D**  
*d* (R1010), **261**, 261, 266–270, 273, 274, 280  
 D edit descriptor, 267  
 data edit descriptor, **260**, 264–274  
 data entity, 5, 6, **7**, 14, 16, 22, 37–39  
 data object, 4, 5, **7**, 7–9, 16–19, 21, 32–34, 37, 39, 40  
 data object designator, **9**, 16, 39, 121  
 data object reference, **16**, 38–40  
 data pointer, **14**, **14**, 40, 475, 487, 515  
 DATA statement, 28, 34, 86, 91, **107**, 118, 292, 407, 506, 509, 517, 532, 533  
 data transfer, 235  
 data transfer input statement, **225**  
 data transfer output statement, **225**  
 data transfer statement, 28, 46, 211–217, 219, **225**, 230, 233–235, 240, 244, 246, 254–257, 259, 260, 271, 276, 278–280, 282, 284, 435, 436, 518, 520, 527, 528, 555, 556, 558, 560  
 data type, **20**, *see* type  
*data-component-def-stmt* (R437), **68**, 68–70  
*data-edit-desc* (R1007), 260, **261**  
*data-i-do-object* (R541), **107**, 107, 108  
*data-i-do-variable* (R542), **107**, 107, 108, 158, 506  
*data-implied-do* (R540), **107**, 107, 108, 158, 506  
*data-pointer-component-name*, 164  
 data-pointer-initialization compatible, **71**  
*data-pointer-object* (R734), 163, **164**, 164, 165, 170, 350, 523  
*data-ref* (R611), 4, **122**, 123–125, 164, 229, 307, 310, 317, 323  
*data-stmt* (R537), 30, **107**, 297, 330, 508  
*data-stmt-constant* (R545), 86, **108**, 108, 109  
*data-stmt-object* (R539), **107**, 107–109  
*data-stmt-repeat* (R544), **108**, 108  
*data-stmt-set* (R538), **107**, 107  
*data-stmt-value* (R543), 107, **108**, 108  
*data-target* (R737), 82, 83, 103, 163, **164**, 164, 165, 170, 317, 330, 350, 514  
 DATE\_AND\_TIME, **368**  
 DBLE, 335, **369**  
 DC edit descriptor, 277  
*dealloc-opt* (R642), **134**, 134, 136

- DEALLOCATE statement, **134**, 136, 137, 198, 434, 492, 526
- deallocate-stmt* (R641), 31, **134**, 523
- decimal edit descriptor, 277
- decimal edit mode, 222
- decimal symbol, **7**, 222, 229, 249, 264–270, 277, 278
- decimal-edit-desc* (R1020), 261, **262**
- DECIMAL= specifier, 221, **222**, 226, 227, **229**, 241, 248, **249**, 277
- declaration, **7**, 33, 89–119
- declaration-construct* (R207), **30**, 30, 176
- declaration-type-spec* (R403), **53**, 53, 60, 68, 69, 89, 91, 112, 156, 304, 305, 323, 326
- declared type, **20**, 54, 71, 83, 87, 89, 123, 124, 131, 132, 134, 152, 154, 159, 162–165, 174, 194, 195, 243, 302, 307, 310, 313, 323, 329, 375, 403, 417, 507
- DEFAULT, **181**, **189**, **193**
- default character, **59**
- default complex, **58**
- default initialization, **8**, 8, 70–73, 82–84, 91, 98, 99, 107, 116, 117, 119, 312, 407, 511, 516, 521
- default real, **57**
- default-char-constant-expr* (R730), 93, **158**, 158, 226, 227
- default-char-expr* (R725), **154**, 154, 158, 197, 220–230
- default-char-variable* (R606), **121**, 121, 130, 221, 248–254
- default-initialized, **8**, 72, 92, 100, 512, 513, 517, 519, 521
- DEFERRED attribute, **75**, 76, 79
- deferred type parameter, **21**, 21, 52, 60, 84, 102, 118, 124, 131–133, 136, 137, 159, 160, 165, 294, 313, 314, 392, 407, 424, 475, 477, 483, 511, 516
- deferred-coshape-spec* (R510), 68, 93, **94**, 94
- deferred-shape array, **2**, 97
- deferred-shape-spec* (R520), 2, 68, 96, **97**, 97, 110
- definable, **8**, 100–103, 128, 160, 175, 231, 310, 312, 314, 319, 514, 523
- defined, **8**, 8, 22, 38, 40
- defined assignment, **8**, 19, 159, 162, 163, 167, 171, 293, 302, 307, 312, 331
- defined assignment statement, **163**, 320, 522
- defined input/output, **8**, 218, 223, 231–233, **237**, **238**, **238**, **239**, **240**, **241**, 237–243, 254, 274, 280, 281, 285, 293, 300, 303, 307, 320, 331, 436, 528, 579
- defined operation, **8**, 143, **152**, 153–155, 181, 293, 301, 307, 320, 331
- defined-binary-op* (R723), 14, 46, **142**, 142, 143, 152, 289
- defined-io-generic-spec* (R1209), 8, 76, 237–239, 243, **296**, 296, 300, 303
- defined-operator* (R309), **46**, 75, 290, 296
- defined-unary-op* (R703), 14, 46, **140**, 140, 143, 152, 289
- definition, **8**, 8
- definition of variables, 517
- deleted features, 25, 27, 28, 531, 532
- DELIM= specifier, 221, **222**, 226, 227, **229**, 241, 248, **250**, 283, 285, 557
- delimiter mode, 222
- derived type, 8, 18, **20**, 36, 37, 41, 51, 62–84, 87, 480, 481
- derived type definition statement, *see* TYPE statement
- derived type determination, 65
- derived-type type specifier, 53
- derived-type-def* (R426), 30, 54, **63**, 64, 67, 481
- derived-type-spec* (R454), 21, 53, 60, **81**, 82, 193, 194, 238, 505
- derived-type-stmt* (R427), **63**, 63, 64, 67, 92, 508
- descendant, **8**, 33, 64, 74, 75, 77, 103, 291, 504
- designator, 5, **8**, 9, 41, 98, 99, 104, 107, 115–117, **125**, 125, 154, 156, 157, 282, 312, 316, 317, 330, 332
- data object, 121
- designator* (R601), 71, 72, 107, 108, **121**, 121, 123–125, 139, 164, 174, 281, 330
- designator*, 139
- digit*, **23**, **43**, 43, 46, 56, 86, 279
- digit-string* (R411), 55, **56**, 56, 57, 265, 266, 272
- digit-string*, 56
- DIGITS, **369**
- DIM, **369**
- DIMENSION attribute, 69, 90, **96**, 96, 109, 118
- DIMENSION statement, **109**, 292
- dimension-spec* (R514), **96**
- dimension-stmt* (R548), 31, **109**, 508
- direct access, **213**
- direct access data transfer statement, **230**
- direct component, **6**, 6, 62, 63, 72, 312, 440, 480
- DIRECT= specifier, 248, **250**
- disassociated, 8, **9**, 22, 40, 54, 71–73, 91, 97, 109, 134, 136, 137, 155, 163, 165, 305, 317, 333, 342, 375, 407, 417, 424, 445, 511, 512, 514, 521

distinguishable, [303](#)

DO CONCURRENT construct, [180](#), [184](#), [187](#), [196](#), [331](#),  
[506](#), [507](#), [513](#), [519](#), [521](#), [526](#), [534](#)

DO CONCURRENT statement, [53](#), [170](#), [180](#)

DO construct, [35](#), [46](#), [87](#), [180](#), [196](#), [232](#), [531](#), [545](#), [546](#)

DO statement, [180](#), [518](#), [532](#), [534](#)

DO WHILE statement, [180](#)

*do-construct* (R819), [31](#), [181](#), [182](#), [184](#), [196](#)

*do-construct-name*, [181](#), [182](#), [184](#)

*do-stmt* (R820), [4](#), [181](#), [181](#), [182](#), [196](#), [522](#)

*do-variable* (R824), [87](#), [107](#), [181](#), [181](#), [182](#), [231](#), [254](#),  
[255](#), [257](#), [279](#), [518](#), [520](#), [522](#), [556](#)

DOT\_PRODUCT, [370](#)

DOUBLE PRECISION, [47](#), [55](#), [57](#), [63](#)

DP edit descriptor, [277](#)

DPROD, [370](#)

DSHIFTL, [370](#)

DSHIFTR, [371](#)

DT edit descriptor, [274](#)

*dtv-type-spec* (R921), [238](#)

dummy argument, [2](#), [3](#), [5](#), [9](#), [9](#), [13](#), [15](#), [21](#), [22](#), [36](#), [40](#),  
[41](#), [44](#), [52–54](#), [59](#), [60](#), [66](#), [69](#), [71](#), [76–78](#), [80](#), [82](#),  
[84](#), [90](#), [92–94](#), [97–102](#), [104](#), [105](#), [107](#), [110](#), [112](#),  
[115](#), [117](#), [131–133](#), [135](#), [136](#), [152](#), [155](#), [156](#), [163](#),  
[199](#), [234](#), [239–241](#), [293–297](#), [299–304](#), [306](#), [307](#),  
[309–317](#), [327](#), [328](#), [330](#), [332](#), [434](#), [483–485](#), [505](#),  
[506](#), [513](#), [522](#), [523](#), [565](#)

characteristics of, [294](#)

restrictions, [317](#)

dummy data object, [4](#), [5](#), [9](#), [54](#), [71](#), [91](#), [98–100](#), [104](#),  
[294](#), [301–303](#)

assumed-rank, [4](#), [54](#), [95](#), [96](#), [129](#), [191](#), [192](#), [294](#), [295](#),  
[303](#), [311–313](#), [318](#), [324](#), [390](#), [391](#), [414](#), [420](#), [422](#),  
[423](#), [430](#), [473](#), [483](#), [485](#), [594](#), [595](#), [597](#)

dummy function, [9](#), [60](#), [90](#)

dummy procedure, [5](#), [9](#), [11](#), [15](#), [99](#), [112](#), [157](#), [165](#), [293](#),  
[294](#), [296](#), [297](#), [299](#), [300](#), [303–305](#), [308](#), [316](#), [322](#),  
[324](#), [326](#), [330](#), [500](#), [504](#), [509](#)

*dummy-arg* (R1237), [326](#), [326–328](#)

*dummy-arg-name* (R1232), [110](#), [112](#), [293](#), [324](#), [325](#), [325](#),  
[326](#), [329](#), [508](#)

dynamic type, [15](#), [20](#), [22](#), [54](#), [78](#), [79](#), [81](#), [84](#), [87](#), [104](#),  
[132](#), [134](#), [136](#), [152](#), [154](#), [160](#), [162](#), [163](#), [165](#),  
[175](#), [193](#), [194](#), [203](#), [243](#), [307](#), [313](#), [323](#), [343](#),  
[375](#), [403](#), [417](#), [425](#), [484](#), [507](#), [511](#), [517](#), [544](#),  
[594](#)

## E

*e* (R1011), [261](#), [261](#), [267–270](#), [273](#), [274](#), [280](#)

E edit descriptor, [267](#)

edit descriptor, [260](#)

/, [275](#)

., [275](#)

A, [272](#)

B, [271](#)

BN, [276](#)

BZ, [276](#)

character string, [260](#), [277](#)

control, [260](#), [274–277](#)

D, [267](#)

data, [260](#), [264–274](#)

DC, [277](#)

decimal, [277](#)

DP, [277](#)

DT, [274](#)

E, [267](#)

EN, [268](#)

ES, [269](#)

EX, [270](#)

F, [266](#)

G, [272](#), [273](#)

H, [531](#)

I, [265](#)

L, [272](#)

O, [271](#)

P, [276](#)

position, [274](#)

RC, [276](#)

RD, [276](#)

RN, [276](#)

round, [276](#)

RP, [276](#)

RU, [276](#)

RZ, [276](#)

S, [276](#)

SP, [276](#)

SS, [276](#)

T, [275](#)

TL, [275](#)

TR, [275](#)

X, [275](#)

Z, [271](#)

effective argument, [2–4](#), [9](#), [21](#), [52](#), [54](#), [60](#), [95–100](#), [199](#),

- 310–312, 314, 316, 319, 322, 414, 476, 484, 507, 513, 516, 518, 521
- effective item, [9](#), [232](#), [235](#), [236](#), [240](#), [241](#), [243](#), [255](#), [262](#), [263](#), [275](#), [278](#), [279](#), [283](#), [528](#)
- effective position, [304](#)
- element sequence, [316](#)
- ELEMENTAL, [10](#), [323](#), [324](#), [328](#), [330](#), [332](#)
- elemental, [9](#), [20](#), [39](#), [60](#), [78](#), [80](#), [152](#), [158](#), [163](#), [166](#), [167](#), [169](#), [293–295](#), [305](#), [312](#), [316](#), [320](#), [322](#), [328](#), [332](#), [333](#), [338](#), [357](#), [358](#), [404](#), [445–447](#)
- elemental array assignment (FORALL), [169](#)
- elemental assignment, [9](#), [163](#)
- elemental operation, [10](#), [145](#), [155](#), [169](#)
- elemental operator, [10](#), [145](#), [440](#)
- elemental procedure, [10](#), [39](#), [155](#), [165](#), [305](#), [308](#), [317](#), [321](#), [323](#), [332](#), [332–334](#)
- elemental reference, [10](#), [169](#), [312](#), [320–323](#), [332](#)
- elemental subprogram, [10](#), [323](#), [324](#), [332](#)
- ELSE IF statement, [47](#), [187](#)
- ELSE statement, [187](#)
- else-if-stmt* (R836), [187](#), [187](#)
- else-stmt* (R837), [187](#), [187](#)
- ELSEWHERE statement, [47](#), [167](#)
- elsewhere-stmt* (R748), [167](#), [167](#), [168](#)
- EN edit descriptor, [268](#)
- ENCODING= specifier, [221](#), [222](#), [248](#), [250](#), [527](#)
- END ASSOCIATE statement, [47](#), [174](#)
- END BLOCK DATA statement, [47](#), [291](#)
- END BLOCK statement, [47](#), [136](#), [176](#)
- END CRITICAL statement, [47](#), [153](#), [179](#), [198](#)
- END DO statement, [47](#), [182](#)
- END ENUM statement, [47](#), [84](#)
- END FORALL statement, [47](#), [170](#)
- END FUNCTION statement, [47](#), [325](#)
- END IF statement, [47](#), [187](#), [531](#)
- END INTERFACE statement, [47](#), [296](#)
- END MODULE statement, [47](#), [288](#)
- END PROCEDURE statement, [47](#), [327](#)
- END PROGRAM statement, [47](#), [287](#)
- END SELECT statement, [47](#), [189](#), [194](#)
- END statement, [10](#), [34](#), [34](#), [35](#), [47](#), [49](#), [78](#), [79](#), [103](#), [118](#), [135](#), [136](#), [198](#), [476](#), [522](#)
- END SUBMODULE statement, [47](#), [291](#)
- END SUBROUTINE statement, [47](#), [326](#)
- END TEAM statement, [35](#), [177](#), [198](#), [337](#), [437](#)
- END TYPE statement, [47](#), [64](#)
- END WHERE statement, [47](#), [167](#)
- end-associate-stmt* (R806), [4](#), [174](#), [174](#), [196](#)
- end-block-data-stmt* (R1122), [10](#), [30](#), [34](#), [291](#), [291](#)
- end-block-stmt* (R810), [4](#), [175](#), [176](#), [176](#), [196](#)
- end-change-team-stmt* (R814), [177](#), [177](#)
- end-critical-stmt* (R818), [4](#), [179](#), [179](#), [196](#)
- end-do* (R831), [181](#), [182](#), [182](#), [184](#)
- end-do-stmt* (R832), [4](#), [182](#), [182](#), [196](#)
- end-enum-stmt* (R463), [84](#), [84](#)
- end-forall-stmt* (R754), [170](#), [170](#)
- end-function-stmt* (R1234), [10](#), [29](#), [31](#), [32](#), [34](#), [188](#), [296](#), [324](#), [325](#), [325](#), [328](#)
- end-if-stmt* (R838), [4](#), [187](#), [187](#), [196](#)
- end-interface-stmt* (R1204), [296](#), [296](#)
- end-module-stmt* (R1106), [10](#), [29](#), [34](#), [288](#), [288](#)
- end-mp-subprogram-stmt* (R1241), [10](#), [30–32](#), [34](#), [188](#), [327](#), [327](#), [328](#)
- end-program-stmt* (R1103), [10](#), [29](#), [31](#), [32](#), [34](#), [36](#), [79](#), [188](#), [197](#), [287](#), [287](#)
- end-select-rank-stmt* (R851), [4](#), [191](#), [191](#), [192](#), [196](#)
- end-select-stmt* (R843), [4](#), [189](#), [189](#), [196](#)
- end-select-type-stmt* (R855), [4](#), [193](#), [194](#), [194–196](#)
- end-submodule-stmt* (R1119), [10](#), [29](#), [34](#), [291](#), [291](#)
- end-subroutine-stmt* (R1238), [10](#), [29](#), [31](#), [32](#), [34](#), [188](#), [296](#), [326](#), [326](#), [328](#)
- end-type-stmt* (R430), [63](#), [64](#)
- end-where-stmt* (R749), [167](#), [167](#), [168](#)
- END= specifier, [4](#), [226](#), [227](#), [244](#), [245](#), [255](#)
- endfile record, [212](#)
- ENDFILE statement, [47](#), [212](#), [213](#), [215](#), [222](#), [241](#), [244](#), [246](#), [555](#)
- endfile-stmt* (R925), [31](#), [245](#), [331](#)
- entity-decl* (R503), [60](#), [69](#), [89](#), [90](#), [90](#), [91](#), [157](#), [158](#), [508](#)
- entity-name*, [106](#), [111](#)
- ENTRY statement, [9](#), [34](#), [152](#), [163](#), [288](#), [293](#), [297](#), [323](#), [324](#), [327](#), [505](#), [515](#), [516](#), [532](#), [534](#)
- entry-name*, [324](#), [327](#), [505](#)
- entry-stmt* (R1242), [30](#), [288](#), [291](#), [297](#), [327](#), [327](#), [505](#), [508](#)
- ENUM statement, [84](#)
- enum-def* (R459), [30](#), [84](#), [85](#)
- enum-def-stmt* (R460), [84](#), [84](#)
- enumeration, [84](#)
- enumerator, [84](#)
- enumerator* (R462), [84](#), [84](#)
- ENUMERATOR statement, [84](#)

- enumerator-def-stmt* (R461), **84**, **84**
- EOR= specifier, **4**, **226**, **227**, **244**, **245**, **255**, **255**, **555**, **556**
- EOSHIFT, **371**
- EPSILON, **372**
- equiv-op* (R721), **45**, **142**, **142**
- equiv-operand* (R716), **142**, **142**
- equivalence association, **116**
- EQUIVALENCE statement, **115**, **115–119**, **176**, **290**, **292**, **515**, **516**, **532**, **534**
- equivalence-object* (R571), **115**, **115–117**, **290**
- equivalence-set* (R570), **115**, **115–117**
- equivalence-stmt* (R569), **31**, **115**, **508**
- ERF, **372**
- ERFC, **373**
- ERFC\_SCALED, **373**
- ERR= specifier, **4**, **221**, **225**, **226**, **244**, **245**, **247–249**, **254**
- errmsg-variable* (R630), **130**, **130**, **132**, **134**, **135**, **137**, **200**, **522**, **526**
- ERRMSG= specifier, **xix**, **130**, **133–135**, **137**, **180**, **200**, **208**, **519**, **526**
- error indicator, **490**
- ERROR STOP statement, **35**, **36**, **197**, **526**
- error termination, **36**, **79**, **133**, **135**, **197**, **209**, **240**, **254**, **255**, **337**, **352–356**, **373**, **374**, **404**, **415**, **525**, **528**
- error-stop-stmt* (R861), **31**, **79**, **197**
- ERROR\_UNIT, **218**, **219**, **223**, **434**
- ES edit descriptor, **269**
- established coarray, **5**, **129**
- evaluation
  - operations, **145**
  - optional, **153**
  - parentheses, **154**
- EVENT POST statement, **198**, **205**, **206**, **209**, **373**, **434**, **435**, **520**, **522**, **526**
- event variable, **22**, **199**, **205**, **206**, **373**, **374**, **434**, **520**
- EVENT WAIT statement, **198**, **205**, **205**, **209**, **434**, **435**, **520**, **522**, **526**
- event-post-stmt* (R870), **31**, **205**
- event-variable* (R871), **205**, **205**, **434**, **522**
- event-wait-spec* (R873), **205**, **205**
- event-wait-stmt* (R872), **31**, **205**, **205**
- EVENT\_QUERY, **xix**, **373**, **585**
- EVENT\_TYPE, **22**, **64**, **100**, **131**, **132**, **205**, **434**
- EX edit descriptor, **270**
- executable construct, **173**
- executable statement, **18**, **18**, **33**
- executable-construct* (R214), **18**, **30**, **31**, **327**
- EXECUTE\_COMMAND\_LINE, **344**, **374**
- execution control, **173**
- execution-part* (R209), **29**, **30**, **30**, **32**, **287**, **324–327**
- execution-part-construct* (R210), **30**, **30**, **173**
- exist, **212**, **219**
- EXIST= specifier, **248**, **250**
- EXIT statement, **174**, **184**, **196**
- exit-stmt* (R856), **31**, **177**, **196**, **196**
- EXP, **374**
- explicit formatting, **259–277**
- explicit initialization, **10**, **72**, **73**, **90**, **91**, **107**, **511**, **516**, **517**
- explicit interface, **10**, **26**, **71**, **75**, **166**, **294–298**, **300**, **305**, **307**, **309**, **316**, **329**, **330**, **504**, **505**, **522**, **564**
- explicit-coshape-spec* (R511), **93**, **94**, **94**
- explicit-shape array, **2**, **54**, **70**, **94**, **97**, **159**, **312**, **317**, **482**
- explicit-shape-spec* (R516), **2**, **68**, **69**, **91**, **96**, **96–98**, **117**
- EXPONENT, **375**, **453**
- exponent* (R417), **57**, **57**
- exponent-letter* (R416), **57**, **57**
- expr* (R722), **24**, **78**, **82**, **83**, **87**, **130**, **139**, **142**, **142**, **154**, **158–165**, **169**, **170**, **174**, **189**, **230**, **307**, **308**, **329**, **330**, **469**, **522**
- expression, **139**, **139–159**
  - constant, **4**, **7**, **21**, **27**, **51**, **52**, **60**, **67**, **69**, **70**, **72**, **87**, **91**, **95**, **97**, **98**, **107**, **108**, **110**, **116**, **156**, **157**, **158**, **158**, **228**, **294**, **295**, **317**, **334**, **346**, **348**, **349**, **360**, **361**, **366**, **376–378**, **382**, **385**, **387**, **388**, **391–393**, **395**, **396**, **398**, **401**, **406**, **414**, **416**, **418**, **420**, **422**, **424**, **430–432**, **482**
  - specification, **17**, **21**, **35**, **67–69**, **78**, **92**, **125**, **156**, **157**, **157**, **176**, **328**, **441**, **525**, **526**, **533**
- extended real model, **336**
- extended type, **3**, **6**, **12**, **20**, **21**, **67**, **74**, **78–80**, **516**, **535**, **541**
- extended-intrinsic-op* (R310), **46**, **46**
- EXTENDS attribute, **20**, **79**, **79**, **481**
- EXTENDS\_TYPE\_OF, **68**, **69**, **375**
- extensible type, **20**, **53**, **63**, **71**, **79**, **238**, **375**, **417**, **543**, **581**
- extension operation, **143**
- extension type, **20**, **54**, **79**, **81**, **194**, **313**, **375**, **581**
- extent, **10**, **39**

- EXTERNAL attribute, [14](#), [25](#), [26](#), [99](#), [99](#), [102](#), [111](#), [112](#), [114](#), [165](#), [288](#), [292](#), [294](#), [297](#), [304](#), [316](#), [321](#), [322](#), [508](#), [509](#), [573](#)
- external file, [10](#), [10](#), [27](#), [211–216](#), [218–220](#), [224](#), [228](#), [247](#), [264](#), [274](#), [283](#), [331](#), [500](#), [527](#), [555](#), [589](#)
- external input/output unit, [10](#), [503](#)
- external linkage, [93](#), [473](#), [498–500](#)
- external procedure, [15](#), [25](#), [32](#), [75](#), [99](#), [112](#), [165](#), [204](#), [293](#), [294](#), [296–300](#), [304](#), [305](#), [308](#), [316](#), [322](#), [323](#), [503](#), [504](#), [508](#), [509](#), [564](#), [565](#), [570](#), [573](#), [574](#)
- EXTERNAL statement, [99](#), [304](#)
- external subprogram, [16](#), [19](#), [32](#), [293](#)
- external unit, [10](#), [218–220](#), [234](#), [240](#), [241](#), [251](#), [256](#), [434–436](#), [527](#)
- external-name*, [304](#)
- external-stmt* (R1212), [31](#), [304](#)
- external-subprogram* (R203), [29](#), [29](#), [298](#), [327](#)
- ## F
- F edit descriptor, [266](#)
- FAIL IMAGE statement, [197](#)
- fail-image-stmt* (R863), [31](#), [197](#)
- failed image, [11](#), [35](#), [36](#), [130](#), [207](#), [338](#), [404](#), [525](#)
- FAILED\_IMAGES, [376](#)
- field, [262](#)
- file
- connected, [219](#)
  - external, [10](#), [10](#), [27](#), [211–216](#), [218–220](#), [224](#), [228](#), [247](#), [264](#), [274](#), [283](#), [331](#), [500](#), [527](#), [555](#), [589](#)
  - internal, [13](#), [13](#), [211](#), [217](#), [218](#), [220](#), [228](#), [232](#), [235](#), [237](#), [240](#), [241](#), [254](#), [256](#), [274](#), [275](#), [518](#), [520](#), [528](#)
- file access method, [212–214](#)
- file connection, [217–225](#)
- file inquiry statement, [247](#)
- file position, [212](#), [215](#)
- file positioning statement, [212](#), [245](#)
- file storage unit, [10](#), [18](#), [211](#), [214–217](#), [223](#), [228](#), [230](#), [236](#), [246](#), [252–254](#), [435](#), [514](#), [527](#)
- file-name-expr* (R906), [221](#), [221](#), [223](#), [248–250](#)
- file-unit-number* (R902), [217](#), [217](#), [218](#), [220](#), [221](#), [225](#), [227](#), [239](#), [244](#), [245](#), [247–254](#), [331](#), [435](#)
- FILE= specifier, [221](#), [223](#), [248](#), [249](#), [249](#), [522](#), [557](#)
- FILE\_STORAGE\_SIZE, [435](#)
- FINAL statement, [10](#), [77](#)
- final subroutine, [4](#), [10](#), [11](#), [26](#), [76–79](#), [128](#), [156](#), [311](#), [540](#), [541](#)
- final-procedure-stmt* (R453), [75](#), [77](#)
- final-subroutine-name*, [77](#)
- finalizable, [11](#), [26](#), [77](#), [78](#), [98](#), [99](#), [135](#)
- finalization, [11](#), [16](#), [77–79](#), [128](#), [170](#), [293](#), [307](#), [320](#), [330](#), [331](#), [526](#)
- FINDLOC, [376](#)
- fixed source form, [48](#), [48](#)
- FLOOR, [377](#)
- FLUSH statement, [213](#), [244](#), [247](#), [256](#)
- flush-spec* (R929), [247](#), [247](#)
- flush-stmt* (R928), [31](#), [247](#), [331](#)
- FMT= specifier, [226](#), [227](#)
- FORALL construct, [169](#), [331](#), [506](#), [507](#), [519](#), [532](#), [534](#)
- FORALL statement, [53](#), [145](#), [171](#), [506](#), [507](#), [518](#)
- forall-assignment-stmt* (R753), [145](#), [170](#), [170](#), [171](#), [331](#)
- forall-body-construct* (R752), [170](#), [170](#), [171](#)
- forall-construct* (R750), [31](#), [169](#), [170](#), [171](#)
- forall-construct-name*, [170](#)
- forall-construct-stmt* (R751), [4](#), [169](#), [170](#), [170](#), [196](#)
- forall-stmt* (R755), [32](#), [170](#), [171](#), [171](#), [196](#)
- FORM TEAM statement, [19](#), [20](#), [35](#), [129](#), [198](#), [206](#), [437](#), [523](#), [526](#)
- form-team-spec* (R877), [206](#), [206](#)
- form-team-stmt* (R875), [31](#), [206](#)
- FORM= specifier, [221](#), [223](#), [248](#), [250](#)
- format* (R915), [225](#), [226](#), [227](#), [227](#), [228](#), [235](#), [259](#), [260](#)
- format control, [262](#)
- format descriptor, *see* edit descriptor
- FORMAT statement, [25](#), [34](#), [46](#), [228](#), [259](#), [259](#), [288](#)
- format-item* (R1004), [260](#), [260](#)
- format-items* (R1003), [259](#), [260](#), [260](#)
- format-specification* (R1002), [259](#), [259](#)
- format-stmt* (R1001), [30](#), [259](#), [259](#), [288](#), [291](#), [297](#)
- FORMATTED, [238](#), [239](#), [296](#)
- formatted data transfer, [236](#)
- formatted input/output statement, [211](#), [227](#)
- formatted record, [211](#)
- FORMATTED= specifier, [248](#), [250](#)
- formatting
- explicit, [259–277](#)
  - list-directed, [237](#), [277–281](#)
  - namelist, [237](#), [281–285](#)
- forms, [212](#)
- Fortran 2003 compatibility, [26](#)
- Fortran 2008 compatibility, [26](#)
- FORTTRAN 77 compatibility, [27](#)
- Fortran 90 compatibility, [27](#)

Fortran 95 compatibility, [27](#)  
 Fortran character set, [43](#), [59](#)  
 FRACTION, [378](#)  
 free source form, [47](#), [47](#)  
 function, [11](#)  
     intrinsic, [333](#)  
     intrinsic elemental, [333](#)  
     intrinsic inquiry, [333](#)  
 function reference, [16](#), [37](#), [38](#), [320](#)  
 function result, [11](#), [26](#), [60](#), [89](#), [113](#), [115](#), [117](#), [135](#), [294](#),  
     [325](#), [327](#), [332](#), [483](#), [505](#), [515](#), [521](#)  
 FUNCTION statement, [9](#), [53](#), [113](#), [152](#), [156](#), [287](#), [323](#),  
     [324](#), [327](#), [328](#), [505](#)  
*function-name*, [90](#), [297](#), [324](#), [325](#), [327](#), [329](#), [505](#), [508](#)  
*function-reference* (R1221), [82](#), [90](#), [121](#), [139](#), [307](#), [310](#),  
     [320](#)  
*function-stmt* (R1230), [29](#), [296](#), [297](#), [324](#), [324](#), [325](#), [505](#),  
     [508](#)  
*function-subprogram* (R1229), [19](#), [29](#), [30](#), [288](#), [324](#), [326](#)

## G

G edit descriptor, [272](#), [273](#)  
 GAMMA, [378](#)  
 generic identifier, [11](#), [12](#), [288](#), [297](#), [300](#), [301](#), [303](#), [322](#),  
     [333](#), [503](#), [508](#)  
 generic interface, [12](#), [76](#), [77](#), [80](#), [84](#), [101](#), [152](#), [163](#), [243](#),  
     [289](#), [290](#), [300](#), [300–302](#), [321](#), [504](#), [583](#)  
 generic interface block, [12](#), [12](#), [297](#), [300](#)  
 generic procedure reference, [303](#)  
 GENERIC statement, [75](#), [76](#), [298](#), [300](#), [321](#)  
*generic-name*, [75](#), [76](#), [296](#), [505](#), [508](#)  
*generic-spec* (R1208), [12](#), [13](#), [75–77](#), [80](#), [105](#), [152](#), [163](#),  
     [289](#), [290](#), [296](#), [296–298](#), [300](#), [321](#), [505](#), [508](#)  
*generic-stmt* (R1210), [30](#), [298](#)  
 GET\_COMMAND, [378](#)  
 GET\_COMMAND\_ARGUMENT, [379](#)  
 GET\_ENVIRONMENT\_VARIABLE, [380](#), [529](#)  
 GET\_TEAM, [381](#), [434–436](#)  
 global entity, [503](#)  
 global identifier, [503](#)  
 GO TO statement, [4](#), [47](#), [196](#), [196](#)  
*goto-stmt* (R857), [31](#), [196](#), [196](#)  
 graphic character, [43](#), [60](#), [284](#)

## H

halting mode, [439](#), [443](#), [443](#), [446](#), [449](#), [461](#), [465](#), [501](#),  
     [530](#)

*hex-constant* (R467), [86](#), [86](#)  
*hex-digit* (R468), [86](#), [86](#), [272](#)  
*hex-digit-string* (R1022), [272](#), [272](#)  
 host, [11](#), [12](#), [32](#), [291](#), [329](#), [505](#), [508](#), [509](#)  
 host association, [3](#), [3](#), [15](#), [33](#), [40](#), [53](#), [60](#), [92](#), [105](#), [107](#),  
     [112](#), [118](#), [156](#), [157](#), [165](#), [291](#), [293](#), [317](#), [329–331](#),  
     [506–510](#), [513](#), [516](#), [607](#)  
 host instance, [11](#), [166](#), [308](#), [309](#), [316](#), [326](#), [508](#), [513](#), [516](#),  
     [522](#)  
 host scoping unit, [11](#), [32](#), [112](#), [298](#), [299](#), [321](#), [509](#), [516](#)  
 HUGE, [382](#)  
 HYPOT, [382](#)

## I

I edit descriptor, [265](#)  
 IACHAR, [62](#), [162](#), [382](#)  
 IALL, [383](#)  
 IAND, [353](#), [383](#)  
 IANY, [384](#)  
 IBCLR, [384](#)  
 IBITS, [384](#)  
 IBSET, [385](#)  
 ICHAR, [61](#), [385](#)  
*id-variable* (R914), [226](#), [226](#)  
 ID= specifier, [226](#), [227](#), [229](#), [244](#), [248](#), [249](#), [250](#), [522](#),  
     [559](#)  
 IEEE infinity, [11](#)  
 IEEE NaN, [11](#)  
 IEEE\_ALL, [440](#)  
 IEEE\_ARITHMETIC, [156](#), [158](#), [346](#), [439–469](#)  
 IEEE\_AWAY, [442](#), [450](#)  
 IEEE\_CLASS, [447](#), [447](#)  
 IEEE\_CLASS\_TYPE, [440](#), [447](#), [468](#)  
 IEEE\_COPY\_SIGN, [444](#), [448](#)  
 IEEE\_DATATYPE, [440](#)  
 IEEE\_DENORMAL, [440](#)  
 IEEE\_DIVIDE, [440](#)  
 IEEE\_DIVIDE\_BY\_ZERO, [440](#)  
 IEEE\_DOWN, [440](#), [442](#)  
 IEEE\_EXCEPTIONS, [156](#), [184](#), [439–469](#)  
 IEEE\_FEATURES, [439–440](#)  
 IEEE\_FEATURES\_TYPE, [440](#)  
 IEEE\_FLAG\_TYPE, [440](#), [448](#), [449](#), [460](#), [461](#), [464](#), [465](#)  
 IEEE\_FMA, [448](#)  
 IEEE\_GET\_FLAG, [184](#), [448](#), [469](#), [470](#), [530](#)  
 IEEE\_GET\_HALTING\_MODE, [184](#), [449](#), [449](#)  
 IEEE\_GET\_MODES, [443](#), [449](#), [449](#), [461](#)

IEEE\_GET\_ROUNDING\_MODE, 442, **450**, 450, 462  
 IEEE\_GET\_STATUS, **450**, 450, 462, 470, 530  
 IEEE\_GET\_UNDERFLOW\_MODE, **450**, 463  
 IEEE\_HALTING, **440**  
 IEEE\_INEXACT, **440**  
 IEEE\_INEXACT\_FLAG, **440**  
 IEEE\_INF, **440**  
 IEEE\_INT, **451**, 530  
 IEEE\_INVALID, **440**  
 IEEE\_INVALID\_FLAG, **440**  
 IEEE\_IS\_FINITE, **451**  
 IEEE\_IS\_NAN, **451**  
 IEEE\_IS\_NEGATIVE, **452**  
 IEEE\_IS\_NORMAL, **452**  
 IEEE\_LOGB, 444, **452**  
 IEEE\_MAX\_NUM, **453**, 530  
 IEEE\_MAX\_NUM\_MAG, **453**, 530  
 IEEE\_MIN\_NUM, **454**, 530  
 IEEE\_MIN\_NUM\_MAG, **454**, 530  
 IEEE\_MODES\_TYPE, **440**, 443, 449, 461  
 IEEE\_NAN, **440**  
 IEEE\_NEAREST, **440**, **442**  
 IEEE\_NEGATIVE\_DENORMAL, **440**  
 IEEE\_NEGATIVE\_INF, **440**  
 IEEE\_NEGATIVE\_NORMAL, **440**  
 IEEE\_NEGATIVE\_SUBNORMAL, **440**, 440, 447, 451, 452  
 IEEE\_NEGATIVE\_ZERO, **440**  
 IEEE\_NEXT\_AFTER, 444, **455**  
 IEEE\_NEXT\_DOWN, **455**, 455  
 IEEE\_NEXT\_UP, **455**  
 IEEE\_OTHER, **440**  
 IEEE\_OTHER\_VALUE, **440**  
 IEEE\_OVERFLOW, **440**  
 IEEE\_POSITIVE\_DENORMAL, **440**  
 IEEE\_POSITIVE\_INF, **440**  
 IEEE\_POSITIVE\_NORMAL, **440**  
 IEEE\_POSITIVE\_SUBNORMAL, **440**, 440, 447, 451  
 IEEE\_POSITIVE\_ZERO, **440**  
 IEEE\_QUIET\_EQ, **456**  
 IEEE\_QUIET\_GE, **456**  
 IEEE\_QUIET\_GT, **457**  
 IEEE\_QUIET\_LE, **457**  
 IEEE\_QUIET\_LT, **457**  
 IEEE\_QUIET\_NAN, **440**  
 IEEE\_QUIET\_NE, **458**  
 IEEE\_REAL, **458**  
 IEEE\_REM, 444, **458**  
 IEEE\_RINT, 444, **459**  
 IEEE\_ROUND\_TYPE, **440**, 450, 451, 459, 461, 462, 466  
 IEEE\_ROUNDING, **440**  
 IEEE\_SCALB, 444, **459**  
 IEEE\_SELECTED\_REAL\_KIND, **460**  
 IEEE\_SET\_FLAG, 450, **460**, 462, 470, 530  
 IEEE\_SET\_HALTING\_MODE, 184, 449, **461**, 465, 470, 530  
 IEEE\_SET\_MODES, 443, 449, **461**, 461  
 IEEE\_SET\_ROUNDING\_MODE, 442, 449, 450, **461**, 461, 462  
 IEEE\_SET\_STATUS, 442, 450, **462**, 462, 470, 530  
 IEEE\_SET\_UNDERFLOW\_MODE, 449, 451, 461, **462**  
 IEEE\_SIGNALING\_NAN, **440**  
 IEEE\_SIGNBIT, **463**  
 IEEE\_SQRT, **440**  
 IEEE\_STATUS\_TYPE, **440**, 443, 450, 462, 470  
 IEEE\_SUBNORMAL, **440**  
 IEEE\_SUPPORT\_DATATYPE, 440, 441, 447, 448, 451, 453–459, 461, 462, **463**, 463, 467, 468  
 IEEE\_SUPPORT\_DENORMAL, **463**  
 IEEE\_SUPPORT\_DIVIDE, **464**, 467  
 IEEE\_SUPPORT\_FLAG, **464**, 467  
 IEEE\_SUPPORT\_HALTING, **465**, 467  
 IEEE\_SUPPORT\_INF, 444, 455, 456, **465**, 467, 468  
 IEEE\_SUPPORT\_IO, **465**  
 IEEE\_SUPPORT\_NAN, 441, 444, **466**, 467, 468  
 IEEE\_SUPPORT\_ROUNDING, 461, 462, **466**, 467  
 IEEE\_SUPPORT\_SQRT, **466**, 467  
 IEEE\_SUPPORT\_STANDARD, **467**  
 IEEE\_SUPPORT\_SUBNORMAL, 444, 445, 447, 455, 464, **467**, 467, 468  
 IEEE\_SUPPORT\_UNDERFLOW\_CONTROL, **468**  
 IEEE\_TO\_ZERO, **440**, **442**  
 IEEE\_UNDERFLOW, **440**  
 IEEE\_UNDERFLOW\_FLAG, **440**  
 IEEE\_UNORDERED, 444, **468**  
 IEEE\_UP, **440**, **442**  
 IEEE\_USUAL, **440**  
 IEEE\_VALUE, 453, 454, **468**  
 IEOR, 356, **385**  
 IF construct, 35, **187**, 531

- IF statement, [145](#), [188](#)
- if-construct* (R834), [31](#), [187](#), [187](#)
- if-construct-name*, [187](#)
- if-stmt* (R839), [31](#), [188](#), [188](#)
- if-then-stmt* (R835), [4](#), [187](#), [187](#), [196](#)
- imag-part* (R420), [58](#), [58](#)
- image, [1](#), [11](#), [11](#), [12](#), [34–36](#), [39](#), [40](#), [78](#), [94](#), [130–133](#), [135–137](#), [162](#), [164](#), [165](#), [179](#), [180](#), [197–204](#), [206–209](#), [212](#), [213](#), [218](#), [219](#), [251](#), [307](#), [311](#), [315](#), [333](#), [337](#), [342–344](#), [352–356](#), [367](#), [368](#), [381](#), [386](#), [387](#), [404](#), [408](#), [412](#), [426–428](#), [431](#), [437](#), [443](#), [503](#), [511](#), [512](#), [520](#), [522](#)
  - active, [11](#)
  - failed, [11](#)
  - stopped, [11](#)
- image control statement, [12](#), [35](#), [153](#), [179](#), [184](#), [198](#), [198–200](#), [203](#), [209](#), [331](#), [337](#), [376](#), [424](#), [437](#)
- image index, [12](#), [34](#), [39](#), [129](#), [130](#), [201](#), [206](#), [212](#), [315](#), [341](#), [343](#), [386](#), [428](#), [431](#), [503](#)
- image-selector* (R624), [5](#), [7](#), [35](#), [123](#), [124](#), [129](#), [129](#), [130](#), [281](#)
- image-selector-spec* (R626), [129](#), [129](#)
- image-set* (R867), [201](#), [201](#)
- IMAGE\_INDEX, [386](#)
- IMAGE\_STATUS, [386](#)
- imaginary part, [58](#)
- implicit interface, [12](#), [69](#), [166](#), [288](#), [305–307](#), [316](#), [477](#), [509](#)
- IMPLICIT NONE statement, [112](#)
- IMPLICIT statement, [34](#), [112](#), [115](#), [292](#)
- implicit-none-spec* (R566), [112](#), [112](#)
- implicit-part* (R205), [30](#), [30](#), [176](#)
- implicit-part-stmt* (R206), [30](#), [30](#)
- implicit-spec* (R564), [112](#), [112](#)
- implicit-stmt* (R563), [30](#), [112](#)
- implied-shape array, [98](#)
- implied-shape-or-assumed-size-spec* (R523), [96](#), [98](#), [98](#), [99](#)
- implied-shape-spec* (R524), [96](#), [99](#), [99](#)
- IMPORT statement, [34](#), [298](#), [509](#)
- import-name*, [298](#), [299](#)
- import-stmt* (R1211), [30](#), [176](#), [298](#)
- IMPURE, [323](#), [324](#), [328](#), [330](#), [332](#)
- IN, [100](#)
- INCLUDE line, [47](#), [49](#)
- inclusive scope, [12](#), [113](#), [176](#), [196](#), [221](#), [225](#), [226](#), [228](#), [244](#), [245](#), [247](#), [249](#), [308](#), [328](#), [503](#), [504](#)
- INDEX, [387](#)
- index-name*, [170](#), [171](#), [181–183](#), [506](#), [507](#), [519](#)
- inherit, [3](#), [6](#), [12](#), [63](#), [76](#), [78–81](#), [516](#), [543](#)
- inheritance association, [3](#), [3](#), [6](#), [80](#), [82](#), [513](#), [516](#)
- initial team, [20](#), [40](#), [129](#), [218](#), [344](#), [381](#), [386](#), [408](#), [427](#), [435](#)
- initial-data-target* (R444), [71](#), [71](#), [72](#), [90](#), [91](#), [108](#), [109](#)
- initial-proc-target* (R1219), [72](#), [305](#), [305](#), [306](#)
- INITIAL\_TEAM, [381](#), [435](#)
- initialization, [91](#)
  - default, [8](#), [8](#), [70–73](#), [82–84](#), [91](#), [98](#), [99](#), [107](#), [116](#), [117](#), [119](#), [312](#), [511](#), [516](#), [521](#)
  - explicit, [10](#), [72](#), [73](#), [90](#), [91](#), [107](#), [511](#), [516](#), [517](#)
- initialization* (R505), [90](#), [90](#), [91](#), [158](#)
- INOUT, [47](#), [100](#)
- input statement, [225](#)
- input-item* (R916), [225](#), [226](#), [230](#), [231](#), [243](#), [257](#), [522](#)
- input/output editing, [259–285](#)
- input/output list, [230](#)
- input/output statement, [518](#)
- input/output statements, [211–256](#)
- input/output unit, [13](#), [22](#), [34](#)
- INPUT\_UNIT, [218](#), [219](#), [223](#), [240](#), [435](#)
- INQUIRE statement, [27](#), [213](#), [214](#), [216](#), [217](#), [219](#), [220](#), [229](#), [230](#), [240](#), [241](#), [244](#), [247](#), [256](#), [257](#), [435](#), [518](#), [520–522](#), [528](#), [554](#)
- inquire-spec* (R931), [248](#), [248](#), [249](#), [257](#)
- inquire-stmt* (R930), [31](#), [248](#), [331](#)
- inquiry function, [12](#), [20](#), [94](#), [97](#), [99](#), [123](#), [133](#), [156](#), [310](#), [311](#), [333–335](#), [338](#), [348](#), [350](#), [359](#), [365](#), [369](#), [372](#), [375](#), [382](#), [386](#), [390–392](#), [397](#), [400](#), [405](#), [410](#), [411](#), [413](#), [417](#), [420](#), [422](#), [424](#), [428](#), [430](#), [431](#), [439](#), [440](#), [442–445](#), [463–468](#), [474](#), [477](#), [478](#)
- inquiry, type parameter, [124](#)
- instance, [326](#)
- INT, [109](#), [161](#), [335](#), [371](#), [383](#), [386](#), [388](#), [388](#), [400](#)
- int-constant* (R307), [45](#), [45](#), [108](#)
- int-constant-expr* (R731), [55](#), [59](#), [66](#), [67](#), [84](#), [85](#), [107](#), [158](#), [158](#), [191](#), [192](#)
- int-constant-name*, [56](#)
- int-constant-subobject* (R546), [108](#), [108](#)
- int-expr* (R726), [34](#), [52](#), [87](#), [122](#), [125](#), [126](#), [129](#), [130](#), [145](#), [154](#), [154](#), [156–158](#), [181](#), [182](#), [196](#), [197](#), [201](#), [205](#), [206](#), [217](#), [218](#), [221](#), [226](#), [231](#), [233](#), [244](#), [248](#), [328](#)

- int-literal-constant* (R408), 45, **55**, 55, 59, 260, 261
- int-variable* (R607), **121**, 121, 130, 221, 225–227, 244, 245, 247, 248, 251–256, 528
- int-variable-name*, 181
- INT16, **435**
- INT32, **435**
- INT64, **435**
- INT8, **435**
- integer constant, 55
- integer editing, 265
- integer model, 336
- integer type, 55–56
- integer-type-spec* (R405), 53, **55**, 55, 66, 87, 107, 181, 506
- INTEGER\_KINDS, **435**
- INTENT (IN) attribute, **99**, 100, 101, 105, 181, 301–303, 311, 314, 316, 318, 330, 331, 334, 352–356, 362–364, 373, 374, 379, 380, 404, 412, 447, 475–477, 497, 576, 592
- INTENT (INOUT) attribute, **99**, 100, 101, 104, 302, 312, 314, 320, 331, 332, 352–356, 362–364, 374, 379, 380, 403, 404, 434, 436, 522, 523, 592
- INTENT (OUT) attribute, 26, 54, 77, 78, 98, **99**, 99–101, 104, 135, 156, 302, 312, 314, 320, 330–332, 352–356, 362–364, 366, 368, 373, 374, 378–380, 403, 404, 412, 426, 448–450, 475, 477, 497, 498, 512, 513, 518, 519, 521–523, 592
- INTENT attribute, **99**, 99–101, 110, 185
- INTENT statement, **110**, 176
- intent-spec* (R526), 89, **100**, 110, 304
- intent-stmt* (R549), 31, **110**
- interface, **12**, 12, 33, 38, 41, 69, 75, 76, 101, 238, 239, 274, 294, **295**, 307, 316, 320, 321, 327, 329, 330, 483–485, 500, 565
- abstract, 288, **295**, 297, 305, 324, 504, 508
- explicit, **10**, 26, 71, 75, 166, 294–298, 300, 305, 307, 309, 316, 329, 330, 504, 505, 522, 564
- generic, **12**, 76, 77, 80, 84, 101, 152, 163, 243, 289, 290, **300**, 300–302, 321, 504
- implicit, **12**, 69, 166, 288, 305–307, 316, 477, 509
- procedure, 295
- specific, **12**, 243, **297**, 297, 300, 305, 321
- interface block, **12**, 33, 238, 243, 289, 296–300, 321, 565
- interface body, 12, 13, 16, 34, 95, 97, 99, 112, 156, **296**, 296, 324, 327, 329, 485, 505, 508, 565
- INTERFACE statement, **296**, 565
- interface-block* (R1201), 30, **296**, 296
- interface-body* (R1205), **296**, 296, 297
- interface-name* (R1217), 75, 76, 304, **305**, 305
- interface-specification* (R1202), **296**, 296, 297
- interface-stmt* (R1203), **296**, 296, 297, 300, 508
- internal file, **13**, 13, 211, 217, 218, 220, 228, 232, 235, 237, 240, 241, 254, 256, 274, 275, 518, 520, 528
- internal procedure, 11, **15**, 32, 165, 293–296, 308, 316, 322, 324, 326, 500, 504, 505, 509
- internal subprogram, **19**, 32, 34, 112, 293, 321, 508
- internal unit, **13**, 13, 218, 220, 234, 240, 249, 256, 435, 436
- internal-file-variable* (R903), **217**, 217, 218, 227, 257, 522
- internal-subprogram* (R212), **30**, 30
- internal-subprogram-part* (R211), 29, **30**, 30, 287, 324–327
- interoperable, **13**, 84, 85, 93, 324, 329, 475, 477–485, 499, 500
- interoperate, **479**
- intrinsic, 6, 9–12, **13**, 15, 19, 20, 37, 39, 41, 52, 54, 78, 86, 99, 295, 321, 322, 432, 504, 506
- intrinsic assignment statement, 83, 92, 135, 137, 155, **159**, 163, 165, 209, 247, 256, 279, 283, 330, 331, 518, 526
- INTRINSIC attribute, 99, **101**, 101, 102, 288, 306, 321, 509
- intrinsic function, 333
- intrinsic operation, 145–152
- intrinsic procedure, 333–432
- INTRINSIC statement, 292, **306**
- intrinsic subroutines, 333
- intrinsic type, 6, **20**, 36, 37, 51, 55–62, 484, 488
- intrinsic-operator* (R308), 14, **45**, 46, 140, 142, 145, 146, 152, 301
- intrinsic-procedure-name*, 306, 508
- intrinsic-stmt* (R1220), 31, **306**, 508
- intrinsic-type-spec* (R404), 53, **55**, 60
- io-control-spec* (R913), 225, **226**, 226, 227, 230, 240, 257
- io-implied-do* (R918), 230, **231**, 231, 232, 235, 257, 518, 520, 522, 556
- io-implied-do-control* (R920), **231**, 231, 233
- io-implied-do-object* (R919), **231**, 231, 235
- io-unit* (R901), 22, **217**, 217, 218, 226, 227, 331
- IOLength= specifier, 216, 247, **254**
- iomsg-variable* (R907), **221**, 221, 225, 226, 244, 245,

- 247, 248, 255, 256, 519
- IOMSG= specifier, 221, 225, 226, 244, 245, 247, 248, 255, **256**, 519
- IOR, 355, **388**
- IOSTAT= specifier, 221, 225, 226, 240, 244, 245, 247, 248, 254, 255, **256**, 390, 435, 518, 528, 555
- IOSTAT\_END, 240, 256, **435**
- IOSTAT\_EOR, 240, 256, **435**
- IOSTAT\_INQUIRE\_INTERNAL\_UNIT, 240, 256, **435**, 438
- IPARITY, **389**
- IS\_CONTIGUOUS, 54, **390**
- IS\_IOSTAT\_END, **390**
- IS\_IOSTAT\_EOR, **390**
- ISHFT, **389**
- ISHFTC, **389**
- ISO 10646 character, **13**, **59**, **62**, 159, 217, 218, 222, 232, 264, 278, 405, 418
- ISO\_C\_BINDING, 54, 68, 79, 93, 99, 123, 131, 132, 156, 162, 331, 432, 473–480, 520–522
- ISO\_Fortran\_binding.h, 485
- ISO\_FORTRAN\_ENV, 22, 26, 64, 100, 130–132, 137, 156, 177, 197, 203, 205, 207, 209, 216, 218, 223, 234, 239, 240, 256, 337, 338, 352–356, 376, 381, 386, 387, 404, 408, 424, 425, 427, 433–437, 529, 556
- ## K
- k* (R1014), **261**, 261, 268, 273, 276
- keyword, **13**
- argument, 9, **13**, 41, 295, 298, 309, 333, 338, 445, 504, **505**, **506**, 565
  - component, **13**, 41, 74, 82, 83, 505
  - statement, **13**, 41
  - type parameter, **13**, 41, 82
- keyword* (R216), **41**, 41, 81, 82, 307
- KIND, 55–59, 62, 85, 125, 161, 162, **390**
- kind type parameter, **21**, 25, 37, 52, 55–60, 62, 67, 77, 78, 85, 87, 157–159, 161, 281, 302, 311, 370, 435, 473, 474, 479
- kind-param* (R409), **55**, 55–57, 59, 60, 62
- kind-selector* (R406), 24, **55**, 55, 62
- ## L
- L edit descriptor, 272
- label, *see* statement label
- label* (R311), 4, **46**, 46, 181, 182, 196, 221, 225–228, 244, 245, 247–249, 255, 307, 308
- label-do-stmt* (R821), **181**, 181, 182
- language-binding-spec* (R508), 89, **93**, 106, 324
- LBOUND, 54, 160, 166, 175, 192, 312, **391**
- lbracket* (R471), 68, 86, **87**, 89, 90, 106, 107, 111, 129, 130
- LCOBOUND, **391**
- LEADZ, **392**
- left tab limit, **275**
- LEN, 125, **392**
- LEN\_TRIM, **392**
- length type parameter, **21**, 21, 37, 52, 62, 71, 87, 102, 160, 311, 392, 478, 479
- length-selector* (R422), 24, **59**, 59, 60
- letter*, **43**, 43, 44, 46, 112, 140, 142
- letter-spec* (R565), **112**, 112
- level-1-expr* (R702), **140**, 140, 143
- level-2-expr* (R706), **140**, 140, 141, 143, 144
- level-3-expr* (R710), **141**, 141
- level-4-expr* (R712), **141**, 141, 142
- level-5-expr* (R717), **142**, 142
- lexical token, 11, **14**, 22, 44, 46
- LGE, 62, **393**
- LGT, 62, **393**
- line, **14**, 47–50
- linkage association, **3**, 3, 499, 507, **510**, 510
- list-directed formatting, 237, 277–281
- list-directed input/output statement, **228**
- literal constant, **7**, 38, 122, 154
- literal-constant* (R305), **45**, 45, 139
- LLE, 62, **394**
- LLT, 62, **394**
- LOCAL, **181**, 185, 187, 506, 513, 521
- local identifier, **503**, 504
- local procedure pointer, **15**, 326
- local variable, **15**, **22**, 38, 40, 91, 92, 95, 105, 133, 135, 156, 308, 326
- local-defined-operator* (R1114), **289**, 289, 290
- local-name*, 289, 290
- LOCAL\_INIT, **181**, 185, 506, 513, 519, 521
- locality, **185**, 506, 513, 519, 521
- locality-spec* (R830), **181**, 181, 182
- LOCK statement, 198, **207**, 209, 436, 437, 520, 523
- lock variable, **22**, 36, 209, 436, 437, 520, 522
- lock-stat* (R879), **207**, 207

*lock-stmt* (R878), 31, **207**  
*lock-variable* (R881), **207**, 207, 436, 523  
 LOCK\_TYPE, 22, 64, 100, 131, 132, 207, **436**  
 LOG, 27, **394**  
 LOG10, **395**  
 LOG\_GAMMA, **395**  
 LOGICAL, **395**  
 logical intrinsic operation, 149  
 logical type, 62  
*logical-expr* (R724), **154**, 154, 167, 181–184, 187, 188, 197  
*logical-literal-constant* (R425), 45, **62**, 140, 142  
*logical-variable* (R604), **121**, 121, 207, 248, 250, 251, 523  
 LOGICAL\_KINDS, **436**  
*loop-control* (R823), **181**, 181–183, 186  
*lower-bound* (R517), **96**, 96–99  
*lower-bound-expr* (R635), **130**, 130, 164  
*lower-cobound* (R512), 94, **95**, 95

## M

*m* (R1009), **261**, 261, 265, 266, 271, 272  
 main program, **14**, 16, 19, 32, 35, 38  
*main-program* (R1101), 29, 32, **287**, 287, 298  
*mask-expr* (R746), **167**, 167–171, 181–183  
 masked array assignment, **14**, **167**, 518  
 masked array assignment (WHERE), 167  
*masked-elsewhere-stmt* (R747), **167**, 167, 168, 171  
 MASKL, **395**  
 MASKR, **396**  
 MATMUL, **396**  
 MAX, 332, 334, **397**  
 MAXEXPONENT, **397**  
 MAXLOC, 334, **397**  
 MAXVAL, **398**  
 MERGE, **399**  
 MERGE\_BITS, **400**  
 MIN, **400**  
 MINEXPONENT, **400**  
 MINLOC, **401**  
 MINVAL, **402**  
 MOD, 27, **402**  
 mode  
   blank interpretation, 222  
   changeable, 218  
   connection, 218  
   decimal edit, 222

  delimiter, 222  
   halting, 439, **443**, 443, 446, 449, 461, 465, 501, 530  
   IEEE rounding, 439, 440, **442**, 443, 444  
   input/output rounding, 218, 224, 230, 253, **271**, 276, 465  
   pad, 223  
   sign, 224, 276  
   underflow, **443**, 443, 446, 450, 462, 468, 530  
 model  
   bit, 335  
   extended real, 336  
   integer, 336  
   real, 336  
 MODULE, **296**, 297, **323**, 324, **327**  
 module, **14**, 14, 16, 19, 32, 33, 38, **287**  
*module* (R1104), 29, **288**, 298  
 module procedure, **15**, 165, 293–297, 299, 300, 305, 308, 316, 322, 324, 326, 327, 330, 432, 504, 505  
 module procedure interface body, **297**, 299  
 module reference, **16**, 288  
 MODULE statement, 287, **288**  
 module subprogram, **19**, 32, 34, 112, 321  
*module-name*, 288, 289, 508  
*module-nature* (R1110), **289**, 289  
*module-stmt* (R1105), 29, **288**, 288  
*module-subprogram* (R1108), 30, **288**, 288, 327  
*module-subprogram-part* (R1107), 29, 77, 81, **288**, 288, 291, 570  
 MODULO, 27, **403**  
 MOLD= specifier, **130**  
 MOVE\_ALLOC, 133, 198, 333, **403**  
*mp-subprogram-stmt* (R1240), 30, **327**, 327  
*mult-op* (R708), 45, **140**, 140  
*mult-operand* (R704), **140**, 140, 143  
 MVBITS, 333, 334, **404**

## N

*n* (R1016), **261**, 261, 262, 275  
 name, **14**, 41, 44, 503  
*name* (R303), 24, 41, **44**, 44, 45, 90, 111, 121, 194, 235, 305, 325  
 name association, **3**, 3, **507**, 513  
 name-value subsequence, **281**, 282  
 NAME= specifier, 89, **93**, 106, 248, 249, **250**, **305**, 305, 324, **499**  
 named constant, **7**, 21, 38, 41, 44, 52, 56, 58–60, 98, 99, 102, 105, 108, 110, 116, 122, 209, 329

*named-constant* (R306), [45](#), [45](#), [49](#), [58](#), [84](#), [110](#), [508](#)  
*named-constant-def* (R552), [110](#), [110](#), [508](#)  
 NAMED= specifier, [248](#), [251](#)  
 namelist formatting, [237](#), [281–285](#)  
 namelist input/output statement, [228](#)  
 NAMELIST statement, [115](#), [176](#), [282](#), [290](#)  
*namelist-group-name*, [115](#), [226–228](#), [234–236](#), [259](#), [282](#), [285](#), [290](#), [508](#), [522](#)  
*namelist-group-object* (R568), [115](#), [115](#), [235](#), [237](#), [243](#), [257](#), [281](#), [282](#), [290](#)  
*namelist-stmt* (R567), [31](#), [115](#), [508](#), [522](#)  
 NaN, [14](#), [266–270](#), [273](#), [346](#), [375](#), [378](#), [416](#), [420](#), [423](#), [441](#), [444](#), [447](#), [448](#), [451](#), [452](#), [466](#)  
 NEAREST, [405](#)  
 NEW\_INDEX= specifier, [206](#), [526](#)  
 NEW\_LINE, [272](#), [273](#), [405](#)  
 NEWUNIT= specifier, [218](#), [221](#), [223](#), [240](#), [519](#), [522](#)  
 NEXTREC= specifier, [248](#), [251](#)  
 NINT, [405](#)  
 NML= specifier, [226](#), [228](#), [522](#)  
 NON\_INTRINSIC, [289](#)  
 NON\_OVERRIDABLE attribute, [75](#), [76](#)  
 NON\_RECURSIVE, [323](#), [324](#)  
 nonadvancing input/output statement, [215](#)  
 nonblock DO construct, [532](#)  
 NONE, [112](#), [181](#), [298](#)  
 nonexecutable statement, [18](#), [33](#)  
*nonlabel-do-stmt* (R822), [181](#), [181](#), [182](#)  
 nonstandard intrinsic, [xviii](#), [13](#), [25](#), [573](#)  
 NOPASS, [69](#), [71](#), [76](#)  
 NOPASS attribute, *see* PASS attribute  
 NORM2, [406](#)  
 normal number, [444](#)  
 normal termination, [11](#), [34](#), [35](#), [36](#), [79](#), [137](#), [197](#), [209](#), [212](#), [224](#), [225](#), [387](#), [424](#), [437](#)  
 NOT, [406](#)  
*not-op* (R718), [45](#), [142](#), [142](#)  
 NULL, [84](#), [90](#), [155](#), [157](#), [158](#), [316](#), [334](#), [407](#), [512](#)  
*null-init* (R506), [71](#), [72](#), [90](#), [90](#), [91](#), [108](#), [109](#), [305](#)  
 NULLIFY statement, [134](#)  
*nullify-stmt* (R639), [32](#), [134](#), [523](#)  
 NUM\_IMAGES, [158](#), [408](#), [431](#)  
 NUMBER= specifier, [248](#), [251](#)  
 numeric conversion, [161](#)  
 numeric editing, [265](#)  
 numeric intrinsic operation, [146](#)

numeric sequence type, [17](#), [65](#), [116](#), [118](#), [516](#), [519](#)  
 numeric storage unit, [18](#), [18](#), [118](#), [436](#), [514](#), [519](#), [520](#)  
 numeric type, [21](#), [55–58](#), [146–148](#), [150](#), [154](#), [161](#), [370](#), [396](#), [397](#), [411](#), [425](#)  
*numeric-expr* (R727), [154](#), [154](#)  
 NUMERIC\_STORAGE\_SIZE, [436](#)

## O

O edit descriptor, [271](#)  
 object, [7](#), [7](#), [37–39](#)  
 object designator, [9](#), [37](#), [38](#), [104](#), [108](#), [122](#), [156](#), [281](#)  
*object-name* (R504), [90](#), [90](#), [106](#), [107](#), [110–112](#), [121](#), [128](#), [129](#), [508](#)  
 obsolescent feature, [24](#), [25](#), [28](#), [532–534](#)  
*octal-constant* (R466), [86](#), [86](#)  
 ONLY, [289](#), [289](#), [290](#), [298](#), [298](#), [299](#), [510](#), [562](#), [563](#)  
*only* (R1112), [289](#), [289](#), [290](#)  
*only-use-name* (R1113), [289](#), [289](#), [290](#)  
 OPEN statement, [28](#), [212](#), [213](#), [217–219](#), [220](#), [220](#), [224](#), [228](#), [236](#), [237](#), [241](#), [250](#), [254](#), [271](#), [283](#), [519](#), [522](#), [527](#), [554](#), [556–558](#)  
*open-stmt* (R904), [32](#), [220](#), [331](#)  
 OPENED= specifier, [248](#), [251](#)  
 operand, [14](#)  
 operation, [51](#)  
   defined, [8](#), [75](#), [143](#), [152](#), [153–155](#), [181](#), [293](#), [301](#), [307](#), [320](#), [331](#)  
   elemental, [10](#), [145](#), [155](#), [169](#)  
   intrinsic, [145–152](#)  
     logical, [149](#)  
     numeric, [146](#)  
     relational, [150](#)  
 OPERATOR, [51](#), [75](#), [152](#), [289](#), [296](#), [296](#), [301](#), [565](#)  
 operator, [14](#), [45](#)  
   character, [141](#)  
   defined binary, [142](#)  
   defined unary, [140](#)  
   elemental, [10](#), [145](#), [440](#)  
   logical, [142](#)  
   numeric, [140](#)  
   relational, [141](#)  
 operator precedence, [143](#)  
 OPTIONAL attribute, [102](#), [102](#), [104](#), [110](#), [156](#), [175](#), [181](#), [295](#)  
 optional dummy argument, [317](#)  
 OPTIONAL statement, [110](#), [176](#)  
*optional-stmt* (R550), [31](#), [110](#)

*or-op* (R720), 45, **142**, 142  
*or-operand* (R715), **142**, 142  
*other-specification-stmt* (R213), **30**, 30  
 OUT, **100**  
 OUT\_OF\_RANGE, **408**  
 output statement, **225**  
*output-item* (R917), 225, 226, **230**, 231, 243, 248  
 OUTPUT\_UNIT, 218, 219, 223, 239, **436**  
 override, 72, 80, 89, 90, 113, 237, 265, 516

## P

P edit descriptor, 276  
 PACK, **409**  
 pad mode, 223  
 PAD= specifier, 27, 28, 221, **223**, 226, 227, **229**, 241, 248, **251**  
 padding, **335**, 335, 388, 414  
 PARAMETER attribute, 7, 38, 85, 91, **102**, 102, 110, 122  
 PARAMETER statement, 34, **110**, 112, 292  
*parameter-stmt* (R551), 30, **110**, 508  
 parent component, 3, **6**, 74, 78, 80, 83, 516, 543  
 parent data transfer statement, 230, **237**, 237–241, 256, 280  
 parent team, **20**, 40, 129, 204, 206, 376, 381, 386, 387, 408, 424, 427, 436  
 parent type, **6**, **21**, 64, 67, 74, 78–80, 303, 543  
*parent-identifier* (R1118), **291**, 291  
*parent-string* (R609), 95, **122**, 122  
*parent-submodule-name*, 291  
*parent-type-name*, 63  
 PARENT\_TEAM, 381, **436**  
 parentheses, 154  
 PARITY, **409**  
*part-name*, 4, 123–125, 129  
*part-ref* (R612), 95, 108, 116, 122, **123**, 123–125, 127, 129  
 partially associated, **515**  
 PASS attribute, 69, **71**, 76, 307  
 passed-object dummy argument, **14**, 71, 75, 76, 80, 303, 304, 309, 310, 579  
 PAUSE statement, **531**  
 pending affector, 92, 229, 234, 501, 502  
 PENDING= specifier, 248, 249, **251**  
 pointer, 3, 8, 9, **14**, 17, 19, 22, 40, 65, 70, 129, 134–136, 158, 294, 295, 312, 330, 424, 473, 485, 511  
     procedure, 477

    pointer assignment, **15**, 97, 99, 134, 162, **163**, **164**, 317, 512  
     pointer assignment statement, 15, 19, 52, 71, 83, 155, **163**, 165, 171, 344, 350  
     pointer association, **3**, 3, 9, 19, 21, 38, 79, 81, 83, 95, 100, 102, 104, 105, 123, 135, 137, 163, 165, 166, 185, 199, 203, 234, 294, 310, 312, 314, 316, 317, 325, 326, 339, 350, 407, 476, 477, 511–514, 516, 517, 523  
     pointer association context, 100, 103, 330, **523**  
     pointer association status, **511**  
     POINTER attribute, 2, 14, 52–54, 62, 68–70, 90, 92, 97, 99, **102**, 102, 104, 109, 111, 123, 126, 134, 164, 175, 185, 191, 192, 293–295, 297, 302, 303, 305, 313, 316–319, 324, 330–332, 478, 482, 497, 511, 513, 516, 517, 579, 583  
     POINTER statement, **110**, 292  
     *pointer-assignment-stmt* (R733), 32, 103, **163**, 170, 330, 523  
     *pointer-decl* (R554), **110**, 110  
     *pointer-object* (R640), **134**, 134, 523  
     *pointer-stmt* (R553), 31, **110**, 508  
     polymorphic, **15**, 26, 54, 71, 98, 99, 123, 134, 154, 159, 160, 165, 174, 175, 193, 195, 231, 237, 294, 295, 307, 310–313, 330, 331, 375, 403, 414, 417, 424, 511, 517  
     POPCNT, **410**  
     POPPAR, **410**  
     POS= specifier, 214–216, 226, 227, **230**, 230, 248, **252**, 528  
     position edit descriptor, 274  
     *position-edit-desc* (R1015), **261**, 261  
     *position-spec* (R927), **245**, 245  
     POSITION= specifier, 221, **223**, 248, **252**, 557  
     positional arguments, 333  
     potential subobject component, **6**, 62–64, 331  
     *power-op* (R707), 45, **140**, 140  
     pre-existing, **516**  
     precedence of operators, 143  
     PRECISION, 56, **410**, 460  
     preconnected, **15**, 213, 218–221, 228, 234, 434–436  
     preconnection, **220**  
     *prefix* (R1227), **323**, 323, 324, 326  
     *prefix-spec* (R1228), **323**, 323, 324, 330, 332  
     PRESENT, 54, 68, 69, 102, 156, 317, 334, **410**, 577  
     present, **317**

- primary, 139
- primary* (R701), 139, 139, 140, 329
- PRINT statement, 213, 222, 225, 234, 239, 241, 244
- print-stmt* (R912), 32, 226, 331
- PRIVATE attribute, 65, 80, 92, 92, 105, 115, 331, 563
- PRIVATE statement, 74, 75, 77, 105, 290
- private-components-stmt* (R445), 64, 74, 74
- private-or-sequence* (R429), 63, 64, 64
- proc-attr-spec* (R1215), 304, 304, 305
- proc-component-attr-spec* (R442), 69, 69
- proc-component-def-stmt* (R441), 68, 69, 69
- proc-component-ref* (R739), 164, 164, 165, 307, 317
- proc-decl* (R1216), 69, 72, 304, 305, 305
- proc-entity-name*, 111
- proc-interface* (R1214), 69, 304, 304, 305
- proc-language-binding-spec* (R1231), 92, 93, 304, 305, 324, 324–326, 329, 483
- proc-pointer-init* (R1218), 305, 305
- proc-pointer-name* (R558), 111, 111, 134, 164
- proc-pointer-object* (R738), 163, 164, 165, 170, 350, 523
- proc-target* (R740), 82–84, 103, 163, 164, 165, 170, 317, 350, 514
- PROCEDURE, 69, 75, 304, 327
- procedure, 8, 15, 16, 42, 102, 296
  - characteristics of, 294
  - dummy, 5, 9, 11, 15, 99, 112, 157, 165, 293, 294, 296, 297, 299, 300, 303–305, 308, 316, 322, 324, 330, 500, 504, 509
  - elemental, 10, 39, 155, 165, 305, 308, 317, 321, 323, 332, 332–334
  - external, 15, 25, 32, 75, 99, 112, 165, 204, 293, 294, 296–300, 304, 305, 308, 316, 322, 323, 503, 504, 508, 509, 564, 565, 570, 573, 574
  - internal, 11, 15, 32, 165, 293–296, 308, 316, 322, 324, 326, 500, 504, 505, 509
  - intrinsic, 333–432
  - module, 15, 165, 293–297, 299, 300, 305, 308, 316, 322, 324, 326, 327, 330, 432, 504, 505
  - non-Fortran, 329
  - pure, 15, 26, 170, 181, 330, 331, 333, 339, 403, 414, 432, 445, 477
  - type-bound, 4, 12, 14, 15, 62–64, 71, 76, 76, 77, 79, 80, 162, 243, 289, 301, 307, 310, 311, 323, 504, 505
- procedure declaration statement, 34, 99, 295, 297, 304, 329, 344, 505
- procedure designator, 9, 16, 39
- procedure interface, 295
- procedure pointer, 5, 11, 14, 14, 15, 33, 90, 99, 102, 117, 165, 166, 293, 294, 300, 304, 308, 310, 316, 317, 322, 326, 477, 508
- procedure reference, 2, 16, 26, 39, 102, 125, 240, 293, 301, 307, 309
  - generic, 303
  - resolving, 320
  - type-bound, 323
- PROCEDURE statement, 296, 300
- procedure-component-name*, 164
- procedure-declaration-stmt* (R1213), 30, 304, 305
- procedure-designator* (R1223), 307, 307, 317, 323
- procedure-entity-name*, 305, 306
- procedure-name*, 75, 76, 165, 166, 296, 297, 305, 307, 308, 327
- procedure-stmt* (R1206), 296, 296, 297
- processor, 15, 24, 25, 41, 42
- processor dependent, 15, 25, 41, 525–530
- PRODUCT, 411
- program, 16, 24, 25, 32
- program* (R201), 29
- PROGRAM statement, 287
- program unit, 14, 16, 16, 19, 24, 25, 29, 32–34, 36, 40, 41, 43, 44, 46–49, 66, 103, 112, 218, 224, 287, 291, 499, 503, 511, 525, 533, 535, 561–563, 565, 569–571, 573, 590
- program-name*, 287
- program-stmt* (R1102), 29, 287, 287
- program-unit* (R202), 24, 29, 29, 32
- PROTECTED attribute, 103, 103, 111, 116, 185, 289
- PROTECTED statement, 111
- protected-stmt* (R555), 31, 111
- PUBLIC attribute, 80, 92, 92, 105, 115, 563
- PUBLIC statement, 105, 290
- PURE, 323, 324, 328, 330
- pure procedure, 15, 26, 170, 181, 330, 331, 333, 339, 403, 414, 432, 445, 477

## Q

QUIET= specifier, 197

## R

- r* (R1006), 260, 260–262
- RADIX, 56, 411, 440, 460
- RANDOM\_INIT, 344, 412, 529

RANDOM\_NUMBER, 344, **412**, 413  
 RANDOM\_SEED, 334, 344, **412**  
 RANGE, 55, 56, **413**, 460  
 RANK, 54, **413**  
 rank, **16**, 17, 37–40, 69, 71, 77, 78, 82, 84, 89, 93, 95–99, 110, 118, 123–127, 129, 131, 133, 152, 154, 155, 159–161, 163–165, 167, 175, 201, 294, 301–303, 312–314, 317, 322, 332, 343, 348, 349, 366–368, 371, 372, 377, 383, 384, 389, 391, 392, 396, 398, 399, 401–403, 406, 407, 409, 411, 412, 415, 420, 422, 423, 425, 428–431, 469, 476, 482, 507, 515, 579, 580  
 RANK ( \* ), 98, **191**  
 RANK DEFAULT, 99, **191**  
*rbracket* (R472), 68, 86, **87**, 89, 90, 106, 107, 111, 129, 130  
 RC edit descriptor, 276  
 RD edit descriptor, 276  
 READ (FORMATTED), 238, **296**  
 READ (UNFORMATTED), 238, **296**  
 READ statement, 28, 38, 214, 218, 222, **225**, 234, 240, 241, 244, 247, 255, 521, 555–557, 559, 560  
*read-stmt* (R910), 32, **225**, 226, 331, 522  
 READ= specifier, 248, **252**  
 READWRITE= specifier, 248, **252**  
 REAL, 27, 161, 335, **414**, 441  
 real and complex editing, 266  
 real model, 336  
 real part, **58**  
 real type, 56–57, 58  
*real-literal-constant* (R414), 45, **57**, 57  
*real-part* (R419), **58**, 58  
 REAL128, **437**  
 REAL32, **437**  
 REAL64, **437**  
 REAL\_KINDS, **437**  
 REC= specifier, 215, 226, 227, **230**  
 RECL= specifier, 221, **223**, 236, 237, 248, **252**, 254, 520, 527  
 record, **16**, 211  
 record file, 10, **16**, **211**, 213, 215–217, 527  
 record number, **213**  
 RECURSIVE, **323**  
 recursive input/output statement, **256**  
 REDUCE, **414**  
 reference, **16**, 39  
     procedure, 26  
*rel-op* (R713), 45, **141**, 141, 151, 441  
 relational intrinsic operation, 150  
*rename* (R1111), **289**, 289, 290, 504  
*rep-char*, **60**, 60, 262, 278, 283  
 REPEAT, **415**  
 repeat specification, **260**  
 representation method, 55, 56, 59, 62  
 RESHAPE, 87, 88, **416**  
 resolving procedure reference, 320  
 resolving procedure references  
     defined input/output, 243  
 restricted expression, **156**  
 RESULT, 324, **325**, 325, 327, 328, 505  
*result-name*, 324, 325, 327, 328, 508  
 RETURN statement, 35, 79, 103, 118, 135, 136, 179, 184, **328**, 476, 522  
*return-stmt* (R1243), 32, 34, **328**, 328  
 REWIND statement, 212, 213, 215, 241, 244, **246**, 246, 555  
*rewind-stmt* (R926), 32, **245**, 331  
 RN edit descriptor, 276  
 round edit descriptor, 276  
*round-edit-desc* (R1019), 261, **262**  
 ROUND= specifier, 221, **224**, 226, 227, **230**, 241, 248, **253**, 277  
 rounding mode  
     IEEE, 439, 440, **442**, 443, 444, 450, 459, 461, 466  
     input/output, 218, 224, 230, 253, **271**, 276, 465  
 RP edit descriptor, 276  
 RRSPACING, **416**  
 RU edit descriptor, 276  
 RZ edit descriptor, 276

## S

S edit descriptor, 276  
 SAME\_TYPE\_AS, 68, 69, **417**  
 SAVE attribute, 16, 22, 28, 72, 79, 91, 93, 94, **103**, 103, 104, 107, 111, 116, 118, 136, 185, 305, 330, 513  
 SAVE statement, **111**, 176, 292, 505  
*save-stmt* (R556), 31, **111**, 508  
 saved, **16**, 511, 517  
*saved-entity* (R557), **111**, 111, 176  
 scalar, **16**, 17, 18, 332  
*scalar-xyz* (R103), 24, 24  
 SCALE, **417**  
 scale factor, **261**, 276

- SCAN, **418**
- scoping unit, **3**, **11**, **12**, **16**, **22**, **32**, **34**, **35**, **38**, **41**, **60**, **65**, **66**, **75**, **78**, **82**, **91–93**, **99**, **101**, **104**, **105**, **111**, **112**, **115**, **117**, **118**, **135**, **157**, **165**, **176**, **181**, **185**, **229**, **232**, **288–290**, **295**, **297–299**, **303**, **320–322**, **324**, **327**, **329**, **439**, **441**, **499**, **504–510**, **513**, **515**, **516**, **520**, **562**, **563**, **573**, **579**
- section subscript, **127**
- section-subscript* (R620), **22**, **123**, **125**, **125**, **127–129**
- segment, **199**
- SELECT CASE construct, **35**, **188**, **533**, **545**
- SELECT CASE statement, **47**, **189**
- SELECT RANK construct, **35**, **98**, **99**, **191**, **317**, **522**
- SELECT RANK statement, **99**, **191**, **510**
- SELECT TYPE construct, **35**, **52**, **53**, **193**, **317**, **506**, **507**, **510**, **522**
- SELECT TYPE statement, **47**, **193**, **510**
- select-case-stmt* (R841), **4**, **188**, **189**, **189**, **196**
- select-construct-name*, **191–194**
- select-rank-case-stmt* (R850), **191**, **191**, **192**
- select-rank-construct* (R848), **31**, **191**, **191**, **192**
- select-rank-stmt* (R849), **4**, **191**, **191**, **192**, **196**
- select-type-construct* (R852), **31**, **193**, **194**
- select-type-stmt* (R853), **4**, **193**, **193**, **194**, **196**
- SELECTED\_CHAR\_KIND, **59**, **418**
- SELECTED\_INT\_KIND, **55**, **419**
- SELECTED\_REAL\_KIND, **56**, **334**, **419**, **535**
- selector, **174**
- selector* (R805), **174**, **174**, **175**, **191**, **193–195**, **317**, **511**, **522**
- separate module procedure, **326**
- separate module subprogram statement, **327**
- separate-module-subprogram* (R1239), **30**, **288**, **326**, **327**, **327**
- sequence, **17**
- sequence association, **316**
- SEQUENCE attribute, **17**, **63**, **64**, **65**, **66**, **79**, **117**, **164**, **165**, **194**, **480**
- SEQUENCE statement, **64**
- sequence structure, **17**
- sequence type, **17**, **17**, **26**, **63**, **65**, **65**, **116**, **436**, **437**, **481**, **514**
- character, **17**, **65**, **116**, **118**, **515**, **516**, **519**
- numeric, **17**, **65**, **116**, **118**, **516**, **519**
- sequence-stmt* (R431), **64**, **64**
- sequential access, **213**
- sequential access data transfer statement, **230**
- SEQUENTIAL= specifier, **248**, **253**
- SET\_EXPONENT, **420**
- SHAPE, **54**, **420**
- shape, **17**, **39**, **312**
- SHARED, **181**, **185**
- SHIFTA, **420**
- SHIFTL, **421**
- SHIFTR, **421**
- SIGN, **27**, **28**, **57**, **421**
- sign* (R412), **55**, **56**, **56**, **57**, **266**
- sign mode, **224**, **265**, **276**
- sign-edit-desc* (R1017), **261**, **262**
- SIGN= specifier, **221**, **224**, **226**, **227**, **230**, **248**, **253**, **276**
- signed-digit-string* (R410), **56**, **57**, **265–267**
- signed-int-literal-constant* (R407), **55**, **55**, **58**, **108**, **261**
- signed-real-literal-constant* (R413), **57**, **58**, **108**
- significand* (R415), **57**, **57**
- simply contiguous, **17**, **128**, **129**, **129**, **165**, **311–315**, **484**, **485**
- SIN, **422**
- SINH, **422**
- SIZE, **54**, **422**
- size, **17**, **39**
- size of a common block, **118**
- SIZE= specifier, **226**, **230**, **248**, **253**, **255**, **518**, **521**, **555**, **556**
- source-expr* (R631), **130**, **130–134**, **331**, **512**, **513**
- SOURCE= specifier, **130**, **133**, **331**, **434**, **519**, **521**
- SP edit descriptor, **276**
- SPACING, **423**
- special character, **43**
- specific interface, **12**, **243**, **297**, **297**, **300**, **305**, **321**
- specific interface block, **12**, **13**, **297**
- specific name, **17**
- specific-procedure* (R1207), **296**, **296**, **298**, **300**
- specification, **89–119**
- specification expression, **17**, **21**, **35**, **67–69**, **78**, **92**, **125**, **156**, **157**, **157**, **176**, **328**, **441**, **525**, **526**, **533**
- specification function, **157**
- specification inquiry, **156**
- specification-construct* (R208), **30**, **30**, **176**
- specification-expr* (R728), **4**, **53**, **59**, **91**, **95**, **96**, **156**, **156**, **332**
- specification-part* (R204), **29**, **30**, **30**, **35**, **75**, **92**, **105**,

- 157, 158, 287, 288, 291, 292, 296, 298, 324, 326, 327, 330, 332
- SPREAD, **423**
- SQRT, **27**, **424**, 445, 466, 467
- SS edit descriptor, **276**
- standard intrinsic, xviii, **13**, 25, 432
- standard-conforming program, **18**, 24
- stat-variable* (R629), 35, 129, **130**, 130, 132, 134, 136, 137, 200, 206, 209, 522, 526
- STAT= specifier, xix, 130, 133–135, **136**, 180, 200, **208**, 437, 519, 526
- STAT\_FAILED\_IMAGE, 36, 130, 338, 376, 404, **437**, 438
- STAT\_LOCKED, 209, **437**, 438
- STAT\_LOCKED\_OTHER\_IMAGE, 209, **437**, 438
- STAT\_STOPPED\_IMAGE, 137, 209, 337, 404, 424, **437**, 438
- STAT\_UNLOCKED, 209, **437**, 438
- statement, **18**, 47
  - accessibility, **105**
  - ALLOCATABLE, **106**
  - ALLOCATE, 52, 53, 60, 92, 95, 97, **130**, 136, 137, 164, 198, 434, 491, 492, 511, 512, 519, 521, 526
  - arithmetic IF, **532**
  - ASSIGN, **531**
  - assigned GO TO, **531**
  - assignment, 14, 15, 38, 52, 78, **159**, 171, 469, 518
  - ASSOCIATE, **174**, 510
  - ASYNCHRONOUS, **106**, 176, 292, 506, 509
  - attribute specification, 105–119
  - BACKSPACE, 212, 215, 241, 244, **245**, 246, 555, 557
  - BIND, **106**, 292, 499, 505
  - BLOCK, 40, 91, 95, 97, **175**, 519
  - BLOCK DATA, 47, 287, **291**
  - CALL, 19, 196, 198, 293, **307**, 320, 328, 404
  - CASE, **189**
  - CHANGE TEAM, 19, 35, 40, **177**, 198, 337, 437, 510, 522
  - CLASS DEFAULT, **193**
  - CLASS IS, **193**
  - CLOSE, 212, 213, 217, 219, 220, **224**, 224, 241, 244, 555
  - COMMON, 6, **117**, 117–119, 176, 290, 292, 505, 515, 516, 532
  - component definition, **68**
  - computed GO TO, 4, **196**, 196, 532, 533
  - CONTAINS, 34, **75**, **328**
  - CONTIGUOUS, **107**
  - CONTINUE, **197**, 531
  - CRITICAL, 153, **179**, 198
  - CYCLE, 174, 180, **184**, 184, 534
  - DATA, 28, 34, 86, 91, **107**, 118, 292, 407, 506, 509, 517, 532, 533
  - data transfer, 28, 46, 211–217, 219, **225**, 230, 233–235, 240, 244, 246, 254–257, 259, 260, 271, 276, 278–280, 282, 284, 435, 436, 518, 520, 527, 528, 555, 556, 558, 560
  - DEALLOCATE, **134**, 136, 137, 198, 434, 492, 526
  - defined assignment, **163**, 163, 320, 522
  - derived type definition, *see* statement, TYPE
  - DIMENSION, **109**, 292
  - DO, **180**, 518, 532, 534
  - DO CONCURRENT, 53, 170, **180**
  - DO WHILE, **180**
  - ELSE, **187**
  - ELSE IF, 47, **187**
  - ELSEWHERE, 47, **167**
  - END, **10**, **34**, 103, 118, 135, 136, 198, 476, 522
  - END ASSOCIATE, 47, **174**
  - END BLOCK, 47, 136, **176**
  - END BLOCK DATA, 47, **291**
  - END CRITICAL, 47, 153, **179**, 198
  - END DO, 47, **182**
  - END ENUM, 47, **84**
  - END FORALL, 47, **170**
  - END FUNCTION, 47, **325**
  - END IF, 47, **187**, 531
  - END INTERFACE, 47, **296**
  - END MODULE, 47, **288**
  - END PROCEDURE, 47, **327**
  - END PROGRAM, 47, **287**
  - END SELECT, 47, **189**, **194**
  - END SUBMODULE, 47, **291**
  - END SUBROUTINE, 47, **326**
  - END TEAM, 35, **177**, 198, 337, 437
  - END TYPE, 47, **64**
  - END WHERE, 47, **167**
  - ENDFILE, 47, 212, 213, 215, 222, 241, 244, **246**, 555
  - ENTRY, 9, 34, 152, 163, 288, 293, 297, 323, 324, **327**, 505, 515, 516, 532, 534

- ENUM, [84](#)
- ENUMERATOR, [84](#)
- EQUIVALENCE, [115](#), [115–119](#), [176](#), [290](#), [292](#), [515](#), [516](#), [532](#), [534](#)
- ERROR STOP, [35](#), [36](#), [197](#), [526](#)
- EVENT POST, [198](#), [205](#), [206](#), [209](#), [373](#), [434](#), [435](#), [520](#), [522](#), [526](#)
- EVENT WAIT, [198](#), [205](#), [205](#), [209](#), [434](#), [435](#), [520](#), [522](#), [526](#)
- executable, [18](#), [18](#), [33](#)
- EXIT, [174](#), [184](#), [196](#)
- EXTERNAL, [99](#), [304](#)
- FAIL IMAGE, [197](#)
- file inquiry, [247](#)
- file positioning, [212](#), [245](#)
- FINAL, [10](#), [77](#)
- FLUSH, [213](#), [244](#), [247](#), [256](#)
- FORALL, [53](#), [145](#), [171](#), [506](#), [507](#), [518](#)
- FORM TEAM, [19](#), [20](#), [35](#), [129](#), [198](#), [206](#), [437](#), [523](#), [526](#)
- FORMAT, [25](#), [34](#), [46](#), [228](#), [259](#), [259](#), [288](#)
- formatted input/output, [211](#), [227](#)
- FUNCTION, [9](#), [53](#), [113](#), [152](#), [156](#), [287](#), [323](#), [324](#), [327](#), [328](#), [505](#)
- GENERIC, [75](#), [76](#), [298](#), [300](#), [321](#)
- GO TO, [4](#), [47](#), [196](#), [196](#)
- IF, [145](#), [188](#)
- IMPLICIT, [34](#), [112](#), [115](#), [292](#)
- IMPLICIT NONE, [112](#)
- IMPORT, [34](#), [298](#), [509](#)
- input/output, [211–256](#), [518](#)
- INQUIRE, [27](#), [213](#), [214](#), [216](#), [217](#), [219](#), [220](#), [229](#), [230](#), [240](#), [241](#), [244](#), [247](#), [256](#), [257](#), [435](#), [518](#), [520–522](#), [528](#), [554](#)
- INTENT, [110](#), [176](#)
- INTERFACE, [296](#), [565](#)
- INTRINSIC, [292](#), [306](#)
- intrinsic assignment, [83](#), [92](#), [135](#), [137](#), [155](#), [159](#), [163](#), [165](#), [209](#), [247](#), [256](#), [279](#), [283](#), [330](#), [331](#), [518](#), [526](#)
- list-directed input/output, [228](#)
- LOCK, [198](#), [207](#), [209](#), [436](#), [437](#), [520](#), [523](#)
- MODULE, [287](#), [288](#)
- NAMELIST, [115](#), [176](#), [282](#), [290](#)
- namelist input/output, [228](#)
- nonexecutable, [18](#), [33](#)
- NULLIFY, [134](#)
- OPEN, [28](#), [212](#), [213](#), [217–219](#), [220](#), [220](#), [224](#), [228](#), [236](#), [237](#), [241](#), [250](#), [254](#), [271](#), [283](#), [519](#), [522](#), [527](#), [554](#), [556–558](#)
- OPTIONAL, [110](#), [176](#)
- PARAMETER, [34](#), [110](#), [112](#), [292](#)
- PAUSE, [531](#)
- POINTER, [110](#), [292](#)
- pointer assignment, [15](#), [19](#), [52](#), [71](#), [83](#), [155](#), [163](#), [165](#), [171](#), [344](#), [350](#)
- PRINT, [213](#), [222](#), [225](#), [234](#), [239](#), [241](#), [244](#)
- PRIVATE, [74](#), [75](#), [77](#), [105](#), [290](#)
- PROCEDURE, [296](#), [300](#)
- procedure declaration, [34](#), [99](#), [295](#), [297](#), [304](#), [329](#), [344](#), [505](#)
- PROGRAM, [287](#)
- PROTECTED, [111](#)
- PUBLIC, [105](#), [290](#)
- READ, [28](#), [38](#), [214](#), [218](#), [222](#), [225](#), [234](#), [240](#), [241](#), [244](#), [247](#), [255](#), [521](#), [555–557](#), [559](#), [560](#)
- RETURN, [35](#), [79](#), [103](#), [118](#), [135](#), [136](#), [179](#), [184](#), [328](#), [476](#), [522](#)
- REWIND, [212](#), [213](#), [215](#), [241](#), [244](#), [246](#), [246](#), [555](#)
- SAVE, [111](#), [176](#), [292](#), [505](#)
- SELECT CASE, [47](#), [189](#)
- SELECT RANK, [99](#), [191](#), [510](#)
- SELECT TYPE, [47](#), [193](#), [510](#)
- separate module subprogram, [327](#)
- SEQUENCE, [64](#)
- statement function, [9](#), [34](#), [60](#), [176](#), [288](#), [321](#), [328](#), [329](#), [506](#), [507](#), [533](#)
- STOP, [35](#), [36](#), [197](#), [198](#), [526](#)
- SUBMODULE, [287](#), [291](#)
- SUBROUTINE, [9](#), [163](#), [287](#), [323](#), [326](#), [328](#)
- SYNC ALL, [198](#), [200](#), [201](#), [202](#), [209](#), [437](#)
- SYNC IMAGES, [198](#), [201](#), [209](#), [437](#)
- SYNC MEMORY, [198](#), [203](#), [209](#), [584](#)
- SYNC TEAM, [198](#), [204](#), [437](#)
- TARGET, [111](#), [292](#)
- TYPE, [63](#), [67](#), [92](#), [508](#)
- type declaration, [34](#), [53](#), [72](#), [89](#), [89–91](#), [99](#), [112](#), [115](#), [118](#), [157](#), [288](#), [292](#), [325](#), [327](#), [329](#), [517](#)
- type guard, [60](#), [193](#)
- TYPE IS, [193](#)
- type parameter definition, [66](#)
- type-bound procedure, [75](#), [76](#)

- unformatted input/output, 212, **227**
- UNLOCK, 198, **207**, 209, 436, 437, 520, 523
- USE, 3, 16, 34, 66, 105, **288**, 292, 321, 322, 504, 508, 510, 562, 563, 569
- VALUE, **112**, 176
- VOLATILE, **112**, 176, 292, 506, 509
- WAIT, 219, 229, **244**, 244, 559
- WHERE, 14, 145, **167**
- WRITE, 213, 218, 222, **225**, 234, 239, 241, 244, 256, 518, 555, 558, 559
- statement entity, **18**, 503, 506
- statement function, 329, 533
- statement function statement, 9, 34, 60, 176, 288, 321, 328, **329**, 506, 507, 533
- statement keyword, **13**, 41
- statement label, 4, **18**, **46**, 46–49, 308, 503
- statement order, 33
- STATUS= specifier, 220, 221, **224**, **225**, 225, 527, 557
- stmt-function-stmt* (R1245), 30, 288, 291, 297, **329**, 508
- STOP statement, 35, 36, **197**, 198, 526
- stop-code* (R862), **197**, 197
- stop-stmt* (R860), 32, 79, **197**
- stopped image, **11**, 35, 337, 404
- STOPPED\_IMAGES, **424**
- storage association, **3**, 3, 40, 41, 115–119, 327, 330, 425, 513–516
- storage sequence, **18**, 63, 65, 116–119, 292, 350, 351, 514, 514, 515
- storage unit, **18**, 18, 115–119, 229, 233, 241, 244, 292, 316, 350, 351, 514–516
  - character, **18**, 18, 98, 116–118, 433, 514, 519, 520
  - file, **10**, 18, 211, 214–217, 223, 228, 230, 236, 246, 252–254, 435, 514, 527
  - numeric, **18**, 18, 118, 436, 514, 519, 520
  - unspecified, **18**, 18, 514, 515, 519, 521
- STORAGE\_SIZE, 86, **424**
- stream access, **214**
- stream access data transfer statement, **230**
- stream file, 10, **18**, **211**, 214, 216, 254, 527
- STREAM= specifier, 248, **253**
- stride* (R622), **126**, 126–128, 233
- structure, 6, **18**, 37, 63
- structure component, **19**, 108, 122, 123, 125, 480, 543
- structure constructor, 6, 13, **19**, 37, 41, 51, 74, 82, 83, 108, 109, 154, 156, 157, 407, 436, 505, 538
- structure-component* (R613), 107, 108, 121, 122, **123**, 129, 130, 134
- structure-constructor* (R456), 19, **82**, 82, 108, 139, 330
- subcomponent, **6**, 8, 72, 82, 164, 512, 513, 517, 519, 521
- submodule, 14, 16, **19**, 19, 32, 33, 38, **291**
- submodule* (R1116), 29, **291**
- submodule identifier, **291**
- SUBMODULE statement, 287, **291**
- submodule-name*, 291
- submodule-stmt* (R1117), 29, **291**, 291
- subobject, 2, 6, 7, **19**, 37–39, 100, **123**, 311, 512
- subprogram, 14, **19**, 32, 34, 35, 38, 40, 112
  - elemental, **10**, 323, 324, 332
  - external, 16, **19**, 32, 293
  - internal, **19**, 32, 34, 293, 508
  - module, **19**, 32, 34
- subroutine, **19**
  - atomic, **19**, 36, 199, 200, 333, 336, 338, 352–356, 373, 433, 437, 522
  - collective, **19**, 333, 337, 338, 361–364, 376, 424, 437
- subroutine reference, 320
- SUBROUTINE statement, 9, 163, 287, 323, **326**, 328
- subroutine-name*, 297, 326, 505
- subroutine-stmt* (R1236), 29, 296, 297, 324, **326**, 326, 505, 508
- subroutine-subprogram* (R1235), 19, 29, 30, 288, **326**, 326
- subroutines
  - intrinsic, 333
- subscript, 125
  - section, 127
- subscript* (R619), 108, **125**, 125, 126, 128, 233
- subscript triplet, 127
- subscript-triplet* (R621), 125, **126**, 126–128
- substring, **122**
- substring* (R608), 115, 116, 121, **122**
- substring ending point., **122**
- substring starting point, **122**
- substring-range* (R610), 95, **122**, 122, 124–126, 129, 233
- suffix* (R1233), 324, **325**, 327
- SUM, **425**
- SYNC ALL statement, 198, **200**, 201, 202, 209, 437
- SYNC IMAGES statement, 198, **201**, 209, 437
- SYNC MEMORY statement, 198, **203**, 209, 584
- SYNC TEAM statement, 198, **204**, 437
- sync-all-stmt* (R864), 32, **200**
- sync-images-stmt* (R866), 32, **201**

*sync-memory-stmt* (R868), 32, 203  
*sync-stat* (R865), 177, 179, 200, 200, 201, 203–207  
*sync-team-stmt* (R869), 32, 204  
synchronous input/output, 222, 228, 230, 233  
SYSTEM\_CLOCK, 426

## T

T edit descriptor, 275  
TAN, 426  
TANH, 426  
target, 9, 19, 38, 40, 54, 71–73, 78, 84, 91, 94, 95, 97, 99–103, 109, 121, 123, 130, 133, 134, 136, 154, 155, 160, 163, 164, 166, 170, 171, 231, 235, 237, 305, 308, 310, 313, 314, 316, 475, 477, 478, 511–514, 516, 519, 521–523  
TARGET attribute, 3, 19, 72, 102, 104, 104, 111, 116, 118, 134, 135, 164, 175, 185, 192, 295, 302, 311, 312, 314, 318, 319, 363, 403, 475, 478, 498, 512, 513, 521, 576, 577  
TARGET statement, 111, 292  
*target-decl* (R560), 111, 111  
*target-stmt* (R559), 31, 111, 508  
team, 12, 19, 19, 39, 40, 129, 130, 136, 177, 178, 200, 201, 204, 206, 333, 337, 386, 387, 408, 424, 427, 428, 526  
    current, 337  
team number, 20, 129, 206  
team variable, 22, 177, 438  
*team-construct-name*, 177  
*team-number* (R876), 206, 206  
*team-variable* (R815), 129, 177, 177–179, 204, 206, 523  
TEAM= specifier, 129, 129  
TEAM\_NUMBER, 427  
TEAM\_NUMBER= specifier, 129, 129  
TEAM\_TYPE, 22, 177, 376, 381, 386, 387, 408, 424, 427, 437  
THEN, 187  
THIS\_IMAGE, 158, 427  
TINY, 428  
TKR compatible, 303  
TL edit descriptor, 275  
totally associated, 515  
TR edit descriptor, 275  
TRAILZ, 428  
TRANSFER, 158, 429  
transfer of control, 173, 196, 255

transformational function, 20, 158, 329, 333, 333, 334, 339, 357, 358, 433, 445  
TRANSPPOSE, 429  
TRIM, 430  
truncation, 335, 388, 414  
TYPE, 53  
type, 20, 36, 51–87  
    abstract, 20, 53, 76, 79, 79, 82, 123, 131  
    character, 59–62  
    complex, 58  
    declared, 20, 54, 71, 83, 87, 89, 123, 124, 131, 132, 134, 152, 154, 159, 162–165, 174, 194, 195, 243, 302, 307, 310, 313, 323, 329, 375, 403, 417, 507  
    derived, 8, 18, 20, 36, 37, 41, 51, 62–84, 87, 480, 481  
    dynamic, 15, 20, 22, 54, 78, 79, 81, 84, 87, 104, 132, 134, 136, 152, 154, 160, 162, 163, 165, 175, 193, 194, 203, 243, 307, 313, 323, 343, 375, 403, 417, 425, 507, 511, 517, 544, 594  
    expression, 154  
    extended, 3, 6, 12, 20, 21, 67, 74, 78–80, 516, 535, 541  
    extensible, 20, 53, 63, 71, 79, 238, 375, 417, 543, 581  
    extension, 20, 54, 79, 81, 194, 313, 375, 581  
    integer, 55–56  
    intrinsic, 6, 20, 36, 37, 51, 55–62  
    logical, 62  
    numeric, 21, 55–58, 146–148, 150, 154, 161, 370, 396, 397, 411, 425  
    operation, 155  
    parent, 6, 21, 64, 67, 74, 78–80, 303, 543  
    primary, 154  
    real, 56–57, 58  
type compatible, 21, 54, 71, 131, 159, 164, 303, 311, 403  
type conformance, 159  
type declaration statement, 34, 53, 72, 89, 89–91, 99, 112, 115, 118, 157, 288, 292, 325, 327, 329, 517  
type equality, 65  
type guard statement, 60, 193  
TYPE IS statement, 193  
type parameter, 2, 4, 12, 13, 17, 21, 26, 36, 52, 54, 55, 62, 65, 67, 71, 87, 89–91, 110, 118, 154, 159, 175, 294, 311, 332, 403, 475, 481, 505, 507  
type parameter definition statement, 66

type parameter inquiry, [21](#), [124](#), [154](#), [156](#)

type parameter keyword, [13](#), [41](#), [82](#)

type parameter order, [21](#), [67](#)

type specifier, [53](#)

CHARACTER, [59](#)

CLASS, [54](#)

COMPLEX, [58](#)

derived type, [53](#)

DOUBLE PRECISION, [57](#)

INTEGER, [55](#)

LOGICAL, [62](#)

REAL, [57](#)

TYPE, [53](#)

TYPE statement, [63](#), [67](#), [92](#), [508](#)

*type-attr-spec* (R428), [63](#), [63](#), [79](#)

type-bound procedure, [4](#), [12](#), [14](#), [15](#), [62–64](#), [71](#), [76](#), [76](#),  
[77](#), [79](#), [80](#), [162](#), [243](#), [289](#), [301](#), [307](#), [310](#), [311](#),  
[323](#), [504](#), [505](#)

type-bound procedure statement, [75](#), [76](#)

*type-bound-generic-stmt* (R451), [75](#), [75](#), [301](#)

*type-bound-proc-binding* (R448), [75](#), [75](#)

*type-bound-proc-decl* (R450), [75](#), [75](#)

*type-bound-procedure-part* (R446), [63](#), [65](#), [75](#), [77](#), [481](#)

*type-bound-procedure-stmt* (R449), [75](#), [75](#)

*type-declaration-stmt* (R501), [30](#), [60](#), [89](#), [89](#), [330](#), [508](#)

*type-guard-stmt* (R854), [193](#), [193](#), [194](#)

*type-name*, [63](#), [64](#), [66](#), [76](#), [81](#)

*type-param-attr-spec* (R434), [66](#), [67](#), [67](#)

*type-param-decl* (R433), [66](#), [66](#), [67](#)

*type-param-def-stmt* (R432), [63](#), [66](#), [67](#)

*type-param-inquiry* (R616), [21](#), [124](#), [125](#), [139](#), [505](#)

*type-param-name*, [63](#), [66](#), [67](#), [69](#), [124](#), [139](#), [505](#), [508](#)

*type-param-spec* (R455), [81](#), [81](#), [82](#)

*type-param-value* (R401), [21](#), [52](#), [52](#), [53](#), [59](#), [60](#), [69](#), [81](#),  
[82](#), [131](#), [132](#), [534](#)

*type-spec* (R402), [53](#), [53](#), [60](#), [87](#), [130–132](#), [193](#), [194](#)

## U

UBOUND, [54](#), [192](#), [312](#), [430](#)

UCOBOUND, [431](#)

ultimate argument, [21](#), [132](#), [136](#), [310](#), [315](#), [337](#)

ultimate component, [6](#), [62–64](#), [68](#), [93](#), [95](#), [98](#), [99](#), [104](#),  
[116](#), [117](#), [131](#), [133](#), [158](#), [237](#), [311](#), [514](#)

undefined, [9](#), [21](#), [22](#), [38](#), [135](#), [511](#), [512](#), [517](#), [518](#)

undefinition of variables, [517](#)

underflow mode, [443](#), [443](#), [446](#), [450](#), [462](#), [468](#), [530](#)

*underscore* (R302), [43](#), [43](#)

UNFORMATTED, [238](#), [239](#), [296](#)

unformatted data transfer, [236](#)

unformatted input/output statement, [212](#), [227](#)

unformatted record, [211](#)

UNFORMATTED= specifier, [248](#), [253](#)

Unicode file, [222](#)

unit, [6](#), [15](#), [22](#), [212–214](#), [217](#), [217–225](#), [229](#), [230](#), [233–](#)  
[235](#), [237](#), [239](#), [240](#), [244](#), [245](#), [247](#), [249–254](#), [256](#),  
[274](#), [280](#), [434](#), [435](#), [518](#), [520](#), [528](#), [555–559](#)

UNIT= specifier, [220](#), [225](#), [226](#), [244](#), [245](#), [247](#), [248](#)

unlimited polymorphic, [22](#), [54](#), [54](#), [87](#), [117](#), [131](#), [164](#),  
[194](#), [313](#), [375](#), [417](#), [424](#), [594](#)

*unlimited-format-item* (R1005), [259](#), [260](#), [260](#), [263](#)

UNLOCK statement, [198](#), [207](#), [209](#), [436](#), [437](#), [520](#), [523](#)

*unlock-stmt* (R880), [32](#), [207](#)

unordered segments, [199](#), [199](#), [336](#), [344](#), [373](#), [412](#)

UNPACK, [431](#)

unsaved, [22](#), [133](#), [135](#), [326](#), [512](#), [513](#), [519](#), [520](#), [522](#)

unspecified storage unit, [18](#), [18](#), [514](#), [515](#), [519](#), [521](#)

*until-spec* (R874), [205](#), [205](#)

UNTIL\_COUNT= specifier, [205](#), [374](#)

*upper-bound* (R518), [96](#), [96](#), [97](#)

*upper-bound-expr* (R636), [130](#), [130](#), [164](#)

*upper-cobound* (R513), [94](#), [95](#), [95](#)

use association, [3](#), [3](#), [15](#), [33](#), [41](#), [60](#), [80](#), [92](#), [103](#), [105](#),  
[106](#), [112](#), [115–117](#), [157](#), [165](#), [288](#), [287–291](#),  
[297](#), [327](#), [330](#), [506–508](#), [511](#)

USE statement, [3](#), [16](#), [34](#), [66](#), [105](#), [288](#), [292](#), [321](#), [322](#),  
[504](#), [508](#), [510](#), [562](#), [563](#), [569](#)

*use-defined-operator* (R1115), [289](#), [289](#), [290](#)

*use-name*, [289](#), [290](#), [504](#)

*use-stmt* (R1109), [30](#), [176](#), [289](#), [289](#), [508](#)

## V

*v* (R1012), [240](#), [261](#), [261](#), [274](#)

VALUE attribute, [54](#), [71](#), [77](#), [99](#), [104](#), [104](#), [112](#), [185](#),  
[234](#), [294](#), [295](#), [297](#), [301](#), [302](#), [310–312](#), [314](#), [330](#),  
[332](#), [363](#), [483](#), [484](#), [501](#), [513](#), [590](#), [592](#)

value separator, [277](#)

VALUE statement, [112](#), [176](#)

*value-stmt* (R561), [31](#), [112](#)

variable, [7](#), [19](#), [22](#), [37](#), [38](#), [41](#), [44](#), [102](#)

definition & undefinition, [517](#)

*variable* (R602), [83](#), [107](#), [108](#), [121](#), [121](#), [129](#), [159–162](#),  
[164](#), [169–171](#), [174](#), [177](#), [193](#), [207](#), [230](#), [307](#), [317](#),  
[469](#), [522](#)

*variable-name* (R603), 115, 117, 118, **121**, 121, 122, 130, 134, 164, 181, 508, 522

vector subscript, **22**, 39, 72, 95, 123, 127, 128, 174, 193, 217, 281, 311, 312, 318, 511, 513, 576

*vector-subscript* (R623), **126**, 126, 127

VERIFY, **432**

VOLATILE attribute, **104**, 104, 105, 112, 164, 165, 175, 185, 288, 290, 294, 295, 312, 313, 315, 508, 509, 514, 519, 522, 542

VOLATILE statement, **112**, 176, 292, 506, 509

*volatile-stmt* (R562), 31, **112**

## W

*w* (R1008), **261**, 261, 265–274, 278, 280, 283

wait operation, 220, 224, 230, 233, 234, **244**, 244–245, 247, 251, 252, 254, 255

WAIT statement, 219, 229, **244**, 244, 559

*wait-spec* (R923), **244**, 244

*wait-stmt* (R922), 32, **244**, 331

WHERE construct, 14, **167**

WHERE statement, 14, 145, **167**

*where-assignment-stmt* (R745), 145, **167**, 167–169, 171

*where-body-construct* (R744), **167**, 167–169

*where-construct* (R742), 31, **167**, 167, 168, 170, 171

*where-construct-name*, 167, 168

*where-construct-stmt* (R743), 4, **167**, 167, 168, 171, 196

*where-stmt* (R741), 32, **167**, 167, 168, 170, 171

WHILE, **181**, 182, 183

whole array, **22**, **125**, 125, 126, 391, 430

WRITE (FORMATTED), 238, 239, **296**

WRITE (UNFORMATTED), 238, 239, **296**

WRITE statement, 213, 218, 222, **225**, 234, 239, 241, 244, 256, 518, 555, 558, 559

*write-stmt* (R911), 32, **225**, 226, 331, 522

WRITE= specifier, 248, **254**

## X

X edit descriptor, 275

*xyz-list* (R101), **24**

*xyz-name* (R102), **24**

## Z

Z edit descriptor, 271

zero-size array, 39, 97, 108

ZZZUTIO16, 136

ZZZUTIO17, 177

ZZZUTIO18, 179

ZZZUTIO19, 178

ZZZUTIO20, 178

ZZZUTIO21, 178

ZZZUTIO22, 200

ZZZUTIO23, 201

ZZZUTIO24, 204

ZZZUTIO25, 208

ZZZUTIO26, 337

ZZZUTIO27, 338

ZZZUTIO28, 376

ZZZUTIO29, 427

ZZZUTIO30, 387