

In the following examples, *ITALIC* denotes text duplicated within the example, ~~STRIKEOUT~~ denotes text removed from the previous example, and UNDERLINE denotes text that is new or changed from the previous example.

Start with the callee and caller in F77:

<pre> FUNCTION F(X) COMPLEX F REAL X F=... END FUNCTION </pre>	<pre> SUBROUTINE S COMPLEX F C=F(Y) END SUBROUTINE </pre>
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5 New procedure attributes and facilities in F90 lead to the creation of the interface block:

<pre> FUNCTION F(X) COMPLEX F REAL, OPTIONAL :: X F=... END FUNCTION </pre>	<pre> SUBROUTINE S COMPLEX F INTERFACE <u>FUNCTION F(X)</u> <u>COMPLEX F</u> <u>REAL, OPTIONAL :: X</u> F=... <u>END FUNCTION</u> <u>END INTERFACE</u> C=F(Y) END SUBROUTINE </pre>
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Little difference for dummy procedures if actual is external procedure:

<pre> FUNCTION F(X) COMPLEX F REAL, OPTIONAL :: X F=... END FUNCTION </pre>	<pre> SUBROUTINE S(D) INTERFACE <u>FUNCTION D(X)</u> <u>COMPLEX D</u> <u>REAL, OPTIONAL :: X</u> <u>END FUNCTION</u> END INTERFACE C=<u>D</u>(Y) END SUBROUTINE </pre>
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Problem when actual is a module procedure. Original F90 "solution" is recursive USE of module.

<pre> MODULE M <u>TYPE T; ... ; END TYPE</u> </pre>	<pre> SUBROUTINE S(D) INTERFACE <u>FUNCTION D(X)</u> <u>USE M</u> <u>TYPE(T) :: D</u> <u>REAL, OPTIONAL :: X</u> <u>END FUNCTION</u> END INTERFACE C=D(Y) END SUBROUTINE </pre>
<pre> FUNCTION F(X) <u>TYPE(T) :: F</u> <u>REAL, OPTIONAL :: X</u> F=... END FUNCTION </pre>	<pre> END MODULE </pre>

That "solution" has been interpreted as not legal. Current F2K is solution is IMPORT:

```

MODULE M
TYPE T; ... ; END TYPE

FUNCTION F(X)
TYPE(T) :: F
REAL, OPTIONAL :: X
F=...
END FUNCTION

SUBROUTINE S(D)
INTERFACE
FUNCTION D(X)
IMPORT T
TYPE(T) :: D
REAL, OPTIONAL :: X
END FUNCTION
END INTERFACE
C=D(Y)
END SUBROUTINE

END MODULE

```

- 1) This put the IMPORT in the middle of text copied from F. (complicates text editing)
- 2) D gets T from S rather than M (as F does), so any declaration of the name T in S could cause D to end up wrong. (Nothing else in S has that effect.)
- 3) Some people object to the necessity of specifically identifying T as imported into D, given that you don't have to do it for F.

The proposed CONTEXT statement addresses these points:

```

MODULE M
TYPE T; ... ; END TYPE

FUNCTION F(X)
TYPE(T) :: F
REAL, OPTIONAL :: X
F=...
END FUNCTION

SUBROUTINE S(D)
INTERFACE
CONTEXT MODULE
FUNCTION D(X)
IMPORT T
TYPE(T) :: D
REAL, OPTIONAL :: X
END FUNCTION
END INTERFACE
C=D(Y)
END SUBROUTINE

END MODULE

```

- 1) CONTEXT MODULE is outside the text of the interface body.
- 2) CONTEXT MODULE is defined to go directly to M, not through S.
- 3) Since this is now host association into M (just like F), T no longer needs to be named explicitly.

01-397 also has CONTEXT EXTERNAL to allow explicit statement of the existing default. The Enhanced Modules TR could introduce CONTEXT SUBMODULE *[name]*.

Notes (from the discussion):

- 1) We might spell CONTEXT MODULE as IMPORT(MODULE) or IMPORT(M).
- 2) The existing IMPORT might be redefined to import from M.
- 3) Some people like having T named in the import list. Presumably, they would like it for F as well as for D. Allow similar import control after the CONTAINS? (But make specifying the list optional!)